

Sorting- useful for a zillion purposes

In previous experiences:

"slow sorts" - selection sort
insertion sort
bubble sort

$$\rightarrow O(n^2)$$

$$1 + 2 + 3 + \dots + \underbrace{n-1}_{ish} \rightarrow O(n^2)$$

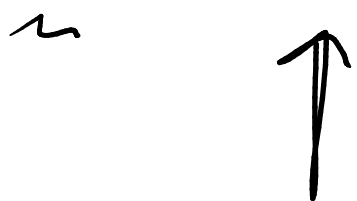
In this class

mergesort
quicksort
heapsort

} all are
 $O(n \log n)$
[with details,
caveats, etc]

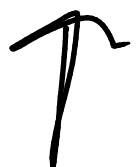
$$O(n^2)$$

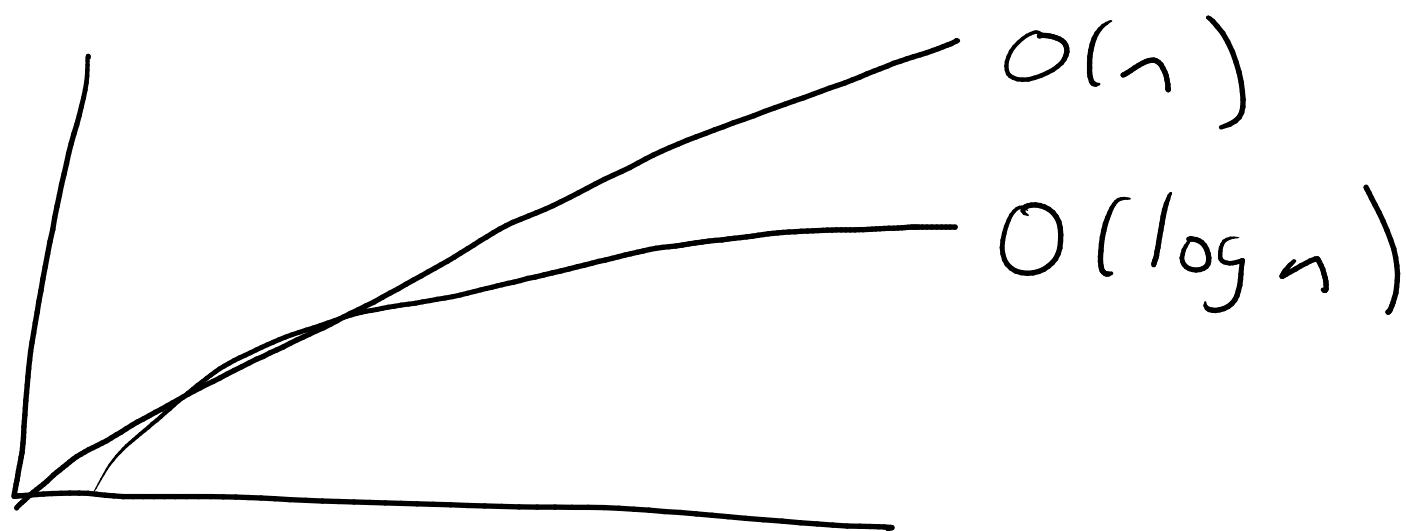
$$O(n * n)$$



$$O(n \log n)$$

$$O(n * \log n)$$





Mergesort is recursive!

- mergesort left half
- mergesort right half
- merge two halves together

6, 2, 4, 3, 9, 8, 1, 7

6 2 4 3

9 8 1 7

6 2 4 3

9 8 1 7

6 2 4 3

9 8 1 7

$\underbrace{2 \quad 6} \quad \underbrace{3 \quad 4} \quad \underbrace{8 \quad 9} \quad \underbrace{1 \quad 7}$
 $\diagdown \quad \diagup \quad \diagdown \quad \diagup$

$\underbrace{2 \quad 3 \quad 4 \quad 6} \quad \underbrace{1 \quad 7 \quad 8 \quad 9}$

