

Math 220-01: Mathematical Reasoning and Proof
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“Notes”

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These are very rough notes for the course, which mostly overlap with the class content.

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MATH 220 HANDOUT 1 - LOGIC

A **statement** is a sentence for which ‘true or false’ is meaningful.

1. Which of these are **statements**?

- (1) Today it is raining.
- (2) What is your name?
- (3) Every student in this class is a math major.
- (4) $2 + 2 = 5$.
- (5) $x + 1 > 0$.
- (6) $x^2 + 1 > 0$.
- (7) If it is raining, then I will wear my raincoat.
- (8) Give me that.
- (9) This sentence is false.
- (10) If x is a real number, then $x^2 > 0$.

2. Which of these are true?

- (1) (T or F) Every student in this class is a math major and a human being.
- (2) (T or F) Every student in this class is a math major or a human being.
- (3) (T or F) $2 + 2 = 5$ or $1 > 0$.
- (4) (T or F) If x is a real number, then $x^2 \geq 0$.
- (5) (T or F) If x is a complex number, then $x^2 \geq 0$.

3. Write the negations of the following.

- (1) $2 + 2 = 5$
- (2) $1 > 0$.
- (3) $2 + 2 = 5$ or $1 > 0$.
- (4) Every student in this class is a math major.
- (5) Every student in this class is a math major or a human being.
- (6) If x is a real number, then $x^2 > 0$.

4. Prove the following using truth tables.

- (1) $P \wedge (Q \vee R) = (P \wedge Q) \vee (P \wedge R)$,
- (2) $(P \vee Q) \vee R = P \vee (Q \vee R)$. (We thus write $P \vee Q \vee R$ for both.)
- (3) $\neg(P \vee Q) = \neg P \wedge \neg Q$,
- (4) $\neg(P \wedge Q) =$ (make a guess similar to problem 3),
- (5) $\neg(\neg P) = P$.

5. In exercise 6, you may use the following variants of exercise 4.

- (1) $P \vee (Q \wedge R) = (P \vee Q) \wedge (P \vee R)$,
- (2) $(P \wedge Q) \wedge R = P \wedge (Q \wedge R)$. (We thus write $P \wedge Q \wedge R$ for both.)
- (3) $P \vee Q = Q \vee P$.
- (4) $P \wedge Q = Q \wedge P$.

6. Prove or disprove the following *without* using truth tables.

- (1) $\neg(P \wedge \neg Q) = \neg P \vee Q$.
- (2) $P \vee ((Q \wedge R) \wedge S) = (P \wedge Q) \vee (P \wedge R) \vee (P \wedge S)$.
- (3) $P \vee (Q \wedge R) \wedge S = (P \vee Q) \wedge (P \vee R) \wedge (P \vee S)$.

7. Write the negations of the following implications.

- (1) If n is even, then n^2 is even.
- (2) If $1 = 0$, then $2 + 2 = 5$.
- (3) If there is free coffee, then DZB will drink it
- (4) If $1 = 0$ and $2 + 2 = 5$, then the sky is blue and kittens are popular on youtube
- (5) If x and y are real numbers such that $xy = 0$, then $x = 0$ or $y = 0$.

8. Which of these are true?

- (1) (T or F) For all $x \in \mathbf{Z}$, x is divisible by 2.
- (2) (T or F) There exists an $x \in \mathbf{Z}$ such that x is divisible by 2.
- (3) (T or F) For all $x \in \mathbf{R}$, if $x \neq 0$, then there exists a $y \in \mathbf{R}$ such that $xy = 1$.
- (4) (T or F) For all $x \in \mathbf{R}$, there exists a $y \in \mathbf{R}$ such that $xy = 1$.

9. Write the negations of the following.

- (1) For all $x \in \mathbf{Z}$, x is divisible by 2.
- (2) There exists an $x \in \mathbf{Z}$ such that x is divisible by 2.
- (3) $\neg(\forall x, P(x))$,
- (4) $\neg(\exists x \text{ s.t. } Q(x))$
- (5) $\forall x, (P(x) \wedge Q(x))$.
- (6) If $\exists x \in \mathbf{R}$ such that $2x = 1$, then for all y , $y^2 < 0$.
- (7) For all $x \in \mathbf{R}$, there exists a $y \in \mathbf{R}$ such that $xy = 1$.

10. Write the converse and contrapositive of the statements from problem 7.

MATH 220 HANDOUT 2 - DIVISIBILITY

- (1) Show that if $d \neq 0$ and $d | a$, then $d | (-a)$ and $-d | a$.
- (2) Show that if $a | b$ and $b | a$, then $a = b$ or $a = -b$.
- (3) Suppose that n is an integer such that $5 | (n+2)$. Which of the following are divisible by 5?
 - (a) $n^2 - 4$
 - (b) $n^2 + 8n + 7$
 - (c) $n^4 - 1$
 - (d) $n^2 - 2n$
- (4) Prove that the square of any integer of the form $5k + 1$ for $k \in \mathbf{Z}$ is of the form $5k' + 1$ for some $k' \in \mathbf{Z}$.
- (5) Show that if $ac | bc$ and $c \neq 0$, then $a | b$.
- (6)
 - (a) Prove that the product of three consecutive integers is divisible by 6.
 - (b) Prove that the product of four consecutive integers is divisible by 24.
 - (c) Prove that the product of n consecutive integers is divisible by $n(n-1)$.
 - (d) (Challenge problem) Prove that the product of n consecutive integers is divisible by $n!$.
- (7) Find all integers $n \geq 1$ so that $n^3 - 1$ is prime. Hint: $n^3 - 1 = (n^2 + n + 1)(n - 1)$.
- (8) Show that for all integers a and b ,

$$a^2b^2(a^2 - b^2)$$

is divisible by 12.

- (9) Suppose that a is an integer greater than 1 and that n is a positive integer. Prove that if $a^n + 1$ is prime, then a is even and n is a power of 2. Primes of the form $2^{2^k} + 1$ are called Fermat primes.
- (10) Suppose that a and n are integers that are both at least 2. Prove that if $a^n - 1$ is prime, then $a = 2$ and n is a prime. (Primes of the form $2^n - 1$ are called Mersenne primes.)
- (11) Let n be an integer greater than 1. Prove that if one of the numbers $2^n - 1, 2^n + 1$ is prime, then the other is composite.
- (12) Show that every integer of the form $4 \cdot 14^k + 1$, $k \geq 1$ is composite. Hint: show that there is a factor of 3 when k is odd and a factor of 5 when k is even.
- (13) Can you find an integer $n > 1$ such that the sum

$$1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n}$$

is an integer?

Week 3: proof by contradiction

Prove that for $x, y \in \mathbb{R}$, P

$$x + y > 2Q \Rightarrow \boxed{x > 10 \text{ or } y > 10.}$$

Proof: Assume $x+y > 2Q$. Assump.

(Q is either true or false)

(either $\rightarrow Q^c$ or $\neg Q$ is true)

Proceed by contradiction. Assume $\neg P$,
i.e., $x \leq 10$ and $y \leq 10$. Adding Post fact
gives Q $x+y \leq 20$. This contradicts $\neg P \Rightarrow Q$
assumption. Thus $\neg P$ must have been
false, so P is true, i.e.,
 $x > 10$ or $y > 10$. \square

Template for proof by contradiction

- We want to prove P .
- There are "2 cases": P is true or P is false.
- If we can rule out " P is false", then P must be true.
- To begin, Assume P is false, i.e. $\neg P$.
- "Argue"; i.e. exhibit some (correct) chain of reasoning end up w/ a statement Q .
(IE write out a proof of " $\neg P \Rightarrow Q$ ")

- Observe (or give a proof that) Q is false.
- Conclude that $\neg P$ is false.

$$(\neg P \Rightarrow Q), \neg \neg Q \Rightarrow P$$

Intuition:

. Chess "If I move here, in
if moves they have
check. So I shouldn't
make that move"

Prove that $x^2 - y^2 = 1$ has no
~~positive~~
integer solutions.

$$(x, y) = (1, 0)$$

y
not positive.

Proof. Proceed by contradiction.

Assume that $[x, y]$ is a positive integer solution to $x^2 - y^2 = 1$. Then $(x-y)(x+y) = 1$. There are 2 possibilities:

$x-y=1$ and $x+y=1$, or $x-y=-1$ and $x+y=-1$.

In the 1st case adding gives $2x = 2$.
Then $x = 1$, but $x - y = 1$, $\boxed{y = 0}$. ^Q This
is a contradiction since y is positive.
In the 2nd case, adding gives $2x = -1$,
so $\boxed{x = -1}$. ^Q This is a contradiction
since x is positive.

In both cases, we get a contradiction.
Thus our assumption was wrong,
and we conclude that there are
no positive integer solutions to
 $x^2 - y^2 = 1$. \square

N.B. Really was a proof that

$$\neg P \Rightarrow (Q_1 \text{ or } Q_2)$$

need both to give a contradiction.

IE if there are cases in concern,
need each case to be false.

Prove that $x^2 = 4y + 3$ has no integer solutions. $\rightarrow p$

Proof: Proceed by contradiction. Assume that there are $x, y \in \mathbb{Z}$ s.t.

$$x^2 = 4y + 3.$$

There are 2 cases: x is even or x is odd.

If x is even then the LHS is even and the RHS is odd. This is a contradiction.

If x is odd, then $x = 2k+1$ for some $k \in \mathbb{Z}$. Plugging in gives

$$(2k+1)^2 = 4k^2 + 4k + 1 \equiv 4k + 3.$$

This is a contradiction, since the LHS has a remainder of 1 and the RHS has a remainder of 3. m

Recall: An integer n is prime if its only divisors are ± 1 and $\pm n$.

2, 3, 5, 7, 11, 13, 17, 19, 23, primes

4, 6, 8, 12 not prime

$$65537 = 2^{10^4} + 1$$

65538 not prime

Euclid's theorem: there exist infinitely many primes.

Proof: Proceed by contradiction. Assume there are only finitely many primes. Let's name them P_1, P_2, \dots, P_r .

($P_1 = 2, P_2 = 3, P_3 = 5, \dots, P_r = ?$)

Label them so that $p_1 = 2$, $p_{i+1} > p_i$
let $N = p_1 p_2 p_3 \cdots p_r + 1$.
(what does the factorization look like?)

Since $N > p_r$, N isn't prime (b/c N is bigger than the biggest prime). By FTA
(Fundamental thm of arithmetic), N

factors into primes. Let g be one of those primes. Since p_1, \dots, p_r are all of the primes, $\exists i$ such that $g = p_i$.

Then $g|N$ and $g|p_1 \cdots \overset{\downarrow}{p_i} \cdots p_r$, thus by the λ root of 3 rule,

$q \mid N - (p_1 - p_r)$, i.e. $q \mid 1$. This
is a contradiction since $q > 1$.
Thus there are infinitely many primes. \square

Joke: There are no uninteresting numbers, positive integers.

Proof: Proceed by contradiction.
Assume that some positive integers are uninteresting. Then there must be

a smallest uninteresting positive integer.
But that's pretty interesting!



(This is a proof technique: think about
the "smallest" counterexample)

if $a|b$ and $a|b+c$ then $a|c$

$$a|(-b)$$

$$a|((b+c)+(-b)) \Rightarrow$$

FTA (Fundamental THM of Arithmetic)

Let $N \in \mathbb{Z}_{>1}$. Then

(i) $\exists P_1 - P_r$ primes s.t.

$$N = P_1 \cdot \dots \cdot P_r$$

(ii) If $N = P_1 \cdots P_r = Q_1 \cdots Q_s$ s.t.
 $P_i \geq P_{i-1}$ and $Q_j \geq Q_{j-1}$, then $r \leq s$
and $\forall i, P_i = Q_i$

$$\begin{array}{c}
 60 = 2 \cdot 5 \cdot 2 \cdot 3 \\
 60 = 2 \cdot 2 \cdot 3 \cdot 5 \\
 \swarrow \qquad \qquad \qquad = 2^3 \cdot 3 \cdot 5 \\
 2 \qquad 30 \\
 \swarrow \\
 5 \cdot 6 \\
 \swarrow \\
 2 \cdot 3
 \end{array}$$

$$\begin{array}{c}
 2 \cdot 2 \cdot 3 \cdot 5 \\
 \swarrow \qquad \qquad \qquad 60 \\
 2 \qquad 30 \\
 \swarrow \qquad \qquad \qquad \swarrow \\
 2 \qquad 15 \\
 \swarrow \qquad \qquad \qquad \swarrow \\
 3 \qquad 5
 \end{array}$$

Prve(i) Proceed by contradiction

Assume that some $N \in \mathbb{Z}_>1$ do not factor into primes. Let N be the smallest such integer. Then N is not prime, otherwise it is already factored.

Thus N is composite. Write

$$N = ab \text{ where } 1 < a, b < N.$$

Since $a, b < N$ and N is the smallest integer that doesn't factor a and b factor. Write $a = p_1 \cdots p_r, b = q_1 \cdots q_s$. Then $N = ab = (p_1 \cdots p_r)(q_1 \cdots q_s)$.

This is a contradiction, since
we just factored N . \square

Lecture 7 is about more contradiction

Want to prove P .

Assume $\neg P$. $\neg P \Rightarrow Q$

Argue.

Conclude Q .

Observe that Q is false

Conclude P .

Prove that ($\forall n \in \mathbb{Z}$, n and $n+1$ have no common prime factors)

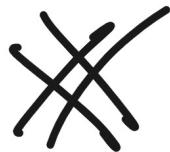
Proof: Proceed by contradiction. Assume $\exists n \in \mathbb{Z}$ s.t. n and $n+1$ have some common prime divisors. Let p be a prime s.t. $\underline{p|n}$ and $\underline{p|n+1}$.

Then $p \mid (n+1) - n$, i.e., $p \mid 1$.

This is a contradiction, since

p is prime, and primes are > 1 .

$\Rightarrow \Leftarrow$



$$\begin{aligned} d \mid a \wedge d \mid b &\Rightarrow d \mid a+b \\ d \mid q &\Rightarrow d \mid -q \\ d \mid ab &\quad d \mid a-b \end{aligned}$$

Let $a, b, c \in \mathbb{Z}$. Suppose $\underline{a^2 + b^2 = c^2}$.
Show that abc is even. Hyp

Proof: Proceed by contradiction.
Assume $\underline{a^2 + b^2 = c^2}$ and abc is odd.
Then a, b , and c are each odd. (B/c if one were even, abc would be even)

Then $a^2, b^2, + c^2$ are odd (b/c products of odd integers are odd). Then
 $a^2 + b^2$ is even, but c^2 is odd.

