NYU Computer Science Bridge to Tandon Course

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Homework 1

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Question 1

$$= 1 + 1 \cdot 2^1 + 1 \cdot 2^3 + 1 \cdot 2^4 + 1 \cdot 2^7$$

$$= 1 + 2 + 8 + 16 + 128$$

$$=(155)_{10}$$

$$= 6 \cdot 7^0 + 5 \cdot 7^1 + 4 \cdot 7^2$$

$$=6+35+196$$

$$=(237)_{10}$$

$$= 10 \cdot 16_0 + 8 \cdot 16^1 + 3 \cdot 16^2$$

$$= 10 + 128 + 768$$

$$= (906)_{10}$$

$$= 4 \cdot 5_0 + 1 \cdot 5^1 + 2 \cdot 5^2 + 2 \cdot 5^3$$

$$=(309)_{10}$$

$$= 64 + 4 + 1$$

$$= 69$$

$$= 1000101_2$$

 $= 111100101_2$

B3) 6D1A₁₆

$$= 10 \cdot 16^0 + 1 \cdot 16^1 + 13 \cdot 16^2 + 6 \cdot 6^3$$

$$= 10 + 16 + 3328 + 24576$$

 $=27930_{10}$

 $= 0110110100011010_2$

C1) 1101011₂

$$\begin{array}{c|cccc}
0110 & 1011 \\
= 2^2 + 2^1 & = 2^3 + 2^1 + 2^0 \\
= 4 + 2 & = 8 + 2 + 1 \\
= 6 & = 11 \Longrightarrow B
\end{array}$$

$$=6B_{16}$$

C2) 895₁0

$$\begin{array}{c|cccc} 895 \div 16 & = 55 \text{ R15} \Longrightarrow F \\ 55 \div 16 & = 3 \text{ R7} \\ 3 \div 16 & = 0 \text{ R3} \end{array}$$

 $= 37F_{16}$

$$1)7266_8 + 4515_8$$

$$7566_8 \\
+4515_8 \\
=14303_8$$

$$2)\ 10110011_2 + 1101_2$$

$$\begin{array}{r}
1111111\\10110011_2\\+00001101_2\\=11000002_2\end{array}$$

3)
$$7A66_16 + 45C5_16$$

$$7 10 06 06_{16} +4 05 12 05_{16} =12 0 02 11_{16}$$

$$\Longrightarrow C02B_{16}$$

4)
$$3022_5 + 2433_5$$

$$\implies 34_5$$

$$A1)124_{10}$$

$$= 0\ 1\ 1\ 1\ 1\ 1\ 0\ 0_{8-bit\ {}_{2^{\cdot s}\ comp}}$$

$$01101100 \\ +1000\ 0100$$

A3) 109₁₀

$$= 0\ 1\ 1\ 0\ 1\ 1\ 0\ 1_{8-bit\ {}_{2`s\ comp}}$$

$$A4) -79_{10}$$

$$\implies$$
 convert to 79_{10}

$$= 0 1 0 0 1 1 1 1$$

$$\begin{array}{c} 0 \ 1 \ 0 \ 0 \ 1 \ 1 \ 1 \ 1 \\ +1 \ 0 \ 1 \ 1 \ 0 \ 0 \ 0 \ 1 \\ \hline 1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \end{array}$$

