

CS 33: Computer Organization

Dis 1B: Week 10 Discussion

Agenda

- Final Review?

or

- Forget about the finals and watch NCAA tournament?



Final

- March 20th, Monday: 3PM ~ 6PM
- Kinsey Pavillion 1220B



Things we covered after the midterm

- Buffer overflow
- Optimization
- Cache
- Processes
- Threading
- Multithreading
- Linking
- Virtual Memory



Final

- Most of the questions will be from the materials after the midterm
- The format of the final will be very similar to the midterm



Practice Problem 1

```
// Return address of get_line is 0x400776
// Register %rbx will contain 0x0123456789ABCDEF
char *get_line()
{
    char buf[4];
    char *result;;
    gets(buf);
    result = malloc(strlen(buf));
    strcpy(result, buf);
    return result;
}

// Our input string is 0123456789012345678901234
0000000000400720 <get_line>:
400720: 53                push    %rbx
400721: 48 83 ec 10       sub     $0x10, %rsp
    // Draw a diagram of the stack at this point
400725: 48 89 e7          mov     %rsp, %rdi
400728: e8 73 ff ff ff   callq  4006a0 <gets>
    // Modify diagram to show stack contents at this point
```

Practice Problem 2

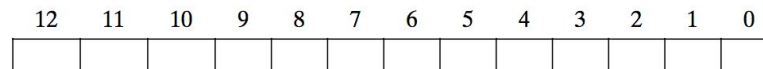
- Assume the following
 - The memory is byte addressable
 - Memory accesses are to 1-byte words
 - Addresses are 13 bits wide
 - The cache is two-way set associative, with 4-byte block size and 8 sets

2-way set associative cache

| Set index | Line 0 | | | | | | Line 1 | | | | | |
|-----------|--------|-------|--------|--------|--------|--------|--------|-------|--------|--------|--------|--------|
| | Tag | Valid | Byte 0 | Byte 1 | Byte 2 | Byte 3 | Tag | Valid | Byte 0 | Byte 1 | Byte 2 | Byte 3 |
| 0 | 09 | 1 | 86 | 30 | 3F | 10 | 00 | 0 | — | — | — | — |
| 1 | 45 | 1 | 60 | 4F | E0 | 23 | 38 | 1 | 00 | BC | 0B | 37 |
| 2 | EB | 0 | — | — | — | — | 0B | 0 | — | — | — | — |
| 3 | 06 | 0 | — | — | — | — | 32 | 1 | 12 | 08 | 7B | AD |
| 4 | C7 | 1 | 06 | 78 | 07 | C5 | 05 | 1 | 40 | 67 | C2 | 3B |
| 5 | 71 | 1 | 0B | DE | 18 | 4B | 6E | 0 | — | — | — | — |
| 6 | 91 | 1 | A0 | B7 | 26 | 2D | F0 | 0 | — | — | — | — |
| 7 | 46 | 0 | — | — | — | — | DE | 1 | 12 | C0 | 88 | 37 |

The following figure shows the format of an address (one bit per box). Indicate (by labeling the diagram) the fields that would be used to determine the following:

- CO* The cache block offset
CI The cache set index
CT The cache tag



Practice Problem 2

2-way set associative cache

| Set index | Line 0 | | | | | | Line 1 | | | | | |
|-----------|--------|-------|--------|--------|--------|--------|--------|-------|--------|--------|--------|--------|
| | Tag | Valid | Byte 0 | Byte 1 | Byte 2 | Byte 3 | Tag | Valid | Byte 0 | Byte 1 | Byte 2 | Byte 3 |
| 0 | 09 | 1 | 86 | 30 | 3F | 10 | 00 | 0 | — | — | — | — |
| 1 | 45 | 1 | 60 | 4F | E0 | 23 | 38 | 1 | 00 | BC | 0B | 37 |
| 2 | EB | 0 | — | — | — | — | 0B | 0 | — | — | — | — |
| 3 | 06 | 0 | — | — | — | — | 32 | 1 | 12 | 08 | 7B | AD |
| 4 | C7 | 1 | 06 | 78 | 07 | C5 | 05 | 1 | 40 | 67 | C2 | 3B |
| 5 | 71 | 1 | 0B | DE | 18 | 4B | 6E | 0 | — | — | — | — |
| 6 | 91 | 1 | A0 | B7 | 26 | 2D | F0 | 0 | — | — | — | — |
| 7 | 46 | 0 | — | — | — | — | DE | 1 | 12 | C0 | 88 | 37 |

- 0x0E34
- 0x0DD5
- 0x1FE4

A. Address format (one bit per box):

| | | | | | | | | | | | | |
|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|--------------------------|
| 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> | <input type="checkbox"/> |

B. Memory reference:

| Parameter | Value |
|-------------------------|-------------------------|
| Cache block offset (CO) | 0x <input type="text"/> |
| Cache set index (CI) | 0x <input type="text"/> |
| Cache tag (CT) | 0x <input type="text"/> |
| Cache hit? (Y/N) | <input type="text"/> |
| Cache byte returned | 0x <input type="text"/> |

Adapted from Bryant and O'Hallaron,
Computer Systems: A Programmer's
Perspective. Third Edition

Practice Problem 3-1

```
int main()
{
    for (int i = 0; i < 3; i++)
    {
        fork();
    }
}
```

// How many processes are created by the program?