This is "even = odd" which
is a contradiction. \square

Defn: A number x is rational if
 $\exists a, b \in \mathbb{Z}$ s.t. $b \neq 0$ and $x = \frac{a}{b}$

A rational # is reduced if a and b have no common prime divisors

Examples: $\frac{2}{3} = \frac{4}{6}$ ~~Fact: Every rational # can be reduced~~
reduced not reduced.

Prove: $\sqrt{2} \notin \mathbb{Q}$. (It $\sqrt{2}$ is not Rational.)

Proof: Proceed by contradiction. Assume $\sqrt{2} \in \mathbb{Q}$.

Then $\exists a, b \in \mathbb{Z}$ s.t. $b \neq 0$ and $\sqrt{2} = a/b$.

Assume that a & b are reduced. In particular at least one of a or b is odd. Then

$$b\sqrt{2} = a. \text{ Then } b^2 \cdot 2 = a^2.$$

Since the LHS is even, the RHS is even, i.e.,
 a^2 is even. Thus a is even. (By how, if
 a is odd a^2 is odd.) Then $4(a^2)$ indeed,
 of course +flg true dt leg). Write $a = 2k$
 for some $k \in \mathbb{Z}$. Then $b^2 \cdot 2 = (2k)^2 = 4k^2$.
 Then $b^2 = 2k^2$. Since the RHS is even, b^2 is even
 so b is even. This is a contradiction,

Since $a+b$ are not both odd.

Hence: $\sqrt{3} \notin \mathbb{Q}$

$$\frac{a}{3} \neq \frac{b}{\sqrt{3}}$$

From 1

$$\pi = \frac{3.141\dots}{e}$$

Prove: $\log_3 2 \notin \mathbb{Q}$.

Recall: $3^{\log_3 2} = 2$

Proof: Proceed by contradiction.

Assume $\log_3 2 \in \mathbb{Q}$. Then $\exists a, b \in \mathbb{Z}$
s.t. $b \neq 0$ and $\log_3 2 = \frac{a}{b}$. WMA
(we may assume) $a+b$ are reduced.

$$\text{Then } 2 = 3^{\log_3 2} = 3^{\frac{a}{b}}.$$

Then $2^b = 3^a \cdot ((3^{a/b})^b \geq 3^a)$

Since $b > 0$, the LHS is even.
But the RHS is odd. This is
a contradiction. \square

Prove: If $a \in \mathbb{Q}$ and $b \notin \mathbb{Q}$, then
 $a+b \notin \mathbb{Q}$.

Proof: Assume $a \in \mathbb{Q}$ and $b \notin \mathbb{Q}$. Proceed by contradiction. Assume $a+b \in \mathbb{Q}$. Then
 $\exists c, d \in \mathbb{Z}$ s.t. $d \neq 0$ and $a = \frac{c}{d}$. Then
 $\exists e, f \in \mathbb{Z}$ s.t. $f \neq 0$ and $a+b = \frac{e}{f}$

$$\text{Then } b = (a+b) - a = \frac{e}{f} - \frac{c}{d} = \frac{ed - cf}{fd}.$$

Since $ed - cf, fd \in \mathbb{Z}$ and $fd \neq 0$,
 $b \in \mathbb{Q}$. Thus contradicting $b \notin \mathbb{Q}$.

We conclude that $a+b \notin \mathbb{Q}$. □

let a, b, c be odd. let x be a solution
to $ax^2 + bx + c = 0$. Prove $x \notin \mathbb{Q}$.

Proof, Assume a, b, c are odd.

Assume $ax^2 + bx + c = 0$.

Proceed by contradiction. Assume $x \in \mathbb{Q}$.

Then $\exists d, e \in \mathbb{Z}$ s.t. $e \neq 0$ and
 $x = d/e$. WMA of d & e are reduced.

Then $a\left(\frac{d}{e}\right)^2 + b\frac{d}{e} + c = 0$. \begin{cases} \text{begin} \\ \text{equation} \end{cases}

Then ad^2 + bde + ce^2 = 0. (1, 1)

Since d, e are reduced, at least one
is odd.

There are 3 cases: d+e is odd, d is even + e is odd,
or d is odd + e is even

In case 1, the LHS of (1.1) is "odd odd + odd
= odd = even"
a contradiction.

In case 2, the LHS is "e + e = e",
Case 3 is similar.

In each case, we have a contradiction

$2, 3, 5, 7, 11, \dots$ $P \neq 1$ and $\begin{cases} \text{if } P = ab, \\ a = \pm 1 \text{ or } b = \pm 1 \end{cases}$

P is prime if $\cancel{P = ab}$

$a = \pm 1 \text{ or } b = \pm 1$

Fermat: $l = l \cdot l = l \cdot l \cdot l = \dots$

Fermat:

$$2^0 + 1 = 1 + 1 = 2$$

$$2^1 + 1 = 2 + 1 = 3$$

$$2^2 + 1 = 4 + 1 = 5$$

$$2^4 + 1 = 16 + 1 = 17$$

$$2^8 + 1 = 256 + 1 = 257$$

$$2^{16} + 1 = 65536 + 1 = 65537$$

Fermat conj'd that $\forall n, 2^n + 1$ is
False prime.
Modern conj $2^n + 1 \nmid 3$ never prime
if $n \geq 5$.

$$2^3 + 1 = 8 + 1 = 9 = 3 \cdot 3$$

$$2^5 + 1 = 32 + 1 = 33 = 3 \cdot 11$$

$$2^7 + 1 = 128 + 1 = \underbrace{129} = 3 \cdot 43$$

Problem: If $2^n + 1$ is prime then n is even.

Proof: Assume $2^n + 1$ is prime.

Proceed by contradiction. Assume n is odd.

$$\begin{aligned} n^2 - 1 &= (n-1)(n+1) \\ x^n - 1 &= (x-1)(x^{n-1} + x^{n-2}y + \dots + y^{n-1}) \end{aligned}$$

Since n is odd, $(-1)^n = -1$.

Thus $2^n + 1 = 2^n + (-1)^n = 2^n - (-1)^n$.

This factors as $(2 - (-1)) (2^{n-1} - 2^{n-2} + \dots \pm 1)$
= $3 \cdot ?$

If $2^n + 1 > 3$, it is not prime b/c
3 is a proper divisor. \square

Induction: Gausg Syo

$$\begin{array}{rcl} 1 & + & 2 + \dots + 100 = S \\ 100 & + & 99 + \dots + 1 = S \\ \hline 101 + 101 \dots + 101 & = & 2S = 101 \cdot 100 \\ 101 & & S = \frac{101 \cdot 100}{2} \end{array}$$

WTP: $\forall n \in \mathbb{Z}_{>0}$
$$1 + 2 + \dots + n = \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$n=1$	$1 = \frac{1(1+1)}{2} = 1$	$P(n)$
$n=2$	$1+2 = \frac{1+2}{2} + 2 = 2 \left(\frac{1}{2} + 1 \right) = 2 \left(\frac{3}{2} \right)$	
$n=3$	$1+2+3 = \frac{2(3)}{2} + 3 = 3 \left(\frac{2}{2} + 1 \right) = 3 \left(\frac{4}{2} \right)$	
\vdots	\vdots	
$(+2+3+4)$	$= 3 \cdot \frac{4}{2} + 4 = 4 \left(\frac{3}{2} + 1 \right) = 4 \left(\frac{5}{2} \right)$	

Proof, Proceed by induction. The statement is already true for $n \geq 1$ b/c $\boxed{1 = \frac{1(d)}{2} \geq 1}$.

"Assume $P(n)$ "
 Assume that we already know the statement for n . IE assume that $(1+d+\dots+n) = \frac{n(n+1)}{2}$.
 Adding $n+1$ to both sides gives
 $1+d+\dots+n+(n+1) = \frac{n(n+1)}{2} + (n+1)$.

The LHS of this is the LHS of what we want to prove. The RHS is

$$(n+1) \left(\frac{1}{2} + 1 \right) = (n+1) \left(\frac{n+1}{2} \right) = (n+1) \underbrace{\frac{(n+1)+1}{2}}_{\text{!}}$$

We conclude that $\forall n \in \mathbb{Z}_{>0}$,
a proof that
 $1+2+\dots+n = \frac{n(n+1)}{2}$. 

R to R B
a proof that
 $P(n) \Rightarrow P(n+1)$.

Let $P(n)$ be a statement which depends on some integer (usually pos) n .

(Ex. $P(n) = "1+2+\dots+n = \frac{n(n+1)}{2}"$)

Goal: Prove $P(n) \quad \forall n \in \mathbb{Z}_{>0}$.

Step 1: Prove $P(1)$ "Base Case"

Step 2: Prove " $P(n) \Rightarrow P(n+1)$ " "Inductive step"

"Induction" = $P(1) \wedge \underline{\underline{P(n) \Rightarrow P(n+1)}} \Rightarrow \forall n \in \mathbb{Z}_{\geq 0}, P(n)$

$P(1), P(2), P(3), P(4), \dots, P(a), P(a+1), \dots$

Warning: $P(n) \neq \boxed{\frac{n(n+1)}{2}}$ ← Not "if or F"

$$P(n) = \boxed{1+2+\dots+n} = \frac{n(n+1)}{2}$$

Note: n is just a variable.

$P(1)$ & $P(n) \Rightarrow P(n+1)$ "prove the following case"
 $P(a) \Rightarrow P(a+1)$
 $P(a-1) \Rightarrow P(a)$

Variant 3: $P(0) \wedge P(n) \Rightarrow P(n+1)$

$$\boxed{P(\alpha) \wedge P(n) \Rightarrow P(n+\alpha)}$$

$$\Rightarrow \forall n \in \mathbb{N}_0, P(n)$$

$$P(1) \wedge P(n) \Rightarrow P(n+1) \quad | \quad P(0), P(1), P(\neg)$$

$$\forall n \in \mathbb{N}_0, P(n)$$

Define: a sequence $a_1, d_1, \dots, d_n, \dots$.

$$a_1 = 2 = 2^1$$

$$a_n = 2 \cdot a_{n-1}$$

"recursive defn"

$$d_n = 2 \cdot d_{n-1}$$

$$d_{n+1} = 2 \cdot d_n$$

$$d_2 = 2 - d_1 = 2 - 2 = 4 = 2^2$$

$$d_3 = 2 - d_2 = 2 - 4 = 8 = 2^3$$

$$d_4 = 2 - d_3 = 2 - 8 = 16 = 2^4$$

Claim: $\forall n \in \mathbb{Z}_{\geq 0},$ $a_n = \lambda^n$

$$P(n) = "a_n = \lambda^n"$$

Proof: Proceed by induction.

Base case: $P(1)$ is " $a_1 = 2$ ", i.e., "2 = 2".

$(P(n) \Rightarrow P(n+1))$ Th β B + ne.

Inductive step: Assume $P(n)$, iE, $a_n = 2^n$,

(wTP $P(n+1)$, i.e., $a_{n+1} = 2^{n+1}$) Then $a_{n+1} = 2 \cdot a_n$

by the defn. of a_n . Then $a_{n+1} = 2 \cdot 2^n = 2^{n+1}$.

Thus $P(n+1)$ is true. \square

Defini: $a_1 = 0$

$$a_n = \overbrace{\sqrt{3 + 2a_{n-1}}}$$

$$a_1 > 0$$

$$a_2 = \overbrace{\sqrt{3 + 2 \cdot 0}} = \sqrt{3}$$

$$a_3 = \overbrace{\sqrt{3 + 2 \cdot a_2}} = \overbrace{\sqrt{3 + 2\sqrt{3}}}$$

$$a_4 = \overbrace{\sqrt{3 + 2a_3}} = \overbrace{\sqrt{3 + 2\sqrt{3 + 2\sqrt{3}}}}$$

Pass: $\forall n \in \mathbb{Z}_{>0}$,

$$a_n < 3$$

$$\text{Pf: } "a_n < 3"$$

Proof: Proceed by induction.

Base case: $P(1)$ is " $a_1 < 3$ ", i.e., " $0 < 3$ ".
This is true.

I.S. Assume $P(n)$, i.e., " $a_n < 3$ ".

(WTP: $P(n+1)$, i.e., " $a_{n+1} < 3$ ") Then, by defn,

$$a_{n+1} = \sqrt{3 + 2a_n}. \text{ Since } a_n < 3,$$

$$a_{n+1} = \sqrt{3 + 2a_n} < \sqrt{3 + 2 \cdot 3} = \sqrt{3+6} = \sqrt{9} = 3$$

Thus $a_{n+1} < 3$. \square

Claim: " $1 + 2 + 4 + 8 + \dots + 2^{n-1} = 2^n - 1$ "

Proof: Proceed by induction. $\forall n \in \mathbb{Z}_{\geq 0}$

BC: $P(1)$ is " $1 = 2 - 1$ ". This is true.

IS: Assume $P(i)$, i.e., $1 + 2 + \dots + 2^{i-1} = 2^i - 1$.

(With $P(i+1)$, i.e., $1 + 2 + \dots + 2^{i+1-1} = 2^{i+1} - 1$)

$$1 + 2 + \dots + 2^{i-1} + 2^i = 2^{i+1} - 1$$

Adding \hat{z}^i to each side gives

$$1 + \hat{z} + \dots + \hat{z}^{i-1} + \hat{z}^i = \hat{z}^{i-1} + \hat{z}^i.$$

The LHS is the LHS of $P(i+1)$.

The RHS is $\hat{z} - \hat{z}^i - 1 = \hat{z}^{i+1} - 1$.

This is the RHS of $P(i+1)$. □

Claim: $\forall n \in \mathbb{Z}_{\geq 0}, 3 \mid 4^n - 1.$

$P(n)$

Proof: Proceed by induction.

Base Case: $P(0) \rightarrow 3 \mid "3 \mid 4^0 - 1",$ i.e., $3 \mid 0$ which is true.

IS: Assume $P(n)$, i.e., $3|4^n - 1$.

(wP $3|4^{n+1} - 1$) Then $\exists m \in \mathbb{Z}$ s.t.

$$4^n - 1 = 3m, \text{ Then } 4^{n+1} = 3m + 1.$$

$$\begin{aligned} \text{Then } 4^{n+1} - 1 &= 4^n \cdot 4 - 1 = (3m+1) \cdot 4 - 1 \\ &= 12m + 4 - 1 = 12m + 3 = 3(4m+1). \end{aligned}$$

Thus $3 \mid 4^{n+1} - 1$



More Induction

WTP $\forall n \in \mathbb{Z}_{>0}, P(n)$

BC = Prove $P(1)$.

