# System Calls

# 1. GOTTA COUNT 'EM ALL

Added array of size 32 to process struct. In the <a href="syscall()">syscall()</a> function in kernel/syscall.c, whenever a syscall is called, add it to the count. In the <a href="syspecsec">syspecsec</a>, extract the required syscall from the mask and return the value of it from the current process if it's valid. While forking, syscall count is copied from parent to child. The user program syscount forks and executes the command and the parent process waits and gets the appropriate count of the syscall and prints it.

### 2. WAKE ME UP WHEN MY TIMER ENDS

Added the following to proc struct in proc.h:

- int interval: duration after which handler called
- int current\_ticks: track ticks passed so far
- int alarm\_status: check if alarm currently set or not
- uint64 alarm handler: pointer to handler function
- **struct trapframe** \***alarm\_tf**: copies trapframe for sigreturn

Initialized/freed these in allocproc()/freeproc() respectively. Added appropriate definitions/definitions in kernel/syscall.c, kernel/syscall.h, user/usys.pl, and user/user.h. Added the new syscalls to the list in user/syscount.c

## In kernel/sysproc.c:

- Defined function uint64 sys\_sigalarm(void) which updates the process with the alarm interval and handler and initializes current\_ticks and alarm\_status to 0.
- Defined function uint64 sys\_sigreturn(void) which restores the original trapframe, frees the copied alarm trapframe, and resets alarm\_status to 0. It also invokes usertrapret() which returns to userspace from the trap and continues execution of the process, now with the updated/restored trapframe.

# In kernel/trap.c:

- Modified the usertrap() function
- Updated the section pertaining to timer interrupts
- If the interval is > 0, that means some interval is set, so now we check the status of the alarm. If it is 0, we can proceed as normal and update the current\_ticks attribute of the process.
- Now, if the current\_ticks reach the interval duration, we reset current\_ticks to 0, and we
  now set alarm\_status = 1. We use kalloc() and memmove() to store the state of the
  trapframe, and then update the program counter to the signal handler.

# Scheduling

#### 1. THE PROCESS POWERBALL

- Added scheduler maco to makefile.
- Added user program settickets that takes input and sets current process tickets accordingly
- Added settickets syscall to list of syscalls
- Modified proc struct in proc.h to have an additional field for tickets
- Initialize tickets to 0 in allocproc in proc.c
- In fork, set child process's tickets to same amount as it's parents tickets
- In the scheduler function, used preprocessor directives to switch between round robin and lottery based scheduling
- Now in the LBS scheduler, we iterate through all runnable processes to calculate total tickets held
- Defined a random generator and seed function in proc.c
- We then initialize the seed based on system tick count and accordingly generate random numbers
- The random number generated is deemed the winning ticket of the lottery
- Now again iterate through all processes, subtracting ticket values from the winning ticket till we get the process that holds the winning ticket
- We set this process to running and a context switch occurs to hand over control

```
xv6 kernel is booting
Using LBS
hart 2 starting
hart 1 starting
init: starting sh
$ schedulertest
Process 9 finished
Process 3 finished
Process 4 finished
Process 8 finished
Process 0 finished
Process 1 finished
Process 2 finished
21 ticket. rtime: 0, wtime: 41
19 ticket. rtime: 0, wtime: 41
17 ticket. rtime: 0, wtime: 41
15 ticket. rtime: 0, wtime: 40
13 ticket. rtime: 0, wtime: 40
11 ticket. rtime: 22, wtime: 19
9 ticket. rtime: 27, wtime: 13
Process 7 finished
7 ticket. rtime: 26, wtime: 17
Process 6 finished
5 ticket. rtime: 24, wtime: 22
Process 5 finished
3 ticket. rtime: 13, wtime: 38
Average rtime 11, wtime 31
```

The output for the lottery based scheduler can be seen above.

```
xv6 kernel is booting
Using RR
hart 2 starting
hart 1 starting
init: starting sh
$ schedulertest
Process 8 finished
21 ticket. rtime: 26, wtime: 19
Process 7 finished
19 ticket. rtime: 25, wtime: 21
Process 6 finished
17 ticket. rtime: 25, wtime: 21
Process 9 finished
15 ticket. rtime: 25, wtime: 21
Process 5 finished
13 ticket. rtime: 15, wtime: 32
PPPrrorcoecsess so 1c efi2s sf inniisshheedd
0 fiPPnrirsohceeoscseds
41 If itfiinniisshheedd
cket. rtime: 0, wtime: 200
9 ticket. rtime: 0, wtime: 200
7 ticket. rtime: 0, wtime: 200
5 ticket. rtime: 0, wtime: 200
3 ticket. rtime: 0, wtime: 200
Average rtime 11, wtime 111
```

Meanwhile, above is the output for RR scheduling.

