# The Hypercubic Manifold in Homotopy Type Theory

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  - A fiber calculation
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## Intro (1): HoTT

## Homotopy Type Theory

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## Homotopy Type Theory

- Vladimir Voevodsky: proposed the "Univalence Axiom"
- Synthetic homotopy
- Computer assisted proofs
- "Univalent Foundations" as a foundational system for maths

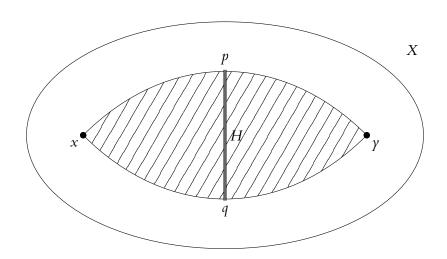
Let X be a topological space and I the unit real interval [0,1]

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- Homotopy: describes the correct notion of "equality" for paths.
- Two paths p and q are "homotopic" if there is "a continuous deformation" of p onto q
- Loops: a loop at point x is a path p such that p(0) = p(1) = x.



#### Paths operations

- inversion : given  $p: x \to y$  there is a path  $p^{-1}: y \to x$
- concatenation : given  $p: x \to y$  and  $q: y \to z$  there is a path  $p \cdot q: x \to z$

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## Fundamental group $\pi_1(X, x)$

Take a point x of X and consider loops at point x.

- Equality: homotopy (of loops)
- Composition: path concatenation (associative)
- Inverse element : inverse path
- Unit element : constant path at *x*

Two (path connected) homeormorphical spaces (more precisely, homotopically equivalent) have the same fundametal groups!

## Higher homotopy groups

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- Is there a (path-connected) space such that  $\pi_1(X) = G$  and  $\pi_n(X) = 1$  for any  $n \ge 2$ ?

## Higher homotopy groups

- Homotopy are themselves paths between paths
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- Is there a (path-connected) space such that  $\pi_1(X) = G$  and  $\pi_n(X) = 1$ for any  $n \ge 2$ ? Answer: yes, the **Eilenberg Mac-Lane** space of G

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- suc :  $\mathbb{N} \to \mathbb{N}$

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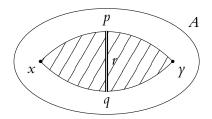
## Propositional Equality

Given a type A and terms a, b: A there is a type  $a =_A b$ . We say that elements a and b are (propositionally) equal when there exists some element:

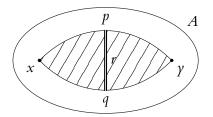
$$p: a =_A b$$

Types	Topology
A	a space A
a : A	a is a point of A
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## Synthetic homotopy

This connection is at the core of synthetic homotopy theory which allows us to define every object we previously talked about in the synthetic framework of type theory.

# Intro (8): The Hypercubic Manifold HM<sup>1</sup>

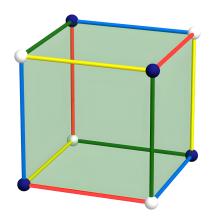


Figure: The hypercubic manifold HM<sup>1</sup> (analysis-situs)

#### Goals

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#### Goals

- We know how to define Eilenberg Mac-Lane spaces for finite groups in HoTT but the construction is not well understood
- We would like to achieve a "nicer" construction that would yield a handy induction principle
- This would open up to define "Higher Group" actions and computer checked cohomology calculations.
- It turns out that  $\pi_1(\mathbb{HM}^1) = Q$  and that by computing the "homotopy fiber" of a certain map my advisors have a way to provide such a "nice" construction in the case of the group Q.

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## Some constructions on types

#### Function types

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## Dependant pair Types

If *B* is a type family over a type *A* (that is a function  $B: A \to \mathcal{U}$ ) we have a type of dependant pairs  $\sum_{x:A} B(x)$ .

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## Coproduct types

Give types A, B one has a coproduct type A + B given by the constructors :

- inl:  $A \rightarrow A + B$
- inr :  $B \rightarrow A + B$

# Manipulating types

#### Introduction rules

They encapsulate how to build an element of a certain type.

- To build an element of  $f: A \to B$  one needs an expression  $\phi(x)$  such that  $a: A \vdash \phi(a): B$  and to set  $f: \equiv \lambda x.\phi(x)$
- To build an element of  $A \times B$  one needs elements a : A and b : B to form  $(a, b) : A \times B$

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- To build an element of A × B one needs elements a : A and b : B to form (a, b) : A × B

#### Induction principles

They encapsulate how to build **dependent** functions from a source type *A*.

• To build a function of type  $f: \prod_{x:A\times B} P(x)$  one only needs to give its value on pairs (a, b).