IS: Prove " $P(\alpha) \Rightarrow P(\alpha+1)$ "

Prove: $\sum_{i=1}^n i^2 = 1^2 + 2^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$

Proof: Proceed by induction.

Base case: $P(1)$ is " $1^2 = 1(2)(3)/6$ "

$$\text{FE } 1 = 1.$$

Inductive step : Assume $P(a)$, i.e.,

$$1^2 + 2^2 + \dots + a^2 = a(a+1)(2a+1)/6. \quad \begin{matrix} \text{WTP} \\ P(a) \Rightarrow P(a+1) \end{matrix}$$

Add $(a+1)^2$ gives

$$1^2 + 2^2 + \dots + a^2 + (a+1)^2 = \frac{a(a+1)(2a+1)}{6} + (a+1)^2.$$

The LHS is the LHS of $P(a+1)$.

The RHS is $a(a+1)(2a+1) + (a+1)^2 =$

$$(a+1) \left(\frac{a(2a+1)}{6} + \frac{6(a+1)}{6} \right) =$$

$$(a+1) \left(\frac{2a^2 + a + 6a + 6}{6} \right) = (a+1) \left(\frac{2a^2 + 7a + 6}{6} \right)$$

$$\frac{a(a+1)(a+\alpha)(2a+3)}{6} =$$

$$\frac{(a+1)(a+1+1)(2(a+1)+1)}{6} \text{ This is } \frac{6}{6} \text{ He}$$

RMS of $\beta(a+1)$. ~~ex~~

Claim! $\forall a \in \mathbb{Z}_{\geq 0}, 3^a$ is odd.

Proof BC: $P(a)$ is " 3^a is odd",
i.e., "1 is odd. This is true,

FS: Assume $P(a)$, i.e., 3^a is odd.

(WTF: $P(a+1)$, i.e. 3^{a+1} is odd)
 $\frac{||}{3^a \cdot 3}$

Then $3^{9+1} = 3^9 \cdot 3$. Since 3^9 is odd
+ 3 is odd, and since the
product of odd integers is odd,

3^{9+1} is odd. ~~Therefore~~

Slogan: "Any time something 'works' for
2 things, it works for many things".

Example: we know that "odd · odd = odd"

Pf: $(2a+1) \cdot (2b+1) = 4ab + 2(a+b) + 1$

rem. 2 | . \rightarrow

Claim: If a_1, \dots, a_n are odd integers
then $a_1 \cdot a_2 \cdot \dots \cdot a_n$ is odd.

QED

Pf: Proceed by induction.

BC is P(1) i.e., "if a, b odd,
then a, b odd". This is true.

$P(2)$ is "if a_1, a_2 are odd, then
 $a_1 a_2$ is odd." This is true
and we previously proved it,
FS: Prove that for $n \geq 2$, $P(n) \Rightarrow P(n+1)$

Assume $P(n)$, i.e., "if a_1, \dots, a_n are odd
then $a_1 \dots a_n$ is odd."

(WTP: $P(n+1) \rightarrow [$ if a_1, \dots, a_{n+1} are odd then
 $a_1 \dots a_{n+1}$ is odd] $]$)

Assume a_1, \dots, a_{n+1} are odd.

Then a_1, \dots, a_n is odd because $P(a)$ is true. Then $(a_1, \dots, a_{n+1}) = (a_1, \dots, a_n) \cdot a_{n+1}$. Since (a_1, \dots, a_n) is odd and a_{n+1} is odd, since we knew $P(a)$, their product is odd. Thus a_1, \dots, a_{n+1} is odd \square

$$a+b = b+a$$

$$\begin{aligned}a+b+c &= a+(b+c) = a+(c+b) \\&= (a+b)+c\end{aligned}$$

2 base cases

$$P(1) \wedge P(2) \wedge (P(1) \Rightarrow P(n+1))$$

The proof that $P(n) \Rightarrow P(n+1)$
doesn't work for $n=1$.

Prove: $n! > 2^n$ for $n \geq 4$.

~~Induction~~

$P(n)$

$$n! = n \cdot (n-1) \cdot (n-2) \cdots 1$$

$$\left| \begin{array}{c} n = \\ 1 \\ 2 \\ 3 \\ 4 \end{array} \right.$$

$$P(n) =$$

$$1! = 1 > 2$$

$$2! = 2 > 4$$

$$3! = 6 > 8$$

$$4! = 24 > 16$$

Proof: Proceed by induction.

BC: $P(4)$ is " $4! > 2^4$ ", i.e., " $24 > 16$ ".
ThB B true.

Assume $n \geq 4$ and assume $P(n)$, i.e., $n! > 2^n$.
(WTP $P(n+1)$, i.e., $(n+1)! > 2^{n+1}$)

Multiplying by $\varphi(n)$ by $n+1$ gives

$$(n+1) n! > (n+1) 2^n.$$

The RHS of this is $(n+1)!$.

Since $n \geq 4$, $n+1 \geq 5 \geq 2$.

Thus $(n+1) 2^n \geq 2 \cdot 2^n = 2^{n+1}.$

We conclude that $(n+1)! > 2^{n+1}$. \square

Fibonacci #'s $F_1 = 1, F_2 = 1, F_3 = 1+1$

$F_4 = 3, \dots$

1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144
Last Perfect Power

Defn: $F_1 = 1, F_2 = 1$

$F_n = F_{n-1} + F_{n-2}$

Same as

$$F_{n+1} = F_n + F_{n-1}$$

$$f_a = f_{a-1} + f_{a-2}$$

$$f_{n+2} = f_{n+1} + f_n$$

$$F_{d,n+2} = F_{d,n+1} + F_{d,n}$$

Claim: " $F_1 + F_3 + F_5 + \dots + F_{2n-1} = F_{2n}$ "

Ex's 1, 2, 3, 5, 8, 13, ... "P(n)"

Proof (B.C) P(1) is " $F_1 = F_2$ ", i.e., " $1 = 1$ ". This is true.

FS: Assume $P(n)$, i.e.,
" $F_1 + F_3 + \dots + F_{2n-1} = F_{2n}$ ".

$$\begin{aligned}2(n+1)-1 &= \\2n+1 &\quad 2n+1\end{aligned}$$

Adding F_{2n+1} to both sides gives

$$F_1 + F_3 + \dots + F_{2n-1} + F_{2n+1} = F_{2n} + F_{2n+1}.$$

By the defn, $F_{2n} + F_{2n+1} = F_{2n+2}$.

Thus $F_1 + F_3 + \dots + F_{2n+1} = F_{2n+2}$ true,

$P(n+1)$ is true. \square

"Lucas"

$$2, 1, 3, 4, 7, 11, 18, 29, \dots$$

$$L_1 = 2 \quad L_2 = 1$$

$$L_n = L_{n-1} + L_{n-2}$$

Claim: $\forall n \in \mathbb{Z}_{\geq 0}, "F_n < \alpha"$

Proof: (BC) $P(0)$ is " $F_0 < \alpha$ ", true,
 $0 < \alpha$. $P(\alpha)$ is " $F_\alpha < \alpha$ ", true,
 $0 < 4$. These are true.

IS: Assume $P(n)$ and $P(n+1)$, i.e.

$$F_n < \hat{\delta}^n \text{ and } F_{n+1} < \hat{\delta}^{n+1}.$$

(WTP $P(n+2)$, i.e., $F_{n+2} < \hat{\delta}^{n+2}$)

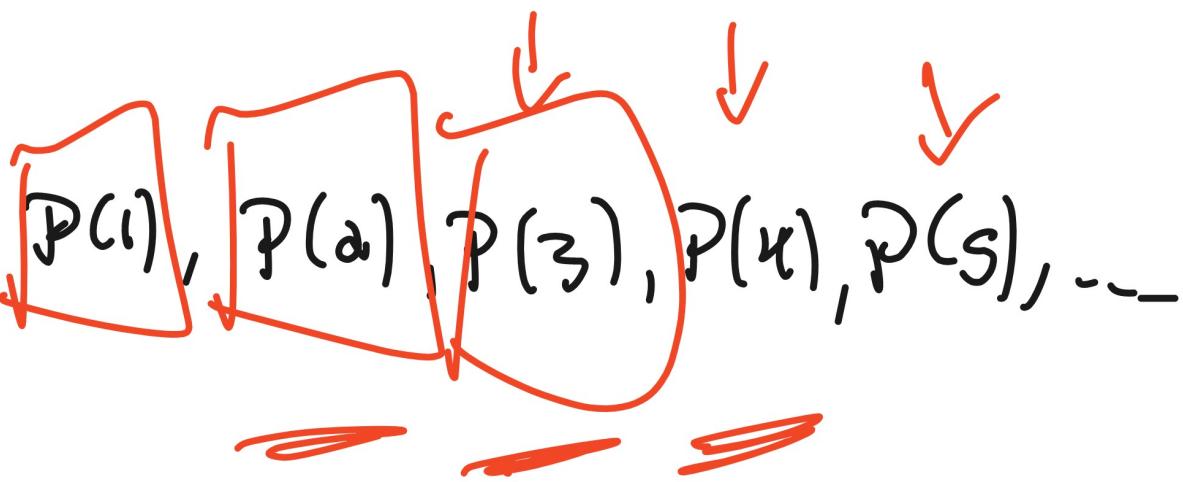
By defn, $F_{n+2} = F_{n+1} + F_n$. Since $P(n)$ and

$$P(n+1) \text{ are true, } F_{n+1} + F_n < \hat{\delta}^{n+1} + \hat{\delta}^n,$$

IE, $F_{n+2} < 2^{n+1} + \delta^n$. (want 2^{n+2})
Since $\delta^n < \delta^{n+1}$, $2^{n+1} + \delta^n < 2^{n+1} + \delta^{n+1}$
 $= \delta \cdot 2^{n+1} = \delta^{n+2}$. Thus

$$F_{n+2} < \delta^{n+2}. \quad \text{By}$$

$$P(1) \wedge P(2) \wedge \left(P(n) \wedge P(n+1) \Rightarrow P(n+a) \right)$$



Claim: " $F_{n-1} \cdot F_{n+1} = F_n^2 + (-1)^n$ " " $= P(0)$

Proof: " $P(0)$ " is the statement $1, 1, 2, 3, 5, 8$

$$\text{" } F_1 \cdot F_3 = F_2^2 + (-1)^2 \text{ "}, \text{ i.e.,}$$

$$1 \cdot 2 = 1^2 + (-1)^2$$

$2 = 2$. This is true.

FS - Assume $P(G)$, i.e., $F_{q-1} \cdot F_{q+1} = F_q^2 + (-1)^q$.
(wrt $P(\text{alt})$, i.e., $F_q - F_{q+1} = F_{q+1}^2 + (-1)^{q+1}$)

By defn, $F_{q-1} + F_q = F_{q+1}$, i.e., $F_{q-1} = F_{q+1} - F_q$.

Sub by gives $(F_{q+1} - F_q) \cdot F_{q+1} = F_q^2 + (-1)^q$.

Then $F_{q+1}^2 - F_q \cdot F_{q+1} = F_q^2 + (-1)^q$.

Then $F_{a+1}^2 - (-1)^a = F_a^2 + F_a \cdot F_{a+1}$,

Then $F_{a+1}^2 + (-1)(-1)^a = F_{a+1}^2 + (-1)^{a+1}$

$= F_a(F_a + F_{a+1}) = F_a F_{a+2}$ by the def.

 Q.E.D.

↑ →
 det of P.b.

Week 6: Sets

Set = "containers", order does not matter
defined by what they contain

Defn: A set is a collection of objects.

An object of a set is called an element.
We write this as $a \in S$.

Examples

$$S = \{1, 2, 3, 4, 5\}$$

\{ ... \} in LaTeX

$$1 \in S, 0 \notin S, \pi \notin S$$

$$T = \{\alpha, 1, 3, 4, 5\} = S$$

$$T = S$$

$$\{1, \sqrt{2}\}, \{\sqrt{2}, \pi\}, \{\text{David, Jenny, Sarah}\}$$

$$\{1, 2, \dots, 10\} \quad \text{use "... to indicate some}$$

\dots vs ... part of

$$\{2, 4, \dots, 20\}$$

Common Sets

$$\mathbb{N} = \{1, 2, 3, \dots\}$$

$$\mathbb{Z} = \{\dots, -3, -1, 0, 1, 2, \dots\}$$

$\mathbb{Q}, \mathbb{R}, \mathbb{C}$ = complex #'s

$$\sqrt{2} \notin \mathbb{Q}, \sqrt{2} \in \mathbb{R}$$

$$\sqrt{-2} \notin \mathbb{R}$$

$$\mathbb{E} = \{\dots, -4, -2, 0, 2, 4, \dots\} = 2\mathbb{Z}$$

$d \in \mathbb{Z}_{>1}$, we define

$$d\mathbb{Z} = \{\text{"multiples of } d\}$$

More detail

$$= \{\dots, -2d, -d, 0, d, 2d, 3d, \dots\}$$

$$= \{dn : n \in \mathbb{Z}\}$$

$$= \{n : n \in \mathbb{Z} \mid d|n\}$$

$$= \{n \in \mathbb{Z} \text{ s.t. } d|n\}$$

General constructor:

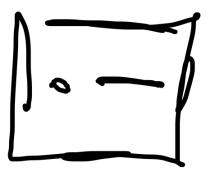
$\{\text{formula : parameters} \mid \text{conditions}\}$

" :" = " | " = "such that" = "s.t."

Example: $a, b \in \mathbb{R}$

$$[a, b] = \{x \in \mathbb{R} \text{ s.t. } a \leq x \leq b\}$$

$$[a, b) = \{x \in \mathbb{R} \text{ s.t. } a \leq x < b\}$$



A set can contain anything

Example:

$$\text{Fun}(\mathbb{R}, \mathbb{R}) = \left\{ \begin{array}{l} \text{functions from } \mathbb{R} \rightarrow \mathbb{R} \\ f : \mathbb{R} \rightarrow \mathbb{R} \end{array} \right\}$$

$$g(x) := x^2, \text{ then } g \in \text{Fun}(\mathbb{R}, \mathbb{R})$$

" $:$ " means "defining"

$$h(x) := \sqrt{x}$$
$$h \in \text{Fun}(\mathbb{R}, \mathbb{R})$$

$$\mathbb{P}[x] := \left\{ a_0 + a_1 x + \dots + a_n x^n : n \in \mathbb{Z}_{\geq 0} \text{ and each } a_i \in \mathbb{Q} \right\}$$

$$\mathbb{R}[x] := \left\{ \sum_{i=0}^n a_i x^i : n \in \mathbb{Z}_{\geq 0}, a_i \in \mathbb{R} \right\}$$



Sets can be elements of sets.

