



LeetCode 704 — Binary Search

1. Problem Title & Link

- **Title:** LeetCode 704 — Binary Search
- **Link:** <https://leetcode.com/problems/binary-search/>

2. Problem Statement (Short Summary)

You are given a **sorted array** `nums` and a `target`.

Return the **index** of `target` if it exists.

Otherwise return **-1**.

You **MUST** use **$O(\log n)$** binary search.

3. Examples (Input → Output)

Example 1

Input: `nums = [-1,0,3,5,9,12]`, `target = 9`

Output: 4

Example 2

Input: `nums = [-1,0,3,5,9,12]`, `target = 2`

Output: -1

Example 3

Input: `nums = [5]`, `target = 5`

Output: 0

4. Constraints

- $1 \leq \text{nums.length} \leq 10^4$
- $-10^4 \leq \text{nums}[i], \text{target} \leq 10^4$
- Array is **sorted in ascending order**
- Must use **$O(\log n)$** time

5. Core Concept (Pattern / Topic)

Binary Search

Technique to search in a sorted array by dividing the search space into half each time.

Key idea:

- Compute $\text{mid} = (\text{left} + \text{right}) // 2$
- If $\text{nums}[\text{mid}] == \text{target} \rightarrow$ return `mid`
- If $\text{target} < \text{nums}[\text{mid}] \rightarrow$ search left half
- Else \rightarrow search right half



6. Thought Process (Step-by-Step Explanation)

1. Start with left = 0, right = n-1
2. While left <= right:
 - Compute middle index
 - Compare with target
3. If mid value > target → reduce right
4. If mid value < target → increase left
5. Else return mid
6. If loop ends → return -1

Binary search guarantees $O(\log n)$ time.

7. Visual / Intuition Diagram

Example:

nums = [-1,0,3,5,9,12], target = 9

Step 1:

left=0, right=5

mid=2 → nums[2]=3 < 9 → search right

Step 2:

left=3, right=5

mid=4 → nums[4]=9 FOUND

Diagram:

[-1	0	3	5	9	12]
L			M		R

[-1	0	3	5	9	12]
			L	M	R

8. Pseudocode

```
left = 0
right = n-1

while left <= right:
    mid = (left + right) // 2
    if nums[mid] == target: return mid
    else if target < nums[mid]: right = mid - 1
    else: left = mid + 1
```



```
return -1
```

9. Code Implementation

✓ Python

```
class Solution:
    def search(self, nums: List[int], target: int) -> int:
        left, right = 0, len(nums) - 1

        while left <= right:
            mid = (left + right) // 2

            if nums[mid] == target:
                return mid
            elif nums[mid] < target:
                left = mid + 1
            else:
                right = mid - 1

        return -1
```

✓ Java

```
class Solution {
    public int search(int[] nums, int target) {
        int left = 0, right = nums.length - 1;

        while (left <= right) {
            int mid = left + (right - left) / 2;

            if (nums[mid] == target)
                return mid;
            else if (nums[mid] < target)
                left = mid + 1;
            else
                right = mid - 1;
        }

        return -1;
    }
}
```



10. Time & Space Complexity

Metric	Complexity
Time	$O(\log n)$
Space	$O(1)$

11. Common Mistakes / Edge Cases

- Forgetting $\text{left} \leq \text{right}$ (should NOT be $<$)
- Wrong mid formula causing overflow in Java $(\text{left} + (\text{right} - \text{left}) / 2)$ is safe)
- Infinite loop due to wrong pointer movement
- Not returning -1 for missing target

Edge cases:

- target at index 0
- target at last index
- array size 1

12. Detailed Dry Run (Step-by-Step Table)

Input:

nums = [-1,0,3,5,9,12]

target = 9

Step	left	right	mid	nums[mid]	Action
1	0	5	2	3	$3 < 9 \rightarrow \text{left} = 3$
2	3	5	4	9	MATCH \rightarrow return 4

Answer = 4

13. Common Use Cases

- Searching in sorted arrays
- Peak finding
- First/last occurrence finding
- Lower bound / upper bound problems
- Binary search on answer space

14. Common Traps

- Using $\text{while left} < \text{right}$ (wrong here)



- Updating left/right incorrectly
- Returning mid without checking

15. Builds To (Related Problems)

- LC 35 — Search Insert Position
- LC 278 — First Bad Version
- LC 374 — Guess Number
- LC 34 — First and Last Position
- LC 875 — Koko Eating Bananas (Binary search on answer)

16. Alternate Approaches + Comparison

Approach	Time	Space	Notes
Linear Search	$O(n)$	$O(1)$	Too slow
Binary Search	$O(\log n)$	$O(1)$	Optimal ✓

17. Why This Solution Works (Short Intuition)

Each comparison removes half of the search space, ensuring logarithmic time complexity.

18. Variations / Follow-Up Questions

- What if array is rotated? → LC 33
- Binary search on descending array
- Binary search on infinite array
- Find first/last position of target