


# Graph Algorithms

CS3104

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Ford-Fulkerson - Method ( $G, S, t$ )

1. Initialize flow,  $f$  to Zero.

2. while there exist an augmenting path,  $P$  in  $G_f$   
do augment flow,  $f$  along  $P$ .

3. return  $f$ .

# The Ford-Fulkerson Algorithm

**FORD-FULKERSON**( $G, s, t$ )

1 **for** each edge  $(u, v) \in G.E$

2      $(u, v).f = 0$

3 **while** there exists a path  $p$  from  $s$  to  $t$  in the residual network  $G_f$

4      $c_f(p) = \min \{c_f(u, v) : (u, v) \text{ is in } p\}$

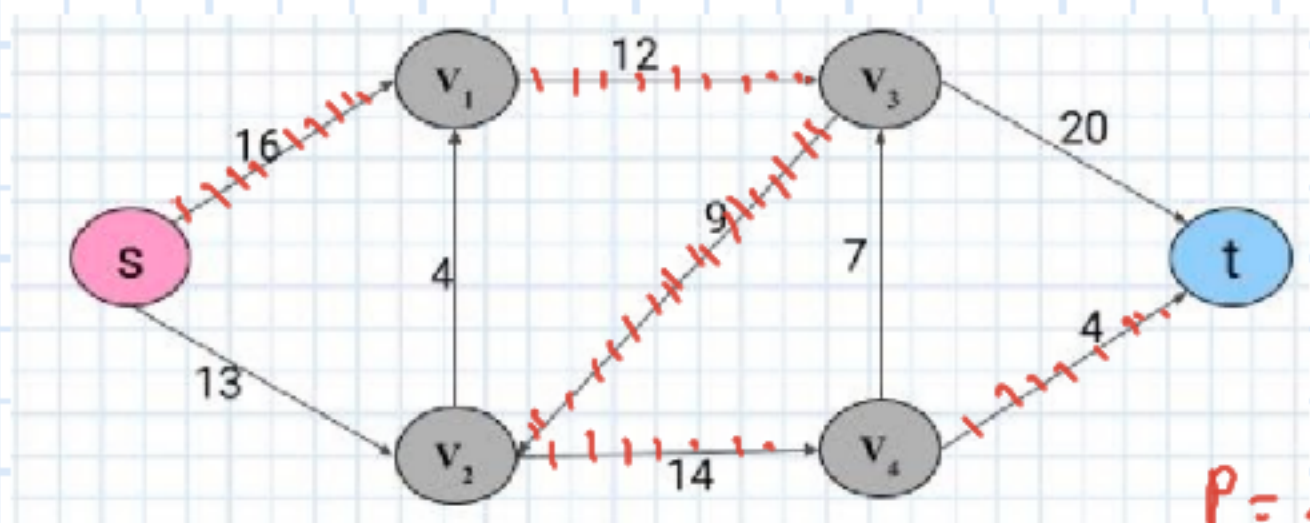
5     **for** each edge  $(u, v)$  in  $p$

