

LLVM IR

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LLVM IR

- LLVM Intermediate Representation
- Better than assembly code
 - Independent of the architecture
 - No headaches reading the x86 code
 - Proper optimizations already implemented
 - Type safety
 - More human readable

LLVM IR Types

- We will discuss only a few of the LLVM IR types, full documentation can be found [here](#).
 - `void` Does not represent a value and has the size 0.
 - `iN` An integer that is N bits long.
 - N can be anywhere from 1 bit to $2^{23} - 1$.
 - `i1` A 1 bit integer
 - `i8` An 8 bit integer
 - `i32` A 32 bit integer
 - `i1000` A 1000 bit integer.
 - Is it unsigned?

LLVM IR Types

- Floating point types.
 - `half` A 16 bit floating point .
 - `float` A 32 bit floating point.
 - `double` A 64 bit floating point.
 - More specific floating point types [here](#).
- Array type [`<count>` x `<type>`] (and yes that is an 'x')
 - [`14` x `i32`] An array of 14 elements of type `i32`.
 - [`1` x `float`] An array of one element of type `float`.

LLVM IR Types

- Function types
 - `i32 (i32)` A function that takes a `i32` and returns a `i32`.
 - `double (float, ...)` A vararg function that takes at least a floating point and returns a `double`.
 - Further details can be found [here](#).
- Pointer types
 - `i32*` A pointer to a 32 bit integer.
 - `i32 (i32)*` A pointer to a function that takes a `i32` and returns a `i32`.
 - `[14 x i32]*` A pointer to an array of 14 elements of type `i32`.
- Label type, represents a code label.
 - `here:` Label at the code.
 - Labels must be placed after a [terminator instruction](#) or at the starting of a block(will be explained further later).

LLVM IR Identifiers

- There are two kinds of identifiers
 - Local identifier names start with %
 - Are only available in the scope.
 - Global identifier names start with @
 - Can be functions or global variables
- Identifiers help handle stack.
 - No need to implement stack.
- Identifiers should fit in the following regex.
 - `[%@] [-a-zA-Z$. _] [-a-zA-Z$. _0-9]*`
- Unnamed value are represented by unsigned integer names.
 - `%1`, `@2`

LLVM IR Functions

- Functions can be defined and declared.
 - The definition specifies what the function does.
 - The declaration specifies what the function signature is.
- A function definition is made of basic execution blocks.
 - These blocks define the control flow graph.
 - Each block may start with a label.**
- Simple function definition
 - `define <type> @<name> (<type> <name>, ...)`
 {
 }
 - Argument list can be empty.
- Read further [here](#).

LLVM IR Instructions

- Again as with the case of the types, we are not going to discuss all the LLVM IR instructions, a full documentation can be found [here](#).
 - `ret <type> <value>` Return a value from inside a function.
 - `ret void` Return for void return type functions.
 - `ret i32 0` Return zero in a function that returns 32 bit integers.
 - `br i1 <cond>, label <iftrue>, label <iffalse>` Conditional branch instruction.
`br label <dest>` Unconditional branch instruction.
 - Infinite loop:
`here:`
`br label %here`
 - Infinite loop (since the condition is a constant 1, it is always true, and branches to `%here`):
`here:`
`br i1 1, label %here, label %there`
`there:`

LLVM IR Instructions

- Read further on flow control instructions [here](#).
- Binary operations
 - Require two operands of the same type.
 - `<result> = add <type> <op1>, <op2>` Adds two operands and the yields the result with type ty (for integers).
 - `<result> = sub <type> <op1>, <op2>` Subtracts op1 from op2 (for integers).
 - `<result> = mul <type> <op1>, <op2>` Multiplies op1 and op2 (for integers).
 - `<result> = udiv <type> <op1>, <op2>` Divedes op1 by op2 (for unsigned integer).
 - `<result> = sdiv <type> <op1>, <op2>` Divedes op1 by op2 (for signed integer).
 - Read further [here](#).
- Memory access instruction
 - You can read further [here](#).

LLVM IR Calling a Functions

- Call instructions calls a function
 - `%<name> = call @<name> (<name>, ...)`
 - Arguments are passed inside the parentheses.
 - Read further on call instruction [here](#).
- Compare instructions
 - These instructions compare two values and yield `i1`.
 - `%<result> = icmp <cond> <type> <op1>, <op2>` Compares two operands with type, `<type>`.
 - The compare condition is specified by `<cond>`.
 - The result is 1 on condition being correct, and 0 on condition being false.
 - Read more [here](#).

Entry Point

- The execution of the program starts from the entry point.
- Each executable should have an entry point.
- Main function is the entry point on LLVM IR.

- `define i32 @main() {
 ret i32 0
}`