

Today: Suppose X scheme/ S , $R \hookrightarrow X \times_S X$ an étale eq. rel. $Y \equiv X/R$

Then: $\Delta: Y \rightarrow Y \times Y$ is representable.

(\Rightarrow from yesterday/earlier today that Y is an algebraic space).

let $j: U \rightarrow X$ open subscheme

let R_U be as in

$$\begin{array}{ccc} R_U & \longrightarrow & U \times U \\ \downarrow & & \downarrow \\ R & \longrightarrow & X \times X \end{array}$$

then $R_U \rightarrow U$ is étale & $\bar{j}: U/R_U \rightarrow X/R$

is open (rep.) immersion \tilde{Y}

$$\begin{array}{ccccc} & & \xrightarrow{\quad} & & \\ R_U & \longrightarrow & U \times U & \longrightarrow & Y \\ \downarrow & & \downarrow & & \downarrow \\ R & \longrightarrow & X \times X & \longrightarrow & X \\ & & \xrightarrow{\quad} & & \end{array}$$

pullback of rep. ét is rep. ét
 $\Rightarrow R_U \rightarrow U$ ét

\bar{j} is a monomorphism ✓

to see \bar{j} open, consider $f: T \rightarrow Y$

$$\begin{array}{ccc} \bigcirc & \xrightarrow{\text{open}} & T \\ \downarrow & & \downarrow \\ \cdot & & \cdot \end{array}$$

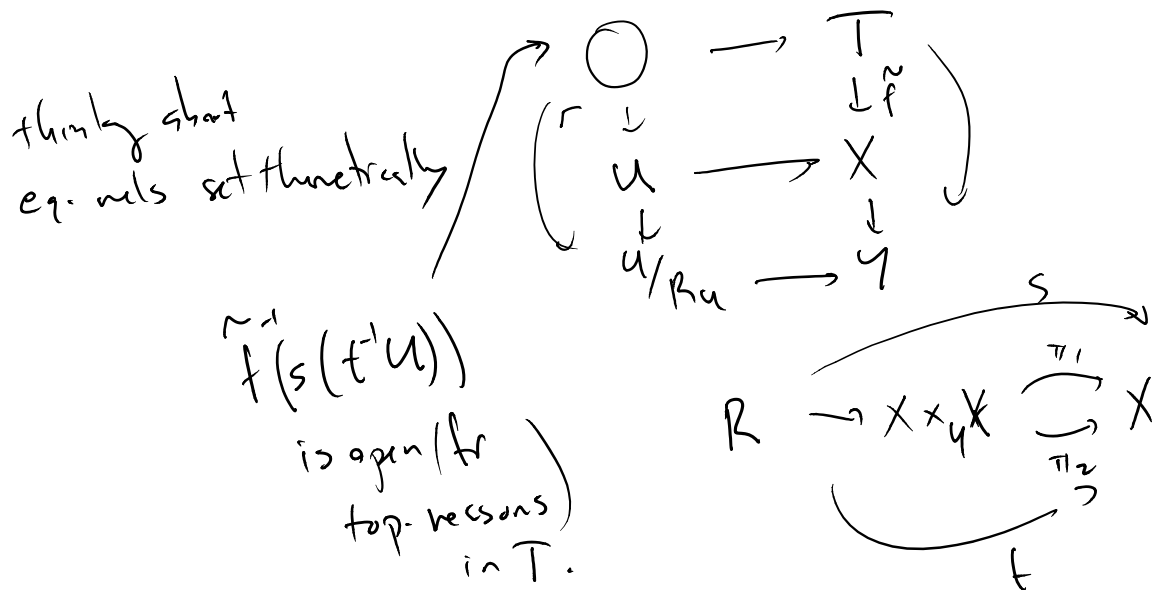
$$\begin{array}{ccc} \odot & \xrightarrow{\text{open}} & T \\ \downarrow & & \downarrow f \\ U/R_u & \xrightarrow{j} & Y = X/R \end{array}$$

to see that $\odot \rightarrow T$ open, can check

étale locally on $T \Rightarrow$ WLOG can assume

$T \xrightarrow{f} Y$ factors through X .

$\tilde{f} \downarrow_X \nearrow$ but now



Let's start to look more directly at
rep. of \mathcal{D}_Y .

$$\begin{array}{ccc} F & \longrightarrow & Y \\ \downarrow & & \downarrow \\ W & \longrightarrow & Y \times Y \end{array}$$

want to show F is a scheme.

"basic strategy:"

$$\begin{array}{ccc} F' & \longrightarrow & F \\ \downarrow r & & \downarrow \\ W' & \longrightarrow & W \end{array}$$

$$W' \rightarrow W \text{ étale}$$

$$W, W' \text{ affine.}$$

$$F' \rightarrow W'$$

$$g. \text{ affine.}$$

\Rightarrow ét. descent, f
 $g.$ affine morphisms
 that $F \rightarrow W$
 is given by $g.$ affine
 morph. of schemes.

$$\begin{array}{ccc} F & \longrightarrow & Y \\ \downarrow & & \downarrow \\ W & \longrightarrow & Y \times Y \end{array}$$

can work Zariski locally on $S \Rightarrow$ can assume S
 affine.

can work \mathbb{A}^1 locally on $W \Rightarrow$ can assume
 W affine.

$X \rightarrow Y$ ét. surj. of schemes

so is $X \times X \rightarrow Y \times Y \Rightarrow$ can find $W' \rightarrow W$
 étale cover s.t.

$$\begin{array}{ccc} W' & \longrightarrow & W \longrightarrow Y \times Y \\ & \searrow & \nearrow \\ & & \end{array}$$

moreover, can assume that

$$W' \times W' \rightarrow W' \times W' \rightarrow Y \times Y$$

$$w \searrow \quad \nearrow \\ X \times X$$

moreover, can also see that w' affine since w is affine & hence \mathbb{A}^1 -compact.

note we have:

$$\begin{array}{ccc} R & \longrightarrow & X \times X \\ \downarrow & \ulcorner & \downarrow \\ Y & \longrightarrow & Y \times Y \end{array}$$

\mathbb{A}^1 so

cart

\Rightarrow

$F' \rightarrow F$ rep étalé (pullback of rep)

$\Rightarrow R' = F' \times_F F'$
this is an étale eq. rel

$\hookrightarrow F'/R' = F$

via fiber product, we see F' is a scheme.

\mathbb{A}^1 $F' \rightarrow W'$ a monomorphism. $\Rightarrow F'$ separated because w' is sep. (affine)

etcetra.

\square

$$\begin{array}{ccccc} & & F' & \xrightarrow{\quad} & W' \\ & \swarrow & \downarrow & \ulcorner & \downarrow \\ & & F & \xrightarrow{\quad} & W \\ R & \xrightarrow{\quad} & X \times X & & \\ \downarrow & \swarrow & \downarrow & \searrow & \\ Y & \xrightarrow{\quad} & Y \times Y & & \end{array}$$