

Parsing: LL vs LR

Parsing assignment

Scoping

Parsing?

Generally, two classes of  
Parsing algorithms

① (We're not doing for assignment)

$\langle B \rangle ::= 0 \langle B \rangle \mid 1 \langle B \rangle \mid \epsilon$

← start

← nothing

Parse 01011

$\langle B \rangle$

0  $\langle B \rangle$

1  $\langle B \rangle$

⋮

building this

top-down,

choose which

rule based

on which

token is

up next

traverse

grammar left

↑ to right


LL algorithm

↓  
left to right (read program)

(2) Approach we're using on our assignment

$$(+ (* 3 4) 6)$$

stack



+

*[Handwritten signature]*

✱

3/

4

6

A hand-drawn diagram representing a list structure. It consists of a large outer rectangle. Inside this rectangle, from left to right, are: a '+' sign, a smaller rectangle containing a '\*' sign, the number '3', and the number '4', and finally the number '6'.

Combined  
later

CS

+

457

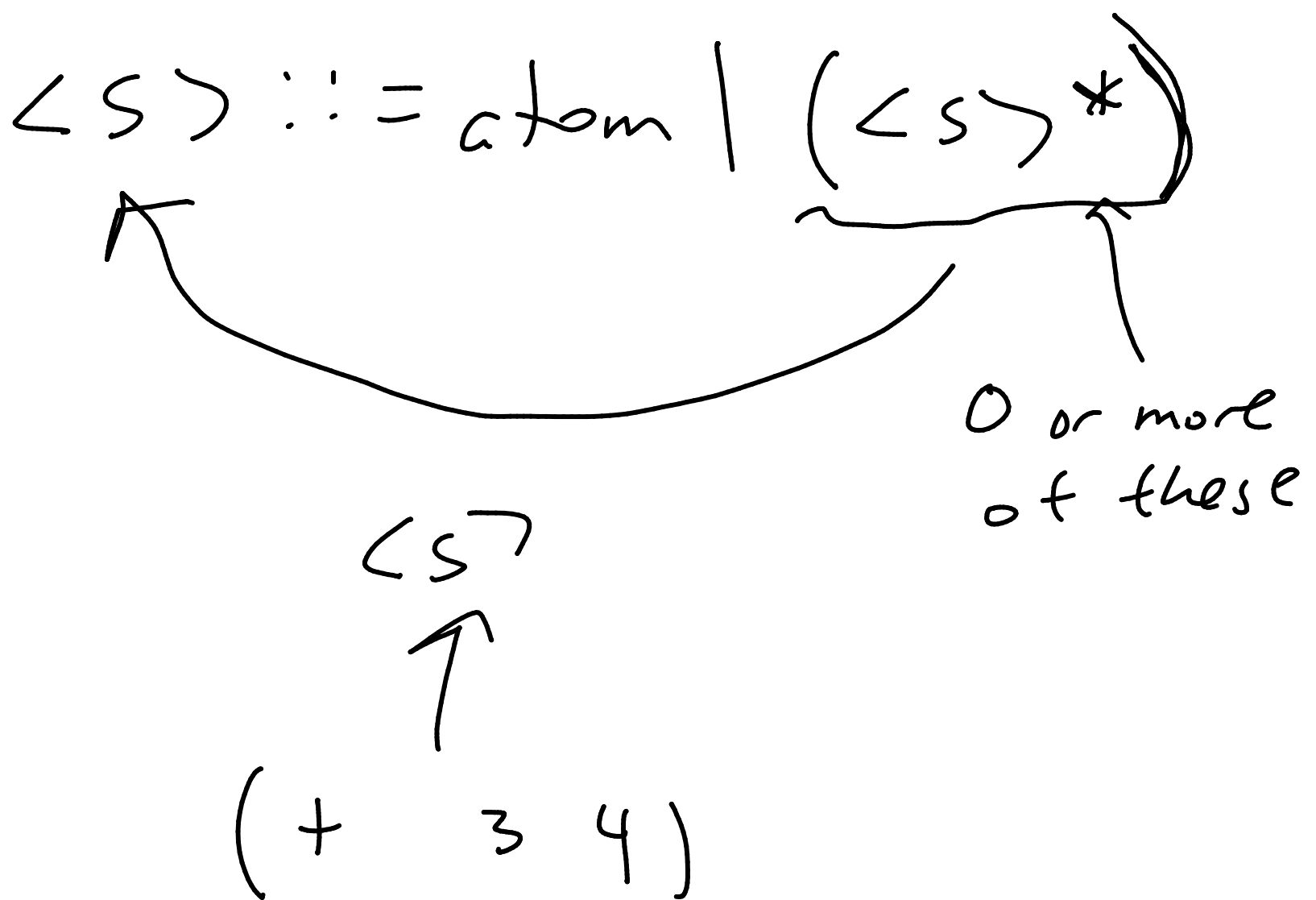
457

16

first

# LR algorithm

↳ traversing grammar from right to left



Pros and cons of each

LR is generally applicable for more languages than LL, but not universally

LL algorithms can (not always) be faster

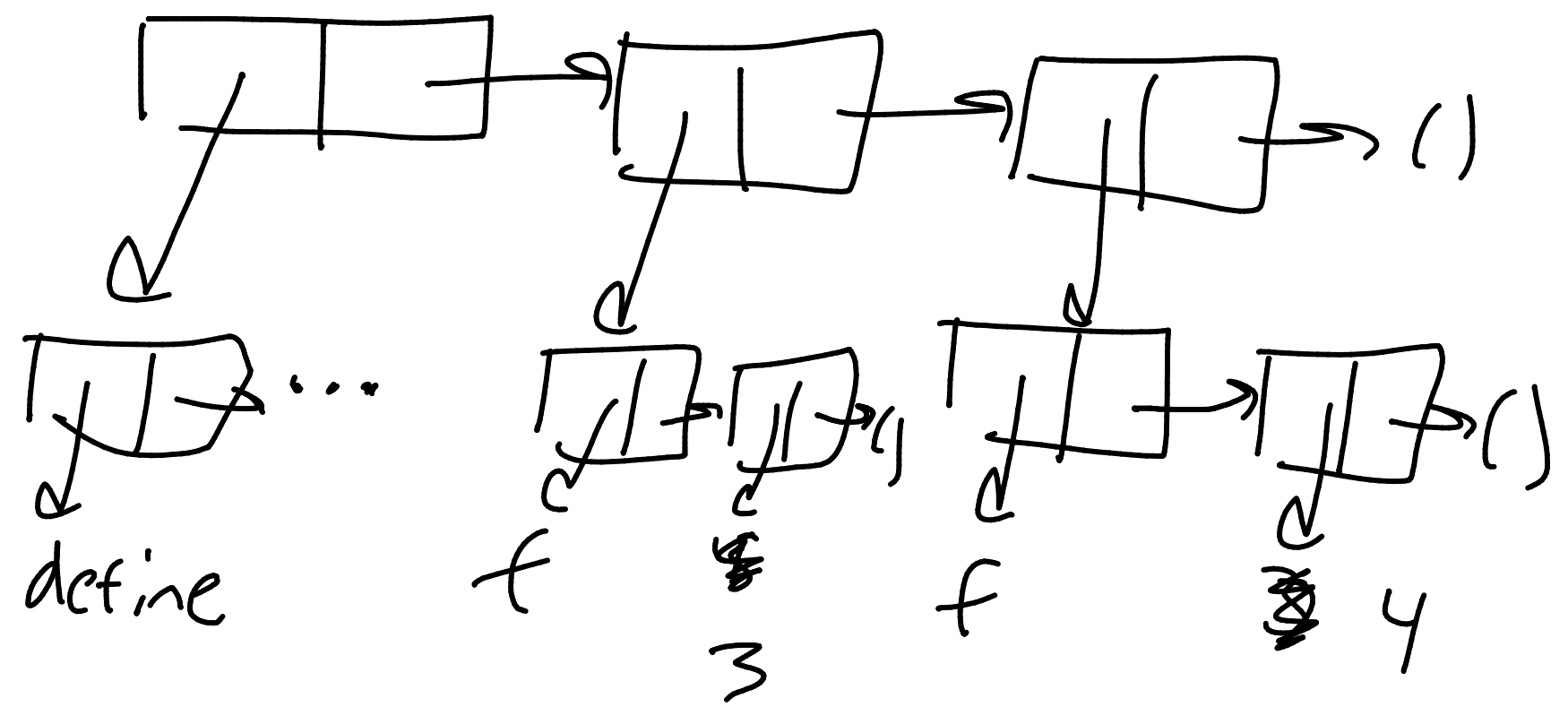
Parsing assignment: parse your tokens!

Uses the Scheme parsing algorithm we've been looking at

Scheme is inconsistent at the very top level of your program

(define f (lambda (x) (+ x 1)))  
(f 3)  
(f 4)  
)  
← these don't exist

A Scheme program is a sequence of possible Scheme expressions  
- you need to have a top level list of these



Single quote.  $\rightarrow$  '(a b)

Also breaks the rules.

Really just an abbreviation for

(quote (a b))

*transform*

"syntactic sugar"

# Scoping

When resolving a variable, how do you find it?

Static scoping = lexical scoping

- find a variable based on structure of code
- typically, by looking outward in surrounding blocks

Dynamic scoping -

- find a variable based on what was last seen based on program execution

