



## LeetCode 347 — Top K Frequent Elements

(Link: <https://leetcode.com/problems/top-k-frequent-elements/>)

### 1. Problem Title & Link

**Title:** LeetCode 347: Top K Frequent Elements

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### 2. Problem Statement (Short Summary)

Given an array of integers, return the **k most frequent elements**.

The output **order does not matter**, only the top k highest-frequency elements must be returned.

### 3. Examples (Input → Output)

#### Example 1

Input: nums = [1,1,1,2,2,3], k = 2

Output: [1,2]

Explanation:

- freq(1) = 3
- freq(2) = 2
- Top 2 frequent → 1 and 2

#### Example 2

Input: nums = [1], k = 1

Output: [1]

#### Example 3

Input: nums = [4,4,4,5,5,6], k=1

Output: [4]

### 4. Constraints

- $1 \leq \text{nums.length} \leq 10^5$
- $-10^4 \leq \text{nums}[i] \leq 10^4$
- $1 \leq k \leq \text{unique elements count}$
- Must run **better than  $O(n \log n)$**  (i.e.,  $O(n)$  or  $O(n \log k)$ )

### 5. Core Concept (Pattern / Topic)

★ Heap / Bucket Sort / Hashmap Counting

### 6. Thought Process (Step-by-Step Explanation)

**Brute Force Approach (Slow)**



1. Count frequency of each element
2. Sort by frequency (descending)  $\rightarrow O(n \log n)$
3. Pick first  $k$   $\rightarrow$  works but too slow for large inputs

## Optimized Thinking

We don't need full sorting.

We only need **top  $k$** , not sorted frequency list.

### Better Approach 1 — Min Heap ( $O(n \log k)$ )

- Push (frequency, element) to a **min-heap**
- Keep heap size  $\leq k$
- Pop smallest when exceeding
- Remaining  $k$  elements = top frequent

### Better Approach 2 — Bucket Sort ( $O(n)$ )

- Maximum frequency =  $n$
- Make buckets: index = frequency
- Fill buckets based on freq
- Traverse from high freq  $\rightarrow$  low freq
- Pick  $k$  elements
- Achieves  **$O(n)$**

## 7. Visual / Intuition Diagram

### Frequency Count

nums = [1,1,1,2,2,3]

freq:

1  $\rightarrow$  3

2  $\rightarrow$  2

3  $\rightarrow$  1

### Bucket Sort

Index: 0 1 2 3

Bucket: \_ [3] [2] [1]

Meaning:

3  $\rightarrow$  [1]

2  $\rightarrow$  [2]

1  $\rightarrow$  [3]

Pick from rightmost until  $k$ .

## 8. Pseudocode (Language Independent)



```

count = frequency map of nums
bucket = array of size n+1

for (num, f) in count:
    bucket[f].append(num)

result = []
for freq from n down to 1:
    for num in bucket[freq]:
        result.append(num)
        if len(result) == k:
            return result

```

## 9. Code Implementation

### ✓ Python

```

from collections import Counter

class Solution:
    def topKFrequent(self, nums, k):
        freq = Counter(nums)
        bucket = [[] for _ in range(len(nums) + 1)]

        for num, f in freq.items():
            bucket[f].append(num)

        res = []
        for f in range(len(bucket) - 1, 0, -1):
            for num in bucket[f]:
                res.append(num)
                if len(res) == k:
                    return res

```

### ✓ Java

```

import java.util.*;

class Solution {
    public int[] topKFrequent(int[] nums, int k) {
        Map<Integer, Integer> freq = new HashMap<>();
        for (int num : nums)
            freq.put(num, freq.getOrDefault(num, 0) + 1);

```



```

List<Integer>[] bucket = new ArrayList[nums.length + 1];
for (int i = 0; i <= nums.length; i++)
    bucket[i] = new ArrayList<>();

for (int num : freq.keySet())
    bucket[freq.get(num)].add(num);

int[] res = new int[k];
int idx = 0;

for (int f = nums.length; f >= 1; f--) {
    for (int num : bucket[f]) {
        res[idx++] = num;
        if (idx == k) return res;
    }
}
return res;
}
}

```

## 10. Time & Space Complexity

### Bucket Sort Approach

- **Time:**  $O(n)$
- **Space:**  $O(n)$

Why?

We store frequencies and buckets up to  $n$ .

## 11. Common Mistakes / Edge Cases

- ✗ Sorting whole list  $\rightarrow O(n \log n)$  (too slow)
- ✗ Forgetting that multiple numbers can have same frequency
- ✗ Using max-heap of full size instead of  $k$
- ✗ Returning sorted order (order doesn't matter)

Edge cases:

- ✓  $k = 1$
- ✓ All elements same
- ✓ All elements unique

## 12. Detailed Dry Run



nums = [1,1,1,2,2,3], k=2

Step	Action	Structure
Count freq	{1:3, 2:2, 3:1}	freq
Bucket fill	bucket[3] = [1], bucket[2] = [2], bucket[1] = [3]	bucket
Traverse	Start from freq=6→0	
freq=3	Add 1	res=[1]
freq=2	Add 2	res=[1,2] → STOP

Output: [1,2]

### 13. Common Use Cases (Real-Life / Interview)

- Finding top-k frequent search queries
- Top-k frequent logs / error codes
- Frequency-based ranking
- Word frequency counters

### 14. Common Traps (Important!)

- ⚠ Using dictionary sorting incorrectly
- ⚠ Forgetting to handle negative numbers
- ⚠ Not considering multiple numbers having same freq
- ⚠ Min-heap popping incorrectly

### 15. Builds To (Related LeetCode Problems)

- LC 451 — Sort Characters By Frequency
- LC 692 — Top K Frequent Words
- LC 973 — K Closest Points
- LC 215 — Kth Largest Element (Heap mastery)

### 16. Alternate Approaches + Comparison

Approach	Time	Space	When to Use
Bucket Sort	$O(n)$	$O(n)$	Best for this problem
Min Heap	$O(n \log k)$	$O(k)$	When k is small
Sorting	$O(n \log n)$	$O(n)$	Simple but slower



## 17. Why This Solution Works (Short Intuition)

Frequencies range from 1 to  $n$ , so we group elements by frequency and pick from highest frequency bucket downward — giving top  $k$  efficiently without sorting.

## 18. Variations / Follow-Up Questions

- What if the output must be **sorted by frequency**?
- What if we need **streaming input**? (Use min-heap)
- What if  $k$  is very large?
- Return top  $k$  least frequent?
- Return top  $k$  frequent words (string version)?