

CS Town Hall: Thurs (tomorrow),  
3:30 PM, AND 329

Today: - end of term scheduling  
- last assignment / letrec  
- types

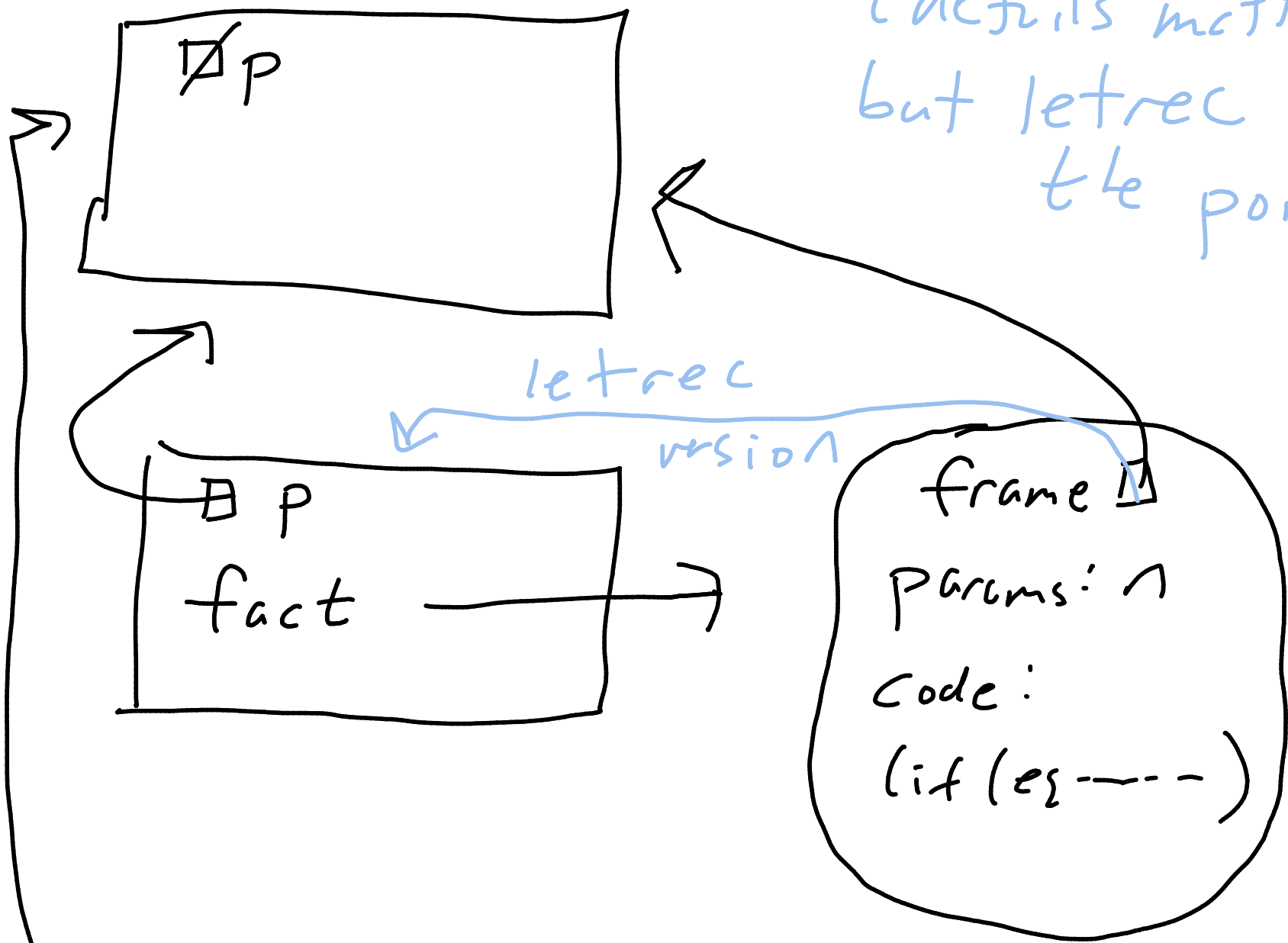
```
;; Try the above with a recursive function
(let ((fact (lambda (n)
               (if (equal? n 1)
                   1
                   (* n (fact (- n 1)))))))
  (fact 3))
```

closure points to global frame

↓  
unable to find fact in recursive calls.

