

Mastic: Private Weighted Heavy-Hitters and Attribute-Based Metrics

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Work at IETF

- Privacy Preserving Measurement (PPM)
 - Use cases
 - aggregation (à la <u>Prio</u>): average, variance, histograms, Bloom filters, linear regression, step of gradient descent
 - heavy hitters (à la <u>Poplar</u>): compute the set of the most frequent inputs (e.g., the most visited websites)
 - Privacy: distribute computation among multiple servers so that no one server sees any measurement in plaintext
 - Robustness: detect invalid measurements uploaded by misbehaving clients
- <u>Verifiable Distributed Aggregation Functions (VDAFs)</u>: Specifications of Prio and Poplar
- <u>Distributed Aggregation Protocol (DAP)</u>: Execution of a VDAF over HTTP



Where does Mastic fit?

- There are things we'd like Poplar or Prio to do:
 - attribute-based aggregation: average, variance, histograms, Bloom filters, linear regression, step of gradient descent grouped by client attributes (e.g., user agent or geolocation)
 - weighted heavy-hitters: compute the set of the most frequentheaviest weight inputs (e.g., the most visited websites websites with the highest engagement)
- Poplar has rough edges
 - Two rounds of communication for input validation (MPC multiplication w/ Beaver triples provided by clients)
 - Different field for leaf versus inner nodes of the IDPF tree, unsafe to use intermediate outputs



Design goals

- Ease of implementation: simplify design, reuse existing components where possible
- Deployability: one round of communication between aggregation servers instead of two
- Squeeze more functionality out of less code: attribute-based aggregation and weighted heavy-hitters



Agenda

- 1 Setting (functionality, architecture, security goals) ~10 minutes
- 2 Primitive #1: VIDPF ~15 minutes
- 3 Primitive #2: FLP ~3 minutes
- 4 Putting it all together ~3 minutes
- 5 Next steps ~5 minutes



Setting



- Each client generates a **measurement** consisting of an **input** (an element of $\{0,1\}^n$) and its **weight** (an element of \mathbb{F}^m)
- For each prefix, compute the total weight of all inputs beginning with the prefix



```
def weighted_heavy_hitters(measurements: list[tuple[Index, Weight]],
                           threshold: Weight,
                           bit_len: int) -> list[Index]:
    1000
   Compute the weighted heavy hitters for the given threshold.
   prefixes = [Index(0), Index(1)]
   for level in range(bit_len):
        next_prefixes = []
        for (x, total_weight) in mastic_func(measurements, prefixes).items():
            if total_weight >= threshold:
                if level < bit len - 1:
                    next_prefixes.append(x.left_child())
                    next_prefixes.append(x.right_child())
                else:
                    next\_prefixes.append(x)
        prefixes = next_prefixes
    return sorted(prefixes)
```

Weighted heavy hitters uses Mastic as a subroutine over multiple rounds of aggregation



Use case #1: Ad attribution

- On-device attribution: browser keeps track of which ads are shown to the user: when the user
 makes a purchase, the browser determines which ad, if any, deserves credit.
 - Input: a unique identifier for the ad campaign (e.g., "Bartlet for America")
 - Weight: whether a purchase was made (Count); how much the user spent (Sum); which category of product was purchased (Histogram); etc.



Use case #2: Browser telemetry

- Many use cases for Prio-style metrics in browsers, e.g.: page load metrics for a fixed set of "benchmark" websites (e.g., google.com or zombo.com)
 - Input: encodes whatever attributes we might want to break down by, eg., software version and geographic location
 - Weight: average (Sum) or distribution (Histogram) of load time



