Note Title

4/10/2022

In the analysis of line search methods, we shall assume that the objective functions $f: \mathbb{R}^n \to \mathbb{R}$ is C', and that its gradient is Lipschitz, i.e.

Il \f(x) - \f(y) || ≤ L || x-y|| \forall x, y ∈ either the whole or part of R^n

choice of norm cannot affect the existence of L, but affects its numerical value.

Typically, L depends on how the domain of f is restricted (and hence on xo).

E.g. $f(x) = x^{\beta} \beta > 2$

|f(x)-f(y)|=|f''(g)||x-y| for some $g\in [x,y]$.

| β(β-1) xβ-2 | ≤ β(β-1) |x0|β-2 for |x| ≤ |x0|.

so L can be chosen to be B(B-1) |xol B-2.

Also, 由L st lf(x)-f'(y)1 < L 1x-y1 \ \x, y \ R.

Write C'L'(U) := {f∈C'(Rn) | 117f(x)-7f(y)11 ≤ L11x-y11, ∀x,y∈U}

E.g. Linear functions $f(x) = a^{T}x + b$ is in $C_0^{1,1}(\mathbb{R}^n)$

Quadratics $f(x) = \frac{1}{2}x^TAx + bx + C$ is C^{∞} , does its gradient have a uniform Lipshitz bound over the whole \mathbb{R}^n ?

Yes it does: $\nabla f(x) = Ax + b$

 $\|\nabla f(x) - \nabla f(y)\| = \|A(x-y)\| \le \|A\|\|x-y\|$

Recall: || All2 = Omax (A)

if A>O, IIAlla = Imax(A).

the matrix norm induced by whichever vector norm used to define Lipshitz continuity ie. ||A||!= max ||Ax||

It's easy to show that, using the fundamental theorem of calculus, that if $f \in C^2(\mathbb{R}^n)$, $f \in C^1(\mathbb{R}^n) \iff \|\nabla^2 f(x)\| \leq L$, $\forall x \in ConvexHull(U)$.

Two global convergence results

Theorem (Zoutendijk) Let $f \in C^1(\mathbb{R}^n)$ be bounded below. Consider any line search method $x_{k+1} = x_k + x_k p_k$, where

• Pk is a descent direction $(\nabla f(x_k)^T p_k < 0)$

• de is a Step size that satisfies the Wolfe conditions $\varphi(\alpha) = \varphi(x_k + \alpha p_k), \quad \varphi(\alpha_k) \leq \varphi(0) + C_1 \alpha_k \varphi'(0), \quad \varphi'(\alpha_k) \geqslant C_2 \varphi'(0), \quad 0 < C_1 < C_2 < 1$

If fecti(x) for x!={xern: f(x)≤f(xo)}.

Then

 $\sum_{k \geqslant 0} \cos^2 \theta_k ||\nabla f_k||_2^2 < \infty$, where $\cos \theta_k = -\nabla f_k^T P_k / ||\nabla f_k||_2 \cdot ||P_k||_2$.

Proof: The 2nd Wolfe condition gives

$$(\nabla f_{RH} - \nabla f_{R})^{T} p_{R} = \phi'(\omega_{R}) - \phi'(0) \ge (c_{2}-1)\phi'(0) = -(1-c_{2})\phi'(0)$$

while Lipschitz condition gives (VfRH - VfR) TPR & [|| MRPR || || PR || = ex [|| PR || 2.

The two inequalities imply: $\alpha \geq \frac{1-c_2}{L} \frac{\phi'(0)}{\|p_k\|_2^2}$

Substitute this lower bound of de into the 1st Wolfe condition,

$$f_{R+1} \leq f_R - c_1 \frac{1-c_2}{L} \frac{\phi'(0)^2}{\|\rho_R\|_2^2} = \left(\nabla f_R^T \rho_R\right)^2$$

$$= f_R - c_1 \frac{1-c_2}{L} \cos^2 \theta_R \|\nabla f_R\|_2^2$$

$$= c$$

So, $f_{RH} \leq f_0 - c \sum_{j=0}^{R} cos^2 \theta_j ||\nabla f_j||^2$

But f is bounded below, so $f_{RH} > -\infty$ and $\sum_{j=0}^{R} \cos^2\theta_j \|\nabla f_j\|^2 < \infty$. Q.E.D.

Called the Zoutendijk condition

Condition (Z) \Rightarrow $\lim_{k \to \infty} \cos^2 \theta_k ||\nabla f_k||^2 = 0$

Under the extra condition that $\cos \theta_k \ge 8 > 0$ $\forall k$, $\lim_{k \to \infty} \|\nabla f_k\| = 0$

For GD, $\cos\theta_k = 1$, $\forall k$, so we can conclude that any cluster point of x_k is a stationary point of f.