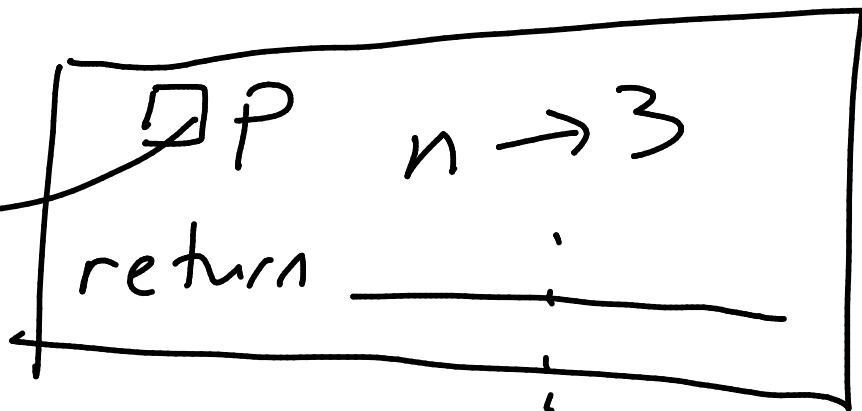
global frame

broadly  
(details matter),  
but letrec moves  
the pointer



(fact 3)

apply closure



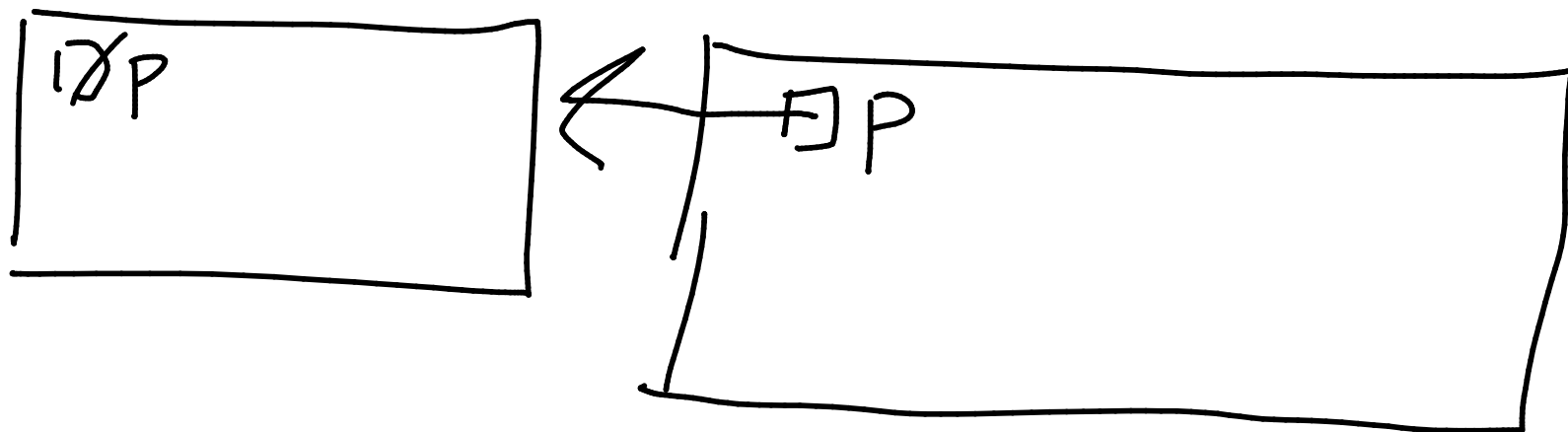
( \* 3 ( fact ( - 3 1 ) ) )