6         **if**  $(u, v) \in E$

7              $(u, v).f = (u, v).f + c_f(p)$

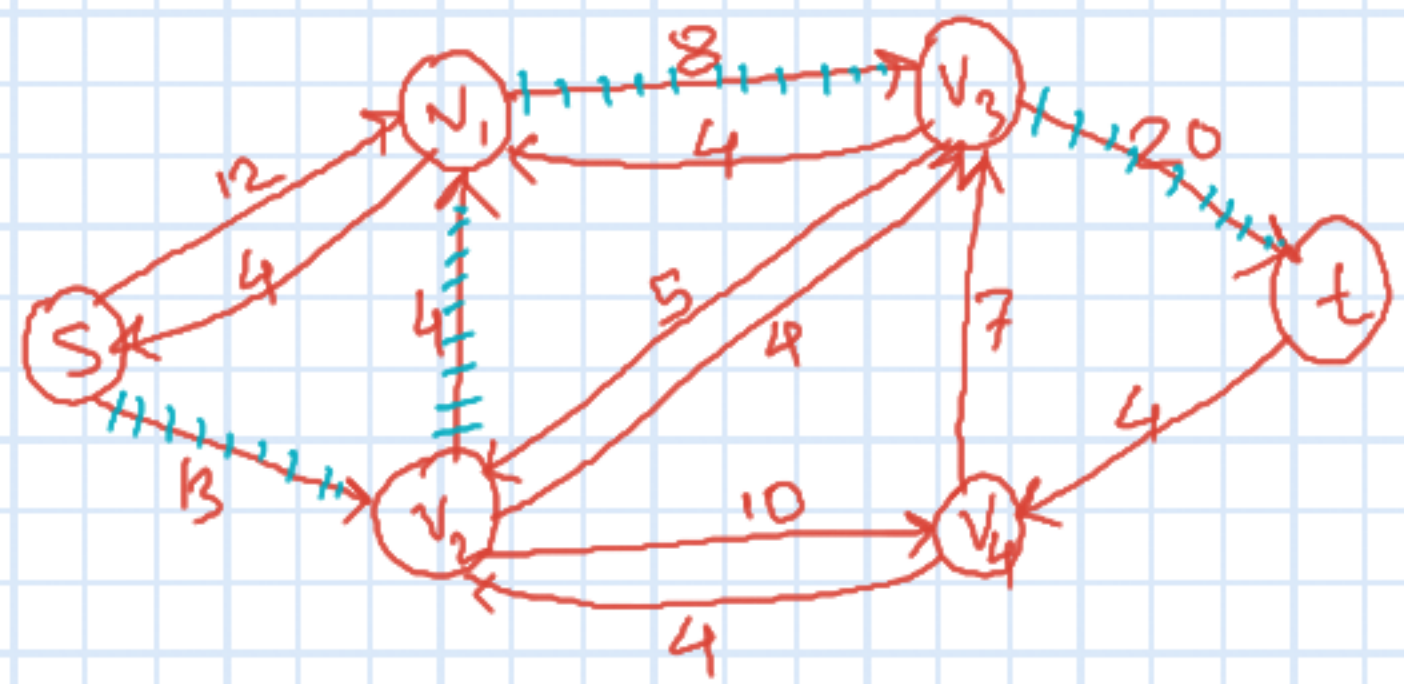
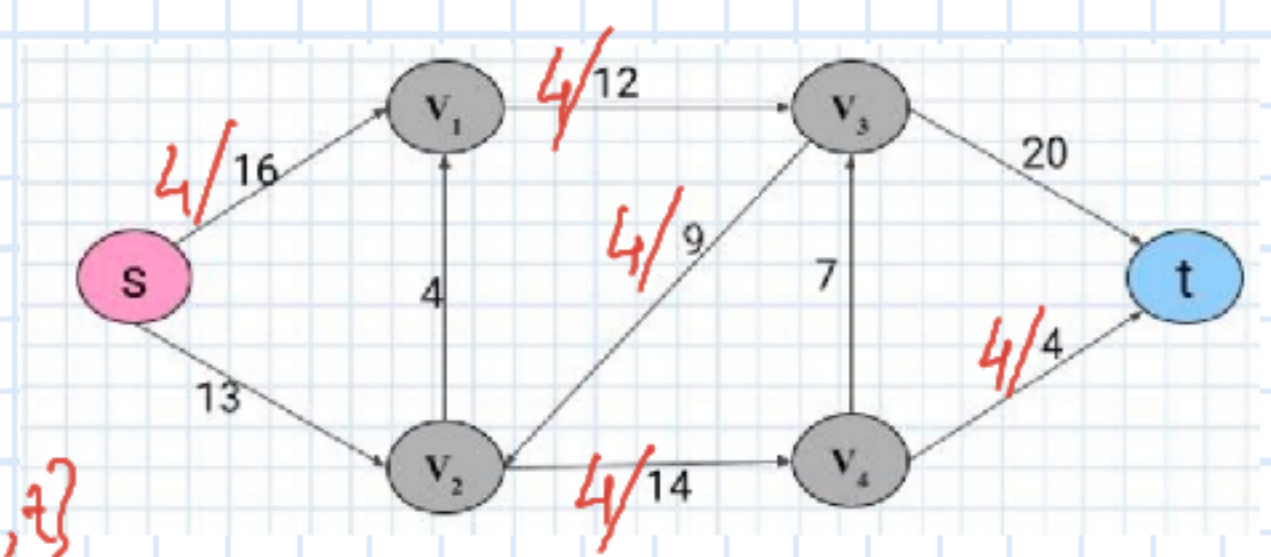
8         **else**  $(v, u).f = (v, u).f - c_f(p)$



