Data Warehousing and OLAP

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Warehousing

- Growing industry: \$8 billion in 1998
- Range from desktop to huge:
 - Walmart: 900-CPU, 2,700 disk, 23TB
 Teradata system
- Lots of buzzwords, hype
 - slice & dice, rollup, MOLAP, pivot, ...

Outline

- What is a data warehouse?
- Why a warehouse?
- Models & operations
- Implementing a warehouse
- Future directions

What is a Warehouse?

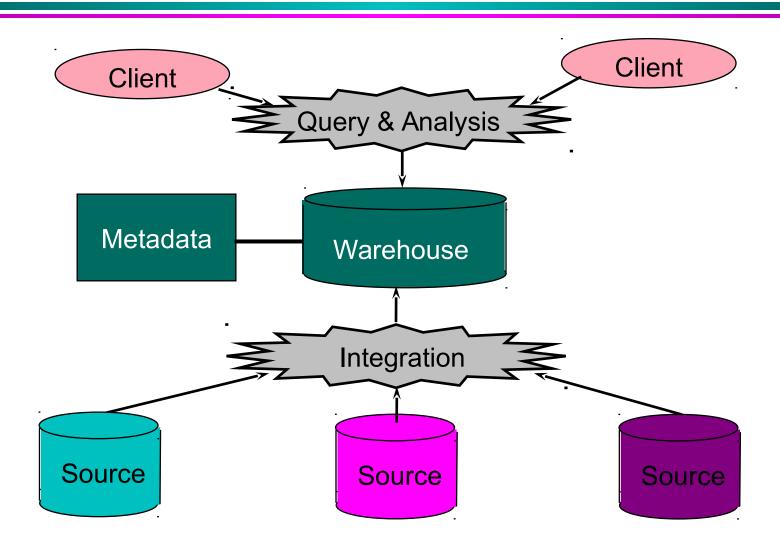
- Collection of diverse data
 - subject oriented
 - aimed at executive, decision maker
 - often a copy of operational data
 - with value-added data (e.g., summaries, history)
 - integrated
 - time-varying
 - non-volatile



What is a Warehouse?

- Collection of tools
 - gathering data
 - cleansing, integrating, ...
 - querying, reporting, analysis
 - data mining
 - monitoring, administering warehouse

Warehouse Architecture

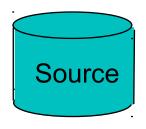


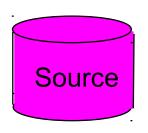
Why a Warehouse?

- Two Approaches:
 - Query-Driven (Lazy)
 - Warehouse (Eager)

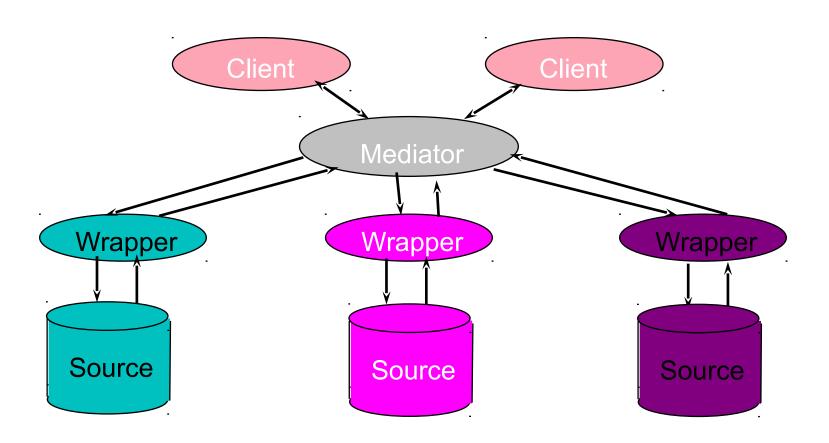








Query-Driven Approach



Advantages of Warehousing

- High query performance
- Queries not visible outside warehouse
- Local processing at sources unaffected
- Can operate when sources unavailable
- Can query data not stored in a DBMS
- Extra information at warehouse
 - Modify, summarize (store aggregates)
 - Add historical information

Advantages of Query-Driven

- No need to copy data
 - less storage
 - no need to purchase data
- More up-to-date data
- Query needs can be unknown
- Only query interface needed at sources
- May be less draining on sources

