Digital Logic Circuits 'Boolean Algebra' ELEC2200 Summer 2009

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Boolean Algebra

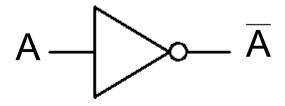
- Developed by George Boole of Bell Labs (1849)
- Originally termed Switching Algebra as it described the math of relay circuits
- Classical Algebra manipulates expressions comprised of real-valued variables and joined by operators such as: Add, Subtract, Multiply, Divide
- Boolean Algebra manipulates expressions comprised of binary-valued variables and joined by operators such as: And, Or, Not
- Other operators can be built from these



Boolean Attributes (Operators)

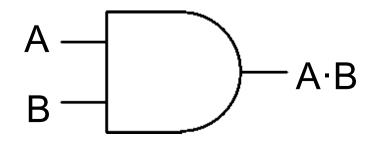
- Complement, Invert,
 Not
- Two notations: A' and

Α	A
0	1
1	0



- And
- Notated as A-B

Α	В	A·B
0	0	0
0	1	0
1	0	0
1	1	1



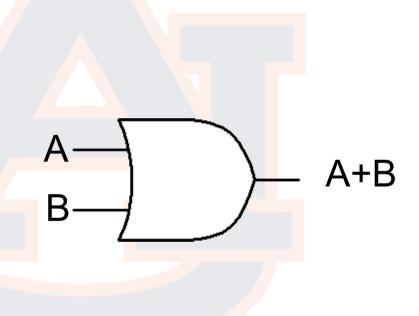


Boolean Attributes (Operators)

Or

Notated as A+B

Α	В	A+B
0	0	0
0	0 1 1	
1	0	1
1	1	1





Operator Precedence

```
Order:
(1)Parentheses
(2)Not
(3)And
(4)Or
```

(5)Left to Right

When in doubt, use parentheses.



Identity

$$A+0=A$$

 $A\cdot 1=A$

Complements

$$A+\overline{A}=1$$
 $A\cdot\overline{A}=0$

Α	В	A+0
0	0	0
0	1 1	
1	0	1
1	1	1

Α	В	A·1
0	0	0
0	1 0	
1	0	0
1	1	1

A	В	A+A
0	0	0
0 1 1		1
1	0	1
1	1	1

Α	В	$A \cdot \overline{A}$
0	0	0
0	1	0
1	0	0
1	1	1



Commutativity

$$A+B=B+A$$

$$A \cdot B = B \cdot A$$

Α	В	A+B	B+A
0	0	0	0
0	1	1	1
1	0	1	1
1	1	1	1

A	В	A·B	B·A
0	0	0	0
0	1	0	0
1	0	0	0
1	1	1	1



Associativity

$$(A+B)+C=A+(B+C)$$

 $(A \cdot B) \cdot C = A \cdot (B \cdot C)$

Α	В	С
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A+B
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A+B	(A+B)+C
0	0	0	0	0
0	0	1	0	1
0	1	0	1	1
0	1	1	1	1
1	0	0	1	1
1	0	1	1	1
1	1	0	1	1
1	1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A+B	(A+B)+C	B+C
0	0	0	0	0	0
0	0	1	0		1
0	1	0	1	1	1
0	1	1	1	1	1
1	0	0	1	1	0
1	0	1	1	1	1
1	1	0	1	1	1
1	1	1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A+B	(A+B)+C	B+C	A+(B+C)
0	0	0	0	0	0	0
0	0	1	0		1	1
0	1	0	1	1	1	1
0	1	1	1	1	1	1
1	0	0	1	1	0	1
1	0	1	1	1	1	1
1	1	0	1	1	1	1
1	1	1	1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A+B	(A+B)+C	B+C	A+(B+C)
0	0	0	0	0	0	0
0	0	1	0	1	1	1
0	1	0	1	1	1	1
0	1	1	1	1	1	1
1	0	0	1	1	0	1
1	0	1	1	1	1	1
1	1	0	1	1	1	1
1	1	1	1	1	1	1



