

# Inter-process Communication using Pipes

#### **Pipes**



- A pipe is a *one-way* communication channel that couples one process to another.
- Used only between processes that have a common ancestor.
  - More specifically, for communication between parent and child in general.
- There are two types of pipes: named and unnamed.
  - Unnamed pipes are used for communication between related processes.
  - Named pipes can be used for communication between unrelated processes.

#### Unnamed Pipes (1)



- int pipe(int fd[2]);
  - To create an unnamed pipe.
  - Needs an array of two int's for two file descriptors (r/w).
- A pipe is another generalization of the UNIX file concept.
  - A pipe is a FIFO file: lseek() does not work on a pipe.
  - Unnamed pipes come and go.
  - Named pipes (also known as FIFOs) are permanent files.
- The size of a pipe is limited.
  - Max is PIPE BUF

#### Unnamed Pipes (2)



- If a read() is issued while the pipe is empty, it will block.
- If a write() is issued when the pipe is full, it will block.
- If all writers of a pipe are closed, a reader will encounter EOF.
- If all readers of a pipe are closed, a writer will face a broken pipe.
- Non-blocking reads and writes:
  - Issue fcntl() with O\_NONBLOCK flag.

#### Example #1: Unnamed Pipe (1)



```
#include <stdio.h>
#define MSGSIZE 16
static char *msg1 = "hello, world #1";
static char *msg2 = "hello, world #2";
static char *msg3 = "hello, world #3";
int main(void)
   char buf[MSGSIZE];
   int fd[2], i, pid;
   /* open unnamed pipe */
   if (pipe(fd) < 0) {
      perror("pipe call");
      exit(1);
```

#### Example #1: Unnamed Pipe (2)

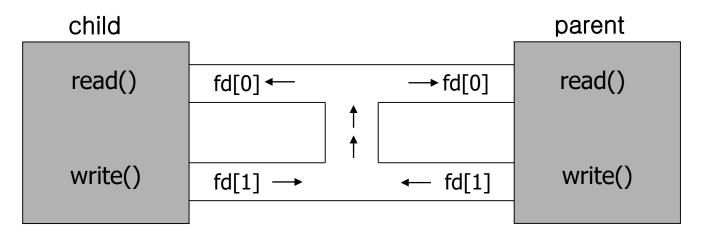


```
if ((pid = fork()) < 0) {
   perror("fork call");
   exit(2);
/* if child, then close read file
 * descriptor and write down pipe
*/
if (pid == 0) {
   close(fd[0]);
   write(fd[1], msg1, MSGSIZE);
   write(fd[1], msg2, MSGSIZE);
   write(fd[1], msg3, MSGSIZE);
```

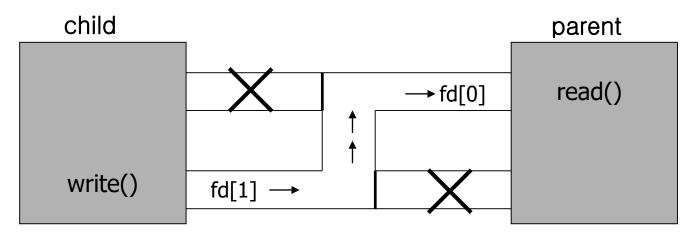
```
/* if parent, then close write file
* descriptor and read from pipe
else {
   close(fd[1]);
   for (i = 0; i < 3; i++) {
      read(fd[0], buf, MSGSIZE);
      printf("%s\n", buf);
  wait(NULL);
exit(0);
```

#### Example #1: Unnamed Pipe (3)





Fig(a). Unnamed pipe: before closing unnecessary file descriptor



Fig(b). unnamed pipe: after closing unnecessary file descriptor

#### Example #2 : Unnamed Pipe



```
#include <signal.h>
#include <unistd.h>
#include <limits.h>
int count;
void alrm_action(int);
main()
     int p[2];
     int pipe_size;
     char c = 'x';
     static struct sigaction act;
     /* 시그널 핸들러를 구축한다. */
     act.sa_handler = alrm_action;
     sigfillset (&(act.sa_mask));
     if (pipe(p) == 1)
          perror ("pipe call");
          exit (1);
```

```
/* 파이프의 크기를 결정한다. */
     pipe_size = fpathconf (p[0], _PC_PIPE_BUF);
     printf ("Maximum size of write to pipe: %d
         bytes □n", pipe_size);
     sigaction (SIGALRM, &act, NULL);
     while (1)
          alarm (20);
          write(p[1], &c, 1);
          alarm(0);
          if ((++count \% 1024) == 0)
                printf ("%d characters in pipe□n",
         count);
void alrm_action (int signo)
     printf ("write blocked after %d characters □n",
         count);
     exit (0);
```

