

Topology Inference for Radial Distribution Feeder based on Power Flow

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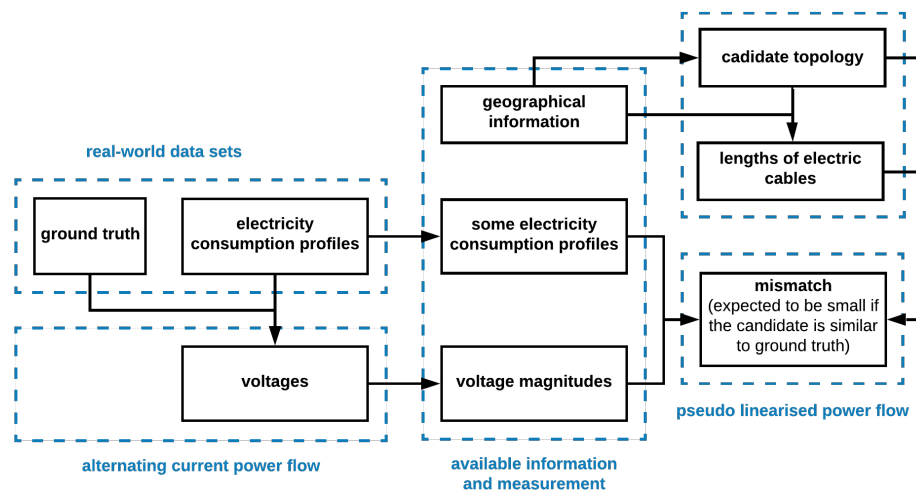
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Chapter 1

Introduction

This website hosts slides for defence of my master graduation project in the Department of Electrical Engineering at Technical University of Denmark. How households are connected to distribution network is always unknown. A framework to infer such connections by utilising all kinds of information is proposed in this project.

A simple flowchart for the project is:



Available information and measurement:

- geographical information about buses
- voltage magnitudes of all the phases of all the buses
- some real power injection profiles

Two batches of computer programs:

- alternating current power flow calculation its linearised version
- three algorithms to handle directed graphs

Chapter 2

Radial Distribution Feeder

2.1 Bus and Edge

There are roughly two types of electrical devices in power grids.

type	definition	examples
delivery element	transport power from one place to another	cable, transformer, capacitor
conversion element	convert power from or to another form	solar panel, battery

- Ignore conversion elements. Not necessary in power flow calculation.
- Delivery element will be called **edge**.

Another concept, **bus**, represent the place where two different delivery elements joint or end of a delivery element, but there is no physical entity corresponding to a bus. There are three common types of buses:

type	know quantities
slack bus	voltage magnitude and phase angle
PQ bus	real power injection and reactive power injection
PV bus	real power injection and voltage magnitude

It is sufficient to model most of RDFs with PQ buses and one kind of edges, cables:

- One slack bus in RDF, corresponding to the **root**.
- Root not in any matrix.
- Ignore other delivery elements.

2.2 Two Special Concepts

Essential for power flow calculation.

2.2.1 Channel

observed active channels

2.2.2 Snapshot

Snapshot is a concept to include power injections and voltages at one time index

- input: real power injections at all channels of PQ buses
- output: voltages, current flow, power flow

Zero-load snapshot is the snapshot where power injections at all the channels are zero and voltages equal to rated voltages in corresponding phases.

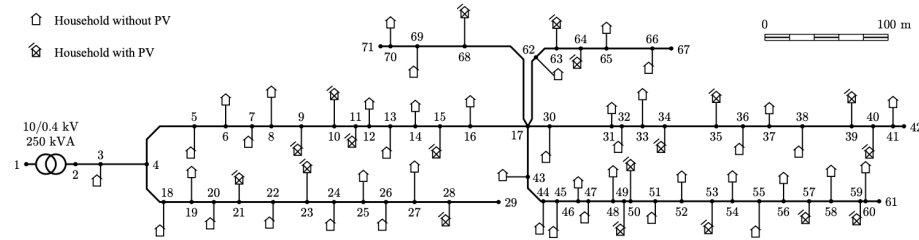
- \bar{V}_{zero} : voltages in zero-load snapshot