For GD at least, we can dispense with the (full) Wolfe condition and use backtracking to attain the first Wolfe (aka Armijo) condition.

Recall:

Armijo condition	satisfied by any small enough o	backtracking alg. easy to implement
Wolfe condition	not satisfied by arb. small oc	alg. complicated to implement

Theorem (convergence of GD) Let $f \in C_L^{ll}(\mathbb{R}^n)$ be bounded below.

Let $X_{RH} = X_R - \alpha_R \nabla f(X_R)$, R > 0, with α_R chosen with one of the following stepsize strategies:

• constant step size $d_R = \overline{d} \in (0, \frac{a}{L})$

(Choice requires knowing L)

· exact line search

(too expensive in general)

• backtracking with parameters a, CE(0,1), PE(0,1). (most practical/robust)

Assume fection, 2:- Ex: fix> f(x)?.

Then $f(x_R)$ is non-increasing and convergent, and $\lim_{R \to \infty} \|\nabla f_R\| = 0$.

Ex: Compare the two theorems. Make sure you understand the similarities and differences of what the two results say.

Under the setting of either theorem,

 $f(x_{RH}) < f(x_R)$, unless $\nabla f(x_R) = 0$ (in which case the iteration terminates.)

Nether theorem (in the way it is stated) says anything about rate of convergence. The analysis of rate of convergence is very subtle in this setting. Recall from last week that it is impossible to get linear convergence for GD.

Proof

Similar to the previous proof, we get the desired result if we can show that

 $f_{k+1} \leq f_k - M \cos^2\theta_i \|\nabla f_k\|^2$ for some constant M>0. — (M)

- [Since f_R is \downarrow and bounded below, f_R must be convergent. But then $M \|\nabla f_R\|^2 \le f_R f_{RH} \to 0$, which also implies $\|\nabla f_R\| \to 0$.]
- 2 Descent Lemma: \(\forall x, y, \) | f(y) f(x) \(\forall f(x)^T(y-x)\) \(\leq \forall 2 \) | x y | \(\forall 2 \)

Proof: By the fundamental thm of calculus,
$$= \frac{d}{dt} f(x+t(y-x))$$

$$f(y) - f(x) = \int_{0}^{1} \left\langle \nabla f(x+t(y-x), y-x) \right\rangle dt$$

$$= \left\langle \nabla f(x), y-x \right\rangle + \int_{0}^{1} \left\langle \nabla f(x+t(y-x), y-x) - \left\langle \nabla f(x), y-x \right\rangle dt$$
Thus,
$$|f(y) - f(x) - \left\langle \nabla f(x), y-x \right\rangle| = |\int_{0}^{1} \left\langle \nabla f(x+t(y-x), y-x) - \left\langle \nabla f(x), y-x \right\rangle dt|$$

$$\leq \int_{0}^{1} |\left\langle \nabla f(x+t(y-x) - \nabla f(x)), y-x \right\rangle| dt$$

$$\leq \int_{0}^{1} |\left\langle \nabla f(x+t(y-x) - \nabla f(x))| ||y-x|| dt$$

$$\leq \int_{0}^{1} |t| ||y-x||^{2} dt = \frac{L}{2} ||y-x||^{2}.$$

3 Sufficient decrease lemma: For any $x \in \mathbb{R}^{n}$, d > 0, (specific for GD) $f(x) - f(x - d \nabla f(x)) \ge d (1 - \underline{L}\underline{d}^{2}) ||\nabla f(x)||^{2} \leftarrow \frac{Proof:}{Proof:} By (2), f(x - d \nabla f(x)) \le f(x) - d ||\nabla f(x)||^{2} + \underline{L}\underline{d}^{2} ||\nabla f(x)||^{2} = f(x) - d (1 - \underline{L}\underline{d}^{2}) ||\nabla f(x)||^{2}.$

4 By 3, we see that a constant step size
$$\overline{a} \in (0, \frac{a}{L})$$
 guarantees

$$f(x_{k}) - f(x_{k+1}) \ge \overline{\alpha}(1 - \frac{L\overline{\alpha}}{2}) ||\nabla f_{k}||^{2}$$

$$\times_{k} - \overline{\alpha} ||\nabla f_{k}||^{2}$$

So in this case, the constant M desired in ① can be set to $M = \overline{\lambda}(1 - \frac{L\overline{\lambda}}{2})$.

Note: $\overline{\lambda} = \frac{1}{L}$ maximizes $\alpha(1-\frac{1}{2}\alpha)$ over $(0,\frac{2}{L})$, and thus a popular choice of step size. In this case

$$f_R - f_{R+1} \ge \frac{1}{2L} ||\nabla f_R||^2$$
.

