

# Foundations of Algorithms

## Homework 5

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1. (a) Consider the following chain of six matrices:  $A_1$ ,  $A_2$ ,  $A_3$ ,  $A_4$ ,  $A_5$ , and  $A_6$ , where  $A_1$  is  $5 \times 10$ ,  $A_2$  is  $10 \times 3$ ,  $A_3$  is  $3 \times 12$ ,  $A_4$  is  $12 \times 5$ ,  $A_5$  is  $5 \times 50$ , and  $A_6$  is  $50 \times 6$ . Find an optimal parenthesization of this matrix-chain. Show both the table containing the optimal number of scalar operations for all slices and the choice table.  
(b) Prove using the strong form of induction that for any  $n \in \mathbb{N}$ , if  $n \geq 1$  then a full parenthesization of an  $n$ -element expression has  $n - 1$  pairs of parentheses.
2. (a) CLRS 15.3-1  
(b) Draw the recursion tree for the merge-sort algorithm on an input sequence of length 16. Explain why memoization fails to speed up a good divide-and-conquer algorithm like merge-sort.  
(c) Consider a variant of the matrix-chain multiplication problem in which the goal is to parenthesize the sequence of matrices so as to *maximize* the number of scalar multiplications. Does this problem exhibit optimal substructure?
3. **(project)** Recall  $F_n$  the recurrence that defines the Fibonacci numbers. Write a function `fibDyn` that computes Fibonacci numbers which implements the naive recurrence via dynamic programming.
4. Consider the 0-1 knapsack problem in CLRS chapter 16.  
(a) Write functional pseudo-code for a recursive solution to the variation on the 0-1 knapsack problem that computes the maximum value that can be placed in the knapsack. The first parameter is a sequence of items (i.e., value-weight pairs), and the second parameter is the knapsack weight capacity.  
(b) **(project)** Give a dynamic programming solution to the 0-1 knapsack problem that is based on the previous problem; this algorithm should take and return the same values as the functional pseudo-code above. Implement this algorithm and call it `knapsack`.  
(c) **(project)** Implement a function called `knapsackContents` that takes the same values but returns a sequence of items that maximize the knapsack's value.
5. Consider the problem of neatly printing a paragraph on the screen (or on a printer). For the project portion, put all functions in the file `printPar`. The input text is a sequence  $S$  of  $n$  words (represented as strings) of lengths  $\ell_1, \dots, \ell_n$  (measured in characters). The input bound  $M$  is the maximum number of characters a line can hold. The key to neatly printing a paragraph is to identify in the text sequence the lines of the paragraph so that new-lines can be placed at the end of each line. We can formalize the notion of the “badness” of a line as the number of extra space characters at the end of the line or  $\infty$  if the bound  $M$  is exceeded. We can formalize the notion of the “badness” of a paragraph as the badness of the worst (i.e., maximum) line of the paragraph not including the last line. Thus to identify the lines for a neat paragraph, we seek to minimize the badness of the paragraph.

- (a) If a given line contains words  $i$  through  $j$ , and we leave exactly one space between words, the number of extra space characters at the end of the line is  $M - j + i - \sum_{k=i}^j \ell_k$ . Write functional pseudo-code for the function  $e(S, M, i, j)$  that computes the number of extra space characters at the end of a line.
- (b) **(project)** Write a procedure `extraSpace` that implements  $e$ .
- (c) Use the function  $e$  to write functional pseudo-code for the function  $bl(S, M, i, j)$  that computes line badness.
- (d) **(project)** Write a procedure `badnessLine` that implements  $bl$ .
- (e) Write functional pseudo-code for the recursive function  $mb(S, M)$  that computes the minimum paragraph badness (using slicing). The base case must be that the sequence  $S$  is the last line (and *not* that  $S$  is empty).
- (f) Write functional pseudo-code for the recursive function  $mb'(S, M, i)$  where  $mb'(S, M, i) = mb(S[i:], M)$ . The base case must be that the sequence characterized by  $S$  and  $i$  is the last line (and *not* that  $S$  is empty).
- (g) **(project)** Write a recursive procedure `minBad` that implements the function  $mb'$ . It should take three parameters:  $S$ ,  $M$ , and  $i$  a slicing index.
- (h) **(project)** Write a procedure `minBadDynamic` that implements the function  $mb'$  using dynamic programming. It should take only two parameters:  $S$ , and  $M$ .
- (i) **(project)** Write a procedure `minBadDynamicChoice` that implements the function  $mb'$  using dynamic programming. In addition to returning  $mb(S, M)$ , it should also return the choices made. Then write a procedure `printParagraph` which takes two parameters:  $S$  and  $M$  that displays the words in  $S$  on the screen using the choices of `minBadDynamicChoice`. What is the asymptotic running time of your algorithm?

## 6. CLRS 24.1-1