Can't find fact

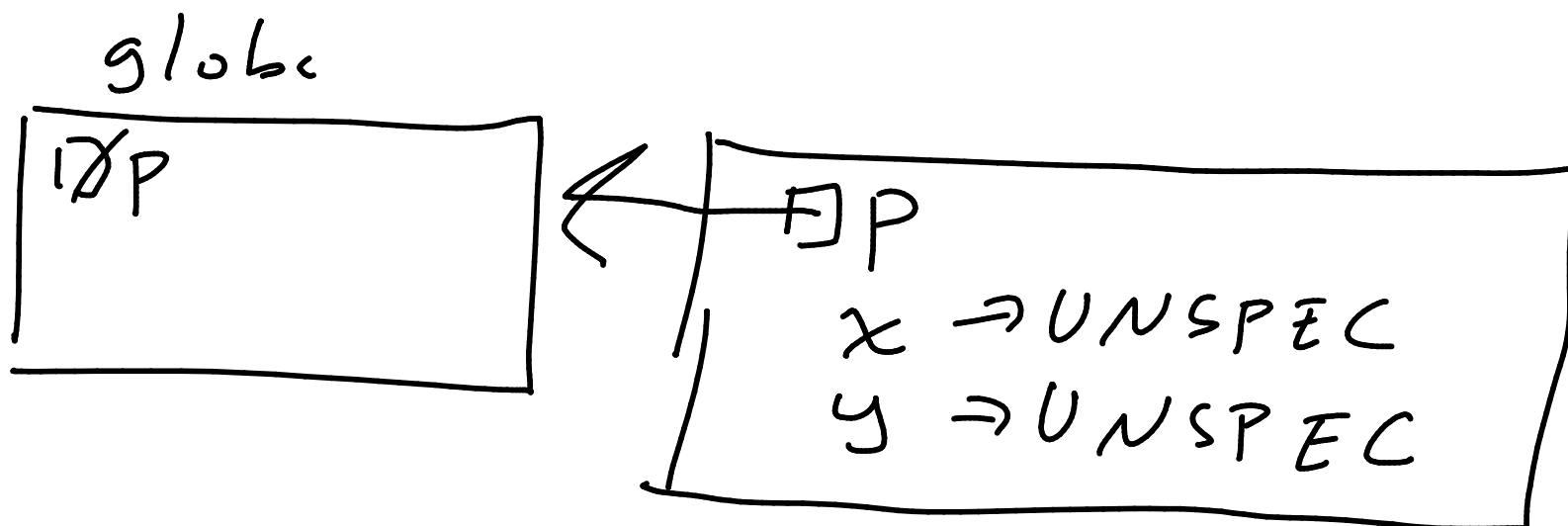
What does letrec actually do?

(let ((x 3) (y 5))  
 (+ x y))

(1) Creates a new frame as usual



(2) Add the vars to the new frame,  
but assign them to  
UNSPECIFIED\_TYPE

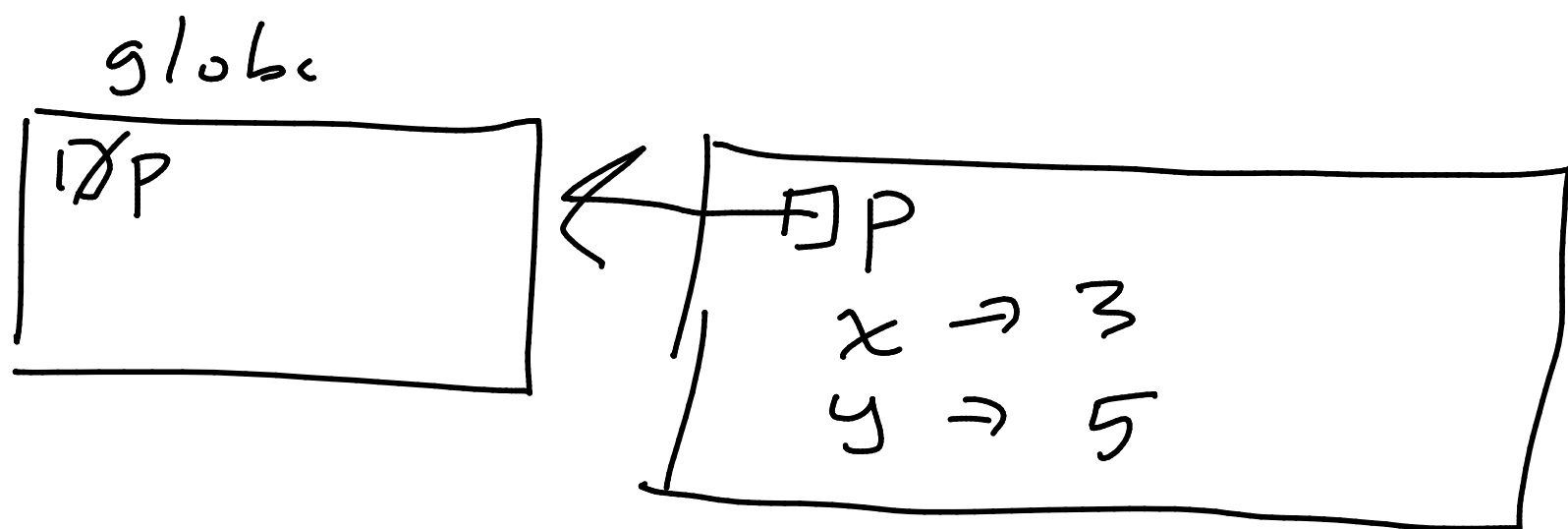


③ Evaluate the right sides of each pair of bindings, and store somewhere temp. <sup>in the new frame</sup> If anything uses an UNSPEC, throw an error.

