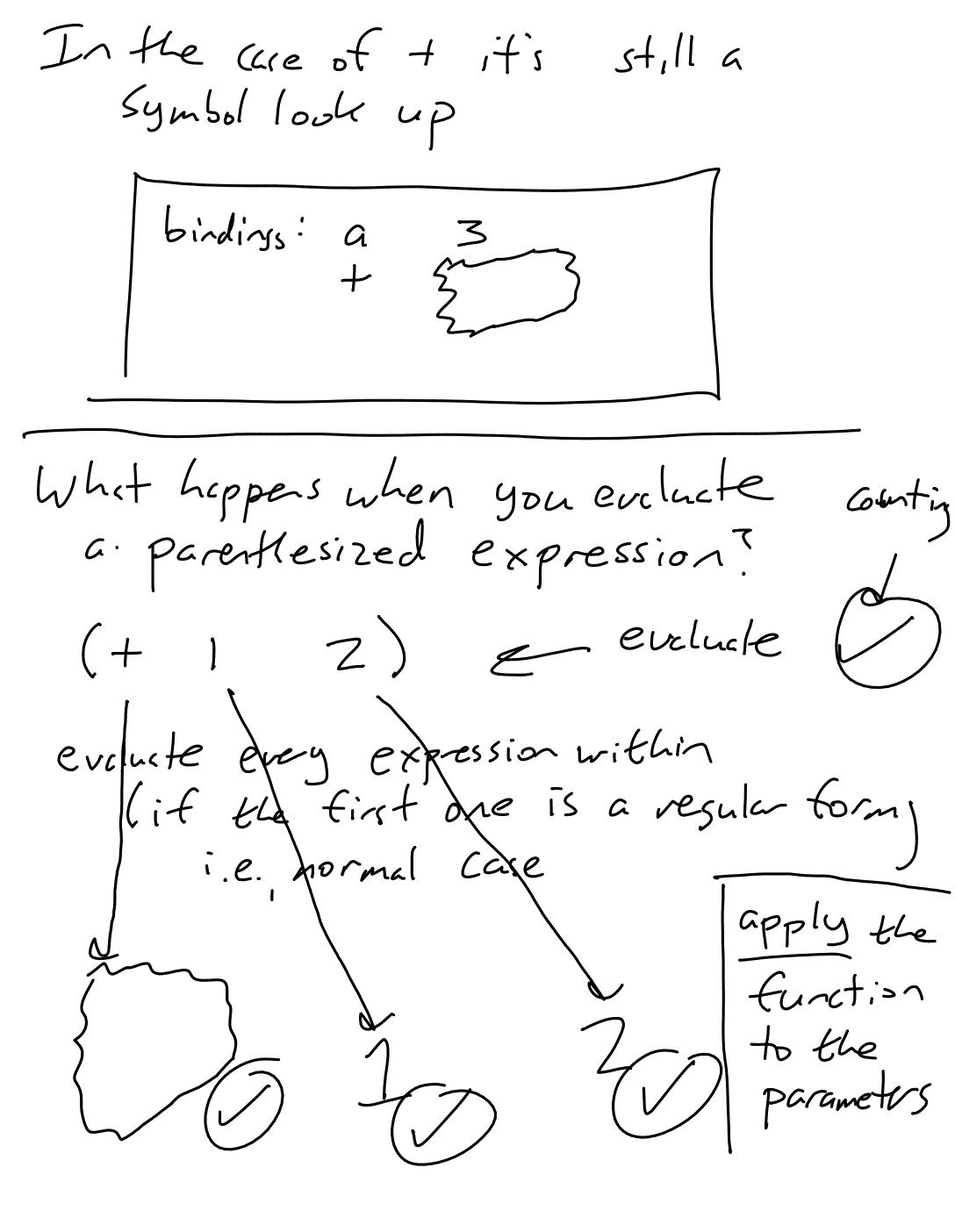
How we evaluate Scheme expressions [we'll get to quote in particular] Landdal closures, may be

When he talk about evaluating a Schene expression he mean exenting it, or interpreting it, but effectively "doing what Scheme does with it." evoluck 3 = 3 evelucting a constant just returns itself, just one call to eval evolucting a symbol looks up the Value and returns it - one call to



evilute (3) Hou if is not a normal Schome Function. it is a special form, and it brecks how Scheme normally works.

Why did the Scheme designers make it weird? it does "short circuit evaluation) (if #t (+ 3 5) (+ 9)) when the condition is true, skip evaluations the alternative when the condition is false, skip evoluctions the consequent That meas that if breaks the whole normal Scheme evoluction process,

(Repeat) How does Scheme evaluable a parænthesized expression? (+ 1 2) evolucke if the first symbol is not a Special form, proceed as + 1 2 d 2 2 2 2 4 cucluctions if the first symbol is a special torm it is entirely in Control of what evals happen. lis our code, we seve eval the Special form itself.)

(if #t 3 5) e eval AACK! special form. never sets evaluated, and instead, Le pass the remaining Parmetes unevaluated to a helper function that does the special form Call note! not eval uneviled

eval ---- #t 3 5) - evel condition #t > #t Since true, evel 3->3

(le+ ((x 3)) 1 (+ x x) ARGHHH special form does evel 3 is when binding of Afte Frame is set let says "evel (+ evals

(let, Code to evaluate list of in the context bindings of the ((x 3)(y 5)) new trame eval Peval bind bind eva | quote is also a special form (quote _____) evel returns evels nothing

special form nooveeoo!

Special form nooveeoo!

Yechhall!

do no evels just return

the parameter, unevel'ed