## **Answer to report question:**

Adding arrival time ensures that if multiple processes have the same number of tickets, the older process will get higher priority. This helps avoid older processes from starvation. However, this approach also brings with it the additional overhead of checking and managing arrival times.

If all processes had the same number of tickets, the scheduler would essentially operate the same as a First Come First Serve scheduler, as it would then just continually choose the oldest process.

# Networking

### 1. XOXO

## TCP:

# Firstly, created header.h:

- Define port for communication, size of each chunk and max chunks allowed
- Defined buffer size, board size, player characters and empty character
- Defined player struct consisting of their socket and their character

## Created client.c:

- Initialize socket for this player to communicate with the server
- In a loop, based on data sent by server, either ask player for a move or ask them if they want to play again
- At the beginning of each iteration of the loop, clear the buffer.
- If did not receive any data, server has disconnected, so break the loop
- Otherwise, print the data from the server.
- Now if server is asking for move, return player's move
- Else, return player's choice to play again
- If loop terminates, close the socket

### Created server.c:

- Initialize board, players array, player count, current player, game running or not as global variables.
- Also initialize mutex lock and cond variable
- Now in the main function, we set up the server and wait for players to connect
- We also set up array tid to hold identifiers for the player threads
- Now we initialize the mutex lock and condition variable
- We also initialize the board such that each tile is blank in the beginning
- We then setup server socket and begin listening for connections
- We loop twice to accept two connections from players, and store the socket descriptor for this
- We create two threads to handle player communication using pthread create()
- The handle client() function is responsible for handling communication with the players
- The main function waits for both threads to finish using pthread\_join, after which it closes the server socket and destroys the mutex lock and condition variable

In handle\_client():

- The function has a single argument which we cast to type Player \* to represent the player structure of the corresponding client
- We maintain a buffer to send/receive messages and row and col variables to represent the players move
- While the game hasn't ended, keep looping
- We first lock the mutex so that both clients wait until their respective turns
- Whenever it is a players turn, we send them an intimation for their move. If no data is found, player has disconnected. Otherwise, update the board state with their move.
- After every move, check for a winner or a draw and send the result if that is reached
- If not, update the current player and send updated board to both
- At the end, broadcast the condition variable so the next player proceeds, and unlock the mutex
- After the game ends, ask player 1 if they want to play again. If they say yes, ask player 2 if they want to play again. If either says no, break the game loop.
   Otherwise, reset the board and current player and start again
- If game loop is over, close both player sockets and return control to main function

# 2. Fake it till you make it

## Firstly created header.h:

- Define port for communication, size of each chunk and max chunks allowed
- Defined timeout interval, 100000 microseconds, or 1 second
- Defined packet struct containing it's sequence number, total chunks and the data itself
- Defined packet info struct which contains the packet, the time it was sent at and if it's been acked or not

### Created server.c:

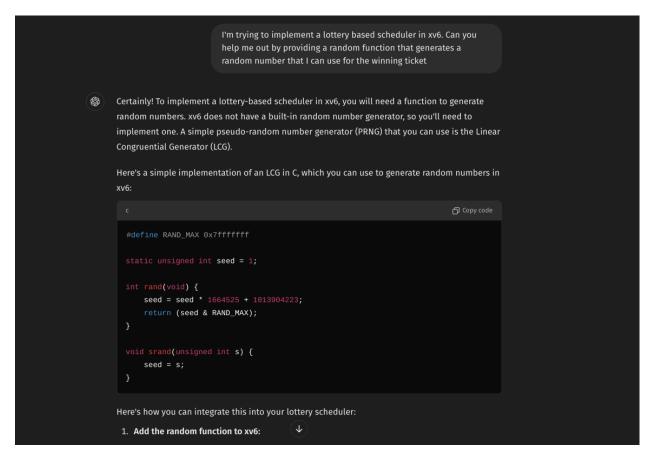
- Initialize socket, Packet to store incoming packet, ack, array for each chunk, total chunks and chunks received
- In an infinite loop, listen for messages
- If we received some data, print we received this chunk
- Send an acknowledgement for that chunk
- For first chunk, get total chunks and set NULL in the array to store upcoming chunks
- If value in array for any chunk seq num is currently NULL, allocate memory for that chunk's message and store it. [Data sequencing]
- If we've received all chunks, loop through array and print values and free memory. Now reset, total, chunks received for next message and start over