OLTP vs. OLAP

- OLTP: On Line Transaction Processing
 - Describes processing at operational sites
- OLAP: On Line Analytical Processing
 - Describes processing at warehouse

OLTP vs. OLAP

OLTP

- Mostly updates
- Many small transactions
- Mb-Tb of data
- Raw data
- Clerical users
- Up-to-date data
- Consistency, recoverability critical

OLAP

- Mostly reads
- Queries long, complex
- Gb-Tb of data
- Summarized, consolidated data
- Decision-makers, analysts as users

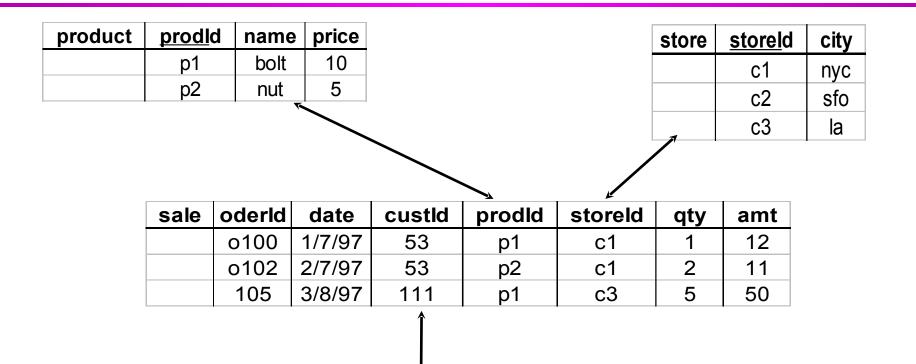
Data Marts

- Smaller warehouses
- Spans part of organization
 - e.g., marketing (customers, products, sales)
- Do not require enterprise-wide consensus
 - but long term integration problems?

Warehouse Models & Operators

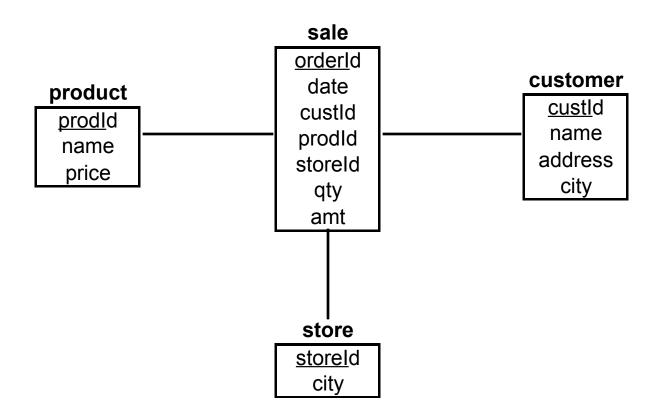
- Data Models
 - relations
 - stars & snowflakes
 - cubes
- Operators
 - slice & dice
 - roll-up, drill down
 - pivoting
 - other

Star



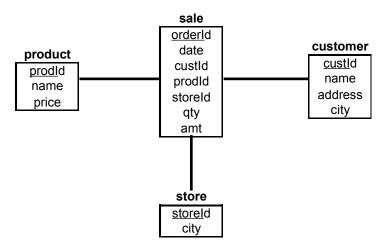
customer	<u>custl</u> d	name	address	city
	53	joe	10 main	sfo
	81	fred	12 main	sfo
	111	sally	80 willow	la

Star Schema



Terms

- Fact table
- Dimension tables
- Measures



Dimension Hierarchies



store	<u>storel</u> d	cityld	tld	mgr
	s5	sfo	t1	joe
	s7	sfo	t2	fred
	s9	la	t1	nancy

sType	<u>tl</u> d	size	location
	t1	small	downtown
	t2	large	suburbs

city	<u>cityl</u> d	pop	regld
	sfo	1M	north
	la	5M	south

- → snowflake schema
- → constellations

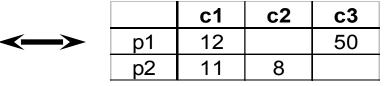
region	<u>regl</u> d	name
	north	cold region
	south	warm region

