

# under Graduate Homework In Mathematics

## Set Theory 4

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General fire extinguisher

**PROBLEM I** Consider  $\mathbb{Q} = \mathbb{Z} \times (\mathbb{Z} \setminus \{0\}) / \sim$ , where  $(a, b) \sim (c, d) \iff ad = bc$ . Define  $+_{\mathbb{Q}}, \cdot_{\mathbb{Q}}$  and  $<_{\mathbb{Q}}$  and verify that your definitions doesn't depend on the choice of representatives.

**SOLUTION**. Let  $[(a, b)] +_{\mathbb{Q}} [(c, d)] = [(ad + bc, bd)]$ ,  $[(a, b)] \cdot_{\mathbb{Q}} [(c, d)] = [(ac, bd)]$ , and  $[(a, b)] <_{\mathbb{Q}} [(c, d)] \iff abd^2 < cdb^2$ . Now we prove they are well-defined, i.e., doesn't depend on the choice of representatives.

For  $+_{\mathbb{Q}}$ , assume  $(a, b) \sim (e, f)$ , we need to prove  $(ad + bc, bd) \sim (ed + fc, df)$ . Since  $af = be$ , we have  $(ad + bc)bf = ad^2f + bdcf = bed^2 + bdcf = (ed + fc)bd$ . So  $+_{\mathbb{Q}}$  is well defined.

For  $\cdot_{\mathbb{Q}}$ , assume  $(a, b) \sim (e, f)$ , we need to prove  $(ac, bd) \sim (ec, fd)$ . Since  $af = be$ , we have  $acfd = bced = bdec$ . So  $\cdot_{\mathbb{Q}}$  is well defined.

For  $<_{\mathbb{Q}}$ , assume  $(a_1, b_1) \sim (a_2, b_2)$ ,  $(c_1, d_1) \sim (c_2, d_2)$  and  $(a_1, b_1) < (c_1, d_1)$ . Now we need to prove  $(a_2, b_2) < (c_2, d_2)$ . Since  $a_1b_2 = a_2b_1$ ,  $c_1d_2 = c_2d_1$  we get  $a_1b_1d_2^2 < c_2d_2b_1^2$   $\square$

**PROBLEM II** The set of all continuous functions  $f : \mathbb{R} \rightarrow \mathbb{R}$  has cardinality  $\mathfrak{c}$  (while the set of all functions has cardinality  $2^{\mathfrak{c}}$ ). [A continuous function on  $\mathbb{R}$  is determined by its values at rational points.]

**SOLUTION**. Consider  $\theta : {}^{\mathbb{R}}\mathbb{R} \rightarrow 2^{\mathbb{Q}}, f \mapsto \{(a, b) \in \mathbb{Q} : f(a) < b\}$ . Now we prove  $f$  is a injection. Assume  $\theta(f) = \theta(g)$ , to prove  $f = g$ . First we prove for  $x \in \mathbb{Q}$  we have  $f(x) = g(x)$ . We have  $f(x) = \sup\{y \in \mathbb{Q} : y < f(x)\} = \sup\{y \in \mathbb{Q} : (x, y) \in \theta(f)\} = \sup\{y \in \mathbb{Q} : (x, y) \in \theta(g)\} = g(x)$ . For  $x \in \mathbb{R}$ , choose a sequence  $x_n \in \mathbb{Q}$  such that  $x_n \rightarrow x$ , then  $f(x) = \lim_{n \rightarrow \infty} f(x_n) = \lim_{n \rightarrow \infty} g(x_n) = g(x)$ . So we get  $f = g$ . So  $\text{card}^{\mathbb{R}}\mathbb{R} \leq \text{card}2^{\mathbb{Q}} = 2^{\aleph_0}$ . Obviously  $\text{card}^{\mathbb{R}}\mathbb{R} \geq 2^{\aleph_0}$ , so we get they are equal.  $\square$

**PROBLEM III** There are at least  $\mathfrak{c}$  countable order-types of linearly ordered sets.

**SOLUTION**. For every sequence  $a = \langle a_n : n \in \mathbb{N} \rangle$  of natural numbers consider the ordertype

$$\tau_a = \{(x, y) \in \mathbb{Z} \times \mathbb{N} : 2 \nmid y \wedge 0 < x < a_{\frac{y}{2}}\}$$

And for  $(x, y), (z, w) \in \tau_a$  we define  $(x, y) < (z, w) \iff y < w \wedge y = w, x < z$ . Now we will show that if  $a \neq b$ , then  $\tau_a \neq \tau_b$ . Assume  $\tau_a \cong \tau_b$ , we need to prove  $a = b$ . assume  $\theta : \tau_a \rightarrow \tau_b$  is the isomorphism.

