

Preparation Activity PA07 – Context-Free Grammars (CFGs)

1. Consider the CFG G_1 of the lecture:

$$E \rightarrow I \mid E+E \mid E \times E \mid (E)$$

$$I \rightarrow a \mid b \mid Ia \mid Ib \mid Io \mid I1$$

- a) Why is this CFG ambiguous?
 - b) Show a sequence of leftmost derivations and the resulting syntax tree when the input string is $((a)+(a \times b))$
2. In order to remove the ambiguity of G_1 we must modify it, but we need to propose a modified CFG representing the same language of G_1 . Trying to remove the ambiguity, members of a team proposed different modified CFGs:
- a) One proposed the CFG G_2 below. Is G_2 ambiguous? Does G_2 represent the same language of G_1 ?
$$E \rightarrow I \mid E+I \mid E \times I \mid (E)$$
$$I \rightarrow a \mid b \mid Ia \mid Ib \mid Io \mid I1$$
 - b) Another one proposed the CFG G_3 below. Is G_3 ambiguous? Does G_3 represent the language of G_1 ?
$$E \rightarrow F \mid E+F \mid E \times F$$
$$F \rightarrow I \mid (E)$$
$$I \rightarrow a \mid b \mid Ia \mid Ib \mid Io \mid I1$$
 - c) Another one proposed the CFG G_4 below. G_4 is not ambiguous but seems to have a problem as it extends the language of G_1 . Indicate the problem.
$$E \rightarrow J \mid E \times J$$
$$J \rightarrow I \mid J+I$$
$$I \rightarrow a \mid b \mid Ia \mid Ib \mid Io \mid I1 \mid (E)$$
 - d) Another one proposed the CFG G_5 below. Besides the fact that G_5 is not ambiguous and represent the same language of G_1 , it also respects the priority of the operators (considering that the input strings represent arithmetic expressions, where the symbol $+$ identifies the addition and the symbol \times identifies the multiplication). Explain why.
$$E \rightarrow T \mid E+T$$
$$T \rightarrow F \mid T \times F$$
$$F \rightarrow I \mid (E)$$
$$I \rightarrow a \mid b \mid Ia \mid Ib \mid Io \mid I1$$