$$= 1 \ 0 \ 1 \ 1 \ 0 \ 0 \ 0 \ 1_{8-bit} \ _{2`s \ comp}$$

B1) 0001111108-bit
$$_{2^{i_s}\ comp}$$

$$= 1 \cdot 2^1 + 1 \cdot 2^2 + 1 \cdot 2^3 + 1 \cdot 2^4$$

$$=2+4=8+16$$

$$=30_{10}$$

B2)
$$11100110_{8-bit}$$
 2's comp

$$0\ 0\ 0\ 1\ 1\ 0\ 1\ 0_{8-bit}\ _{2`s\ comp} =$$

$$= 1 \cdot 2^1 + 1 \cdot 2^3 + 1 \cdot 2^4$$

$$= 2 + 8 + 16$$

$$=26_{10}$$

$$1\ 1\ 1\ 0\ 0\ 1\ 1\ 0_{8-bit\ 2's\ comp} =$$

$$= -26_{10}$$

B3)
$$00101101_{8-bit\ 2^{\circ}s\ comp}$$

$$= 1 \cdot 2^0 + 1 \cdot 2^2 + 1 \cdot 2^3 + 1 \cdot 2^5$$

$$= 1 + 4 + 8 + 32$$

$$=45_{10}$$

$$1\; 0\; 0\; 1\; 1\; 1\; 1\; 0$$

$$+0\ 1\ 1\ 0\ 0\ 1\ 0$$

$$0\ 1\ 1\ 0\ 0\ 0\ 1\ 0_{8-bit\ 2's\ comp} =$$

$$= 1 \cdot 2^1 + 1 \cdot 2^5 + 1 \cdot 2^6$$

$$=2+32+64$$

$$=98_{10}$$

$$1\ 0\ 0\ 1\ 1\ 1\ 1\ 0_{8-bit\ 2's\ comp} =$$

$$= -98_{10}$$

1b)
$$\neg (p \lor q)$$

| p | q | $\neg (p \lor q)$ |
|---|---|-----------------------------------|
| T | Т | $\neg T \lor T \Longrightarrow F$ |
| T | F | $\neg T \lor F \Longrightarrow F$ |
| F | T | $\neg F \lor T \Longrightarrow F$ |
| F | F | $\neg F \lor F \Longrightarrow T$ |

1c) r
$$\lor (p \land \neg q)$$

| p | q | r | $(p \land \neg q)$ | $r \lor (p \land \neg q)$ |
|---|----------|----------|--------------------|---------------------------|
| T | Т | Т | F | T |
| T | Γ | F | \mathbf{F} | F |
| T | F | Γ | ${ m T}$ | ight] ight. |
| T | F | F | ${ m T}$ | ight] T |
| F | Γ | Γ | \mathbf{F} | ight] T |
| F | Γ | F | \mathbf{F} | F |
| F | F | Γ | \mathbf{F} | ight] T |
| F | F | F | \mathbf{F} | F |

2b)
$$(p \to q) \to (q \to p)$$

| p | q | $(p \rightarrow q)$ | $(q \rightarrow p)$ | $(p \rightarrow q) \rightarrow (q \rightarrow p)$ |
|---|----------|---------------------|---------------------|---|
| T | Т | Τ | Т | T |
| T | F | \mathbf{F} | T | m T |
| F | Γ | ${ m T}$ | F | F |
| F | F | ${ m T}$ | Τ | T |

2d) (p
$$\leftrightarrow$$
 q) \oplus (p \leftrightarrow \neg q)

| p | q | $(p \leftrightarrow q)$ | $(\mathbf{p} \leftrightarrow \neg q)$ | $(p \leftrightarrow q) \oplus (p \leftrightarrow \neg q)$ |
|---|----------|-------------------------|---------------------------------------|---|
| T | Т | Т | F | T |
| T | F | F | T | T |
| F | Γ | F | ${ m T}$ | ${ m T}$ |
| F | F | T | F | Т |

1b) The applicant must present at least two of the following forms of identification: birth certificate, driver's license, marriage license.

$$=(B\wedge D)\vee (B\wedge M)\vee (D\wedge M)$$

1c) Applicant must present either a birth certificate or both a driver's license and a marriage license.

$$= B \vee (D \wedge M)$$

2b) A person can park in the school parking lot if they are a senior or at least seventeen years of age.

$$= p \rightarrow (s \lor y)$$

2c) Being 17 years of age is a necessary condition for being able to park in the school parking lot.

$$= p \rightarrow y$$

2d) A person can park in the school parking lot if and only if the person is a senior and at least 17 years of age.

$$= p \leftrightarrow (s \land y)$$

2e) Being able to park in the school parking lot implies that the person is either a senior or at least 17 years old.

$$=\mathbf{p}\rightarrow (s\vee y)$$

3c) The applicant can enroll in the course only if the applicant has parental permission.

$$= c \rightarrow p$$

3d) Having parental permission is a necessary condition for enrolling in the course.

$$= c \rightarrow p$$

1b) Maintaining a B average is necessary for Joe to be eligible for the honors program.

If Joe wants to be eligible for the program, then maintaining a B average is necessary.

1c) Rajiv can go on the roller coaster only if he is at least four feet tall.

= If Rajiv goes on the roller coaster, then he is at least four feet tall.

1d) Rajiv can go on the roller coaster if he is at-least four feet tall.

= If Rajiv is at-least four feet tall, then he can go on the roller coaster.

2c)
$$(p \lor r) \leftrightarrow (q \land r)$$

p:T

q:F

r: Unknown

| p | q | r | $(p \lor r)$ | $(q \land r)$ | $(p \lor r) \leftrightarrow (q \land r)$ |
|---|----------|----------|--------------|---------------|--|
| T | Т | Т | Τ | Т | Т |
| T | Γ | F | Τ | F | F |
| T | F | Γ | ${ m T}$ | \mathbf{F} | F |
| T | F | F | ${ m T}$ | F | F |
| F | Γ | T | ${ m T}$ | Τ | Т |
| F | Γ | F | ${ m T}$ | F | F |
| F | F | Т | Τ | \mathbf{F} | F |
| F | F | F | ${ m T}$ | F | F |

False.