We know  $(x, 0)$  can be defined as  $\phi(p) = \exists_{k=1}^{x-1} t_k, \wedge_{1 \leq i < j \leq x-1} t_i \neq t_j, \forall k = 1, \dots, x-1, t_k < p$ . And  $\theta$  is isomorphism. So  $\theta(x, 0) = (x, 0)$ . For  $(x, 1)$ , we let  $b_0$  satisfy  $\theta(0, 1) = (b_0, m)$ . Since the set  $\{(x, y) : y = 1\}$  can be defined by  $\psi(p) = \forall r, s(r, s < p \wedge \tau(r) \wedge \tau(s) \rightarrow \text{card}[r, s] < \infty)$ , where  $\tau(r) := \{s : s < r\}$  and  $[r, s] = \{y : r < y < s\}$ . we get  $\theta[\{(x, y) : y = 1\}] = \{(x, y) : y = 1\}$ . So we can delete the element whose second coordinary is 0, 1, and  $\theta$  is isomorphism, too. Do this repeatedly, we get  $\theta(x, 2n+1) = (x, 2n+1)$ . So  $a_n = \text{card}\{(x, 2n+1) \in \tau_a\} = \text{card}\{(x, 2n+1) \in \tau_b\} = b_n$  and thus  $a = b$ .  $\square$

**PROBLEM IV** The set of all algebraic reals is countable.

**SOLUTION**. Assume  $\{f_n : n \in \mathbb{N}\}$  is the set of all integral coefficient polynomial. Consider  $A_n := \{x \in \mathbb{C} : f_n(x) = 0\}$  is finite set. Then we get  $\bigcup_{n \in \mathbb{N}} A_n$  is at most countable. Obviously  $\bigcup_{n \in \mathbb{N}} A_n$  is infinite, so it's countable.  $\square$

**PROBLEM V** If  $S$  is a countable set of reals, then  $|\mathbb{R} - S| = \mathfrak{c}$ . [Use  $\mathbb{R} \times \mathbb{R}$  rather than  $\mathbb{R}$  (because  $|\mathbb{R} \times \mathbb{R}| = 2^{\aleph_0}$ ).]

**SOLUTION.** Assume  $\theta : \mathbb{R} \rightarrow \mathbb{R} \times \mathbb{R}$  is bijection, and  $T = \theta(S)$ . Then  $T$  is countable. And  $\text{card}(\mathbb{R} \setminus S) = \text{card}(\mathbb{R} \times \mathbb{R} \setminus T)$ . So we only need to prove  $\mathbb{R} \times \mathbb{R} \approx \mathbb{R} \times \mathbb{R} \setminus T$ . Obviously  $\text{card} \mathbb{R} \times \mathbb{R} \setminus T \leq \text{card} \mathbb{R} \times \mathbb{R}$ , so we only need  $\mathbb{R} \times \mathbb{R} \setminus T \geq \mathbb{R}$ . Since  $T$  is countable, we get  $\{x : \exists y, (x, y) \in T\}$  is countable. Choose  $t \notin \{x : \exists y, (x, y) \in T\}$ . Let  $f : \mathbb{R} \rightarrow \mathbb{R} \times \mathbb{R} \setminus T, x \mapsto (t, x)$ . Easily we get  $f$  is injection. So  $\text{card} \mathbb{R} \times \mathbb{R} \setminus T = \mathfrak{c}$ .  $\square$

**PROBLEM VI** Assume  $T$  is a tree.

1. If  $s, t, u \in T$ , then  $R_{stu} := \{\delta_{st}, \delta_{tu}, \delta_{us}\}$  has at most 2 elements. And if  $p, q \in R_{stu}$ , then  $p \subset q \vee q \subset p$ .
2.  $\prec$  is a linear ordering of  $T$  which extends  $\sqsubset$ .
3. For every  $t \in T$ , Prove  $T^t := \{s \in T : t \sqsubset s\}$  is an interval in  $(T, \prec)$ .

**SOLUTION.** 1. First we prove for  $p, q \in R_{stu}$  we have  $p \subset q \vee q \subset p$ . Without loss of generality assume  $p = \delta_{st}, q = \delta_{tu}$ . We have  $p, q \subset (\cdot, t)$ . Since  $(\cdot, t)$  is well ordered, and easily  $p, q$  are initial segment, so  $p \subset q \vee q \subset p$ .

Now we prove there are at most two elements. From above we know  $(R_{stu}, \subset)$  is linear order set, and it's finite. Without loss of generality we assume  $\delta_{st} \subset \delta_{tu} \subset \delta_{us}$ . Then we get  $\delta_{tu} = \delta_{tu} \cap \delta_{us} = (\cdot, t) \cap (\cdot, u) \cap (\cdot, s) \subset \delta_{st}$ . That means  $\delta_{st} = \delta_{tu}$ , so there is at most two elements.

