# Implementation of Custom Routing Algorithm in Cloud

By Disha Bhattacharya Biswajeet Chakraborty Palash Dey Ananya Laha

#### **Basic Idea**

Algorithm for distributing data to all nodes in a network based on the concept of percolation centrality (PC) or betweenness centrality (BC)

- Enhancement of the controlled flooding algorithm
  - Adding concept of percolation centrality

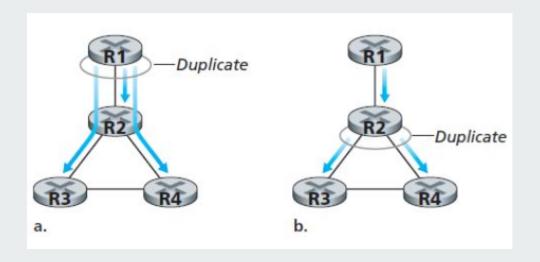
## Purpose

- The purpose of the routing algorithm is to make decisions for the router concerning the optimal paths for data distribution
- The router uses the routing algorithm to get the path that would best serve to transport the data throughout the network
- The routing algorithm that our protocol uses is a major factor in the performance of our routing environment

## **Broadcasting**

- Message is destined to all network devices
- Most straightforward way: N-way-unicast
- Broadcast Algorithms:
  - ➤ 1. Uncontrolled Flooding
  - > 2. Controlled Flooding
  - > 3. Spanning Tree Broadcast

## **Broadcasting**



Source duplication vs in-network duplication

## **Brief Description**

- Similar to the controlled flooding algorithm
- Enhance the flooding algorithm using the concept of percolation centrality
- Send a message that percolates via the nodes of the network
- Time taken will be the least by using the concept of percolation centrality

## **Implementation**

Tech Stack:





## **Algorithm Implementation**

#### Algorithm 1

To start routing from node with highest Betweenness Centrality

- 1. procedure
- 2. graphPC = descending\_PercCentrality(G)
- 3. for i < -0, n-1 do
- 4. graphPC[i].MARK = False
- 5. for i < -0, n-1 do
- 6. Call enhanced\_flooding(graphPC[i])

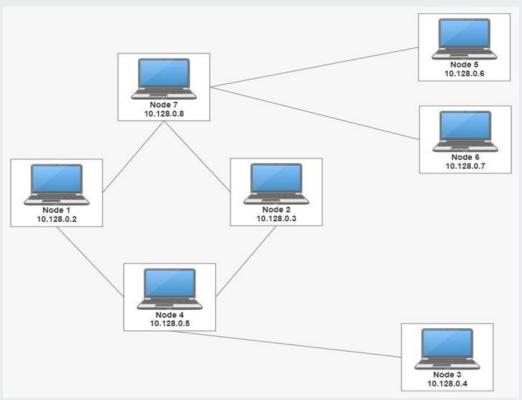
## **Algorithm Implementation**

#### Algorithm 2

Algorithm for controlled flooding mechanism

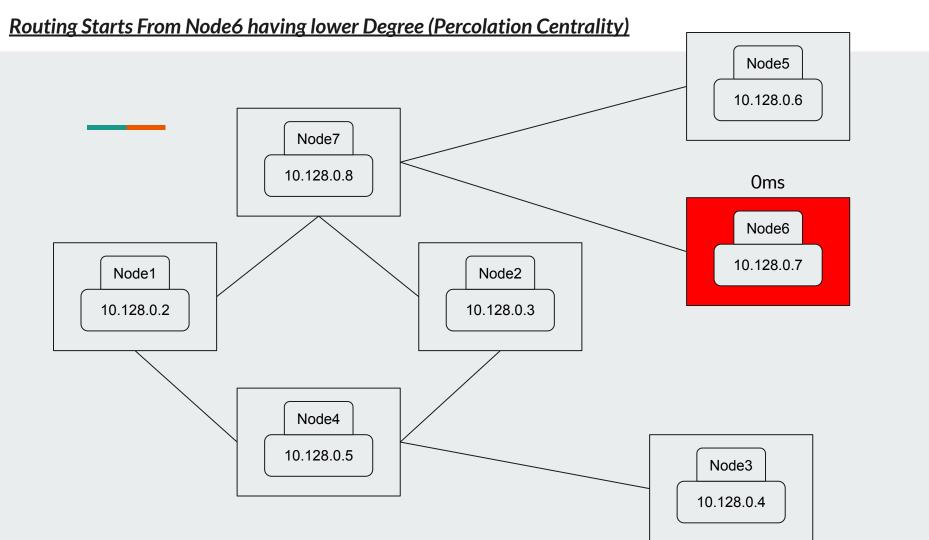
- procedure enhanced\_flooding(v)
- 2. if v.MARK = False then
- 3. v.MARK = True
- 4. Accept message in v
- 5. parfor each node k E v.adjacent() do
- 6. Call enhanced\_flooding(k)
- 7. end parfor

