Database Administration

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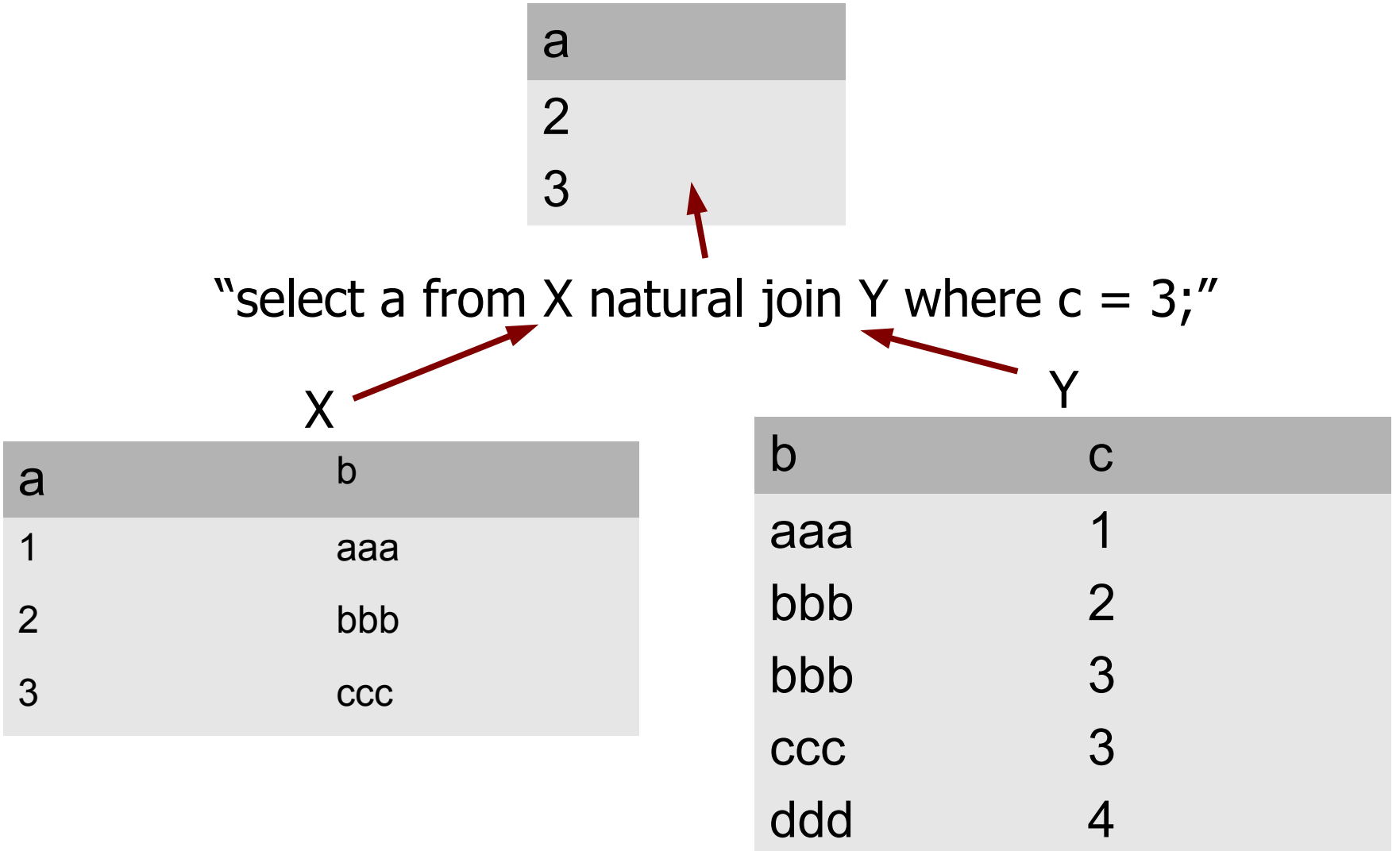
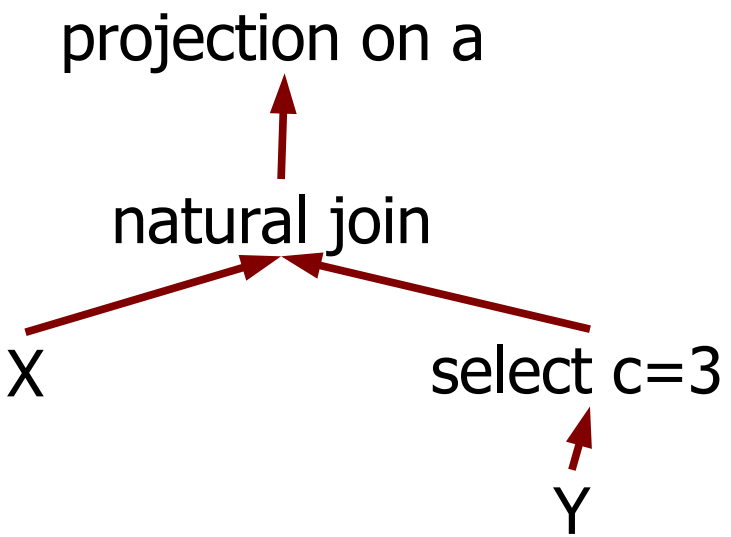
Query processing



