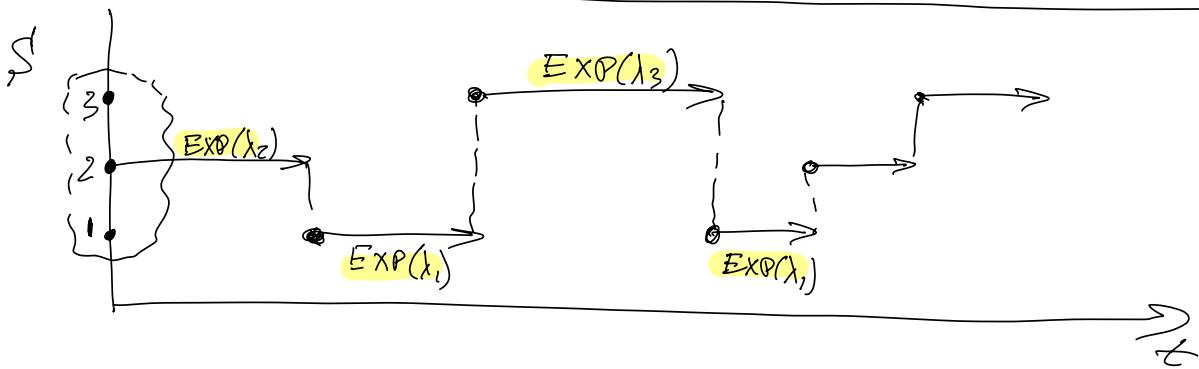


CONTINUOUS TIME MARKOV CHAINS (CTMC)



$$\begin{aligned}
 R_{12} &= \frac{3}{4} & \lambda_2 \Rightarrow & R_{21} = \frac{2}{7} \\
 R_{13} &= \frac{1}{4} & \lambda_1 = 4 & R_{23} = \frac{5}{7} \\
 && \lambda_3 = 8 & \\
 && R_{31} = \frac{5}{8}, R_{32} = \frac{3}{8} &
 \end{aligned}$$

DEFINITION: CTMC $X(t)$, $t \geq 0$.

$X(t) \in S$ (FINITE OR COUNTABLE).

$\text{EXP}(\lambda_i)$ is THE DISTR OF TIME THE PROC. $X(\cdot)$ "SITS" IN STATE $i \in S$.

$\forall i: R_{ij} = \mathbb{P}\{ \text{JUMP TO } j \text{ AFTER LEAVING } i \}$,
 $j \neq i$.

$$\sum_{j \neq i} R_{ij} = 1.$$

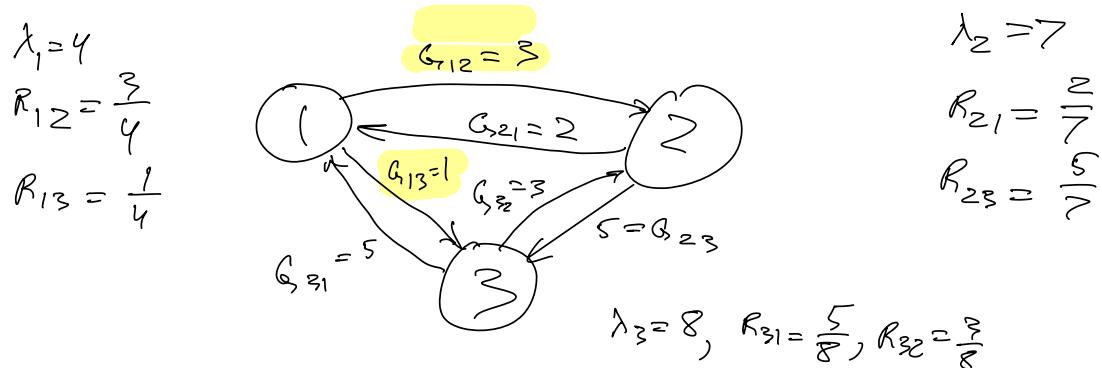
EQUIVALENT DEF. OF CTMC.