Cube

Fact table view:

sale	prodld	storeld	amt
	p1	c1	12
	p2	c1	11
	p1	с3	50
	p2	c2	8

Multi-dimensional cube:



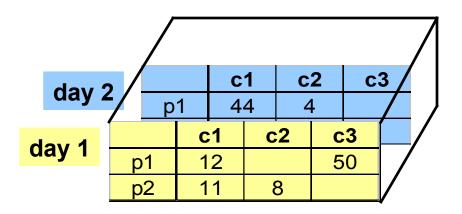
dimensions = 2

3-D Cube

Fact table view:

sale	prodld	storeld	date	amt
	p1	c1	1	12
	p2	c1	1	11
	p1	c3 c2	1	50
	p2	c2	1	8
	p1	c1	2	44
	p1	c2	2	4

Multi-dimensional cube:



dimensions = 3

ROLAP vs. MOLAP

- ROLAP:
 Relational On-Line Analytical Processing
- MOLAP: Multi-Dimensional On-Line Analytical Processing

Aggregates

- Add up amounts for day 1
- In SQL: SELECT sum(amt) FROM SALE WHERE date = 1

sale	prodld	storeld	date	amt
	p1	c1	1	12
	p2	c1	1	11
	p1	с3	1	50
	p2	c2	1	8
	p1	c1	2	44
	p1	c2	2	4



81

Aggregates

- Add up amounts by day
- In SQL: SELECT date, sum(amt) FROM SALE GROUP BY date

sale	prodld	storeld	date	amt
	p1	c1	1	12
	p2	c1	1	11
	p1	с3	1	50
	p2	c2	1	8
	p1	c1	2	44
	p1	c2	2	4



ans	date	sum
	1	81
	2	48

Another Example

- Add up amounts by day, product
- In SQL: SELECT date, sum(amt) FROM SALE GROUP BY date, prodId

sale	prodld	storeld	date	amt
	p1	c1	1	12
	p2	c1	1	11
	p1	c3	1	50
	p2	c2	1	8
	p1	c1	2	44
	p1	c2	2	4



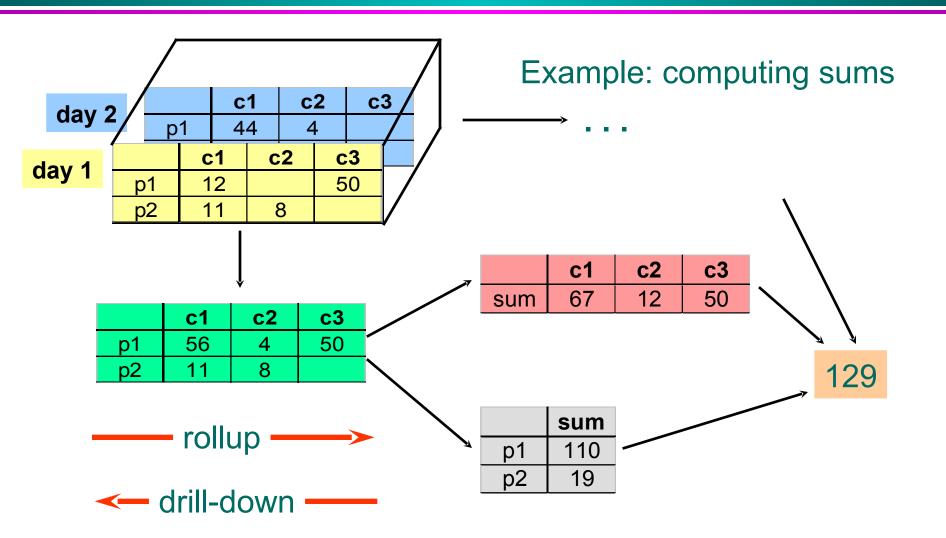
sale	prodld	date	amt
	p1	1	62
	p2	1	19
	p1	2	48