# A working example: product types (1)

### Projections

We can define projections  $\operatorname{pr}_1: A \times B \to A$  and  $\operatorname{pr}_2: A \times B \to B$  by the **induction** principle for product types by setting  $\operatorname{pr}_1((a,b)) :\equiv a$  and  $\operatorname{pr}_2((a,b)) :\equiv b$ .

#### Definition

Given a type A and elements a, b : A there exists a type  $a =_A b$  called the identity type of a and b. We say that "a is equal to b" if there exists some element  $p : a =_A b$ 

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### Reflexivity

For any type A and point x : A there is an element refl<sub>x</sub> :  $x =_A x$ 

#### Path induction

To prove a predicate that depends on x, y: A and p:  $x =_A y$  one only needs to prove it in the case where  $x \equiv y$  and  $p \equiv \text{refl.}$ 

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## Example (Functions preserve equals)

For any  $f: A \rightarrow B$ , x, y: A and  $p: a =_A b$  there is an element

$$f(p): f(a) =_B f(b)$$

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- We then have to provide an element of  $f(x) =_B f(x)$
- We conclude by setting  $f(\operatorname{refl}_x) := \operatorname{refl}_{f(x)}$



# HoTT: The homotopical interpretation of Type Theory (1)

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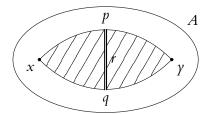


Figure: two points x, y : A, two paths  $p, q : x =_A y$  and a homotopy  $r : p =_{x =_A y} q$ 

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- **Symmetry**: Given a path  $p: x =_A y$ , one has an inverse path  $p^{-1}: y =_A x$  such that for any x,  $\operatorname{refl}_x^{-1} \equiv \operatorname{refl}_x$ .
- Transitivity: Given paths  $p: x =_A y$  and  $q: y =_A z$  one has a concatenated path  $p \cdot q: x =_A z$  such that for any x,  $\operatorname{refl}_x \cdot \operatorname{refl}_x \equiv \operatorname{refl}_x$ .

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There are additional properties that one can prove such as

$$p \cdot p^{-1} = \text{refl}$$

that justify this whole interpretation.



## Transports

## Equalities and dependent functions

Take  $f: \prod_{x:A} P(x), x, y:A$  and p:x=y.

Note that we can't have a path between f(x) : P(x) and f(y) : P(y).

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#### Transport

By path induction, for any x, y and p : x = y one can define :

transport<sup>P</sup>: 
$$\prod_{a,b:A} \prod_{p:a=A} P(a) \rightarrow P(b)$$

We may write transport P(a, b, p) as  $p_*$  if the context is clear.

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## Example

The map  $p_*$  encapsulates how are transported elements of the "fiber" P(x) to the fiber P(y) by going along the path p: x = y.

# Fibrations (1)

#### **Fibration**

Take a continuous map  $p: E \to B$ . It is called a fibration if it has a certain path lifting property from the **base space** B to the **total space** E.

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#### Fibrations in HoTT

Take a type family B over A (that is  $B: A \to \mathcal{U}$ ).

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More precisely:

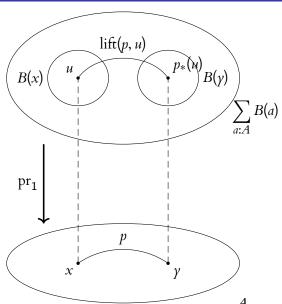
Given x : A and u : P(x), for any y : A and p : x = y, one has a path

$$lift(p, u) : (x, u) = (\gamma, p_*(u))$$

such that  $pr_1(lift(p, u)) = p$ .



# Fibrations (2)



# Fibrations (3)

## Paths over paths

Now, given  $f: \prod_{x:A} B(x)$  and p: x = y there is a natural identification between points f(x) and  $p_*(f(x))$ .

# Fibrations (3)

## Paths over paths

Now, given  $f: \prod_{x:A} B(x)$  and p: x = y there is a natural identification between points f(x) and  $p_*(f(x))$ .

A path lying over p should then factor through this path and consist of a path between  $p_*(f(x))$  and a certain point u: P(y).

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Dependent functions yield such paths:

$$\operatorname{apd}_f: \prod_{p:x=A} p_*(f(x)) = f(y)$$

## Equivalences and Univalence

#### Equivalences

We can build them from quasi-equivalences. Two types are quasi-equivalent if there exists maps  $f:A\to B$  and  $g:B\to A$  such that :

$$\prod_{x:A} g(f(x)) = x \quad \text{and} \quad \prod_{y:B} f(g(y)) = y$$

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#### <u>Univalence</u>

One has:

$$(A \simeq B) \simeq (A = B)$$



# Higher Inductive Types

#### Idea

When dealing with HITs we allow ourselves to use paths constructors, paths between paths constructors and so on.

