# **Report: Assignment 2**

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# Signal Flow and Key Handling

## 1. Ctrl + C (SIGINT):

- When the user presses Ctrl
   + C, consoleintr() detects
- The SIGINT signal is passed to keyinterrupt() in proc.c.
- Inside keyinterrupt(), the kernel marks all processes (with pid > 2) as killed.
- The trap.c function checks
   if the killed flag is set. If it is, the exit() function is called, which
   terminates the process and cleans up its memory.
- The parent process of the terminated process receives -1 from the wait() system call, indicating that the child process has been killed.

## 2. Ctrl + B (SIGBG):

- When Ctrl + B is pressed, consoleintr() sends the SIGBG signal to keyinterrupt().
- In keyinterrupt(), the kernel marks processes with pid > 2 as suspended. The suspended processes are not scheduled.
- These suspended processes return -2 from wait(), which tells the parent not to wait for them.
- The scheduler checks for the suspended flag and skips these processes until they are resumed.
- Now the parent, ie the shell (pid 2) is woken up.

```
case C('C'): // Ctrl+C (SIGINT)
  consputs("Ctrl-C is detected by xv6\n");
  keyinterrupt(SIGINT);
  break;
case C('B'):
consputs("Ctrl-B is detected by xv6\n");
  keyinterrupt(SIGBG);
  break:
case C('F'):
  consputs("Ctrl-F is detected by xv6\n");
  keyinterrupt(SIGFG);
  break;
case C('G'):
  consputs("Ctrl-G is detected by xv6\n");
  keyinterrupt(SIGCUSTOM);
  break;
```

```
if(myproc() && myproc()->suspended)
wakeup(myproc()->parent);
```

## 3. Ctrl + F (SIGFG):

- Pressing Ctrl + F sends the SIGFG signal to the kernel via consoleintr() and keyinterrupt().
- Inside keyinterrupt(), the suspended flag for the relevant processes is cleared, allowing them to continue running.
- The trap.c function is used to bring the suspended process back into the foreground.

## 4. Ctrl + G (SIGCUSTOM):

 Pressing Ctrl + G triggers the SIGCUSTOM signal, which is detected by consoleintr().

```
int
sys_signal(void)

sighandler_t handler;

if(argptr(0, (void*)&handler, sizeof(handler)) < 0)
    return -1;
myproc()->sighandler = handler;
return 0;
}
```

case SIGCUSTOM:

if (p->sighandler) {

p->pending custom signal = SIGCUSTOM;

- The pending\_custom\_signal is set to the SIGCUSTOM inside only if the handler is registered.
- o In trap.c, if the pending\_custom\_signal is set, it is invoked to handle the custom signal.

```
if(myproc() && (myproc()->pending_custom_signal == SIGCUSTOM)){
  myproc()->tf->esp -= 4;
  *(uint*)(myproc()->tf->esp) = myproc()->tf->eip;
  myproc()->tf->eip = (uint)myproc()->sighandler;
  myproc()->pending_custom_signal = 0;
}
```

#### Xv6 Scheduler

#### Custom fork Creation (custom\_fork)

The **custom** fork function creates processes with a specified execution time:

```
//custom fork
if((p->pid > 2) && (p->exec_time != -1) && (p->cpu_ticks >= p->exec_time)) {
    p->killed = 1;
}
```

- If **exec\_time** is **-1**, the process runs to completion.
- If **exec\_time** is a positive value, the process is forcibly terminated in the scheduler once it exceeds the allowed time.

```
if(p->state == RUNNABLE) {
    | p->dynamic_priority = p->init_priority - (ALPHA * p->cpu_ticks) + (BETA * p->wait_ticks);
    |
}
```

#### Effect of α (ALPHA) on Scheduling

- ALPHA helps control how long a process stays on the CPU. If ALPHA is higher:
  - Processes that have already used more CPU time are less likely to be rescheduled.
  - The scheduler gives priority to processes that have been waiting for CPU time, allowing fairer scheduling between CPU-bound and I/O-bound processes.
- Scheduling Behavior: The higher the ALPHA value, the less likely CPU-heavy processes are to be selected for execution again. This ensures that waiting processes are given a chance to run.

#### Effect of $\beta$ (BETA) on Scheduling

- **BETA** prioritizes processes that have been **waiting** for a long time.
- Higher BETA values make long-waiting processes more likely to be scheduled next.
- **Scheduling Behavior**: This helps prevent starvation of processes that have been waiting for a long time in the queue. It gives them a higher chance of running, ensuring fairness in resource allocation.

#### Test results:

- 1. When ALPHA was set large (10), and BETA was small(1), the processes which were scheduled once, only got a chance in cyclic order, ie, the first one scheduled was rescheduled after all the processes were scheduled once.
- 2. When ALPHA was set to a negative large value, then the process that was running initially was the one which got rescheduled.
- 3. When BETA was set negative and ALPHA was also negative then the process which waited a lot was frequently scheduled.

## **Overall Scheduler Behavior**

#### **CPU-Bound Jobs**

When **ALPHA** is higher, the scheduler deprioritizes processes that have already consumed more CPU time.

Higher ALPHA = Less Likely to be Rescheduled: This means that
 CPU-bound jobs, which tend to consume a lot of CPU time, are less likely to be selected for execution when ALPHA is set to a higher value.

**BETA** prioritizes processes that have been waiting for CPU time, i.e., processes with higher waiting times are more likely to be scheduled next.

Higher BETA = Less Likely to Favor CPU-Bound Jobs: Since BETA increases
the chances of long-waiting processes (like I/O-bound jobs) getting scheduled,
CPU-bound jobs are less likely to be selected if they have not been waiting for a
long time.

#### I/O-Bound Jobs:

- ALPHA impacts how much CPU time a process has already consumed. If ALPHA is set to a higher value, the scheduler deprioritizes processes that have used more CPU time.
  - Lower ALPHA = More Likely to be Rescheduled: I/O-bound jobs
    typically consume less CPU time since they spend most of their time
    waiting for I/O. If ALPHA is lower, these processes are more likely to be
    scheduled because they haven't consumed much CPU time.
- BETA prioritizes processes that have waited the longest for CPU time.

Higher BETA = More Likely to be Scheduled: Since I/O-bound jobs typically spend a lot of time waiting for I/O, they are likely to have high waiting times in the scheduler queue. Higher BETA ensures that these processes will have a better chance of getting scheduled, as their waiting time will be taken into account. This allows I/O-bound jobs to execute when the CPU is available, rather than being overlooked in favor of CPU-bound jobs.