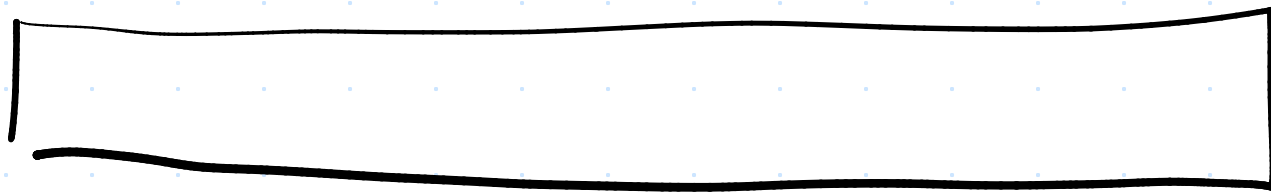


# Queues

- First In, First Out (FIFO)

# Stacks

- Last In, First Out (LIFO)



front  
remove  
(pop)  
dequeue



rear  
add  
(push)  
enqueue

Implement strategies

- array
- list
- linked chain



With a list, same as a stack  
(almost)

0	1	2	3	4
2	9	3	7	5

enqueue

2  
9  
3  
7  
5

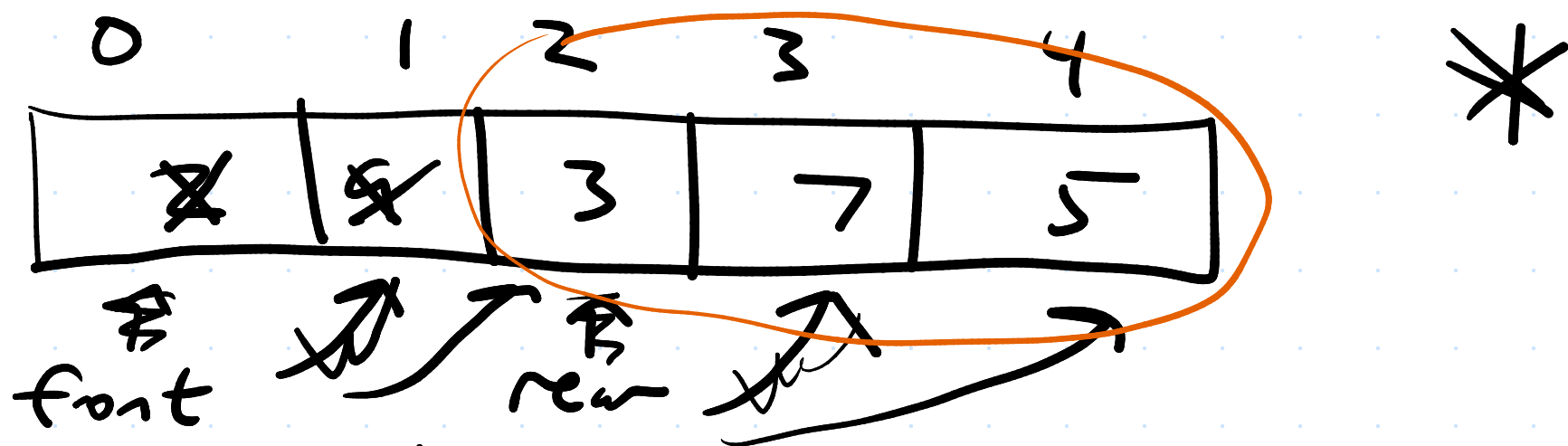
Naive approach

- add on end  
list.add  $O(1)$

- remove from  
location 0

list.remove(0)  
            $O(n)$

# Smarter approach: "snake"



enqueue(2)

enqueue(9)

enqueue(3)

enqueue(7)

dequeue()  $\Rightarrow$  2

enqueue(5)