INSTEAD OF $\{\lambda_i\}$, $\{R_{ij}\}$, CTMC CAN BE EQUIVALENTLY DEFINED BY THE SET OF PARAMETERS

$$G_{ij} = \lambda_i R_{ij}, \quad j \neq i.$$

G_{ij} IS THE TRANSITION RATE FROM i TO j .

CTMC TRANSITION RATE DIAGRAM (GRAPH):



WHEN PROCESS ENTERS STATE i , FOR EACH $j \neq i$, AN INDEPENDENT "ALARM CLOCK", $\text{EXP}(G_{ij})$, IS SET. IF (ij) ALARM CLOCK RINGS FIRST, THE TRANSITION ("JUMP") FROM i TO j OCCURS.

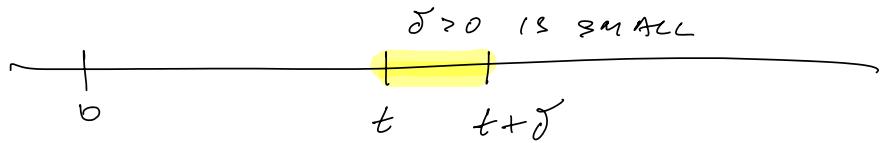
EQUIVALENT DEFINITION, BECAUSE:

TIME TO "SIT" IN i IS $\text{EXP}(\lambda_i)$, $\lambda_i = \sum_j G_{ij}$, $\forall i$

$$\text{P}\{\text{TRANS. FROM } i \text{ TO } j \neq i\} = R_{ij} = \frac{G_{ij}}{\sum_k G_{ik}}.$$

INTERPRETATION OF TRANSITION RATES G_{ij}

SUPPOSE $X(t) = i$. RECALL $\lambda_i = \sum_{j \neq i} G_{ij}$, $R_{ij} = \frac{G_{ij}}{\sum_{k \neq i} G_{ik}}$



$$\mathbb{P}\{\text{NO JUMP IN } [t, t + \delta]\} = e^{-\lambda_i \delta} = (1 - \lambda_i \delta + o(\delta)) \approx (1 - \left[\sum_{j \neq i} G_{ij} \right] \delta)$$

$$\mathbb{P}\{\text{JUMP IN } [t, t + \delta]\} \approx \left[\sum_{j \neq i} G_{ij} \right] \delta$$

$$\mathbb{P}\{\text{JUMP } i \rightarrow j \text{ IN } [t, t + \delta]\} = \mathbb{P}\{\text{JUMP IN } [t, t + \delta]\} \cdot R_{ij} \approx G_{ij} \delta$$

CTMC GENERATOR (MATRIX)

$$G_{ij}, \quad j \neq i$$

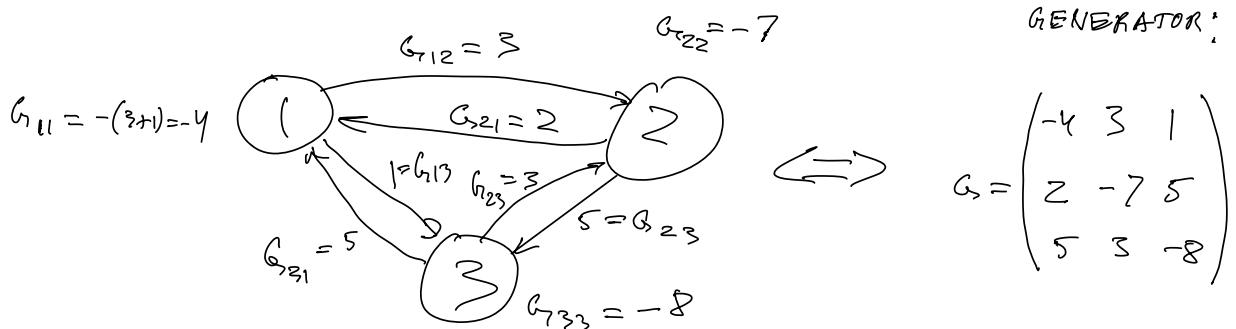
$$G_{ii} = - \sum_{j \neq i} G_{ij} = -\lambda_i$$

$n = |S| = \text{TOTAL NUMBER OF STATES (MAYBE } \infty)$

GENERATOR (MATRIX)

