

From last time, we have two questions. First, why is the dimension of a motive well defined? Second, how does it behave under \otimes , \oplus ? We answer the 2nd first.

Direct Sum

Prop: Let M, N be motives. Then if $M + N$ are evenly (oddly) f.d., then so is $M \otimes N$.
 If M, N are f.d., then $\dim M \otimes N \leq \dim M + \dim N$.

Proof: Suppose $\dim_+ M = m$, $\dim_- N = n$. Then we aim to show $\Lambda^{n+m} (M \otimes N) = 0$. By a direct computation shows:

$$\Lambda^{n+m} (M \otimes N) \cong \bigoplus_{\substack{r+s \\ = n+m+1}} \Lambda^r M \otimes \Lambda^s N = 0$$

by an index count. An identical argument proves the converse. \blacksquare

Tensor Product

Def: For λ a partition of n , $M = (X, p, m)$, define $\pi_\lambda M = (X, d_\lambda(M), \text{op}, nm)$.

Lemma (Vanishing): Let $g \geq n$ and λ a partition of g . Then

- 1) $\text{Sym}^{n+1}(M) = 0 \Rightarrow \lambda_1 > n \Rightarrow \pi_\lambda M = 0$
- 2) $\Lambda^{n+1}(M) = 0 \Rightarrow \lambda_{n+1} \neq 0 \Rightarrow \pi_\lambda M = 0$.

Let T be the Young diagram according to λ . Define $R_\lambda(T) = \{\sigma \in S_n \mid \sigma \text{ permutes only } T\}$, $C_\lambda(T) = \{\sigma \in S_n \mid \sigma \text{ fixes } T\}$.

Now define $a_\lambda(T) = \sum_{\sigma \in R_\lambda(T)} \sigma$, $b_\lambda(T) = \sum_{\sigma \in C_\lambda(T)} \text{sign}(\sigma) \sigma$, $c_\lambda(T) = a_\lambda(T) b_\lambda(T)$.

Recall from Rep. Theory:

- 1) $V \in \text{Vect}_\mathbb{Q}$, and $S_n \subset V^{\otimes n}$ naturally. Then $\text{Im}(a_\lambda(T)) = \text{Sym}^{\lambda_1}(V) \otimes \cdots \otimes \text{Sym}^{\lambda_n}(V)$ and $\text{Im}(b_\lambda(T)) = \Lambda^{n+1} V \otimes \cdots \otimes \Lambda^{n+1} V$.
- 2) $c_\lambda(T) \circ c_\lambda(T) = \gamma_\lambda(T) c_\lambda(T)$, for $\gamma_\lambda(T) \neq 0$.
- 3) $R_{S_n} c_\lambda(T)$ is an irred. R_{S_n} -module.
- 4) $R_{S_n} c_\lambda(T) \cong R_{S_n} c_\mu(T) \Leftrightarrow \lambda = \mu$.
- 5) e_λ is a linear combination of the $c_\lambda(T)$.
- 6) $e_{(1, \dots, 1)} \cdot (e_\lambda \otimes e_\mu) = \begin{cases} 0 & \text{if } \lambda \neq \mu \\ e_{(1, \dots, 1)} & \text{else.} \end{cases}$

Now we have the following theorem:

Thm: Let M, N be f.d. motives. Then $\dim M \otimes N \leq \dim M \cdot \dim N$.

Proof: See Mumford. \blacksquare

Now we have the following goal: We had defined $\dim M = \dim_+ M + \dim_- M$ for some decomposition $M = M_+ \oplus M_-$, yet we said nothing about this decomposition. We now aim to show it's well-defined.

Before we do this however, let's take a brief detour to setup some tools. Recall that if X is Sm. Proj. / k , and $Z \subset X$ is a cycle, then we say Z is smash nilpotent, $Z \sim_0 0$ if $Z^n \sim_{\mathbb{Q}} 0$ on X^n for some $n \geq 1$.

