

# *Introduction to Database Systems*

## *Chapter 1*

# What Is a DBMS?

비트 Bit(b)	바이트 Byte(B)	킬로바이트 Kilobyte(KB)	메가바이트 Megabyte(MB)
기가바이트 Gigabyte(GB)	테라바이트 Terabyte(TB)	<b>페타바이트 Petabyte(PB)</b>	엑사바이트 Exabyte(EB)
제타바이트 Zettabyte(ZB)	요타바이트 Yottabyte(YB)	론나바이트 Ronnabyte(RB)	퀘타바이트 Quettabyte(QB)

1 KByte =  $10^3$  byte

- v Data vs Information vs Knowledge
  - Data: raw facts that are described, observed, or measured
  - Information: data that have been prepared or organized
  - Knowledge: data/information/rules that are used for actual decision making 패턴
- v A (large, integrated) collection of (related) data. ⇒ DB
- v Models real-world enterprise.
  - Entities (e.g., students, courses)
  - Relationships (e.g., Madonna is taking CS564)
- v A Database Management System (DBMS) is a software package designed to store and manage databases.

cf) M4 Access

강제!



# Why Use a DBMS?



- v <sup>\*</sup> **Data independence** and efficient access.
- v Reduced application development time.
- v **Data integrity** and **security**.
- v <sup>무결성</sup> **Uniform data** administration. <sup>관리</sup>
- v Concurrent access, recovery from crashes. <sup>동시 액세스</sup>

(f) 모델링: 현실세계를 묘사하기 위한 가상의 것을 만드는 것

# Data Models

schema와 instance 구분하기!

ex) schema: DB의 구조 설명

ex) DB 'A' 는 relational data model이고 'SID' row가 있다

instance: DB의 스냅샷

ex) 9/12 일자로 보니 2016~ 학생데이터가 왔다

- v A data model is a collection of concepts for describing data. → DB를 만들기 위함 (DB: collection of data)
- v A schema is a description of a particular collection of data, using the given data model.  
DB의 구조를 묘사 (늘, 고침, ...)
- v The relational model of data is the most widely used model today. ⇒ chap2.3에서 계속
  - Main concept: relation, basically a table with rows and columns.
  - Every relation has a schema, which describes the columns, or fields.

# History of Data Models

백만칩의 구조비자: 금융업계, 군부  
(오라클 → 은행) 데이터관리 중요  
OLTP: Online Transaction Process  
↓  
DB

비행기 여행회사

## v Early 60's : Network Model

- By Charles Bachman at GE
- CODASYL (Conference On Data Systems Language)

## v Late 60's : Hierarchical Model

- IMS by IBM (IBM의 DBMS) → 대중적, 선두주자
- SABRE airline reservation system

자료구조에서의 트리 / 그래프 (네트워크)

node & edge (vertex & link)  
edge에 weight를 줌

graph tree graph: cycle 有

## v 70's: Relational Model

비대칭 relation (table)으로 바뀜

By Edgar Codd at IBM BCNF 카드

SQL from IBM System R : prototype model & language (시작)

file: collection of record (data)

## Late 80's & early 90's: Extended(Object) Relational Model or Object-Oriented Model → 객체지향

Late 90's: Data Warehousing, Data Mining, Distributed DB, Web Database, ERP(Enterprise Resource Planning), 전사적자원관리 (SAP, SCM) → 구글