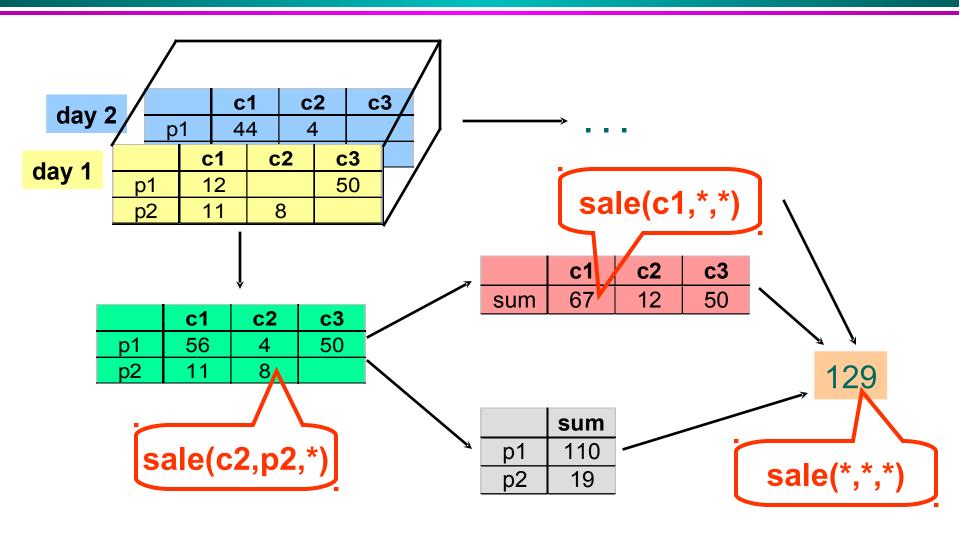
Aggregates

- Operators: sum, count, max, min, median, ave
- "Having" clause
- Using dimension hierarchy
 - average by region (within store)
 - maximum by month (within date)

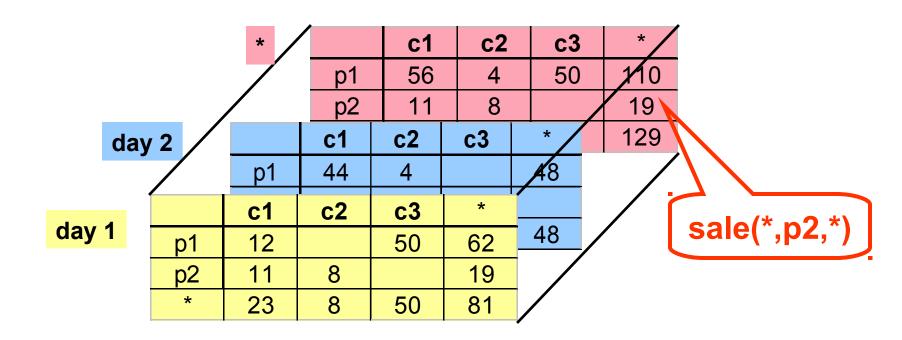
Cube Aggregation



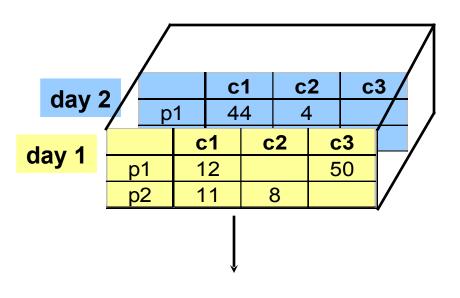
Cube Operators



Extended Cube



Aggregation Using Hierarchies





	region A	region B
p1	56	54
p2	11	8

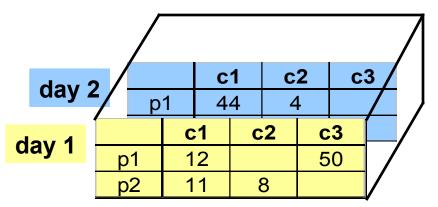
(customer c1 in Region A; customers c2, c3 in Region B)

Pivoting

Fact table view:

sale	prodld	storeld	date	amt
	p1	c1	1	12
	p2	c1	1	11
	p1	c3	1	50
	p2	c2	1	8
	p1	c1	2	44
	p1	c2	2	4

Multi-dimensional cube:





	с1	c2	с3
p1	56	4	50
p2	11	8	

Implementing a Warehouse

- Monitoring: Sending data from sources
- Integrating: Loading, cleansing,...
- Processing: Query processing, indexing, ...
- Managing: Metadata, Design, ...

