

# Today

How do we add capability to the interpreter for functions we can't write, like  $+$ ?

Why not

big eval loop

if (car(—))

- "if" —

- "lambda" —

- "define" —

- "+" —

- "-" —

+ "\*" —

- "/" —

- "%"

} scales poorly

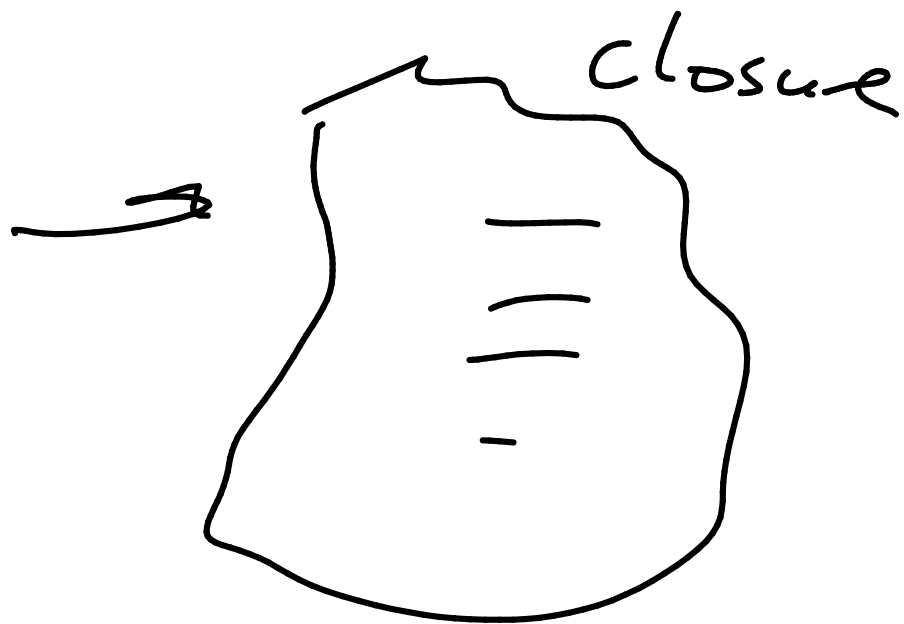
(define myadd +)  
(myadd 3 5)

Code above would fail on this.

You need a way to treat functions  
as data so you can point at them

You are doing this right now w/  
closures for functions defined in  
Scheme.

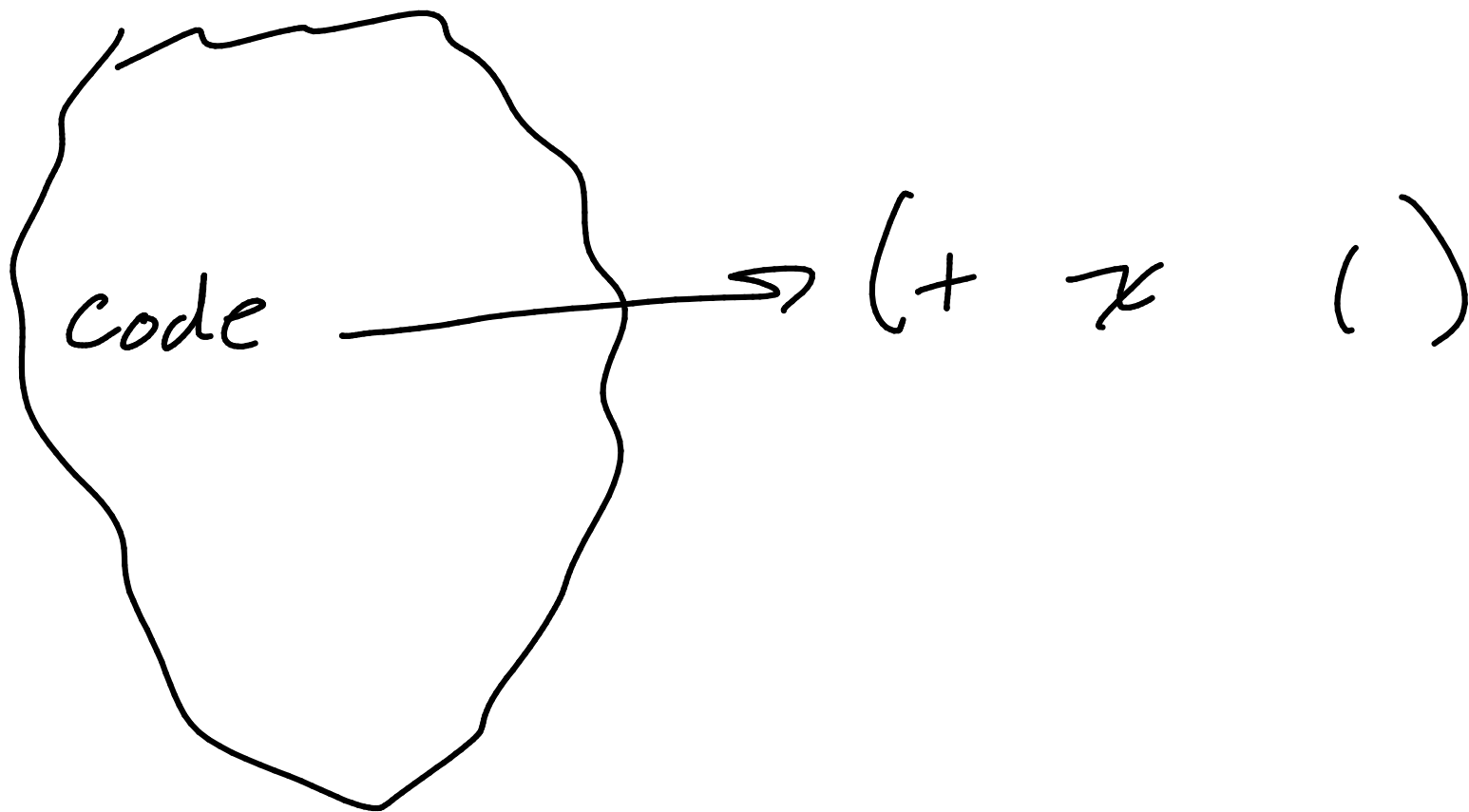
(lambda (x)  
 (+ x 1))



