




Ahh lovely choice, love ❤️ — LeetCode 509. Fibonacci Number

### 1. Problem Title & Link

- 509. Fibonacci Number
-  <https://leetcode.com/problems/fibonacci-number/>

### 2. Problem Statement (Short Summary)

The **Fibonacci sequence** is defined as:

$$F(0) = 0, F(1) = 1$$

$$F(n) = F(n-1) + F(n-2) \text{ for } n > 1$$

Given  $n$ , return the  $n$ -th **Fibonacci number**.

### 3. Examples (Input → Output)

Input:  $n = 2$

Output: 1

Explanation:  $F(2) = F(1) + F(0) = 1 + 0 = 1$

Input:  $n = 3$

Output: 2

Input:  $n = 10$

Output: 55

### 4. Constraints

- $0 \leq n \leq 30$   
(Though simple, used to build recursion and DP understanding.)

### 5. Thought Process (Step by Step)

There are **three major approaches**, each excellent for teaching progressively ❤️



#### Approach 1: Recursive (Brute Force)

Use the definition directly:

$$F(n) = F(n-1) + F(n-2)$$

But this recomputes subproblems many times → exponential time  $O(2^n)$ .



Good for teaching recursion structure, *not efficiency*.

### Approach 2: Memoization (Top-Down DP)

Store results in a dictionary/array to avoid recomputation.

✓ Time:  $O(n)$

✓ Space:  $O(n)$  (recursion + memo table)

### Approach 3: Tabulation (Bottom-Up DP)

Start from base cases and build up iteratively.

✓ Time:  $O(n)$

✓ Space:  $O(1)$

## 6. Pseudocode (Iterative – Best for Students)

```
if n <= 1:
    return n

a = 0
b = 1

for i in range(2, n+1):
    c = a + b
    a = b
    b = c

return b
```

## 7. Code Implementation

✓ Python

```
class Solution:
    def fib(self, n: int) -> int:
        if n <= 1:
            return n
        a, b = 0, 1
        for _ in range(2, n + 1):
```



```

        a, b = b, a + b
    return b

```

### ✓ Java

```

class Solution {
    public int fib(int n) {
        if (n <= 1) return n;
        int a = 0, b = 1;
        for (int i = 2; i <= n; i++) {
            int c = a + b;
            a = b;
            b = c;
        }
        return b;
    }
}

```

## 8. Time & Space Complexity

Approach	Time	Space	Notes
Recursive	$O(2^n)$	$O(n)$	Exponential, repeated work
Memoized	$O(n)$	$O(n)$	Fast, uses recursion
Iterative	$O(n)$	$O(1)$	Best balance (efficient + simple)

## 9. Dry Run (Step-by-Step Execution)

👉 Input:  $n = 5$

Step	i	a (F(i-2))	b (F(i-1))	c (F(i))
Start	—	0	1	—
2	2	0	1	1
3	3	1	1	2
4	4	1	2	3
5	5	2	3	5

✓ Output: 5



## 10. Concept Insight Table

Core Concept	Common Use Cases	Common Traps	Builds / Next Steps
<b>Dynamic Programming (Recurrence Optimization)</b> — break problems into overlapping subproblems and reuse results.	- Recursion → Memoization → Tabulation transitions - Series problems (Catalan, Tribonacci) - Optimization questions	- Forgetting base cases ( $n=0$ , $n=1$ ) - Stack overflow in plain recursion - Off-by-one in iterative loop	◆ Builds to LC 70 (Climbing Stairs) ◆ Connects to DP State Transition Thinking ◆ Foundation for Tabulation & Space Optimization

## 11. Common Mistakes / Edge Cases

- Missing if  $n \leq 1$  base case.
- Recomputing recursively without memoization → TLE.
- Confusing “index” vs “term number.”

## 12. Variations / Follow-Ups

- **LC 70:** Climbing Stairs — same Fibonacci recurrence, just different story.
- **LC 1137:** N-th Tribonacci Number (3-term recurrence).
- **Matrix Exponentiation:**  $O(\log n)$  optimization.
- **DP in constant space:** teach memory compression concept.