Monitoring

- Source Types: relational, flat file, IMS, VSAM, IDMS, WWW, news-wire, ...
- Incremental vs. Refresh

customer	<u>i</u> d	name	address	city
	53	joe	10 main	sfo
	81	fred	12 main	sfo
	111	sally	80 willow	la



Disadvantages Advantages

Monitoring Techniques

- Periodic snapshots
- Database triggers
- Log shipping
- Data shipping (replication service)
- Transaction shipping
- Polling (queries to source)
- Screen scraping
- Application level monitoring

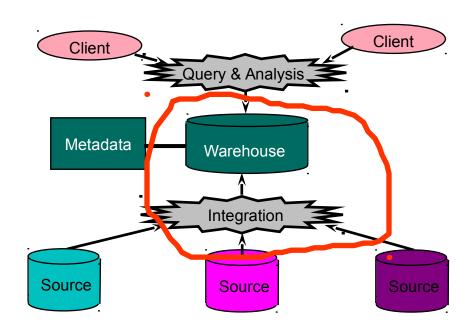


Monitoring Issues

- Frequency
 - periodic: daily, weekly, ...
 - triggered: on "big" change, lots of changes, ...
- Data transformation
 - convert data to uniform format
 - remove & add fields (e.g., add date to get history)
- Standards (e.g., ODBC)
- Gateways

Integration

- Data Cleaning
- Data Loading
- Derived Data



Data Cleaning

- Migration (e.g., yen ⇒ dollars)
- Scrubbing: use domain-specific knowledge (e.g., social security numbers)
- Fusion (e.g., mail list, customer merging)

```
billing DB → customer1(Joe) ← merged_customer(Joe) service DB → customer2(Joe)
```

 Auditing: discover rules & relationships (like data mining)

Loading Data

- Incremental vs. refresh
- Off-line vs. on-line
- Frequency of loading
 - At night, 1x a week/month, continuously
- Parallel/Partitioned load

Derived Data

- Derived Warehouse Data
 - indexes
 - aggregates
 - materialized views (next slide)
- When to update derived data?
- Incremental vs. refresh

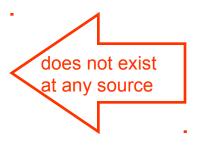
Materialized Views

Define new warehouse relations using SQL expressions

sale	prodld	storeld	date	amt
	p1	c1	1	12
	p2	c1	1	11
	p1	c3	1	50
	p2	c2	1	8
	p1	c1	2	44
	p1	c2	2	4

product	id	name	price	
	p1	bolt	10	
	p2	nut	5	

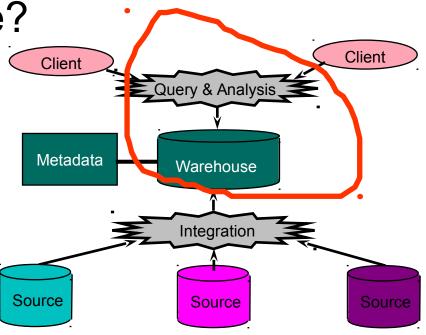
joinTb	prodld	name	price	storeld	date	amt
	p1	bolt	10	c1	1	12
	p2	nut	5	c1	1	11
	p1	bolt	10	с3	1	50
	p2	nut	5	c2	1	8
	p1	bolt	10	c1	2	44
	p1	bolt	10	c2	2	4



Processing

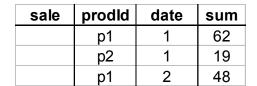
- ROLAP servers vs. MOLAP servers
- Index Structures
- What to Materialize?