Practice Problem 3-2

```
int main()
{
    int foo = fork();
    int bar = fork();

    if (foo != 0 && bar != 0)
    {
        int tricky = fork();
        if (tricky == 0)
            printf("Wow this is tricky!\n");
        else
            printf("What\n");
    }
    else if (foo != 0 && bar == 0)
        printf("Huh?\n");
    else
        printf("Hmm...\n");
}

// How many processes are created by the program?
```

Practice Problem 4

- P locks the variable passed in as an argument
- V unlocks the variable passed in as an argument
- Will this program deadlock?

Thread 1:

```
P(a);  
P(b);  
V(b);  
P(c);  
V(c);  
V(a);
```

Thread 2:

```
P(c);  
P(b);  
V(b);  
V(c);  
P(a);  
V(a);
```

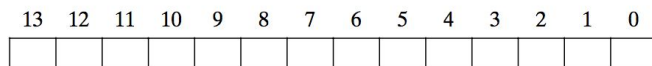
Thread 3:

```
P(c);  
V(c);  
P(b);  
P(a);  
V(a);  
V(b);
```

Practice Problem 5

- 0x03d4
- 0x03d7
- 0x03a9
- 0x0040

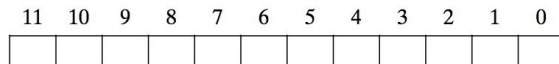
A. Virtual address format



B. Address translation

| Parameter | Value |
|-------------------|-------------------|
| VPN | <u> </u> |
| TLB index | <u> </u> |
| TLB tag | <u> </u> |
| TLB hit? (Y/N) | <u> </u> |
| Page fault? (Y/N) | <u> </u> |
| PPN | <u> </u> |

C. Physical address format



D. Physical memory reference

| Parameter | Value |
|---------------------|-------------------|
| Byte offset | <u> </u> |
| Cache index | <u> </u> |
| Cache tag | <u> </u> |
| Cache hit? (Y/N) | <u> </u> |
| Cache byte returned | <u> </u> |

TLBT

← TLBI →

13

12

11

10

9

8

7

6

5

4

3

2

1

0

Virtual address

VPN

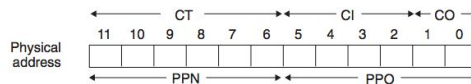
VPO

| Set | Tag | PPN | Valid | Tag | PPN | Valid | Tag | PPN | Valid | Tag | PPN | Valid |
|-----|-----|-----|-------|-----|-----|-------|-----|-----|-------|-----|-----|-------|
| 0 | 03 | — | 0 | 09 | 0D | 1 | 00 | — | 0 | 07 | 02 | 1 |
| 1 | 03 | 2D | 1 | 02 | — | 0 | 04 | — | 0 | 0A | — | 0 |
| 2 | 02 | — | 0 | 08 | — | 0 | 06 | — | 0 | 03 | — | 0 |
| 3 | 07 | — | 0 | 03 | 0D | 1 | 0A | 34 | 1 | 02 | — | 0 |

(a) TLB: Four sets, 16 entries, four-way set associative

| | | | | | |
|-----|-----|-------|-----|-----|-------|
| VPN | PPN | Valid | VPN | PPN | Valid |
| 00 | 28 | 1 | 08 | 13 | 1 |
| 01 | — | 0 | 09 | 17 | 1 |
| 02 | 33 | 1 | 0A | 09 | 1 |
| 03 | 02 | 1 | 0B | — | 0 |
| 04 | — | 0 | 0C | — | 0 |
| 05 | 16 | 1 | 0D | 2D | 1 |
| 06 | — | 0 | 0E | 11 | 1 |
| 07 | — | 0 | 0F | 0D | 1 |

(b) Page table: Only the first 16 PTEs are shown



| | | | | | | |
|-----|-----|-------|-------|-------|-------|-------|
| Idx | Tag | Valid | Blk 0 | Blk 1 | Blk 2 | Blk 3 |
| 0 | 19 | 1 | 99 | 11 | 23 | 11 |
| 1 | 15 | 0 | — | — | — | — |
| 2 | 1B | 1 | 00 | 02 | 04 | 08 |
| 3 | 36 | 0 | — | — | — | — |
| 4 | 32 | 1 | 43 | 6D | 8F | 09 |
| 5 | 0D | 1 | 36 | 72 | F0 | 1D |
| 6 | 31 | 0 | — | — | — | — |
| 7 | 16 | 1 | 11 | C2 | DF | 03 |
| 8 | 24 | 1 | 3A | 00 | 51 | 89 |
| 9 | 2D | 0 | — | — | — | — |
| A | 2D | 1 | 93 | 15 | DA | 3B |
| B | 0B | 0 | — | — | — | — |
| C | 12 | 0 | — | — | — | — |
| D | 16 | 1 | 04 | 96 | 34 | 15 |
| E | 13 | 1 | 83 | 77 | 1B | D3 |
| F | 14 | 0 | — | — | — | — |

(c) Cache: Sixteen sets, 4-byte blocks, direct mapped

Adapted from Bryant
and O'Hallaron,
Computer Systems: A
Programmer's
Perspective, Third
Edition



Good luck on your finals!

Thanks!