

# Assignment 1

## Randomized Algorithms

Due: Monday 12<sup>th</sup>, by 11:55pm

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**Assignments are to be done individually. Submit a single pdf file to the classroom.**

**Problem 1** (2 points).

There are  $n$  new year greeting cards that you have written up for  $n$  different people. There are  $n$  envelopes for the  $n$  cards to go into. You give your nephew the task of putting one card into each envelope, in return for a gift. The nephew, too eager to get his gift randomly places each card into an empty envelope. What is the expected number of people who receive the card meant for them ?

**Problem 2** (2 points).

There are  $n$  new year greeting cards that you have written up for  $n$  different people. There are  $n$  envelopes for the  $n$  cards to go into. You give your nephew the task of putting one card into each envelope, in return for a gift. The nephew, too eager to get his gift randomly places each card into an empty envelope. What is the probability that exactly  $r$  of the cards are in their right envelope? As  $n \rightarrow \infty$  what is the probability that none of the cards are in the correct envelope?

**Problem 3** (2 points).

Design a randomized algorithm that finds the  $k^{th}$  smallest element in an array of  $n$  numbers. The expected running time of the algorithm should be  $O(n)$ .

**Problem 4** (2 points).

Suppose we have designed a new algorithm that can sort a given permutation of numbers. However, the algorithm is randomized, and we don't know if the output is indeed a permutation. Design a randomized algorithm to check if the output is indeed a permutation, for an input of size  $n$ , and a given parameter  $\epsilon > 0$ . The algorithm must succeed with probability at least  $\epsilon > 0$ . Assume that you have access to a prime number generator that outputs a prime larger than  $n/\epsilon$  and the elements of the input and output arrays in  $O(1)$  time.

**Problem 5.** Consider the following variant of the quicksort algorithm: We repeat the pivot step until we find a pivot such that the number of elements on either side of the partition has at most  $2/3$  fraction of the elements. That is, we pick a pivot uniformly at random and compute the partition. If either side of the array has more than  $2/3$  fraction of the elements, we go back to picking a pivot uniformly at random. This process continues until we find a pivot such that either side has at most  $2/3$  fraction of the elements. What is the expected running time of this algorithm?