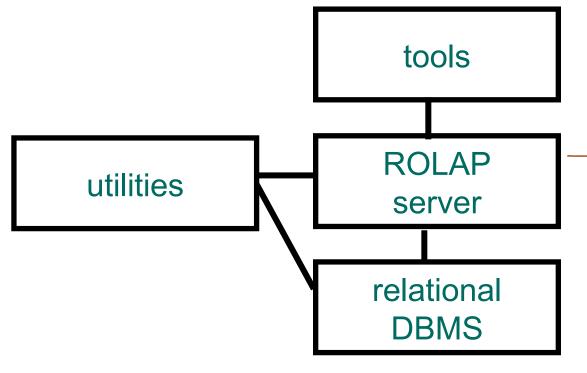
Algorithms



ROLAP Server

Relational OLAP Server

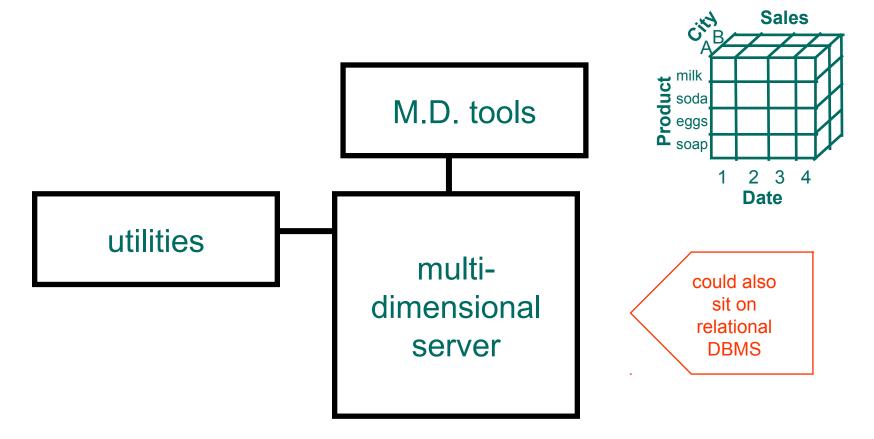




Special indices, tuning;
Schema is "denormalized"

MOLAP Server

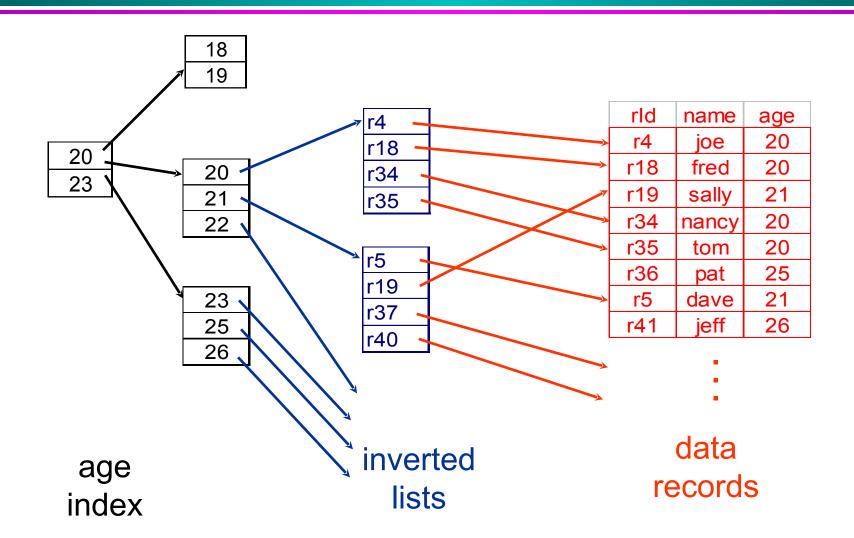
• Multi-Dimensional OLAP Server



Index Structures

- Traditional Access Methods
 - ◆ B-trees, hash tables, R-trees, grids, ...
- Popular in Warehouses
 - inverted lists
 - bit map indexes
 - join indexes
 - text indexes

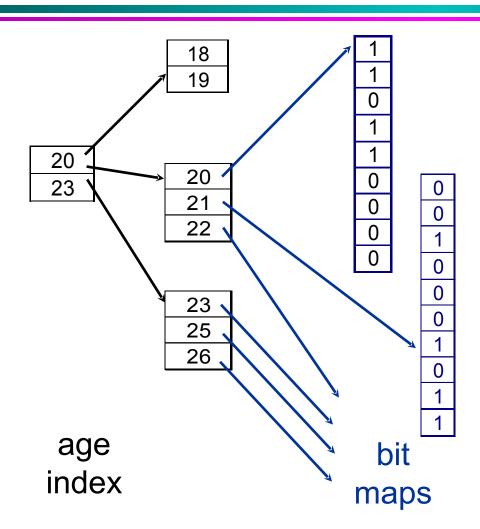
Inverted Lists



Using Inverted Lists

- Query:
 - Get people with age = 20 and name = "fred"
- List for age = 20: r4, r18, r34, r35
- List for name = "fred": r18, r52
- Answer is intersection: r18