### Created client.c:

- Initialize ack, message, packet info, socket
- Set the socket to non-blocking mode

- In an infinite loop, get message from user
- Calculate the total chunks needed to transmit the message based on the input
- Set up packet info by setting sequence num, total chunks and message value. Also define acked as 0 by default and get the current time to be the time it is sent
- While all values haven't been acked, loop through them and check if they've been acked, If any hasn't all acked becomes false. Now check the time that the non acked packet was sent at. If it's over our time definition, resend it and update the send time. We also add a timeout to decide when to next resend it
- In this ack loop we also check if anything has been received from the server in the form of an ACK
- If yes, update that packet's info with acked and print that we have received the ack

# **ChatGPT Commands**



#### 1. Add the random function to xv6:

 Add the rand and srand functions to a suitable file in the xv6 source code. You can create a new file called rand.c and add these functions there.

#### 2. Include the random functions in your scheduler:

 In your scheduler code, you can include the rand.h header (which you'll need to create) to access the rand and srand functions.

#### 3. Initialize the seed:

 Initialize the seed using srand at some point during the boot process or at the start of the scheduler.

### 4. Generate a random number to pick the winning ticket:

 Use the rand function to generate a random number each time you need to pick a winning ticket.

Here's how you can create the rand.c and rand.h files:

#### rand.c:

```
c
#include "rand.h"

#define RAND_MAX 0x7fffffff

static unsigned int seed = 1;

int rand(void) {
    seed = seed * 1664525 + 1013904223;
    return (seed & RAND_MAX);
}
```

This implementation uses a simple LCG to ge.  $\checkmark$  ate random numbers, and the scheduler uses

```
void scheduler(void) {
    ...
    srand(ticks); // Initialize the seed with the system tick count or any oth
    int total_tickets = get_total_tickets(); // Function to calculate total tic
    int winning_ticket = rand() % total_tickets;

struct proc *p;
    int current_ticket = 0;
    for(p = ptable.proc; p < &ptable.proc[NPROC]; p++) {
        if(p->state != RUNNABLE)
            continue;
        current_ticket += p->tickets;
        if(current_ticket > winning_ticket) {
            // p is the chosen process
            // Context switch to p
            ...
            break;
        }
    }
    ...
}
```

This implementation uses a simple LCG to generate random numbers, and the scheduler uses these random numbers to select a winning ticket for the lottery. This is a straightforward approach and should work well for basic lottery scheduling in xv6.



what are the general steps to add a syscall to xv6



Adding a system call (syscall) to the XV6 operating system involves several steps. Here's a general guide:

# Step 1: Choose a Syscall Name and Function

- Decide on the functionality: Determine what the syscall will do (e.g., read a file, get system information).
- 2. Name the syscall: Choose a descriptive name for your syscall, such as <code>mysyscall</code>.

## Step 2: Define the Syscall

- 1. Implement the function:
  - Create the implementation of your syscall in a .c file, usually in the kernel directory.
     For example, sys\_mysyscall.c.

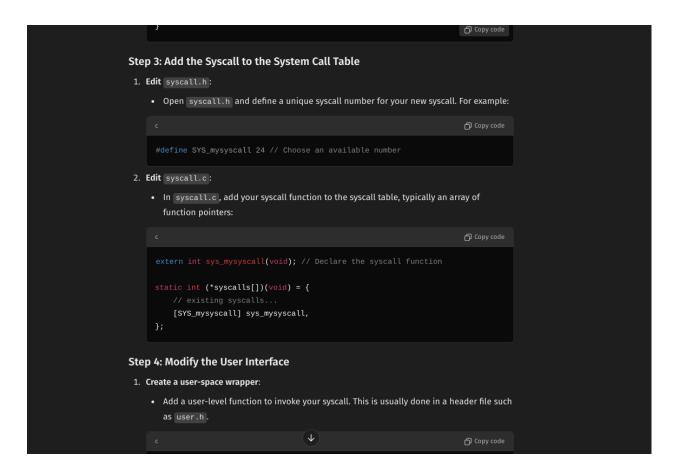
```
c
int sys_mysyscall(void) {
    // Implementation of the syscall
    return 0; // Return an integer as a status or result
}
```