2d) (p
$$\wedge r$$
) \leftrightarrow ($q \wedge r$)

| p | q | r | $(p \land r)$ | $(q \land r)$ | $(p \land r) \leftrightarrow (q \land r)$ |
|---|----------|--------------|---------------|---------------|---|
| T | Т | Т | Т | Т | Т |
| T | Γ | F | \mathbf{F} | F | F |
| T | F | \mathbf{T} | ${ m T}$ | F | F |
| T | F | F | F | F | F |
| F | Т | Т | F | Т | T |
| F | Γ | F | F | F | F |
| F | F | Т | F | F | F |
| F | F | F | F | F | F |

Unknown. When "r" is T, then conclusion is F; When "r" is F, then the conclusion is T.

2e) p
$$\rightarrow (r \lor q)$$

| p | q | r | $(r \land q)$ | $p \to (r \lor q)$ |
|---|----------|--------------|---------------|--------------------|
| T | Т | Т | Τ | T |
| T | Γ | F | F | \mathbf{F} |
| T | F | \mathbf{T} | \mathbf{F} | F |
| T | F | F | \mathbf{F} | F |
| F | Γ | Γ | ${ m T}$ | T |
| F | Γ | F | \mathbf{F} | T |
| F | F | Т | \mathbf{F} | T |
| F | F | F | \mathbf{F} | T |

Unknown. When "r" is T, then conclusion is T; When "r" is F, then the conclusion is T.

2f) $(p \land q) \rightarrow r$

| p | q | r | $(p \land q)$ | $(p \land q) \to r$ |
|---|----------|----------|---------------|---------------------|
| T | Т | Т | Т | Т |
| T | Γ | F | Т | F |
| T | F | Γ | F | T |
| T | F | F | F | ${ m T}$ |
| F | Γ | Γ | F | ${ m T}$ |
| F | Γ | F | F | ${ m T}$ |
| F | F | Γ | F | ${ m T}$ |
| F | F | F | F | ${ m T}$ |

True. The hypothesis is true and the conclusion is True.

• j: Sally got the job.

• 1: Sally was late for her interview

• r: Sally updated her resume.

1b) If Sally did not get the job, then she was late for interview or did not update her resume. $\neg j \rightarrow (l \lor \neg r)$

If Sally updated her resume and was not late for her interview, then she got the job. $(r \land \neg l) \rightarrow j$

| j | l | r | $\neg j \to (l \lor \neg r)$ | $(\mathbf{r} \land \neg l) \to j$ |
|---|----------|--------------|------------------------------|-----------------------------------|
| T | Т | Т | Τ | T |
| T | Γ | F | ${ m T}$ | T |
| T | F | \mathbf{T} | ${ m T}$ | ${ m T}$ |
| T | F | F | ${ m T}$ | ${ m T}$ |
| F | Γ | Γ | ${ m T}$ | ${ m T}$ |
| F | Γ | F | ${ m T}$ | ${ m T}$ |
| F | F | Т | \mathbf{F} | F |
| F | F | F | ${ m T}$ | T |

$$\neg j \to (l \vee \neg r) \equiv (r \wedge \neg l) \to j$$

1c) If Sally got the job then she was not late for her interview. (j $\rightarrow \neg l)$

If Sally did not get the job, then she was late for her interview. $(\neg j \rightarrow l)$

| j | l | $(j \rightarrow \neg l)$ | $(\neg j \to l)$ |
|---------------|---|--------------------------|------------------|
| T | Т | F | ${ m T}$ |
| $\mid T \mid$ | F | ${ m T}$ | ${ m T}$ |
| F | T | ${ m T}$ | ${ m T}$ |
| F | F | ${ m T}$ | \mathbf{F} |

$$(j \to \neg l) \not\equiv (\neg j \to l)$$

1c) If Sally updated her resume or she was not late for her interview, then she got the job. $(r \lor \neg l) \to j$

