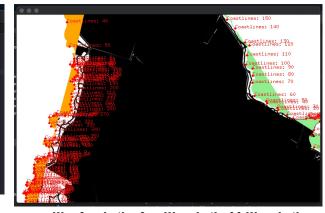


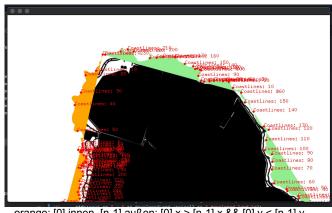
orange: [0] außen, [n-1] außen; [0].x < [n-1].x && [0].y < [n-1].y grün: -



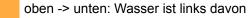
orange: [0] außen, [n-1] innen; [0].x > [n-1].x && [0].y < [n-1].y grün: [0] innen, [n-1] außen; [0].x > [n-1].x && [0].y > [n-1].y



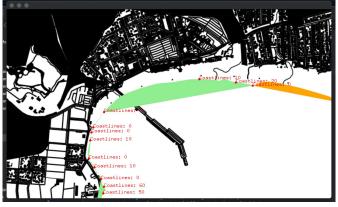
orange: [0] außen, [n-1] außen; [0].x > [n-1].x && [0].y < [n-1].y grün: [0] außen, [n-1] außen; [0].x > [n-1].x && [0].y > [n-1].y



orange: [0] innen, [n-1] außen; [0].x > [n-1].x && [0].y < [n-1].y grün: [0] außen, [n-1] innen; [0].x > [n-1].x && [0].y > [n-1].y



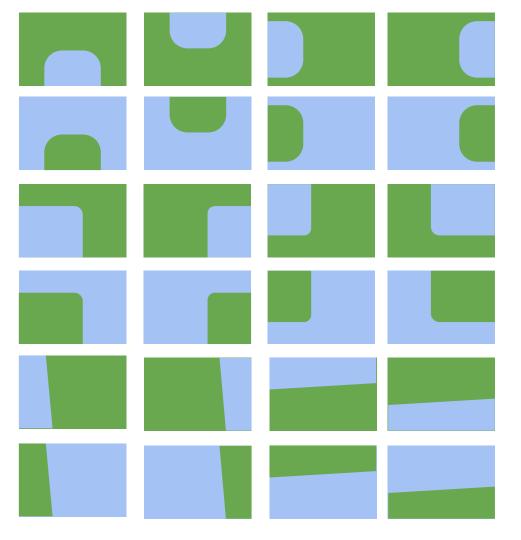
unten -> oben: Wasser ist rechts davon



orange: [0] innen, [n-1] außen; [0].x < [n-1].x && [0].y < [n-1].y grün: [0] außen, [n-1] innen; [0].x < [n-1].x && [0].y > [n-1].y

(0,0) = links oben; (x_max, y_max) = rechts unten

		[0] auße n	[0] innen	[n-1] außen	[n-1] innen	[0].x < [n-1].x	[0].x > [n-1].x	[0].y < [n-1].y	[0].y > [n-1]. y	polygon corners clockwise from top-left
orange: [0] außen, [n-1] außen; [0].x < [n-1].x && [0].y < [n-1].y grün: - orange: - grün: [0] außen, [n-1] außen; [0].x < [n-1].x && [0].y > [n-1].y		х		х		х		х		(0,0)-([0].xy)-([n-1].xy)-(0,max _y)
, manada, julia de la		х		х		х			х	([n-1].xy)-(max_x,0)-(max_x,m ax_y)-([0].xy)
orange: [0] außen, [n-1] innen; [0]. $x > [n-1].x && [0].y < [n-1].y$ grün: [0] innen, [n-1] außen; [0]. $x > [n-1].x && [0].y > [n-1].y$		x			х		х	х		([n-1].x,0)-(max_x,0)-([0].xy)-([n-1.xy])
orange: [0] außen, [n-1] außen; [0].x > [n-1].x && [0].y < [n-1].y			х	x			x		x	(0,0)-([0].x,0)-([0].xy)-([n-1].xy)
		x		х			х	х		(0,0)-([0].xy)-([n-1].xy)-(0,max _y)
grün: [0] außen, [n-1] außen; [0].x > [n-1].x && [0].y > [n-1].y		х		х			х		х	([n-1].xy)-(max_x,0)-(max_x,m ax_y)-([0].xy)
			х	х			х	х		(0,0)-([0].x,0)-([0].xy)-([n-1].xy) -(0,max_y) 5 PUNKTE !!!
orange: [0] innen, [n-1] außen; [0].x > [n-1].x && [0].y < [n-1].y grün: [0] außen, [n-1] innen; [0].x > [n-1].x && [0].y > [n-1].y		x			x		x		x	([n-1].x,0)-(max_x,0)-(max_x, max_y)-([0].x,max_y)-([n-1].xy) 5 PUNKTE !!!!
orange: [0] innen, [n-1] außen; [0].x < [n-1].x && [0].y < [n-1].y grün: [0] außen, [n-1] innen; [0].x < [n-1].x && [0].y > [n-1].y			х	х		х		х		([0].xy)-([n-1].xy)-(max_x,max _y)-([0].x,max_y)
		х			х	х			x	from bottom left: ([0].xy)-([n-1].xy)-([n-1].x,max _y)
oben -> unten: Wasser ist links davon unten -> oben: Wasser ist rechts davon (0,0) = links oben; (x_max, y_max) = rechts unten										