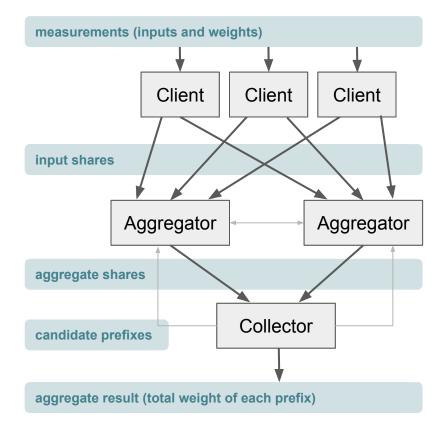
Use case #3: Network error logging

- <u>NEL</u>: To detect configuration issues, content-delivery networks like Cloudflare want to know the types of errors preventing eyeballs from connecting to their customers (DNS, TCP, TLS, etc.)
 - Input: the customer identified by a DNS name
 - Weight: indication of which error occurred (Histogram)



Architecture

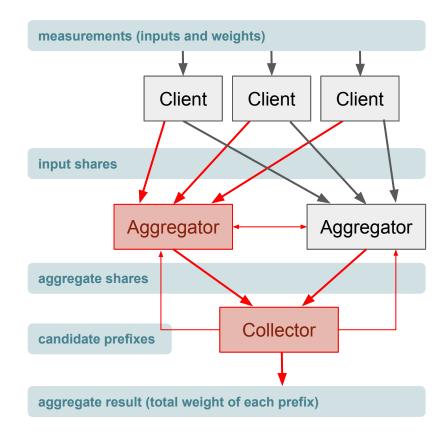
- Clients shard their measurements
- Aggregators interact to validate the measurements, then compute shares of the total weight for each candidate prefix
- Collector unshards the aggregate result





Security goal #1: Privacy

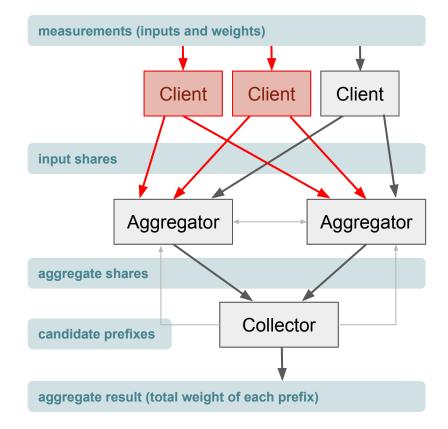
- Attacker controls one of the Aggregators and the Collector
- Attacker's view is efficiently simulatable given only aggregate results





Security goal #2: Robustness

- Attacker controls a subset of the clients and eavesdrops on the network
- Aggregate result computed by the collector is efficiently extractable from the input shares





Primitive #1: VIDPF



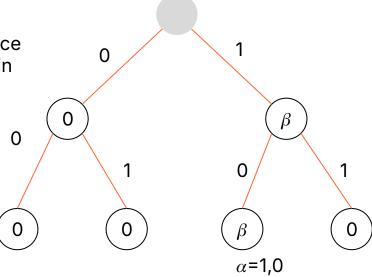
Verifiable Incremental Distributed Point Function (VIDPF)

• point function: $F(x) = \beta$ if $x = \alpha$ else 0

• incremental: $F(x) = \beta$ if x.is_prefix(α) else 0

• **distributed:** client generates secret shares of F such that each aggregator computes a secret share of F(x) for any x without learning α or β

• **verifiable:** aggregators verify that, for any sequence of prefixes prefixes, F(x)!=0 for at most one x in prefixes and the others are 0 (**onehot**)





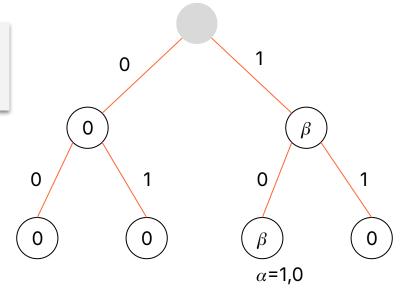
Verifiable Incremental Distributed Point Function (VIDPF)

```
Client: (correction_words, [key0, key1]) = vidpf.gen(\alpha, \beta)
```