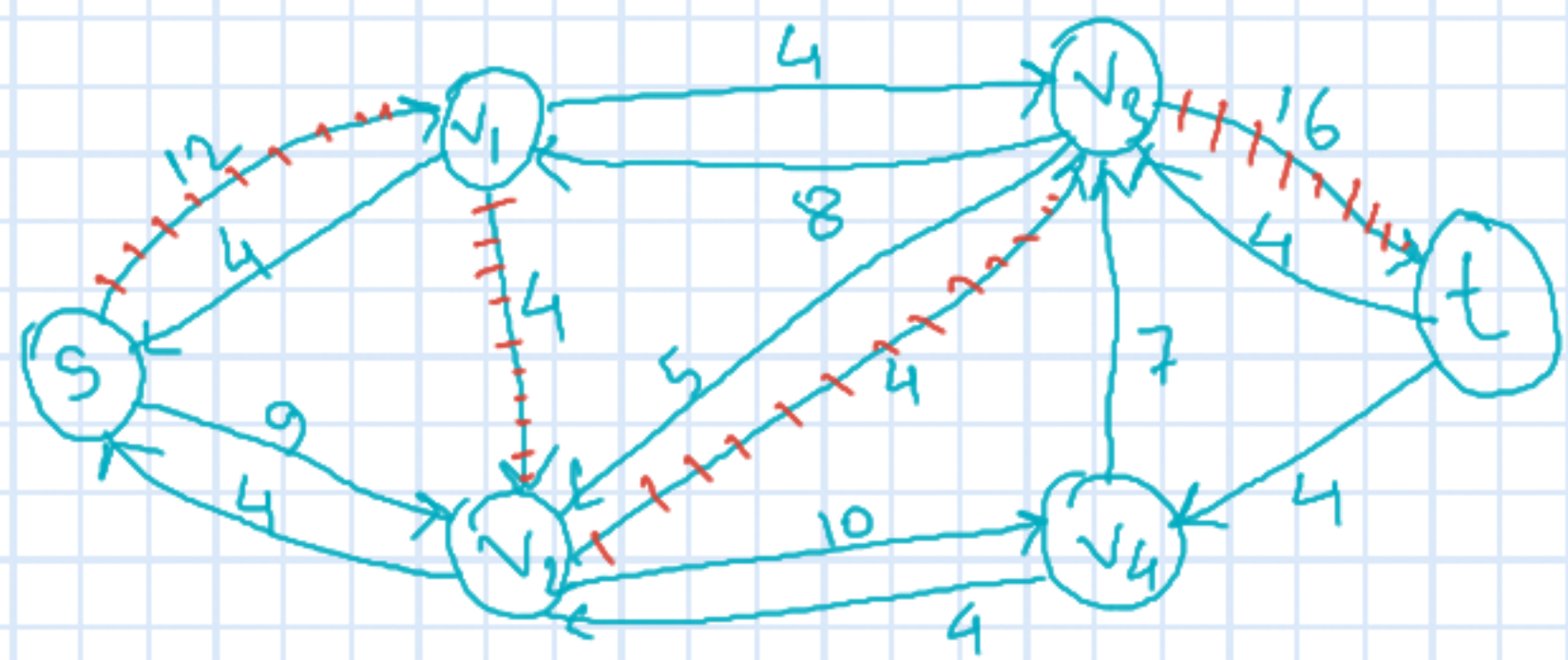
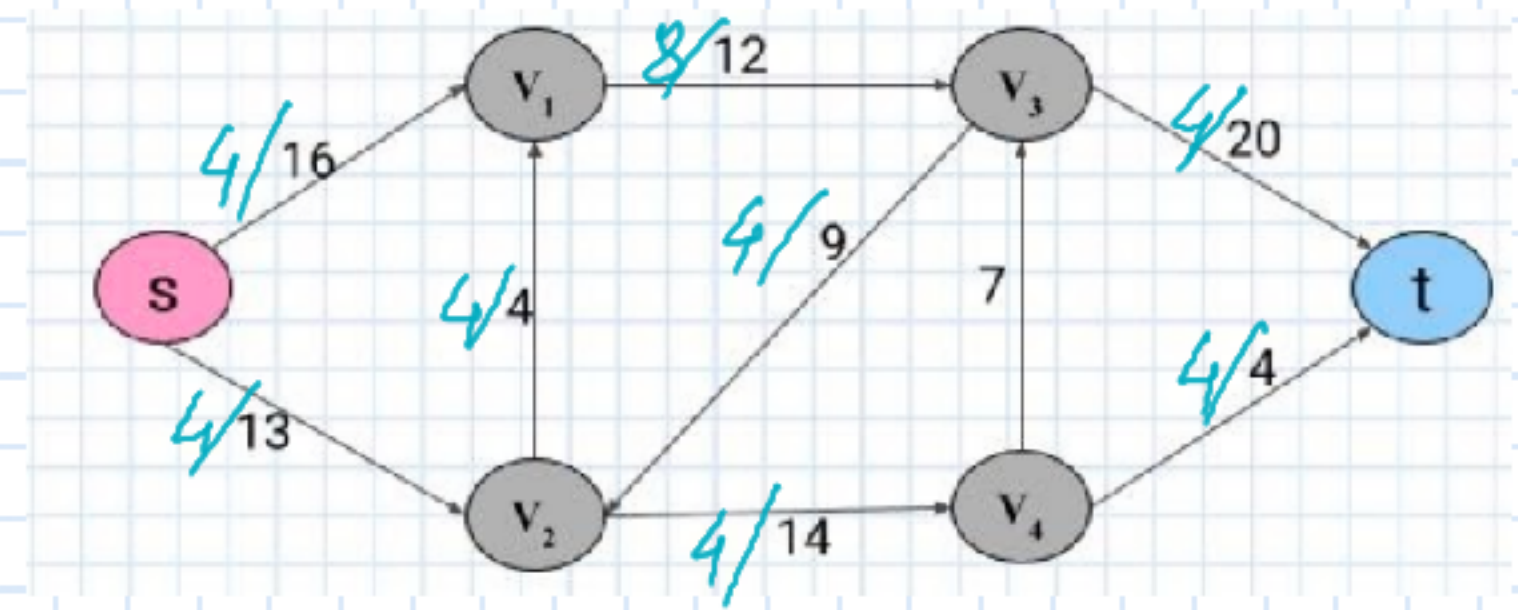