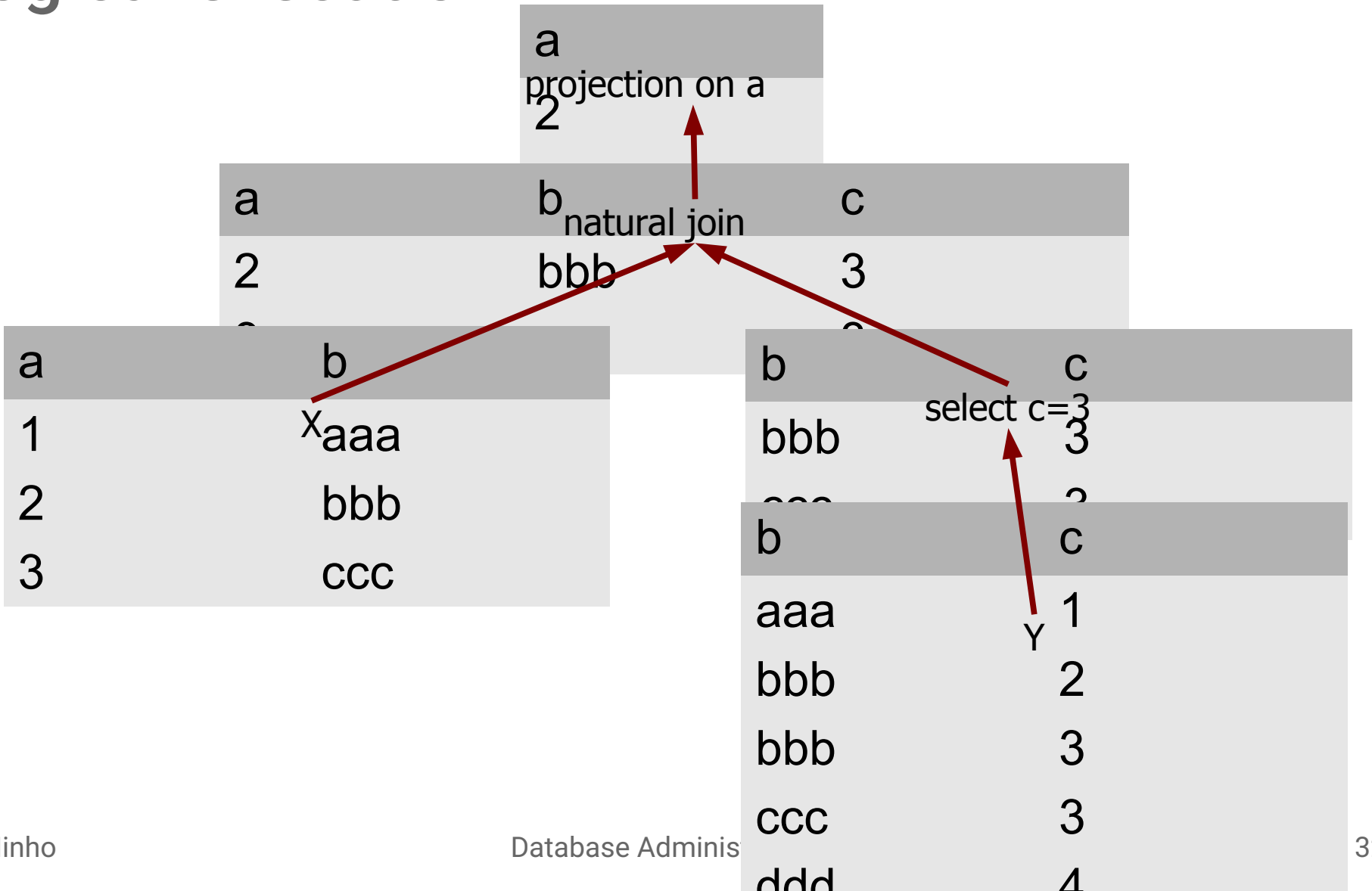
Compilation

SQL { "select a from X natural join Y where c = 3;"