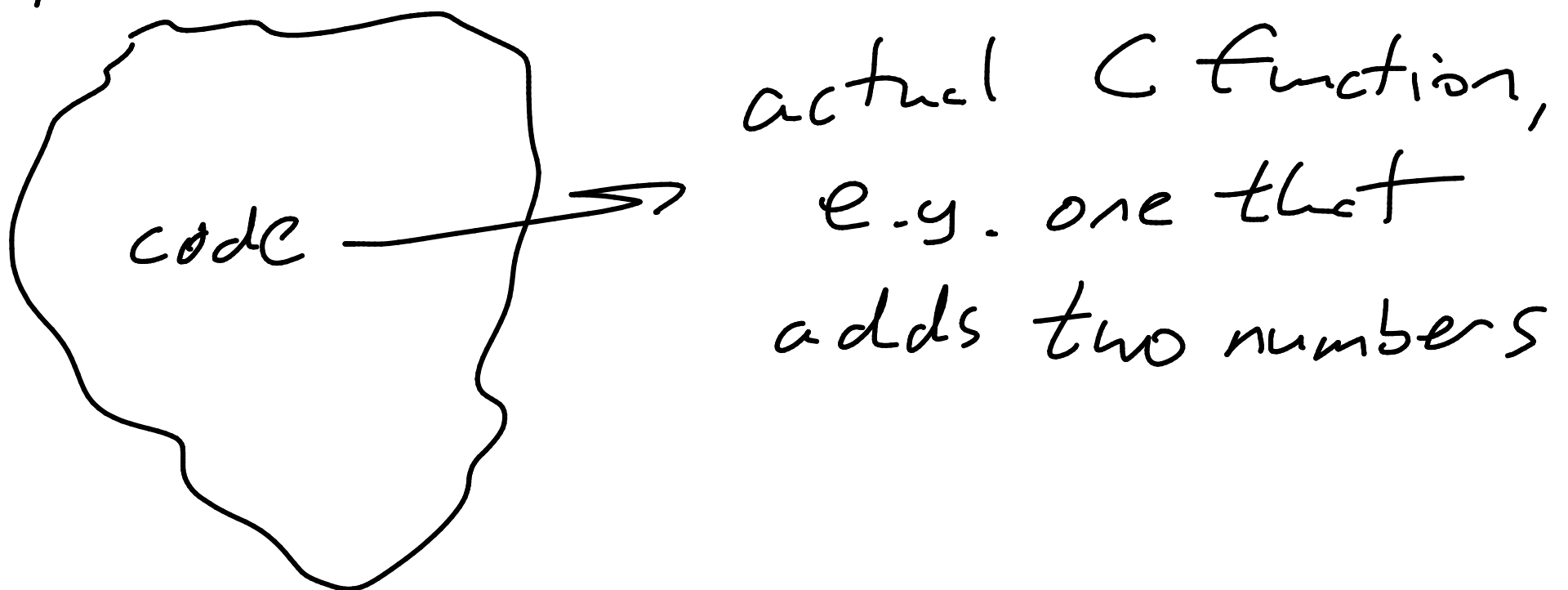
Implement something similar to  
closures, that we'll call primitives.

A primitive (`PRIMITIVE_TYPE`) is another SchemeVal that points to actual C code.

closure



Primitive



Look at function pointers in C.

```
typedef struct SchemeVal {  
    objectType type;  
    union {
```

```
        ...  
        struct SchemeVal *(*pf)(struct SchemeVal *);  
    };  
} SchemeVal;
```

pf is a pointer

return type:  
it returns a  
ptr to  
a SchemeVal

name of  
the function  
pointer  
(pf)

one  
param -  
ptr to  
a SchemeVal  
(i.e. a list  
of SchemeVals)

SchemeVal \*primitiveExp(SchemeVal \*args) {  
 // check that args has length one and car(args) is numerical  
 // assume that car(args) is of type double, should check that as we  
 SchemeVal \*result = malloc(sizeof(SchemeVal));

result->type = DOUBLE\_TYPE;

// Uses built-in C function exp

result->d = exp(car(args) → d);  
return result;

\* Error checking  
missing  
frame

exp

void bind(char \*name, SchemeVal \*(\*function)(SchemeVal \*), Frame \*frame)  
 // Add primitive functions to top-level bindings list  
 SchemeVal \*value = malloc(sizeof(SchemeVal));

value->type = PRIMITIVE\_TYPE;

value->pf = function;

// Your code will differ from each other some on the following;  
// it adds a binding to a frame; can leave that out here  
frame->bindings = ....;

Yai've done this

void interpret(SchemeVal \*tree) {

// Make top-level bindings list

Frame \*frame = malloc(sizeof(Frame));  
frame->bindings = makeEmpty();  
frame->parent = NULL;

bind("exp", primitiveExp, frame);

// ...

x=3, y=5  
exp → primitiveExp

Preview of what's next:

types in programming languages

Languages generally fall into 2 categories regarding where the type ~~of~~ goes.

C

int x = 3;

↑  
type

x (int)

3

x = 9.6;  
// error

Python

x = 3

x →

3  
(int)

x = 9.6

x →

9.6  
(double)