How about ER(Entity-Relationship Model) → 학부생 수준

Semantic Data Model vs Record-based Model

의미적 개념적

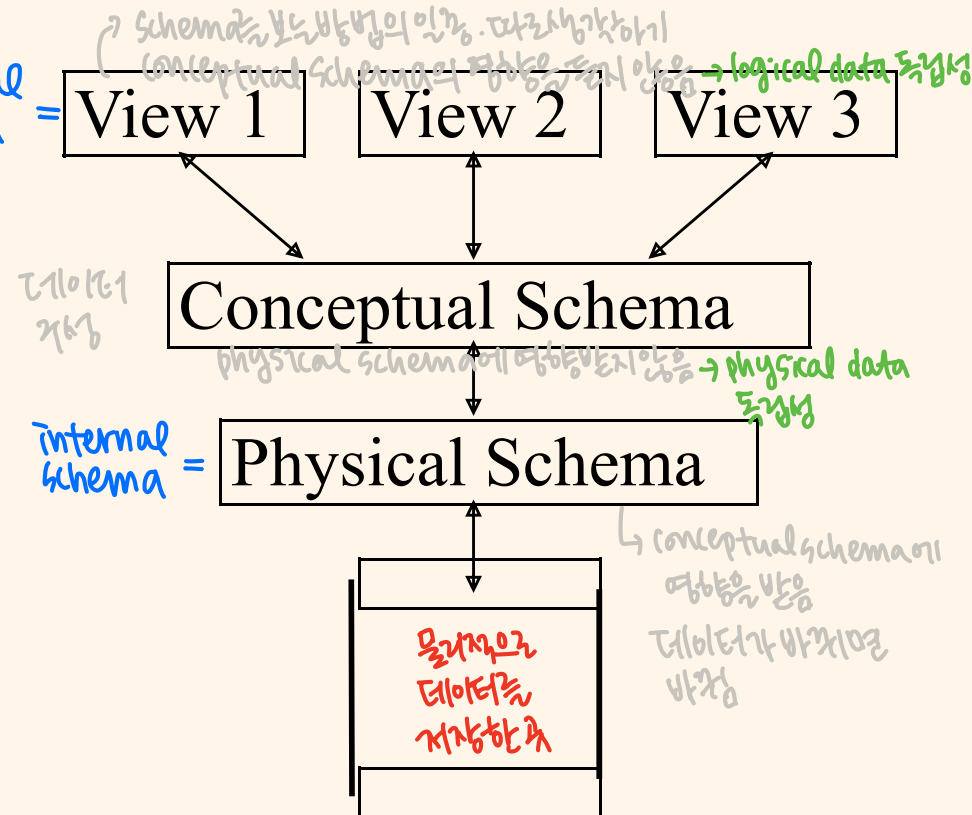
물리적(저장)

# Levels of Abstraction

→ data independence를 위해

v Many views, single conceptual (logical) schema and physical schema.

- Views describe how users see the data.
- Conceptual schema defines a entire logical structure of the data
- Physical schema describes the files and indexes used.



\* Schemas are defined using DDL; data is modified/queried using DML.

conceptual schema는  
DDL language

instance는 DML

# Example: University Database

## v Conceptual schema:

- *Students*(*sid*: string, *name*: string, *login*: string, *age*: integer, *gpa*: real)
- *Courses*(*cid*: string, *cname*: string, *credits*: integer)
- *Enrolled*(*sid*: string, *cid*: string, *grade*: string)

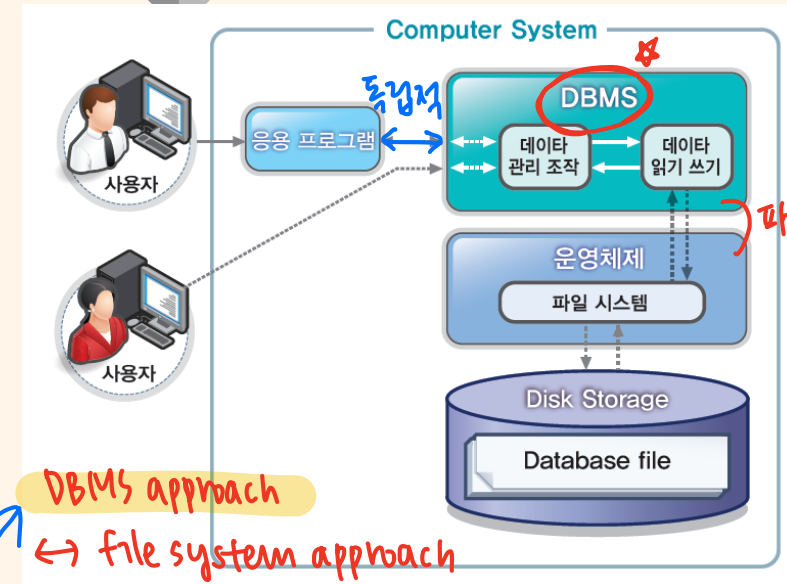
## v Physical schema: (internal schema)