{

#### Example #3: Unnamed Pipe (1)



```
#include <fcntl.h>
#include <errno.h>
#define MSGSIZE 6
int parent (int *);
int child (int *);
char *msg1 = "hello";
char *msq2 = "bye!!";
main()
    int pfd[2];
    /* 파이프를 개방한다 */
    if(pipe(pfd) == -1)
         fatal ("pipe call");
```

```
/* p[0]의 O_NONBLOCK 플래그를 1로 설정한다 */
    if (fcntl (pfd[0], F SETFL, O NONBLOCK) ==
       -1)
        fatal ("fcntl call");
    switch(fork()){
    case -1: /* 오류 */
        fatal("fork call");
    case 0: /* 자식 */
        child(pfd);
    default: /* 부모 */
        parent (pfd);
int parent (int p[2]) /* 부모의 코드 */
 int nread;
 char buf[MSGSIZE];
```

## Example #3: Unnamed Pipe (2)



```
close (p[1]);
                                                  int child(int p[2])
for(;;)
                                                   int count;
  switch (nread = read(p[0], buf, MSGSIZE)){
  case -1:
        /* 파이프에 아무것도 없는지 검사한다. */ close (p[0]);
        if (errno == EAGAIN)
                                                   for (count = 0; count < 3; count++)
              printf ("(pipe empty)\squaren");
             sleep (1);
                                                       write (p[1], msg1, MSGSIZE);
             break;
                                                       sleep(3);
        else fatal ("read call");
  case 0:
        /* 파이프가 닫혔음. */
                                                   /* 마지막 메시지를 보낸다 */
        printf ("End of conversation\squaren");
                                                   write (p[1], msg2, MSGSIZE);
        exit (0);
  default:
                                                   exit (0);
        printf ("MSG=%s\squaren", buf);
```

## select (1)



- #include < sys/types.h>
   #include <sys/time.h>
   #include <unistd.h>
   int select(int nfds, fd\_set \*readfds, fd\_set \*writefds, fd\_set \*errorfds,struct timeval\_\*timeout);
  - Handles multiple pipes simultaneously
  - Allows device polling.
  - nfds gives the number of the descriptors being selected.
  - readfds, writefds and errorfds point to bit masks, each bit representing a file descriptor.
    - If a bit is turned on, it denotes interest in the relevant file descriptor

## select (2)



#### Bit manipulation macros

```
#include <sys/time.h>
/* initialize the mask pointed to by fdset*/
void FD_ZERO(fd_set *fdset);
/* set the bit, fd in the mask pointed to by fdset */
void FD_SET(int fd, fd_set *fdset);
/* is the bit, fd, set in the mask pointed to by fdset */
int FD_ISSET(int fd, fd_set *fdset);
/* turn of the bit, fd, in the mask pointed to by fdset */
void FD_CLR(int fd, fd_set *fdset);
```

# select (3)



- timeout indicates how long select() should sleep, waiting for data to arrive.
  - If data arrives for any file descriptors and the *timeout* value has not expired, select() return, indicating in the bit masks which file descriptors were selected.
- For instance, if a user wished to sleep until receiving input on file descriptors 0, 1 or 2, readfds would point to the bit mask 7; when select() returns, the bit mask would be overwritten with a mask indicating which file descriptors had data ready.
- The bit mask writefds dose a similar function for write file descriptors, and the bit mask errorfds indicates when exceptional conditions exist for particular file descriptors, useful in networking.

## select (4)



#### timeout

```
    struct timeval {
        long tv_sec; /*seconds*/
        long tv_usec; /*and microseconds*/
        };
```

#### Three condition

- timeout == NULL
  - » Wait without a time limit.
- timeout ->tv\_sec == 0 && timeout->tvusec == 0
  - » No wait
- timeout ->tv\_sec != 0 || timeout->tvusec != 0
  - » Wait until the time has expired.