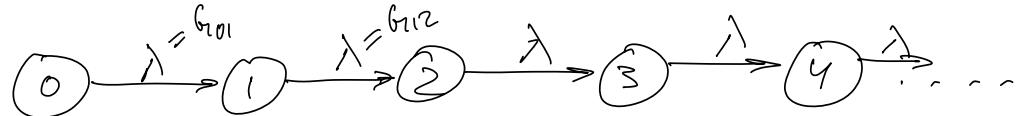
$$G = (G_{ij}) = \left(\begin{array}{c} \\ \\ \end{array} \right) \Bigg\} n$$

n



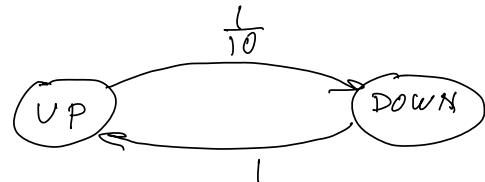
EXAMPLES OF CTMC

Ex 0. POISSON PROCESS OF RATE λ .



$$S = \{0, 1, 2, 3, \dots\}.$$

Ex 1.



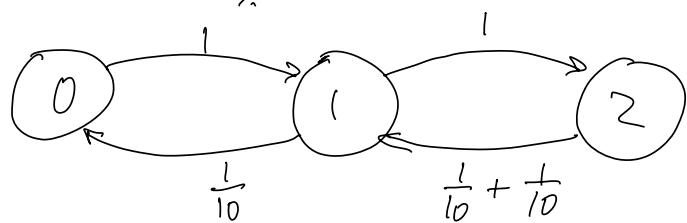
$$S = \{\text{UP}, \text{DOWN}\}$$

Ex 2.



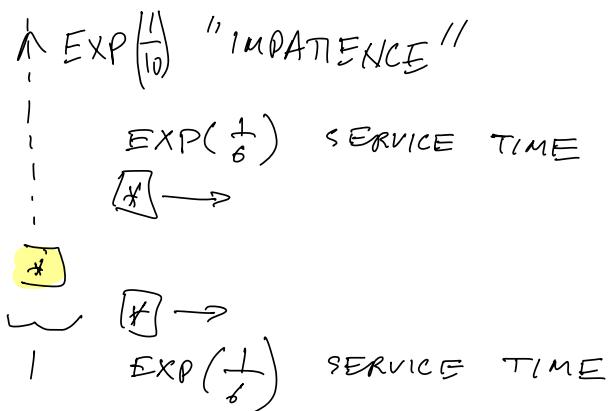
$X(t) = \# \text{ OF MACHINES WHICH ARE 'UP'}$
AT TIME t .

$$S = \{0, 1, 2\}.$$



EX 3.

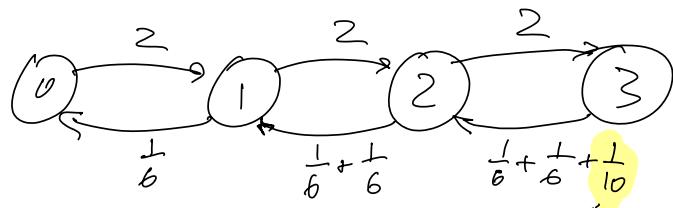
$$\lambda=2 \quad \text{POIS-PROC}(\lambda) \rightarrow$$



$X(t) = \# \text{ OF } \text{CUST. } \text{IN THE SYSTEM}$

IS A CTMC,

$$S = \{0, 1, 2, 3\}$$

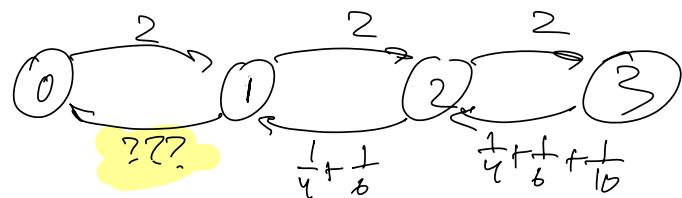


"IMPATIENCE" OF CUSTOMERS WAITING IN THE QUEUE.

Ex 4 (extension of Ex 3).