Bit Maps



name	age
joe	20
fred	20
sally	21
nancy	20
tom	20
pat	25
dave	21
jeff	26
	joe fred sally nancy tom pat dave

-

•

data records

Using Bit Maps

- Query:
 - Get people with age = 20 and name = "fred"
- List for age = 20: 1101100000
- List for name = "fred": 0100000001
- Answer is intersection: 01000000000
- Good if domain cardinality small
- Bit vectors can be compressed

Join

- "Combine" SALE, PRODUCT relations
- In SQL: SELECT * FROM SALE, PRODUCT

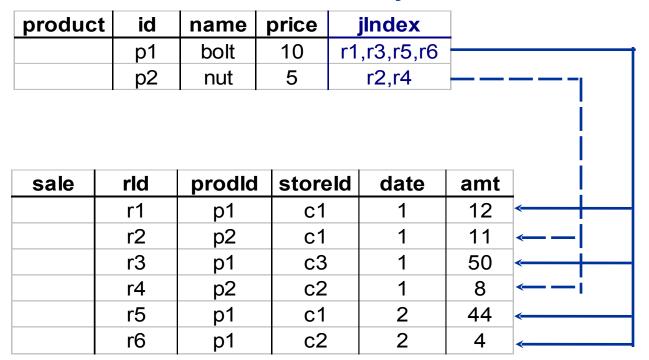
sale	prodld	storeld	date	amt
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product	id	name	price
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joinTb	prodld	name	price	storeld	date	amt
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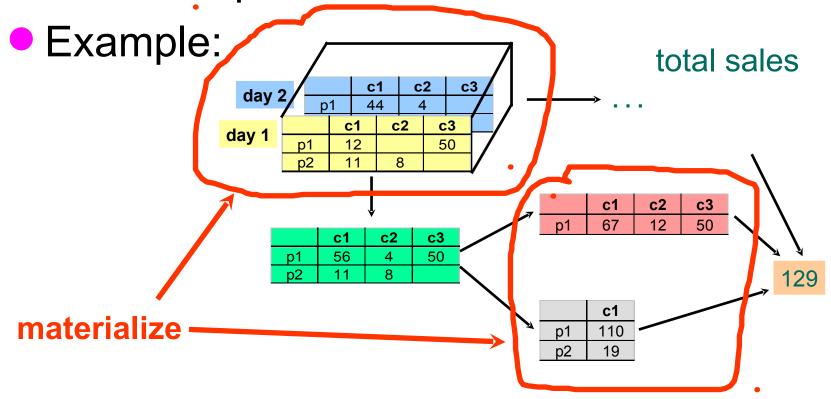
Join Indexes

join index



What to Materialize?

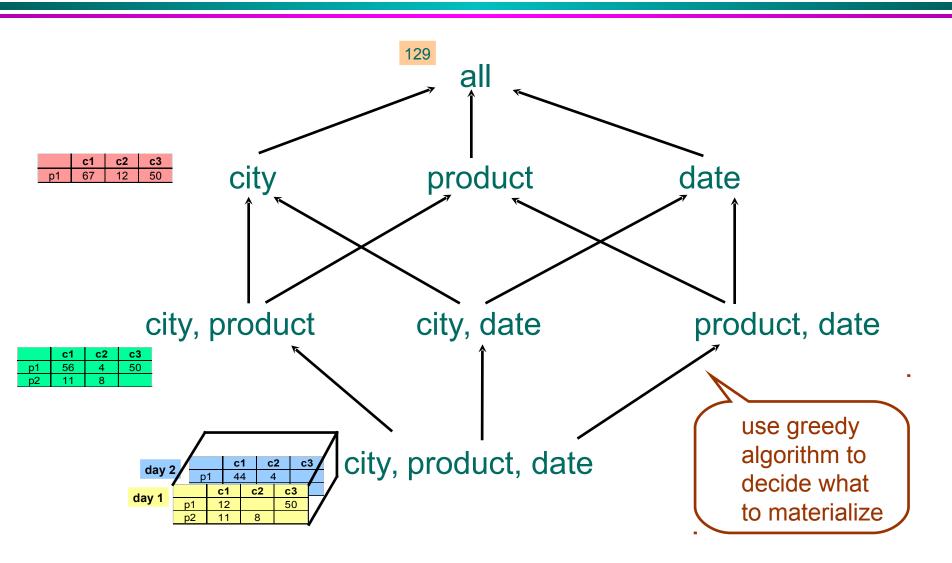
Store in warehouse results useful for common queries