Now in Mot_n , we have $\text{Hom}((X, p, m), (Y, q, n)) = g \circ \text{Corr}_n^{n-m}(X, Y) \circ p$, and the space of correspondences $\text{Corr}_n^r(X_d, Y) = C_n^{\text{dtr}}(X_d \times Y; \mathbb{Q})$, extended linearly. Since morphisms between motives are themselves cycles, we make the following definition:

Def: A morphism $f: M \rightarrow N$ in $\text{Mot}_n(k)$ is called smash-nilpotent if for some $n \geq 1$, the associated correspondence is smash-nilpotent w.r.t. \sim (may not be $\sim_{\mathbb{Q}}$!).

This means that if Γ_f is the correspondence for f , $\Gamma_f^n \times \Gamma_f^n \sim 0$ in $C_n(X^n \times Y^n)$. This is exactly the same as $f^{\otimes n} = f \times \dots \times f$ vanishing in the n^{th} tensor product of motives. We have the following obvious lemma:

Lemma: Let $f, g: M \rightarrow N$ be smash-nilpotent. Then so are $f \circ g, f - g$.

Proof: Let Γ_f, Γ_g be the associated correspondences. Then one checks that

$$(\Gamma_f + \Gamma_g)^n = \sum \binom{n}{r} \Gamma_f^r \times \Gamma_g^{n-r}$$

which can be made ~ 0 for sufficiently large n . The other is similar. \square

While this result was simple, the real focus of this detour is the following:

Theorem: Let $f: M \rightarrow M$ be a smash-nilpotent morphism in $\text{Mot}_n(k)$. Then $f^{(n)} = f \circ \dots \circ f = 0$. That is, smash-nilpotence \Rightarrow nilpotence.

This in turn is implied by

Prop: Let $f: M \rightarrow N$ in $\text{Mot}_n(k)$ be smash nilpotent of order n , and let $g_i: N \rightarrow M$, $i=1, \dots, n-1$ be morphisms. Then $f \circ g_{n-1} \circ f \circ \dots \circ f \circ g_1$ vanishes.

Clearly by taking $N = M$, $g_i = \text{id}$ we recover the theorem above, so now we prove the proposition.

Proof: To illustrate how this is proven, consider just $f \circ g_1 \circ f$. Denote their correspondences by Γ_f, Γ_{g_1} . Then by definition, if $M = (X, p, -)$ and $N = (Y, q, -)$, $f = g \circ \Gamma_f \circ p$ and $g_1 = p \circ \Gamma_{g_1} \circ q$. Hence $f \circ g_1 = g \circ \Gamma_f \circ p \circ \Gamma_{g_1} \circ q \in \text{Corr}(X, X)$. If we omit the projectors for a moment, and set π_{ijk} the projections from $X \times Y \times X \times Y$, s_{ij} the projections from $X \times Y \times X$, and p_{ij} projections from $X \times X \times Y$, then consider the cycles:

$$\alpha = \pi_{123}^*(s_{12}^* \Gamma_f \cdot s_{23}^* \Gamma_f) (= ((\Gamma_f \times X) \cdot (X \times \Gamma_f)) \times Y) \in C_n(X \times Y \times X \times Y)$$

$$\beta = p_{23}^* \Gamma_f (= X \times \Gamma_f) \in C_n(X \times X \times Y).$$

$$\text{Now } \alpha \cdot \pi_{134}^*(\beta) = (\Gamma_f \times X \times Y) \cdot (X \times \Gamma_f \times Y) \cdot (X \times Y \times \Gamma_f) = (\Gamma_f \times \Gamma_f) \cdot (X \times \Gamma_f \times Y) = 0 \text{ as } \Gamma_f \times \Gamma_f = 0.$$

Now use the projection formula: $O = (\pi_{134})_*(\alpha \cdot \pi_{134}^*(\beta)) = (\pi_{134})_* \alpha \cdot \beta$ and note that since $\pi_{134} = S_3 \times \text{id}_Y$, $(\pi_{134})_*(\alpha) = \pi_{123}^*(\Gamma_g \circ \Gamma_f) = (\Gamma_g \circ \Gamma_f) \times Y$. Then on $X \times X \times Y$, we have:

$$(\pi_{134})_* \alpha \cdot \beta = \{(\Gamma_g \circ \Gamma_f) \times Y\} \cdot (X \times \Gamma_f) = O.$$

finally, apply $(S_{13})_*$ to the above ($\pi_{134} = S_{13} \times \text{id}_Y$) to get $\Gamma_f \circ \Gamma_g \circ \Gamma_f = O$. \blacksquare

Now let us turn back to the finite dimensionality of motives, we begin with a crucial result.

