

Declaration: Few results and derivations are referred from Trefethen & Bau, Numerical Linear Algebra; the primary reference for this course.

1. $A, B \in \mathbb{C}^{m \times m}$ $C = AB$

$$\left. \begin{array}{l} a_1 \geq \dots \geq a_m \\ b_1 \geq \dots \geq b_m \\ c_1 \geq \dots \geq c_m \end{array} \right\} \text{ singular values}$$

Before proving what is asked, we will prove some lemmas we'll use.

Lemma 1: $c_i \leq a_i b_1$ & $c_i \leq a_1 b_i$

Proof: By using min-max characterization of singular values, we have

$$\begin{aligned} c_i &= \max_{S: \dim(S)=i} \min_{x \in S, \|x\|=1} \|ABx\| \\ &\leq \|A\| \cdot \max_{S: \dim(S)=i} \min_{x \in S, \|x\|=1} \|Bx\| \\ &= a_1 b_i \end{aligned}$$

$$\Rightarrow c_i \leq a_1 b_i \quad \text{--- (a)}$$

$\therefore \sigma_i(AB) = \sigma_i(BA)$ [where σ denotes singular value]
[\because they have identical spectra]

$$\text{Similarly } \Rightarrow c_i \leq a_i b_1 \quad \text{--- (b)}$$

From (a) & (b),

$$\Rightarrow \boxed{c_i \leq \min(a_1 b_i, a_i b_1)}$$

1. Lemma 2: $c_m \geq a_m b_m$

Proof: Note that the m^{th} singular value is the smallest/minimum value of singular value for those particular matrices.

$$\begin{aligned}
 c_m = c_{\min} &= \min_{x \neq 0} \frac{\|ABx\|_2}{\|x\|_2} \\
 &= \min_{x \neq 0} \frac{\|ABx\|_2}{\|Bx\|_2} \cdot \frac{\|Bx\|_2}{\|x\|_2} \\
 &\geq \min_{y \neq 0} \frac{\|Ay\|_2}{\|y\|_2} \min_{x \neq 0} \frac{\|Bx\|_2}{\|x\|_2} \\
 &= a_m b_m
 \end{aligned}$$

$$\Rightarrow \boxed{c_m \geq a_m b_m}$$

Lemma 3: $c_1 \geq a_1 b_m$, $c_1 \geq a_m b_1$

Proof: Note that the 1st singular value is the the largest/maximum value of singular value for those particular matrices.

$$\begin{aligned}
 c_1 = c_{\max} &= \max_{x \neq 0} \frac{\|ABx\|_2}{\|x\|_2} \\
 &= \max_{x \neq 0} \frac{\|ABx\|_2}{\|Bx\|_2} \cdot \frac{\|Bx\|_2}{\|x\|_2} \\
 &\geq \max_{y \neq 0} \frac{\|Ay\|_2}{\|y\|_2} \max_{x \neq 0} \frac{\|Bx\|_2}{\|x\|_2} \\
 &\geq \max_{y \neq 0} \frac{\|Ay\|_2}{\|y\|_2} \min_{x \neq 0} \frac{\|Bx\|_2}{\|x\|_2} \\
 &= a_1 b_m
 \end{aligned}$$

1. Contd-3

$$\therefore c_1 \geq a, b_m$$

$$\text{Illy } c_1 \geq a m b_1$$

$$\Rightarrow \boxed{c_1 \geq \max(a, b_m, a m b_1)}$$

Now,

$$\text{Using lemma 1, } a, b_1 \geq c_1$$

$$\text{Using lemma 3, } c_1 \geq \max(a, b_m, a m b_1)$$

$$\Rightarrow \boxed{a, b_1 \geq c_1 \geq \max(a, b_m, a m b_1)}$$

Again,

$$\text{Using lemma 1, } \min(a, b_m, a m b_1) \geq c_m$$

$$\text{Using lemma 2, } c_m \geq a m b_m$$

$$\Rightarrow \boxed{\min(a, b_m, a m b_1) \geq c_m \geq a m b_m}$$

Hence, proved.

2. $S \cap T = \{0\}$
 $S + T = \mathbb{C}^m$
 $S \perp T$

We know that $I-P$ is a complementary projector of P such that $\text{subspaces of } (I-P) \cap (P) = \{0\}$
 $(\because \text{range}(P) \cap \text{null}(P) = \{0\})$.