Analogy Big Amazon box containing
Many smaller boxes.

Examples:

$$T := \left\{ \left\{ 1, 2 \right\}, \left\{ 2 \right\}, \left\{ 3, 4 \right\} \right\}$$

T has 3 elements, not 4

$$\left\{ 1 \right\} \in T, \left\{ 2 \right\} \in T, \left\{ 3, 4 \right\} \in T$$

$$\left\{ 1, 2 \right\} \notin T$$

$$S = \left\{ \left\{ 2 \right\}, 3 \right\}$$

$$3 \in S$$

$$\left\{ 2 \right\} \in S$$

$$2 \notin S$$

$$2 \neq \left\{ 2 \right\}$$

" \in " is not transitive
 $x \in y \wedge y \in z \not\Rightarrow x \in z$

$$R = \left\{ 1, \left\{ 1 \right\} \right\}$$

$$1 \in R$$

$$\left\{ 1 \right\} \in R$$

Empty set \rightarrow "empty box"

Defn: The empty set \emptyset is the

set with the property that

$\forall x. x \notin \emptyset$. ($T \cdot E$ · $x \in \emptyset$ is

always false.)

Defn: We say that 2 sets S and T are equal if $x \in S \iff x \in T$.

(i.e., S & T have the same elements)

Ex. $\{1, 2, 3\} = \{2, 1, 3\}$

$\{1, 2, 3\} \neq \{2, 3\}$ b/c $1 \in \{1, 2, 3\}$ and $1 \notin \{2, 3\}$

$$\mathbb{Z} \neq \mathbb{Q}$$

b/c $\frac{1}{2} \in \mathbb{Q}$ but $\frac{1}{2} \notin \mathbb{Z}$.

Defn: let S and T be sets. we say that S is a subset of T if $x \in S \Rightarrow x \in T$. In this case we write $S \subseteq T$. (OR $S \subset T$)

(Equivalently: $\forall x \in S, x \in T$)

To show $S \not\subseteq T$

Find $x \in S$ s.t. $x \notin T$

$\exists x \in S$ s.t. $x \notin T$

Example: $\{1, 2\} \subseteq \{1, 2, 3\}$

$$\begin{matrix} \{1, 2, 3\} \\ \Downarrow \\ 3 \end{matrix} \neq \begin{matrix} \{1, 2\} \\ \Updownarrow \\ 3 \end{matrix}$$

Remarks: $S = T \iff$

$S \subseteq T$ and $T \subseteq S$

$$\begin{array}{ccccccc} \mathbb{N} & \subseteq & \mathbb{Z} & \subseteq & \mathbb{Q} & \subseteq & \mathbb{R} \\ \text{not } = & & \text{not } = & & \text{not } = & & \text{not } = \\ -1 & & -1 & & \sqrt{2} & & \pi \end{array}$$

Proofs w/ sets

$$P \Rightarrow Q$$

Recall " $A \subseteq B$ " means $x \in A \Rightarrow x \in B$



An implication

Start by "assuming the assumption"

- ① "Assume $x \in A$ "
- ② Write out what " $x \in A$ " means
(IE write out the defn)
- ③ "Argue" or "do calculations"



- ④ Conclude that $x \in B$.

has some defn +

in step 3, you verify this

$$d\mathbb{Z} = \{n : n \in \mathbb{Z} \mid d|n\}$$

Prove or disprove:

$$(i) 6\mathbb{Z} \subseteq 2\mathbb{Z}$$

$$(ii) 2\mathbb{Z} \subseteq 6\mathbb{Z}$$

Proof: (ii) This is false b/c $2 \in 2\mathbb{Z}$,
but $2 \notin 6\mathbb{Z}$ (b/c $6 \nmid 2$).

(i) Let $x \in 6\mathbb{Z}$. Then $x \in \mathbb{Z}$ and $6|x$.

Since $2|6$, by transitivity, $2|x$. Thus

$$x \in 2\mathbb{Z}, \quad \square$$

$$A = \{4^n - 1 : n \in \mathbb{Z}_{\geq 0}\} = \{0, 3, 15, 63, \dots\}$$

$$B = 3\mathbb{Z}_{\geq 0}$$

$$:= \{n \in \mathbb{Z}_{\geq 0} \text{ s.t. } 3|n\}$$

We know from week 2 that $3|4^n - 1$. $\forall x \in A \subseteq B$

Claim: $A \subseteq B$.

Proof: Let $x \in A$. Then $\exists n \in \mathbb{Z}_{\geq 0}$ s.t. $x = 4^n - 1$,

By week 2, $3|4^n - 1$. Thus $4^n - 1 \in 3\mathbb{Z}$. \blacksquare

Converse? Is $B \subseteq A$? NO!

$6 \in B$ but $6 \notin A$.

Lemma: Let A, B, C be sets.

Suppose that $A \subseteq B$ and $B \subseteq C$.

Then $A \subseteq C$.

$$x \in A \rightarrow x \in C$$

Proof: Assume $A \subseteq B$ and $B \subseteq C$.

Let $x \in A$. Since $A \subseteq B$, $x \in B$. Since

$x \in B$, and $B \subseteq C$, $x \in C$. \square

\emptyset is the set s.t. " $x \in \emptyset$ " is false $\forall x$.

Claim: If set A , $\emptyset \subseteq A$,

Proof: "There is nothing to check" \square

Is every $x \in \emptyset$ also $x \in A$? Yes...

Contradiction: Suppose $\emptyset \not\subseteq A$.

($\emptyset \subseteq A$ means)
 $x \in \emptyset \Rightarrow x \in A$

It suppose that $\exists x \in \emptyset$ s.t. $x \notin A$.

Since $x \in \emptyset$ is always false, we found a contradiction.

You can't disprove $\emptyset \subseteq A$.

Contrapositive: $x \notin A \Rightarrow x \notin \emptyset$.

Suppose $x \notin A$. Well.... $x \notin \emptyset$ is true. \square

Week 7: More parts of sets

$A \subseteq B$ means $x \in A \Rightarrow x \in B$

$\forall x \in A, x \in B$

"let $x \in A$

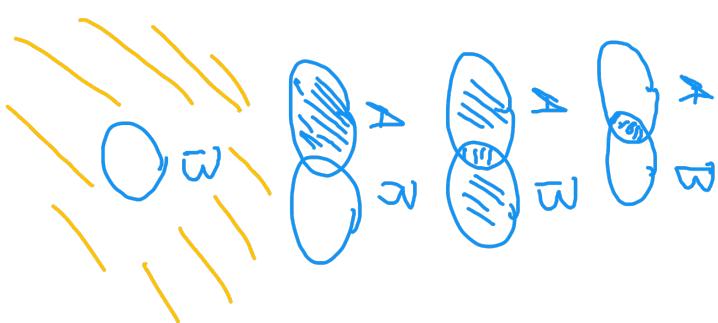
Thus $x \in B$."

$$A \cap B = \{x : x \in A \text{ and } x \in B\}$$

$$A \cup B = \{x : x \in A \text{ or } x \in B\}$$

$$A - B = \{x \in A \mid x \notin B\}$$

$$\overline{B} = \{x : x \notin B\}$$



7

$U = "everything"$

Prove or disprove:

(i) $6\mathbb{Z} = 2\mathbb{Z} \cup 3\mathbb{Z}$ F

(ii) $6\mathbb{Z} = 2\mathbb{Z} \cap 3\mathbb{Z}$ T

$$d\mathbb{Z} = \{n \in \mathbb{Z} \text{ s.t. } d|n\}$$

Proof

(i) This is false: $2 \in 2\mathbb{Z} \cup 3\mathbb{Z}$, but $2 \notin 6\mathbb{Z}$.

(ii) " $=$ " is " \leq " and " \geq "

" \leq " Let $x \in 6\mathbb{Z}$. Then $x \in \mathbb{Z}$ and $6|x$.

(WTS: $x \in 2\mathbb{Z} \cap 3\mathbb{Z}$, i.e., $x \in 2\mathbb{Z}$ and $x \in 3\mathbb{Z}$, i.e., $x \in \mathbb{Z}$, $2|x$ and $3|x$)

Since $2|6$ and $3|6$, by transitivity of division, $2|x$ and $3|x$.

Thus $x \in 2\mathbb{Z}$ and $x \in 3\mathbb{Z}$, so $x \in 2\mathbb{Z} \cap 3\mathbb{Z}$.

" \geq " Let $x \in 2\mathbb{Z} \cap 3\mathbb{Z}$. Then $x \in 2\mathbb{Z}$ and $x \in 3\mathbb{Z}$. Then
 $x \in \mathbb{Z}$ and $2|x$ and $3|x$.

(WTS: $x \in 6\mathbb{Z}$, i.e., $6|x$).

Since $\gcd(2, 3) = 1$, $2 \cdot 3|x$. Thus $x \in 6\mathbb{Z}$. \square

$$P \Rightarrow (Q \Rightarrow R)$$

negative

$$P \Rightarrow Q$$

$$x \in A - C \Rightarrow x \in A - B$$

$$B \subseteq C \wedge \exists x \in A - C \text{ s.t.}$$

claim: $(B \subseteq C) \Rightarrow (A - C \subseteq A - B)$

Proof: Let $x \in B$.

\downarrow

$\#$ wrong answer

Assume $B \subseteq C$. Let $x \in A - C$. Then $x \in A$ and $x \notin C$.

(WTS: $x \in A - B$, i.e., $x \in A$ and $x \notin B$).

Proceed by contradiction. Assume $x \in B$.

Since $B \subseteq C$, $x \in C$. This contradicts $x \notin C$.

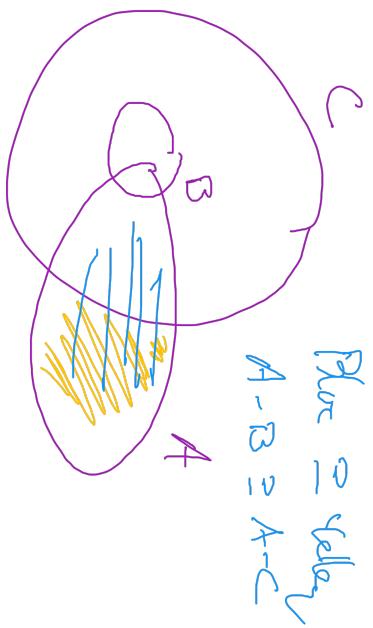
thus $x \notin B$. thus $x \in A - B$. \square

$$(x \in B \Rightarrow x \in C)$$

$$B \subseteq C, x \notin B.$$

Note: The contrapositive of " $B \subseteq C$ " is

$$x \notin C \Rightarrow x \notin B$$



$$\text{Plot} \cong \text{yellow}$$

$$A - B \cong A - C$$

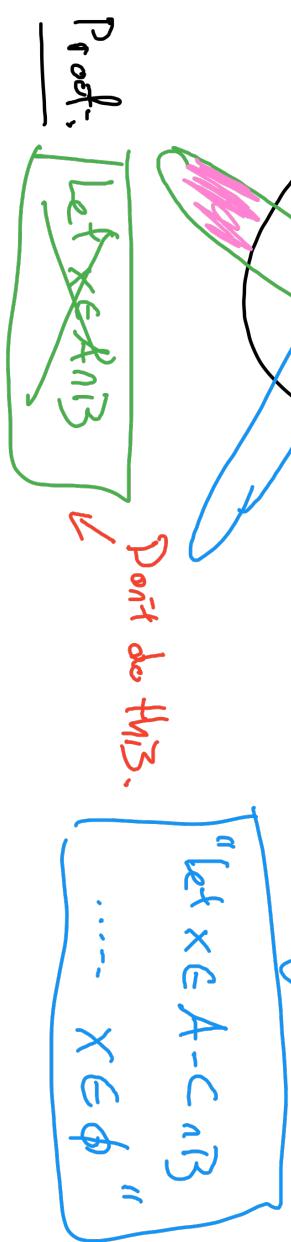
$$P \Rightarrow Q$$

Claim: $A \cap B \subseteq C \Rightarrow (A - C) \cap B = \emptyset$

To prove " $D = \emptyset$ ", do proof by contradiction.

" $\emptyset \neq \emptyset$ ", i.e. $\exists x \in D$.

Pink dog not in set B
For, but added



Suppose $A \cap B \subseteq C$. Proceed by contradiction. Assume $(A - C) \cap B \neq \emptyset$.

Thus $\exists x \in (A - C) \cap B$. Then $x \in A - C$ and $x \in B$, then $x \in A$ and $x \notin C$.

thus $x \in A \cap B$, so since $A \cap B \subseteq C$, $x \in C$, this is a contradiction,
thus $A - C \cap B = \emptyset$. \square

De Morgan's Laws

$$\overline{C} = \{x : x \notin C\}$$

$$\begin{aligned}
 \text{(a)} \quad \overline{A \cap B} &= \overline{A} \cup \overline{B} \\
 \overline{A \cup B} &= \overline{A} \cap \overline{B}
 \end{aligned}$$



Proof: " \subseteq " let $x \in A \cap B$. Then $x \notin \overline{A \cap B}$.

(WTS $x \in \overline{A \cup \overline{B}}$ i.e.

$x \in \overline{A} \text{ or } x \in \overline{B}$ i.e.

$x \notin A \text{ or } x \notin \overline{B}$

$$= \neg(x \in A \text{ and } x \in \overline{B}) = \neg(x \in A \cap B)$$

Thus $x \notin A$ or $x \notin \overline{B}$. Thus $x \in \overline{A} \text{ or } x \in \overline{B}$.

Thus $x \in \overline{A} \cup \overline{B}$. \blacksquare

\$ A times B \$

Products + Power sets

Defn: Let A and B be sets. The Cartesian Product $A \times B$ is the set $\{ (ab) : a \in A \text{ and } b \in B \}$.