## Step 3: Add the Syscall to the System Call Table

- 1. Edit syscall.h:
  - Open syscall.h and define a unique syscall number for your new syscall. For example:



Copy code



```
c
int mysyscall(void) {
   return syscall(SYS_mysyscall);
}
```

# Step 5: Update the Makefile

- 1. Include your new syscall file:
  - Ensure that the newly created <code>.c</code> file is included in the compilation process. Update the <code>Makefile</code> in the <code>kernel</code> directory:

```
makefile

OBJS = sys_mysyscall.o ... // Add your syscall object file here
```

### Step 6: Test the Syscall

- 1. Create a user-level test program:
  - Write a simple user program that calls your new syscall to test its functionality.

```
c

#include "user.h"

int main(int arge, char *argv[]) {
   mysyscall(); // Call your syscall
   exit(0);
}
```

2. Compile and run the test:

Copy code

## 2. Compile and run the test:

• Compile your user program and run it to see if it works as expected.

## **Step 7: Debugging and Refinement**

- 1. **Check for errors**: If the syscall does not work as intended, use print statements or a debugger to identify issues.
- 2. Refine the implementation: Make any necessary adjustments based on your testing.

#### **Example Summary**

This example briefly summarizes the process for adding a syscall named | mysyscall :

- Implement sys\_mysyscall in sys\_mysyscall.c.
- Define it in syscall.h.
- Add it to the syscall table in syscall.c.
- Create a user-level wrapper in user.h.
- Update the Makefile.
- Write a test program to validate your syscall.

### Note

Make sure you consult the existing syscalls in XV6 as references to ensure consistency in your implementation.

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how can i use non blocking sockets and udp to send data

in C

Using non-blocking sockets with UDP in C involves several steps. Here's a complete example of how to set up a non-blocking UDP socket, send data, and handle sending and receiving without blocking. This example includes creating a socket, setting it to non-blocking mode, and sending data.

# Step 1: Include Required Headers

```
#include <stdio.h>
#include <stdib.h>
#include <string.h>
#include <unistd.h>
#include <arpa/inet.h>
#include <fcntl.h>
#include <errno.h>
```

# Step 2: Create a Non-Blocking UDP Socket

You will create a UDP socket and set it to non-blocking mode.

```
int create_non_blocking_udp_socket() {
  int sockfd = socket(AF_INET, SOCK_DGRAM, 0);
  if (sockfd < 0) {
     perror("socket creation failed");
}</pre>
```

```
#include <unistd.h>
#include <arpa/inet.h>
#include <fcntl.h>
#include <errno.h>
```

# Step 2: Create a Non-Blocking UDP Socket

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```
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int create_non_blocking_udp_socket() {
   int sockfd = socket(AF_INET, SOCK_DGRAM, 0);
   if (sockfd < 0) {</pre>
        perror("socket creation failed");
       exit(EXIT_FAILURE);
   // Set the socket to non-blocking mode
   int flags = fcntl(sockfd, F_GETFL, 0);
   if (flags < 0) {</pre>
       perror("fcntl get flags failed");
       close(sockfd);
       exit(EXIT_FAILURE);
   if (fcntl(sockfd, F_SETFL, flags | 0_NONBLOCK) < 0) {</pre>
        perror("fcntl set non-blocking failed");
       close(sockfd);
       exit(EXIT_FAILURE);
   return sockfd;
```

```
return sockfd;

[7] Copy code

}
```

# Step 3: Send Data Using Non-Blocking Socket

Now you can implement a function to send data. If sending fails due to the socket being non-blocking, you can handle that error appropriately.

```
Copy code
void send_data(int sockfd, const char *message, const char *ip, int port) {
   struct sockaddr_in dest_addr;
   memset(&dest_addr, 0, sizeof(dest_addr));
   dest_addr.sin_family = AF_INET;
   dest_addr.sin_port = htons(port);
   inet_pton(AF_INET, ip, &dest_addr.sin_addr);
   ssize_t sent_bytes = sendto(sockfd, message, strlen(message), 0,
                                 (struct sockaddr *)&dest_addr, sizeof(dest_addr));
   if (sent_bytes < 0) {</pre>
       if (errno == EAGAIN || errno == EWOULDBLOCK) {
           printf("Send would block, try again later\n");
       } else {
           perror("sendto failed");
   } else {
       printf("Sent %ld bytes: %s\n", sent_bytes, message);
```



# Step 4: Example Usage

Now you can put everything together and create a simple program that sends data over UDP.

