

Coping With Hardness

As you have already covered in class, there are many problems that we do not think to be solvable in polynomial time. Is it then the end of the world? Not at all. It turns out that many real world problems that we care about are NP-hard.

This week, we will solve the 0-1 Knapsack problem once again. Here there are n objects each with a weight w_i and a value v_i . We want to maximize the values of the items we take, but we are also bounded by a capacity W . Formally, we want to find the subset S of the universe of objects such that $\sum_{i \in S} v_i$ is maximized subject to the constraint that $\sum_{i \in S} w_i \leq W$.

In this assignment, you will first implement Algorithm-II in the attached PDF. Augment your algorithm so that you can output the **index of the values** you **take in your final solution**.

Consider the input file to contain n and W in two separate lines. Then each of the following n lines will contain v_i and w_i .

Example Input:

```
8
65
→ 30 15
→ 45 15
→ 140 17
→ 15 50
→ 154 50
→ 7 25
→ 77 23
→ 15 50
```

You will print the total value of the items as well as the indices of the items in your final solution.

Example Output (only for the original instance):

```
Answer: 262
Used Weight: 55
Indices: 7 3 2
```

Note that $v[2] + v[3] + v[7] = 45 + 140 + 77 = 262$.

Knapsack is such a problem that admits an **FPTAS**. In easy terms, it can be arbitrarily approximated within a threshold ϵ . The attached PDF contains the details. For each $\epsilon \in [0.5, 0.2, 0.1, 0.05]$, you will run your FPTAS algorithm for knapsack. For each ϵ , print the following info.

Example Output (For attached in.txt):

```
Original Instance:
✓ Answer: 12701
Used Weight: 659
Indices: 41 37 34 29 24 17 13 12 1
```

$v_i \geq \frac{w_i}{\theta}$

$$n=4^2$$

$$m \geq 14/2$$

$$\theta = \frac{\epsilon \cdot V_{\text{max}}}{2^n}$$

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Rounded Instance with Eps: 0.5
 Theta: 8.404761904761905
 Answer of reduced instance: 1514 *← sum of values of taken obj.*
 Answer of reduced instance multiplied by theta: 12724.809523809525
 Indices: 41 39 27 24 19 12 11 10 4
 Answer of original instance (rounded up): 12680 *← (taking same indices)*
 Used Weight: 649 *✓*
 Ratio: 1.0016561514195583 → ~~Woriginal~~ *WE*

Rounded Instance with Eps: 0.2
 Theta: 3.3619047619047624
 Answer of reduced instance: 3782
 Answer of reduced instance multiplied by theta: 12714.723809523812
 Indices: 41 37 34 30 28 24 21 12 10
 Answer of original instance (rounded up): 12694
 Used Weight: 654
 Ratio: 1.000551441625965

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Rounded Instance with Eps: 0.1
 Theta: 1.6809523809523812
 Answer of reduced instance: 7562
 Answer of reduced instance multiplied by theta: 12711.361904761907
 Indices: 41 37 34 29 24 17 13 12 1
 Answer of original instance (rounded up): 12701
 Used Weight: 659
 Ratio: 1.0

Rounded Instance with Eps: 0.05
 Theta: 0.8404761904761906
 Answer of reduced instance: 15115
 Answer of reduced instance multiplied by theta: 12703.79761904762
 Indices: 41 37 34 29 24 17 13 12 1
 Answer of original instance (rounded up): 12701
 Used Weight: 659
 Ratio: 1.0

Try to generate test cases such that the ratio becomes closer to $(1 + \epsilon)$.

Submissions Guidelines

Submit your source files (.c/.cpp/.java/...) as 2105xyz.zip. In addition, attach some (at least three) input and output files you experimented with. Do not copy from anyone or anywhere. Usual guidelines will apply.

Marks Distribution

Item	Marks
Solving for Original Instance	40%
Solving for Rounded Instance	50%
Experimenting for Several Test Cases	10%

If you can generate test cases where ratio is larger than $(1 + 0.7\epsilon)$, you can claim a 10% bonus.

Submission Deadline

Monday, March 04, 2024, 11:55 PM.