$$c_f(p) = 4$$

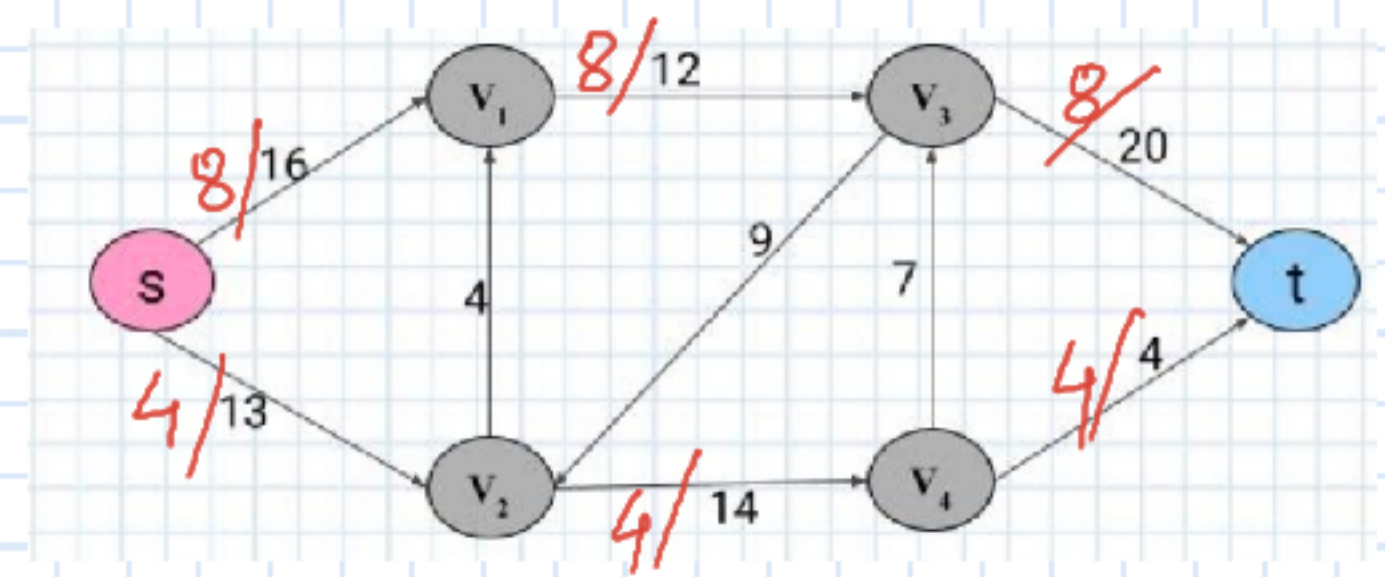
$$P = \{s, v_1, v_3, v_2, v_4, t\}$$



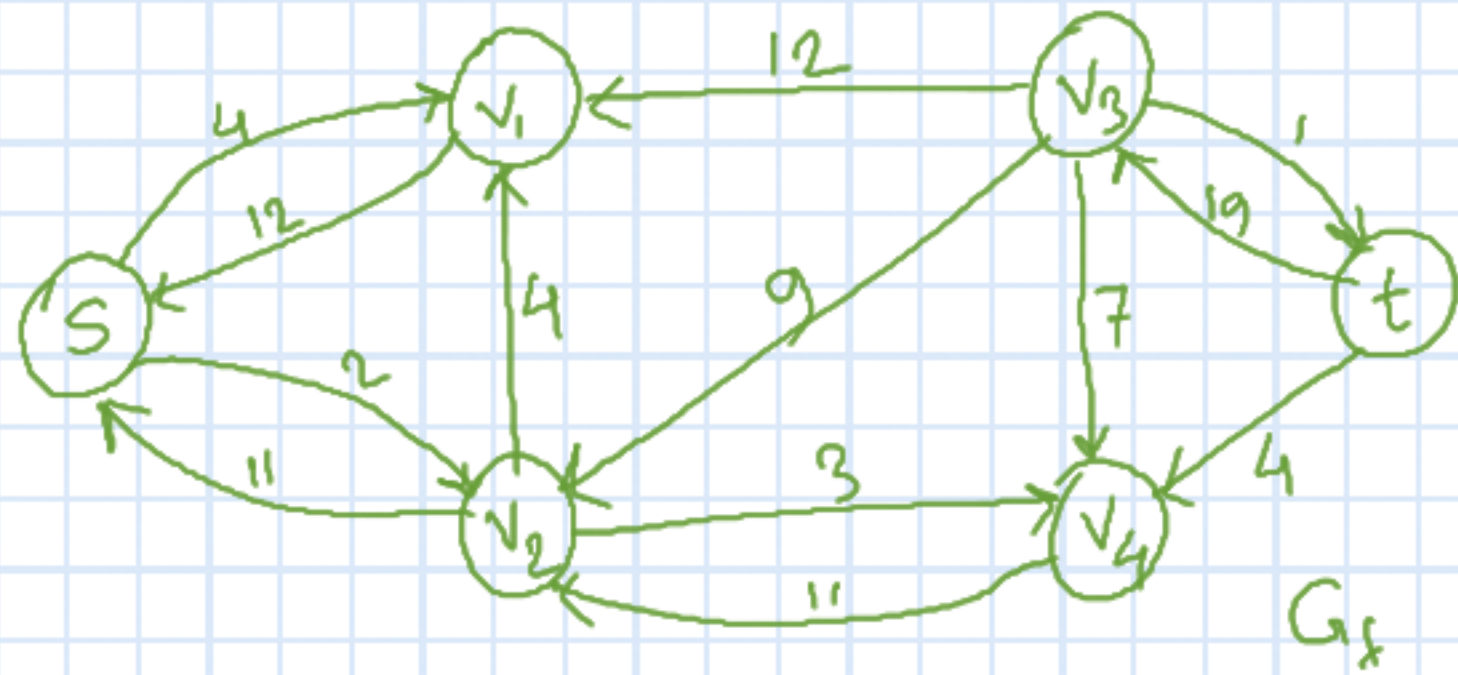