## select (5)



- Return value
  - -1: Error
  - 0: No descriptor is ready.
  - Positive value: The total count of the number of descriptors that are ready.

#### Example #4: select (1)



```
#include <sys/time.h>
#include <sys/wait.h>
#define MSGSIZE 6
char *msg1 = "hello";
char *msg2 = "bye!!";
void parent(int [] []);
int child(int []);
main()
     int pip[3] [2];
     int i;
```

```
/* 세 개의 통신 파이프를 생성하고, 세 개의 자식을 낳
   는다. */
    for (i = 0; i < 3; i++)
         if (pipe(pip[i]) == -1)
               fatal ("pipe call");
         switch (fork()){
         case -1:
                                /* 오류 */
               fatal ("fork call");
          case 0:
                                 /* 자식 */
               child (pip[i]);
         }
     parent (pip);
    exit (0);
}
```

#### Example #4: select (2)



```
/* 부모는 세 개의 파이프에 전부 귀를 기울이고 있다. */
void parent(int p[3] [2]) /* 부모의 코드 */
    char buf[MSGSIZE], ch;
    fd_set set, master;
    int i;
    /* 모든 원하지 않는 화일 기술자를 닫는다 */
    for (i = 0; i < 3; i++)
        close (p[i] [1]);
    /* select 시스템 호출의 비트 마스크를 설정한다. */
    FD_ZERO (&master);
    FD_SET (0, &master);
    for (i = 0; i < 3; i++)
        FD_SET (p[i] [0], &master);
```

## Example #4: select (3)



```
/* 타임아웃 없이 select를 호출한다. 사건이 발생할 때까지 select는 봉쇄될 것이다 */
 while (set = master, select (p[2][0]+1, &set, NULL, NULL, NULL) > 0)
 {
     /* 표준 입력, 즉 화일 기술자 0에 있는 정보를 잊어버리면 안됨. */
     if (FD_ISSET(0, &set))
          printf ("From standard input..."); read (0, \&ch, 1); printf ("%c\Boxn", ch);
     for (i = 0; i < 3; i++)
           if (FD_ISSET(p[i] [0], & set))
                if (read(p[i] [0], buf MSGSIZE)>0)
                   printf ("Message from child%d\squaren", i);
                   printf ("MSG=%s\squaren",buf);
```

#### Example #4: select (4)



```
/* 서버는 모든 자식이 죽으면 주 프로그램으로 복귀한다. */
        if (waitpid (-1, NULL, WNOHANG) == -1)
             return;
int child(int p[2])
    int count;
    close (p[0]);
    for (count = 0; count < 2; count++)
        write (p[1], msg1, MSGSIZE);
        /* 임의의 시간 동안 중지한다. */
        sleep (getpid() % 4);
    /* 최종 메시지를 보낸다. */
    write (p[1], msg2, MSGSIZE);
    exit (0);
```

#### Pipes and exec



- Pipe can be set up between two programs at shell level
  - \$ Is | wc
  - Open file descriptors are kept open across exec calls
  - Two pipe file descriptors opened prior to a fork/exec will still be open when the child process begins execution of the new program
  - Shell couples the standard output of Is to the write end of the pipe, and the standard input of wc to the read end

#### Example #5: pipe and exec (1)



```
/* join -- 두 명령을 파이프로 결합한다. */
int join (char *com1[], char *com2[])
{
    int p[2], status;
    /* 명령을 수행할 자식을 생성한다. */
    switch (fork()){
    case -1: /* 오류 */
        fatal ("1st fork call in join");
    case 0: /* 자식 */
        break;
    default: /* 부모 */
        wait(&status);
        return (status);
    /* 루틴의 나머지 부분으로 자식에 의해 수행된다. */
    /* 파이프를 만든다. */
    if (pipe(p) == -1)
         fatal ("pipe call in join");
```

#### Example #5: pipe and exec (2)



```
/* 다른 프로세스를 생성한다. */
    switch (fork()){
    case -1:
        /* 오류 */
        fatal ("2nd fork call in join");
    case 0:
        /* 쓰는 프로세스 */
        dup2 (p[1],1); /* 표준 출력이 파이프로 가게 한다. */
        close (p[0]); /* 화일 기술자를 절약한다. */
        close (p[1]);
        execvp (com1[0], com1);
        /* execvp가 복귀하면, 오류가 발생한 것임. */
        fatal("1st execvp call in join");
    default:
        /* 읽는 프로세스 */
        dup2(p[0], 0); /* 표준 입력이 파이프로부터 오게 한다 */
        close (p[0]);
        close (p[1]);
        execvp (com2[0], com2);
        fatal ("2nd execvp call in join");
```

