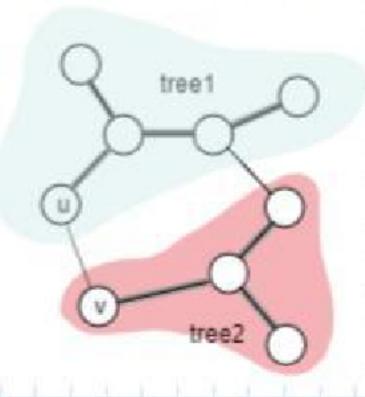
# Graph Algorithms

Dr. Samit Biswas, *Assistant Professor*, Department of Computer Sc. and Technology, Indian Institute of Engineering Science and Technology, Shibpur

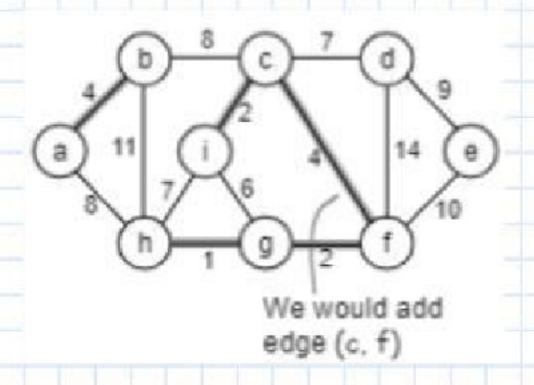
## Kruskal's Algorithm

- How is it different from Prim's algorithm?
  - Prim's algorithm grows one tree all the time
  - Kruskal's algorithm grows multiple trees (i.e., a forest) at the same time.
  - Trees are merged together using safe edges
  - Since an MST has exactly |V| 1 edges, after |V| 1 merges, we would have only one component

