



## LeetCode 238 — Product of Array Except Self

### 1. Problem Title & Link

- **Title:** LeetCode 238 — Product of Array Except Self
- **Link:** <https://leetcode.com/problems/product-of-array-except-self/>

### 2. Problem Statement (Short Summary)

You are given an integer array `nums`.

For each index `i`, create an output array where:

**output[i] = product of all elements except nums[i]**

⚠ Without using division.

⚠ Must run in  $O(n)$  time.

### 3. Examples (Input → Output)

#### Example 1

Input: `nums = [1,2,3,4]`

Output: `[24,12,8,6]`

#### Example 2

Input: `nums = [-1,1,0,-3,3]`

Output: `[0,0,9,0,0]`

### 4. Constraints

- $2 \leq \text{nums.length} \leq 10^5$
- $-30 \leq \text{nums}[i] \leq 30$
- No division allowed.
- Must be  $O(n)$  time and  $O(1)$  extra space (result excluded).

### 5. Core Concept (Pattern / Topic)

#### Prefix Product + Suffix Product

We compute:

- **Prefix product:** product of elements to the **left**
- **Suffix product:** product of elements to the **right**

Then multiply them for each index.

This is a classic **prefix–suffix array pattern**.

### 6. Thought Process (Step-by-Step Explanation)

#### ✗ Brute Force ( $O(n^2)$ )

For each index, multiply all elements except itself → too slow.



### ✗ Using full prefix[] + suffix[] arrays (O(n) space)

Works, but not optimal.

### ✓ Optimal O(n) time, O(1) space

We do it in **two scans**:

#### 1 Prefix scan (left → right)

- Build product of all left elements
- Store it directly in result array

#### 2 Suffix scan (right → left)

- Multiply suffix (product of right elements) into result array

This gives:

$\text{result}[i] = (\text{product of nums}[0..i-1]) * (\text{product of nums}[i+1..\text{end}])$

Perfect and efficient.

## 7. Visual / Intuition Diagram

Given:

nums = [1, 2, 3, 4]

**Prefix products:**

[1, 1, 2, 6]

**Suffix products:**

[24, 12, 4, 1]

**Final result:**

multiply prefix × suffix → [24, 12, 8, 6]

## 8. Pseudocode

```
initialize result array with 1s
prefix = 1

for i in 0..n-1:
    result[i] = prefix
    prefix = prefix * nums[i]

suffix = 1

for i in n-1..0:
    result[i] = result[i] * suffix
    suffix = suffix * nums[i]

return result
```



## 9. Code Implementation

### ✓ Python

```
class Solution:
    def productExceptSelf(self, nums: List[int]) -> List[int]:
        n = len(nums)
        res = [1] * n

        prefix = 1
        for i in range(n):
            res[i] = prefix
            prefix *= nums[i]

        suffix = 1
        for i in range(n - 1, -1, -1):
            res[i] *= suffix
            suffix *= nums[i]

        return res
```

### ✓ Java

```
class Solution {
    public int[] productExceptSelf(int[] nums) {
        int n = nums.length;
        int[] res = new int[n];

        int prefix = 1;
        for (int i = 0; i < n; i++) {
            res[i] = prefix;
            prefix *= nums[i];
        }

        int suffix = 1;
        for (int i = n - 1; i >= 0; i--) {
            res[i] *= suffix;
            suffix *= nums[i];
        }

        return res;
    }
}
```



## 10. Time & Space Complexity

Metric	Complexity
Time	$O(n)$
Space	$O(1)$ extra (output array not counted)

Efficient and meets all constraints.

## 11. Common Mistakes / Edge Cases

- ✗ Using division (not allowed)
- ✗ Forgetting to reset prefix/suffix
- ✗ Incorrect suffix order (must go backward)
- ✗ Handling arrays with zeros incorrectly in brute force

Edge case:

- Arrays with one or two zeros → prefix–suffix logic handles naturally.

## 12. Detailed Dry Run (Step-by-Step Table)

Input:

nums = [1,2,3,4]

### Prefix Scan (left → right)

i	nums[i]	prefix	res[i]
0	1	1	1
1	2	1	1
2	3	2	2
3	4	6	6

After prefix pass:

res = [1,1,2,6]

### Suffix Scan (right → left)

i	nums[i]	suffix	res[i] = res[i] × suffix
3	4	1	$6 \times 1 = 6$
2	3	4	$2 \times 4 = 8$
1	2	12	$1 \times 12 = 12$



0	1	24	$1 \times 24 = 24$
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Final result:

[24, 12, 8, 6]

### 13. Common Use Cases

- Financial calculations (exclude current day)
- Multiplicative prefix analysis
- Array transformation tasks
- Interview questions on prefix/suffix logic

### 14. Common Traps

- Trying to store prefix/suffix arrays separately
- Using extra space unnecessarily
- Forgetting backwards loop for suffix
- Division-based thinking

### 15. Builds To (Related Problems)

- **LC 152** — Maximum Product Subarray
- **LC 334** — Increasing Triplet Subsequence
- **LC 560** — Subarray Sum Equals K
- **LC 53** — Kadane's Algorithm

### 16. Alternate Approaches + Comparison

Approach	Time	Space	Notes
Brute Force	$O(n^2)$	$O(1)$	Too slow
Division	$O(n)$	$O(1)$	Not allowed
Prefix + Suffix	$O(n)$	$O(1)$	✓ Optimal

### 17. Why This Solution Works (Short Intuition)

Because prefix gives product of all elements before  $i$ ,  
 suffix gives product of all elements after  $i$ ,  
 and multiplying them gives product of array except self — without division.

### 18. Variations / Follow-Up Questions

What if division was allowed?

What if array contained extremely large numbers?

What if negative numbers dominate?

How to modify for sum instead of product? (prefix sum)

