

The Hypercubic Manifold in Homotopy Type Theory

Dylan Laird

04/09/2023

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 - Defining $\mathbb{H}\mathbb{M}^1$
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Homotopy Type Theory

- Vladimir Voevodsky : proposed the "Univalence Axiom"

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- Synthetic homotopy
- Computer assisted proofs
- "Univalent Foundations" as a foundational system for maths

Intro (2) : Homotopy theory

Let X be a topological space and I the unit real interval $[0, 1]$

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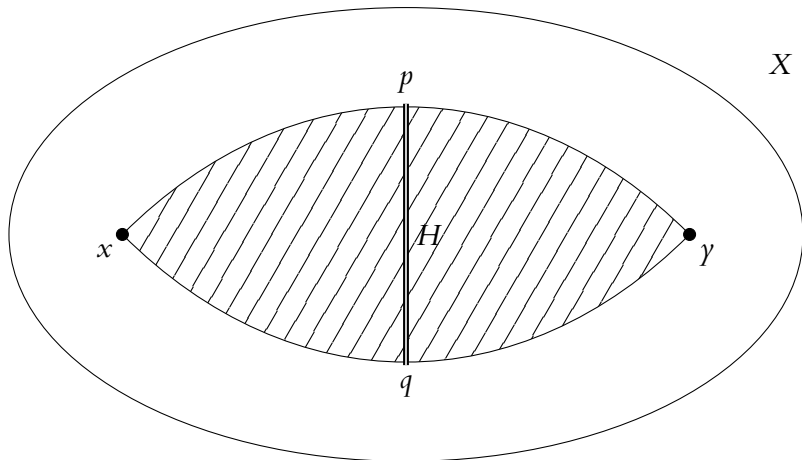
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- Two paths p and q are "homotopic" if there is "a continuous deformation" of p onto q
- Loops: a loop at point x is a path p such that $p(0) = p(1) = x$.

Intro (3) : Homotopy theory



Intro (4) : Homotopy theory

Paths operations

- inversion : given $p : x \rightarrow y$ there is a path $p^{-1} : y \rightarrow x$
- concatenation : given $p : x \rightarrow y$ and $q : y \rightarrow z$ there is a path $p \cdot q : x \rightarrow z$

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Fundamental group $\pi_1(X, x)$

Take a point x of X and consider loops at point x .

- Equality: homotopy (of loops)
- Composition : path concatenation (associative)
- Inverse element : inverse path
- Unit element : constant path at x

Two (path connected) homeomorphical spaces (more precisely, homotopically equivalent) have the same fundamental groups !

Higher homotopy groups

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Intro (5) : Homotopy theory

Higher homotopy groups

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Natural questions

- Given a finite group G is there a (path-connected) space X such that $\pi_1(X) = G$?

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- Is there a (path-connected) space such that $\pi_1(X) = G$ and $\pi_n(X) = 1$ for any $n \geq 2$? Answer : yes, the **Eilenberg Mac-Lane** space of G

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We manipulate types A, B and terms $a : A$ of some certain type.

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We distinguish between **definitional equality** \equiv used when defining objects and **propositional equality** which is a type theoretical concept.

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Propositional Equality

Given a type A and terms $a, b : A$ there is a type $a =_A b$. We say that elements a and b are (propositionally) equal when there exists some element :

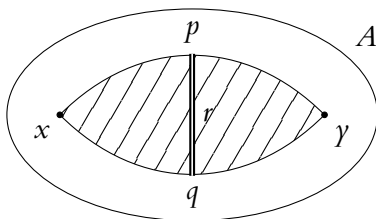
$$p : a =_A b$$

Intro (7) : What's the connection with Type Theory ?

| Types | Topology |
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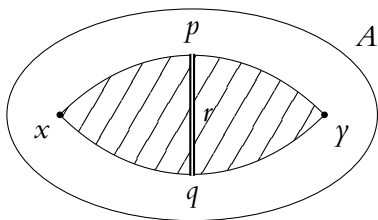
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Synthetic homotopy

This connection is at the core of synthetic homotopy theory which allows us to define every object we previously talked about in the synthetic framework of type theory.