Let $u = Px \in S$
 and $v = (I-P)y \in T \in \mathbb{C}^m$

Since $S \perp T$, ~~$\langle u, v \rangle = 0$~~

From the orthogonal projectors' definition, we know that for S & T to be orthogonal $P^* = P$ (From definition of Bau) (Also below)

$$u^* v = x^* P^* (I-P)y = x^* (P - P^2)y = 0$$

($\because P = P^2$ for projectors)

$\therefore T$ is the set $\{v \in \mathbb{C}^m : u^* v = 0 \ \forall u \in S\}$

Proof of $P^* = P$ for orthogonal subspaces:

Let $\{q_1, q_2, \dots, q_m\}$ be orthonormal basis for \mathbb{C}^m
 where $\{q_1, \dots, q_n\}$ is basis for S & $\{q_{n+1}, \dots, q_m\}$ is basis for T . For $j \leq n$, we have $Pq_j = q_j$
 & for $j > n$ we have $Pq_j = 0$.

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2. Now let Q be the unitary matrix whose j th column is q_j . We then have

$$PQ = \left[q_1 \mid \dots \mid q_n \mid 0 \mid \dots \right]$$

$$Q^* PQ = \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ & & 1 & & \\ & & & 0 & \ddots \\ & & & & \ddots \end{bmatrix} = \Sigma$$

first n -entries are 1

$$P = Q \Sigma Q^*$$

$$P^* = (Q \Sigma Q^*)^* = Q \Sigma^* Q^* = Q \Sigma Q^* = P$$

$$\Rightarrow \boxed{P^* = P}$$

3. (a) Let us look at this problem geometrically. (6)

$\|x_i\|_2 = \|y_i\|_2$ denotes that the vectors have same length for each i .

$\|x_i - x_j\|_2 = \|y_i - y_j\|_2$ denotes that the vectors have same angle between them for each i & j .

For any full column rank matrix $A = [a_1, \dots, a_n]$
and $P_u a = \frac{\langle u, a \rangle}{\langle u, u \rangle} u$ (P is projection)

Using Gram-Schmidt process, we have

$$\hat{R} = \begin{pmatrix} \langle e_1, a_1 \rangle & \langle e_1, a_2 \rangle & \langle e_1, a_3 \rangle & \dots \\ 0 & \langle e_2, a_2 \rangle & \langle e_2, a_3 \rangle & \dots \\ 0 & 0 & \langle e_3, a_3 \rangle & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

where $e_k = \frac{u_k}{\|u_k\|}$ & $u_k = a_k - \sum_{j=1}^{k-1} \text{proj}_{u_j} a_k$

$$\& a_k = \sum_{j=1}^k \langle e_j, a_k \rangle e_j$$

e_k denote the orthonormal basis of A .
 $\forall k=1 \text{ to } n$

Now, for x & y ,

$$\hat{R}_x = \begin{pmatrix} \langle e_1^n, x_1 \rangle & \langle e_1^n, x_2 \rangle & \langle e_1^n, x_3 \rangle \dots \\ 0 & \langle e_2^n, x_2 \rangle & \langle e_2^n, x_3 \rangle \dots \\ 0 & 0 & \langle e_3^n, x_3 \rangle \dots \\ \vdots & \vdots & \vdots \end{pmatrix}$$

Similarly for y .

inner product denotes dot product or how much of the first vector points towards the second vector.

$$\langle a, b \rangle = \|a\|_2 \|b\|_2 \cos \theta$$

$$\langle e_i, x_j \rangle = \|e_i\|_2 \|x_j\|_2 \cos \theta_{ij} = \|x_j\|_2 \cos \theta_{ij}$$

where θ_{ij} is the angle b/w $\|e_i\|_2$ & $\|x_j\|_2$

$$\langle e_i, x_j \rangle = \|x_j\|_2 \cos \theta_{ij} = \|y_j\|_2 \cos \theta_{ij} = \langle e_i, y_j \rangle \quad \text{--- (1)}$$

(since they have same lengths & angles between them)

\therefore All the entries in \hat{R}_x & \hat{R}_y are equal by (1).

$\therefore x$ & y have same \hat{R} .

3. (b)

$$Qx_i = y_i$$

$$\Rightarrow Q = y_i x_i^{-1}$$

$$\Rightarrow Q = \frac{y_i x_i^*}{x_i^* \cdot x_i}$$

$$\Rightarrow Q = \frac{y_i x_i^*}{\|x_i\|^2}$$

$$\left[\begin{array}{l} VV^{-1} = I \\ V^* V^{-1} = V^* \\ V^{-1} = \frac{V^*}{V^* V} \end{array} \right]$$

$$r = \|x_i\| = \|x_i\|$$

For $m = 1$ to i For $n = 1$ to i

$$q_{mn} = y_{mi}^* x_{in} / r$$

Assuming that the vectors are non-zero (otherwise Q could be anything trivially) vectors have full rank.

$\Rightarrow Q$ & R in the QR factorisation of a vector are both invertible.

