Report

Laboratory Work 4

Dmitry Ladutsko

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Task 1: Auto Trace configuration training

3.7-	A		
Nº	Auto Trace Configuration Options	Expected Results	Description
	set autotrace off	No AUTOTRACE report is generated.	This is the default.
	set autotrace on	The AUTOTRACE report	Includes both the optimizer execution path and the SQL statement execution statistics.
	set autotrace traceonly	Like SET AUTOTRACE ON	Suppresses the printing of the user's query output, if any. This option is useful when you are tuning a large query, but do not want to see the query report.
	set autotrace on explain	The AUTOTRACE report	Shows only the optimizer execution path.
	set autotrace on statistics	The AUTOTRACE report	Shows only the SQL statement execution statistics.
	set autotrace on explain statistics	Shows only the SQL statement execution statistics.	= set autotrace on
	set autotrace traceonly explain	Execution plan w/o query result	
	set autotrace traceonly statistics	Statistics w/o query result	These options should be used when a large result set is expected.
	set autotrace traceonly explain statistics	Execution plan + statistics w/o query result	
	set autotrace off explain		
	set autotrace off statistics set autotrace	Disables aut	cotrace utility
	off explain statistics		

Task 2: Nested Loop Joins

Example:

SELECT /*+ gather_plan_statistics */ *

FROM scott.emp e, scott.dept d

WHERE e.deptno = d.deptno

AND d.deptno = 10

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Picture 2.1 - Query Result

Note. I used /*+ gather_plan_statistics */ hint because for some reason explain plan statement did not give me full needed stats (only id, operation and name columns). Nested loops join compare every row of table B with current row in table A. If it is a row, which satisfies the specified condition, it is selected, otherwise goes next and compare second raw from table A with all rows from table B. This method is quite useful with small table, but not with big.