Prop: Let M, N be two f.d. motives of different parity (M evenly, N oddly for example), and $f: M \rightarrow N$. Then f is smash nilpotent, $f^{\otimes \lambda} = O$ if $\lambda > \dim M \cdot \dim N$.

Proof: Set $m = \dim M$, $n = \dim N$, and $\lambda > m \cdot n$. Let λ, μ be two partitions of λ , and consider the composition:

$$M^{\otimes \lambda} \xrightarrow{d_\lambda} M^{\otimes \lambda} \xrightarrow{f^{\otimes \lambda}} N^{\otimes \lambda} \xrightarrow{d_\mu} N^{\otimes \lambda}.$$

where $d_\lambda = \Gamma_{e_\lambda}$ is the graph of the idempotent e_λ on $(X \times Y)$. Recall further that the projectors d_λ commute with morphisms, so the map above is equal to $f^{\otimes \lambda} \circ d_\lambda \circ d_\mu$. Since $e_\lambda \cdot e_\mu = O$ if $\lambda \neq \mu$, we see that we get

$$f^{\otimes \lambda} d_\lambda d_\mu = \begin{cases} O & \lambda \neq \mu \\ f^{\otimes \lambda} d_\lambda & \text{else.} \end{cases} \quad (\text{as } d_\lambda \text{ are idempotents}).$$

Hence it's enough to show the above composition vanishes for $\lambda = n+m+1$ and $\lambda = \mu$. Now suppose (for example) that $\Lambda^{n+1} M = O$ any $\text{Sym}^{n+1} N = O$. Then the vanishing lemma proves the claim. \blacksquare

Corollary: Suppose $M = (X, p, m)$ is both evenly and oddly f.d.. Then $M = O$.

Proof: Apply the above to p . \blacksquare

Now we can accomplish our goal.

Proposition: Let $M = (X, p, m) \cong M_+ \oplus M_-$ be a f.d. motive. If $M \cong M'_+ \oplus M'_-$ of even and odd f.d. motives, then $M'_+ \cong M_+$ and $M'_- \cong M_-$.

Proof: Suppose we had two decompositions. Write $p = p_+ + p_- = p'_+ + p'_-$. Then $p - p'_+$ is a projector which maps to M'_- , and hence $(p - p'_+) \circ p_+$ is smash-nilpotent, but goes from X to X , so must be nilpotent. Noting that $p = \text{id}_M$:

$$((p - p'_+) \circ p_+)^n = (p_+ - p'_+ \circ p_+)^n = O \Rightarrow p_+ = p_+^n = \underbrace{f(p_+, p'_+)}_{\text{Some expression.}} p'_+ \circ p_+$$

Since $p'_+ \circ p_+: M_+ \rightarrow M'_+$, we must have $f: M'_+ \rightarrow M_+$, i.e., a section. But this implies that $M'_+ = M_+ \oplus K$, and so $\dim M'_+ \geq \dim M_+$. Proving the opposite equality shows $\dim K = 0$, but this implies that $K = O$, so we're done \blacksquare

Surjections in Mot_n(k)

Let us work with Chow motives for now.

Def: Let $f: M \rightarrow N$ be a morphism of motives. Then f is surjective if for all smooth projective varieties Z the induced map $\text{CH}(M \otimes \text{ch}(Z))_{\mathbb{Q}} \rightarrow \text{CH}(N \otimes \text{ch}(Z))_{\mathbb{Q}}$ is surjective.