dequeue()  $\Rightarrow$  9

front

~~X~~

X

2

rear

~~X~~

~~X~~

4

$O(1)$  for everything

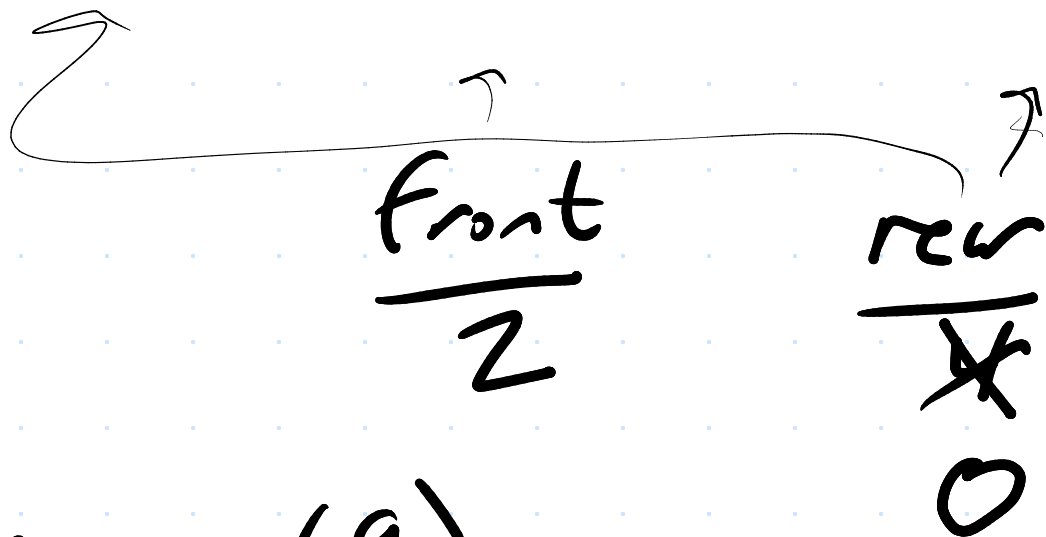
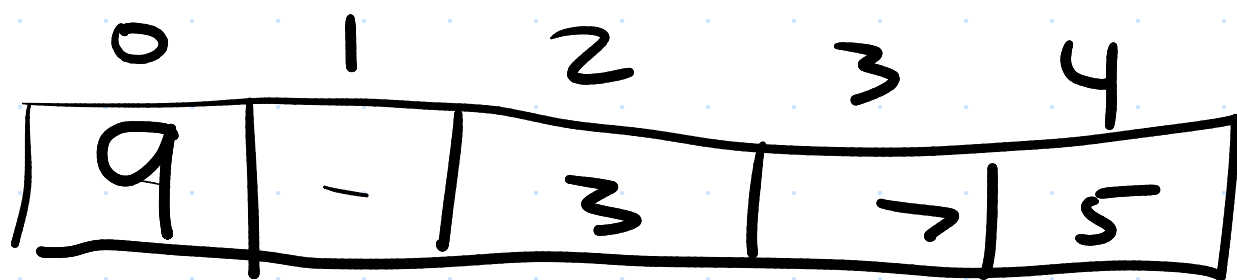
Empty memory  
at front just  
piles up

-eventually eat up all the  
memory you have

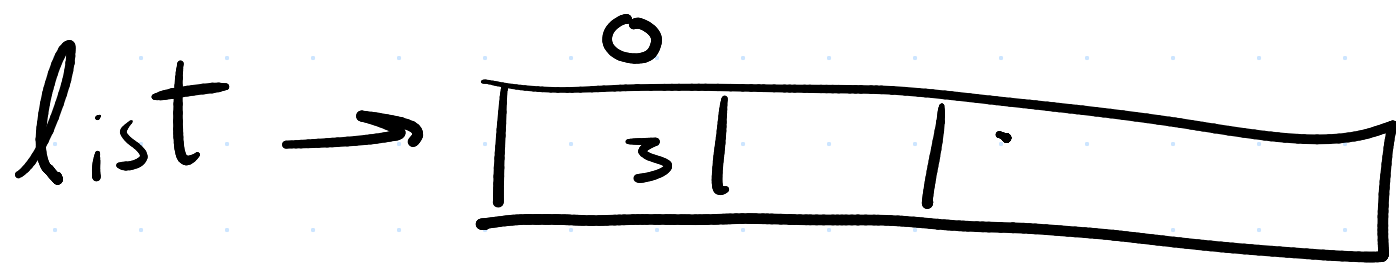
Fix: slide everything down  
? every now and then  
... ok ...  $O(n)$  when it  
happens