Associativity

$$(A+B)+C=A+(B+C)$$

 $(A \cdot B) \cdot C = A \cdot (B \cdot C)$

Α	В	С
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A·B
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A·B	(A·B)·C
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	0	0
1	0	0	0	0
1	0	1	0	0
1	1	0	1	0
1	1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A· <mark>B</mark>	(A·B)·C	B·C
0	0	0	0	0	0
0	0	1	0	0	0
0	1	0	0	0	0
0	1	1	0	0	1
1	0	0	0	0	0
1	0	1	0	0	0
1	1	0	1	0	0
1	1	1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A· <mark>B</mark>	(A·B)·C	B·C	A·(B·C)
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	0	1	0
1	0	0	0	0	0	0
1	0	1	0	0	0	0
1	1	0	1	0	0	0
1	1	1	1	1	1	1



$$(A+B)+C=A+(B+C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

Α	В	С	A·B	(A·B)·C	B·C	A·(B·C)
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	0	1	0
1	0	0	0	0	0	0
1	0	1	0	0	0	0
1	1	0	1	0	0	0
1	1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

		1 [
Α	В	C	A+B
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A+B	A+C
0	0	0	0	0
0	0	1	0	1
0	1	0	1	0
0	1	1	1	1
1	0	0	1	1
1	0	1	1	1
1	1	0	1	1
1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A+B	A+C	(A+B)·(A+C)
0	0	0	0	0	0
0	0	1	0	1	0
0	1	0	1	0	0
0	1	1	1	1	1
1	0	0	1	1	1
1	0	1	1	1	1
1	1	0	1	1	1
1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A+B	A+C	(A+B)·(A+C)	B·C
0	0	0	0	0	0	0
0	0	1	0	1 _	0	0
0	1	0	1	0	0	0
0	1	1	1	1	1	1
1	0	0	1	1	1	0
1	0	1	1	1	1	0
1	1	0	1	1	1	0
1	1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A+B	A+C	(A+B)·(A+C)	B·C	A+(B·C)
0	0	0	0	0	0	0	0
0	0	1	0	1 4	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	1	1	1	1
1	0	0	1	1	1	0	1
1	0	1	1	1	1	0	1
1	1	0	1	1	1	0	1
1	1	1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A+B	A+C	(A+B)·(A+C)	B·C	A+(B·C)
0	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	1	1	1	1
1	0	0	1	1	1	0	1
1	0	1	1	1	1	0	1
1	1	0	1	1	1	0	1
1	1	1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A·B
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A·B	A·C
0	0	0	0	0
0	0	1	0	0
0	1	0	0	0
0	1	1	0	0
1	0	0	0	0
1	0	1	0	1
1	1	0	1	0
1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A·B	A·C	(A·B)+ (A·C)
0	0	0	0	0	0
0	0	1	0	0	0
0	1	0	0	0	0
0	1	1	0	0	0
1	0	0	0	0	0
1	0	1	0	1	1
1	1	0	1	0	1
1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A·B	A·C	(A·B)+	B+C
					(A·C)	
0	0	0	0	0	0	0
0	0	1	0	0	0	1
0	1	0	0	0	0	1
0	1	1	0	0	0	1
1	0	0	0	0	0	0
1	0	1	0	1	1	1
1	1	0	1	0	1	1
1	1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

Α	В	С	A·B	A·C	(A·B)+ (A·C)	В+С	A·(B+C)
0	0	0	0	0	0	0	0
0	0	1	0	0	0	1	0
0	1	0	0	0	0	1	0
0	1	1	0	0	0	1	0
1	0	0	0	0	0	0	0
1	0	1	0	1	1	1	1
1	1	0	1	0	1	1	1
1	1	1	1	1	1	1	1



$$A+(B\cdot C)=(A+B)\cdot (A+C)$$

$$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$$

A	В	С	A·B	A·C	(A·B)+ (A·C)	B+C	A·(B+C)
0	0	0	0	0	0	0	0
0	0	1	0	0	0	1	0
0	1	0	0	0	0	1	0
0	1	1	0	0	0	1	0
1	0	0	0	0	0	0	0
1	0	1	0	1	1	1	1
1	1	0	1	0	1	1	1
1	1	1	1	1	1	1	1



Boolean Theorems

Idempotency

$$A + A = A$$

$$A \cdot A = A$$

Null Elements

$$A+1=1$$

$$A \cdot 0 = 0$$

Involution

$$\overline{A}=(A')'=A$$

Α	A+A	A·A	A+1	A·0	Ā	Ā
0	0	0	1	0	1	0
1	1	1	1	0	0	1



Boolean Theorems

Absorption (Covering)

$$A \cdot (A+B)=A$$

$$A+(A\cdot B)=A$$

Α	В	A+B	A·B	A·(A+B)	A+(A·B)
0	0	0	0	0	0
0	1	1	0	0	0
1	0	1	0	1	1
1	1	1	1	1	1



Boolean Theorems

Absorption (Covering)

$$A \cdot (\overline{A} + B) = A \cdot B$$

 $A + (\overline{A} \cdot B) = A + B$

A	В	Ā	A+B	$A \cdot (\overline{A} + B)$	A·	A·B	$A+(\overline{A}\cdot B)$	A+B
					В			
0	0	1	0	0	0	0	0	0
0	1	1	1	0	0	1	1	1
1	0	0	0	0	0	0	1	1
1	1	0	1	1	1	0	1	1