2.3 Case with 70 Buses

Assumptions about feeders:

- arborescence
- step-down transformer is not considered
- three-phase four-wire cable
- one phase star connection

A case with 70 buses is primarily used here:



- located in Belgium
- bus 1 is omitted
- Houses associated with buses 3, 7, 10, 13, 16, 20, 23, 26, 30, 33, 36, 39, 43, 46, 49, 52, 55, 58, 62, 65, 69 are connected through phase A.
- Houses associated with buses 5, 8, 11, 14, 18, 21, 24, 27, 31, 34, 37, 40, 44, 47, 50, 53, 56, 59, 63, 66, 70 are connected through phase B.

Chapter 3

Problem Formulation

3.1 Directed Graph

weighted directed graph $G = (\mathcal{N}, \mathcal{E}, \sigma, \tau, \omega)$

- set of nodes: \mathcal{N}
- set of edges: \mathcal{E}
- incidence functions: source σ , target τ
- (edge) weighting function, $\omega : E \rightarrow \mathbb{R}$.
- 2-D Euclidean Distance as Weight

complete graph for a set of nodes

- all edges are potential edges
- some are impossible to exist
- association network inference cannot be used

spanning arborescence

- subgraph of a directed graph
- root
- include every node (bus, in our case)

feasible region

- Composed of all the spanning arborescences of the directed graph.
- Number of spanning arborescences is finite, making it a combinatorial optimisation problem.
- To count number of spanning arborescences.

3.2 IP Formulation

Symbols and definitions of sets:

symbol	definition
\mathcal{E}	all the potential edges (edges in the complete graph)
\mathcal{C}	available measurements of voltages and power injections
$\mathcal{E}_{\text{impossible}}$	potential edges that are impossible to exist

Symbols, definitions, types and sets of variables:

symbol	definition	type	set
x_{ij}	if edge from i to j is in the solution	$\{0, 1\}$	\mathcal{E}

Symbols, definitions, sets of constants:

symbol	definition	set
$d_{i,j}$	weight of directed edge from i to j based on distance	\mathcal{E}

The **integer programming formulation** is:

$$\begin{aligned}
& \min_{x_{ij} \forall (i,j) \in \mathcal{E}} (1 - \alpha) \sum_{(i,j) \in \mathcal{E}} d_{ij} x_{ij} + \alpha \mathcal{H}(\{x_{ij} \forall (i,j) \in \mathcal{E}\}, \mathcal{C}) \\
& \text{s.t.} \quad \sum_{(i,j) \in \delta^-(j)} x_{ij} = 1 \quad \forall j \in V' \quad (\text{a directed forest}) \\
& \quad \sum_{(i,j) \in \delta^-(S)} x_{ij} \geq 1 \quad \forall S \subseteq V', |S| \geq 2 \quad (\text{a connected graph}) \\
& \quad x_{ij} = 0 \quad \forall (i,j) \in \mathcal{E}_{\text{impossible}} \quad (\text{remove impossible potential edges})
\end{aligned}$$

Two terms in the objective function:

term	definition	coefficient
$(1 - \alpha) \sum_{(i,j) \in \mathcal{E}} d_{ij} x_{ij}$	weight of candidate arborescence	$1 - \alpha$
$\alpha \mathcal{H}(\{x_{ij} \forall (i,j) \in \mathcal{E}\}, \mathcal{C})$	assessment of candidate arborescence	α

Three sets of constraints:

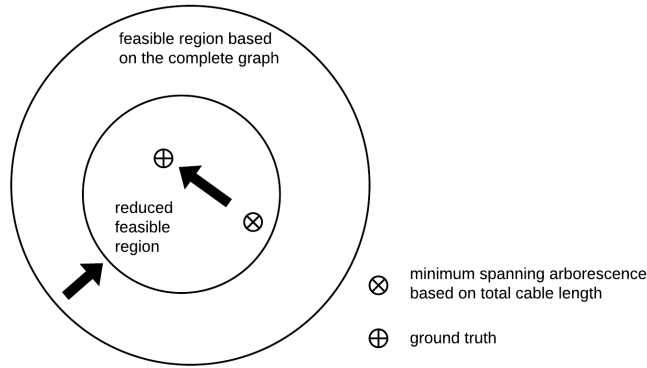
- First two sets ensure arborescence.
- Last set removes impossible potential edges.