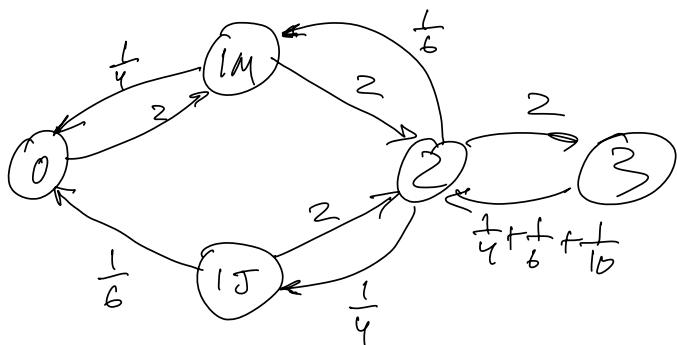
SAME AS EX 3, EXCEPT TWO AGENTS
ARE DIFFERENT

AGENT JOHN SERVICE TIME $\sim \text{EXP}(\frac{1}{6})$
— 1 — MARY ————— 1 — $\sim \text{EXP}(\frac{1}{4})$



DOES NOT
WORK.

NOT A CTMC



CTMC

MARKOV PROPERTY.

A process $(X(t), t \geq 0)$, $X(t) \in S$,
 is a CTMC if and only if
 if $t_0 < t_1 < \dots < t_{n-1} < t$, $s \geq 0$,
 if $i_0, i_1, \dots, i_{n-1}, i_j, j \in S$
 $\mathbb{P}\{X(t+s) = j | X(t_0) = i_0, \dots, X(t_{n-1}) = i_{n-1}, X(t) = i\} =$
 $= \mathbb{P}\{X(t+s) = j | X(t) = i\} \equiv P_{ij}(s)$

$P(t) = (P_{ij}(t))$, $t \geq 0$ TRANSITION PROBABILITIES' MATRIX

$$\begin{array}{l}
 \parallel P(t+s) = P(t) P(s) \\
 \quad \quad \quad \uparrow \\
 \parallel P_{ij}(t+s) = \sum_k P_{ik}(t) P_{kj}(s)
 \end{array}$$

HOW TO COMPUTE $P(t)$?

$$P(0) = I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$P'(t+s) = P(t) \cdot P'(s) \underset{s \downarrow 0}{\in} \frac{d}{ds} P(t+s) = \frac{d}{ds} [P(t)P(s)]$$

$$P'(t) = P(t) \cdot P'(0)$$

$$P'(0) = \lim_{\delta \downarrow 0} \frac{P(\delta) - \overbrace{P(0)}^I}{\delta} = G$$

PROOF:

$$P(\delta) \approx \begin{pmatrix} 1 + G_{11} \cdot \delta & \dots \\ \vdots & \ddots \\ \dots & \dots \end{pmatrix}$$

$$P(\delta) - I \approx G \cdot \delta$$

$$P'(t) = P(t)G$$

SOLUTION FOR $P(t)$:

$$\begin{cases} P'(t) = P(t) G \\ P(0) = I \end{cases} \implies P(t) = e^{Gt}$$

WHAT IS e^A ? A IS A MATRIX

FOR A REAL NUMBER x :

$$e^x = 1 + x + \frac{1}{2}x^2 + \frac{1}{3!}x^3 + \dots$$

FOR A SQUARE MATRIX A :

$$e^A = I + A + \frac{1}{2}A^2 + \frac{1}{3!}A^3 + \dots$$

VERIFICATION THAT e^{Gt} IS INDEED THE SOLUTION OF $P'(t) = P(t) G$

$$P(t) = e^{Gt} = I + Gt + \frac{1}{2}G^2t^2 + \frac{1}{3!}G^3t^3 + \dots$$

$$P'(t) = G + G^2t + \frac{1}{2}G^3t^2 + \dots =$$

$$= \underbrace{\left[I + Gt + \frac{1}{2}G^2t^2 + \dots \right]}_{P(t)} G =$$

$$= P(t)G$$

$\mu(t) = (\mu_i(t))$, ROW-VECTOR

$$\mu_i(t) = \mathbb{P}\{X(t) = i\}, \quad i \in S$$

$\mu(t)$: DISTRIBUTION AT TIME $t \geq 0$.