Now, let $x_i = Qx_R$ & $y_i = Qy_R$

1. $QQ^*R = Qy_R$ (calculating using gram-schmidt)

2. $QQ^*R = Qy_R$ ($\because R$ is invertible)

3. $Q = Qy_R Q^*{}^{-1}$ ($\because Q^* & Qy_R$ are invertible)

Algorithm.

1. calculate Q_n & Q_y using modified gram-schmidt.

2. calculate Q_n^{-1}

3. Find $Q = Q_y Q_n^{-1}$

Here, Q will be orthogonal since both Q_y & Q_n^{-1} are orthogonal & multiplication of orthogonal matrices is orthogonal.

4. $A \in \mathbb{C}^{m \times n}$

$v \in \mathbb{C}^m$

(a)

$$F = I - \frac{2vv^*}{v^*v}$$

$$FA = \left[I - \frac{2vv^*}{v^*v} \right] A$$

$$= A - \frac{2vv^*A}{v^*v}$$

$$= A - v \left[\frac{2v^*A}{v^*v} \right]$$

Comparing it with $FA = A + vw^*$, we see that

$$w^* = \frac{2v^*A}{v^*v} \quad \& \quad w \in \mathbb{C}^n$$

$\therefore w$ is a vector.

(i) computing F then performing matrix multiplication.

(a) $F = I \ominus \frac{2vv^*}{v^*v}$

\downarrow $\frac{2vv^*}{v^*v}$ $\rightarrow m^2$ multiplications (flops)

\downarrow v^*v $\rightarrow m$ multiplications (flops)

m^2 subtractions $m-1$ additions (flops)

$$\sim 2m^2 + 3m - 1 \text{ operations (flops)}$$

(b) FA $F \in \mathbb{C}^{m \times m}$ $A \in \mathbb{C}^{m \times n}$

$$\sim (2m-1)mn \text{ operations (flops)}$$

(a) + (b) $= \sim 2m^2n - mn + 2m^2 + 3m - 1$

$\sim 2m^2n$ operation count (flops)

(ii) computing w , then vw^* then matrix addition

(a) $w = \frac{2(V^*A)}{V^*V} \rightarrow n(2m-1) \text{ operations (flops)}$
 m multiplications
 $m-1$ additions

$$\sim (2m-1)n + (2m-1)$$

$$\sim (2m-1)(n+1) \text{ operations}$$

(b) vw^* $v \in \mathbb{C}^m$ $w^* \in \mathbb{C}^{1 \times n}$

$$\sim mn \text{ operations (flops)}$$

(c) Add A & vw^* of size $m \times n$
 $\sim mn \text{ operations (flops)}$

$$\begin{aligned} (a) + (b) + (c) &= (2m-1)(n+1) + mn + mn \\ &= 2mn + 2m - n - 1 + mn + mn \\ &= 4mn + 2m - n - 1 \end{aligned}$$

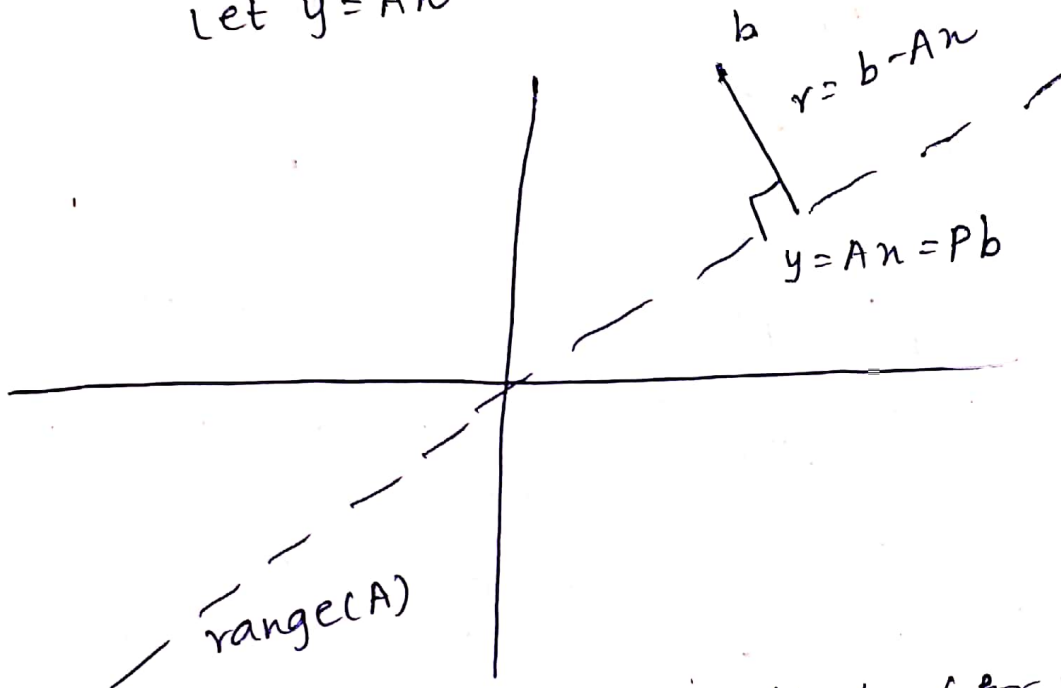
$$\boxed{\sim 4mn \text{ operations}} \text{ (flops)}$$

Method (ii) is more efficient in terms of time complexity.