Exact line search $dR \in argmin_{0>0} f(x_R - \alpha \nabla f_R)$ guarantees

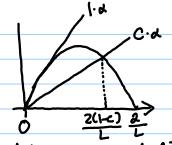
So in this case, the constant M desired in 1 can be set to M= 1/2L.

In backtracking, we seek a small enough de for which

$$f(x_R) - f(x_R - \alpha_R \nabla f_R) > C \alpha_R ||\nabla f_R||^2$$
, $C \in (0,1)$. — (SD)

By 3), f(x- d \(\nabla f(x) \ge \tau (1-\frac{1}{2}) \) | \(\nabla f(x)\)|^2.

So (SD) is satisfied as long as $\alpha \in [0, \frac{2(1-c)}{L}]$



Backtracking guarantees that d_k is either \overline{d} (when no backtracking is needed), $d_k / \partial v > \frac{\partial (1-c)}{L}$, i.e. $d_k > \frac{2(1-c)p}{L}$.

So $\forall k \geq \min\{\overline{a}, \frac{2(1-c)p}{L}\}$

By (SD), the constant M desired in 1 can be set to

 $M = C \min \{ \overline{\alpha}, \frac{2(1-c)p}{L} \}$

Q.E.D.

Note: For the backtracking case, it is easy to see that the proof goes through if we weaken the " $f \in C^{ij}(\mathbb{R}^n)''$ assumption to

 $f \in C_L^{(1)}(\mathcal{L})$, $\mathcal{L} := \{x : f(x) \leq f(x_0)\}$ (as in Zoutendijk's theorem).

But for the constant step size case, the weaker condition does not seem to be sufficient.

Thm (ROC of gradient norms) Under the setting of the previous theorem,

write $f^* = \lim_{k \to \infty} f(x_k)$. Then $\forall n > 0$, min $\|\nabla f(x_k)\| \le \sqrt{\frac{f(x_0) - f^*}{M(n+1)}}$,

where $M = \begin{cases} \overline{a}(1-\overline{a}) & \text{constant step size} \\ \frac{1}{2}L & \text{exact line search} \\ C \min\{\overline{a}, \frac{a(1-c)p}{L}\} & \text{backtracking} \end{cases}$

Proof: We proved that for < for - M 117fol12

result.

Q.E.D.

Note: This O(元) bound is probably not tight. In the convex case (see below),

it can be shown that $\|\nabla f_n\| = O(\frac{1}{n})$.

The following example tells you something subtle about GD applied to a non-convex objective.

It is possible that GD converges to a saddle point [Recall that our theorem only promises GD, when converges, converges to a Critical point. It does not promise the limit must be a local minimizer.

 $f(x) = ax_1^2 - bx_2^2$ The origin is a saddle point. Any initial point $x^0 = \begin{bmatrix} x_1^0 \\ x_2^0 \end{bmatrix}_{\neq 0}$ would be sent to ∞ by aD. But any initial point $x^0 = \begin{bmatrix} x^0 \end{bmatrix}$ stays on the x_1 -axis and is

attracted to the saddle point.

Note:
$$\nabla f(x) = \begin{bmatrix} 2ax_1 \\ -2bx_2 \end{bmatrix} + \begin{bmatrix} 2ax_1 \\ 0 \end{bmatrix}$$

However, as this and the next example suggest, for most initial points GD converges (when it converges) to a local minimizer. This has been proved rigorously.

A slightly more interesting example

Ex: (i) Prove that the only critical points are [8], [9], [9]

- (ii) Are these critical points local min, local max, or saddle points?
- (iii) what happens if GD (with any step size selection method) is applied to f if x^0 is on the x_1 -axis?
- (iv) Is $f \in C_L^{U}(\mathbb{R}^2)$ for some L>0?

Is f bounded below?

Is f coercive?

(v) with any step size selection method, is it possible to get global convergence for GD applied to this objective?

Two rate of convergence results for GD

(I) Suppose that $f \in C^2(\mathbb{R}^n)$ and that the iterates generated by the GD method with exact line searches converge to a point x^* at which the Hessian matrix $\nabla^2 f(x^*)$ is positive definite. Let

$$r \in (\frac{\lambda_n - \lambda_1}{\lambda_n + \lambda_1}, 1),$$

where $0<\lambda_1<\dots<\lambda_n$ are the eigenvalues of $\nabla^2 f(x^*)$. The for all k sufficiently large, we have

$$f(x_{k+1}) - f(x^*) \leq r^2 \left[f(x_k) - f(x^*) \right].$$

(II) Let $f \in C_L^{U}(\mathbb{R}^n)$ and convex, and admit a minimizer x^* . Then the GD method with constant step size $\overline{A} = V_L$ generates a sequence x_R such that $|\nabla f(x_R)||_2 \le \frac{2}{\sqrt{k_L(R+1)}} ||x_0 - x^*||_2$

and

$$f(x_{k}) - f(x^{*}) \leq \frac{L}{2(k+1)} \|x_{0} - x^{*}\|_{2}$$

recall the computation example at the end of UO-2.