---

FIX! Wrap around the  
end of the list



enqueue(9)



one item

front = 0

rear = 0

} front and  
rear  
same  
means

one item

In this simple

version,

front = -1

rear = -1

empty queue

For adding 2 possibilities

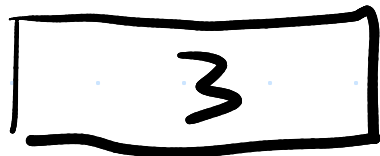


front

rear

room in  
list

list[rear] =  
item

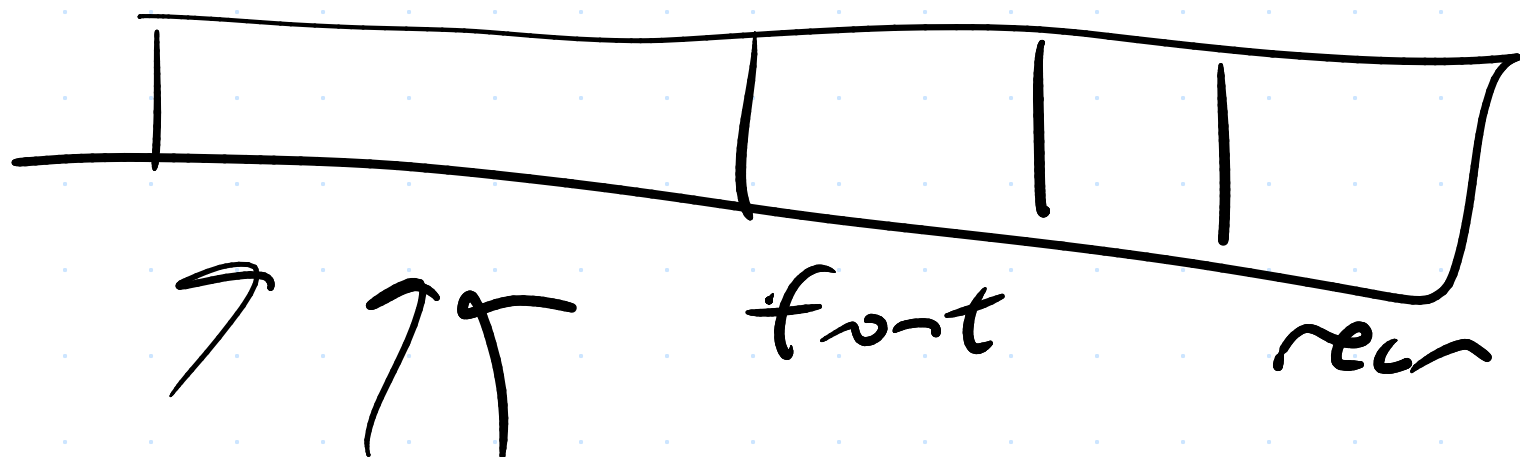