Task 3: Sort-Merge Joins

E.DEPTNO=10

```
set autotrace on explain;

SELECT /*+ use_merge(e d) gather_plan_statistics */ *

FROM scott.emp e, scott.dept d

WHERE e.deptno = d.deptno
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Picture 3.1 - Query Result

Note. Using this type of join we need to be sure that all used tables are sorted by columns, mentioned in query condition. Oracle scans each row once. This gives us nice benefits, but we should keep in mind that sorting loads the system a lot if we use big tables.

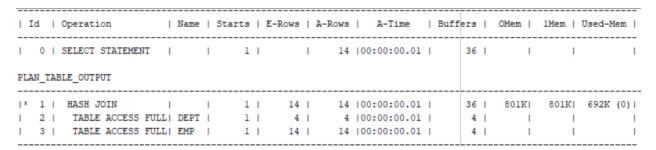
Task 4: Hash Joins

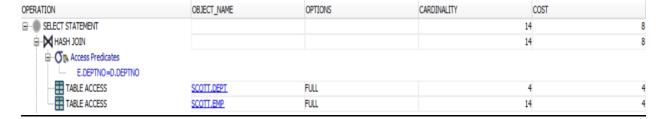
set autotrace on explain;

SELECT /*+ gather_plan_statistics USE_HASH(e d) */ *

FROM scott.emp e, scott.dept d

WHERE e.deptno = d.deptno





Picture 4.1 - Query Result

Note. Hash join is used when projections of the joined tables are not sorted on the join columns. In this case, the optimizer builds an in-memory hash table on the inner table's join column. The optimizer then scans the outer table for matches to the hash table, and joins data from the two tables accordingly.

Task 5: Cartesian Joins

set autotrace on explain;

SELECT /*+ gather_plan_statistics */ *

FROM scott.emp, scott.dept

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Picture 5.1 - Query Result

Note. Cartesian joins are used when all rows in all tables listed in a query: each row in the first table is paired with all the rows in the second table. This happens when there is no relationship defined between the two tables.

Task 6: Left/Right Outer Joins

1. Left outer JOIN

set autotrace on explain;

select /*+ gather_plan_statistics */ e.ename , e.deptno, e.job, d.dname

from scott.emp e, scott.dept d

where e.deptno = d.deptno(+);

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Picture 6.1 - Query Result

Note. The specific of this type of join is that result includes unmatched rows from only the table that is specified before the LEFT OUTER JOIN clause. If we are joining two tables and want the result set to include unmatched rows from only one table, we use a LEFT OUTER JOIN clause (or a RIGHT OUTER JOIN) clause.

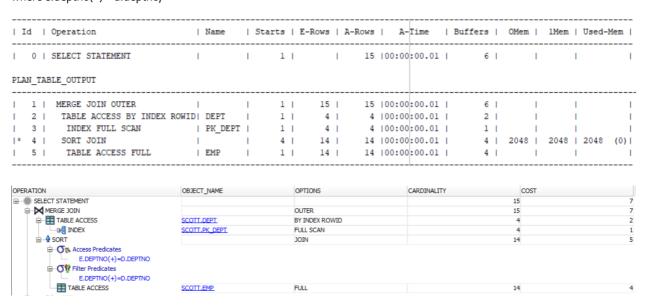
2. Right outer JOIN

set autotrace on explain;

select /*+ gather_plan_statistics */ e.ename , e.deptno, e.job, d.dname

from scott.emp e, scott.dept d

where e.deptno(+) = d.deptno;



Picture 6.2 - Result

Task 7: Full Outer Join

set autotrace on explain;

select /*+ gather_plan_statistics */ e.ename, e.deptno, e.job, d.dname

from scott.emp e

scott.dept d

full outer join

on(e.deptno = d.deptno);

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Picture 7.1 - Result

Note. We use Full outer join if we want the result set to include unmatched rows from both tables.

Task 8: Semi Joins

1. Using /*semijoin*/:

set autotrace on explain;

select /*+ gather_plan_statistics semijoin */ dname

from scott.dept d

where deptno in (select deptno from scott.emp e);

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Picture 8.1 - Query Result

Note. A semi-join returns one copy of each row in first table for which at least one match is found. Semi-joins are also written using the EXISTS construct.

2. Using /*no semijoin*/:

set autotrace on explain;

select /*+ gather_plan_statistics no_semijoin */ dname

from scott.dept d

where deptno in (select deptno from scott.emp e);

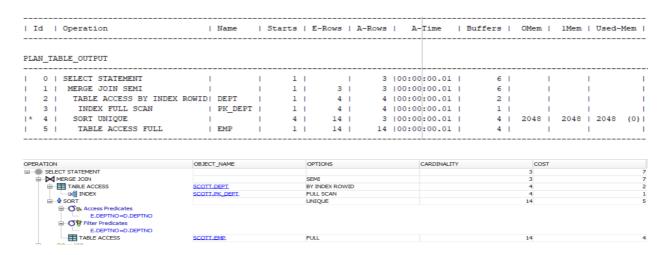
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Picture 8.2 -Query Result

Note. As a consequence, NO vital difference

3. Using /*using exists*/

set autotrace on explain;
select /*+ gather_plan_statistics */ dname
from scott.dept d
where exists (select null from scott.emp e
where e.deptno = d.deptno);



Picture 8.3 -Query Result

Note. The EXISTS function checks to find a single matching row to return the result in a subquery. Because the IN function retrieves and checks all rows, so it is slower.

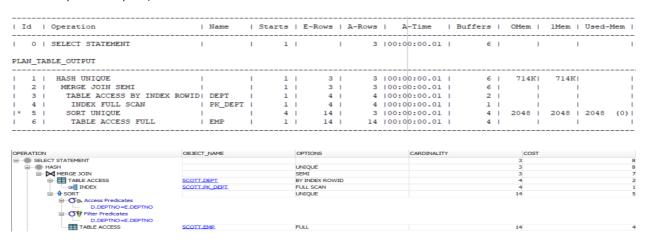
4. Select distinct using /* inner join with distinct */ hint

set autotrace on explain;

select /*+ gather_plan_statistics inner join with distinct */ distinct dname

from scott.dept d, scott.emp e

where d.deptno = e.deptno;



Picture 8.4 - Query Result

Note. The DISTINCT clause is used in a SELECT statement to filter duplicate rows in the result set. This way we can ensure that returned rows are unique for the column or columns specified in the SELECT clause.

Task 9: Anti Joins

1. Not exists using /*not exists*/ hint

set autotrace on explain;

select /*+ gather_plan_statistics not exists ANTIJOIN */ dname

from scott.dept d

where not exists (select null from scott.emp e

where e.deptno = d.deptno);

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Note. The NOT EXISTS operator returns true if the subquery returns no row. Otherwise, it returns false. We should keep in mind that NOT EXISTS operator returns false if the subquery returns any rows with a NULL value.

2. Not in using /*not in*/ hint

set autotrace on explain;

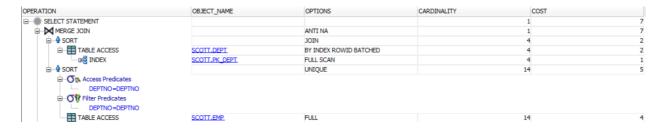
select /*+ gather_plan_statistics not in */ dname

 $from\ scott.dept\ d$

where deptno not in

(select deptno from scott.emp e)

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Note. Using NOT IN operator can be risky because it does not return null values

NOT IN can be just as efficient as NOT EXISTS even if an anti-join can be used (if the subquery is known to not return nulls)

3. Minus using /*minus*/ hint

set autotrace on explain;

select /*+ gather_plan_statistics minus*/ dname

from scott.dept

where deptno in

(select deptno from scott.dept minus

select deptno from scott.dept);



Note. The Oracle MINUS operator is used to return all rows in the first SELECT statement that are not returned by the second SELECT statement. Each SELECT statement will define a dataset. The MINUS operator retrieve all records from the first dataset and then remove from the results all records from the second dataset.

4. Left outer anti join

set autotrace on explain;
select /*+ gather_plan_statistics left outer */ dname
from scott.dept d, scott.emp e
where d.deptno = e.empno(+)

and e.deptno is null;

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Picture 9.1 -Query Result

Task 10: Summary table

Join Access "A"	Join Access "B"	Nested	Hash Join	Sort-	Anti-	Semi-	
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Small Table	Small Table(Indexed)	Good	Ineffective	menective			
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Laboratory work summary:

We used many join methods to see which has better performance in different situations.

At the table above and insights from each task (or sub – task) you can see some results and conclusion.

