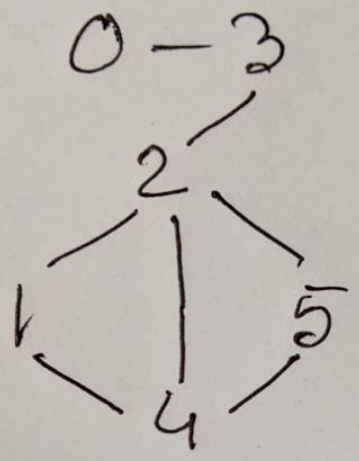
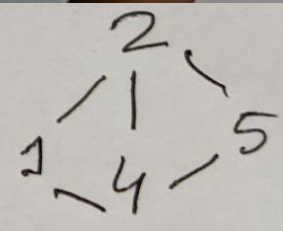


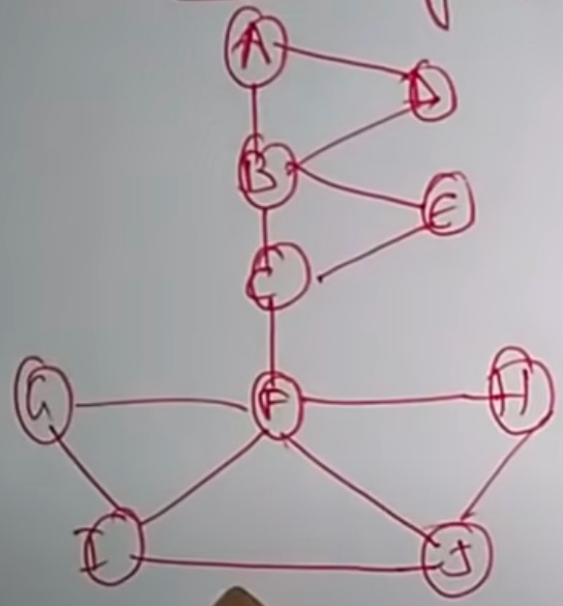
Bcc:



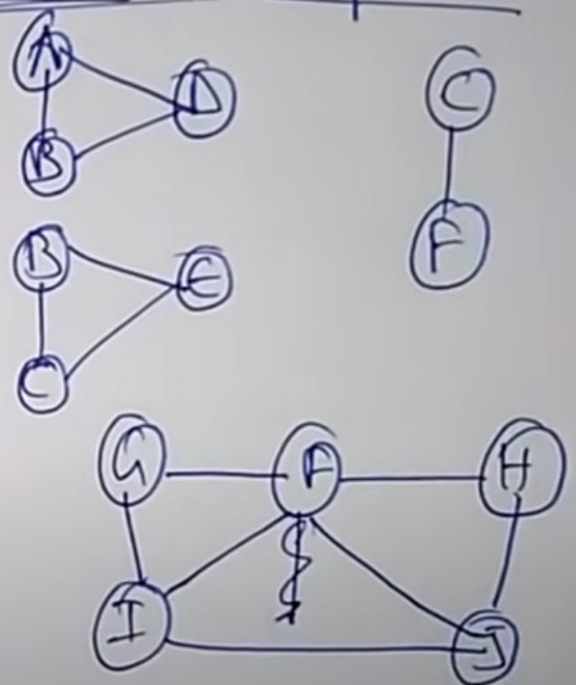
0-3      3-2

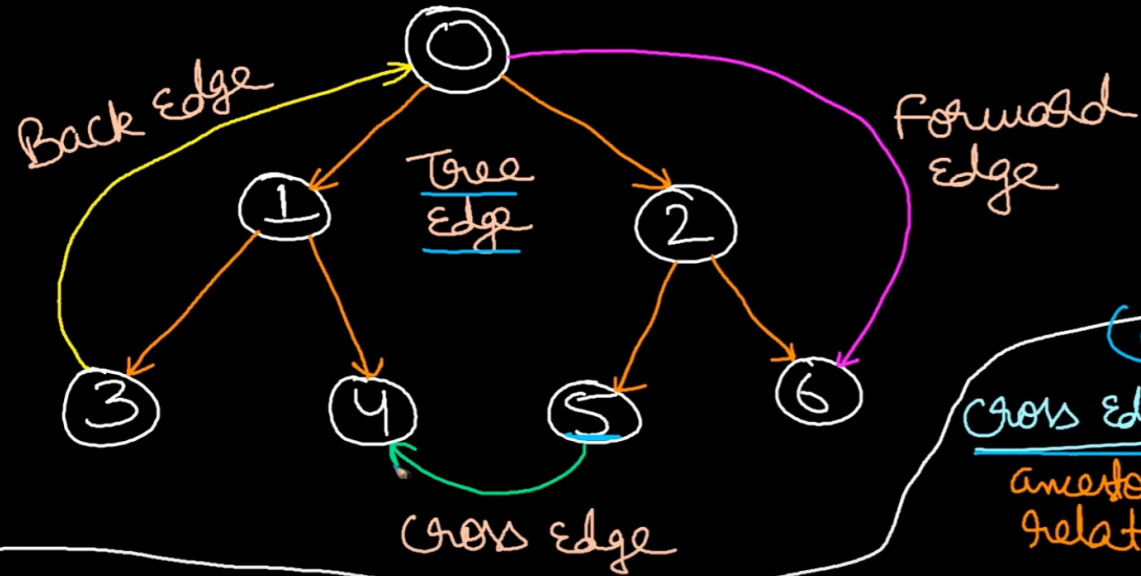


Undirected graph



biconnected components





Cross edge: don't have ancestor - descendent relationship.

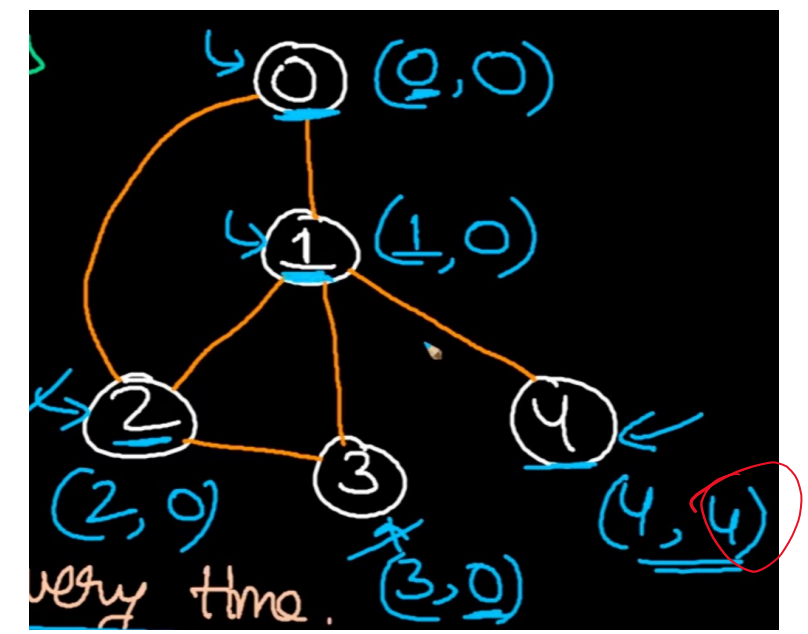
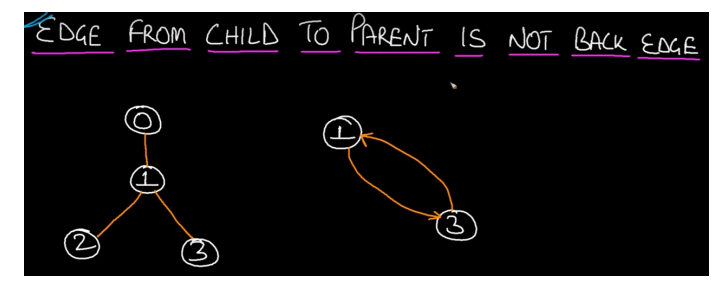
Tree Edge:  $(U, V)$   
 Parent (U) → child (V)  
 Ancestor (U) → descendent (V)

Forward Edge:  $(U, V)$   
 Ancestor (U) → descendent (V)

Back Edge:  $(U, V)$   
 Descendent (U) → Ancestor (V)