1

2

static A

The diagram illustrates the difference between static and dynamic scoping using Scheme code. A vertical arrow labeled 'static' points to the first line of code, `(define x 1)`. A curved arrow labeled 'dynamic' points from the `(fun2)` call inside `fun1` to the `(define fun2)` definition. The code is as follows:

```
(define x 1)

(define fun1
  (lambda ()
    (let ((x 2))
      (fun2))))

(define fun2
  (lambda ()
    (display x)))

(fun1)
```

Python, Java, Kotlin, Scheme  
all use static scoping

Perl uses dynamic scoping

Why the difference, and what  
are pros/cons?

Many people find dynamic scoping  
to be more intuitive.

Most early languages did dynamic scoping. Early developers thought it made sense.

History begged to differ.

Why?

```
(define x 1)

(define fun1
  (lambda ()
    (let ((x 2))
      (fun2))))

(define fun2
  (lambda ()
    (display x)))

(fun1)
```

```
x = "1"

function fun1
{
    local x = "2"
    fun2;
}

function fun2
{
    echo $x;
}

fun1
```

Task: debug fun2

So... which x?

Static Scoping: look at code  
(or use editor tools)



Dynamic:

I have to consider all  
1000 places in code that  
call fun2.

E.g. we call `talloc` from  
everywhere.

```
void *talloc( — ) {
```

```
    ==  
    ==  
    ==
```

```
    list = - - - -
```

```
        ↑  
    global
```

```
        - - - -  
        malloc
```

```
        - - - -
```

```
    }
```

---

but buried, deep in code

```
void evilmcker() {
```

```
    Node *list = bad thing  
    talloc()
```

```
}
```

// if dynamically scoped,  
talloc uses wrong list

In hindsight, dynamic scoping, though seemingly intuitive was a bad idea, mostly.

- still here in langs w/ historical baggage
  - arguments made that it still makes sense in highly interactive langs, designed to be used one command at a time, like shell scripting
- 

Going forward, we need to implement static scoping in Scheme.