3.3 Combinatorial Optimisation

At least two possible values for α :

value	term lefted	to find	disadvantage
1	$\mathcal{H}(\{x_{ij} \mid \forall (i, j) \in \mathcal{E}\}, \mathcal{C})$	ground truth	NP-hard and non-linear
0	$\sum_{(i,j) \in \mathcal{E}} d_{ij} x_{ij}$	topology with min total cable length	cannot find ground truth

Such two situations can be visualised:



In this project, **local search heuristic algorithm** is proposed to to move from \otimes to \oplus :

function	what it does	in this project
objective	assess candidate	pseudo linearised power flow
neighbourhood	generate candidate	rank spanning arborescence

- The starting point is found by minimum spanning arborescence
- Every candidate is reachable from the starting point.
- Ground truth should be found before long.
- Not in parallel.

Chapter 4

AC Power Flow

Discussion is based on one-line model. Can be generalised for multi-phase model.

- Model RDF with one bus impedance matrix.
- Model power flow using direct impedance method.
- Calculate power flow using fixed point method.

4.1 Bus Impedance Matrix

Bus impedance matrix (BIM), \bar{Z} , is defined as:

$$\begin{aligned}\bar{Z} &= K^\top \bar{Z}_{\text{edge}} K \\ &= R + jX\end{aligned}$$

where:

- **edge path incidence matrix (EPI)**, K : current injection to current flow
- **edge impedance diagonal block matrix (EIDB)**, \bar{Z}_{edge} : voltage drop to voltage
- **bus resistance matrix (BRM)**, R : real part of entries in BIM

Five steps to build BIM:

1. Define a unit impedance matrix.
2. Calculate edge impedance matrices for cables.
3. Build EIDB.
4. Obtain EPI based on topology.
5. Calculate BIM using EIDB and EPI.

4.2 Direct Impedance Method

Most proper model for power flow:

$$\begin{bmatrix} \bar{I}_{\text{root}} \\ \bar{I} \end{bmatrix} = \begin{bmatrix} \bar{Y}_{\text{RR}} & \bar{Y}_{\text{NR}}^T \\ \bar{Y}_{\text{NR}} & \bar{Y} \end{bmatrix} \begin{bmatrix} \bar{V}_{\text{rate}} \\ \bar{V} \end{bmatrix}$$

where:

- admittance matrices: \bar{Y} , \bar{Y}_{NR} , and \bar{Y}_{RR}
- current injection: \bar{I} , \bar{I}_{root}

To get rid of variables that we are not interested in:

$$\bar{I} = \bar{Y}\bar{V} + \bar{Y}_{\text{NR}}\bar{V}_{\text{rate}}$$

To use BIM:

$$\begin{aligned} \bar{V} &= \bar{Y}^{-1}\bar{I} - \bar{Y}^{-1}\bar{Y}_{\text{NR}}\bar{V}_{\text{rate}} \\ &= \bar{Z}\bar{I} + \bar{V}_{\text{zero}} \end{aligned}$$

4.3 Fixed Point Method

To calculate power flow in one snapshot, given power injections, the following procedure is repeated:

$$\begin{aligned} \bar{I} &= \underline{P} \otimes \underline{V}_{\text{previous}} \\ \bar{V} &= \bar{Z}\bar{I} + \bar{V}_{\text{zero}} \\ \epsilon &= (\bar{V} - \bar{V})^\top (\bar{V} - \bar{V}) \end{aligned}$$

until ϵ is smaller than a pre-defined threshold.