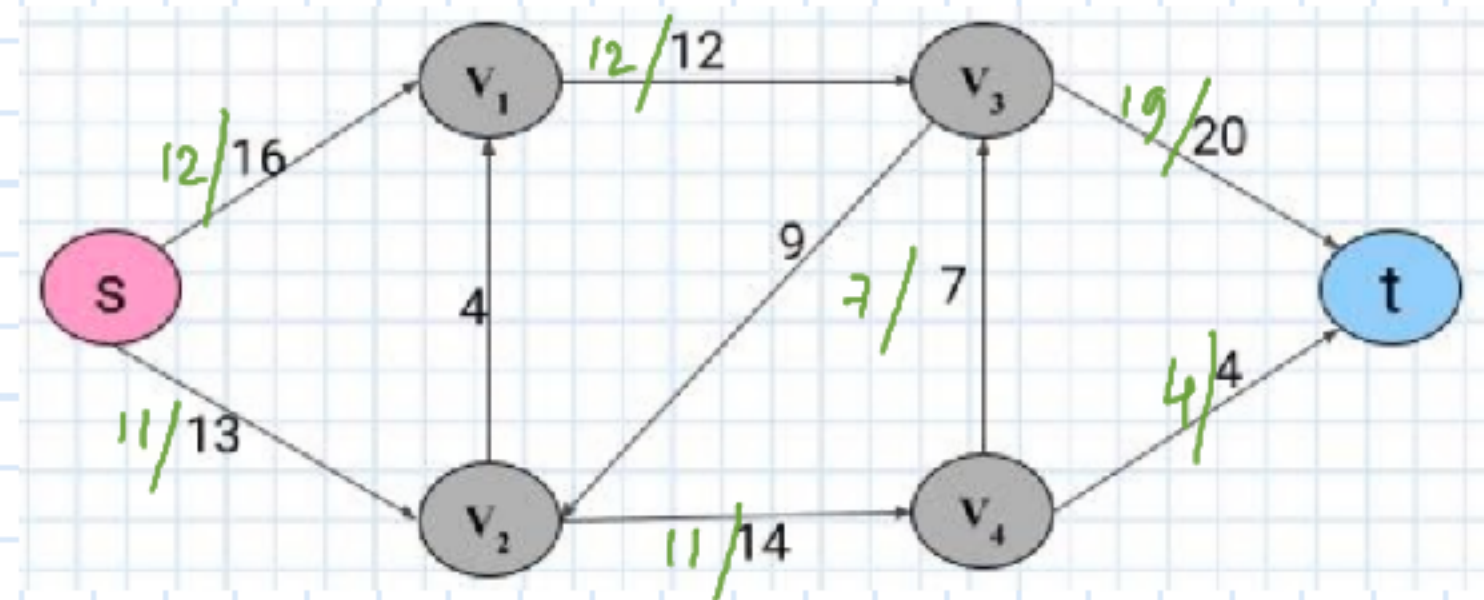
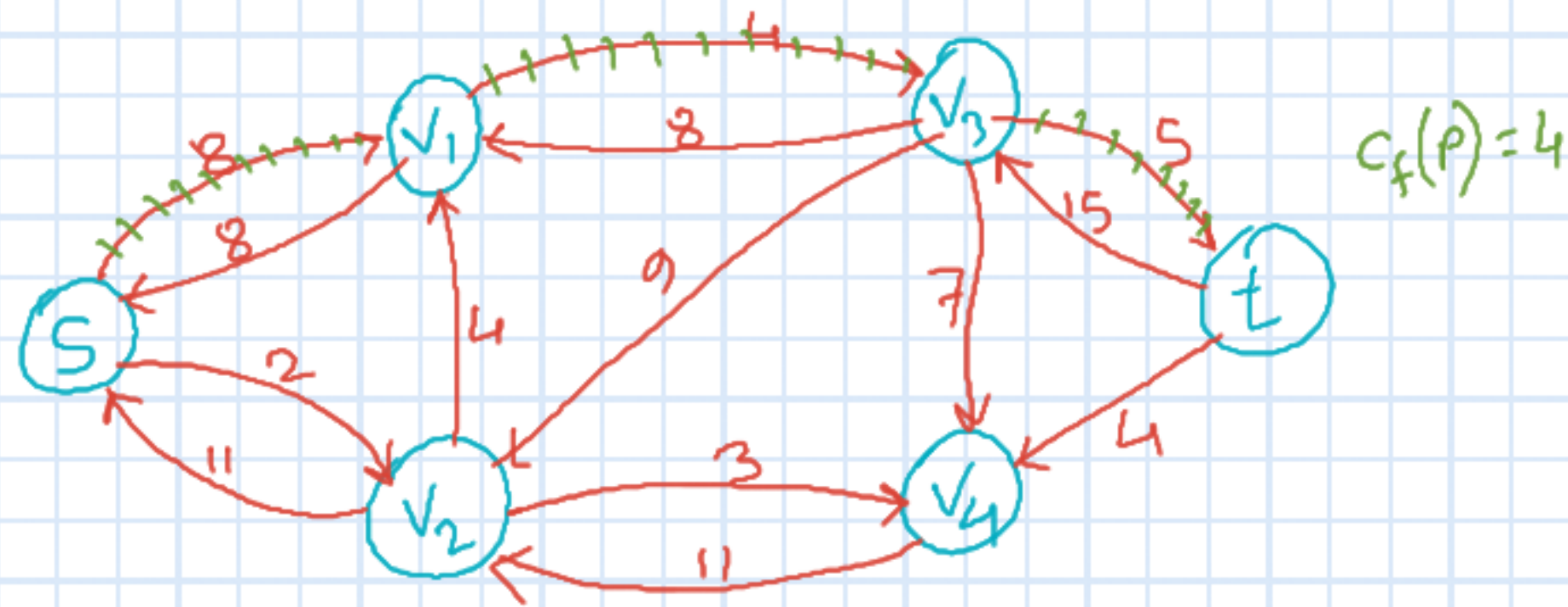
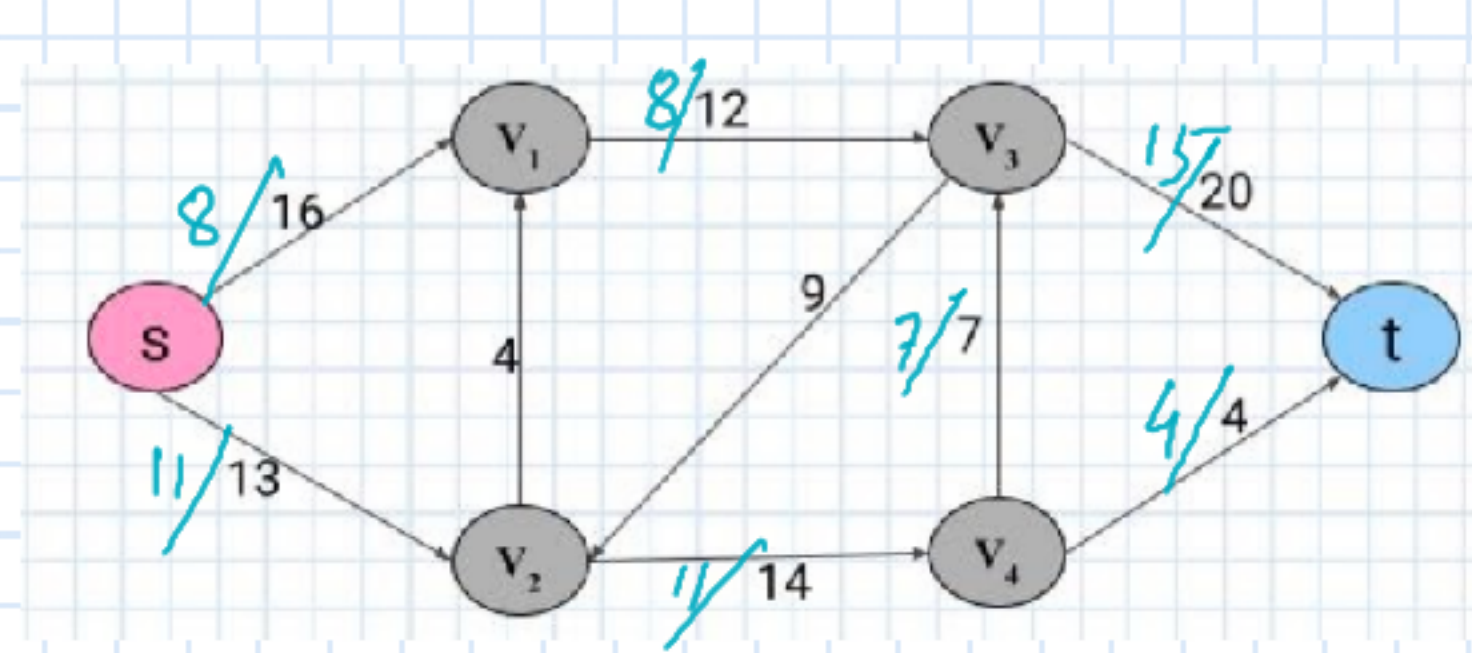