$\mu(0)$: INITIAL DISTRIBUTION

$$\mu(t) = \mu(0) P(t)$$



$$\mu_i(t) = \sum_k \mu_k(0) \cdot P_{ki}(t), \quad \forall i \in S$$

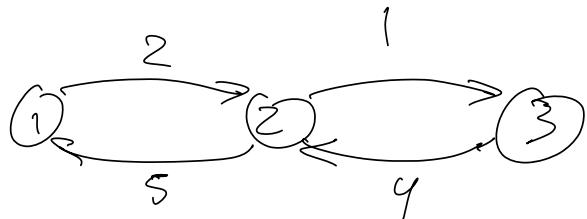
STATIONARY DISTRIBUTION

| π is STAT. DISTR. IF $\pi = \pi P(t)$, $\forall t > 0$

$$[\pi]' = \pi P'(0) \Rightarrow 0 = \pi G$$

$$\left\{ \begin{array}{l} 0 = \pi G \\ \sum_i \pi_i = 1 \end{array} \right.$$

Ex.



$$G = \begin{pmatrix} -2 & 2 & 0 \\ 5 & -6 & 1 \\ 0 & 4 & -4 \end{pmatrix}$$

$$(0, 0, 0) = (\pi_1, \pi_2, \pi_3) \cdot G$$

$$\left. \begin{array}{l} 0 = \pi_1 \cdot (-2) + \pi_2 \cdot 5 + \pi_3 \cdot 0 \\ 0 = \pi_1 \cdot 2 + \pi_2 \cdot (-6) + \pi_3 \cdot 4 \\ 0 = \pi_1 \cdot 0 + \pi_2 \cdot 1 + \pi_3 \cdot (-4) \end{array} \right\} \Rightarrow \pi$$

$\pi_1 + \pi_2 + \pi_3 = 1$

EQUATIONS FOR A STATIONARY DISTRIBUTION
USUALLY EASIER TO WRITE AS BALANCE EQ'S.
(SEE ALSO NEXT PAGE)

$$\left. \begin{array}{l} \pi_1 \cdot 2 = \pi_2 \cdot 5 \\ \pi_2 \cdot (5+1) = \pi_1 \cdot 2 + \pi_3 \cdot 4 \\ \pi_3 \cdot 4 = \pi_2 \cdot 1 \end{array} \right\} \quad \pi_1 + \pi_2 + \pi_3 = 1$$

STATIONARY DISTRIBUTION EQUATIONS

IN SCALAR FORM
(BALANCE EQUATIONS)

$$0 = \pi_i G_i$$

$$0 = \sum_j \pi_j G_{ji} \quad , \quad \forall i$$

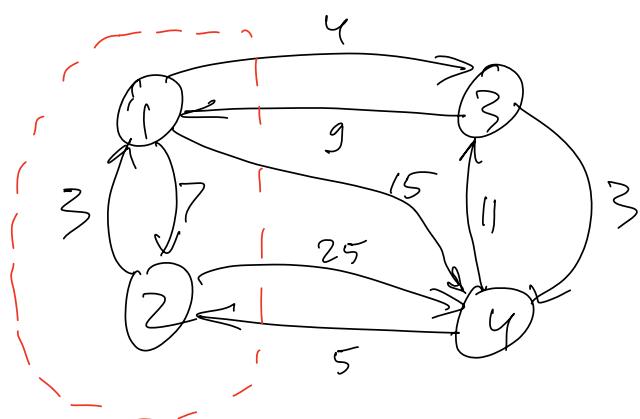
$$-\pi_i G_{ii} = \sum_{j \neq i} \pi_j G_{ji}$$

$$\pi_i \left[\sum_{j \neq i} G_{ij} \right] = \sum_{j \neq i} \pi_j G_{ji}$$

$$\sum_{j \neq i} \pi_j G_{ij} = \sum_{j \neq i} \pi_i G_{ji}, \quad \forall i \in S$$

PROB. FLOW BALANCE EQUATION FOR STATE i

BALANCE EQ. FOR A SUBSET (CUT):



$$\pi_1 \cdot 4 + \pi_1 \cdot 15 + \pi_2 \cdot 25 = \\ = \pi_3 \cdot 9 + \pi_4 \cdot 5$$