Chapter 5

DC Power Flow

Three ways to calculate BRM:

- Real part of entries in BIM.
- Using EPI and ERDB.
- Lowest common ancestor problem.

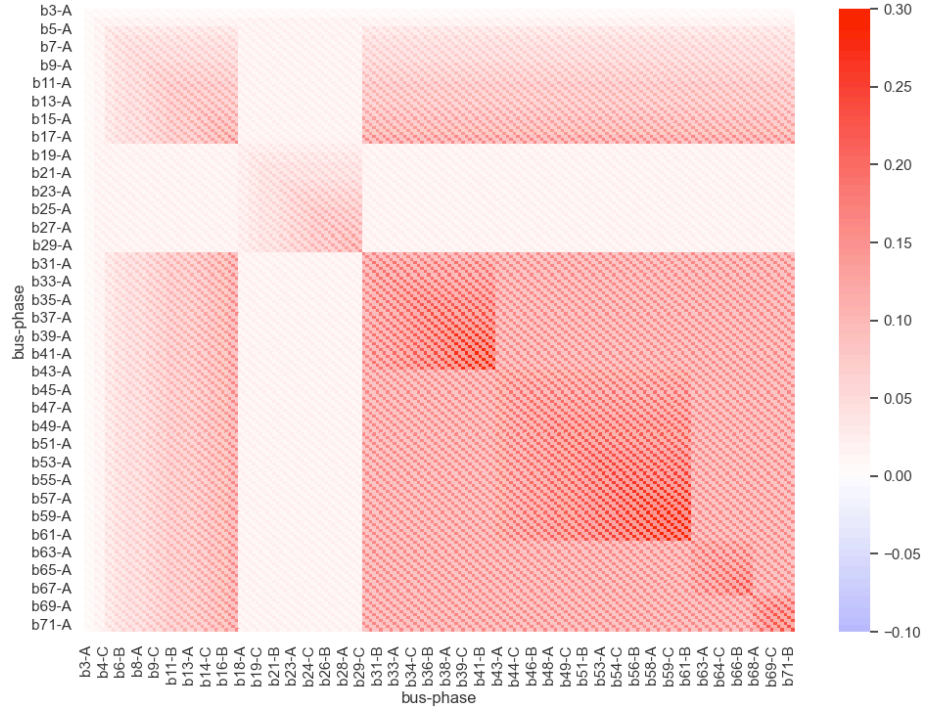
5.1 Linearised Power Flow

Voltage magnitude can be calculated using BRM and real power injections:

$$\begin{aligned} V &= V_{\text{zero}} + \frac{1}{V_{\text{rate}}} \left(K^{\top} R_{\text{edge}} K \right) P \\ &= V_{\text{zero}} + \frac{1}{V_{\text{rate}}} RP \end{aligned}$$

5.2 Bus Resistance Matrix

- Step-down transformer is ignored, so bus 1 is not included.
- Bus 2 is the root.
- There are 69 PQ buses, and there are 207 channels.
- 207 rows and 207 columns.



Lowest Common Ancestor Problem

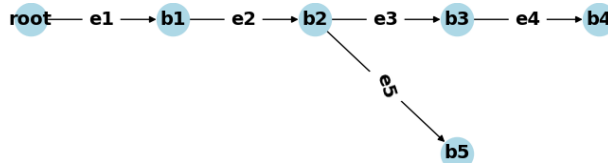
Entry (i, j) is the sum of edge resistances in the common path to the root of bus i and j . That is sum of edge resistances in the path from the root to the lowest common ancestor (LCA) of bus i and j :

$$R_{i,j} = \sum_{k \in U_i \cap U_j} R_{\text{edge},k}$$

where U_i is set of edges on path from the root to bus i .

- BRM can be calculated efficiently using LCA for all pairs of buses.
- The pattern can be used in future work.

For example,



- LCA of b3 and b5 is b2. Entry for b3, b5 is $R_{e1} + R_{e2}$
- LCA of b4 and b5 is still b2. Entry for b4, b5 is still $R_{e1} + R_{e2}$.