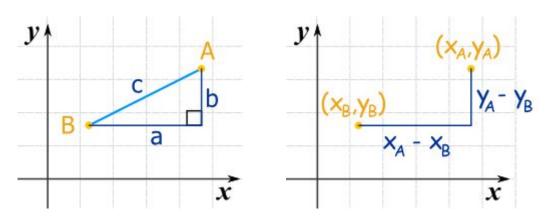
Land is always to the left of coastline ways, water is to the right. Coastline ways can never just stop in a dead-end, they always have to be connected to another coastline

Connecting the Lines.

- Connect lines of which the startpoint and endpoint are not both outside of the bbox (so called "through" lines).
 - Find a line where the startpoint is not equal to any endpoint. This would be the start of the connected line.
- Connect lines where the endpoint is the startpoint of another line.
- If there are still unconnected lines which are not a loop (island) or connected with another line, or is not a "through" line, then:
 - Take the startpoint of this unconnected line and look for a close endpoint in all the other lines.
 - Take the endpoint of this unconnected line and look for a close startpoint in all the other lines.

Determining the Distance between Points.

Pythagoras doesn't work on a round earth but that is ok here, we only need a rough measure anyway.



$$c = sqrt((x_a - x_b)^2 + (y_a - y_b)^2)$$

Determining the Distance between Points (in km).

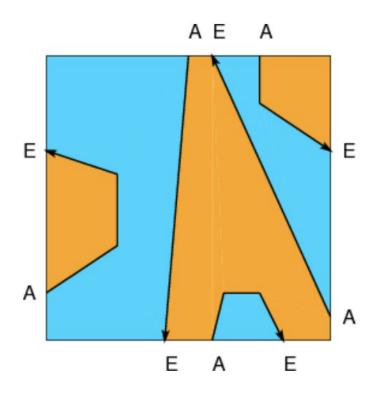
```
#define d2r (M_PI / 180.0)
//calculate haversine distance for linear distance
double haversine_km(double lat1, double long1, double lat2, double long2)
{
    double dlong = (long2 - long1) * d2r;
    double dlat = (lat2 - lat1) * d2r;
    double a = pow(sin(dlat/2.0), 2) + cos(lat1*d2r) * cos(lat2*d2r) * pow(sin(dlong/2.0), 2);
    double c = 2 * atan2(sqrt(a), sqrt(1-a));
    double d = 6367 * c;
    return d;
}
```

You can replace (M_PI / 180.0) with 0.0174532925199433 for better performance.

General information about geo and lines: https://www.movable-type.co.uk/scripts/latlong.html

Finding the Algorithm...

Some inspiration (maybe)

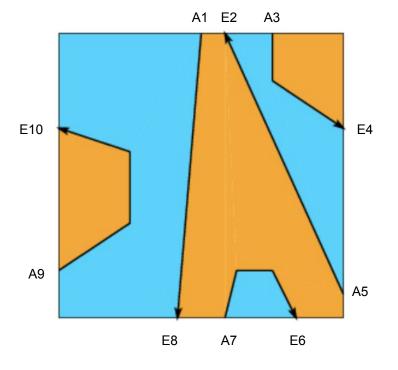


```
sl = []
Für alle offenen Küstenlinien k:
       Berechne Schnittpunkt s(A), s(E) mit der BoundingBox
       sl += (s(A), k, A)
       sl += (s(E), k, E)
m := Mittelpunkt der BoundingBox
Sortiere sl nach Winkel der Strecke [m, s]
0 = 0
s0, k0, ty0 = sl[i]
Falls ty0 == 'A':
      i += 1
       s0, k0, ty0 = sl[i]
Solange nicht alle Punkte aus sI verarbeitet:
       s1, k1, ty1 = sl[(i+1) \mod n]
       Verbinde k0 mit k1 durch die Strecke [s0, s1]
      i := (i + 2) \mod n
```

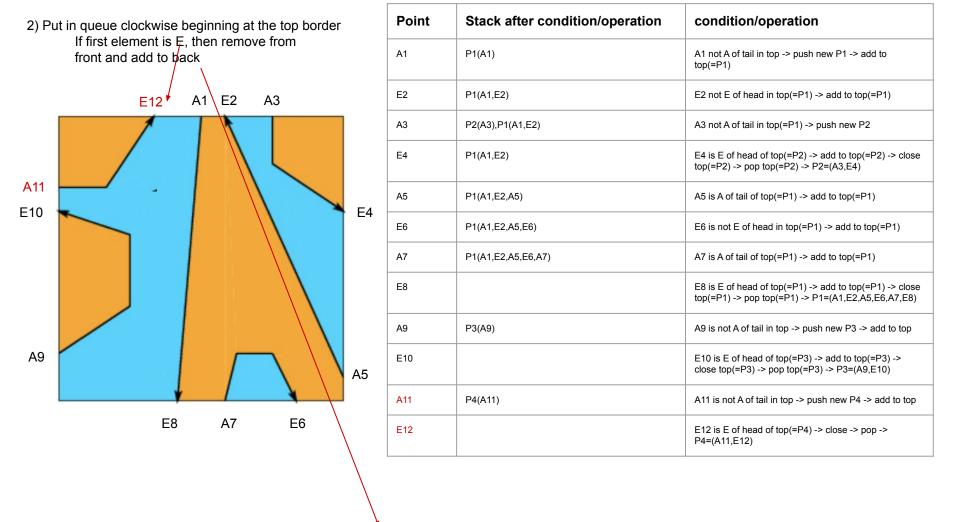
Source: https://monarch.gucosa.de/api/gucosa%3A19708/attachment/ATT-0/