front

past  
end of  
list  
rear

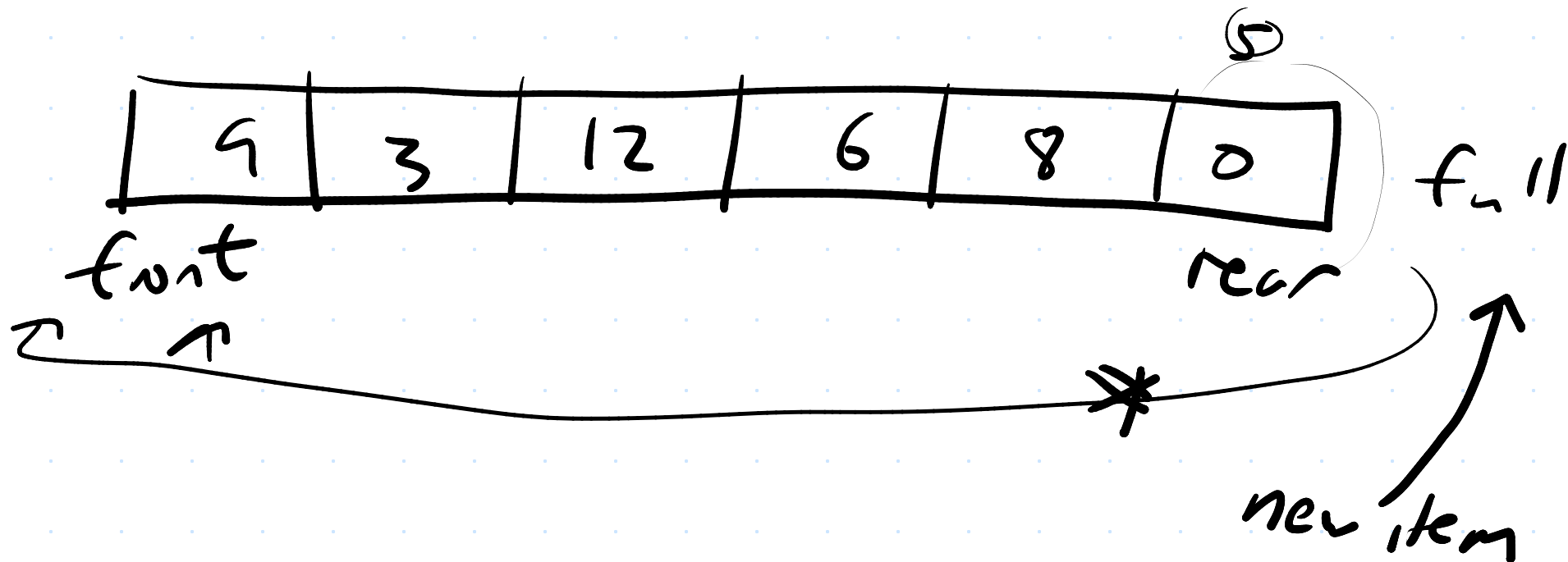
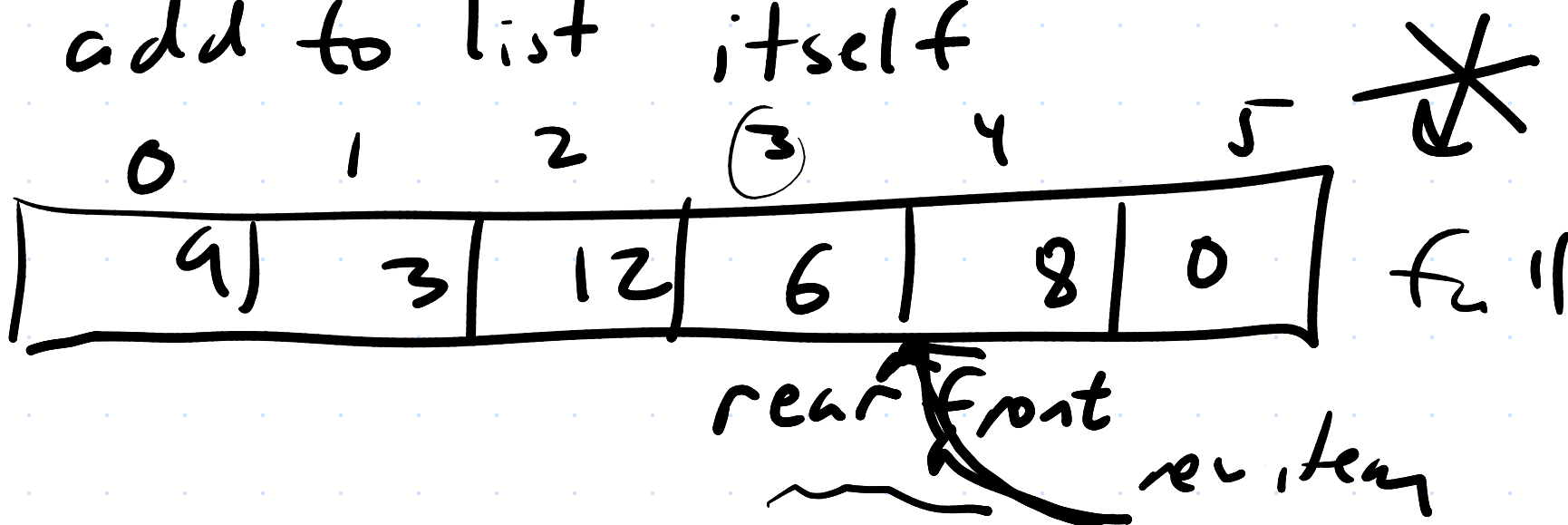
list.add(item)

Fix enqueue to instead wrap around end of list if you are at the end of list and there is room at beginning



What is condition for a full list?

If list is completely full, add to list itself



have room