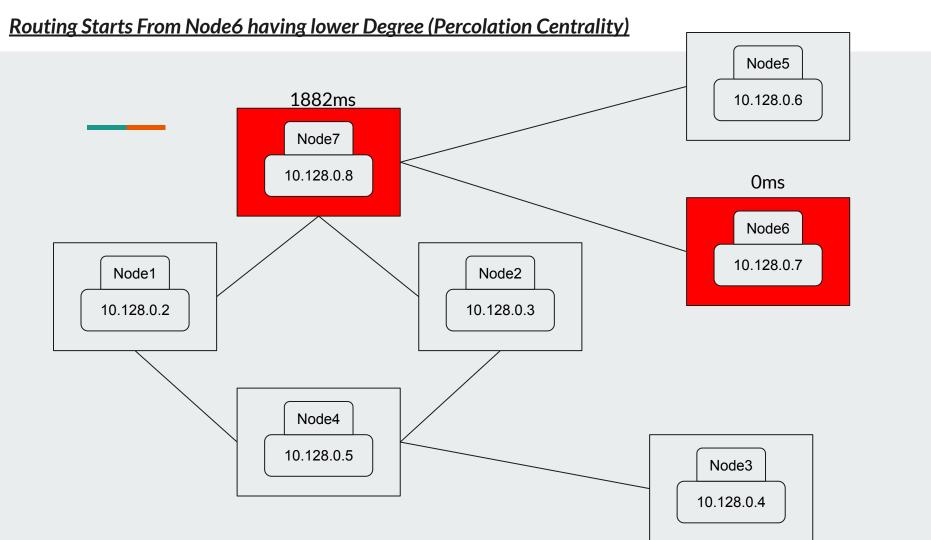
#### **Network Formation**

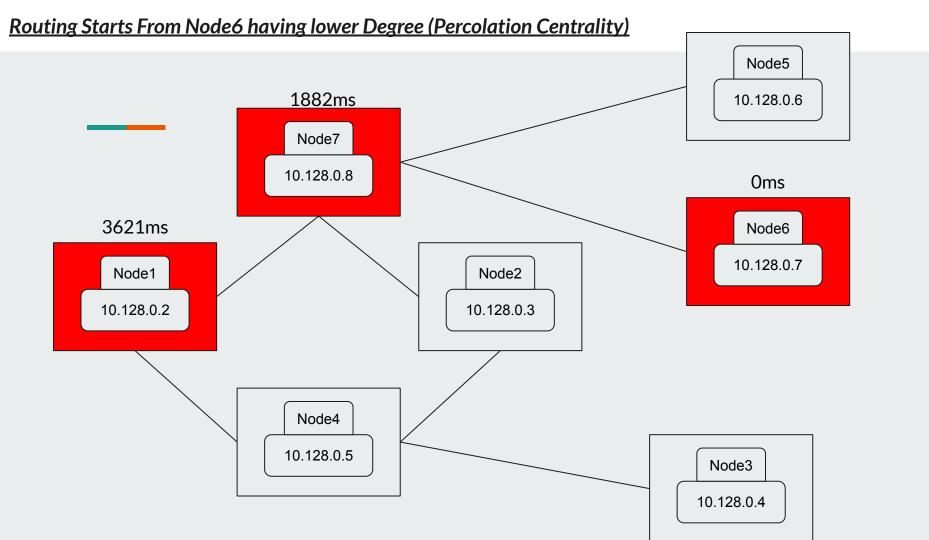


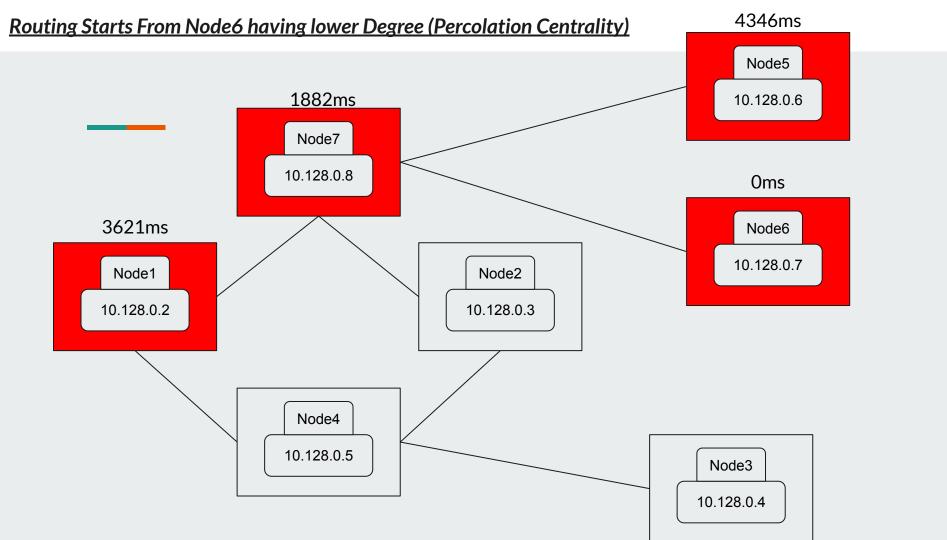