Aggregator 0: $(out0, proof0) = vidpf.eval(0, correction_words, key0, prefixes)$

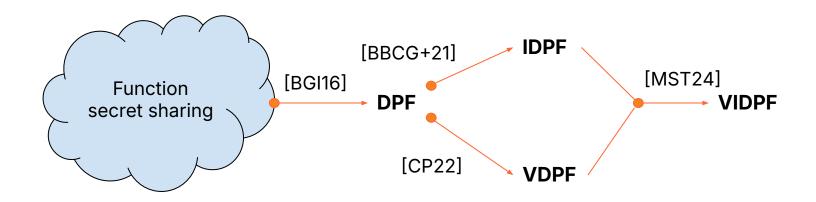
Aggregator 1: (out1, proof1) = vidpf.eval(1, correction_words, key1, prefixes)

verifiability: It's computationally infeasible to
find prefixes for which proof0 == proof1
but out0 + out1 is not onehot.





A bit of history



[BGI16] Boyle et al. "Function Secret Sharing: Improvements and Extensions." CCS 2016.

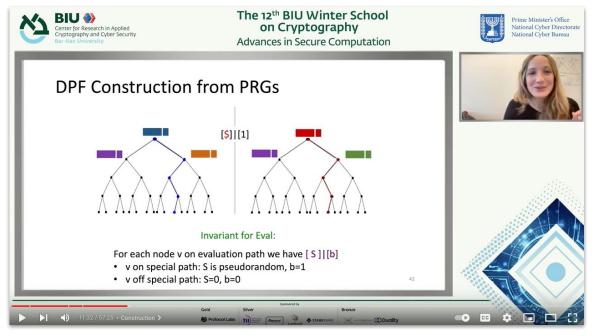
[BBCG+21] Boneh et al. "Lightweight Techniques for Private Heavy Hitters." IEEE S&P 2021.

[CP22] de Castro and Polychroniadou. "Lightweight, Maliciously Secure Verifiable Function Secret Sharing." Eurocrypt 2022.

[MST24] Mouris et al. "PLASMA: Private, Lightweight Aggregated Statistics against Malicious Adversaries." PETS 2024.



The DPF of [BGI16]

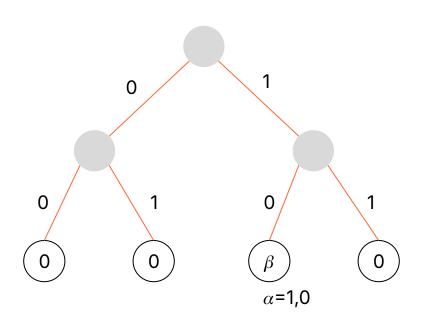


For a complete tutorial, check out <u>Elette Boyle's series</u> on function secret sharing for the 12th BIU Winter School



The DPF of [BGI16]

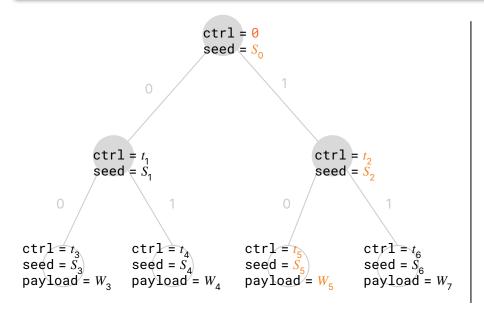
- Evaluation of input x:
 - Initialization:
 - ctrl = aggregator's ID (0 or 1)
 - seed = aggregator's key share
 - for each level i:
 - Extend the current seed into a ctrl/seed for each child node
 - If x[i] == 0 then select the left child;
 otherwise select the right child.
 - Share of the output is derived from the seed for the last level

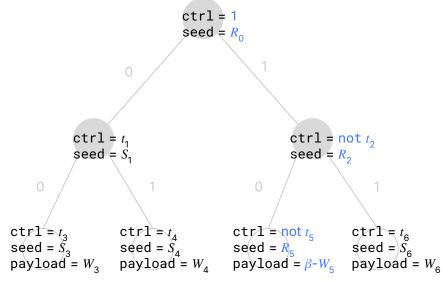




The DPF of [BGI16]

DPF invariant: if $x[:i].is_prefix(\alpha)$, then the ctrls are shares of 1 and the seeds are distinct; otherwise, the ctrls are shares of 0 and the seeds are the same. (If ctrl==1, then "correct" with the level's **correction word**.)

