Characteristics of Sorting Algorithms

	Worst-Case Time	Average-Case Time	Best-Case Time	Worst-Case Space
BubbleSort				
	$O(n^2)$	$O(n^2)$	O(n)	O(n) total
* more writes to memory than SelectionSort	comparisons	comparisons	comparisons	
* adaptive				
* stable	O(n²) swaps	O(n²) swaps	O(1) swaps	O(1) auxiliary
* in-place	\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \			J
InsertionSort				
* more writes to memory than SelectionSort	O(n²) comparisons	O(n²) comparisons	O(n) comparisons	O(n) total
* adaptive				
* stable	$O(n^2)$	$O(n^2)$	O(1)	O(1)
* in-place	swaps	swaps	swaps	auxiliary
* online				
SelectionSort				
* double comparisons compared with InsertionSort	O(n²) comparisons	O(n²) comparisons	O(n²) comparisons	O(n) total
* not stable but can be modified				
* not adaptive	O(n) swaps	O(n) swaps	O(1) swaps	O(1) auxiliary
* in-place	зтарз	зтирэ	этирэ	иилнигу
HeapSort				
* not stable but can be modified	O(nlogn)	O(nlogn)	O(nlogn) distinct keys	O(n) total
* not adaptive			O(n)	O(1)
* in-place			equal keys	auxiliary

	Worst-Case Time	Average-Case Time	Best-Case Time	Worst-Case Space
MergeSort Divide&Conquer			O(nlogn)	O(n)
* not adaptive	O(nlogn)	O(nlogn)	typical	total
* stable			<i>O(n)</i>	O(n)
* not in-place in typical implementation (can be for linked lists)			natural variant	auxiliary with arrays
* accesses data sequentially and the need of random access is low				O(1) auxiliary with linked lists
* used for external sorting				
QuickSort Divide&Conquer				
* adaptive	$O(n^2)$	O(nlogn)	O(nlogn)	O(n)
* not stable but can be modified (typically not implemented stable)	comparisons		simple partition	auxiliary (naive)
* in-place	$O(n^2)$		<i>O</i> (<i>n</i>)	O(logn)
* with better partition strategy (choosing pivot) worst-case time is rare!	swaps		3-way partition with equal keys	(Sedgewick 1978)
* needs random access				