Intro (8) : The Hypercubic Manifold $\mathbb{H}\mathbb{M}^1$

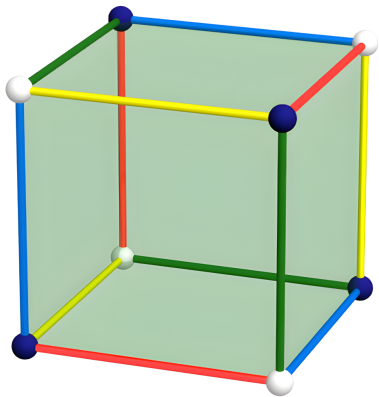


Figure: The hypercubic manifold $\mathbb{H}\mathbb{M}^1$ (analysis-situs)

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Goals

- We know how to define Eilenberg Mac-Lane spaces for finite groups in HoTT but the construction is not well understood
- We would like to achieve a "nicer" construction that would yield a handy induction principle
- This would open up to define "Higher Group" actions and computer checked cohomology calculations.
- It turns out that $\pi_1(\mathbf{HM}^1) = Q$ and that by computing the "homotopy fiber" of a certain map my advisors have a way to provide such a "nice" construction in the case of the group Q .

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Some constructions on types

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Given types A, B one has a type $A \rightarrow B$ of functions from A to B .

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Dependant pair Types

If B is a type family over a type A (that is a function $B : A \rightarrow \mathcal{U}$) we have a type of dependant pairs $\sum_{x:A} B(x)$.

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Coproduct types

Give types A, B one has a coproduct type $A + B$ given by the constructors :

- $\text{inl} : A \rightarrow A + B$
- $\text{inr} : B \rightarrow A + B$

Manipulating types

Introduction rules

They encapsulate how to build an element of a certain type.

- To build an element of $f : A \rightarrow B$ one needs an expression $\phi(x)$ such that $a : A \vdash \phi(a) : B$ and to set $f \equiv \lambda x. \phi(x)$
- To build an element of $A \times B$ one needs elements $a : A$ and $b : B$ to form $(a, b) : A \times B$

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Induction principles

They encapsulate how to build **dependent** functions from a source type A .

- To build a function of type $f : \prod_{x:A \times B} P(x)$ one only needs to give its value on pairs (a, b) .

A working example : product types (1)

Projections

We can define projections $\text{pr}_1 : A \times B \rightarrow A$ and $\text{pr}_2 : A \times B \rightarrow B$ by the **induction** principle for product types by setting $\text{pr}_1((a, b)) \equiv a$ and $\text{pr}_2((a, b)) \equiv b$.

Identity types (1)

Definition

Given a type A and elements $a, b : A$ there exists a type $a =_A b$ called the identity type of a and b . We say that " a is equal to b " if there exists some element $p : a =_A b$

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Reflexivity

For any type A and point $x : A$ there is an element $\text{refl}_x : x =_A x$

Identity types (2)

Path induction

To prove a predicate that depends on $x, y : A$ and $p : x =_A y$ one only needs to prove it in the case where $x \equiv y$ and $p \equiv \text{refl}$.

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Example (Functions preserve equals)

For any $f : A \rightarrow B$, $x, y : A$ and $p : x =_A y$ there is an element

$$f(p) : f(x) =_B f(y)$$

Identity types (3)

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- By path induction suppose that $x \equiv y$ and $p \equiv \text{refl}$
- We then have to provide an element of $f(x) =_B f(x)$
- We conclude by setting $f(\text{refl}_x) :\equiv \text{refl}_{f(x)}$

□

HoTT : The homotopical interpretation of Type Theory (1)

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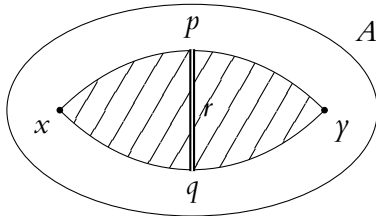