Let $n \in \mathbb{N}$. Then $A^n = \{ (a_1, \dots, a_n) \text{ s.t. } \forall i \in \{1, \dots, n\}, a_i \in A \}$
 $(A \times B \times C \rightarrow (a, b, c), a \in A, b \in B, c \in C)$

Order matters $A \times B \neq B \times A$
But A and B can be different

$$\textcircled{1} \quad \mathbb{R} \times \mathbb{R} = \mathbb{R}^2 \ni (x, y), x, y \in \mathbb{R}$$

$$\textcircled{2} \quad A = \{1, 2\}, B = \{3, 4\}$$

$$A \times B = \left\{ (1, 3), (1, 4), (2, 3), (2, 4) \right\} \quad \text{b/c } (1, 2) \notin A \times B$$

$$\textcircled{3} \quad (a, b) \in A \times B \iff (a, b) \in A \text{ and } b \in B$$

$a \in A$ and $b \in B$

$$\textcircled{4} \quad \{1, 2, 3\} \times \mathbb{Z} \rightarrow \{1, 2, 3\}$$

$$\textcircled{5} \quad \mathbb{R} \times \text{Fun}(\mathbb{R}, \mathbb{R}) \rightarrow C^{\pi, f}$$

$$\boxed{I \in \mathbb{Z}} \quad \boxed{\mathbb{Z} \ni I}$$

$$\begin{array}{c} \boxed{1} \\ \boxed{2} \\ \vdots \\ \boxed{n} \end{array}$$

$$\begin{array}{c} A \\ \cap \\ B \\ \supset \\ A \supset B \end{array}$$

Lemma: Suppose $A \subseteq B$ and $C \subseteq D$.

Then $A \times C \subseteq B \times D$.

Proof: Suppose $A \subseteq B$ and $C \subseteq D$. Let $(a, c) \in A \times C$.

Then $a \in A$ and $c \in C$.

($\because a \in B$ and $c \in D$)

Since $A \subseteq B$, $a \in B$. Since $C \subseteq D$, $c \in D$.

Thus $(a, c) \in B \times D$. \blacksquare

Power Set: Let A be a set. Then we define the powerset $P(A)$ to be

$$P(A) = \{B \text{ s.t. } B \subseteq A\}$$

Example: $A = \{1, 2\}$ $P(A) = \left\{ \begin{array}{l} \{1\}, \{2\}, \emptyset, \{1, 2\} \\ \{1\} \subseteq \{1, 2\} \\ \{2\} \subseteq \{1, 2\} \\ \emptyset \subseteq \{1, 2\} \\ \{1, 2\} \subseteq \{1, 2\} \end{array} \right\}$

Rule: $B \in P(A) \iff B \subseteq A$

$1 \notin P(\{1, 2\})$ b/c $1 \notin \{1, 2\}$

Claim: $\# P(A) = 2^{\#A}$

Ex: $A = \{1, 2\}$ $P(A) = \{\emptyset, \{1\}, \{2\}, \{1, 2\}\}$

Ex: $P(\emptyset) = \{B : B \subseteq \emptyset\}$
 $= \{\emptyset\}$

$\#\emptyset = 0$

$\#\{\emptyset\} = 1 = 2^0$

Ex: $\emptyset, A \in P(A)$ b/c $\emptyset \subseteq A$
 $A \subseteq A$

$P(\mathbb{Z}) = \left\{ \emptyset, \mathbb{Z}, \mathbb{Z}, \mathbb{Z}, \mathbb{Z}, \dots \right\}$
 $\left\{ \begin{array}{l} \{1\}, \{2\}, \{3\}, \dots \\ \{1, 2\}, \{1, 2, 3\}, \dots \end{array} \right\}$

$\mathbb{Z} \in P(\mathbb{Z})$

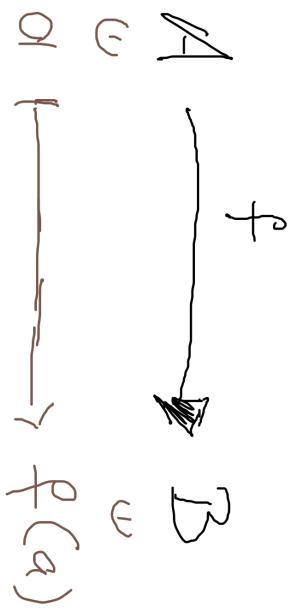
$\mathbb{Z} \in P(\mathbb{Z})$

$\mathbb{R} \notin P(\mathbb{Z})$ b/c $\mathbb{R} \notin \mathbb{Z}$

Week 9: Functions, injectivity, & surjectivity

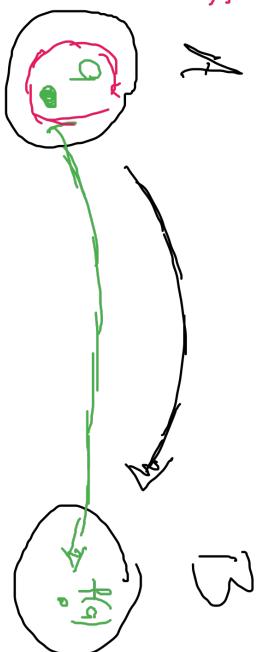
Let A, B be sets.

A function is a "rule" that associates, to each $a \in A$, some $b \in B$



Can't graph..

$$w \subseteq A$$



Examples

$$f: \mathbb{R}^2 \rightarrow \mathbb{R}^2$$

$$(x_1, y_1) \mapsto (x_1 + y_1, x_1)$$

$$f(x, y) = (x+y, x)$$



can't graph

"unambiguous" := $\forall a \in A, \exists$ exactly one output $f(a) \in B$.

"not multivalued"

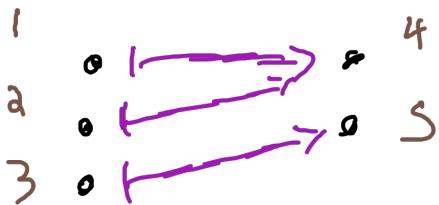
"passes the VLT"

If $A=B=\mathbb{R}$

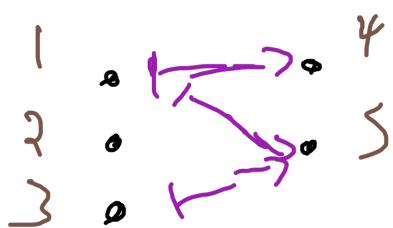
$$\begin{aligned} f(0,0) &= (0,0) \\ f(1,1) &= (d,1) \end{aligned}$$

Let go of the formulas

$$A = \{1, 2, 3\}, B = \{4, 5\}$$

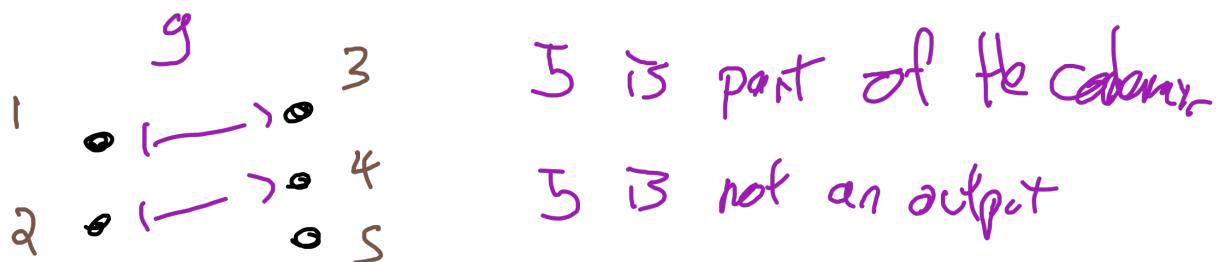
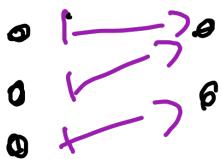


$$f(1) = 4 \quad f(2) = 4 \quad f(3) = 5$$



2 problems
Ambig vars (what is $f(1)$?
Didn't define $f(2)$...

From far away ... just dots



Domain

codomain

range \subset image

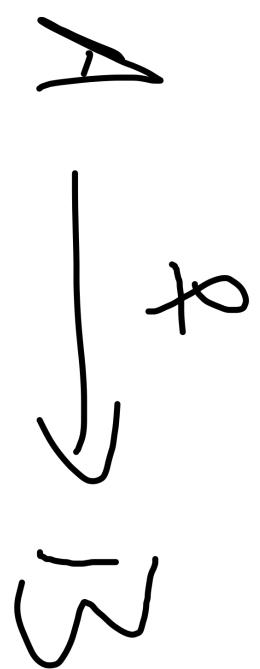
"
inputs

Potential outputs

Actual outputs

$A = \{x \mid x \text{ is a student in math class}\}$

$B = \{y \mid y \text{ is yes, no}\}$



$f(x) = \text{Answer to "Is } x \text{ wearing glasses?"}$

$f(\text{Angela}) = \text{no}$

$f(\text{Tucker}) = \text{no}$

$$\mathbb{Z} \xrightarrow{g} \mathbb{Z}$$

$$x \mapsto \begin{cases} x/2 & \text{if } x \in \mathbb{E} \\ 3x+1 & \text{otherwise} \end{cases}$$

$$g(1) = 4$$

$$g(2) = 1$$

↑
other wise

Caution

$$\mathbb{Z} \xrightarrow{h} \mathbb{Z}$$

$$x \mapsto x/2$$

invalid b/c $x/2 \notin \mathbb{Z}$

if $x=1$

$$\mathbb{E} \rightarrow \mathbb{Z} \quad \underline{\text{ok}}$$

Indicator fn of G

$$\mathbb{R} \xrightarrow{f} \mathbb{R}$$

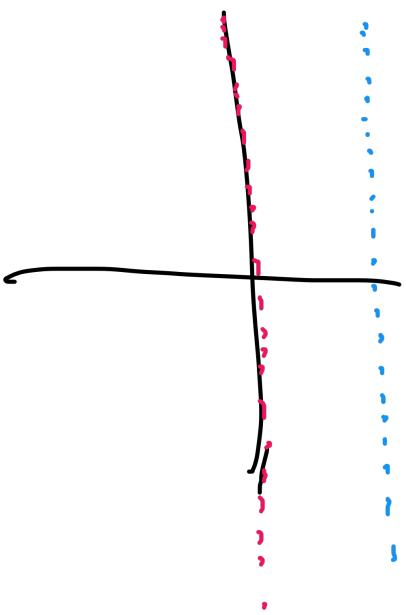
$$x \mapsto \begin{cases} 0 & \text{if } x \notin Q \\ 1 & \text{if } x \in Q \end{cases}$$

On ambiguous,

$$g(0) = 1 \quad g(\sqrt{d}) = 0 \quad g(\tan 1) = ?$$

$$g(\alpha/\pi) = 1 \quad g(\pi) = 0$$

$$(x^5 + 7x + 1 = 0)$$



What does it mean for $f = g$?

$$f: A \rightarrow B$$

$$g: C \rightarrow D$$

Def'n: We say that $f = g$ if

$$\forall a \in A, f(a) = g(a)$$

$$\forall a, b \in A, f(a) = f(b) \quad \text{if } a = b$$

↳ *constant*
↳ *unambiguity*

To show $f \neq g$, show $A \neq C$, $B \neq D$, or
 $\exists a \in A, f(a) \neq g(a)$.

Examples:

$$\mathbb{R} \xrightarrow{f} \mathbb{R}$$

$$x \mapsto 2x$$

$$\mathbb{Z} \xrightarrow{g} \mathbb{Z}$$

$$x \mapsto 2x$$

$$\mathbb{Z} \xrightarrow{h} \mathbb{E}$$

$$x \mapsto 2x$$

$$\mathbb{Z} \xrightarrow{h_2} \mathbb{E}$$

$$n \mapsto 2n$$

$h = h_2$

$$\begin{matrix} 1 & \xrightarrow{f} & 2 \\ & \bullet \longmapsto \bullet & \\ & & 3 \end{matrix}$$

$f \neq g$ b/c different domains
and codomains.

$g(\sqrt{2})$ is undefined w/c

$$\begin{matrix} \sqrt{2} \notin \mathbb{Z} \\ h \neq g \end{matrix}$$

b/c different codomains.

$$\forall b, h(b) = h_2(b)$$

$$\begin{matrix} \text{||} & \text{||} \\ 2b & 2b \end{matrix}$$

$$\begin{matrix} 1 & 2 \\ & \bullet \xrightarrow{g} \bullet \\ & 3 \end{matrix}$$

$f \neq g$

$$f(1) \neq g(1)$$

$$\begin{matrix} \text{||} & \text{||} \\ 2 & 3 \end{matrix}$$

$$f(n) = n+1$$

$$f(n) = 2n$$

involuted

\downarrow

$$S \rightarrow S^{\text{inv}} \quad P(B) = \{A \subseteq B\}$$

$$P(\mathbb{Z}) \longrightarrow P(\mathbb{Z}) \quad A \in P(B) \iff A \subseteq B$$

$$S \xrightarrow{f} S \cup \{1\} = f(S)$$

$$S \xrightarrow{g} S \cap E$$

$$S \xrightarrow{h} S \cup \{\pi\}$$

h is involuted

$$\text{blk } S \cup \{\pi\} \notin P(\mathbb{Z})$$

$$f(E) = E \cup \{3\}$$

$$f(\mathbb{Z}) = \mathbb{Z} \cup \{3\} = \mathbb{Z}$$

$$f(\{1,2,3\}) = \{1,2,3\} \cup \{3\} = \{1,2,3\}$$

$$f(\{2,3\}) = \{2,3\} \cup \{E\} = \{2,3\}$$

$$f(\mathbb{Z}) = \mathbb{Z} \cup E = E$$

$$\mathbb{R} \rightarrow P(\mathbb{R})$$

$x \rightarrow (\infty, \infty)$

X
—
↓
X
—

$$(x_{100}) = \left\{ a \in \mathbb{R} \text{ s.t. } x > a \right\}$$

X

λ — $y + \ln u$ — $y + P(R)$

Common functions

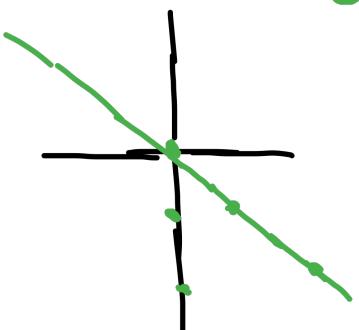
$$A \xrightarrow{\text{id}_A} A \quad \text{"do nothing"}$$

$$x \mapsto x$$

$$\text{id}_{\mathbb{R}}(x) = x$$

$$\text{id}(x) = x$$

$$A = \mathbb{R}$$

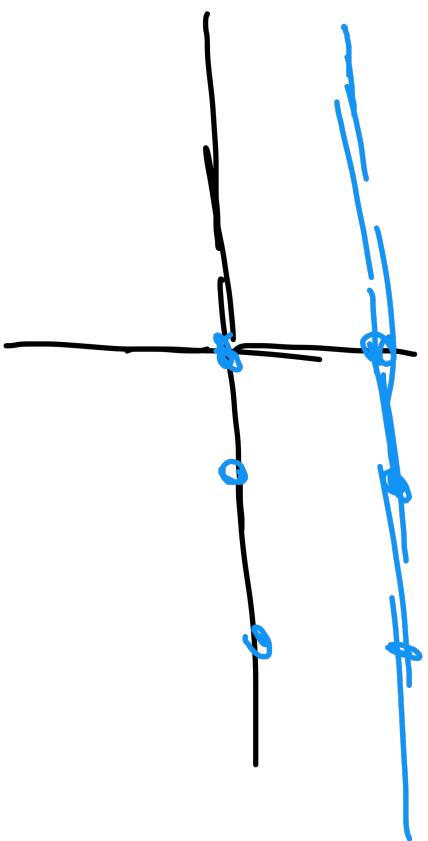


Sps $B \neq \emptyset$ and let $b \in B$.

$$A \xrightarrow{c_b} B \quad A = B = \mathbb{R} \quad b = 1$$

$$a \mapsto b$$

$$c_b(a) = b$$



Defn. Let A and B be sets and $f: A \rightarrow B$ be a function.