5. $Ax = b$ $A \in \mathbb{C}^{m \times n}$
 $m > n$
 $\text{rank}(A) = r < n$
 SVD of A $A = U \Sigma V^*$

(a) $y \in \text{range}(A)$

Let $y = Ax$



Geometrically, we can see that (for $r = b - Ax$)
 $\|r\|_2 = \|b - Ax\|_2$ is minimized only when
 $r \perp \text{range}(A)$

$\Rightarrow A^* r = 0$

From the properties of orthogonal projectors, we know that we can have P st $Pb = Ax$ & $P \in \mathbb{C}^{m \times m}$ is the orthogonal projector onto $\text{range}(A)$.

$\therefore y = Pb$ \rightarrow The explicit formula

5. (b) From the previous problem,

$$\begin{aligned}
 r &= b - An \\
 A^* r &= 0 \\
 A^* (b - An) &= 0 \\
 A^* b - A^* A n &= 0 \\
 A^* A n &= A^* b \\
 n &= (A^* A)^{-1} A^* b
 \end{aligned}$$

$$\begin{aligned}
 A^* r &= 0 \\
 r &= b - An \\
 A^* (b - An) &= 0 \\
 A^* b - A^* A n &= 0 \\
 A^* A n &= A^* b \\
 n &= (A^* A)^{-1} A^* b
 \end{aligned}$$

It is given by $z = A^+ b + (I - A^+ A) w$

where $A^+ = (A^* A)^{-1} A^*$ & w is an arbitrary n -dimensional vector.

From above, $A^+ b$ is already a solution. If we can show that $(I - A^+ A) w$ gives every vector in $K(A)$ & no other then, we have all solutions given by z .

Let $k = (I - A^+ A) w$, then:

$$Ak = A(I - A^+ A) w = (A - AA^+ A) w = (A - A) w = 0$$

\therefore Every vector of the form $(I - A^+ A) w$ is a member of $K(A)$

suppose k is a vector in the kernel of A , st $Ak = 0$. Then, as $(I - A^+ A)k = (k - A^+ A k) = (k - A^+ 0) = k$, every vector in $K(A)$ may be written as $(I - A^+ A) w$

$z = A^+ b$ minimizes $\|x\|_2$.

6. (a)

$$p(x) = \sum_{k=0}^n p_k x^k$$

$$q(x) = \sum_{j=0}^n q_j x^j$$

$$(p(x), q(x)) = \sum_{k,j=0}^n \bar{p}_k q_j (x^k, x^j) \quad \text{--- (1)}$$

$$\text{From} \quad (p, q) = [p]^* G [q] \quad \text{--- (2)}$$

From (1) & (2),
we have G of the form

$$G = \begin{bmatrix} (x^1, x^1) & (x^2, x^1) & \dots & (x^n, x^1) \\ \vdots & \ddots & & \vdots \\ (x^1, x^n) & & & (x^n, x^n) \end{bmatrix}$$

$$g_{mn} = (x^m, x^n)$$

$$g_{mn}^* = (x^n, x^m)$$

$$\text{Since } (x^m, x^n) = (x^n, x^m)$$

$$G = G^*$$

Hence G is hermitian.

6. (b) It means finding an orthogonal 143
 projection of q in the subspace $\langle p \rangle$
 If it is r , r is given by

$$r = \frac{(q, p)}{(p, p)} p$$

where $(p, q) = \int_{-1}^1 \overline{p(x)} q(x) dx$

Since r is just a projector, an orthogonal projector, (from Ref. 2 Bau)

$$r = (p(p^* p)^{-1} p^*) q$$

\therefore The matrix that acts on q is $(p(p^* p)^{-1} p^*)$.

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4. (b)

$$\tilde{F} = I - 2 \frac{\tilde{V} \tilde{V}^*}{\tilde{V}^* \tilde{V}}$$

$$\frac{\|\tilde{V} - V\|_2}{\|V\|_2} = O(\epsilon_m)$$

$$\tilde{V} = V + \delta V$$

$$\frac{\|\delta V\|_2}{\|V\|_2} = O(\epsilon_m)$$