Let me remind you that for $M = (X, p, m)$, $\text{CH}^i(M) = \text{Im}(p_*: \text{CH}^{i+m}(X)_{\mathbb{Q}} \rightarrow \text{CH}^{i+m}(X)_{\mathbb{Q}}) \cong \text{Hom}_{\text{Mot}}(\mathbb{L}^{\otimes i}, M)$, $\mathbb{L} = (\text{Spec } k, \text{id}, -1)$.

Example: Let $\phi: X \rightarrow Y$ be a generically finite morphism of degree r . Then on motives, we have morphisms $\phi_* \dashv \phi^*$ s.t. $\phi_* \circ \phi^* = r \text{id}$. \Rightarrow surjective.

Example: Consider the inverse of a blow-up $X \xrightarrow{\delta} Y = \text{Bl}_p X$ of a sm. proj. X at a point. Then $\text{CH}^i(Y) = \text{CH}^i(X) \oplus \mathbb{Z}[E]$, and $E \notin \text{Im } \phi_*$ \Rightarrow not surjective. In general dominant morphisms are surjective, but not dominant rational maps.

Lemma: Let $f: (X, p, m) \rightarrow (Y, q, n)$ be a morphism. Then TFAE:

- i) f is surjective,
- ii) \exists a right inverse to f ,
- iii) $q = f \circ s$ for some $s \in \text{Corr}^0(Y, X)$.

Theorem: Let $f: M \rightarrow N$ be a surjective morphism of motives. If M is f.d., so is N .

Proof:

Step I: Suppose M is evenly (oddly) f.d. then the above lemma guarantees us a right inverse, $g: N \rightarrow M$, such that $fg = \text{id}_N$. This induces a decomposition $M = N \oplus K \Rightarrow N$ and K are evenly (oddly) f.d.

Step II: Write $M = M_+ \oplus M_-$. One needs to show existence of $N = N_+ \oplus N_-$ such that $M_+ \rightarrow N_+$ and $M_- \rightarrow N_-$. Since the degree doesn't matter in the definition, we may take degrees zero, and regard f as a correspondence. Using the above lemma and M 's decomposition, we get two endomorphisms $g'_\pm: N \rightarrow N$ ($M = (X, p, 0)$, $N = (Y, q, 0)$).

Step III: Show that there is a polynomial $P(t)$ such that $P(g'_\pm)$ are (almost) projectors. We set $g_\pm = g'^k_\pm \circ r_\pm$, $r_+ = P^k(g'_+)$, $r_- = P(g'^k_-)$.

Step IV: Show M_\pm surjects onto $(Y, g_\pm, 0)$. \square

This implies the following:

Corollary:

- 1) If $f: X \rightarrow Y$ is a dominant morphism with $\text{ch}(X)$ f.d., then $\text{ch}(Y)$ is also.
- 2) $M \oplus N$ f.d. \Rightarrow M and N are f.d.
- 3) A motive which is dominated by a morphism from a finite product of curves is f.d. In particular the motive of an abelian variety is f.d.
- 4) Every summand of a tensor product of curves is f.d. They form a full tensor subcategory.

Now we have the following theorem:

Thm: Let $M = (X, P, m)$ and $f: M \rightarrow M$ be a morphism of Chow motives. Assume M is evenly (oddly) finite dimensional. We have:

- 1) There exists a nonzero polynomial $g(t) \in \mathbb{Q}[t]$ of degree $n-1$ with $g(f)=0$.
- 2) If $f \sim_{\text{num}} 0$, then $g(t) = t^{n-1}$.

Proof omitted.

Applications and Conjectures

We have functors $\text{Mot}_{\text{rat}} \rightarrow \text{Mot}_{\text{hom}} \rightarrow \text{Mot}_{\text{num}}$ taking $(X, P, m) \rightarrow (X, P_{\text{num}}, m)$ $\rightarrow (X, P_{\text{num}}, m)$. The first functor is not fully faithful, as if $Z \in Z^i(X)$ is a cycle, then Z is given by a morphism $f: \mathbb{P}^{n-i} \rightarrow X$. If $Z \not\sim_{\text{rat}} 0$ but $Z \sim_{\text{num}} 0$, then this is an example. Explicitly, on an elliptic curve E , $Z = a - b$ is not rationally zero, but is algebraically (hence homologically) zero.