Figure: two points $x, \gamma : A$, two paths $p, q : x =_A \gamma$ and a homotopy $r : p =_{x=_A \gamma} q$

HoTT : The homotopical interpretation of Type Theory (2)

| Equality | Homotopy |
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- **Symmetry** : Given a path $p : x =_A y$, one has an inverse path $p^{-1} : y =_A x$ such that for any x , $\text{refl}_x^{-1} \equiv \text{refl}_x$.
- **Transitivity** : Given paths $p : x =_A y$ and $q : y =_A z$ one has a concatenated path $p \cdot q : x =_A z$ such that for any x , $\text{refl}_x \cdot \text{refl}_x \equiv \text{refl}_x$.

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There are additional properties that one can prove such as

$$p \cdot p^{-1} = \text{refl}$$

that justify this whole interpretation.

Equalities and dependent functions

Take $f : \prod_{x:A} P(x)$, $x, y : A$ and $p : x = y$.

Note that we can't have a path between $f(x) : P(x)$ and $f(y) : P(y)$.

Transports

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Transport

By path induction, for any x, y and $p : x = y$ one can define :

$$\text{transport}^P : \prod_{a,b:A} \prod_{p:a=A}^b P(a) \rightarrow P(b)$$

We may write $\text{transport}^P(a, b, p)$ as p_* if the context is clear.

Transports

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Example

The map p_* encapsulates how are transported elements of the "fiber" $P(x)$ to the fiber $P(\gamma)$ by going along the path $p : x = \gamma$.

Fibrations (1)

Fibration

Take a continuous map $p : E \rightarrow B$. It is called a fibration if it has a certain path lifting property from the **base space** B to the **total space** E .

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Fibrations in HoTT

Take a type family B over A (that is $B : A \rightarrow \mathcal{U}$).
Then, the type $\sum_{x:A} B(x)$ equipped with its first projection pr_1 is a fibration.

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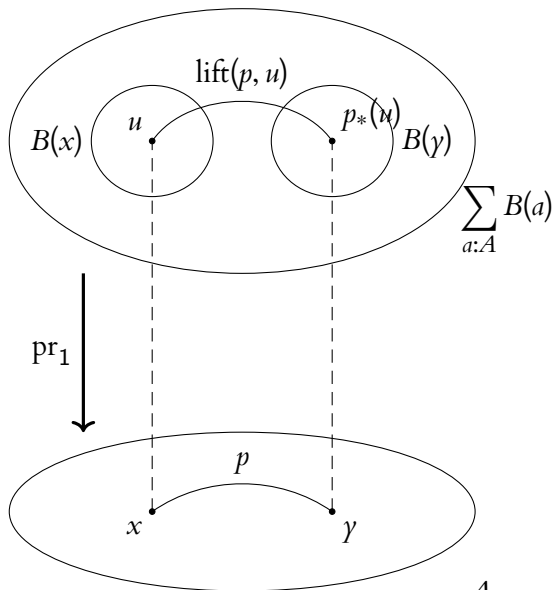
More precisely :

Given $x : A$ and $u : B(x)$, for any $y : A$ and $p : x = y$, one has a path

$$\text{lift}(p, u) : (x, u) = (y, p_*(u))$$

such that $\text{pr}_1(\text{lift}(p, u)) = p$.

Fibrations (2)



Paths over paths

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Dependent functions yield such paths :

$$\mathrm{apd}_f : \prod_{p:x=Ay} p_*(f(x)) = f(y)$$

Equivalences and Univalence

Equivalences

We can build them from quasi-equivalences. Two types are quasi-equivalent if there exists maps $f : A \rightarrow B$ and $g : B \rightarrow A$ such that :

$$\prod_{x:A} g(f(x)) = x \quad \text{and} \quad \prod_{y:B} f(g(y)) = y$$

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Univalence

One has :

$$(A \simeq B) \simeq (A = B)$$

Higher Inductive Types

Idea

When dealing with HITs we allow ourselves to use paths constructors, paths between paths constructors and so on.