Relational algebra {



Logical execution



Roadmap

- What physical operators exist for each logical operation?
- How are physical operators implemented and composed?
- Later: How are physical operators selected?

One-pass, record-at-a-time

- Operators:
  - Sequential scan
  - Selection
  - Projection
- Memory requirements:
  - No more than one record required
  - Always possible

User defined functions (UDFs)

- Functions can be defined in various languages
  - Python
- Scalar functions used in projections/selections:

  SELECT a, f1(a) FROM t;

  SELECT * FROM t WHERE f2(a)
- Table functions can be used in sequential scans:

  SELECT * FROM f3(…);
- User defined functions can access external services:
  - Web services
  - GenAI

One-pass, full relation, unary

- Duplicate elimination:
  - Cache unique records
  - “select distinct * from X;”
- Grouping and aggregation:
  - Cache groups
  - “select count(*) from X group by b;”
- Sorting:
  - Cache all records and sort in memory
  - “select * from X order by b;”

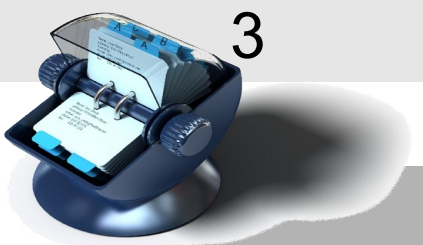
One-pass, full relation, binary

- Union, difference, intersection, product, join:
  - Read and cache the smallest relation
  - Organize for fast look-up (e.g. hash)
  - Read and operate on each record from the largest relation

One-pass, full relation, binary

- Load smaller table into memory and add search structure:

a	b	c
2	bbb	3
3	ccc	3



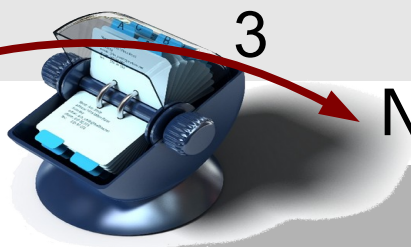
a	b
1	aaa
2	bbb
3	ccc

b	c
bbb	3
ccc	3

One-pass, full relation, binary

- Test each record from the largest relation:

a	b	c
2	bbb	3
3	ccc	3



Not found

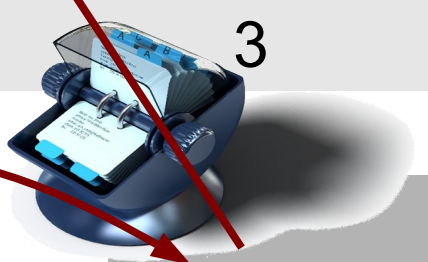
a	b
1	aaa
2	bbb
3	ccc

b	c
bbb	3
ccc	3

One-pass, full relation, binary

- Test each record from the largest relation:

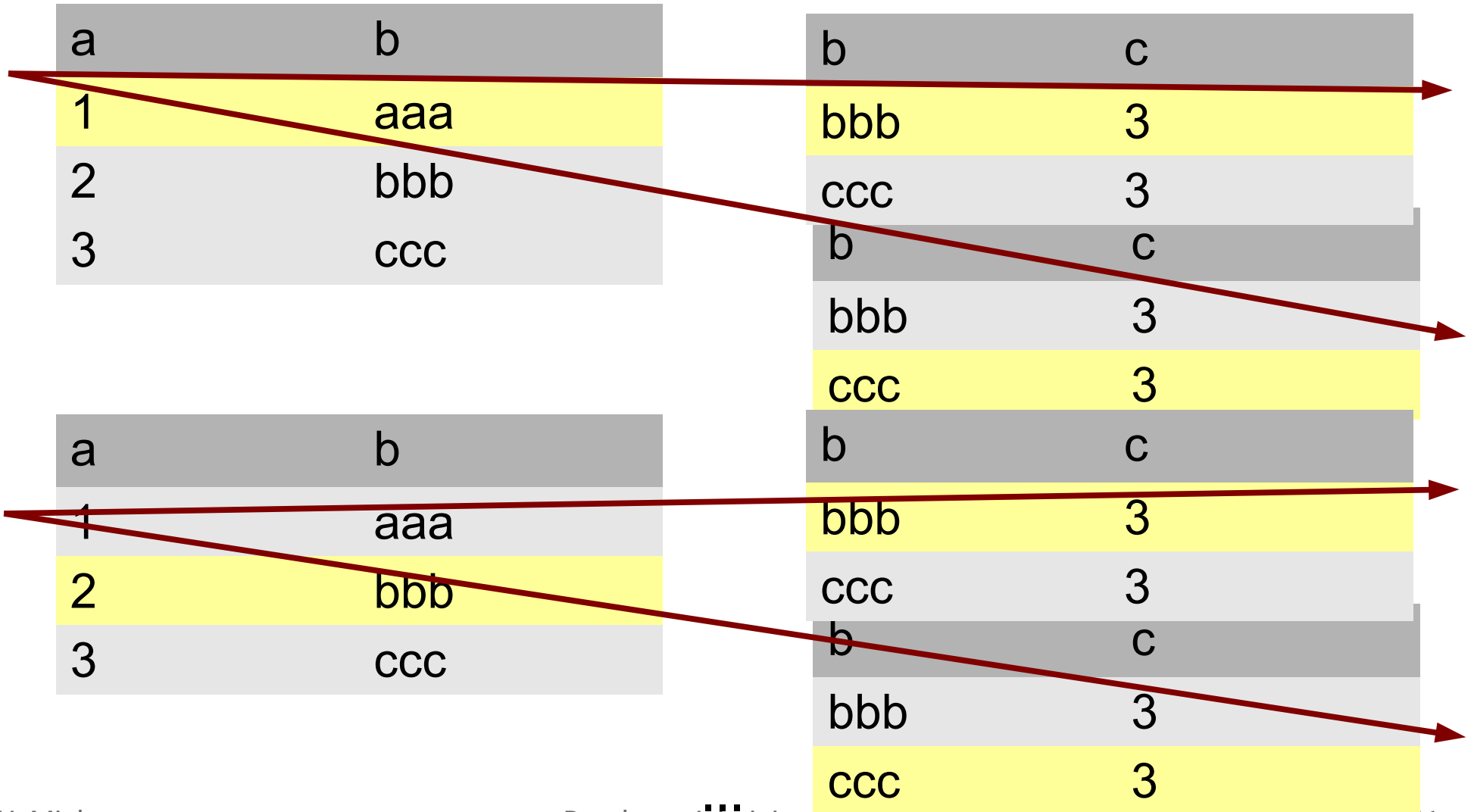
a	b	c
2	bbb	3
3	ccc	3



a	b
1	aaa
2	bbb
3	ccc

b	c
bbb	3
ccc	3

Nested-loop join (NLJ)



Nested-loop join (NLJ)

- Memory requirements:
  - One record from each relation
- Operations:
  - If outer loop has N records
  - Reads inner relation N times