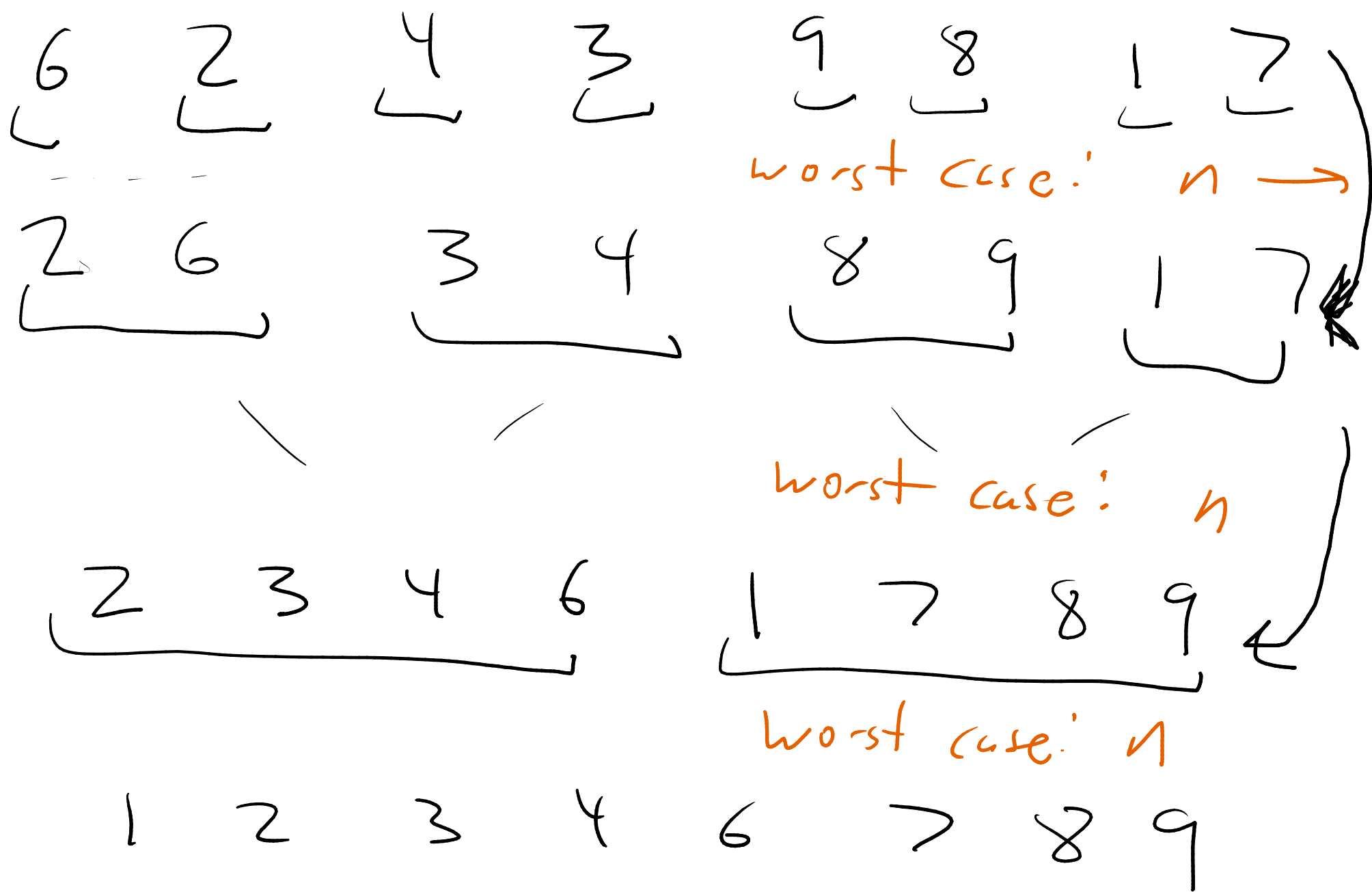
1 2 3 4 6 7 8 9

How much work?

Splitting takes essentially no work at all, because all we really do is accounting

- "mergesort list indices $0 \rightarrow \text{half}$ "
- "mergesort list indices $\text{half} + 1 \rightarrow \text{end}$ "

Work is in the merging.

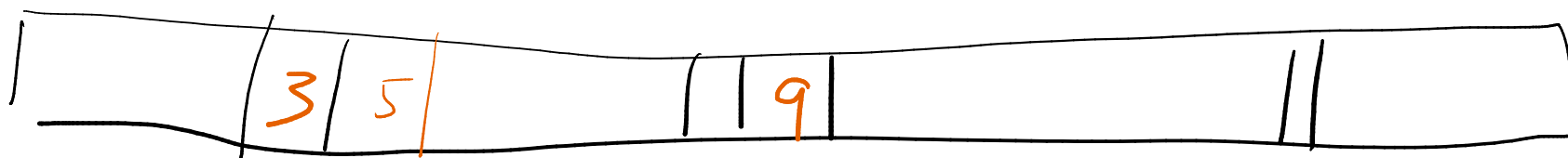


How many numbers were rewritten?

Performance: $n * \text{\# of rows above}$
 $\log n$

$$n * \log n$$

list



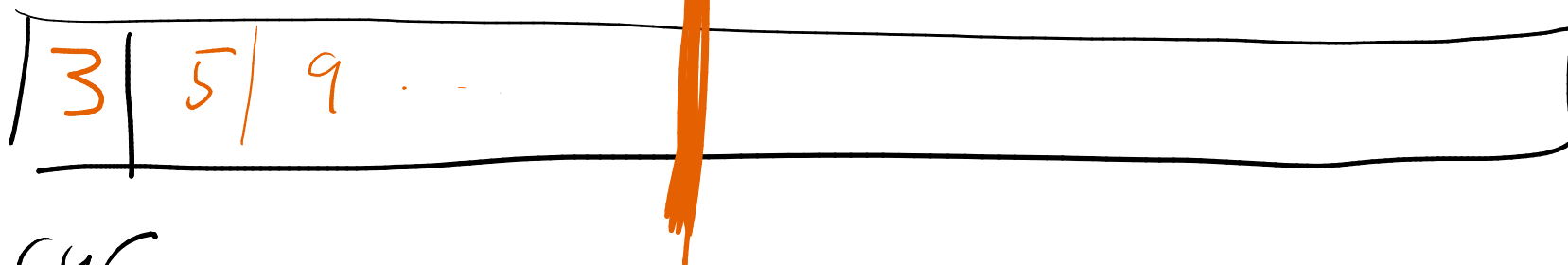
leftStart
left →

left right
End Start

right →

right
End

temp



cur

0



started at 0 in temp list, whereas
leftStart in original list
was elsewhere