



Operating Systems  
Course Code: **71203002004**  
*Threads*

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# What is a Thread?

- A **thread** is a lightweight part of a process that runs independently.
- Multiple threads can exist in a single process, sharing the same code, data, and OS resources, but each has its own stack, registers, and program counter.
- In systems that support **multithreading**, threads help improve performance, especially on multi-core CPUs.
- On a single CPU, threads take turns using the processor (context switching).

# Why Use Threads?

Threads allow **concurrent execution**, which boosts application performance.

For example, in apps like Microsoft Word or Google Docs, typing, auto-saving, and formatting can happen at the same time.

## Key features:

- Threads share memory, so no need for inter-process communication.
- Each thread has its own **Thread Control Block (TCB)**.
- Threads can have different priorities.
- Synchronization is needed to manage shared data.

# Components of Thread

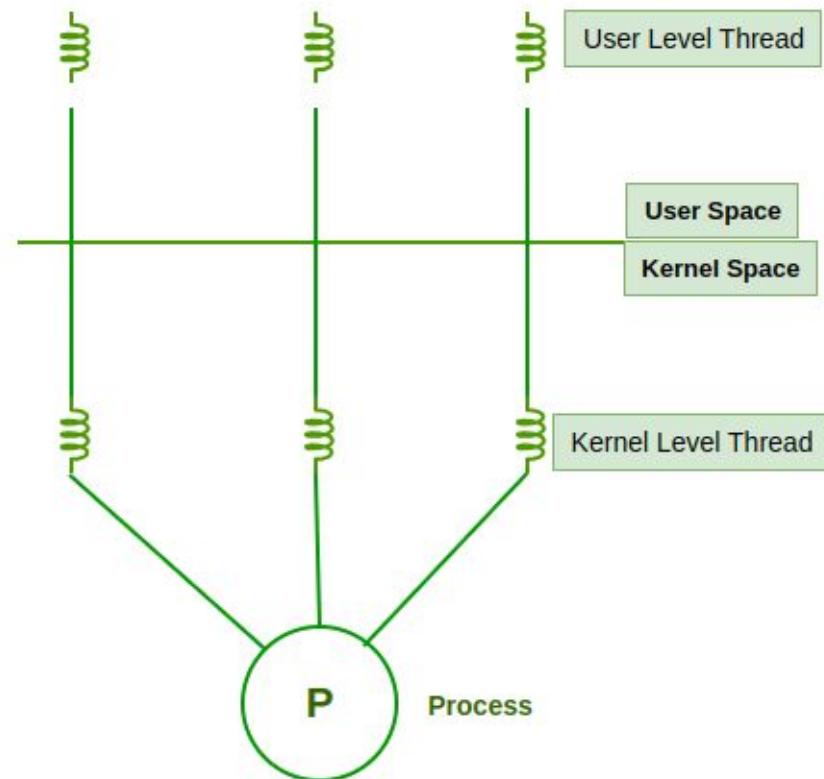
- These are the basic components of the Operating System.
  - **Stack Space:** Stores local variables, function calls, and return addresses specific to the thread.
  - **Register Set:** Hold temporary data and intermediate results for the thread's execution.
  - **Program Counter:** Tracks the current instruction being executed by the thread.



# Types of thread in OS

Threads are of two types:

1. User level
2. Kernel level



## User-Level Threads (ULT)

- Managed by the application, not the OS.
- **Advantages:**
  - Easy to implement
  - Fast context switching
  - More efficient in simple tasks
- **Disadvantages:**
  - OS doesn't manage them, so no CPU optimization
  - One blocking thread can pause the whole process
  - Complex scheduling for multiple threads

## **Kernel-Level Threads (KLT)**

- Managed by the OS kernel.

### **Advantages:**

- Can run on multiple cores
- OS manages scheduling and load balancing
- Better for frequent blocking tasks