Two implications of the ROC results

1. In finding critical point of a function $f:\mathbb{R}^n\to\mathbb{R}$, instead of solving $\nabla f(x)=0$, one might consider solving $\min_{x}\frac{1}{2}\|\nabla f(x)\|^2=h(x)$

An appealing feature of this approach is that as long as x is a critical point of f, be it a local min, local max, a saddle point etc, x is always a global minimizer of h, with minimum value O.

A problem of this method is that the conditioning of the original problem is squared, making it susceptible to very slow convergence if GD is used.

If
$$\overline{x}$$
 is a critical point of f (ie. $\nabla f(\overline{x}) = 0$), then
$$f(x) = f(\overline{x}) + \nabla f(\overline{x}) f(x-\overline{x}) + \frac{1}{2}(x-\overline{x})^T \nabla^2 f(\overline{x})(x-\overline{x}) + \cdots$$

 $\nabla f(x) = \nabla^2 f(\overline{x})(x-\overline{x}) + (higher order terms in (x-\overline{x}))$

$$h(x) = \frac{1}{2} \nabla f(x)^{T} \nabla f(x) = \frac{1}{2} (x - \overline{x})^{T} \left[\nabla^{2} f(\overline{x}) \right]^{2} (x - \overline{x}) + \text{h.o.t}$$

$$\nabla^{2} h(\overline{x}) = \left[\nabla^{2} f(\overline{x}) \right]^{2} \Rightarrow K(\nabla^{2} h(\overline{x})) = \frac{\Lambda_{max} (\nabla^{2} h(\overline{x}))}{\Lambda_{min} (\nabla^{2} h(\overline{x}))} = \frac{\Lambda_{max} (\nabla^{2} f(\overline{x}))}{\Lambda_{min} (\nabla^{2} f(\overline{x}))}$$

$$= \frac{\Lambda_{max} (\nabla^{2} f(\overline{x}))}{\Lambda_{min} (\nabla^{2} h(\overline{x}))} = \frac{\Lambda_{min} (\nabla^{2} f(\overline{x}))}{\Lambda_{min} (\nabla^{2} f(\overline{x}))} = \frac{\Lambda_{min} (\nabla^{2} f(\overline{x}))}{\Lambda_{min} (\nabla^{2} f(\overline{x}))}$$

2. Later in the course we shall study methods for Constrained optimization:

Consider (P) min f(x) st. h(x) = 0 $f, h: \mathbb{R}^n \to \mathbb{R}$

We may turn it into an unconstraint optimization problem by solving

(Pu) min $f(x) + \frac{1}{2}h(x)^2$ for a large $\mu>0$. (called a penalty method)

The larger the μ the better. (In fact, if χ_{μ}^{*} is a solution of (P_{μ}) , then we expect that $\lim_{\mu \to \infty} \chi_{\mu}^{*}$, if exists, is a solution of (P).)

Problem is that the bigger the u, the more ill-conditioned (Pw) is at a minimizer.

For example, consider a quadratic program (QP) min $\frac{1}{2}x^TAx + b^Tx$ st. $C^Tx = Co$ (Pu) is min $\frac{1}{2}x^TAx + b^Tx + \mu(c^Tx - c_0)^2$ (assume A>D for simplicity) $= \frac{1}{2}x^T(A + \mu cc^T)x + (b-2\mu c_0c)^Tx + \mu c_0^2$ The Hessian of the objective of (Pu) is A+ucc^T.

Some juicy linear algebra:

the eigenvalues of A and A+ $(\sqrt{\mu}c)(\sqrt{\mu}c)^T$ are interlaced: $2\mu \leq \cdots \leq 2\mu$

λ, ε λ, μ < λ = λ = λ = - - · · ελη ε λ, μ.

moreover, (for most vectors c) $\lambda_n^{\mu} \rightarrow \infty$ as $\mu \uparrow \infty$.

SO $\mathcal{K}(A + \mu cc^T) = \frac{\lambda_{max}(A + \mu cc^T)}{\lambda_{min}(A + \mu cc^T)} \uparrow \infty \text{ as } \mu \uparrow \infty.$

This means, to the very least, it is unsuitable to apply GD to solve (Pu) for a large penalty parameter μ .