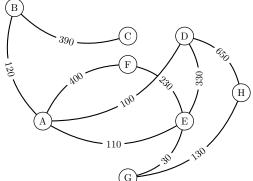
Graph Theory, Spring 2016, Homework 2

- 1. Show that G is connected if and only if we cannot find nonempty subgraphs H_1, H_2 such that G is a disjoint union of H_1 and H_2 .
- 2. Recall that c(G) is the number of components of the graph G, and for a vertex $v \in V_G$, G v is the graph obtained by removing the vertex v and all its incident edges (more formall, $G v = G[V_G \setminus \{v\}]$. Suppose G is a graph and $v \in V_G$ with $\deg(v) = 1$. Then show that c(G) = c(G v).
- 3. Recall that a bridge is an edge e in a graph G such that c(G e) > c(G). Show that if G has 7 vertices and is connected, then it must have at least 6 edges.
- 4. Show that if a simple graph G has 7 vertices, and at least 16 edges, then it must be connected, and if it has 7 vertices and at least 17 edges, then it cannot have any bridges.
- 5. (6000 level) More generally, show that if a simple graph G satisfies $e(G) > \binom{v(G)}{2}$ then it must be connected, and if $e(G) > \binom{v(G)}{2} + 1$ then it cannot have any bridges.
- 6. Consider the following graph, with weights assigned to the edges:



Apply Dijkstra's algorithm to find a minimal path from the vertex H to the vertex B. Illustrate the sequence of graphs T_0, T_1, \ldots which arise during the process.

7. (6000 level) Suppose that one uses Dijkstra's algorithm to, starting from a vertex v in a graph G, construct a tree T < G containing a minimal length path from the vertex v to some other vertex $v \in V_G$. Show that if $v' \in V_T$, and if W is a (v', w) path in T which does not pass through v, then $\ell(W) = d(v', w)$. That is, W is a minimal path from v' to w.