Materialization Factors

- Type/frequency of queries
- Query response time
- Storage cost
- Update cost

Cube Aggregates Lattice

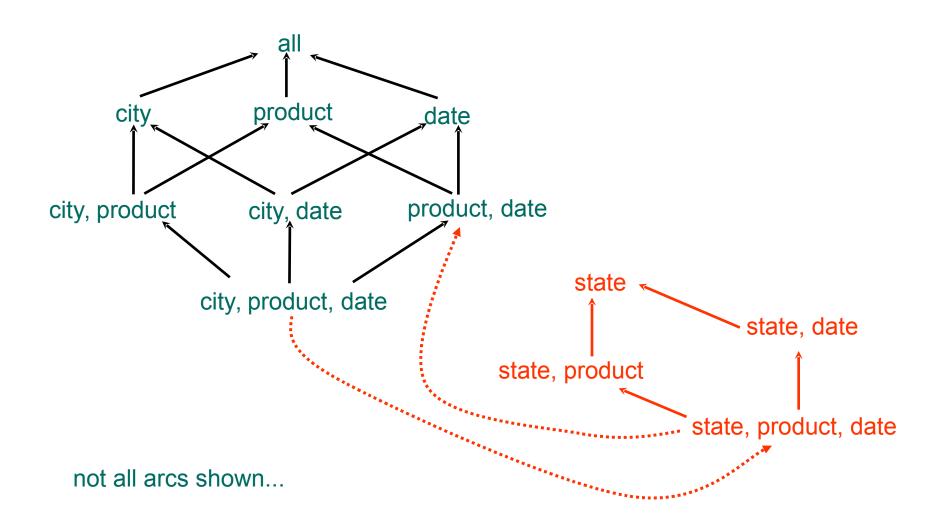


Dimension Hierarchies

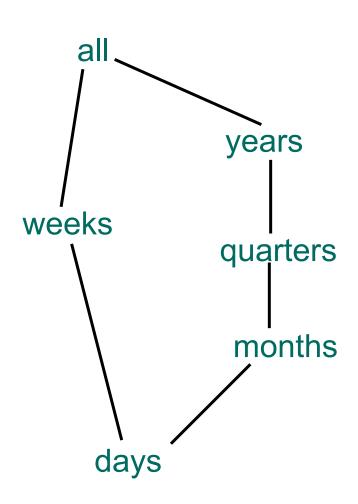


cities	city	state
	c1	CA
	c2	NY

Dimension Hierarchies



Interesting Hierarchy



time	day	week	month	quarter	year
	1	1	1	1	2000
	2	1	1	1	2000
	3	1	1	1	2000
	4	1	1	1	2000
	5	1	1	1	2000
	6	1	1	1	2000
	7	1	1	1	2000
	8	2	1	1	2000

conceptual dimension table

Design

- What data is needed?
- Where does it come from?
- How to clean data?
- How to represent in warehouse (schema)?
- What to summarize?
- What to materialize?
- What to index?

Tools

- Development
 - design & edit: schemas, views, scripts, rules, queries, reports
- Planning & Analysis
 - what-if scenarios (schema changes, refresh rates), capacity planning
- Warehouse Management
 - performance monitoring, usage patterns, exception reporting
- System & Network Management
 - measure traffic (sources, warehouse, clients)
- Workflow Management
 - "reliable scripts" for cleaning & analyzing data

Current State of Industry

- Extraction and integration done off-line
 - Usually in large, time-consuming, batches
- Everything copied at warehouse
 - Not selective about what is stored
 - Query benefit vs storage & update cost
- Query optimization aimed at OLTP
 - High throughput instead of fast response
 - Process whole query before displaying anything