

# LOGIC

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## PROPOSITIONAL LOGIC (Deals with propositions)

Proposition: A declarative statement that is either true or false but not both.

Example

- ① Ahmedabad is the capital of Gujarat. F } Proposition
- ②  $5+2=4$  F }
- ③  $1+1=2$  T }

④ Please pass the paper. (not a proposition as not declaring anything)

⑤  $x+5=7$  (not a proposition but a predicate logic)

• Requests, orders are not propositions and expressions that cannot be evaluated as T or F are not propositions.

- Propositional variable or statement variables are denoted by  $p, q, r, s, \dots$

- Truth values are denoted by T or F.

From simple propositions, compound propositions can be formed using logical connectives or operators.

① Let  $p$  be a proposition. The negation of  $p$ , denoted by " $\neg p$ ", is the statement that, "It is not the case that  $p$ ". The truth value of  $\neg p$  is the opposite of the truth value of  $p$ .

Example  $p$ : Today is Friday.  
 $\neg p$ : "It is not the case that today is Friday" or "Today is not Friday".

Truth Table

$\neg p$	$p$	$\neg p$
	T	F
	F	T

②  $p \wedge q$  (conjunction/and) is the statement " $p$  and  $q$ ".

Example  $p$ : Today is Friday  
 $q$ : It is raining today

$p \wedge q$ : Today is Friday and it is raining today.

The statement is T on rainy Fridays and is F on any day that is not a Friday & on Fridays when it does not rain.

Truth Table (T, T)

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F



③ **Disjunction**  $p \vee q$  : "p or q" (Inclusive OR) ②

**Example** Same as the previous

T.T

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

Today is Friday or it is raining today.  
True on any day that is either Friday or a rainy day (including the rainy Fridays), false on days that are not Fridays when it also does not rain.

④ **Exclusive OR**  $p \oplus q$  : "either p or q but not both"

True when exactly one of them is true and is false otherwise.

T.T

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

⑤ **Conditional Statements**  $p \rightarrow q$  : "If p then q"

$p \rightarrow q$  is ~~not~~ false when p is T & q is F & true otherwise.

hypothesis → conclusion/consequence

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

- "If p then q"
- "If p, q"
- "p is sufficient for q"
- "q if p"
- "q when p"
- "a necessary condition for p is q"
- "q unless  $\neg p$ "
- "p implies q"
- "p only if q"
- "a ~~not~~ sufficient condition for q is p"
- "q whenever p"
- "q is necessary for p"
- "q follows from p"



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## Converse, Inverse & Contrapositive

- $q \rightarrow p$  is the converse of  $p \rightarrow q$
- $\neg p \rightarrow \neg q$  is the inverse of  $p \rightarrow q$
- $\neg q \rightarrow \neg p$  is the contrapositive of  $p \rightarrow q$ .

## Biconditionals

$p \leftrightarrow q$  : "p if and only if q"

True only when both are true or both are false & false otherwise.

T.T.

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

## T.T. of compound propositions

### Example

$$(p \vee \neg q) \rightarrow (p \wedge q)$$

p	q	$\neg q$	$p \vee \neg q$	$p \wedge q$	$(p \vee \neg q) \rightarrow (p \wedge q)$
T	T	F	T	T	T
T	F	T	T	F	F
F	T	F	F	F	T
F	F	T	T	F	F

## Precedence of logical operators

operator	Precedence
$\neg$	1
$\wedge$	2
$\vee$	3
$\rightarrow$	4
$\leftrightarrow$	5

## Propositional Equivalence

A compound proposition that is always true no matter what the truth values of the propositions that occur in it is called Tautology.

A compound proposition that is always false is called contradiction.

Contingency : Compound prop. that is neither a tautology nor a contradiction.



**Example**  $p \vee \neg p$  is a tautology &  $p \wedge \neg p$  is a contradiction.  
T. T.

$p$	$\neg p$	$p \vee \neg p$	$p \wedge \neg p$
T	F	T	F
F	T	T	F
		always true	always false

## Logical Equivalences

Compound propositions  $p$  &  $q$  are logically equivalent if  $p \leftrightarrow q$  is a tautology and is denoted by  $p \equiv q$  or " $p \Leftrightarrow q$ ".