```
int main() {
   int sockfd = create_non_blocking_udp_socket();
   const char *message = "Hello, UDP!";
   const char *ip = "127.0.0.1"; // Change this to the target IP
   int port = 8080; // Change this to the target port

   // Attempt to send data
   send_data(sockfd, message, ip, port);

   // Close the socket
   close(sockfd);
   return 0;
}
```

# **Explanation**

- 1. **Socket Creation**: The create\_non\_blocking\_udp\_socket function creates a UDP socket and sets it to non-blocking mode using fcntl.
- Sending Data: The send\_data function prepares the destination address and sends the
  message using sendto. If sending would block (since the socket is non-blocking), it prints a
  message and continues.
- 3. Error Handling: If an error occurs (other than would-block), it prints the error.

# **Compilation and Running**





# **Explanation**

- 1. **Socket Creation**: The create\_non\_blocking\_udp\_socket function creates a UDP socket and sets it to non-blocking mode using fcntl.
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  message using sendto. If sending would block (since the socket is non-blocking), it prints a
  message and continues.
- 3. Error Handling: If an error occurs (other than would-block), it prints the error.

# **Compilation and Running**

You can compile the code with gcc:



# **Note on Receiving Data**

Receiving data with a non-blocking socket is similar; you would use recvfrom and check for EAGAIN or EWOULDBLOCK errors to handle the non-blocking nature.

# **Example of Receiving Data**



./udp\_nonblocking

# Note on Receiving Data

Receiving data with a non-blocking socket is similar; you would use recyfrom and check for EAGAIN or EWOULDBLOCK errors to handle the non-blocking nature.

🖒 Copy code

# **Example of Receiving Data**

Here's a simple example of how to receive data with a non-blocking socket:

You can call receive\_data(sockfd) in a loop or after sending data to continuously check for incoming messages.

This code provides a solid foundation for working with non-blocking UDP sockets in C. You can expand upon this as needed for your application!



© Creating a Tic Tac Toe game in C is a fun project that can help you practice basic programming concepts such as arrays, loops, and conditionals. Below is a simple implementation of a command-line Tic Tac Toe game for two players.

## Tic Tac Toe Game in C

Here's a complete C program for a simple Tic Tac Toe game:

```
🖒 Copy code
void initializeBoard(char board[3][3]) {
    for (int i = 0; i < 3; i++) {
        for (int j = 0; j < 3; j++) {
            board[i][j] = ' ';
void printBoard(char board[3][3]) {
    printf(" 0 1 2\n"); // Column headers
    for (int i = 0; i < 3; i++) {
        printf("%d ", i); // Row header
        for (int j = 0; j < 3; j++) {
            printf("%c", board[i][j]);
            if (j < 2) {
                                       \downarrow
```

```
if (i < 2) {
            printf(" ----\n"); // Horizontal separator
int checkWin(char board[3][3]) {
        if ((board[i][0] == board[i][1] && board[i][1] == board[i][2] && board[i][0]
            (board[0][i] == board[1][i] \&\& board[1][i] == board[2][i] \&\& board[0][i]
    if ((board[0][0] == board[1][1] && board[1][1] == board[2][2] && board[0][0] !=
        (board[0][2] == board[1][1] && board[1][1] == board[2][0] && board[0][2] !=
int checkDraw(char board[3][3]) {
        for (int j = 0; j < 3; j++) {
            if (board[i][j] == ' ') {
```

```
char board[3][3];
int row, col;
int player = 0; // Player 0 starts
char players[2] = {'X', '0'}; // Player symbols
initializeBoard(board);
while (1) {
    printBoard(board);
    printf("Player %d (%c), enter your move (row and column): ", player + 1, pl
    scanf("%d %d", &row, &col);
    if (row < 0 || row > 2 || col < 0 || col > 2 || board[row][col] != ' ') {
    board[row][col] = players[player];
    if (checkWin(board)) {
        printBoard(board);
        printf("Player %d (%c) wins!\n", player + 1, players[player]);
    if (checkDraw(board)) {
        printBoard(board);
```