### **Betweenness Centrality of Nodes**

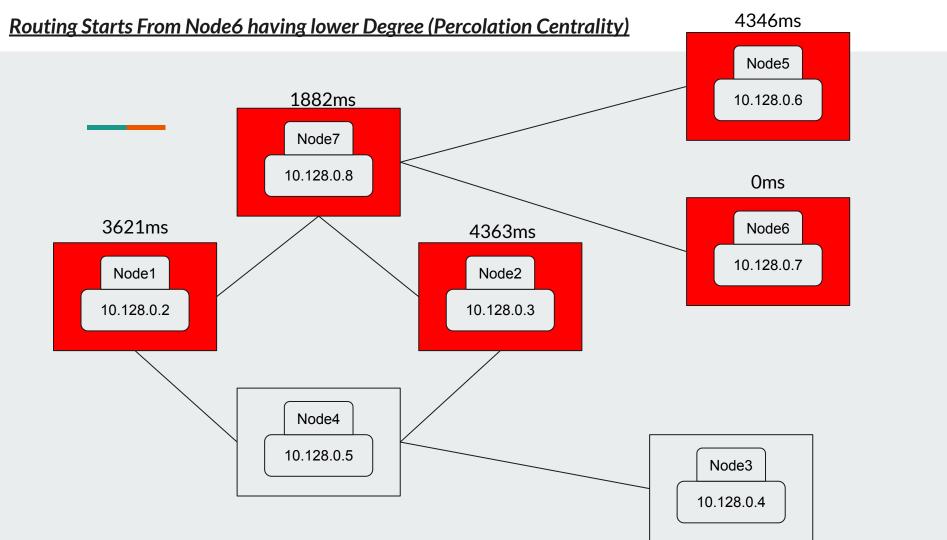
Node ID	0.2 0.2 0				
1					
2					
3					
4	0.366667				
5	0				
6	0				
7	0.633333				

