



## LeetCode 39 — Combination Sum

### 1. Problem Title & Link

- **Title:** LeetCode 39 — Combination Sum
- **Link:** <https://leetcode.com/problems/combination-sum/>

### 2. Problem Statement (Short Summary)

You are given:

- An array of **distinct** integers candidates
- A target integer target

Find **all unique combinations** of candidates where the chosen numbers **sum to target**.

You may use **each number unlimited times**.

Order of numbers inside a combination does NOT matter.

### 3. Examples (Input → Output)

#### Example 1

Input: candidates = [2,3,6,7], target = 7

Output: [[2,2,3],[7]]

#### Example 2

Input: candidates = [2,3,5], target = 8

Output: [[2,2,2,2],[2,3,3],[3,5]]

#### Example 3

Input: candidates = [2], target = 1

Output: []

### 4. Constraints

- $1 \leq \text{candidates.length} \leq 30$
- $1 \leq \text{candidates}[i] \leq 200$
- All numbers are unique
- Unlimited usage of each number
- Must return unique combinations

### 5. Core Concept (Pattern / Topic)

#### Backtracking – Choose, Explore, Unchoose

- Explore combinations by picking candidates starting from a given index (so no duplicates).
- Allow re-use of same element by **not moving to next index** after picking.



## 6. Thought Process (Step-by-Step Explanation)

We use **DFS + backtracking**.

At any recursive call:

**We track:**

- current combination (path)
- remaining target
- start index (prevents permutations)

**Steps:**

1. If target == 0 → valid combination → add path to answer.
2. If target < 0 → invalid → stop exploring.
3. For each index from start to end:
  - Choose candidates[i]
  - Recurse with target - candidates[i]
  - After returning, remove last element (unchoose)

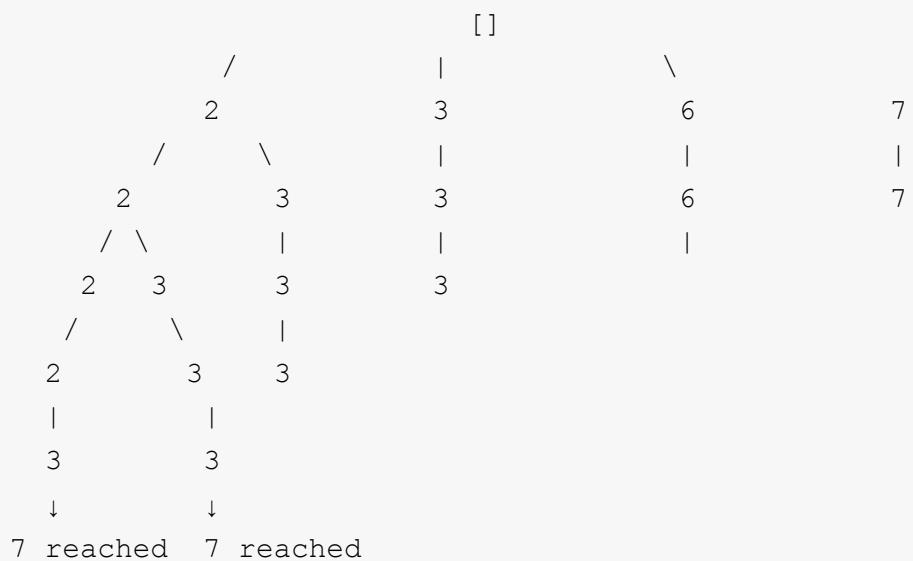
**Why i (same index) instead of i+1?**

Because numbers can be used unlimited times.