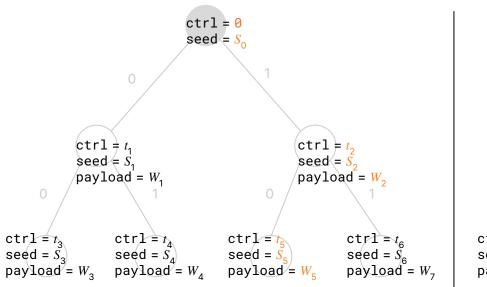


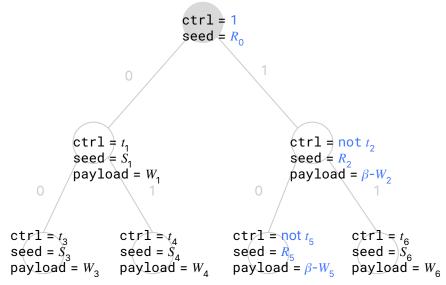




DPF → IDPF [BBCG+21]

Generate payloads for intermediate levels: use seed to derive the next seed and a share of the payload (after correction).

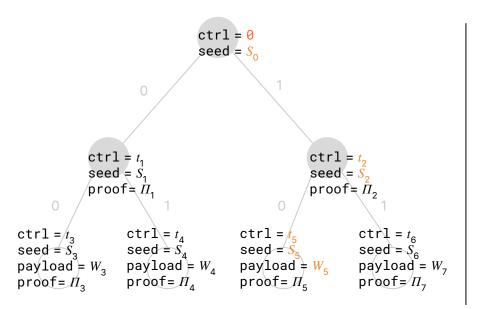


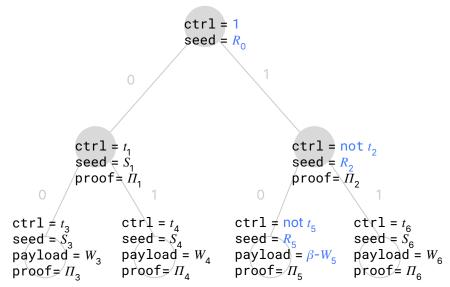




DPF → VDPF [CP22]

Verify the DPF invariant: let proof = hash(seed, x[:i]) be the **node proof**; if the DPF invariant holds, then each aggregator computes the same node proof (after correction).



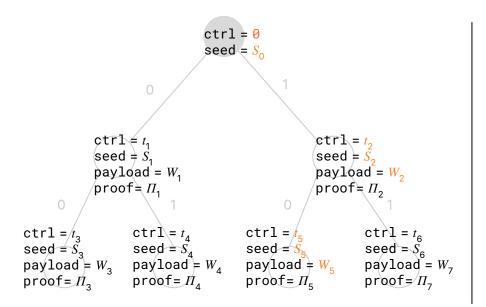


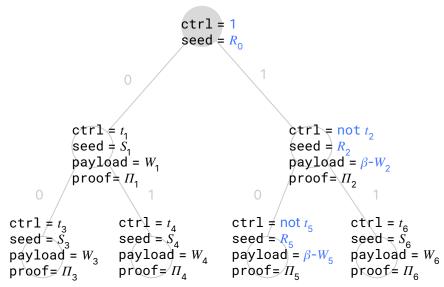
[CP22] de Castro and Polychroniadou. "Lightweight, Maliciously Secure Verifiable Function Secret Sharing." Eurocrypt 2022.



IDPF, VDPF → VIDPF [MST24]

Generate intermediate payloads, verify the DPF invariant, and verify **path consistency:** the payload of each node traversed is equal to the sum of its children $\Rightarrow \beta$ is the output for each prefix of α .