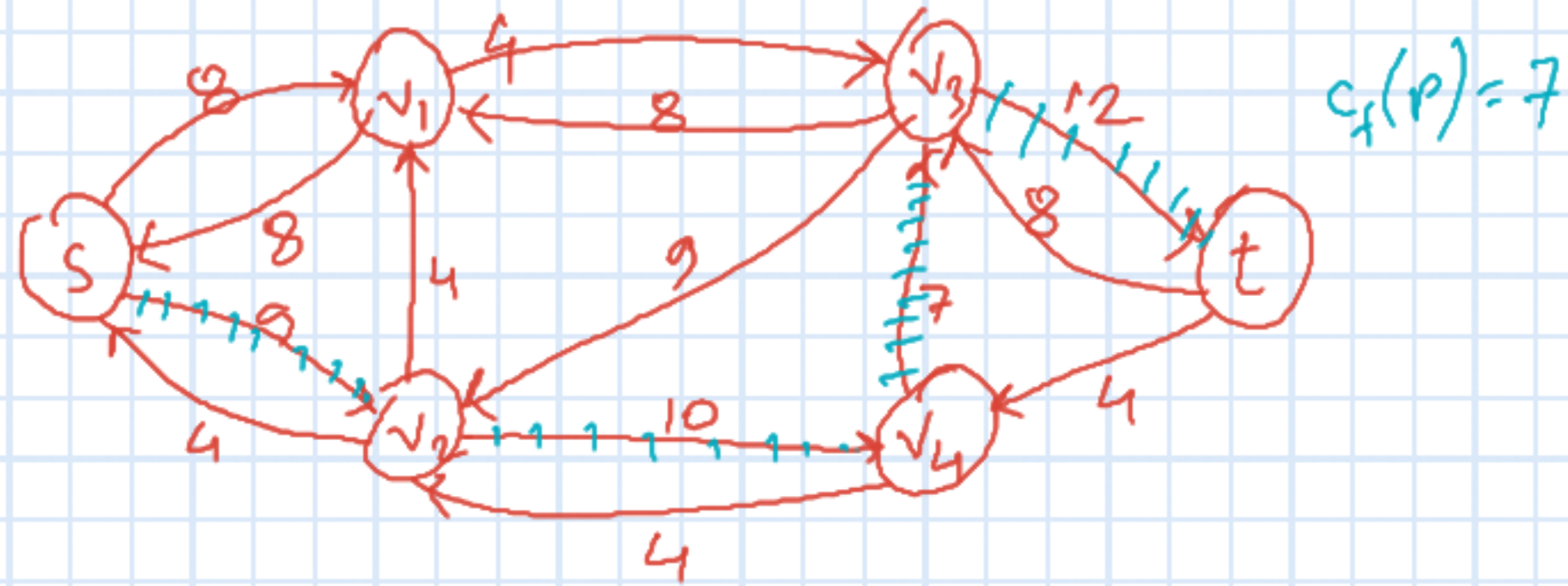
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After this in  $G_f$  no augmenting path is possible.