Block-based NLJ

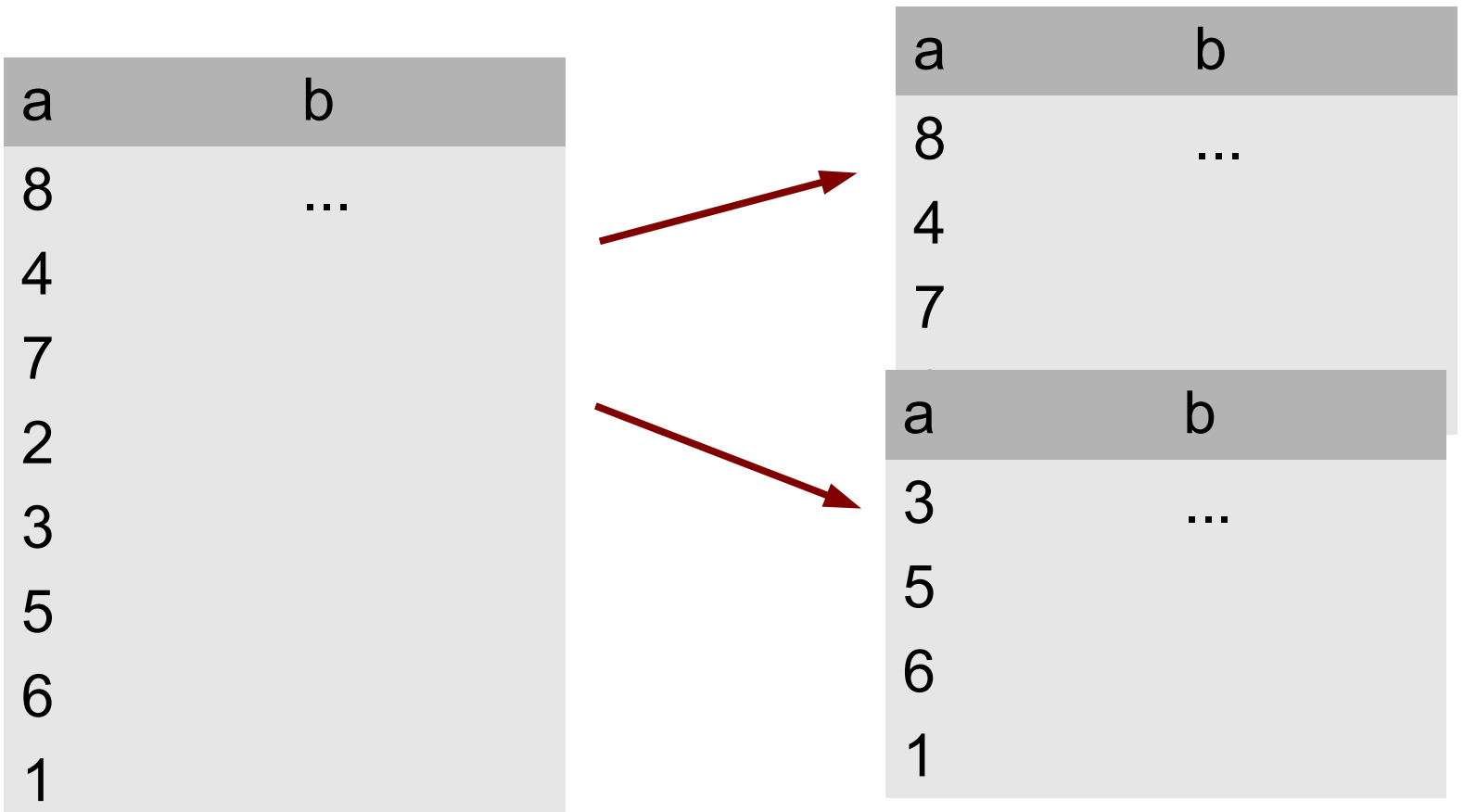
- Much smarter: Execute NLJ by blocks
- Memory requirements:
  - One block from each relation
- Operations:
  - If outer loop has N records / B blocks (B << N)
  - Reads inner relation B times (B << N!)

Large relations and sorting

- Algorithms using sorted data are more efficient (e.g. than nested loops)
- How to sort data that does not fit in memory?

Merge-sort

- Split data in chunks that fit in memory:



Merge-sort

- Load and sort each of them:

a	b
2	...
4	
7	
8	

a	b
3	...
5	
6	
1	

Merge-sort

- Load and sort each of them:

a	b
2	...
4	
7	
8	

a	b
1	...
3	
5	
6	

Merge-sort

- When iterating, select the next record from the fragment with the next key:

a	b
2	...
4	
7	
8	

a	b
1	...
3	
5	
6	

Merge-sort

- When iterating, select the next record from the fragment with the next key:

a	b
2	...
4	
7	
8	

a	b
1	...
3	
5	
6	

Merge-sort

- When iterating, select the next record from the fragment with the next key:

a	b
2	...
4	
7	
8	

a	b
1	...
3	
5	
6	

Two-pass, full relation, unary

- First pass is sorting
- Duplicate elimination:
  - Cache last record
  - “select distinct * from X;”
- Grouping and aggregation:
  - Cache last group
  - “select count(*) from X group by b;”

Two-pass, full relation, binary

- Union, difference, intersection, product, join:
  - Read record R1 from sorted relation T1
  - Read record R2 from sorted relation T2
  - If R1 = R2:
    - Use R1 and R2
  - If R1 < R2:
    - R1 does not exist in T2
    - Skip R1
  - If R2 < R1
    - ....