Absorption (Combining)

$$(A+B)\cdot(A+\overline{B})=A$$

 $(A\cdot B)+(A\cdot \overline{B})=A$

Α	В	\overline{B}	A+B	A+B	$(A+B)\cdot(A+\overline{B})$	A·B	$A \cdot \overline{B}$	$(A \cdot B) + (A \cdot \overline{B})$
0	0	1	0	1	0	0	0	0
0	1	0	1	0	0	0	0	0
1	0	1	1	1	1	0	1	1
1	1	0	1	1	1	1	0	1



Absorption (Combining)

$$(A \cdot B)+(A \cdot \overline{B} \cdot C)=(A \cdot B)+(A \cdot C)$$

 $(A+B)\cdot(A+\overline{B}+C)=(A+B)\cdot(A+C)$

Α	В	С	\overline{B}	A-B	A·B·C	A.	$(A \cdot B) + (A \cdot \overline{B} \cdot C)$	(A·B)+(A·C)
						С		
0	0	0	1	0	0	0	0	0
0	0	1	1	0	0	0	0	0
0	1	0	0	0	0	0	0	0
0	1	1	0	0	0	0	0	0
1	0	0	1	0	0	0	0	0
1	0	1	1	0	1	1	1	1
1	1	0	0	1	0	0	1	1
1	1	1	0	1	0	1	1	1



Absorption (Combining)

$$(A \cdot B) + (A \cdot \overline{B} \cdot C) = (A \cdot B) + (A \cdot C)$$

$$(A+B)\cdot(A+\overline{B}+C)=(A+B)\cdot(A+C)$$

Α	В	С	\overline{B}	A+B	A+B+C	A+C	$(A+B)\cdot(A+\overline{B}+C)$	(A+B)·(A+C)
0	0	0	1	0	1	0	0	0
0	0	1	1	0	1	1	0	0
0	1	0	0	1	0	0	0	0
0	1	1	0	1	1	1	1	1
1	0	0	1	1	1	1	1	1
1	0	1	1	1	1	1	1	1
1	1	0	0	1	1	1	1	1
1	1	1	0	1	1	1	1	1



DeMorgan's Theorem

$$\overline{A+B} = \overline{A} \cdot \overline{B}$$

 $\overline{A \cdot B} = \overline{A} + \overline{B}$

- Break the bar and change the operator
- Connect the bar and change the operator
- Expands to include more terms(Generalized DeMorgan's):

$$\overline{A+B+C+...}=\overline{A}\cdot\overline{B}\cdot\overline{C}\cdot...$$

 $\overline{A\cdot B\cdot C\cdot ...}=\overline{A}+\overline{B}+\overline{C}+...$



DeMorgan's Theorem

$$\overline{A+B} = \overline{A} \cdot \overline{B}$$

$$\overline{A \cdot B} = \overline{A} + \overline{B}$$

Α	В	A+B	A·B	A+B	A·B	A	В	$\overline{A} \cdot \overline{B}$	Ā+B
0	0	0	0	1	1	1	7	1	1
0	1	1	0	0	1	1	0	0	1
1	0	1	0	0	1	0	1	0	1
1	1	1	1	0	0	0	0	0	0



- Consensus Theorem
- $(A \cdot B) + (\overline{A} \cdot C) + (B \cdot C) = (A \cdot B) + (\overline{A} \cdot C)$
- $(A+B)\cdot(\overline{A}+C)\cdot(B+C)=(A+B)\cdot(\overline{A}+C)$

Α	В	C	$ \overline{A} $	A٠	Ā·	B-C	$(A \cdot B) + (\overline{A} \cdot C) +$	$(A \cdot B) + (\overline{A} \cdot C)$
				В	С		(B·C)	
0	0	0	1	0	0	0	0	0
0	0	1	1	0	1	0	1	1
0	1	0	1	0	0	0	0	0
0	1	1	1	0	1	1	1	1
1	0	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0
1	1	0	0	1	0	0	1	1
1	1	1	0	1	1	1	1	1