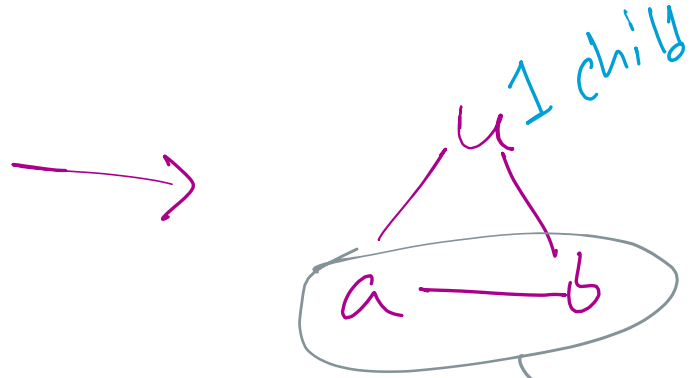
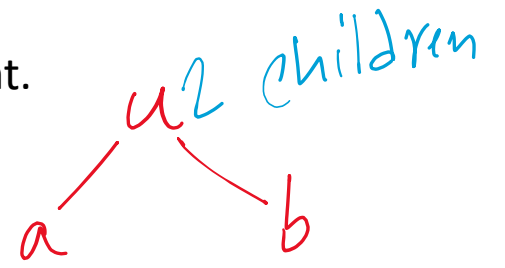
Low value is also discovery time .curr node theke accessible shokol node er moddhe je node er disc time min

For directed graph:

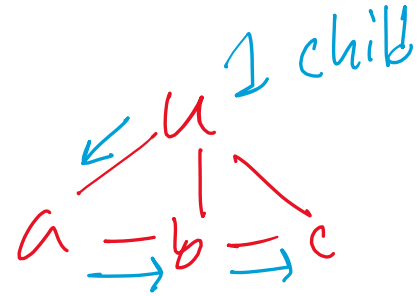


Case 1:

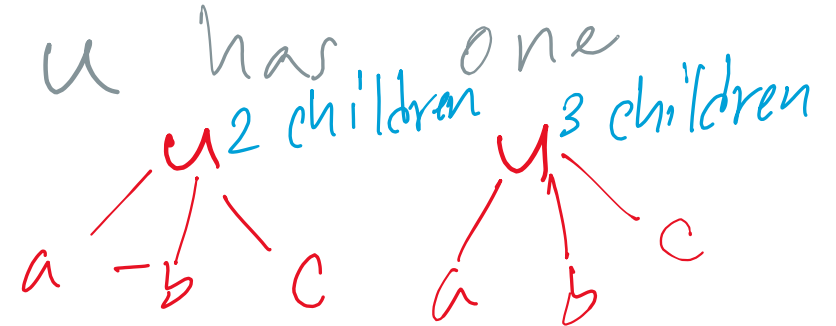
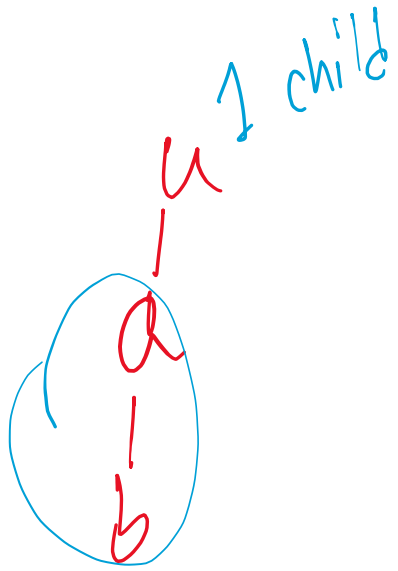
If  $u$  is the root node of DFS tree and has at least two children ( $a, b$ ) then it is an articulation point.



This is the only subgraph of  $a$  here.

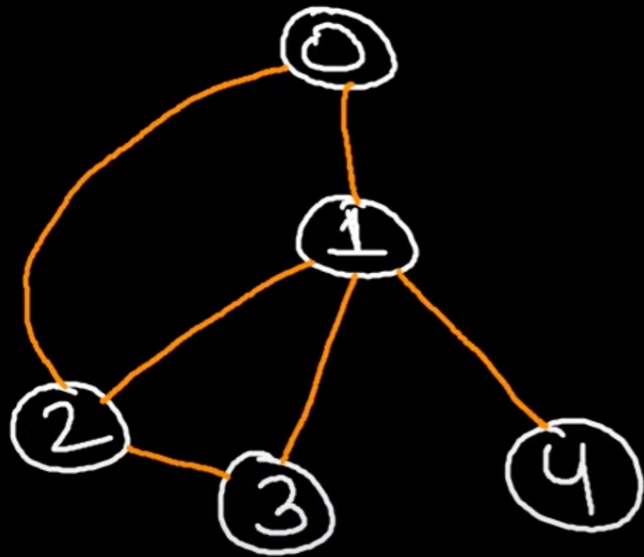


According to DFS, after visiting  $a$ ,  $b$  will be visited.  $b$  will not visit as the child of  $a$ .  $a-b$  will be visited as the subgraph of  $u$ .  $u$  has one child in the DFS tree.