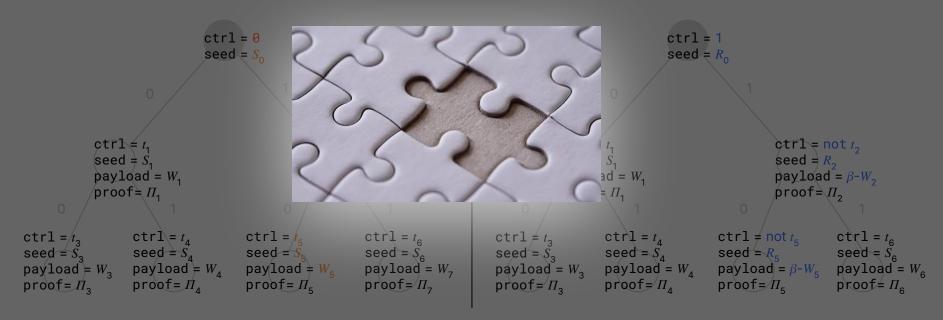


[MST24] Mouris et al. "PLASMA: Private, Lightweight Aggregated Statistics against Malicious Adversaries." PETS 2024.



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[MST24] Mouris et al. "PLASMA: Private, Lightweight Aggregated Statistics against Malicious Adversaries." PETS 2024.



Primitive #2: FLP



Fully Linear Proof (FLP)

Zero-knowledge proof system on distributed data [BBCG+19]. Syntax here is as presented in [draft-irtf-cfrg-vdaf]:

```
Client: proof = flp.prove(\beta); secret share proof into
```

proof0 + proof1 and β into beta0 + beta1

Aggregator 0: verifier0 = flp.query(beta0, proof0)

Aggregator 1: verifier1 = flp.query(beta1, proof1)

```
proof verification: if flp.decide(verifier0 + verifier1)
then the weight is valid with high probability.
```

[BBCG+19] Boneh et al. "Zero-Knowledge Proofs on Secret-Shared Data via Fully Linear PCPs". Crypto 2019.



Putting it all together



VIDPF + FLP =

Weight type is determined by the FLP:

- Count: weight is 0 or 1 (same functionality as Poplar)
- Sum: weight is an integer in a bounded range
- Histogram: weight is a histogram with a fixed number of buckets

Client runs the sharding algorithm to get the *public share* (the correction words) sent to both aggregators and one *input share* for each aggregator.

```
def shard(self,
          measurement: tuple[Index, Weight],
          nonce: bytes):
    # Encode the weight according to the FLP.
    (\alpha, \text{ weight}) = \text{measurement}
    \beta = self.flp.encode(weight)
    # Generate VIDPF keys (depends on the nonce).
    (correction_words, keys) = self.vidpf.gen(\alpha, \beta, nonce)
    # Generate FLP and split it into shares.
    proof = self.flp.prove(\beta)
    helper_seed = gen_rand(32)
    helper_proof_share = self.helper_proof_share(helper_seed)
    leader_proof_share = vec_sub(proof, helper_proof_share)
    return (correction words.
            [(keys[0], leader_proof_share), # aggregator 0
              (keys[1], helper_seed)])
                                              # aggregator 1
```



VIDPF + FLP =

Aggregator, on input of the report from the client (correcton words, VIDPF key share, FLP proof share) and prefixes from the collector:

- Evaluate the VIDPF at prefixes to get output output_share, the secret shares of F(x) for each x in prefixes
- Interact with the co-Aggregator to check (one roundtrip over the network)
 - **onehotness:** Check that DPF invariant holds for prefixes \Rightarrow at most one of the outputs is β (the rest are 0)
 - **path consistency:** For each node traversed in the prefix tree, check that the payload is equal to the sum of its children $\Rightarrow F(x) = \beta$ for each prefix x of α
 - **weight validity**: Verify the FLP $\Rightarrow \beta$ encodes a valid weight
- Aggregate output_share.