- Consensus Theorem
- $(A \cdot B) + (\overline{A} \cdot C) + (B \cdot C) = (A \cdot B) + (\overline{A} \cdot C)$
- $(A+B)\cdot(\overline{A}+C)\cdot(B+C)=(A+B)\cdot(\overline{A}+C)$

Α	В	С	Ā	A+B	Ā+C	B+C	$(A+B)\cdot(\overline{A}+C)\cdot(B+C)$	$(A+B)\cdot(\overline{A}+C)$
0	0	0	1	0	1	0	0	0
0	0	1	1	0	1 1 0		0	
0	1	0	1	1	1	1	1	1
0	1	1	1	1	1	1	1	1
1	0	0	0	1	0	0	0	0
1	0	1	0	1	1	1	1	1
1	1	0	0	1	0	1	0	0
1	1	1	0	1	1	1	1	1



Shannon's Expansion

$$f(A,B,C)=(A\cdot f(1,B,C))+(\overline{A}\cdot f(0,B,C))$$

 $f(A,B,C)=(A+f(0,B,C))\cdot(\overline{A}+f(1,B,C))$

Example:

$$f(A,B,C)=A\cdot B\cdot C+A\cdot \overline{B}\cdot C+\overline{A}\cdot \overline{B}\cdot C$$

$$f(A,B,C)=A\cdot f(1,B,C)+\overline{A}\cdot f(0,B,C)$$

$$f(1,B,C)=B\cdot C+\overline{B}\cdot C$$

$$f(0,B,C)=\overline{B}\cdot C$$

$$f(A,B,C)=A\cdot (B\cdot C+\overline{B}\cdot C)+\overline{A}\cdot (\overline{B}\cdot C)$$

Important for circuit minimization



Duality Principle

- Any theorem or postulate in Boolean Algebra remains true if
 - 0 and 1 are swapped and
 - AND and OR are swapped

Postulates:

Postulate	Dual Pairs							
Identity	A+0=A	A·1=A						
Complements	$A+\overline{A}=1$	$A \cdot \overline{A} = 0$						
Commutativity	A+B=B+A	A·B=B·A						
Associativity	(A+B)+C=A+(B+C)	$(A \cdot B) \cdot C = A \cdot (B \cdot C)$						
Distributivity	$A+(B\cdot C)=(A+B)\cdot (A+C)$	$A \cdot (B+C) = (A \cdot B) + (A \cdot C)$						

Duality Principle

Theorems

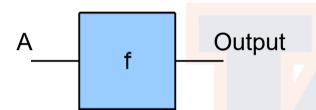
Theorem	Dual	Pairs
Idempotency	A+A=A	A·A=A
Null Elements	A+1=1	A·0=0
Involution	$\overline{A}=A$	$\overline{A}=A$
Covering	A+(A·B)=A	A·(A+B)=A
Covering	$A+(\overline{A}\cdot B)=A+B$	$A \cdot (\overline{A} + B) = A \cdot B$
Combining	$(A \cdot B) + (A \cdot \overline{B}) = A$	$(A+B)\cdot(A+\overline{B})=A$
Combining	$(A \cdot B) + (A \cdot \overline{B} \cdot C) = (A \cdot B) + (A \cdot C)$	$(A+B)\cdot (A+\overline{B}+C)=(A+B)\cdot (A+C)$
DeMorgan's	$\overline{A+B}=\overline{A}\cdot\overline{B}$	$\overline{A \cdot B} = \overline{A} + \overline{B}$
Consensus	$(A \cdot B) + (\overline{A} \cdot C) + (B \cdot C) = (A \cdot B) + (\overline{A} \cdot C)$	$(A+B)\cdot(\overline{A}+C)\cdot(B+C)=(A+B)\cdot$ $(\overline{A}+C)$
Shannon's	$f(A,B,C)=A\cdot f(1,B,C) + \overline{A}\cdot f(0,B,C)$	$f(A,B,C)=A+f(0,B,C)\cdot\overline{A}+f(1,B,C)$

- We've shown proof that two expressions are equivalent by testing every possible input
- We can also generate every expression given the # of inputs, n
- The simplest (trivial) case, n=0, or no inputs:



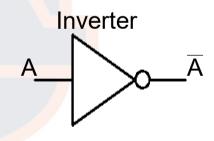


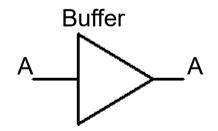
• 1 input, n=1



Α	f0	f1	f2	f3
0	0	1	0	1
1	0	0	1	1

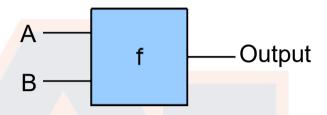
Function	Expression	Description
fO	0	Logic 0
f1	Ā	Inverter
f2	А	Buffer
f3	1	Logic 1