EXISTENCE AND UNIQUENESS OF STAT. DISTRIBUTIONS

CTMC $(X(t), t \geq 0)$, $X(t) \in \mathcal{S}$ (STATE SPACE)

$X(\cdot)$ IS IRREDUCIBLE IF IT CAN GET FROM ANY STATE i TO ANY OTHER STATE j
(ALL STATES COMMUNICATE)

STATE i IS POSITIVE RECURRENT IF THE AVERAGE TIME TO RETURN TO IT (AFTER LEAVING) IS FINITE

STATE i IS NULL RECURRENT IF THE PROBABILITY OF RETURNING TO IT (AFTER LEAVING) IS 1, BUT THE AVERAGE TIME TO RETURN IS INFINITE

STATE i IS TRANSIENT IF THE PROBABILITY OF RETURNING TO IT (AFTER LEAVING) IS LESS THAN 1.

IF $X(\cdot)$ IS IRREDUCIBLE, ALL STATES HAVE THE SAME RECURRENCE STATUS. IN PARTICULAR,

IF ONE STATE IS POSITIVE RECURRENT, THEN ALL STATES ARE POSITIVE RECURRENT

$X(\cdot)$ IS POSITIVE RECURRENT IF ALL STATES ARE POSITIVE RECURRENT

THERE IS NO SUCH THING AS PERIODICITY

OR APERIODICITY FOR CONTINUOUS TIME MC

TH. CONSIDER CTMC $(X(t), t \geq 0)$.

(i) IF X IS IRREDUCIBLE, THEN AT MOST ONE STAT. DISTR. EXISTS.

(ii) IF X IS FINITE (i.e. \mathcal{S} IS FINITE), THEN AT LEAST ONE STAT. DISTR. EXISTS.

(iii) SUPPOSE X IS IRREDUCIBLE. THEN

$$\begin{aligned} & \{\text{STAT. DSTR. } \pi \text{ EXISTS}\} \Leftrightarrow \\ & \Leftrightarrow \{\forall \mu(0), \mu(t) = \mu(0)P(t) \rightarrow \pi\} \Leftrightarrow \\ & \Leftrightarrow X(\cdot) \text{ IS POSITIVE RECURRENT} \end{aligned}$$

(iv) SUPPOSE X IS IRREDUCIBLE AND STAT.

DISTR. π EXISTS, THEN, $\forall \mu(0)$,

$\pi_i =$ AVERAGE LONG-TERM FRACTION OF TIME

$X(t)$ SPENDS IN STATE i ,

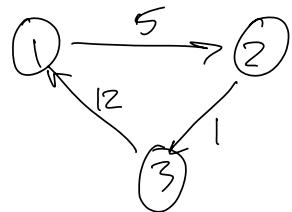
$$\text{i.e. } \lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t I\{X(u) = i\} du = \pi_i, \forall i.$$

COROLLARY. IF CTMC IS IRREDUCIBLE AND FINITE,

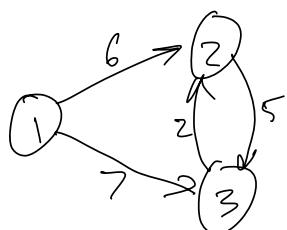
THEN IT HAS UNIQUE STAT. DISTR.

(AND THEN ALL THE NICE PROPERTIES

IN (iii), (iv) IN ABOVE TH HOLD).

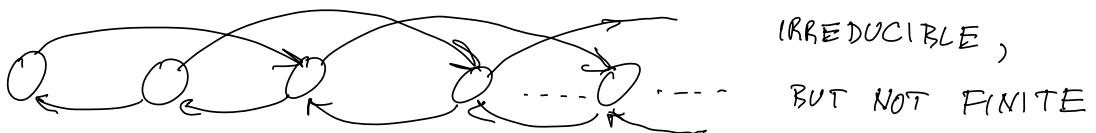


IRREDUCIBLE, FINITE
UNIQUE STAT. DISTR.