The image (range) of f is

(write as $\text{im } f$ or $f(A)$)

$$\text{im } f = \{f(a) : a \in A\}$$

If $W \subseteq A$, define

$$f(W) = \{f(a) : a \in W\} \quad (\text{onto})$$

We say that f is surjective if

$$f(A) = \text{im } f = B \quad (\text{IE, "f takes every possible value"})$$

$a \in A, f(a) \in B$ elements

$f(A) \subseteq B$ a set

A
and an
elt

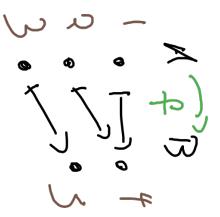
$f(w) \in$

$f(A) \subseteq B$

$\{f(a) : a \in A\}$

To prove $f(A) = B$,
only need to prove
 $B \subseteq f(A)$.

$f: B \rightarrow S^3$.



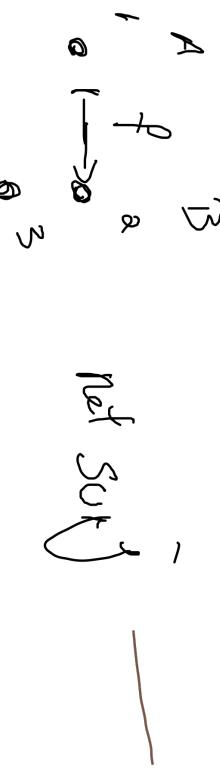
$$f(A) = \{ f(a) : a \in A \}$$

$$= \{ f(a) : a \in \{1, 2, 3\} \}$$

$$= \{ f(1), f(2), f(3) \} = \{ 4, 5, 6 \} = \{ 4, 5 \} = B$$

$$w = \{1, 2\}$$

$$f(w) = \{ f(a) : a \in w \} = \{ f(1), f(2) \} = \{ 4, 5 \} = \{ 4 \} = B$$



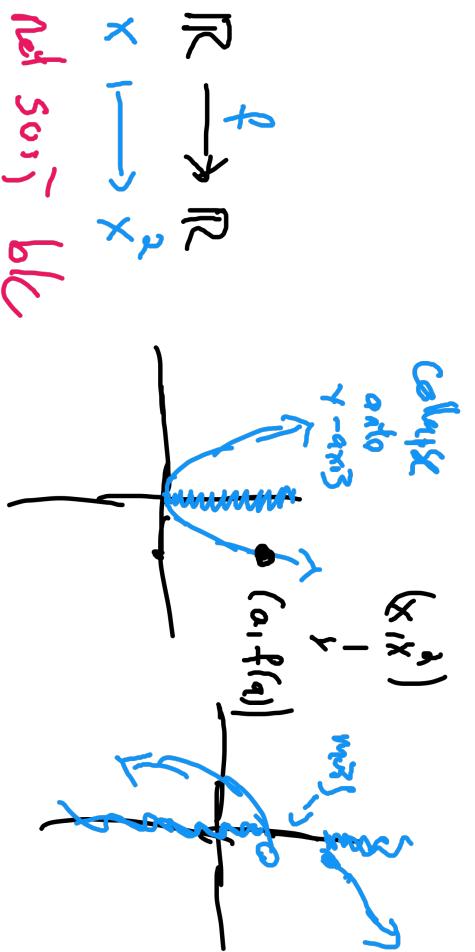
$$f(A) = \{ f(a) \} = \{ 2 \} \neq \{ 1, 3 \} = B.$$

the defn of $b \in f(A)$

$$B \neq f(A) \text{ b/c } 3 \in B, \text{ but } 3 \notin f(A)$$

To prove $f(A) = B$, need to show $\forall b \in B, \exists a \in A \text{ s.t. } f(a) = b$

To prove $f(A) \neq B$, need to show $\exists b \in B \text{ s.t. } \forall a \in A, f(a) \neq b$



$\mathbb{R} \xrightarrow{f} \mathbb{R}$

$$x \mapsto x^2$$

not so nice

-1 & in f

TE, $\forall a \in \mathbb{R}$ s.t. $f(a) = -1$

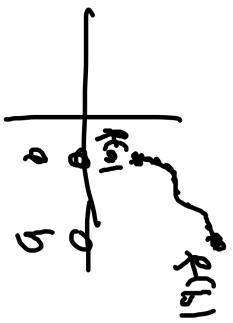
TE $\exists a \in \mathbb{R}$ s.t. $a^2 = -1$

$b/c \nexists \forall a \in \mathbb{R}$

$f(\mathbb{R}) = \text{im } f = \mathbb{R}_{\geq 0}$

(by "continuity")

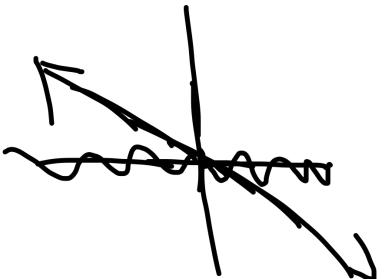
$f(0) = 0$ AND $\lim_{x \rightarrow \infty} f(x) = \infty$



$$\mathbb{R} \xrightarrow{f} \mathbb{R}$$

$$x \mapsto dx + 1$$

$$\inf f = \mathbb{R}$$



Pf: $\exists \sqrt{\epsilon}$, use continuity + limits.

Pf2: claim: $f(\mathbb{R}) = \mathbb{R}$.

Automatically: $f(\mathbb{R}) \subseteq \mathbb{R}$.

Goal: $\mathbb{R} \subseteq f(\mathbb{R})$.

Let $b \in \mathbb{R}$. (wts $b \in f(\mathbb{R})$).

(Need $\exists a \in \mathbb{R}$ s.t. $f(a) = b$, i.e.,
 $a = \frac{b-1}{2}$)

let $a = \frac{b-1}{2}$. Then $a \in \mathbb{R}$ and $f(a) = 2\left(\frac{b-1}{2}\right) + 1 = b-1+1 = b$.

Thus $b \in f(\mathbb{R})$. \square

Week 10, preimages

$$f: A \rightarrow B$$

$$w \subseteq A$$

$$f(w) = \{ f(a) : a \in w \}$$

element

$$f(w) \subseteq B$$

set

$$\forall b \in f(w) \exists a \in w \text{ such that } f(a) = b$$

Useful: $\forall a \in w \mid f(a) \in f(w)$

$$P(A) \longrightarrow P(B)$$

$$w \longmapsto f(w)$$

"Pre image" or "inverse image"

Defn: Let $A \xrightarrow{f} B$ be a fn. Let $w \in B$. We define the preimage of w under f to be

$$f^{-1}(w) = \{a \in A \text{ s.t. } f(a) \in w\}$$

Caution!

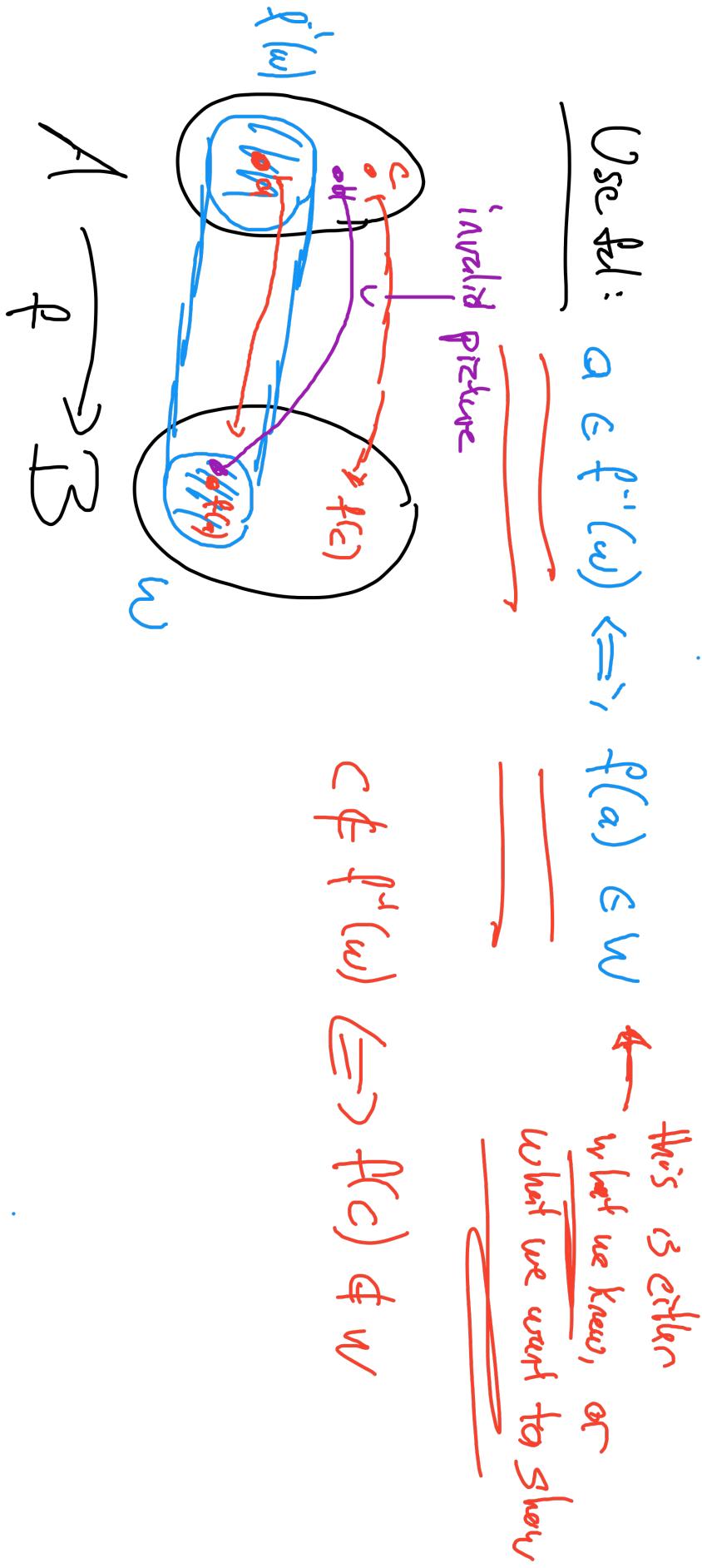
THIS IS NOT RELATED TO THE INVERSE FUNCTION

Example: $f: \mathbb{R} \rightarrow \mathbb{R}_{\geq 0}$ $x \mapsto x^2$

$W = \mathbb{R}_{\geq 0}$ $f^{-1}(w) = \{a \in \mathbb{R} \text{ s.t. } f(a) \in \mathbb{R}_{\geq 0}\}$
= \mathbb{R}

$$W = [1, 4]$$

$$f^{-1}(w) = [1, 2] \cup [-2, -1]$$



$x \xrightarrow{f} B$ $w = B$?

Comment: $f^{-1}(B) = \{a \in A \text{ s.t. } f(a) \in B\} = A$

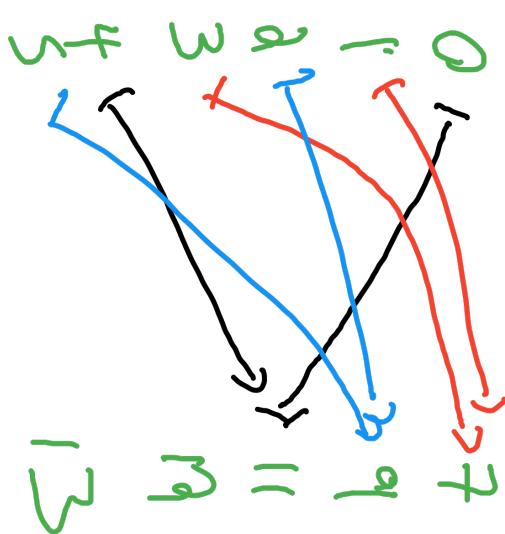
Always true

(By the defn of f^{-1})
Co-domain

Examples:

$$A \xrightarrow{f} B$$

$$\begin{aligned} f^{-1}(\{q, 1\}) &= \{a, 1, 3\} \\ &= \{a, 4, 2, 5\} \end{aligned}$$



$$\begin{aligned} f^{-1}(\{a, b\}) &= \{1, 3\} \\ f^{-1}(\{c\}) &= \{3\} \end{aligned}$$

$$f = \left(\begin{matrix} 1 & 2 & 3 \\ a & b & c \end{matrix} \right)$$

$$\begin{aligned} f(0) &= f(4) = b \\ f(1) &= f(2) = a \\ f(3) &= f(5) = d \end{aligned}$$

Example: $A \xrightarrow{f} B$

$$w = \phi \in B$$

$$f^{-1}(\phi) = \{ a \in A \text{ s.t. }$$

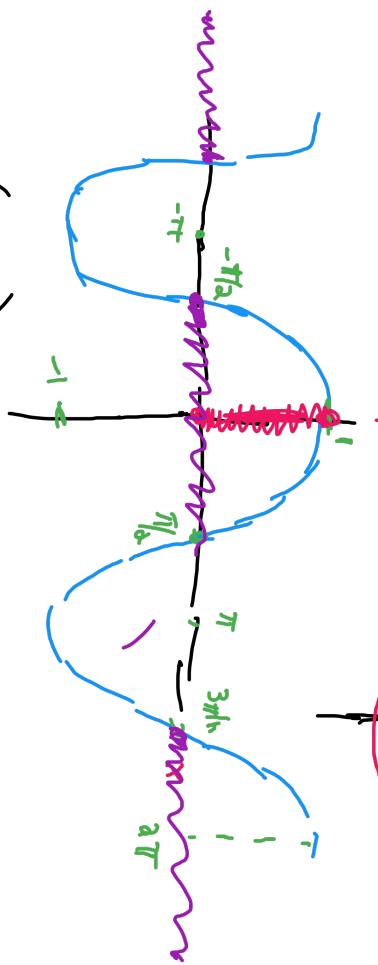
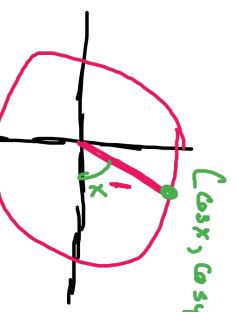
$f(a) \in \phi$

