Telco Dt.: Pg.:

Experiment-4

AIM:- Wap to write merge sort and find its complexity

Software used :- Code blocks

Merge sort algo :-

```
MERGE_SORT (A, b, z)
    if b < z
        q = FLOOR[(b+z)/2]
        merge_sort [A, b, q]
        Merge_sort [A, q+1, z]
        Merge (A, b, q, z)

Merge (A, b, q, z)
n1 = b-q+1
n2 = z-q
create arrays L[1... n1+1] and R[1... n2+1]
for i=1 to n1
    L[i] = A[b+i-1]
for j=1 to n2
    R[j] = A[q+j]
L1[n1+1] = ∞
L2[n2+1] = ∞
i=1
j=1
```

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```
for K = p to r
    if L[i] <= R[j]
        A[K] = L[i]
        i = i + 1
    else
        A[K] = R[j]
        j = j + 1
```

$$T(n) = \begin{cases} O(1) & \text{if } n = 1 \\ 2T(n/2) + n & \text{if } n > 1 \end{cases}$$

Since in each call to mergeSort two calls are made again to mergeSort that divides the array into two parts

↳ merge fun^n which iterate over the loop.

Using master's Theorem $a = 2, b = 2$

$T(n) = a\,T(n/b) + f(n)$

$n^{\log_b a} = n^{\log_2 2} = n$

$f(n) = n$ $\therefore$ $n^{\log_b a} = f(n)$

Using 2nd case of master's theorem

$T(n) = \Theta(n \log n)$