- Relations stored as unordered files.
  - Index on first column of *Students*.
- 만약치면 오버헤드↑  
변경의 structure →
- 파일구조

## v External Schema (View):

- *Course\_info*(*cid*: string, *enrollment*: integer)

↳ 필요한 일부 column만 추출  
(schema를 보는 방법의 일종)



DBMS approach

↔ file system approach

(각접접근, DBMS는 거치지 않음)

schema 바뀌어도  
다들 알고차지 않아도 됨

```
typedef struct
{
    char sid[10];
    char name[300];
    char login[20];
    int age;
    float gpa;
    char address[300];
} students;
```

```
int main()
{
    int counter;
    FILE *ptr_myfile;
    students my_record;
```

```
    ptr_myfile=fopen("c:\MyDB\test.db","rb");
    if (!ptr_myfile)
    {
        printf("Unable to open file!");
        return 1;
    }
```

```
    for ( counter=1; counter <= 10; counter++)
    {
        fread(&my_record,sizeof(struct rec),1,ptr_myfile);
        printf("%d\n",my_record.x);
    }
    fclose(ptr_myfile);
    return 0;
}
```



가장 먼저 구성한 이유!

# Data Independence

- v Applications insulated from how data is structured and stored.
- v Logical data independence: Protection from changes in *logical* structure of data. → conceptual schema가 바뀌면 view에 영향X
- v Physical data independence: Protection from changes in *physical* structure of data.  
→ physical schema가 바뀌면 logical schema(+view)에 영향X  
예) 인덱스 안다나느새 새로 만들면 → physical schema의 변화

\* One of the most important benefits of using a DBMS!

transaction : 거래  
cf) ATM - deposit, withdraw

DB의 처리단위 (operation)

# Transaction: An Execution of a DB Program

- v Key concept is transaction, which is an atomic sequence of database actions (reads/writes).  
→ operation  
SELECT INSERT, DELETE, UPDATE  
↗ 각각 독립적인 logical unit
- v Each transaction, executed completely, must leave the DB in a consistent state if DB is consistent when the transaction begins.  
transaction 실행 중일 때의 의미가 아님!  
시작과 끝에 consistent하게 유지하면 됨

DBMS가 도맡는다  
commit, rollback  
→ consistent state를 유지  
DB

오스 내용!

→ DBMS가 알아서

# Concurrency Control

동시성 제어

throughput을 높이기 위해 동시작업을 처리하면 동시성 제어가 필요

- v Concurrent execution of user programs is essential for good DBMS performance.
  - Because disk accesses are frequent, and relatively slow, it is important to keep the cpu humming by working on several user programs concurrently.
- (멀티프로세싱)  
v 교차실행 병렬처리 x → Interleaving actions of different user programs can lead to inconsistency: e.g., check is cleared while account balance is being computed. 원시제어 관리 문제
- v DBMS ensures such problems don't arise: users can pretend they are using a single-user system.
- v *Scheduling, Locking, , , , , , ,* +) semaphore

transaction fail 난 경우? (transaction 종료에) ⇒ 없던 것으로 처리해야 함!