# Higher Inductive Types

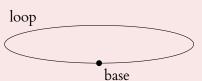
#### Idea

When dealing with HITs we allow ourselves to use paths constructors, paths between paths constructors and so on.

## The circle $\mathbb{S}^1$

The circle  $\mathbb{S}^1$  is defined by the following HIT:

- base :  $\mathbb{S}^1$
- loop: base =<sub>S1</sub> base



## HITs: Induction Principle

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To define a map  $f: \prod_{x:\mathbb{S}^1} P(x)$ , one needs to specify its value f (base): P(base). Then, what should you send loop to?

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#### The case of $\mathbb{S}^1$

To define a map  $f: \prod_{x:\mathbb{S}^1} P(x)$ , one needs to specify its value f (base) : P(base). Then, what should you send loop to ? Recall that by transport :

$$P(\text{base}) \stackrel{\text{loop}_*}{\simeq} P(\text{base})$$

Since f is a dependent function, it should send loop to a dependent path between loop, (f(base)) and f(base).

## Some other HITs (1)

#### Coequalizers

Given types B, A and maps f,  $g: B \to A$  the (homotopy) **coequalizer** type CoEq(f, g) is given by the HIT:

- $c: A \to CoEq(f, g)$
- $\bullet \ p: \prod_{b:B} c(f(b)) = c(g(b))$

$$B \xrightarrow{g \atop f} A \xrightarrow{c} \operatorname{CoEq}(f,g)$$

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#### The circle as a coequalizer

$$1 \xrightarrow{1 \atop 1} 1 \xrightarrow{c} CoEq(1,1) \simeq \mathbb{S}^1$$



## Cubical Type Theory: Completing Squares

#### Filling property

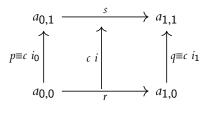
Given an incomplete square of paths

$$\begin{array}{ccc}
 a & \xrightarrow{p} & b \\
 q \downarrow & & \downarrow r \\
 c & & d
\end{array}$$

Then there is a unique path s from c to d that completes the square by making it commute, that is:

$$q \cdot s = p \cdot r$$

## Why do we complete squares?





## Some other HITs (2): The torus $\mathbb{T}^2$

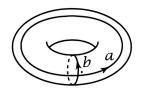




Figure: The torus  $\mathbb{T}^2$  and its presentation as a cell complex (Hatcher)

## Some other HITs (2): The torus $\mathbb{T}^2$

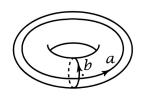




Figure: The torus  $\mathbb{T}^2$  and its presentation as a cell complex (Hatcher)

This leads to the following HIT definition of  $\mathbb{T}^2$ :

- $p: \mathbb{T}^2$
- $\bullet$  a:p=p
- $\bullet$  b: p = p
- fill :  $a \cdot b = b \cdot a$

## Some other HITs (3): The Torus $\mathbb{T}^2$

In cubical type theory (and cubical) agda, we fill squares!

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In cubical type theory (and cubical) agda, we fill squares! This gives another definition of the Torus with an already built-in "filled square Type".

```
data Torus2 : Type where

p : Torus2

a : p \equiv p

b : p \equiv p

fill : Square b b a a
```

#### Fundamental group presentation

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Given a presentation of G we can build a space H with  $\pi_1(H) = G$ .

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# The Hypercubic Manifold HM<sup>1</sup>

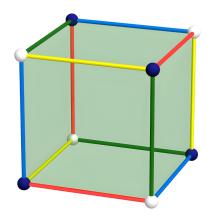


Figure: The hypercubic manifold HM<sup>1</sup> (analysis-situs)

## Defining $\mathbb{HM}^1$ (1)

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To begin with, we will name the vertex constructors  $b^V$  and  $w^V$  (for blue and white vertex) and the edge constructors  $b^E$ ,  $r^E$ ,  $g^E$ ,  $\gamma^E$ .

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#### Challenges

- Identifying opposite sides under a quater of a turn rotation
- Filling the cube

# The Hypercubic Manifold HM<sup>1</sup>

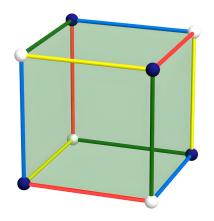


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$$r^{-1}$$
  $\xrightarrow{p}$   $\downarrow$   $s^{-1}$ 

We have just defined a map:

rot : Square  $p \neq r s \rightarrow \text{Square } \overline{r} \ \overline{s} \neq p$ 

# Defining HM<sup>1</sup> (3): Result

Figure: Synthetic description of the hypercubic manifold as a cubical HIT in cubical agda

### Homotopy fibers

#### Homotopy fiber

Given a map  $f:A\to B$  and b:B the fiber of f over b is given by :

$$\operatorname{fib}_{f}(b) := \sum_{x:A} f(x) =_{B} b$$

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The final goal of the internship was to compute such a fiber, it has not yet been computed but we will look at an instructive example instead.