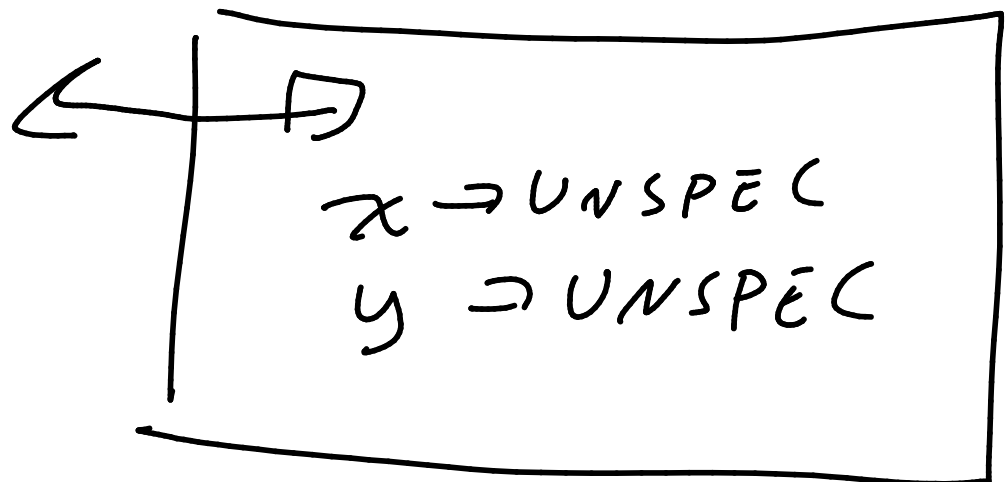
eval 3  $\rightarrow$  3 remember

eval 5  $\rightarrow$  5 remember.

No errors? Great! Now write these into new frame.



$(\text{letrec } (x \ 3) \ (y \ (+ \ x \ 1))$   
 $(+ \ x \ y))$



eval 3  $\rightarrow$  3

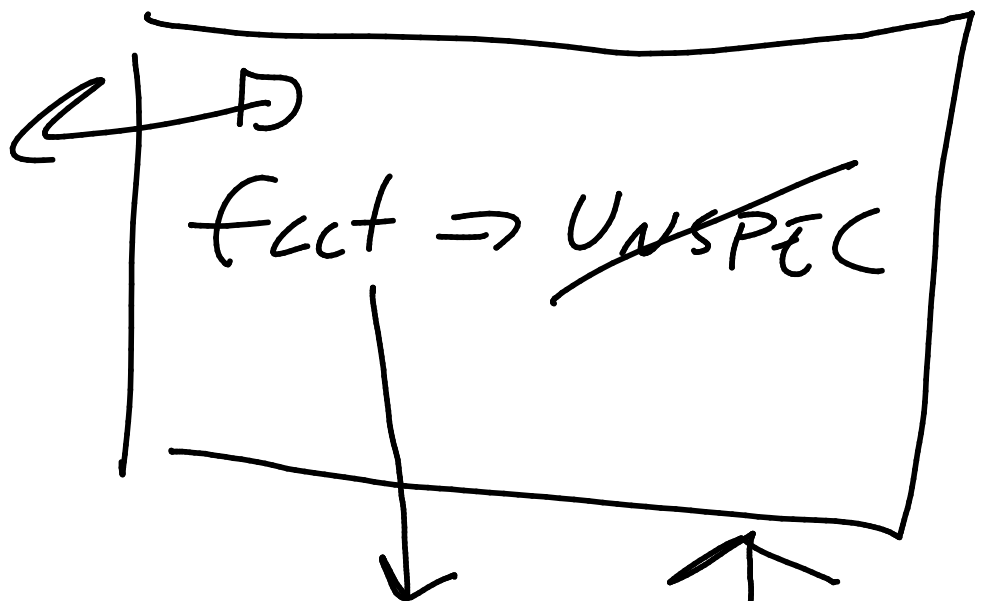
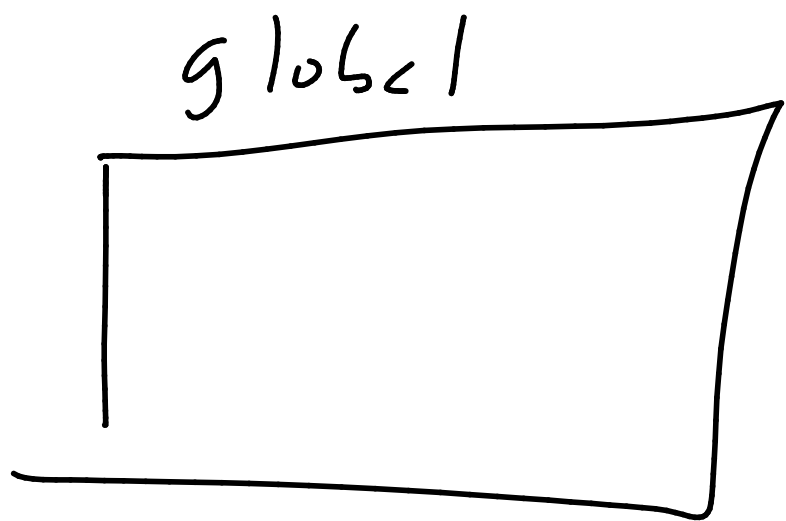
eval  $(+ \ x \ 1)$