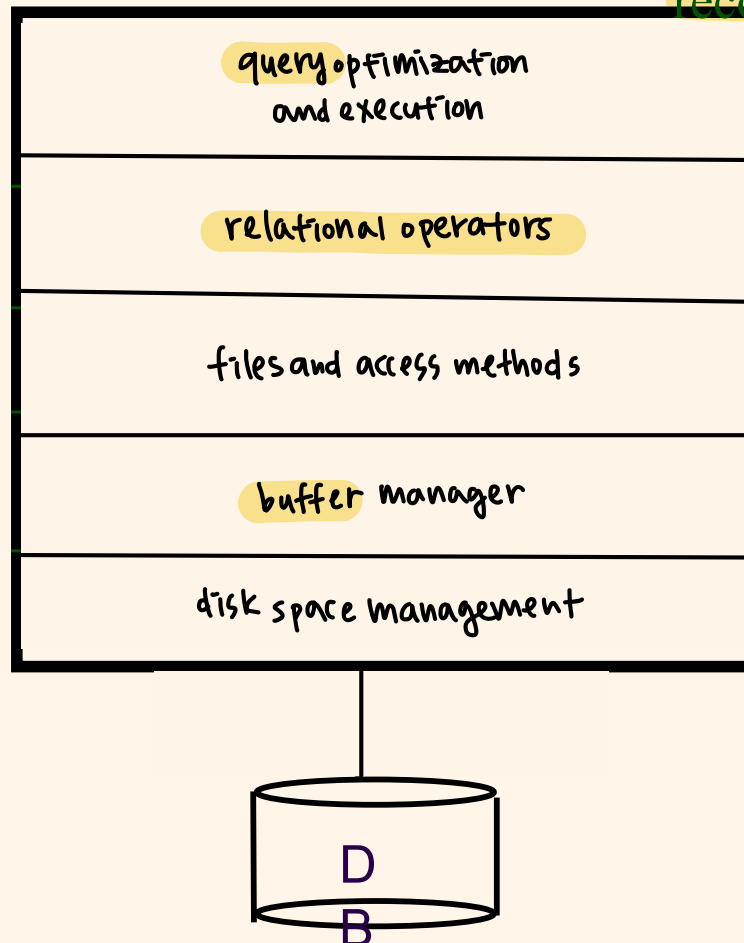
read/write 후 commit을 call해야 하나의 transaction이 끝남

로그에 commit이 없는 경우 transaction fail → recovery 작업 필요

# Structure of a DBMS

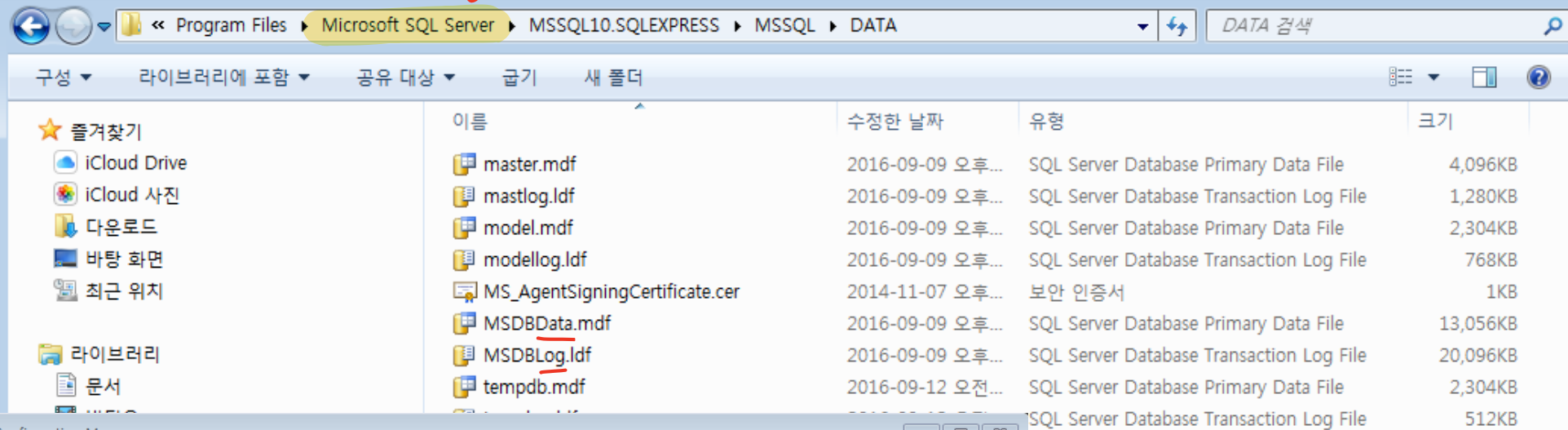
하나의 예시!

- ✓ A typical DBMS has a layered architecture.
- ✓ The figure does not show the concurrency control and recovery components.
- ✓ This is one of several possible architectures; each system has its own variations.