## 7. Visual / Intuition Diagram

Example: candidates = [2,3,6,7], target = 7

Backtracking tree (major branches) :



Valid:

[2,2,3]

[7]

## 8. Pseudocode

```
result = []
```



```

function backtrack(start, path, target):
    if target == 0:
        add copy(path) to result
        return
    if target < 0:
        return

    for i from start to len(candidates)-1:
        path.append(candidates[i])
        backtrack(i, path, target - candidates[i]) # reuse allowed
        path.pop()

call backtrack(0, [], target)
return result

```

## 9. Code Implementation

### Python

```

class Solution:
    def combinationSum(self, candidates: List[int], target: int) ->
List[List[int]]:
    res = []

    def backtrack(start, path, total):
        if total == target:
            res.append(path[:])
            return
        if total > target:
            return

        for i in range(start, len(candidates)):
            path.append(candidates[i])
            backtrack(i, path, total + candidates[i])
            path.pop()

    backtrack(0, [], 0)
    return res

```

### Java

```
class Solution {
```



```

public List<List<Integer>> combinationSum(int[] candidates, int target) {
    List<List<Integer>> res = new ArrayList<>();
    backtrack(0, candidates, target, new ArrayList<>(), res);
    return res;
}

private void backtrack(int start, int[] candidates, int remaining,
                      List<Integer> path, List<List<Integer>> res) {
    if (remaining == 0) {
        res.add(new ArrayList<>(path));
        return;
    }
    if (remaining < 0) return;

    for (int i = start; i < candidates.length; i++) {
        path.add(candidates[i]);
        backtrack(i, candidates, remaining - candidates[i], path, res);
        path.remove(path.size() - 1);
    }
}
}

```

## 10. Time & Space Complexity

| Metric | Complexity                                       |
|--------|--|
| Time   | $O(2^t)$ worst case (tree branching)             |
| Space  | $O(t)$ recursion stack + $O(\text{answer size})$ |

## 11. Common Mistakes / Edge Cases

- ✗ Using  $i+1$  instead of  $i \rightarrow$  prevents reuse
- ✗ Sorting candidates (unnecessary but ok)
- ✗ Forgetting to pop the path after recursion
- ✗ Adding path directly instead of making a copy

Edge cases:

- target smaller than all candidates  $\rightarrow$  return []
- one candidate equal to target



## 12. Detailed Dry Run (Step-by-Step)

Input:

candidates = [2,3,6,7]

target = 7

Start:

path=[], start=0, total=0

**Try 2:**

path=[2], total=2

Try 2 again:

path=[2,2], total=4

Try 2 again:

path=[2,2,2], total=6

Try 2 again:

path=[2,2,2,2], total=8 → >7 → backtrack

Try 3:

path=[2,2,3], total=7 → valid

Store → [2,2,3]

Backtrack to try next options.

**Try 3 as first:**

path=[3], total=3

path=[3,3], total=6

path=[3,3,3], total=9 → invalid

**Try 6:**

path=[6], total=6

path=[6,6], total=12 → invalid

**Try 7:**

path=[7], total=7 → valid

Final result:

[[2,2,3], [7]]

## 13. Common Use Cases

- Money/change formation
- Combination of items with repetition
- Unbounded knapsack background
- Recipe/mix formation problems

## 14. Common Traps



- Not passing correct start index
- Letting combinations get duplicated
- Forgetting to pop the last element
- Overcounting due to considering permutations

## 15. Builds To (Related Problems)

- **LC 40** — Combination Sum II (no repeats, duplicates allowed → harder)
- **LC 216** — Combination Sum III
- **LC 377** — Combination Sum IV (DP version)
- **LC 17** — Letter combinations of phone number (backtracking)
- **LC 78/90** — Subsets / Subsets with duplicates

## 16. Alternate Approaches + Comparison

| Approach        | Time                   | Space                           | Notes                         |
|-----------------|------------------------|---------------------------------|-------------------------------|
| Backtracking    | exponential            | $O(\text{target}/\text{depth})$ | ✓ Best                        |
| DP (count ways) | $O(n * \text{target})$ | $O(\text{target})$              | Counts ways, not combinations |
| BFS             | exponential            | high                            | Rarely used                   |

## 17. Why This Solution Works (Short Intuition)

The combination search tree explores adding each candidate unlimited times, but by always moving forward from start → end, we avoid permutations and generate only unique combinations.

## 18. Variations / Follow-Up Questions

- What if candidates have duplicates? (use LC 40 logic)
- What if numbers can only be used ONCE? (subset problem)
- What if order matters? (turn into DP counting problem)
- How to optimize with pruning? (sort + break early)