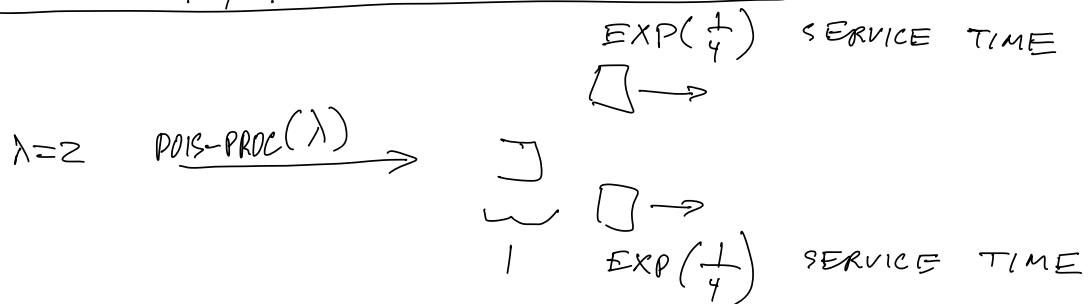
NOT IRREDUCIBLE, FINITE.
STILL HAS UNIQUE STAT. DISTR.
AND ALL THE NICE PROPERTIES
(iii), (iv) HOLD.

So, IRREDUCIBILITY IS NOT
A NECESSARY CONDITION FOR
THE EXISTENCE/UNIQUENESS OF
STAT. DISTR.



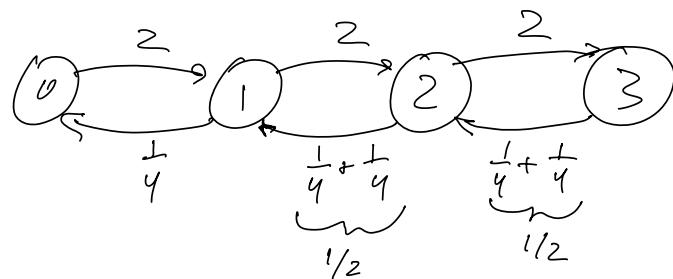
IRREDUCIBLE,
BUT NOT FINITE

Ex. M/M/2/3 QUEUING SYSTEM



$$X(t) = \frac{\# \text{ OF } \text{ CUST. } (\text{IN THE SYSTEM})}{\text{IS A CTMC}},$$

$$S = \{0, 1, 2, 3\}$$



$(\pi_0, \pi_1, \pi_2, \pi_3)$ UNIQUE STAT. DISTR,

$$\begin{cases} \pi_0 \cdot 2 = \pi_1 \cdot \frac{1}{4} & \Rightarrow \pi_1 = 8\pi_0 \\ \pi_1 \cdot 2 = \pi_2 \cdot \frac{1}{2} & \Rightarrow \pi_2 = 4\pi_1 = 32\pi_0 \\ \pi_2 \cdot 2 = \pi_3 \cdot \frac{1}{2} & \Rightarrow \pi_3 = 4\pi_2 = 128\pi_0 \\ \sum \pi_i = 1 & \Rightarrow \pi_0 + 8\pi_0 + 32\pi_0 + 128\pi_0 = 1 \Rightarrow \pi_0 = \frac{1}{169} \end{cases}$$

$$\pi = \left(\frac{1}{169}, \frac{8}{169}, \frac{32}{169}, \frac{128}{169} \right)$$

WHAT KIND OF QUANTITIES CAN WE COMPUTE?

CONSIDER THE ABOVE M/M/2/3 SYSTEM AS AN EXAMPLE.

π_0 = LONG-RUN FRAC. OF TIME WHEN THERE ARE NO COST.

$\pi_2 + \pi_3$ = —— | — | — | — BOTH SERVERS ARE BUSY

π_3 = — | — | — | — NO PLACE TO TAKE ADDL. COST.

$\lambda \pi_3$ = RATE (CUST/UNIT TIME) AT WHICH CUSTOMERS ARE BLOCKED (LOST)

$\lambda(1 - \pi_3) = \pi_1 \cdot \frac{1}{4} + \pi_2 \cdot \frac{1}{2} + \pi_3 \cdot \frac{1}{2}$ = RATE AT WHICH CUSTOMERS ARE ACCEPTED

= RATE AT WHICH CUSTOMERS LEAVE SYSTEM

= SYSTEM THROUGHPUT