Higher Inductive Types

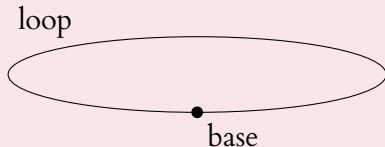
Idea

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The circle \mathbb{S}^1

The circle \mathbb{S}^1 is defined by the following HIT :

- $\text{base} : \mathbb{S}^1$
- $\text{loop} : \text{base} =_{\mathbb{S}^1} \text{base}$



HITs : Induction Principle

We have a problem

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The case of \mathbb{S}^1

To define a map $f : \prod_{x:\mathbb{S}^1} P(x)$, one needs to specify its value $f(\text{base}) : P(\text{base})$. Then, what should you send loop to ?

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Recall that by transport :

$$P(\text{base}) \stackrel{\text{loop}_*}{\simeq} P(\text{base})$$

Since f is a dependent function, it should send loop to a dependent path between $\text{loop}_*(f(\text{base}))$ and $f(\text{base})$.

Some other HITs (1)

Coequalizers

Given types B, A and maps $f, g : B \rightarrow A$ the (homotopy) **coequalizer** type $\text{CoEq}(f, g)$ is given by the HIT:

- $c : A \rightarrow \text{CoEq}(f, g)$
- $p : \prod_{b:B} c(f(b)) = c(g(b))$

$$B \begin{array}{c} \xrightarrow{g} \\ \xrightarrow{f} \end{array} A \xrightarrow{c} \text{CoEq}(f, g)$$

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The circle as a coequalizer

$$1 \begin{array}{c} \xrightarrow{1} \\ \xrightarrow{1} \end{array} 1 \xrightarrow{c} \text{CoEq}(1, 1) \simeq \mathbb{S}^1$$

Cubical Type Theory : Completing Squares

Filling property

Given an incomplete square of paths

$$\begin{array}{ccc} a & \xrightarrow{p} & b \\ q \downarrow & & \downarrow r \\ c & & d \end{array}$$

Then there is a unique path s from c to d that completes the square by making it commute, that is :

$$q \cdot s = p \cdot r$$

Some other HITs (2) : The torus \mathbb{T}^2

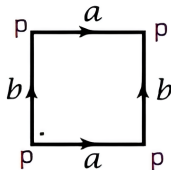
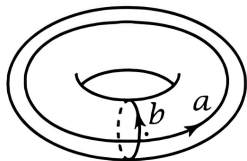


Figure: The torus \mathbb{T}^2 and its presentation as a cell complex (Hatcher)

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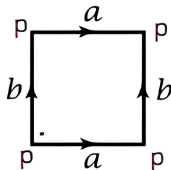
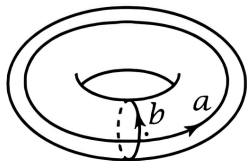


Figure: The torus \mathbb{T}^2 and its presentation as a cell complex (Hatcher)

This leads to the following HIT definition of \mathbb{T}^2 :

- $p : \mathbb{T}^2$
- $a : p = p$
- $b : p = p$
- $\text{fill} : a \cdot b = b \cdot a$

Some other HITs (3) : The Torus \mathbb{T}^2

In cubical type theory (and cubical) agda, we fill squares !