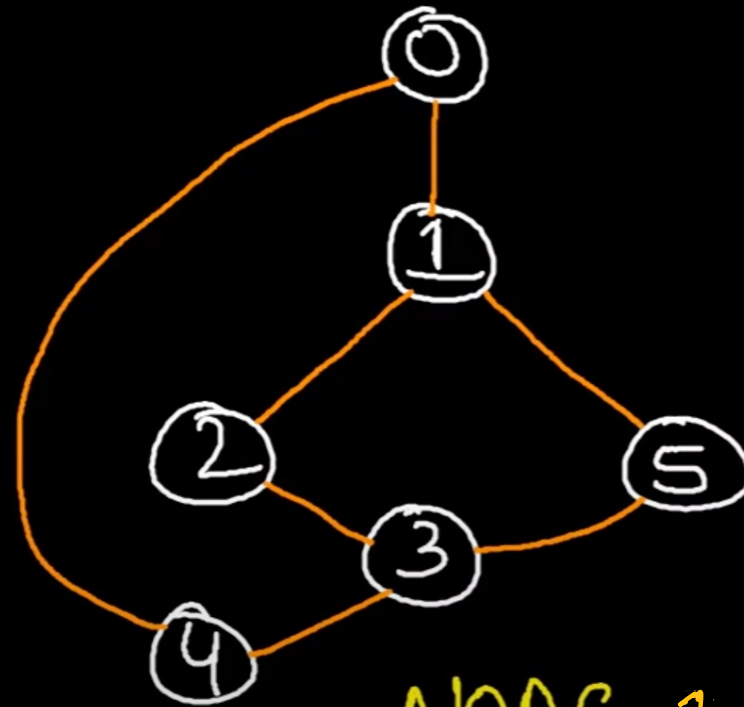


Case 2

{ 2) Node  $U$  is not root of DFS tree and it has a child  $v$  such that no vertex in subtree rooted with  $v$  has a back edge to one of the ancestors of  $U$ .



Node-1 is AP

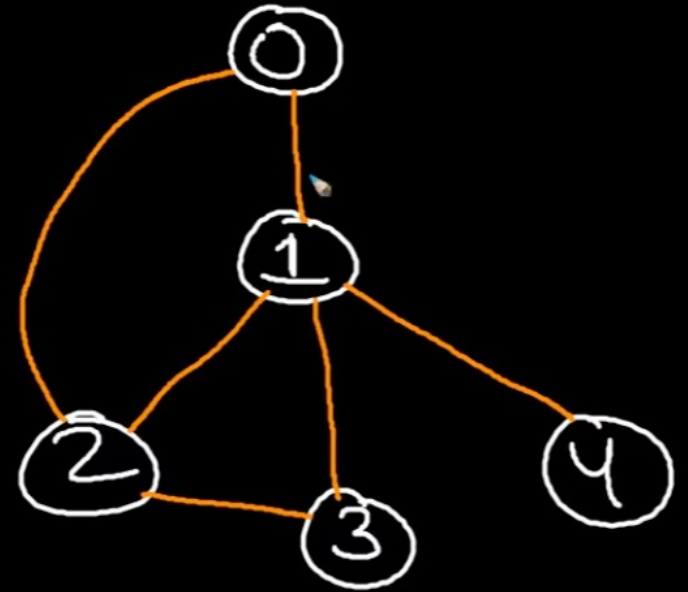


NODE-1 is NOT AP



## HOW TO DETECT AP FOR CASE-2

① we need to find the order of vertices from earliest to latest to detect back-edges.



∴ we use timeStamp to mark Nodes with increasing value.

∴ we can do this by assigning discovery time.

② we need to maintain the earliest possible node accessible for a given node which will indicate if we have any back edge.

For that we will assign Low value to each node.