2. Easily to prove  $\sqsubset \subset \prec$ . Now we prove  $\prec$  is linear ordered. Consider a bigger linear ordered set  $Y$  is obtained by adding a minimum,  $-\infty$ , in  $X$ . Consider the tree  $U := {}^{<\alpha}Y$ . We try to

make a map from  $T$  to  $B_U$ . Let  $\theta : T \rightarrow B_U, \theta(f)(\beta) := \begin{cases} f(\beta), \beta \in \text{dom } f \\ -\infty, \beta \notin \text{dom } f \end{cases} \quad \forall \beta \in \alpha, f \in T$ .

Then we it's easily to prove  $\theta$  is injective and  $f \prec g \iff \theta(f)(\beta) < \theta(g)(\beta)$ , where  $\beta = \min\{t \in \alpha : \theta(f)(t) \neq \theta(g)(t)\}$ . We define  $f, g \in B_U, f < g \iff f(\beta) < g(\beta)$ , where  $\beta = \min\{t \in \alpha : f(t) \neq g(t)\}$ . Now we only need to prove  $(B_U, <)$  is linear ordered. Easily  $f \not\prec f, \forall f \in B_U$ . And for  $f \neq g, f < g \vee g < f$ . Assume  $f < g < h$ , to prove  $f < h$ .

If  $n_{fg} < n_{gh}$  then we get  $f(n_{fg}) < g(n_{fg}) = h(n_{fg})$ . So  $n_{fh} \leq n_{fg}$ . From VI.1 we get  $n_{fh} = n_{fg} \vee n_{fh} = n_{gh}$ . So  $n_{fh} = n_{fg}$ , and thus  $f < h$ .

If  $n_{fg} > n_{gh}$ , then we get  $h(n_{gh}) > g(n_{gh}) = f(n_{gh})$ . Same as above we get  $n_{fh} = n_{gh}$ , so  $f < h$ .

If  $n_{fg} = n_{gh}$ , it's obvious  $f < h$ .

So we have proved  $B_U$  is linear ordered, and thus  $(T, \prec)$  is linear orderd.

3. Only need to prove if  $t \sqsubset u, t \sqsubset v, u \prec v$ , then  $\forall s : u \prec s \prec v, t \sqsubset s$ . If  $u \sqsubseteq s$  then  $t \sqsubset u \sqsubseteq s$ . Else we get  $u \not\sqsubseteq s$ . So we get  $s \not\sqsubseteq u \wedge s(n_{su}) > u(n_{su})$ . From VI.2 we get  $t \prec s$ . So if  $t \not\sqsubseteq s$

then  $s \not\sqsubseteq t \wedge s(n_{st}) > t(n_{st})$ . Since  $t \sqsubseteq v$  we get  $s(n_{st}) > t(n_{st}) = v(n_{st})$ . Since  $t \sqsubset v$  we get  $n_{st} = n_{sv}$ , so  $v \prec s$ , contradiction! So  $t \sqsubset s$ .

□

### PROBLEM VII

1. Prove that  $\prec$  is linear ordered on  $T \cup B_T$ .
2. For every  $t \in T$ , prove that  $B_t = \{f \in B_T : t \in f\} \cup \{f \in T : t \sqsubset f\}$  is interval in  $(T \cup B_T, \prec)$ .

*SOLUTION.* 1. consider a bigger tree  $U$ . For  $f \in T$  let  $\theta(f) = f$ , for  $f \in B_T$  we let  $\theta(f)$  is a map from  $ordertype(\text{dom}(f))$  to  $X$ , and  $\theta(f)(\beta) := g(\beta)$ , where  $g \in f$  and  $\beta \in \text{dom}(g)$ . Let  $U = \theta(T \cup B_T)$ . Then easily  $T \subset U$ . Now we prove  $\theta$  is isomorphic from  $(T \cup B_T, \prec)$  to  $(U, \prec)$ . Easily for  $f, g \in T$  we have  $f \sqsubset g \iff \theta(f) \sqsubset \theta(g)$ .

And for  $f \in T, g \in B_T$  we have  $f \in g \iff \theta(f) \sqsubset \theta(g)$ . So from the definition of  $\prec$  we get  $\theta$  is isomorphic. Since we have proved  $(U, \prec)$  is linear order(VI.2), we get  $(T \cup B_T, \prec)$  is linear order, too.

2. Since  $\theta(B_t) = U^{\theta(t)}$  is an interval(VI.3), we get  $B_t$  is interval, too.

□

### PROBLEM VIII