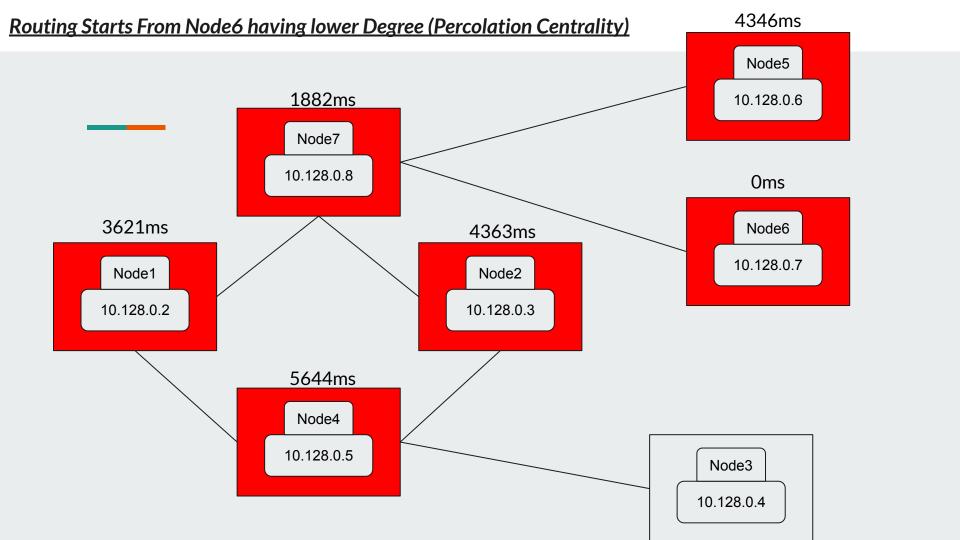


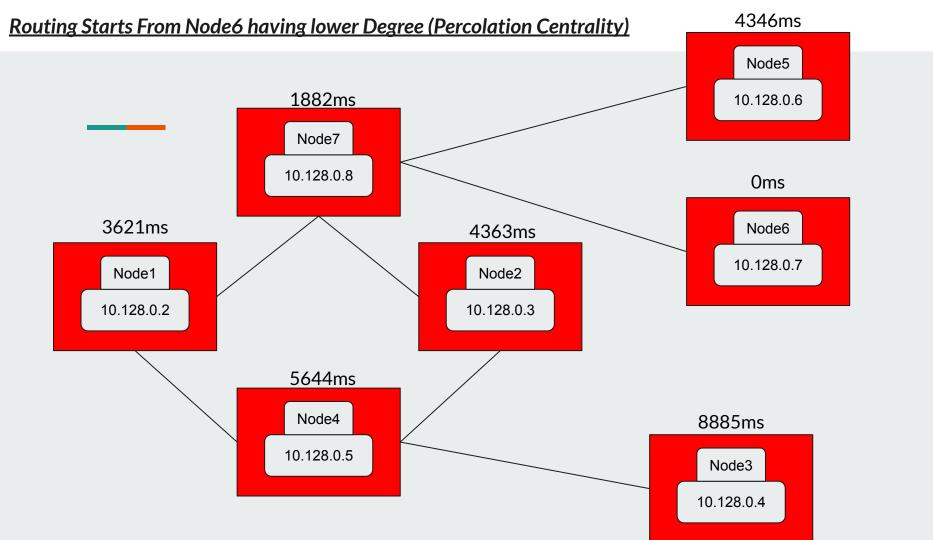


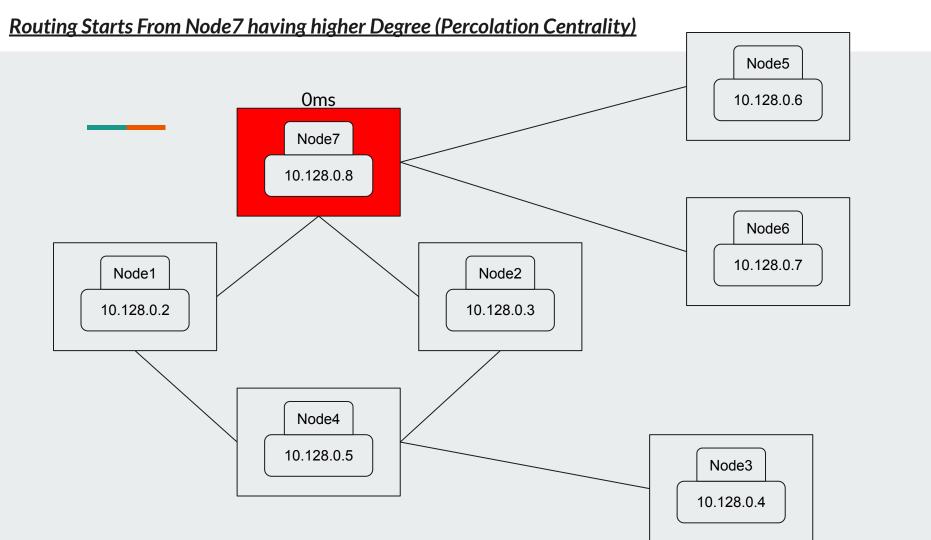


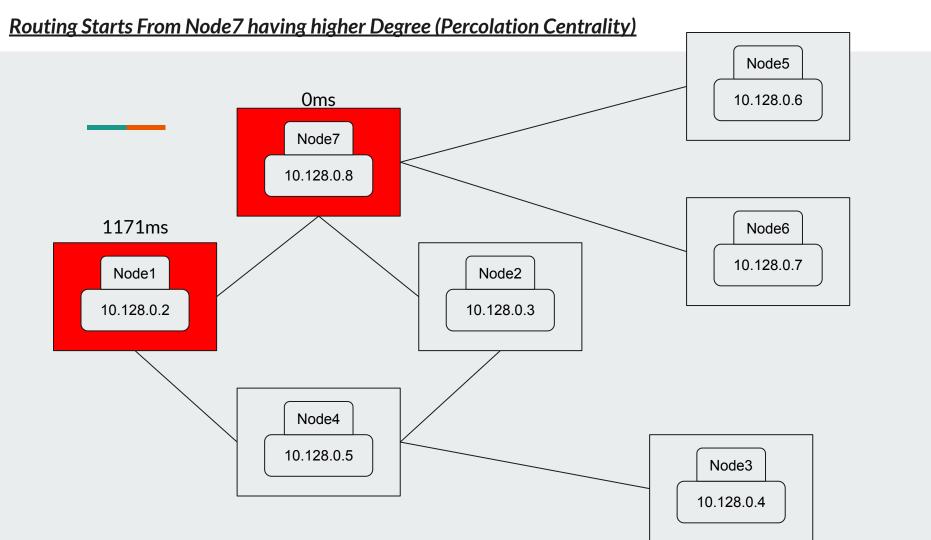


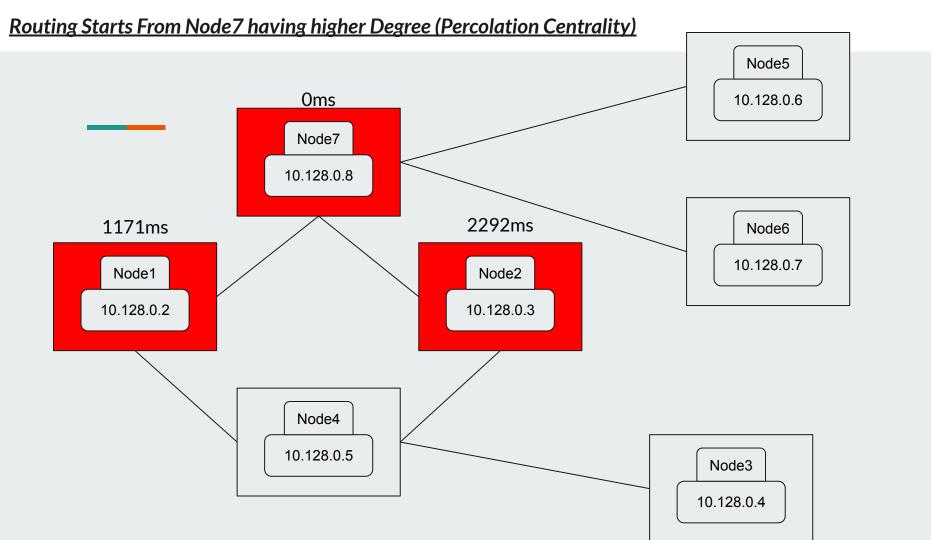


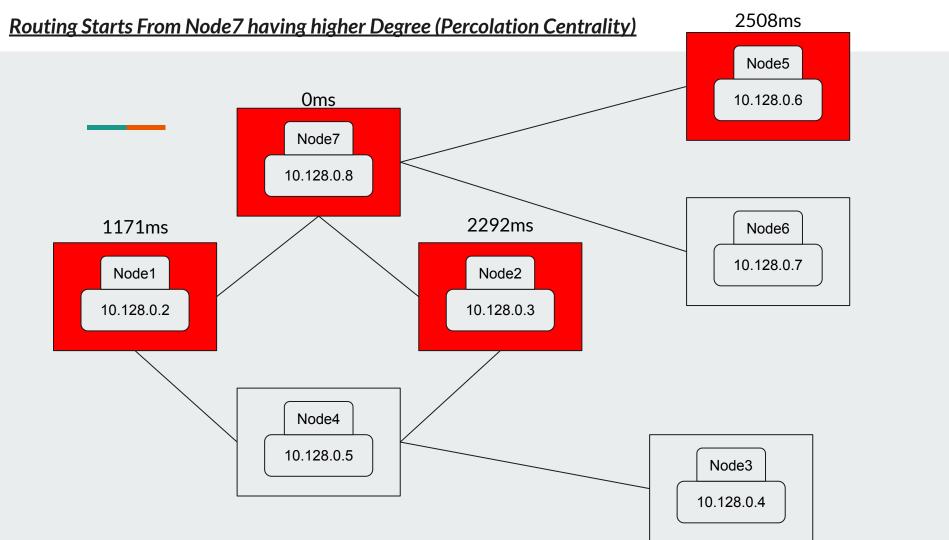


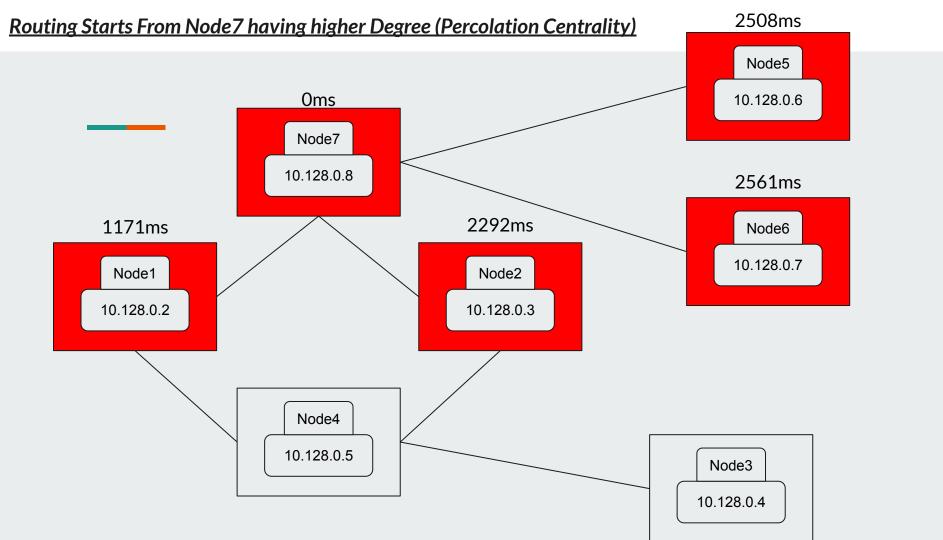


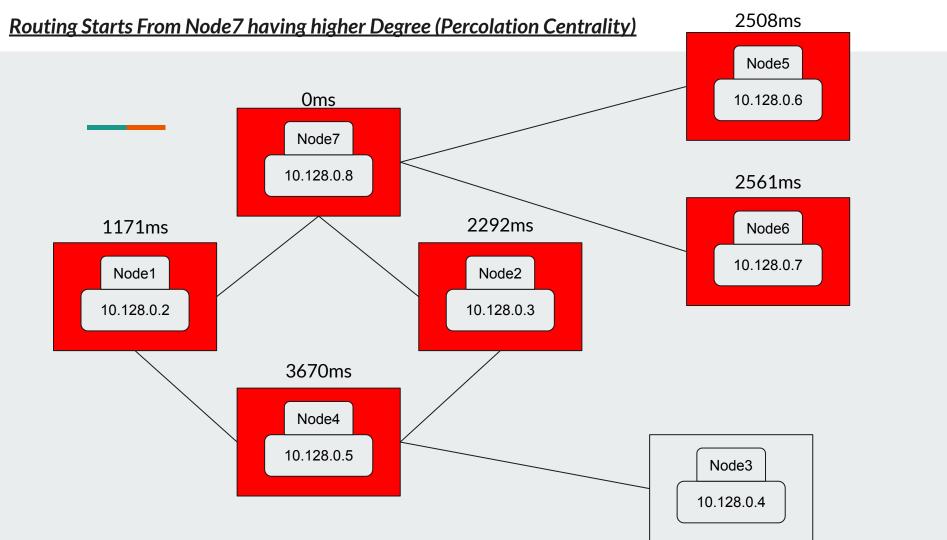