#### Example #5: pipe and exec (3)



```
#include <stdio.h>
main()
    char *one[4] = {"ls", "-l", "/usr/lib", NULL};
    char *two[3] = {"grep", "^d", NULL};
    int ret;
    ret = join (one, two);
    printf ("join returned %d□n", ret);
    exit (0);
```

#### FIFOs or Named Pipes



#### Drawbacks of unnamed pipes

- Only be used to connect processes that share a common ancestry, such as a parent and its child process
- Not permanently

#### FIFO or named pipes

- \$mkfifo channel (or \$mknod channel p)
- \$\s -\la \channel
- prw-r--r-- 1 dhlee adm 0 11월 3 14:51 channel
- \$cat < channel /\* this command would be blocked \*/</p>
- \$cat < channel &</p>
- \$ls -la >! channel; wait (or \$ls -la >> channel; wait )

## mkfifo (1)



- #include <sys/types.h>
   #include <sys/stat.h>
   int mkfifo (const char \*pathname, mode\_t mode)
- Create a FIFO file (named pipe) named by the first parameter pathname with mode permissions
- Can be used between unrelated processes for data exchange.
- Once created, a FIFO must be opened using open
  - mkfifo("/tmp/fifo", 0666)
  - Fd = open("/tmp/fifo", O\_WRONLY);
- The open will block until another process opens the FIFO for reading
- Non-blocking open calls are possible with O\_NONBLOCK flag
  - fd = open("/tmp/fifo", O\_WRONLY|O\_NONBLOCK);

## mkfifo (2)



- Named pipe can be also created with mknod().
  - A value of octal 010000 must be added to the mode value to signify a FIFO.
  - if (mknod("fifo", 010600, 0) < 0)</li>perror("mknod failed");

#### Example #6: FIFO (1)



```
#include <fcntl.h>
#include <stdio.h>
#include <errno.h>
#define MSGSIZ
                      63
char *fifo = "fifo";
main (int argc, char **argv)
     int fd, j, nwrite;
     char msgbuf[MSGSIZ+1];
     if (argc < 2)
          fprintf (stderr, "Usage: sendmessage msg ... \Box n");
          exit(1);
```

#### Example #6: FIFO (2)



```
/* O_NONBLOCK을 설정하여 fifo를 개방한다. */
if ((fd = open(fifo, O WRONLY | O NONBLOCK)) < 0)
    fatal ("fifo open failed");
/* 메시지를 보낸다. */
for (j = 1; j < argc; j++)
    if (strlen(argv[j]) > MSGSIZ)
         fprintf (stderr, "message too long %s□n", argv[j]);
         continue;
    strcpy (msgbuf, argv[j]);
    if ((nwrite = write (fd, msgbuf, MSGSIZ+1)) == -1)
         fatal ("message write failed");
exit (0);
```

#### Example #6: FIFO (3)



```
/* rcvmessage -- fifo를 통해 메시지를 받는다. */
#include <fcntl.h>
#include <stdio.h>
#include <errno.h>
#define MSGSIZ
                    63
char *fifo = "fifo";
main (int argc, char **argv)
    int fd;
    char msgbuf[MSGSIZ+1];
    /* fifo가 이미 존재하지 않으면, 생성한다 */
    if (mkfifo(fifo, 0666) == -1)
         if (errno != EEXIST)
              fatal ("receiver: mkfifo");
```

#### Example #6: FIFO (3)



```
/* fifo를 읽기와 쓰기용으로 개방한다. */
if ((fd = open(fifo, O_RDWR)) < 0)
    fatal ("fifo open failed");
/* 메시지를 받는다 */
for(;;)
    if (read(fd, msgbuf, MSGSIZ+1) <0)
        fatal ("message read failed");
/*
 * 메시지를 프린트한다 ; 실제로는 보다 흥미 있는 일이 수행된다.
 */
    printf ("message received:%s□n", msgbuf);
```