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Let now m be a nonnegative integer and  $r \in \mathbb{Z}_m^{\times}$  (the integers mod m). Then we have a short exact sequence of groups :

$$\mathbb{Z} \xrightarrow{\times m} \mathbb{Z} \xrightarrow{1 \mapsto \bar{r}} \mathbb{Z}_m$$

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Which, for topological reasons, yields a fibration of Eilenberg Mac-Lane spaces:

$$\mathbb{S}^1 \xrightarrow{\text{base} \mapsto \text{base}, \quad \text{loop} \mapsto \text{loop}^m} \mathbb{S}^1 \xrightarrow{g} \mathbf{B}\mathbb{Z}_m$$

This is in fact a fibration whose base space is  $B\mathbb{Z}_m$ , of total space  $\mathbb{S}^1$ , and for which the fiber of g over the canonical element \* of  $B\mathbb{Z}_m$  should be  $\mathbb{S}^1$ .

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By definition, the fundamental group at \* in  $\mathbb{B}\mathbb{Z}_m$  is  $\mathbb{Z}_m$ , we choose a presentation of it as the powers of  $s_m$  the successor mod m. Then the map g is given by :

$$g(\text{base}) := *, \quad \text{ap}_g(\text{loop}) := \text{ua}(s_m^r)$$

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$$g(base) := *, ap_g(loop) := ua(s_m^r)$$

We can now begin!

## A fiber calculation (3) : diagrams

If  $P : \mathbb{BZ}_m \to \mathcal{U}$  is defined by  $P :\equiv x \mapsto x = s_m$  then we have a fibration over  $\mathbb{S}^1$  given by  $P \circ g$  whose total space is by definition:

$$\sum_{x:\mathbb{S}^1} P \circ g(x) := \sum_{x:\mathbb{S}^1} (g(x) = s_m) := \mathrm{fib}_g(s_m)$$

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Recalling that  $\mathbb{S}^1$  is obtained as a coequalizer type, we obtain the following commutative diagram :

$$1 \longrightarrow 1 \xrightarrow{\text{base}} \mathbb{S}^1 \xrightarrow{g} \mathbb{B}\mathbb{Z}_m$$

$$s_m = s_m \qquad p \xrightarrow{p \circ g} \qquad P$$

$$\mathcal{U}$$

## A fiber calculation (4): The flattening lemma

#### Flattening Lemma

Suppose given a fibration over a coequalizer type:

$$B \xrightarrow{f} A \xrightarrow{c} W \xrightarrow{P} \mathcal{U}$$

Then, one has a coequalizer diagram between the total spaces:

$$\sum_{b:B} P \circ c \circ f(a) \xrightarrow{(b,x) \mapsto (g(b),p_b^*(x))} \sum_{a:A} P \circ c(a) \longrightarrow \sum_{w:W} P(w)$$

In our case, the lemma yields the following coequalizer diagram:

$$s_m = s_m \xrightarrow{\text{Id}} s_m = s_m \longrightarrow \text{fib}_g(s_m)$$

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We get that  $fib_g(s_m)$  is equivalent to the HIT W defined by the following constructors :

- $c : \operatorname{Fin} m \mapsto W$
- $p: \prod_{x: \text{Fin } m} (c(x) = c(s_m^r(x)))$

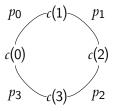


#### A fiber calculation (6)

Since r is coprime to m, it turns out that  $s_m^r$  is of order m so the orbit of c(0) under the action of the group  $\langle s_m^r \rangle$  is the whole type W.

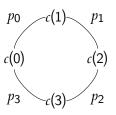
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Hence:

$$\operatorname{fib}_g(s_m)\simeq \mathbb{S}^1$$

### A fiber calculation (7): sketch of proof

The idea is that the general case should come from the case r = 1.

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The idea is that the general case should come from the case r = 1. For r = 1, we have that :

$$p_0 \cdot p_1 \cdot \ldots \cdot p_m : c(0) = c(0)$$

and we can successively contract each path:

$$p_i:c(i)\to c(\overline{i+1}_r)$$

to end up with a type with a constructor c(0) and a path  $p_0: c(0) = c(0)$ , which is the circle.

#### A fiber calculation (8): An extension

Let's consider the case where r and m are not coprime. In that case, let  $\delta$  be the gcd of r and m, then  $s_m^r$  is now only of order  $\frac{m}{\delta}$ .

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We then naturally obtain:

$$\operatorname{fib}_g(s_m) = \bigoplus_{i=1}^{\delta} \mathbb{S}^1$$

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# Towards $\mathbb{HM}^1$ as $\mathbb{S}^3/Q \curvearrowright \mathbb{S}^3$

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