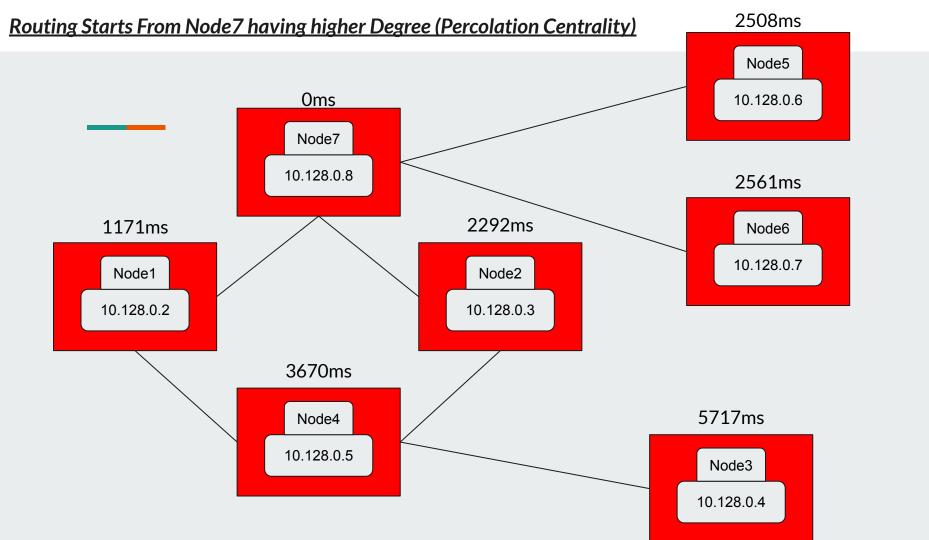










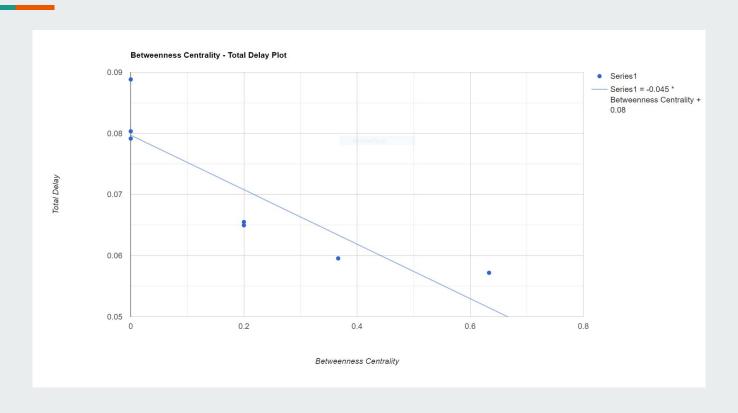


## **Observation**

#### Time taken for distribution:

Starting Node ↓		Time in msec →							
	1	2	3	4	5	6	7	Total Time	
1	0	0.03729	0.03861	0.01765	0.05247	0.06495	0.02108	0.06495	
2	0.03818	0	0.03931	0.0181	0.05266	0.06549	0.022	0.06549	
3	0.03671	0.04063	0	0.01897	0.07665	0.07916	0.05724	0.07916	
4	0.01862	0.02246	0.02237	0	0.05953	0.0583	0.03896	0.05953	
5	0.03961	0.04219	0.08035	0.05755	0	0.04908	0.02033	0.08035	
6	0.03621	0.04363	0.08885	0.05644	0.04346	0	0.01882	0.08885	
7	0.01711	0.02292	0.05717	0.0367	0.02508	0.02561	0	0.05717	

#### Betweenness Centrality vs. Total Delay



#### Conclusion

It can be concluded from the scatter plot that the node having higher value of betweenness centrality will distribute the data/file in lesser time than the one having lower value of betweenness centrality.

## **Future Scope**

- Introduction of automation to minimize the number of manual steps to run the scripts for each routing.
- Distributing resources with complex format

## Thank You