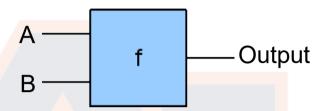
• 2 inputs, n=2



Α	В	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f
		0	1	2	3	4	5	6	7	8	9	Α	В	С	D	E	F
0	0	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
0	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	0	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	1	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1

 Looking closely, we've seen some of these before

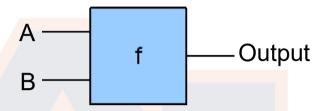




Α	В	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f
		0	1	2	3	4	5	6	7	8	9	Α	В	C	D	E	F
0	0	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
0	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	0	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	1	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1

- These are the two trivial logic sources
 - $f_0 = 0$
 - f₌= 1

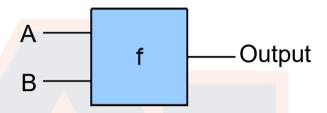




Α	В	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f
		0	1	2	3	4	5	6	7	8	9	Α	В	С	D	E	F
0	0	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
0	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	0	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	1	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1

- f_c and f_A are buffers for A and B respectively
- f₃ and f₅ are inverters for A and B respectively

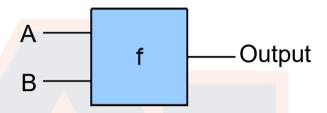




Α	В	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f
		0	1	2	3	4	5	6	7	8	9	Α	В	C	D	Е	F
0	0	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
0	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	0	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	1	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1

- f_E is A+B
 f₈ is A·B





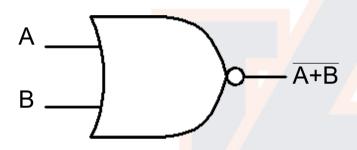
Α	В	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f
		0	1	2	3	4	5	6	7	8	9	Α	В	C	D	Е	F
0	0	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
0	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	0	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	1	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1

- f_1 is $\overline{A+B}$, NOR
- f_7 is $\overline{A \cdot B}$, NAND



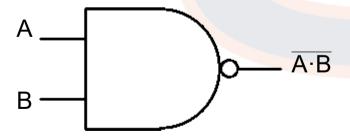
New Operators

• NOR



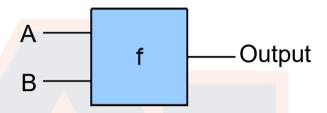
Α	В	A+B				
0	0	1				
0	1	0				
1	0	0				
1	1	0				

NAND



A	В	A·B			
0	0	1			
0	1	1			
1	0	1			
1	1	0			





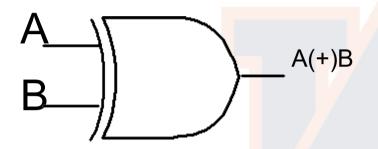
Α	В	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f	f
		0	1	2	3	4	5	6	7	8	9	Α	В	C	D	E	F
0	0	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1
0	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	0	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	1	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1

- f₆, Exclusive OR, XOR
- f₉, Exclusive NOR, XNOR



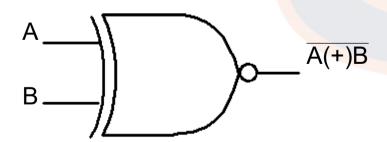
New Operators

• XOR



Α	В	A(+)B
0	0	0
0	1	1
1	0	1
1	1	0

XNOR



Α	В	A(+)B
0	0	1
0	1	0
1	0	0
1	1	1



Switching Function, n=2

Function	Expression	Description				
f0	0	Logic 0				
f1	A+B	NOR				
f2	Ā·B					
f3	Ā	Inverter A				
f4	$A \cdot \overline{B}$					
f5	B	Inverter B				
f6	A(+)B	XOR				
f7	A·B	NAND				
f8	A·B	AND				
f9	A(+)B	XNOR				
fA	В	Buffer B				
fB	Ā+B					
fC	Α	Buffer A				
fD	A+B					
fE	A+B	OR				
fF	1	Logic 1				

- How many possible functions given # of inputs,
 n?
 - n=0, 2 functions
 - n=1, 4 functions
 - n=2, 16 functions
- In general there are 2^{2ⁿ} functions
- So for n=3 there are 256 possible functions