If Sally got the job, then she updated her resume and was not late for her interview. $j \to (r \land \neg l)$

| j | l | r | $(\mathbf{r} \vee \neg l) \to j$ | $\mathbf{j} \to (r \land \neg l)$ |
|---------------|----------|----------|----------------------------------|-----------------------------------|
| T | Т | Т | T | F |
| T | Γ | F | \mathbf{F} | T |
| $\mid T \mid$ | F | Γ | ${ m T}$ | Γ |
| T | F | F | \mathbf{F} | ${ m T}$ |
| F | Γ | Γ | ${ m T}$ | F |
| F | Γ | F | ${ m F}$ | Γ |
| F | F | Γ | ${ m T}$ | F |
| F | F | F | ${ m T}$ | ${ m T}$ |

$$(\mathbf{r} \vee \neg l) \not\equiv \mathbf{j} \to (r \wedge \neg l)$$

1c)
$$(p \to q) \land (p \to r) \equiv p \to (q \land r)$$

$$\begin{array}{c|c} (\neg p \lor q) \land (\neg p \lor r) & \text{Conditional Identity} \\ (\neg p \lor (q \land r)) & \text{Distributive Law} \\ p \to (q \land r) & \text{Conditional Identity} \end{array}$$

1f)
$$\neg (p \lor (\neg p \land q)) \equiv \neg p \land \neg r$$

$$\begin{array}{c|c} \neg p \wedge \neg (\neg p \wedge q) \\ \neg p \wedge (\neg \neg p \vee \neg q) \\ \neg p \wedge (p \vee \neg q) \\ (\neg p \wedge p) \vee (\neg p \wedge \neg q) \\ \text{False } \vee (\neg p \wedge \neg q) \\ (\neg p \wedge \neg q) \vee F \\ (\neg p \wedge \neg q) \end{array} \begin{array}{c} \text{De Morgan's Law} \\ \text{Double Negation} \\ \text{Distributive Law} \\ \text{Complement Law} \\ \text{Commutative Law} \\ \text{Identity Law} \end{array}$$

1i)
$$(p \land q) \rightarrow r \equiv (p \land \neg r) \rightarrow \neg q$$

$$\neg (p \land q) \lor r$$

$$(\neg p \land \neg q) \lor r$$

$$r \lor (\neg p \lor \neg q)$$

$$(r \lor \neg p) \lor \neg q$$

$$\neg \neg (r \lor \neg p) \lor \neg q$$

$$\neg (r \lor \neg p) \to \neg q$$

$$\neg (\neg p \lor r) \to \neg q$$

$$(\neg \neg p \land \neg r) \to \neg q$$

$$(p \land \neg r) \to \neg q$$

Conditional Identity
De Morgan's Law
Commutative Law
Associative Law
Double Negation Law
Conditional Identity
Commutative Law
De Morgan's Law
Double Negation Law

2c)
$$\neg r \lor (\neg r \to p)$$

$$\neg r \lor (\neg \neg r \to p)
 \neg r \lor (r \to p)
 ((\neg r \lor r) \lor p)
 ((r \lor \neg r) \lor p)
 T \lor p
 p \lor T
 T$$

Conditional Identity
Double Negation Law
Associative Law
Commutative Law
Complement Law
Commutative Law
Domination Law

2d)
$$\neg (p \to q) \to \neg q$$

| $\neg(\neg p \lor q) \to \neg q$ |
|------------------------------------|
| $\neg\neg(\neg p\vee q)\vee\neg q$ |
| $(\neg p \lor q) \lor \neg q$ |
| $\neg p \lor (q \lor \neg q)$ |
| $\neg p \lor T$ |
| ${ m T}$ |

Conditional Identity
Conditional Identity
Double Negation Law
Associative Law
Complement Law
Domination Law

1c) There is a number that is equal to its square.

$$\exists x(x=x^2)$$

1d) Every number is less than or equal to its square.

$$\forall x (x \le x^2)$$

- S(x): x was sick yesterday
- W(x): x went to work yesterday
- V(x): x was on vacation yesterday

2b) Everyone was well and went to work yesterday.

$$\forall x (\neg S(x) \land W(x))$$

2c) Everyone who was sick yesterday did not go to work.

$$\forall x (S(x) \to \neg W(x))$$

2d) Yesterday someone was sick and went to work.

$$\exists x (S(x) \land W(x))$$

1c)
$$\exists x ((x = c) \rightarrow P(x))$$

False.

1d)
$$\exists x (Q(x) \land R(x))$$

TRUE.

1e) Q(a)
$$\wedge P(d)$$

TRUE.

1f)
$$\forall x ((x \neq b) \rightarrow Q(x))$$

TRUE.

1g)
$$\forall x (P(x) \lor R(x))$$

False. Counter-example: c.

1h)
$$\forall x (R(x) \to P(x))$$

TRUE.

1i)
$$\exists x (Q(x) \lor R(x))$$

TRUE.