oh no,  $x$  is  
UNSPEC

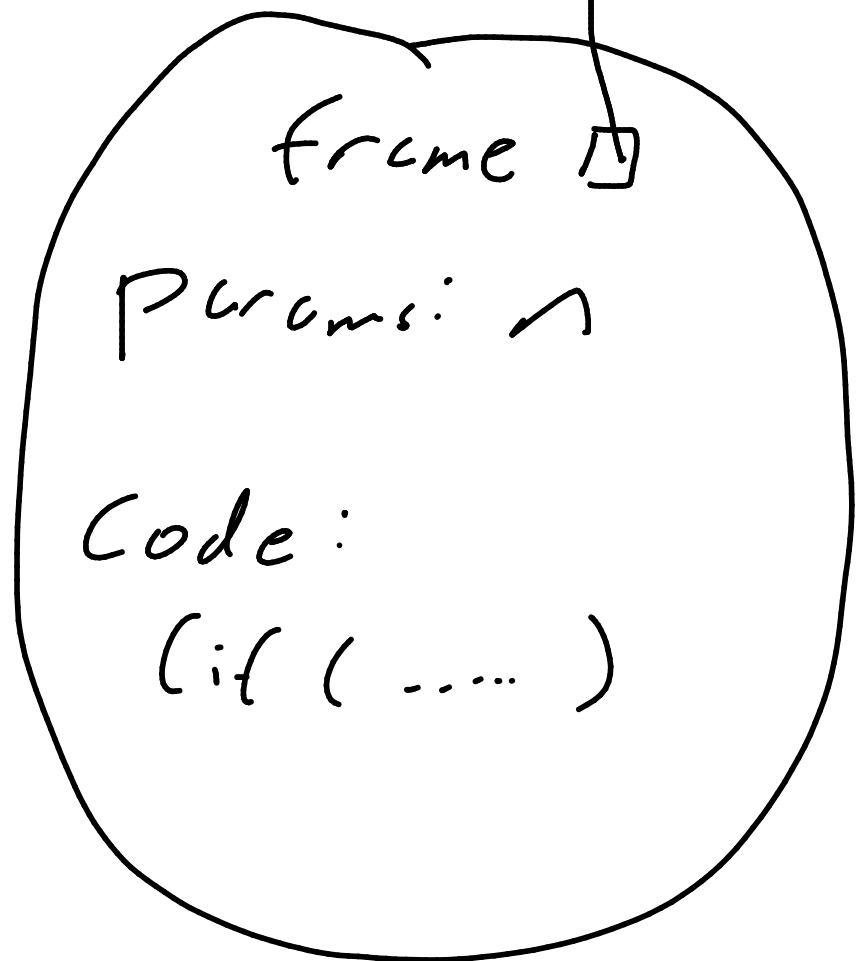
$\rightarrow$  error.

try the above with a recursive function

```
(let ((fact (lambda (n)
  letrec
    (if (equal? n 1)
      1
      (* n (fact (- n 1))))))
  (fact 3))
```



eval the lambda



Test 61, scm (E-tests)  
due to Donald Knuth

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Way, way back he wrote a letter  
to the editor of a PL journal  
and said (in different words)

- all these ALGOL interpreters stink.  
They don't do eval right.

Here's a test I wrote to  
distinguish the good ones  
from the bad ones.

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One specific example Guile gets  
wrong(!)

(letrec ((x 3) (y x))  
x)  
not in frame