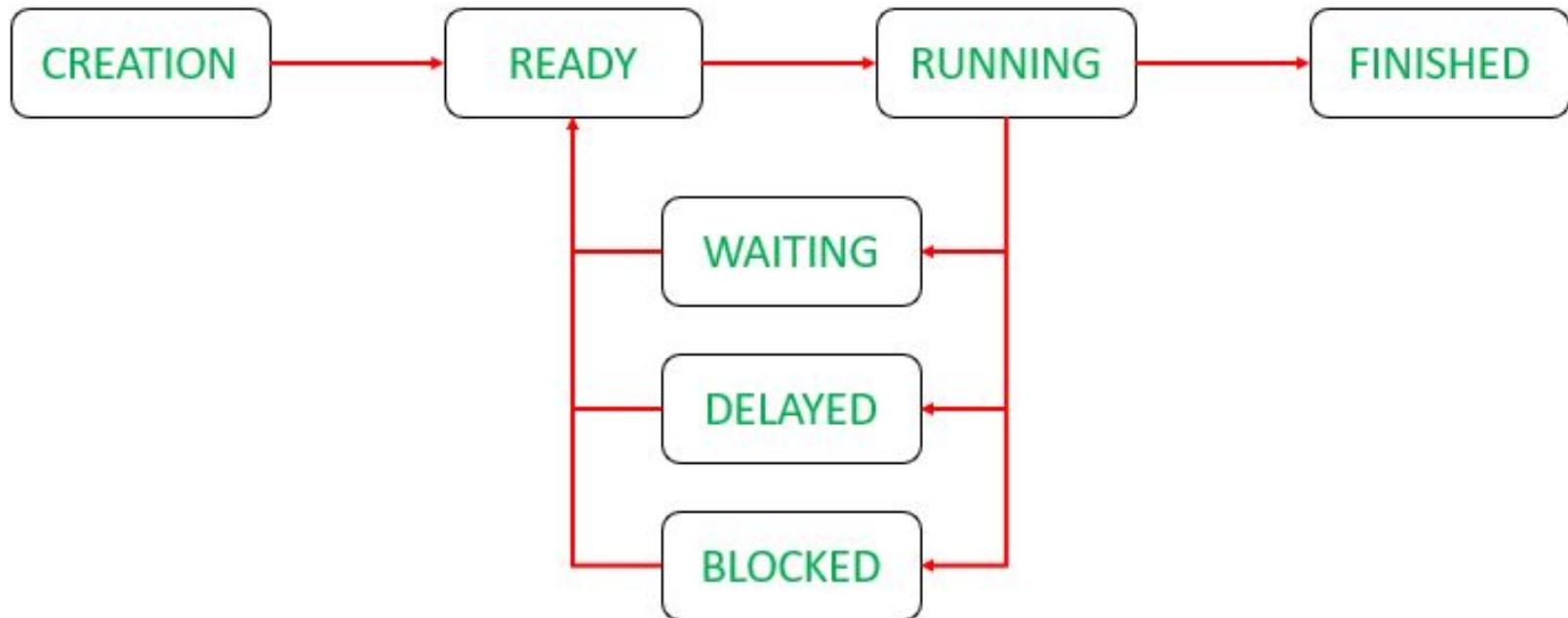
### **Disadvantages:**

- Slower context switching
- Higher CPU overhead due to system calls
- More complex to implement

# Thread States in OS

A thread in an operating system can be in one of the five-main states (excluding creation and finished states):

1. **Ready** : thread is prepared to run and waiting for CPU assignment.
2. **Running**: thread id currently being executed by the CPU.
3. **Waiting**: thread is paused , waiting for an event or signal from another thread/ process.
4. **Delayed**: thread is intentionally put to sleep for a fixed time (e.g., snoozed alarm).
5. **Blocked**: thread is waiting for an I/O operation to complete (e.g., user input).



# Thread State Transition

## 1. Ready > Running

When the CPU scheduler assigns a thread to the processor.

## 1. Running > Waiting

When the thread needs to wait for another event or process to complete.

## 1. Running > Delayed

When the thread is programmatically paused for a set time.

## 1. Running > Blocked

When the thread performs an I/O and must wait for it to finish

## 1. Running > Finished

When the thread completes its task.

## Key Differences

- **Waiting:** Waits for an external signal or event with **known burst time** (e.g., waiting for another thread).
- **Blocked:** Waits for **unknown time**, often due to user input or I/O operations.

## Thread Control Block (TCB)

To track thread states and manage scheduling, the **OS uses a Thread Control Block (TCB)**, which stores:

- Thread ID
- State
- Program counter
- Register values
- Stack pointer
- Scheduling information



# Difference between Thread and Process

Feature	Process	Thread
<b>Memory</b>	Has its own separate memory space	Shares memory with other threads
<b>Independence</b>	Independent from other processes	Dependent on other threads in the process
<b>Communication</b>	Requires inter-process communication	Can communicate directly via shared memory
<b>Resources</b>	Has its own OS resources	Shares OS resources (e.g., files, signals)
<b>Components</b>	Has its own PC, registers, and stack	Has its own PC, registers and stack
<b>Overhead</b>	Higher(more resource intensive)	Lower(lightweight)



Feature	Process	Thread
<b>Context Switching</b>	Slower (more data to save/load)	Faster (less data to manage)
<b>Creation Time</b>	Slow to create	Faster to create
<b>Crash Impact</b>	One process crashing doesn't affect others	A crashing thread can affect the whole process.