Always False

$$= \emptyset$$

Example: $f: \mathbb{R} \xrightarrow{\cos} \mathbb{R}$

$$x \mapsto \cos x$$



$$f^{-1}([0, 1]) = \left\{ a \in \mathbb{R} \text{ s.t. } \cos a \in [0, 1] \right\}$$

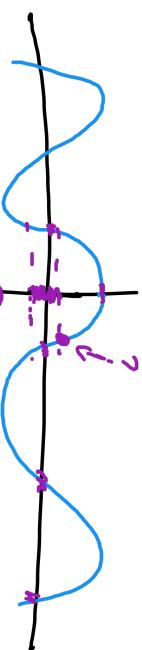
$$\dots [-5\pi/2, 3\pi/2] \cup [-\pi/2, \pi/2] \cup [3\pi/2, 5\pi/2] \cup [\pi/2, 7\pi/2] \dots$$

$$(\Sigma) \quad \bigcup_{n=-\infty}^{\infty} [-\pi/2 + n\pi, \pi/2 + n\pi]$$

$$n = -\infty$$

$$f^{-1}(\mathbb{R}_{\geq 0}) = f^{-1}([0, 1])$$

$$\begin{aligned} & \text{solve } \cos a = 0.17 \\ & a = \cos^{-1} 0.17 \end{aligned}$$



$$f^{-1}([-1, 1]) = \mathbb{R}$$

$$f^{-1}([2, 3]) = \emptyset$$

$$f^{-1}([-0.23, 0.17]) = \text{bunch of intervals}$$

(and no nice formula)

$$\mathbb{R} \xrightarrow{x^2} \mathbb{R} \quad f^{-1}([1,4]) = [1,2] \cup [2,3]$$

$$[a,b] = \{x \in \mathbb{R} \text{ s.t. } a \leq x \leq b\}$$

Proof: " \supseteq " Let $x \in [1,2] \cup [-2,-1]$. Then $x \in [1,2]$ or $x \in [-2,-1]$

Then $1 \leq x \leq 2$ or $-2 \leq x \leq -1$.

(WTS: $x \in f^{-1}([1,4])$ i.e., $f(x) \in [1,4]$ i.e., $1 \leq x^2 \leq 4$)

Case 1: $1 \leq x \leq 2$. Then Squaring gives $1 \leq x^2 \leq 4$. Thus $x^2 = f(x) \in [1,4]$, thus $x \in f^{-1}([1,4])$.

Case 2: $-2 \leq x \leq -1$. Then Squaring gives $1 \leq x^2 \leq 4$. Thus $x^2 = f(x) \in [1,4]$.

Thus $x \in f^{-1}([1,4])$.

" \subseteq " Let $x \in f^{-1}([1,4])$. Then $f(x) = x^2 \in [1,4]$. Then $1 \leq x^2 \leq 4$.

Then $1 \leq x \leq 2$ or $-2 \leq x \leq -1$. Thus $x \in [1,2]$ or $x \in [-2,-1]$

Thus $x \in [1,2] \cup [-2,-1]$ \square

18. For the following functions, compute the *inverse* image of the given subsets of the codomain. (No proofs are necessary.)

- $f: \mathbf{Z} \rightarrow \mathbf{Z}, f(n) = 3n + 1; W_1 = E, \text{ the set of even integers, } W_2 = \{4\}, W_3 = \{1, 5, 8\}$
- $f: \mathbf{R} \rightarrow \mathbf{R}, f(x) = 3x + 1; W_1 = \{4\}, W_2 = \{1, 5, 8\}, W_3 = (4, \infty)$,
 $W_4 = (2, 4), W_5 = \mathbf{Z}, W_6 = E, \text{ the set of even integers}$
- $f: \mathbf{R} \rightarrow \mathbf{R}, f(x) = \cos x; W_1 = [-1, 1], W_2 = \{x \in \mathbf{R} \mid x \geq 0\}, W_3 = \mathbf{Z}$
- $f: \mathbf{R} \rightarrow \mathbf{R}, f(x) = e^x; W_1 = [-1, 0], W_2 = \{x \in \mathbf{R} \mid x \geq 0\}, W_3 = \{1\}$
- $f: \mathbf{Z} \rightarrow \mathbf{Z}, f(n) = \begin{cases} n & \text{if } n \text{ is even} \\ n - 1 & \text{if } n \text{ is odd} \end{cases}; W_1 = E, W_2 = \{1\}, W_3 = \{6\}, W_4 = \mathbf{O}, \text{ the set of odd integers}$

$$(a) f^{-1}(E) = \emptyset$$

" \supseteq " Let $a \in \mathbb{D}$. (wts $a \in f^{-1}(E)$. $\exists b \in E, f(b) \in E$)

Then $f(a) = 3a + 1$. Since a is odd, $3a$ is odd,

so $3a + 1$ is even. Thus $f(a) \in E$, thus $a \in f^{-1}(E)$.

" \subseteq " Let $a \in f^{-1}(E)$. Then $f(a) \in E$. Thus $3a + 1$ is even.

Thus $3a$ is odd, so $a \in \mathbb{D}$. \square

$$e) f(n) = \sum_{n=1}^{\infty} n \quad n \in E$$

$$\begin{aligned} f(a) &= 0 & f(0) &= 0 \\ f(1) &= 0 & f(3) &= 0 \end{aligned}$$

$$f^{-1}\left(\{1, 3\}\right) = \emptyset$$

$$f^{-1}\left(\{6, 7\}\right) = \{6, 7\}$$

$$f^{-1}(E) = \emptyset$$

$$f^{-1}(x \cup y) = f^{-1}(x) \cap f^{-1}(y)$$

Abstract proof: $f: A \rightarrow B$

$$x, y \subseteq B$$

$$f^{-1}(x \cup y) \subseteq f^{-1}(x) \cup f^{-1}(y)$$

Proof: Let $a \in f^{-1}(x \cup y)$. Then $f(a) \in x \cup y$.

$$\text{Then } f(a) \in x \text{ or } f(a) \in y.$$

WTS: $a \in f^{-1}(x) \cup f^{-1}(y)$. If $a \in f^{-1}(x)$ or $a \in f^{-1}(y)$

"Let $a \in f^{-1}(x \cup y)$ then $f(a) \in x \cup y$.
 Then $f(a) \in x$ and
 $f(a) \in y$.
 Then $a \in f^{-1}(x)$ and $a \in f^{-1}(y)$.
 Then $a \in f^{-1}(x) \cap f^{-1}(y)$.

$$\vdash f(a) \in x \text{ or } f(a) \in y$$

Thus $a \in f^{-1}(x)$ or $a \in f^{-1}(y)$.

thus $a \in f^{-1}(x) \cup f^{-1}(y)$

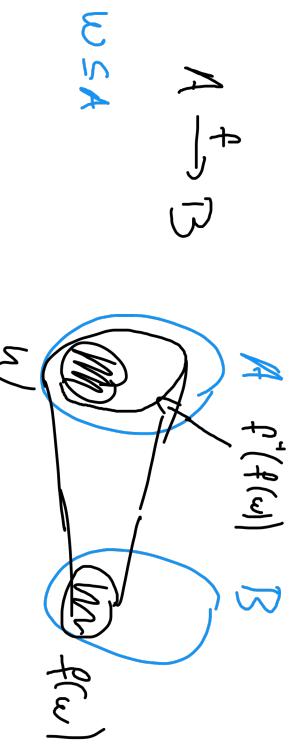
every step is reversible i.e., an "if and only if" (\iff)

TE de the same proof brackets

Switch "or" with "and" in prev. proof

(and $\cup \cup$)

$A \xrightarrow{f} B$



$$R^x \xrightarrow{x} R^x \quad f^{-1}(f(w)) = f^{-1}(R_{20}) = R$$

$w = R_{20} \quad f(w) = R_{20}$



claim $w \subseteq f^{-1}(f(w))$

Let $a \in w$. (i.e. $a \in f^{-1}(f(w))$) $\exists E$

Then $f(a) \in f(w)$ by def of image.

Then $a \in f^{-1}(f(w))$, \blacksquare
(def of pre-im.)

$a \in w$ \downarrow
 \downarrow $\circ f$
 $f(a) \in f(w)$

$f^{-1}(f(w)) \subseteq w \circ f \circ f^{-1}$

$R \subseteq R_{20}$

for $w = R_{20} \quad f(x) = x$

Week 4: Injective functions

or "one-to-one"

Defn: we say that a function $f: A \rightarrow B$ is injective if

$$\forall a, b \in A, a \neq b \Rightarrow f(a) \neq f(b)$$

"Slogan": "Distinct inputs give distinct outputs"

Cont' rep' position: $\forall a, b \in A, f(a) = f(b) \Rightarrow a = b$

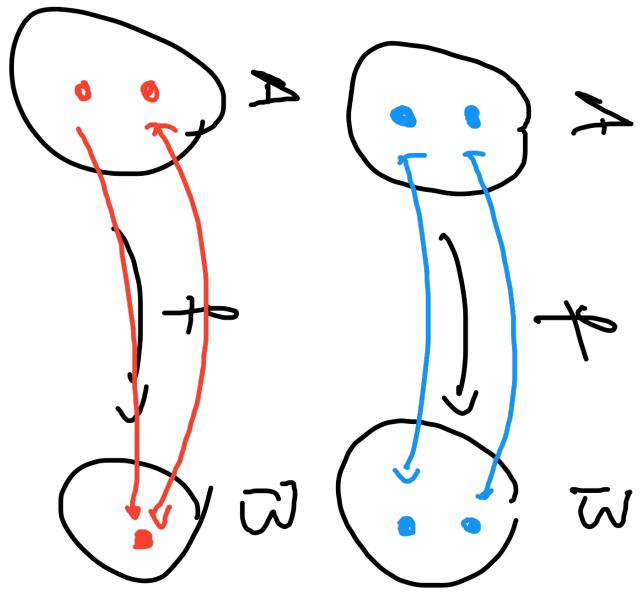
(often easier to do proofs with)

Repetition: $\exists a, b \in A$ s.t. $a \neq b$ AND $f(a) = f(b)$

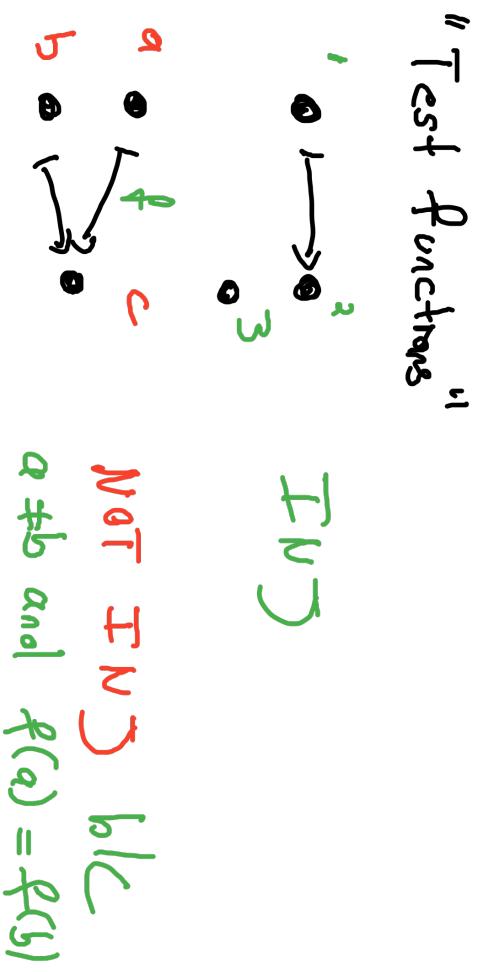
#1 Mistake: $\forall a, b \in A, a = b \Rightarrow f(a) = f(b)$

wrong

just.... the dom of a f'm



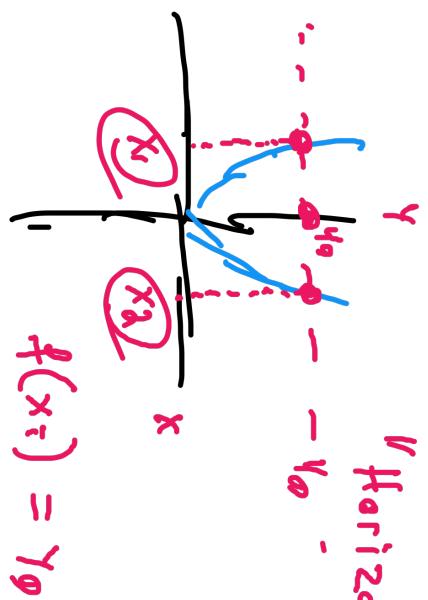
f
 \bar{f}_{ij}



Examples:

$$g: \mathbb{R} \rightarrow \mathbb{R}$$

$$x \mapsto x^q$$



$$\mathbb{R} \rightarrow \mathbb{R}$$

"Horizon. line test":

At L with the graph, if
at most one point

f is inj.

NOT INJ

$$1 \neq -1 \text{ and } g(1) = g(-1)$$

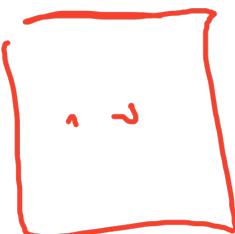
$$1 \neq -1 \quad g(1) = g(-1)$$

Note: To prove a function $\boxed{f: A \rightarrow B}$ is inj. , need an argument,

" $\forall a, b \in A$, $a \neq b \Rightarrow f(a) \neq f(b)$ "

" $\forall a, b \in A$, $f(a) = f(b) \Rightarrow a = b$ "

"Let $a, b \in A$. Assume $f(a) = f(b)$.



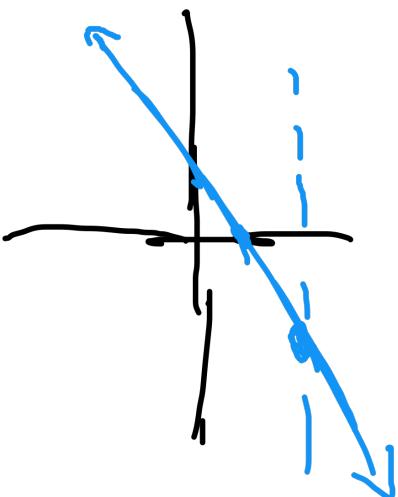
... we conclude "
that $a = b$ ".