## Cuts of the flow Network.

Cut  $(S, t)$  of flow network  $G = (V, E)$  is a partition of vertices,  $V$  into  $S$  &  $T = V - S$

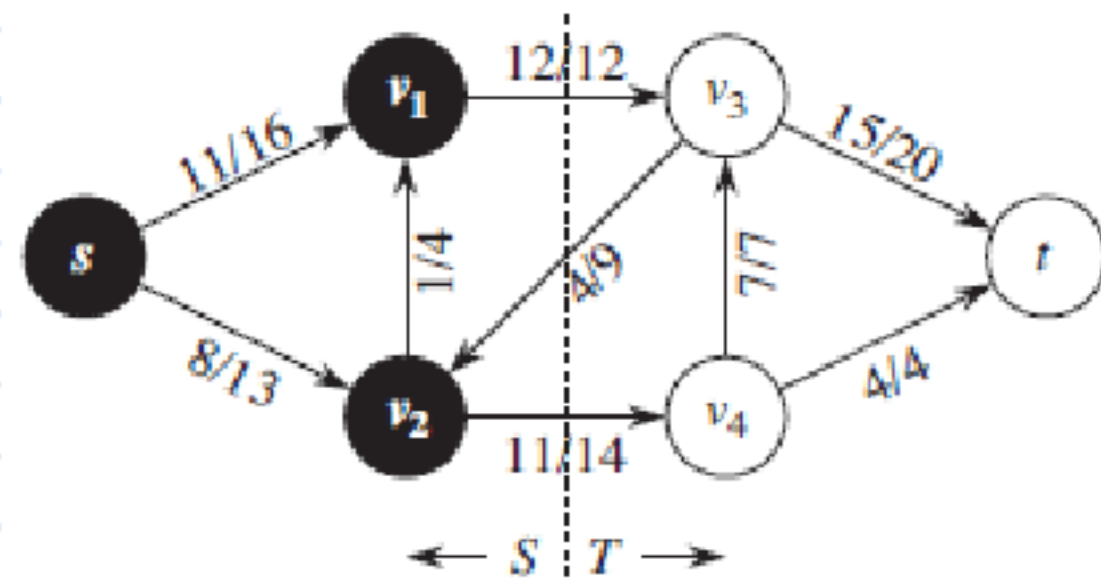
$$s \in S, t \in T$$

$$f(S, T) = \sum_{u \in S} \sum_{v \in T} f(u, v) - \sum_{u \in S} \sum_{v \in T} f(v, u)$$