```
if (checkDraw(board)) {
    printBoard(board);
    printf("It's a draw!\n");
    break;
}

// Switch players
    player = 1 - player; // Toggle between 0 and 1
}

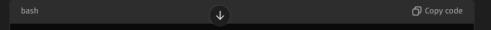
return 0;
}
```

## **Explanation**

- 1. **Initialize the Board**: The initializeBoard function fills the board with spaces ('') to represent empty spots.
- 2. **Print the Board**: The printBoard function displays the current state of the game board.
- 3. **Check for a Win**: The checkWin function evaluates all possible winning combinations (rows, columns, and diagonals).
- 4. Check for a Draw: The checkDraw function checks if all spots are filled without a winner.
- 5. **Main Function**: The main function contains the game loop, where players take turns to enter their moves, and the game state is checked after each move.

## **How to Compile and Run**

- 1. Save the code in a file named tic\_tac\_toe.c.
- 2. Open your terminal or command prompt and navigate to the directory where the file is saved.
- 3. Compile the code using a C compiler, such as GCC:

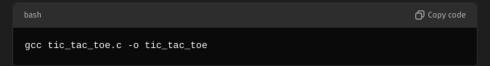


# **Explanation**

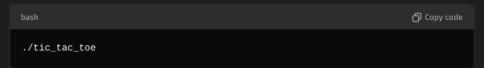
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# **How to Compile and Run**

- 1. Save the code in a file named tic\_tac\_toe.c.
- 2. Open your terminal or command prompt and navigate to the directory where the file is saved.
- 3. Compile the code using a C compiler, such as GCC:



4. Run the compiled program:



This will start the Tic Tac Toe game, allowing two players to enter their moves and see the game progress in the terminal. Enjoy playing!





how do threads work in networking, what does pthread do, how can i use it, what is cond and mutex

**\$** 

Threads in networking allow programs to perform multiple tasks simultaneously, which is particularly useful for handling multiple connections in a client-server architecture. Below is an explanation of how threads work in networking, an overview of pthread, and a detailed description of condition variables and mutexes.

# **How Threads Work in Networking**

### 1. Concurrency

 Concurrency allows a program to manage multiple tasks at once. In the context of networking, this means that a server can handle multiple client requests simultaneously, improving responsiveness and efficiency.

## 2. Threading Models

- Multithreaded Server: Each incoming client connection is handled by a separate thread. This
  way, the server can process multiple requests concurrently without blocking others.
- Thread Pool: A pool of worker threads is maintained. Incoming requests are assigned to an
  available thread, reducing the overhead of creating and destroying threads for each request.

### 3. Advantages

- Improved responsiveness: The server can process multiple connections simultaneously.
- Better resource utilization: Threads can be managed efficiently, sharing the same process resources.
- Simplified handling of asynchronous events: Each thread can handle I/O operations independently.

 Simplified handling of asynchronous events: Each thread can handle I/O operations independently.

# What is pthread?

pthread stands for "POSIX threads," which is a standard for threading in UNIX-like operating systems. It provides a set of APIs to create and manage threads, synchronize their execution, and perform inter-thread communication.

# Common pthread Functions

### 1. Creating Threads

• pthread\_create(): Used to create a new thread. It takes parameters for the thread identifier, attributes, the function to run, and arguments to pass to that function.

```
pthread_t thread;
pthread_create(&thread, NULL, thread_function, arg);
```

# 2. Joining Threads

 pthread\_join(): Waits for a thread to terminate. It can be used to retrieve the exit status of the thread.

