

EXERCISES ABOUT CONTEXT-FREE LANGUAGES (CFLS)

- 1 Draw a deterministic push-down automaton (DPDA) for the language:
 $\{0^n 1^m 0^n \mid n \geq 0 \text{ and } m \geq 0\}$.
- 2 Test the determinism in the automata of previous exercises. Verify if they are deterministic automata and if they are not, identify transitions which make them non-deterministic.
- 3 “Given a non-deterministic PDA, there exists always an equivalent deterministic PDA, i.e., which recognizes exactly the same language.” Is this sentence true? Justify.
- 4 Consider the CFG below. Convert it to the Chomsky Normal Form (CNF), showing the intermediate simplification steps.

$$\begin{aligned} S &\rightarrow A1B \\ A &\rightarrow 0A \mid \varepsilon \\ B &\rightarrow 0B \mid 1B \mid \varepsilon \end{aligned}$$
- 5 Convert the following CFG into the Chomsky Normal Form, showing intermediate steps.

$$\begin{aligned} S &\rightarrow 0S00 \mid 0B0 \mid B \\ B &\rightarrow 11B22 \mid 12 \mid C \\ C &\rightarrow 0 \mid \varepsilon \end{aligned}$$
- 6 Show, using the pumping lemma for context-free languages, that the language of the strings $a^n b^n c^i$, in which $n \leq i \leq 2n$, is not a context-free language.
- 7 Show, using the pumping lemma for context-free languages, that the language of the strings 0^p , with p a prime number, is not a context-free language. [Note: we have already proven that this language does not satisfy the pumping lemma for regular languages; let us use a similar strategy to prove that this language does not satisfy the pumping lemma for context-free languages.]
- 8 Show, using the pumping lemma for context-free languages, that the language of the strings $0^i 1^j$, in which $j=i^2$, is not a context-free language.
- 9 When we try to apply the pumping lemma for context-free languages to a language of this family, trying to show that it does not belong to the context-free languages, it happens that the “adversary wins” always and we cannot complete the prove. Illustrate this process with the languages $L1 = \{00, 11\}$ and $L2 = \{0^n 1^n, n \geq 1\}$.
- 10 Given two string w and x , let us call $inter(w,x)$ to the set of strings obtained interchanging symbols of w and x by the order they occur in w and in x . We can extend the operation to two languages $L1$ and $L2$, naming $inter(L1, L2)$ to the union, for all the pairs of strings w from $L1$ and x from $L2$, of $inter(w,x)$.
 - a) Determine the value of $inter(00,111)$.
 - b) Determine the value of $inter(L1, L2)$ with $L1 = L(0^*)$ and $L2 = \{0^n 1^n, n \geq 0\}$.
 - c) Show that if $L1$ and $L2$ are both regular languages then $inter(L1, L2)$ is a regular language, as well. [Suggestion: consider the DFAs of $L1$ and $L2$.]
 - d) Show that if L is a context-free language and R a regular language then $inter(L, R)$ is a context-free language. [Suggestion: consider a PDA for L and a DFA for R .]
- 11 Show that in a grammar in the Chomsky Normal Form (CNF), all the analysis trees of length n have $2n-1$ inner nodes.