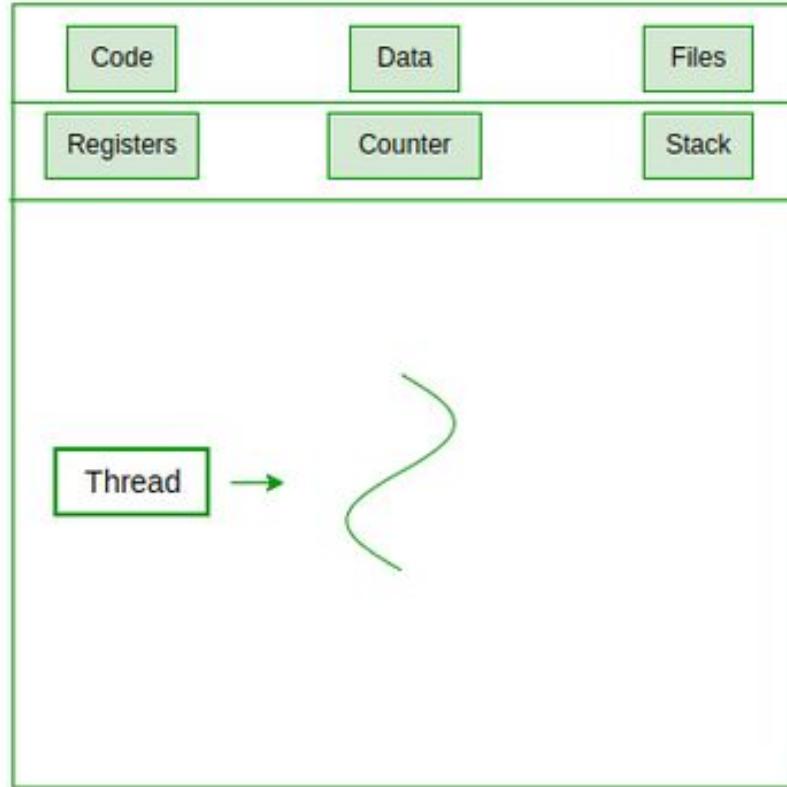
# What is a Multi-Threading?

**Multithreading** is the ability of a process to run multiple threads concurrently to achieve **parallelism** and improve performance.

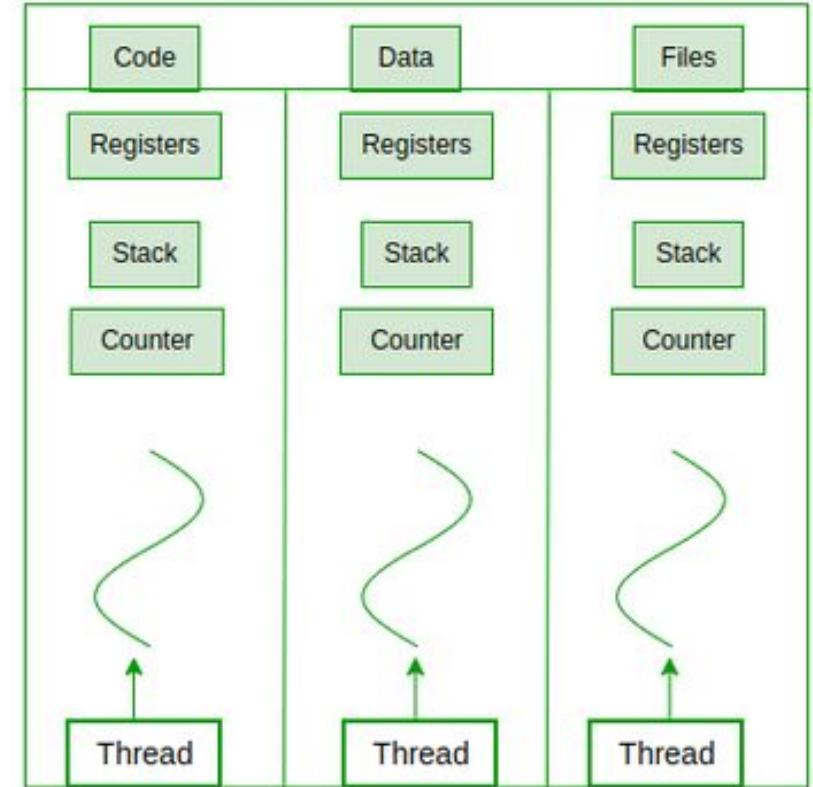
- A **thread** is a lightweight unit of a process.
- All threads in a process share the same memory and resources (code, data, files).
- Each thread has its own **stack**, **program counter**, and **registers**.

## Examples:

- In a browser, each tab can be a thread.
- In MS Word, one thread handles typing, another handles formatting, another handles autosaving, etc.



Single Threaded Process



Multi Threaded Process



## How It Works

- Threads can be managed by the **operating system (kernel-level)** or by the **application itself (user-level)**.
- For example, **Java** uses the JVM to manage threads without depending on the OS.
- User-level threads are handled by a **thread library** within the application; the OS is unaware of them.

## Benefits of Multithreading

- **Responsiveness:** threads can return result as they finish, improving user experience.
- **Faster Context Switching:** Switching between threads is quicker than between full process.
- **Better Use of Multiprocessors :** Threads can run in parallel on multiple CPU cores, speeding up execution.
- **Resource Sharing:** Threads share code, data and files, reducing duplication.



- **Easier Communication:** threads use shared memory, avoiding complex inter process communication.
- **Higher Throughout:** More work (threads) done in less time increases system efficiency.



## DISCUSSION & REVISION

1. What does OS assign to each new process?
2. Which component decides which process runs next?
3. What happens to a process when it finishes execution?
4. What do processes use to communicate with each other?
5. What prevents processes from waiting forever in a loop?