CS 430 Introduction to Algorithms

Spring Semester, 2021

Homework 5 Version 1.2

Assigned: April 7 Due: April 21

Please respect the following guidelines for writing pseudocode:

- 1. C/Java/Python instructions are fine. But do not write object-oriented additions. Do not declare or use any class. Declare only procedures (if necessary) and explain in words what each procedure does, and what is the use of each parameter.
- 2. One instruction per line
- 3. Match the brackets with a horizontal line
- 4. Number your lines
- 5. Write down if your array is indexed $0 \dots n-1$ or $1 \dots n$.

Problem 1 Given a directed graph G = (V, E) we define the graph $G^2 = (V, E^2)$, such that $(u, w) \in E^2$ if and only if for some $v \in V$, both $(u, v) \in E$ and $(v, w) \in E$. That is, G^2 contains an edge between u and w whenever G contains a path with exactly two edges between u and w. Describe an efficient algorithm for computing the adjacency matrix of G^2 given the adjacency matrix of G. Present the pseudocode and analyze the running time in terms of |V| and |E|.

(This was on a previous final exam)

Problem 2 Exercise 22.4-1 page 614 from the textbook. It has the same number (but different page) in the second edition.

Note: the DFS-based algorithm must be used, not the one from the next problem.

Problem 3 There is another O(|V|+|E|) method for topological sort. Repeatedly find a vertex of in-degree 0, print it, and "remove it from the graph" by adjusting the in-degree of the other nodes as if this node was removed. Give pseudocode and analyze the running time. Argue correctness. Discuss what happens if the input graph is cyclic.

Problem 4 Consider the following divide-and-conquer algorithm for computing minimum spanning trees. Given a complete graph G = (V, E), partition the set V of vertices into two sets V_1 and V_2 such that $|V_1|$ and $|V_2|$ differ by at most 1. Let E_1 be the set of edges that are incident only on vertices in V_1 , and let E_2 be the set of edges that are incident only on vertices in V_2 . Recursevely solve a minimum spanning tree problem on each of the two subgraphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$. Finally, select the minimum-weight edge in e that crosses the cut (V_1, V_2) , and use this edge to unite the resulting two minimum spanning trees into a single spanning tree.

Either argue the algorithm correctly computes a minimum spanning tree of G (regardless of how the partition is done), or provide an example for which the algorithm fails, showing the run of the algorithm (where you can choose the partition) and a better spanning tree.

(This was on a previous final exam)