2b)
$$\exists x \forall y Q(x,y)$$

TRUE.

2c)
$$\exists x \forall y P(y, x)$$

TRUE.

2d)
$$\exists x \exists y S(x, y)$$

False.

2e)
$$\forall x \exists y Q(x,y)$$

False.

2f) $\forall x \exists y P(x, y)$

TRUE.

2g) $\forall x \forall y P(x, y)$

False.

2h) $\exists x \exists y Q(x,y)$

TRUE.

2i) $\forall x \forall y \neg S(x,y)$

TRUE.

1c) There are two numbers whose sum is equal to their product.

$$\exists x \exists y (x + y = xy)$$

1d) The ratio of every two positive numbers is also positive.

$$\forall x \forall y (((x>0) \land (y>0)) \rightarrow (x/y>0)$$

1e) The reciprocal of every positive number less than one is greater than one.

$$\forall x (((x>0) \land (x<1)) \rightarrow (1/x>1))$$

1f) There is no smallest number.

$$\neg \exists x \forall y (x < y)$$

1g) Every number besides 0 has a multiplicative inverse.

$$\forall x \exists y ((x \neq 0) \rightarrow (xy = 1))$$

- P(x, y): x knows y's phone number. (A person may or may not know their own phone number.)
- D(x): x missed the deadline.
- N(x): x is a new employee.
- 2c) There is at least one new employee who missed the deadline.

$$\exists x (N(x) \land D(x))$$

2d) Sam knows the phone number of everyone who missed the deadline.

$$\forall y(D(y) \to P(Sam, y))$$

2e) There is a new employee who knows everyone's phone number.

$$\exists x \forall y (N(x) \land P(x,y))$$

2f) Exactly one new employee missed the deadline.

$$\exists x \forall y ((N(x) \land D(x)) \land (((x \neq y \land N(y)) \rightarrow \neg D(y)))$$

3c) Every student has taken at least one class besides Math 101.

$$\forall x \exists y (T(x,y) \land (y \neq Math101))$$

3d) There is a student who has taken every math class besides Math 101.

$$\exists x \forall y ((y \neq Math101) \rightarrow T(x, y))$$

3e) Everyone besides Sam has taken at least two different math classes.

$$\forall x\exists y\exists z(((x\neq Sam))\rightarrow ((y\neq z)\wedge T(x,y)\wedge T(x,z)))$$

3f) Sam has taken exactly two math classes.

$$\exists y \exists z \forall w (((y \neq z) \land T(Sam, y) \land T(Sam, z)) \land (((w \neq y) \land (w \neq z)) \rightarrow \neg T(Sam, w)$$

• P(x): x was given the placebo

• D(x): x was given the medication

• M(x): x had migraines

1b) Every patient was given the medication or the placebo or both.

 $\forall x (D(x) \lor P(x))$

Negation: $\neg \forall x (D(x) \lor P(x))$

Applying De Morgan's Law: $\exists x (\neg D(x) \land P(x))$

English: Some patient was not given medication and not given placebo

1c) There is a patient who took the medication and had migraines.

 $\exists x (D(x) \land M(x))$

Negation: $\neg \exists x (D(x) \land M(x))$

Applying De Morgan's Law: $\forall x (\neg D(x) \lor \neg M(x))$

English: Every patient was either not given medication or not had migraines(or both).

1d) Every patient who took the place bo had migraines. (Hint: you will need to apply the conditional identity, $p \to q \equiv \neg p \lor q$.)

 $\forall x (P(x) \to M(x))$

Negation: $\neg \forall x (P(x) \to M(x))$

Applying De Morgan's Law: $\exists x (\neg P(x) \land \neg M(x))$

English: Some patient who took the placebo and did not have migraines.

1e) There is a patient who had migraines and was given the placebo.

 $\exists x (M(x) \land P(x))$

Negation: $\neg \exists x (M(x) \land P(x))$

Applying De Morgan's Law: $\forall x (\neg M(x) \lor \neg P(x))$

English: Every patient either not had migraines or was not given the placebo (or both).

2c) $\exists x \forall y (P(x,y) \rightarrow Q(x,y))$

 $\forall x \exists y (P(x,y) \land \neg Q(x,y))$

2d) $\exists x \forall y (P(x,y) \leftrightarrow P(y,x))$

 $\forall x \exists y ((P(x,y) \land \neg P(y,x)) \lor (P(y,x) \land \neg P(x,y)))$

2e) $\exists x \exists y P(x,y) \land \forall x \forall y Q(x,y)$

 $\forall x \forall y \neg P(x,y) \lor \exists x \exists y \neg Q(x,y)$