Next steps



Planned work

Status	Task
Done 🔽	Initial design and analysis [MPD+25].
In progress	Implementation (github.com/divviup/libprio-rs) in collaboration with ISRG (Divvi Up), who maintains open source implementations of Prio and Poplar (Rust).
In progress	Specification [draft-mouris-cfrg-mastic]. Many improvements are being considered (github.com/jimouris/draft-mouris-cfrg-mastic/issues).
Planned	End-to-end analysis of the final specification, especially concrete security of VIDPF.

[MPD+25] Mouris et al. "Mastic: Private Weighted Heavy-Hitters and Attribute-Based Metrics." PETS 2025.



Status of Mastic at IETF

- [draft-mouris-cfrg-mastic] is an *individual draft* and has not been adopted by a working group
- We will pursue adoption if/when there is sufficient interest to deploy Mastic. Currently in a holding pattern:
 - Top priority right now is completing the core documents [draft-ietf-ppm-dap, draft-irtf-cfrg-vdaf]
 - CFRG currently has low bandwidth for new work
 - PPM is chartered to do protocol design, not develop new cryptographic algorithms (this
 is meant to be delegated to CFRG)
 - We considered replacing Poplar with Mastic, but the <u>current consensus is to keep Poplar</u> in the core VDAF draft.

[draft-mouris-cfrg-mastic] Mouris et al. "The Mastic VDAF." Individual draft, version 03.

[draft-ietf-ppm-dap] Geoghegan et al. "Distributed Aggregation Protocol for Privacy Preserving Measurement." PPM working group draft, version 11.



Thank you

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datatracker.ietf.org/doc/draft-mouris-cfrg-mastic



prefix tree: Label each node with the 0 total weight of the prefix given by the path from the root 0 0 0

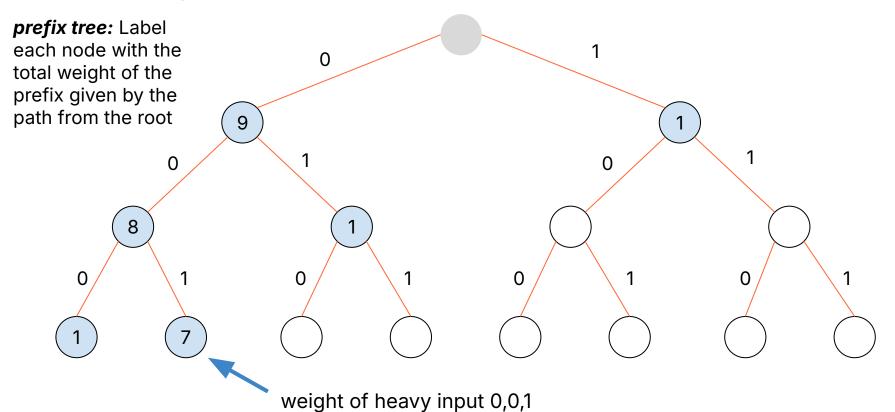


prefix tree: Label each node with the 0 total weight of the prefix given by the path from the root 0 0 0



prefix tree: Label each node with the 0 total weight of the prefix given by the path from the root 0 0 0







```
def prep_init(self,
              verify_key: bytes,
              agg_id: int,
              agg_param: MasticAggParam,
              nonce: bytes,
              correction_words: list[CorrectionWord],
              input_share: MasticInputShare.
              ) -> tuple[MasticPrepState, MasticPrepShare]:
    (level, prefixes, do_weight_check) = agg_param
    (key, proof_share) = \
        self.expand_input_share(agg_id, input_share)
    # Evaluate the VIDPF.
    (beta_share, out_share, eval_proof) = \
        self.vidpf.eval(agg_id, correction_words, key,
                        level, prefixes, nonce)
    # Query the FLP if applicable.
    verifier_share: Optional[list[F]] = None
    if do_weight_check:
        query_rand = self.query_rand(verify_key, nonce, level)
        verifier share = \
        self.flp.query(beta_share, proof_share, query_rand)
    prep_share = (eval_proof, verifier_share)
    return (out_share, prep_share)
```