```
pthread_join(thread, NULL);
```

## 3. Exiting Threads

 pthread\_exit(): Terminates a thread. The exit status can be retrieved by other threads using pthread\_join().

### 4. Thread Attributes



#### 4. Thread Attributes

 pthread\_attr\_init(), pthread\_attr\_setdetachstate(), etc.: Functions to set various attributes for threads, such as whether they are joinable or detached.

# **Using** pthread

Here's a basic example of using pthread in a multithreaded server that handles client connections:

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
#include <qtring.h>
#include <arpa/inet.h>

#define PORT 8080
#define BACKLOG 5

void *handle_client(void *arg) {
   int client_socket = *((int *)arg);
   free(arg); // Free dynamically allocated memory
   char buffer[1024];

// Communicate with the client
   recv(client_socket, buffer, sizeof(buffer), 0);
   printf("Received message: %s\n", buffer);
   send(client_socket, "Hello from server", 17, 0);

close(client_socket);
   return NULL;
}
```

```
🖒 Copy code
int main() {
   int server_socket, *client_socket;
    socklen_t addr_size;
    server_socket = socket(AF_INET, SOCK_STREAM, 0);
    server_addr.sin_family = AF_INET;
    server_addr.sin_addr.s_addr = INADDR_ANY;
    server_addr.sin_port = htons(PORT);
    bind(server_socket, (struct sockaddr *)&server_addr, sizeof(server_addr));
    listen(server_socket, BACKLOG);
    printf("Server is listening on port %d\n", PORT);
   while (1) {
       addr_size = sizeof(client_addr);
       client_socket = malloc(sizeof(int)); // Allocate memory for the client sock
       *client_socket = accept(server_socket, (struct sockaddr *)&client_addr, &ad
       pthread_t thread;
       pthread_create(&thread, NULL, handle_client, client_socket);
       pthread_detach(thread); // Detach the thread to avoid memory leaks
   close(server_socket);
```

# **Explanation of Example**

• Socket Creation: A TCP socket is created \( \frac{1}{2} \) sten for incoming connections.

#### Explanation of Example

- Socket Creation: A TCP socket is created to listen for incoming connections.
- Accepting Connections: For each incoming connection, a new thread is created to handle the client using pthread\_create().
- . Thread Function: The handle\_client function handles the communication with the client.
- **Memory Management**: Memory is dynamically allocated for the client socket descriptor, which is freed in the thread function after use.

# What are Condition Variables and Mutexes?

# 1. Mutex (Mutual Exclusion)

- Purpose: A mutex is used to protect shared resources from concurrent access. Only one thread
  can lock a mutex at a time, ensuring that critical sections of code are not executed
  simultaneously by multiple threads.
- · Usage:

```
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL); // Initialize the mutex

pthread_mutex_lock(&mutex); // Lock the mutex

// Critical section
pthread_mutex_unlock(&mutex); // Unlock the mutex

pthread_mutex_destroy(&mutex); // Destroy the mutex
```

# 2. Condition Variables

• **Purpose**: Condition variables are used for signaling between threads. They allow a thread to wait for a certain condition to be met wh ↓ eleasing the associated mutex.

- Purpose: Condition variables are used for signaling between threads. They allow a thread to
  wait for a certain condition to be met while releasing the associated mutex.
- Usage:

```
pthread_cond_t cond;
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL);
pthread_cond_init(&cond, NULL);

// Thread that signals
pthread_mutex_lock(&mutex);
// Change some condition
pthread_cond_signal(&cond); // Signal waiting threads
pthread_mutex_unlock(&mutex);

// Thread that waits
pthread_mutex_lock(&mutex);
while (/* condition not met */) {
    pthread_cond_wait(&cond, &mutex); // Wait for the condition
}
pthread_mutex_unlock(&mutex);

pthread_mutex_unlock(&mutex);

pthread_cond_destroy(&cond);
pthread_mutex_destroy(&mutex);
```

# **Key Points**

- Mutexes ensure that only one thread accesses a shared resource at a time, preventing race conditions.
- Condition Variables allow threads to wait for certain conditions and to signal other threads when those conditions change, facilitating coordination between threads.

 $\downarrow$ 

```
pthread_mutex_lock(&mutex);

// Change some condition
pthread_cond_signal(&cond); // Signal waiting threads
pthread_mutex_unlock(&mutex);

// Thread that waits
pthread_mutex_lock(&mutex);
while (/* condition not met */) {
    pthread_cond_wait(&cond, &mutex); // Wait for the condition
}
pthread_mutex_unlock(&mutex);

pthread_mutex_unlock(&mutex);

pthread_cond_destroy(&cond);
pthread_mutex_destroy(&mutex);
```

# **Key Points**

- Mutexes ensure that only one thread accesses a shared resource at a time, preventing race conditions.
- Condition Variables allow threads to wait for certain conditions and to signal other threads when those conditions change, facilitating coordination between threads.

# Conclusion

Threads play a crucial role in networking by enabling concurrent processing, improving performance, and enhancing responsiveness. The pthread library in C provides robust tools for managing threads, while mutexes and condition variables are essential for synchronization and coordination among threads. Understanding these concepts is fundamental for developing efficient and safe multithreaded applications in network programming.