Capacity of the cut  $(S, T)$  is

$$c(S, T) = \sum_{u \in S} \sum_{v \in T} c(u, v)$$

minimum Cut  $\Rightarrow$  min over all cuts of the network



$$\text{Cut}(\{s, v_1, v_2\}, \{v_3, v_4, t\})$$

Net flow

$$f(v_1, v_3) + f(v_2, v_4) - f(v_3, v_2) = 12 + 11 - 4 = 19$$

Capacity of this cut

$$c(v_1, v_3) + c(v_2, v_4) = 12 + 14 = 26.$$



Flow across any cut is the same and equals the value of the flow

## Max-Flow min-cut Theorem.

Let  $f$  is a flow in network  $G$ . Then the following cond<sup>s</sup> are equivalent

(a)  $f$  is a maximum flow in  $G$ .

(b) Residual network  $G_f$  contains no augmenting path.

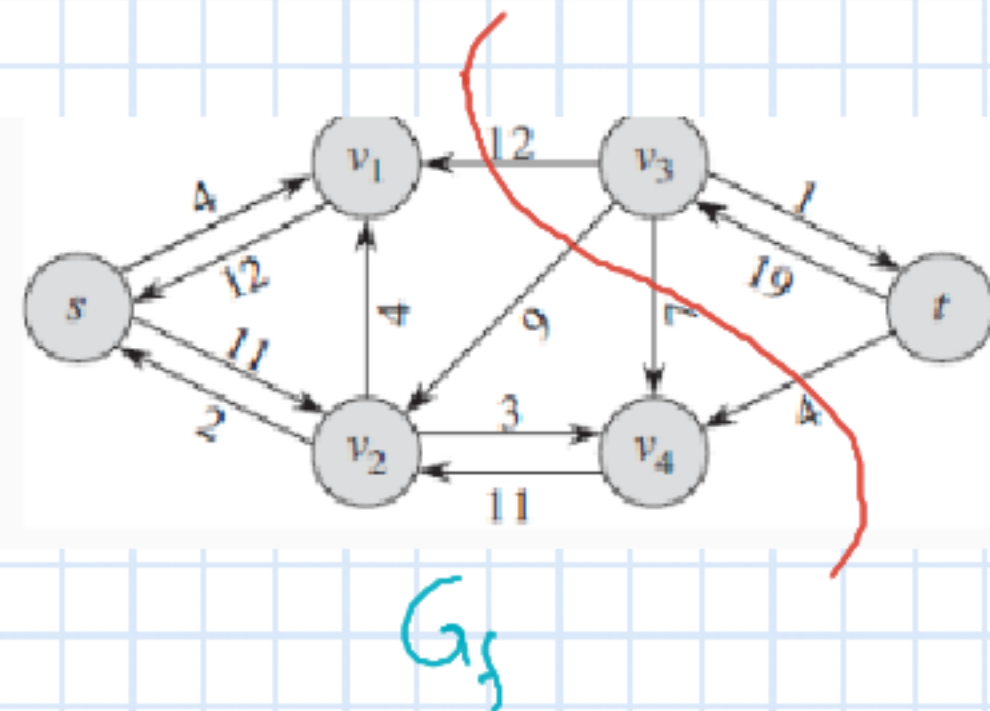
(c) there is a cut  $(S, T)$  with  $|f| = C(S, T)$

maximum flow = capacity of minimum cut

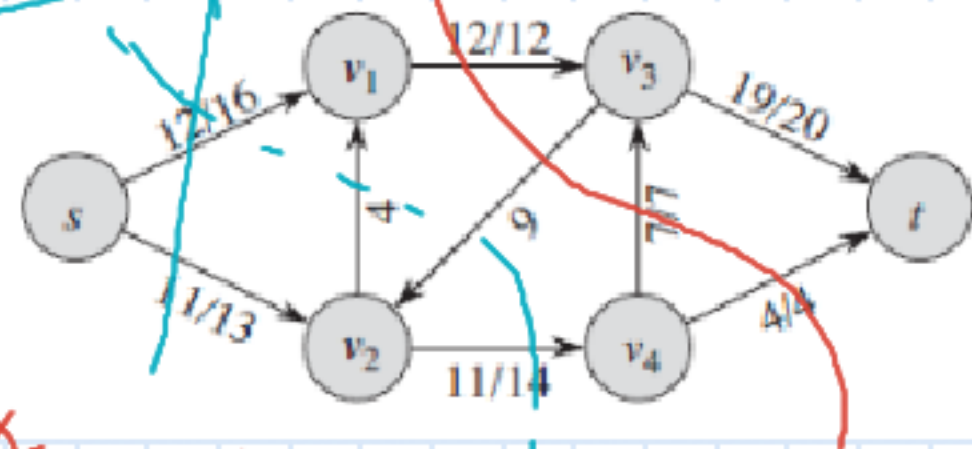
$$S = \{s, v_1, v_2, v_4\}$$

$$T = \{v_3, t\}$$

$$C(S, T)$$



Augmented flow.



Cut capacity  
 $C(S, T) = 23$