Some other HITs (3) : The Torus \mathbb{T}^2

In cubical type theory (and cubical) agda, we fill squares ! This gives another definition of the Torus with an already built-in "filled square Type".

```
data Torus2 : Type where
p  : Torus2
a  :  $p \equiv p$ 
b  :  $p \equiv p$ 
fill : Square b b a a
```

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The Hypercubic Manifold $\mathbb{H}\mathbb{M}^1$

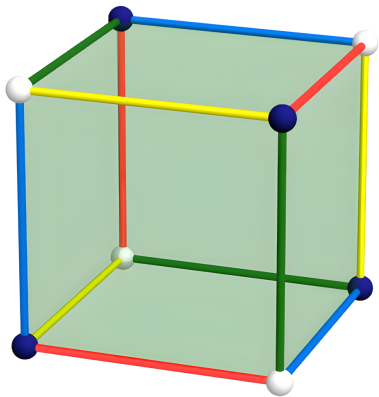


Figure: The hypercubic manifold $\mathbb{H}\mathbb{M}^1$ (analysis-situs)

The HIT/CW-complex correspondence

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To begin with, we will name the vertex constructors b^V and w^V (for blue and white vertex) and the edge constructors b^E, r^E, g^E, γ^E .

Defining \mathbf{HM}^1 (1)

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Challenges

- Identifying opposite sides under a quarter of a turn rotation
- Filling the cube

The Hypercubic Manifold $\mathbb{H}\mathbb{M}^1$

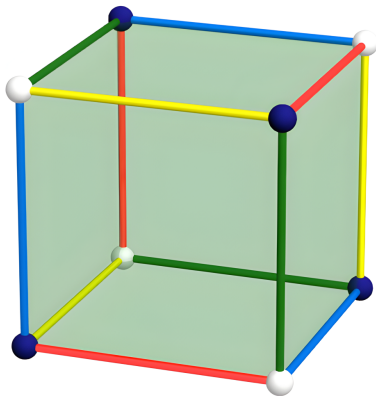
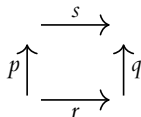


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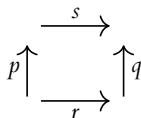
Defining $\mathbb{H}\mathbb{M}^1$ (2) : Rotation

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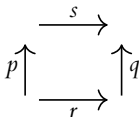
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We have just defined a map :

$\text{rot} : \text{Square } p \ q \ r \ s \rightarrow \text{Square } \bar{r} \ \bar{s} \ q \ p$

Defining $\text{HM}^1(3)$: Result

```
data Hypercubic : Type where
  ··· blueV : Hypercubic
  ··· whiteV : Hypercubic
  ··· yellowE : whiteV ≡ blueV
  ··· greenE : whiteV ≡ blueV
  ··· redE : whiteV ≡ blueV
  ··· blueE : whiteV ≡ blueV
  ··· f1 : Square (sym yellowE) greenE (sym blueE) redE
  ··· f3 : Square (sym yellowE) blueE (sym redE) greenE
  ··· f5 : Square (sym blueE) greenE (sym redE) yellowE
  ··· 3-cell : Cube f1 (rot f1) f3 (anti-rot f3) f5 (rot f5)
```

Figure: Synthetic description of the hypercubic manifold as a cubical HIT in cubical agda

Homotopy fibers

Homotopy fiber

Given a map $f : A \rightarrow B$ and $b : B$ the fiber of f over b is given by :

$$\mathrm{fib}_f(b) := \sum_{x:A} f(x) =_B b$$

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The final goal of the internship was to compute such a fiber, it has not yet been computed but we will look at an instructive example instead.

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Let now m be a nonnegative integer and $r \in \mathbb{Z}_m^\times$ (the integers mod m).
Then we have a short exact sequence of groups :

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Which, for topological reasons, yields a fibration of Eilenberg Mac-Lane spaces:

$$\mathbb{S}^1 \xrightarrow{\text{base} \mapsto \text{base}, \quad \text{loop} \mapsto \text{loop}^m} \mathbb{S}^1 \xrightarrow{g} B\mathbb{Z}_m$$

This is in fact a fibration whose base space is $B\mathbb{Z}_m$, of total space \mathbb{S}^1 , and for which the fiber of g over the canonical element $*$ of $B\mathbb{Z}_m$ should be \mathbb{S}^1 .

A fiber calculation (2)

By definition, the fundamental group at $*$ in $B\mathbb{Z}_m$ is \mathbb{Z}_m , we choose a presentation of it as the powers of s_m the successor mod m . Then the map g is given by :

$$g(\text{base}) :\equiv *, \quad \text{ap}_g(\text{loop}) :\equiv \text{ua}(s_m^r)$$

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We can now begin !