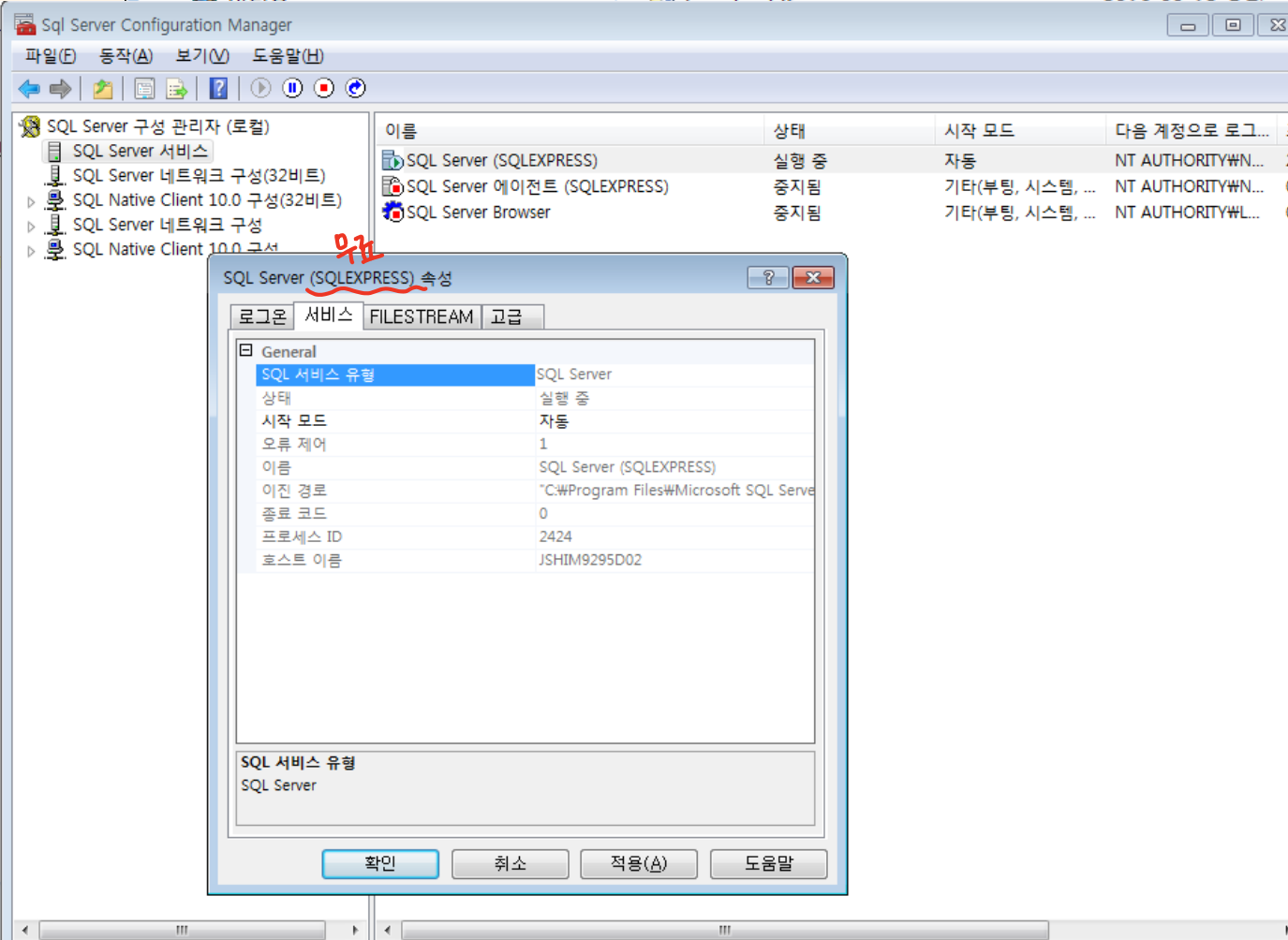


These layers must consider concurrency control and recovery

orade의가이벌!



↳ 디스크에서 file를 저장.관리됨





## Summary

- v DBMS used to maintain, query large datasets.
- v Benefits include recovery from system crashes, concurrent access, quick application development, data integrity and security.
- v Levels of abstraction give data independence. → schema
- v A DBMS typically has a layered architecture. + component 기반
- v DBAs hold responsible jobs and are well-paid!
- v DBMS R&D is one of the broadest, most exciting areas in CS.