Use a stack and queue, and go through the points

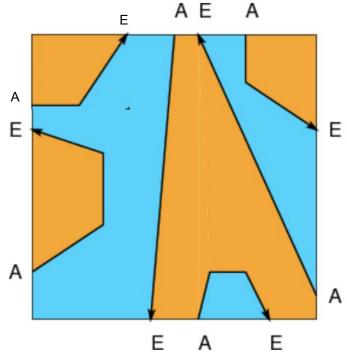
clockwise.



Point	Stack after condition/operation	condition/operation
A1	P1(A1)	A1 not A of tail in top -> push new P1 -> add to top(=P1)
E2	P1(A1,E2)	E2 not E of head in top(=P1) -> add to top
A3	P2(A3),P1(A1,E2)	A3 not A of tail in top(=P1) -> push new P2
E4	P1(A1,E2)	E4 is E of head of top(=P2) -> add to top(=P2) -> close top(=P2) -> pop top(=P2) -> P2=(A3,E4)
A5	P1(A1,E2,A5)	A5 is A of tail of top(=P1) -> add to top(=P1)
E6	P1(A1,E2,A5,E6)	E6 is not E of head in top(=P1) -> add to top(=P1)
A7	P1(A1,E2,A5,E6,A7)	A7 is A of tail of top(=P1) -> add to top(=P1)
E8		E8 is E of head of top(=P1) -> add to top(=P1) -> close top(=P1) -> pop top(=P1) -> P1=(A1,E2,A5,E6,A7,E8)
A9	P3(A9)	A9 is not A of tail in top -> push new P3 -> add to top
E10		E10 is E of head of top(=P3) -> add to top(=P3) -> close top(=P3) -> pop top(=P3) -> P3=(A9,E10)



The Algorithm.



Observations:

- Between E and A there is only water if they don't belong to the same coastline; If E and A are from different coastlines, then there is just water in between
- Consecutive A and E belong to the same landmass
- If you go clockwise only close path if E is end of the first A of the path

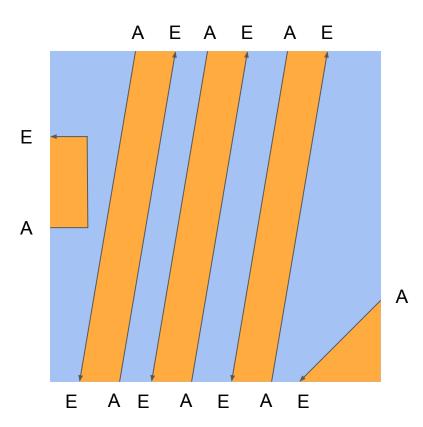
- 1) Calculate the border intersection points BIPs (DONE)
- 2) Put in queue q clockwise beginning at the top border If first element is E, then remove from front and add to back

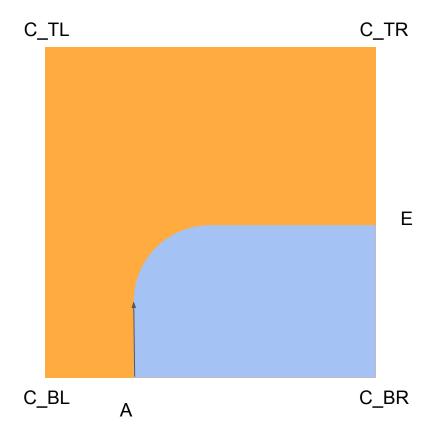
queue q; for all BIP on top: add BIP to q for all BIP on right: add BIP to q for all BIP on bottom: add BIP to q for all BIP on left: add BIP to q if q.front is E: remove q.front and add to q.back

3) Then run the algorithm:

```
stack<path> s;
while (!q.isempty())
    BIP = q.front
    check if BIP is A of tail of s.top
        no: s.push(new path); s.top.append(BIP)
        yes: s.top.append(BIP)
    check if BIP is E of head of s.top
        no: s.top.append(BIP)
        yes: s.top.append(BIP), close s.top, pop s.top
```

path is just a vector of BIPs





Check this only between A and E, because only between A and E (clockwise) there is landmass. Between E and A (clockwise) is water. And only if A and E are on different borders (no matter which borders).

A C_BL C_TL C_TR E

Add the corner points as long as the E is not between the current and the next corner point.