**Example**  $\neg(p \wedge q) \equiv \neg p \vee \neg q$   
 $\neg(p \vee q) \equiv \neg p \wedge \neg q$  } De Morgan's Law

T. T.

$p$	$q$	$p \vee q$	$\neg(p \vee q)$	$\neg p$	$\neg q$	$\neg p \wedge \neg q$
T	T	T	F	F	F	F
T	F	T	F	F	T	F
F	T	T	F	T	F	F
F	F	F	T	T	T	T

**Example**  $p \rightarrow q \equiv \neg p \vee q$

T. T.

$p$	$q$	$\neg p$	$\neg p \vee q$	$p \rightarrow q$
T	T	F	T	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T



Example  $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$  Distributive Law ⑤

p	q	r	$q \wedge r$	$p \vee (q \wedge r)$	$p \vee q$	$p \vee r$	$(p \vee q) \wedge (p \vee r)$
T	T	T	T	T	T	T	T
T	T	F	F	T	T	T	T
T	F	T	F	T	T	T	T
T	F	F	F	T	T	T	T
F	T	T	T	T	T	T	T
F	T	F	F	F	T	F	F
F	F	T	F	F	F	T	F
F	F	F	F	F	F	F	F

$$\neg (p_1 \vee p_2 \vee \dots \vee p_n) \equiv (\neg p_1 \wedge \neg p_2 \wedge \dots \wedge \neg p_n)$$

$$\neg (p_1 \wedge p_2 \wedge \dots \wedge p_n) \equiv (\neg p_1 \vee \neg p_2 \vee \dots \vee \neg p_n)$$

### Some Logical Equivalences

$$\left. \begin{array}{l} p \wedge T \equiv p \\ p \vee F \equiv p \end{array} \right\} \text{Identity Laws}$$

$$\left. \begin{array}{l} p \vee T \equiv T \\ p \wedge F \equiv F \end{array} \right\} \text{Domination Laws}$$

$$\left. \begin{array}{l} p \vee p \equiv p \\ p \wedge p \equiv p \end{array} \right\} \text{Idempotent Laws}$$

$$\neg (\neg p) \equiv p \quad \left. \begin{array}{l} \text{Double} \\ \text{Negation Law} \end{array} \right\}$$

$$\left. \begin{array}{l} p \vee q \equiv q \vee p \\ p \wedge q \equiv q \wedge p \end{array} \right\} \text{Commutative Laws}$$

$$\left. \begin{array}{l} (p \vee q) \vee r \equiv p \vee (q \vee r) \\ (p \wedge q) \wedge r \equiv p \wedge (q \wedge r) \end{array} \right\} \text{Associative Laws}$$

$$\begin{array}{l} p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r) \\ p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r) \end{array}$$

Distributive Laws

$$\begin{array}{l} \neg (p \wedge q) \equiv \neg p \vee \neg q \\ \neg (p \vee q) \equiv \neg p \wedge \neg q \end{array}$$

De Morgan's Laws

$$\begin{array}{l} p \vee (p \wedge q) \equiv p \\ p \wedge (p \vee q) \equiv p \end{array}$$

Absorption Laws

$$\begin{array}{l} p \vee \neg p \equiv T \\ p \wedge \neg p \equiv F \end{array}$$

Negation Laws



## Conditional Logical Equi.

$$p \rightarrow q \equiv \neg p \vee q$$

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

$$p \vee q \equiv \neg p \rightarrow q$$

$$p \wedge q \equiv \neg (q \rightarrow \neg p)$$

$$\neg (p \rightarrow q) \equiv p \wedge \neg q$$

$$(p \rightarrow q) \wedge (p \rightarrow r) \equiv p \rightarrow (q \wedge r)$$

$$(p \rightarrow r) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$$

$$(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$$

$$(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$$

## Biconditional Equi.

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$

$$p \leftrightarrow q \equiv \neg p \leftrightarrow \neg q$$

$$p \leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)$$

$$\neg (p \leftrightarrow q) \equiv p \leftrightarrow \neg q$$

[Example] Show that  $\neg (p \rightarrow q) \equiv p \wedge \neg q$

→ We know that,  $p \rightarrow q \equiv \neg p \vee q$

Then,  $\neg (p \rightarrow q) \equiv \neg (\neg p \vee q)$   
 $\equiv p \wedge \neg q$  (using De Morgan's Laws)

## Bitwise operations

Bitwise OR, Bitwise AND, Bitwise XOR

Consider two bit string  $x=10110$  and  $y=11011$ .

x	y	$x \vee y$	$x \wedge y$	$x \oplus y$
1	1	1	1	0
0	1	1	0	1
1	0	1	0	1
1	1	1	1	0
0	1	1	0	1

Consider 0 as F  
and 1 as T.  
and computation  
follows bitwise.