Def: Let M be a Chow motive and M_{hom} be its image in Mot_{hom} . If $M_{\text{hom}} = 0$ yet $M \neq 0$, we say M is a phantom motive.

Thm: If M is a f.d. Chow motive, it is not a phantom motive.

Proof: Suppose it was. Then $P_{\text{num}} \sim 0$, which implies $P_{\text{num}} \sim 0$. Since $M = M_+ \oplus M_-$, $P = P_+ + P_-$ are each numerically trivial, so $P_{\pm}^{n-1} = 0 \Rightarrow P_{\pm} = 0$, hence $M = 0$. \blacksquare

Recall that if $H_+(M) = 0$ then $H(\text{Sym}^n M) = \Lambda^n H_-(M)$, hence $H(\text{Sym}^n M) = 0$ if $n > \dim H(M)$, and if $H_-(M) = 0$ then $H(\Lambda^n M) = \Lambda^n H_+(M)$, hence $H(\Lambda^n M) = 0$ if $n > \dim H(M)$.

Corollary: If M is a f.d. motive, $\dim M = \dim H(M)$.

Proof: It's enough to see it for M evenly f.d., so set $n = \dim M$ and note that since $\Lambda^n M \neq 0$, $H(\Lambda^n M) \neq 0$. Indeed motivic cohomology is defined via the image of the projector, hence nonzero. The above remarks show $\dim M \geq \dim H(M) = \dim H_+(M)$, and $H_+(\Lambda^n M) = \Lambda^n H_+(M)$, so $\dim H_+(\Lambda^n M) \geq \dim M$. This shows equality. \blacksquare

Conjecture (Kimura, O'Sullivan): Every Chow motive is f.d.

Conjecture $N(M), N'(M)$: The ideal $J_M = \text{Ker}(End_{\text{rat}}(M) \rightarrow End_{\text{num}}(M))$ is nilpotent, resp a nil-ideal (every element is nilpotent).

Thm: Kimura + O'Sullivan $\Rightarrow N(M) \neq 0$ no phantom motives

Proof

It is a ^{non}commutative algebra fact that if a nil-ideal has the order of nilpotency uniformly bounded, then it is nilpotent. Now let f be an endomorphism of a f.d. motive which is numerically trivial. We aim to show its nilpotent.

We have $f = (p_+ + p_-) \circ f \circ (p_+ + p_-) = f_+ + f_- + f_{\text{mix}}$, f_{mix} does not preserve parity $\Rightarrow f^r \sim 0$, r independent of f . Then expanding f^n , we get as an average monomial:

$$f_{\pm}^{k_1} \circ f_{\text{mix}}^{\lambda_1} \circ \dots \circ f_{\text{mix}}^{\lambda_r} \circ f_{\pm}^{k_r} \quad i \leq \lambda_1 + \dots + \lambda_r \leq r-1.$$

Since f (hence f_{\pm}) are numerically trivial gives $f_{\pm}^s = 0$. Then $f^n = 0$ for $n \geq r_s$.

Now recall that Voevodsky's conjecture asks whether a numerically trivial $f: M \rightarrow M$ is smash nilpotent (note this implies actual nilpotence). We claim this implies Kimura-O'Sullivan. Indeed Voevodsky implies $D(X)$ ($\sim_{\text{hom}} = \sim_{\text{num}}$), but $D(X) \Rightarrow A(X, L)$. If this holds for all X , we get $B(X) \Rightarrow C(X)$. But Voevodsky (by def) implies $N'(X)$, and $C(X)$ implies $S(X)$ (the sign conjecture we skipped). It is now a theorem of Jannsen that this gives $\text{ch}(X)$ f.d. for all X .