#### Sample Run: FIFO



dhlee@kde:~/Course/SP/example>rcvmsg &

[1] 32615

dhlee@kde:~/Course/SP/example>sndmsg "msg1" "msg2"

dhlee@kde:~/Course/SP/example>message received:msg1

message received:msg2

dhlee@kde:~/Course/SP/example>sndmsg "msg3"

message received:msg3

dhlee@kde:~/Course/SP/example>

#### More Examples #1: Pipe (1)



```
/* pipe.c
  This program illustrates how to use the pipe() system call.
  This example shows a way of implementing who | sort using pipe().
  The output from who is redirected to a system pipe and sort gets its
  input from the system pipe.
 Algorithm outline:
   a. to create a pipe
   b. to fork a child
   c. to duplicate a file descriptor
   d. to close unused ends of the pipe
   e. to exec a process to execute the command
*/
#include <stdio.h>
int main(void)
  int fd[2];
  pipe(fd);
            /* a pipe is created: fd[0] for read; fd[1] for write */
```

#### More Examples #1: Pipe (1)



```
if (fork() == 0) { /* 1st child - output redirection */
  dup2(fd[1], 1); /* connect stdout to pipe */
  close(fd[0]);
                     /* close unneeded file descriptors */
  close(fd[1]);
  printf("\nThis is the write end of the pipe:\n");
  fflush(stdout); /* write out the msg to terminal */
  execl("/bin/who", "who", (char *)0);
  printf("\nThis line shouldn't be here");
if (fork() == 0) { /* 2nd child - input redirection */
  dup2(fd[0], 0); /* connect stdin to pipe */
  close(fd[0]); /* close unneeded file descriptors */
  close(fd[1]);
  printf("\nThis is the read end of the pipe:\n");
  execl("/bin/sort", "sort", (char *)0);
  printf("\nThis line shouldn't be here");
                      /* parent is taking it easy */
close(fd[0]);
close(fd[1]);
printf("\nParent is terminating\n");
return(0);
```

#### More Examples #2: Pipe (1)



```
/* np.c
 This program shows an example of creating named pipes.
 A named pipe is known as FIFO. It is implemented as a disk file.
  Therefore, it can outlive the program that creates it.
 Also, most of the file manipulation commands apply to FIFOs.
  However, the length of a FIFO file is alway shown as 0, and all reads
 from FIFO are destructive.
 Any of the following two sytem calls can be used to create FIFO:
   makefifo() or mknod()
*/
#include <stdio.h>
#include <fcntl.h>
#define NPNAM "fifo"
#define BSIZ BUFSIZ
#define ERR -1
#define EMPTY 0
char *mesg[] = {"yesterday", "today", "tomorrow"};
```

#### More Examples #2: Pipe (2)



```
int main(void)
  int npid;
  int npcreate(char *, int), nperror(int, char *),
     npread(int, char *), npwrite(int, int, char *[]);
  if ((npcreate(NPNAM, 0664)) == ERR) /* to create a fifo */
     nperror(1, "Can't create named pipe\n");
  if (fork() == 0) { /* writer */
     if ((npid = open(NPNAM, O_WRONLY)) == ERR)
       nperror(2, "Can't open named pipe\n");
     if (npwrite(npid, sizeof(mesg)/sizeof(char *), mesg) == ERR)
       nperror(3, "Can't write named pipe\n");
     exit(0);
  if (fork() == 0) { /* reader */
     char rbuf[BSIZ];
     if ((npid = open(NPNAM, O_RDONLY)) == ERR)
```

#### More Examples #2: Pipe (3)



```
while(1) {
       switch(npread(npid, rbuf)) {
          case ERR:
             nperror(5, "Can't read named pipe\n");
          case EMPTY:
             break;
          default:
             printf("Got: %s\n", rbuf);
             continue;
       break;
     exit(0);
  wait((int *)0);
  return(0);
} /* main */
```

#### More Examples #2: Pipe (4)



```
int npcreate(char *npnam, int perms)
             /* permissions to be determined by perms */
  umask(0);
  return(mkfifo(npnam, perms)); /* either one of these two works */
  /* return(mknod(npnam, 010000 | perms)); 010000: fifo special */
} /* npcreat */
int nperror(int errno, char *errmsg)
  write(2, errmsg, strlen(errmsg));
  exit(errno);
} /* nperror */
int npread(int npid, char *rbuf)
  return(read(npid, rbuf, BSIZ));
} /* npread */
```

#### More Examples #2: Pipe (5)