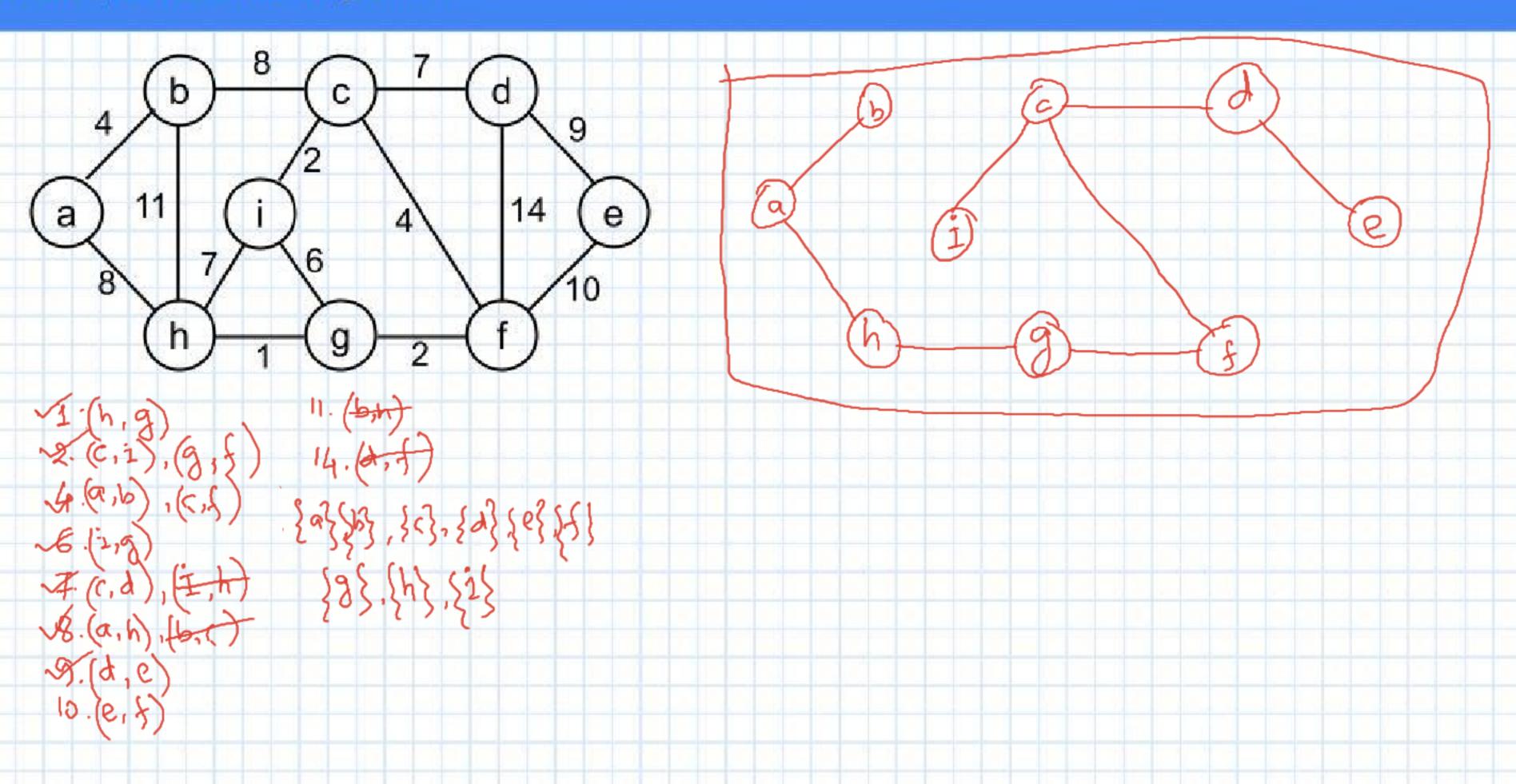


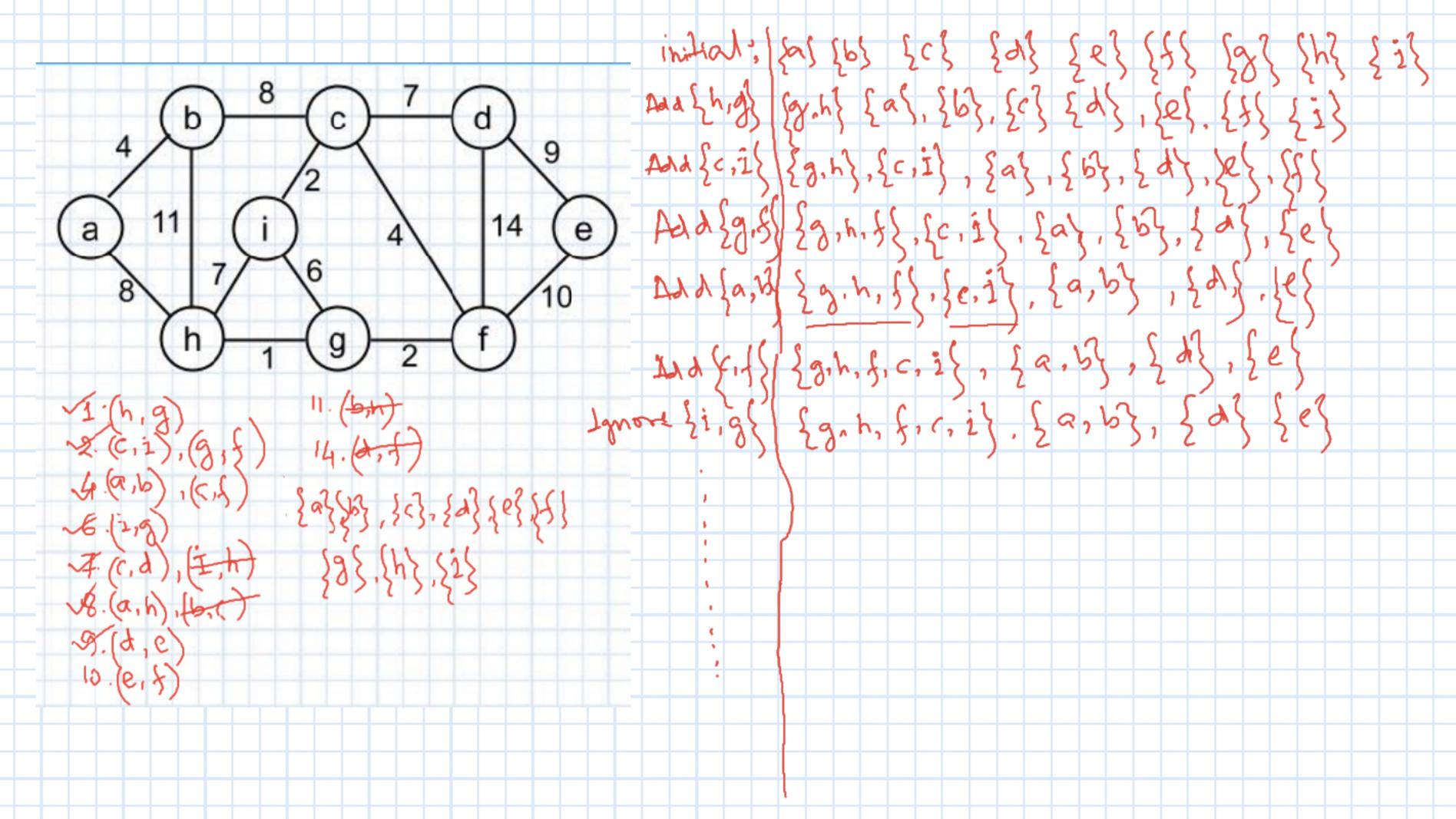
# Kruskal's Algorithm

- Start with each vertex being its own component
- Repeatedly merge two components into one by choosing the light edge that connects them
- Which components to consider at each iteration?
  - Scan the set of edges in monotonically increasing order by weight



# Example: Kruskal's Algorithm





## Implementation of Kruskal's Algorithm

- Uses a disjoint-set data structure to determine whether an edge connects vertices in different components
- Disjoint Set Operations
  - MAKE-SET(x): creates a new set whose only member (and thus representative) is x.
  - UNION(x, y): unites the dynamic sets that contain x and y, say S<sub>x</sub> and S<sub>y</sub>, into a new set that is the union of these two sets. We assume that the two sets are disjoint prior to the operation.
  - FIND-SET(x): returns a pointer to the representative of the (unique) set containing x.

## Kruskal's Algorithm



- 1. A ← Ø
- for each vertex v ∈ V
- do MAKE-SET(v)
- 4. sort E into non-decreasing order by w
- 5. for each (u, v) taken from the sorted list
- 6. do if FIND-SET(u) ≠ FIND-SET(v)
- 7. then  $A \leftarrow A \cup \{(u, v)\}$
- 8. UNION(u, v)
- 9. return A