A fiber calculation (3) : diagrams

If $P : \mathbf{B}\mathbb{Z}_m \rightarrow \mathcal{U}$ is defined by $P := x \mapsto x = s_m$ then we have a fibration over \mathbb{S}^1 given by $P \circ g$ whose total space is by definition:

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Recalling that \mathbb{S}^1 is obtained as a coequalizer type, we obtain the following commutative diagram :

$$\begin{array}{ccccccc} 1 & \rightrightarrows & 1 & \xrightarrow{\text{base}} & \mathbb{S}^1 & \xrightarrow{g} & \mathbf{BZ}_m \\ & & \searrow & & \downarrow & \swarrow & \\ & & s_m = s_m & & P \circ g & & P \\ & & & & \downarrow & & \\ & & & & \mathcal{U} & & \end{array}$$

A fiber calculation (4) : The flattening lemma

Flattening Lemma

Suppose given a fibration over a coequalizer type:

$$B \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} A \xrightarrow{c} W \xrightarrow{P} \mathcal{U}$$

Then, one has a coequalizer diagram between the total spaces:

$$\sum_{b:B} P \circ c \circ f(a) \begin{array}{c} \xrightarrow{(b,x) \mapsto (g(b), p_b^*(x))} \\ \xrightarrow{(b,x) \mapsto (f(b), x)} \end{array} \sum_{a:A} P \circ c(a) \longrightarrow \sum_{w:W} P(w)$$

A fiber calculation (5)

In our case, the lemma yields the following coequalizer diagram :

$$s_m = s_m \begin{array}{c} \xrightarrow{\text{Id}} \\ \xrightarrow[p \mapsto p^r]{} \end{array} s_m = s_m \longrightarrow \text{fib}_g(s_m)$$

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We get that $\text{fib}_g(s_m)$ is equivalent to the HIT W defined by the following constructors :

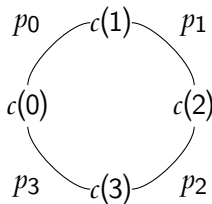
- $c : \text{Fin } m \mapsto W$
- $p : \prod_{x:\text{Fin } m} (c(x) = c(s_m^r(x)))$

A fiber calculation (6)

Since r is coprime to m , it turns out that s_m^r is of order m so the orbit of $c(0)$ under the action of the group $\langle s_m^r \rangle$ is the whole type W .

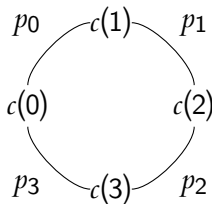
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Hence :

$$\boxed{\mathrm{fib}_g(s_m) \simeq \mathbb{S}^1}$$

A fiber calculation (7) : sketch of proof

The idea is that the general case should come from the case $r = 1$.

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For $r = 1$, we have that :

$$p_0 \cdot p_1 \cdot \dots \cdot p_m : c(0) = c(0)$$

and we can successively contract each path :

$$p_i : c(i) \rightarrow c(\overline{i+1}_r)$$

to end up with a type with a constructor $c(0)$ and a path $p_0 : c(0) = c(0)$,
which is the circle.

A fiber calculation (8) : An extension

Let's consider the case where r and m are not coprime.

In that case, let δ be the gcd of r and m , then s_m^r is now only of order $\frac{m}{\delta}$.

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We then naturally obtain :

$$\mathrm{fib}_g(s_m) = \bigoplus_{i=1}^{\delta} \mathbb{S}^1$$

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Towards $\mathbb{H}\mathbb{M}^1$ as $\mathbb{S}^3/Q \simeq \mathbb{S}^3$

The final bit of the work would consist in an other fiber calculation that expresses $\mathbb{H}\mathbb{M}^1$ as a homotopy quotient of \mathbb{S}^3 under an action of Q .

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