## **Data Structures**

3. Recursion

### 3. Recursion

- 1. Recursion
- 2. Factorial
- 3. Binary Search
- 4. Permutation

#### Recursion

- A function is called by itself while running
- possible when the same task is repeated with **exit** condition
- "Divide & Conquer" strategy
- Iterative function with if-else or while can be rewritten in recursion
- In compiler's view, recursive call is the same as another new function is called
- Advantages
  - Clear representation of algorithm
- Disadvantage
  - Frequent function call overhead

### **Factorial**

```
• n! = 1 where n = 0
• n! = n * (n - 1)! where n \ge 1
```

```
h'_{1} = 1 * 2* ... * (n-1) * n
(n-1)!
f(4) = 4* f(3)
= 3* f(2)
```

Program 3.1

```
long factorial (int n)
{
   if (n == 0)
      return 1;
   else
      return _factorial (n-1)
   ;
}
```

# GCD

Greatest Common Divisor

- gcd(x, y) = gcd(y, x % y) if y > 0
- gcd(x, 0) = x
- ex) gcd(48, 8) = gcd(8, 48.1.8) = gcd(8, 0) = 8
- Program 3.2

### Binary Search

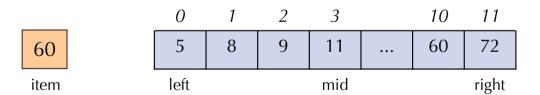
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- Problem) search a specified key in a \_\_\_\_\_\_\_ list
  - S = [3, 7, 8, 11, 13, 19, 25]
  - No overlapping number
- Algorithm

```
while (there are any item to compare)
{
    find Median position;
    if (item < median)
        change the range with left sub-list
    else if (item > median)
        change the range with right sub-list
    else
        return median position;
}
```

# Binary Search

- Program 3.3 194
  - Iterative Binary Search



## Recursion - Binary Search

- Program 3.4 "মানাস ামুহচার" P5T
  - Recursive Binary Search

### Hanoi Tower

- Goal of Problem
  - move some discs in a pole to another pole without changing the order of discs
- Condition
  - 1. Only single disc can be moved at a time
  - 2. A disc cannot be placed over smaller one



#### Hanoi Tower

- Algorithm
  - Assuming that there are three poles A, B, C, and all discs of pole A have to be moved to pole B
  - 1. Move n 1 discs from pole A to pole \_\_\_ -> চপা কুম্বান মুধ্যান মু
  - 2. Move single disc from pole A to pole B
  - 3. Move n 1 discs from C to pole \_\_\_\_\_\_\_\_
    - Step 1 and 3 should be solved recursively
- Program 3.5

### Hanoi Tower P57

- # of movements
- time to move single disk : t(1)t(1) = 1
- time to move n disks : t(n)

$$t(n) = 2 t(n - 1) + 1$$

$$= 2 (2 t(n - 2) + 1) + 1$$

$$= 2 (2 (2 t(n - 3) + 1) + 1) + 1 = 2^{3} * t(n - 3) + 2^{2} + 2^{1} + 2^{0}$$

$$= ....$$

$$= 2^{n-1} + 2^{n-2} + ... + 2^{1} + 2^{0}$$

$$= \frac{2^{n-1}}{2^{n-2}}$$

$$t(n) = O(2^{n}) // Time complexity / 7 + 2^{n} + 2^{n}$$

### Permutation (5%)

- Find all possible permutations with n elements (n > 1)
- ex) For a set {a, b, c}, find all permutations
  - $P = \{(a, b, c), (a, c, b), (b, a, c), (b, c, a), (c, a, b), (c, b, a)\}$
  - $_{3}P_{3} = 3 \times 2 \times 1 = 3! = 6$

#### **Permutation**

- Permutations of {a, b, c, d}
  - 24 permutations are divided into 4 groups
    - a)  $(a, b, c, d), (a, b, d, c), ... => (a, Perm(b, c, d)) \rightarrow 6$ b) (b, Perm(a, c, d))c) (c, Perm(a, b, d))d) (d, Perm(a, b, c))
  - *Perm(b, c, d)* can be solved in the same way
    - (b , Perm(c, d)), (c , Perm(b, d)), (d , Perm (b, c))
  - => a clue for recursive solution

#### **Permutation**

```
• char list[] = { 'a', 'b', 'c', 'd' };
• i : start index of range
• n : end index of range
• ex)
   • perm (list, 0, 3)
    • perm (list, 0, 2)
                              PCI,37
```

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```
void perm(char list[], int i, int n)
  int j;
   if (i == n)
     for (j = 0; j \le n; j++)
        printf ("%c", list[j]);
     printf ("
   else
     for (j = i; j \le n; j++)
       Swap (list[i], list[j]);
      perm (list, it , n); y azulete
         Swap (list[i], list[j]);
            6 9 1 मित्रामिता
```

### Permutation

Perm(0, 3)	Perm(1, 3)	Perm(2, 3)	Perm(3, 3)
a(bcd)	/ a (b (c d))	a(b(c(d))) —	a, b, c, d
		$\sim a(b(d(c))) -$	— a, b, d, c
	a(c(b d))	a(c(b(d)))	— a, c, b, d
		a(c(d(b))) —	— a, c, d, b
	a(d(c b))	-a(d(c(b)))	— a, d, c, b
	\	a(d(b(c))) —	— a, d, b, c
b (a c d)	$\int b(a(c d)) <$	_ b(a(c(d)))	— b, a, c, d
	$\int b(c(ad))$	b(a(d(c)))	— b, a, d, c
	$\sqrt{b(d(c a))}$	:	
c (b a d)	c ( b ( a d ) )	c ( b ( a ( d ) ) )	c, b, a, d
	•••••	•••••	
d (b c a)	d(b(ca))	d(b(c(a)))	d, b, c, a
	•••••	•••••	

• Exercise 3.3, 3.4

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