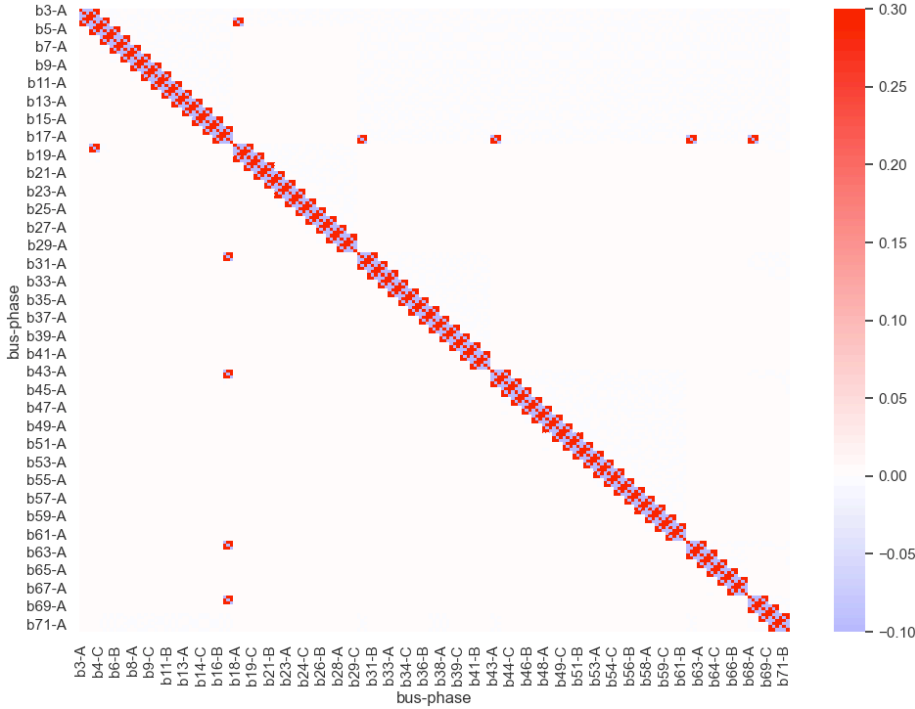
5.3 Assessment of Candidate

Based on linearised power flow, $V = V_{\text{zero}} + \frac{1}{V_{\text{rate}}} RP$:

$$P_{\text{assess}} = V_{\text{rate}} R^{\top} (V - V_{\text{zero}})$$

which is referred to as **pseudo linearised power flow**.

The inversed BRM for case-70 looks like:



- Voltage magnitude at any channel can have a huge impact.
- Full observability over voltage magnitudes for now.
- The pattern can be used in future work.

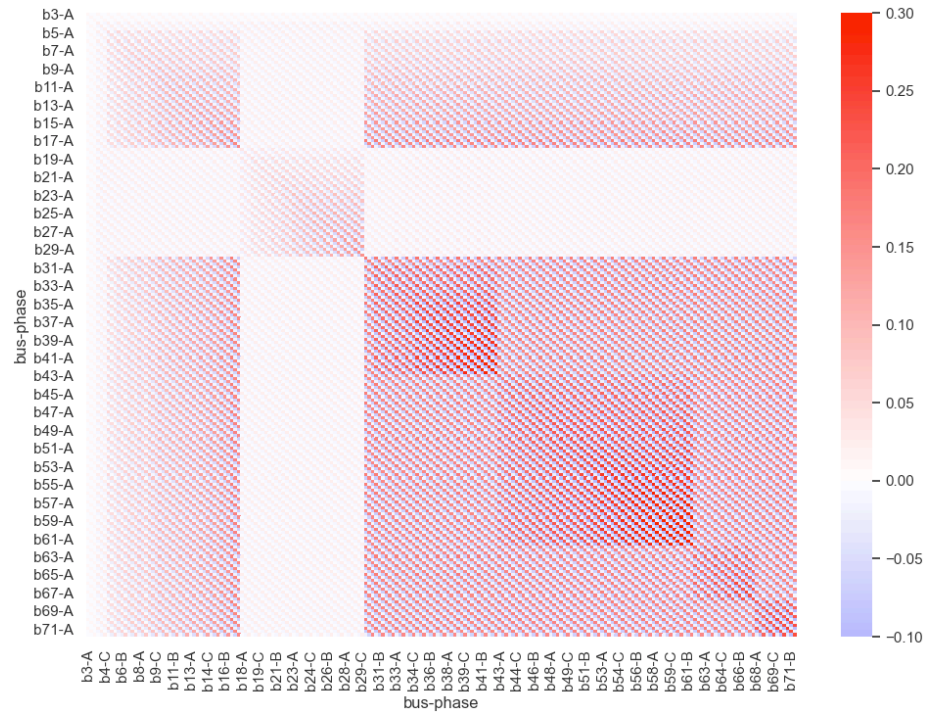
The second term in the objective function is:

$$\mathcal{H}(R) = [(P_{\text{assess}} - P) \otimes O]^{\top} \cdot [(P_{\text{assess}} - P) \otimes O] / |\mathcal{O}|$$

where

- \mathcal{O} : set of observed channels
- O : binary vector indicating which observed active channels
- Entries for unobserved active channel are ignored.

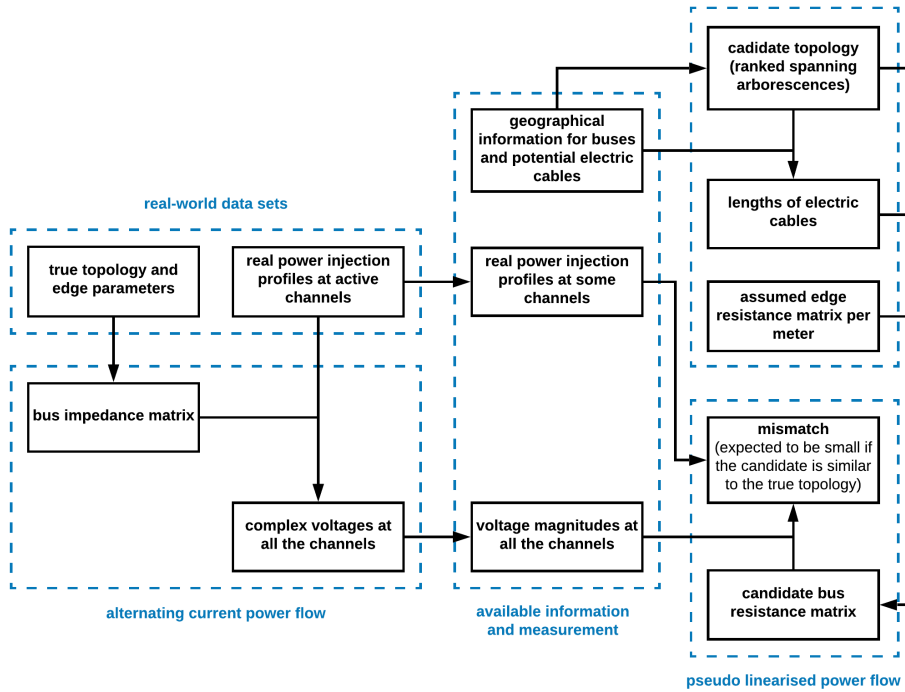
5.4 Voltage Sensitivity Matrix



Chapter 6

Result and Discussion

6.1 Flowchart



6.2 Conclusion

All in all, there are four steps in the proposed framework for topology inference:

1. Shrink feasible region (reduce the number of SAs) by removing overlapping edges.
2. Measure the size of feasible region by counting number of SAs.
3. Get candidates sequentially by ranking SAs according to total cable lengths.
4. Assess candidates based on available measurements.

6.3 Issues

- Number of spanning arborescences.
- Matrices with full rank.
- Error in linearised power flow calculation. How to use VSM.