Argue

Examp:

$$\begin{array}{ccc} \mathbb{R} & \xrightarrow{f} & \mathbb{R} \\ x & \mapsto & \frac{x+1}{x-1} \end{array}$$

Claim: f is inj,



Proof: Let $a, b \in \mathbb{R}$. Suppose $f(a) = f(b)$. $\exists E$,

$$\frac{a+1}{2} = \frac{b+1}{2} . \quad (\text{wts: } a=b) \quad \text{Multiplying by 2 gives } a+1 = b+1.$$

Subtracting 1 gives $a = b$. \square

$$\text{E.i. } \mathbb{R}_{\geq 0} \xrightarrow{f} \mathbb{R}$$

$$x \mapsto \frac{x-1}{x+1}$$

I

Proof: Let $a, b \in \mathbb{R}_{\geq 0}$. Suppose $f(a) = f(b)$, i.e., $\frac{a-1}{a+1} = \frac{b-1}{b+1}$.

Clearing denominators gives $(a-1)(b+1) = (b-1)(a+1)$, i.e.,
 $ab - b - 1 = ab - a + b - 1$. Thus $-b + a = -a + b$, so $2a = 2b$, thus

$a = b$. \square

$$\frac{x-1}{x+1} = \frac{x+1-2}{x+1} = 1 - \frac{2}{x+1} = 1 - \frac{2}{a+b}$$

$$\mathbb{R}^3 \xrightarrow{g} \mathbb{R}^2$$

$$(x, y, z) \mapsto (x, y)$$

$$g(1, d, 3) = (1, d)$$

$$g(1, d, 4) = (1, d)$$

but $(1, d, 3) \neq (1, d, 4)$

g B NOT $\tilde{t}_{\alpha j}$

$$\mathbb{R}^2 \xrightarrow{f} \mathbb{R}^3 \quad \boxed{f}$$

$$(x,y) \mapsto (x+y, x-y, x^2+y^2)$$

$$(0,0) \mapsto (0,0,0)$$

$$(1,1) \mapsto (2,0,2)$$

⋮

Proof: Sps $(a,b), (c,d) \in \mathbb{R}^2$. Sps $f(a,b) = f(c,d)$.

$$\text{Thus } (a+b, a-b, a^2+b^2) = (c+d, c-d, c^2+d^2).$$

$$\text{Thus } a+b = c+d, \quad a-b = c-d, \quad a^2+b^2 = c^2+d^2.$$

$$\left. \begin{array}{l} a+b = c+d \\ a-b = c-d \\ a=c \\ b=d \end{array} \right\} \Rightarrow (a,b) = (c,d)$$

Adding the last two gives $2a = 2c$, so $a=c$.

Subtracting gives $2b = 2d$, thus $b=d$. We conclude that

$$(a,b) = (c,d)$$

Example: $\mathbb{R} \rightarrow \mathbb{R}$

$$x \mapsto x^3 + x$$

$$x^3 + x = x(x^2 + 1)$$

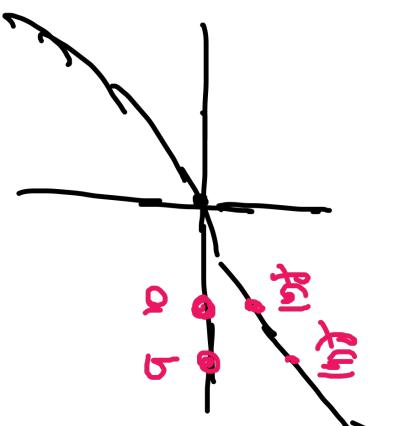
Try "usual" proof: Let $a, b \in \mathbb{R}$. If $f(a) = f(b)$. Then

$$a^3 + a = b^3 + b \quad \dots$$

f acts like ~~flourish~~ to go...

Calc I: $f'(x) > 0 \Rightarrow f$ is increasing

Proof: Since $f'(x) = 3x^2 + 1 \geq 1 \forall x \in \mathbb{R}$. Thus f is increasing and therefore injective. \square



b > a $\Rightarrow f(b) > f(a)$

or decreasing \Rightarrow increasing

$a < b \Rightarrow f(a) < f(b) \Rightarrow f(a) \neq f(b)$

or
 $a > b \Rightarrow f(a) > f(b) \Rightarrow f(a) \neq f(b)$.

$\int x^s$:

$$f(x) = x^5 + 7x^3 + 3x$$

$$f'(x) =$$

$$\boxed{3x^4 + 7x^2 + 3 \geq 3}$$

$$\Rightarrow f(13) \text{, } \rightarrow$$

might need

a small argument

to show $f' > 0$.

Not Inv



$\cos \theta = \cos \alpha\pi$
but $\theta \neq \alpha\pi$.

Different funcs b/c different domains

$[0, \pi] \rightarrow \mathbb{R}$

$x \mapsto \cos x$

This is injective.

Pf: The der. of $\cos x$ is $-\sin x$. Since $-\sin x \leq 0 \forall x \in [0, \pi]$,

\cos is decreasing, and thus injective. \square

$$\begin{array}{l} \mathbb{C} \xrightarrow{\quad} \mathbb{C} \\ x \mapsto x^2 \end{array} \quad \text{NI } (i)^2 = (-i)^2$$

$$\begin{array}{c} " \\ -1 \\ " \\ (-1)^2 = i^2 \\ " \\ 1(-1) = -1 \end{array}$$

$$P(\mathbb{R}) \longrightarrow P(\mathbb{Z})$$

$$S \xrightarrow{\quad} S \cap \mathbb{Z}$$

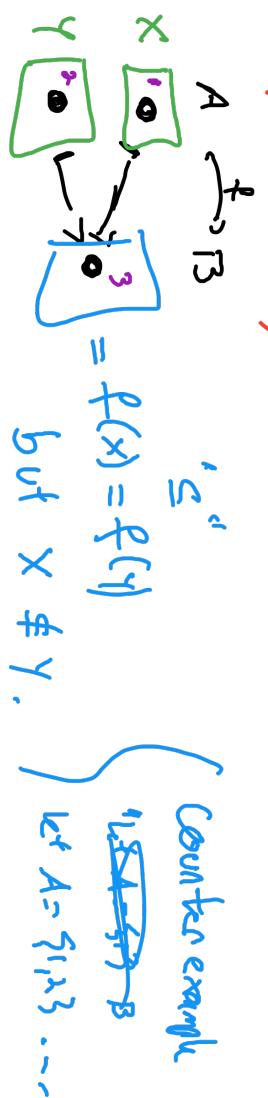
Not I: Are there sets $S_1, S_2 \subseteq \mathbb{R}$
 s.t. $S_1 \neq S_2$ but $S_1 \cap \mathbb{Z} = S_2 \cap \mathbb{Z}$?

$$\boxed{\{1, \pi\}} \cap \mathbb{Z} = \{1\} = \boxed{\{1, e\}} \cap \mathbb{Z}$$

$$A \xrightarrow{f} B$$

$$x, y \in A$$

Recall: " $f(x) \leq f(y) \Rightarrow x \leq y$ " is false.



"Find the "correct" hypothesis"

Not INJ.

L: \exists f is inj and $f(x) \leq f(y)$. Then $x \leq y$.

THE DEFN OF
 $f(x)$

P: \exists f is inj and $f(x) \leq f(y)$. Let $a \in X$. Then $f(a) \in f(X)$.

(To use the hypothesis " $f(x) \leq f(y)$ ", need an $c \in f(X)$)

Since $f(a) \in f(X)$, and $f(x) \leq f(y)$, $\boxed{f(a)} \in f(Y)$.

Thus $\exists c \in Y$ s.t. $f(a) = f(c)$. Since f is

injective, $a=c$. Since $c \in Y$, $a \in Y$. \blacksquare

injective, $a=c$. Since $c \in Y$, $a \in Y$. \blacksquare

Week 1d: Compositions of function

Defn: Let $f: A \rightarrow B$ and $g: B \rightarrow C$ be fns.

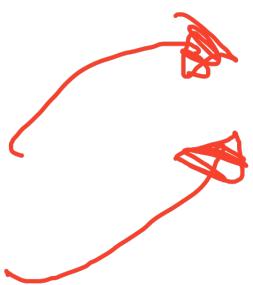
We define the composition of g and f to be

the fn

$g \circ f : A \rightarrow C$ defined by

$$(\text{def}) \quad a \mapsto (g \circ f)(a) := g(f(a)).$$

ALT notation $gf = \underline{g \circ f}$



Any time you do a proof w/ compositions, use the defn

Example:

$$f: \mathbb{R} \rightarrow \mathbb{R}$$
$$x \mapsto x^2$$
$$x \mapsto x+1$$

$$(g \circ f)(1) = g(f(1)) = g(1) = 2$$

$$(f \circ g)(1) = f(g(1)) = f(2) = 4$$

Note: $g \circ f \neq f \circ g$ i.e., composition is not commutative.

Same domain & codomain \mathbb{R} But

$$(g \circ f)(1) \neq (f \circ g)(1)$$

Sometimes there is a "diff" formula for $g \circ f$

$$(g \circ f)(x) = g(f(x)) = g(x^2) = (x^2) + 1$$

$$(f \circ g)(x) = f(g(x)) = f(x+1) = (x+1)^2 = x^2 + 2x + 1$$

Warning: usually $f \circ g$ and $g \circ f$ don't both make sense.

$$f: A \rightarrow B \quad g: B \rightarrow C$$

$g \circ f$ is ok but

$f \circ g$ is not defined!



$$(f \circ g)(b) = f(g(b)) \quad \text{but}$$

$$b \in B \rightsquigarrow g(b) \in C$$

But domain of $f = A$

If: $C \subseteq A$ we can fix this.

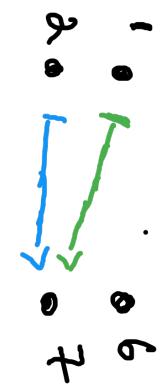
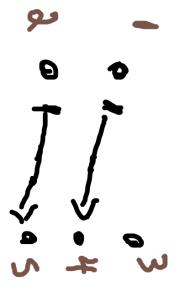
Example:

$$A \xrightarrow{f} B$$

$$B \xrightarrow{g} C$$

$$A \xrightarrow{f} B \xrightarrow{g} C$$

$\circ g \circ f$



$$(g \circ f)(1) = g(f(1)) = g(4) = 7$$

$$(g \circ f)(2) = g(f(2)) = g(5) = 7$$

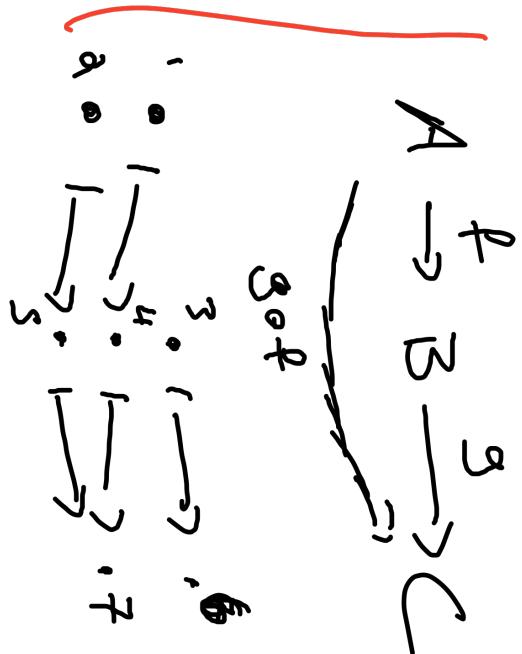
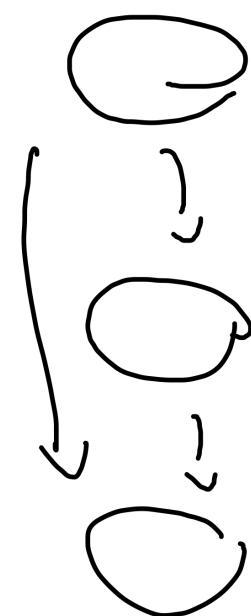
$$(g \circ f)(3) = g(f(3)) = g(6) = 7$$

$$(g \circ f)(4) = g(f(4)) = g(7) = 7$$

$$(g \circ f)(5) = g(f(5)) = g(8) = 7$$

Note: f is inj + g is surj, but $g \circ f$ is neither inj or surj

E.g. $f: \mathbb{N} \rightarrow \mathbb{N}$ and $g: \mathbb{N} \rightarrow \mathbb{N}$



"Simplity" (TEST Functions)

Smallest

"Smallest"

inj,

ref

surj

- f
- g
- \circ
- \rightarrow
- \hookleftarrow
- \hookrightarrow

fun

- f'
- g'
- \circ'
- \rightarrow'

gof

USEFUL FOR COUNTER EXAMPLES

(1) Let $f: \mathbf{R} \rightarrow \mathbf{R}$ be the function $f(x) = \frac{1}{1+x^2}$ and let $g: \mathbf{R} \rightarrow \mathbf{R}$ be the function $g(x) = e^x$.

- (a) What is $\underline{g \circ f}(0)$?
(b) What is $\underline{(f \circ g)}(0)$?
(c) Give a formula for $f \circ g$ and $g \circ f$.

$$\underline{\text{got } (a)} = (g \circ f)(0)$$

$$(g \circ f)(0) = g(f(0)) = g(1) = e$$
$$(f \circ g)(0) = f(g(0)) = f(1) = \frac{1}{2}$$

$$(f \circ g)(x) = f(g(x)) = f(e^x) = \frac{1}{1+(e^x)^2} = \frac{1}{1+e^{2x}}$$
$$(g \circ f)(x) = g(f(x)) = g\left(\frac{1}{1+x^2}\right) = e^{\frac{1}{1+x^2}}$$

- (2) Let $f: \mathbf{R} \rightarrow \mathbf{Z}$ be the function $f(x) = \lfloor x \rfloor$ (i.e., round x down to the nearest integer) and let $g: \mathbf{Z} \rightarrow \mathbf{Z}$ be the function $g(n) = \underline{\text{the number of distinct prime factors of } n}$. (So $g(0) = g(1) = 0$, $g(4) = 1$, $g(6) = 2$)

- (a) What is $g \circ f(\pi)$?
- (b) What is $g \circ f(91.1023124)$?
- (c) Is $g \circ f$ injective? Surjective?

$$f(1.1) = 1, f(1.0) = 1$$

$f(x) = \underline{\text{the largest integer } y \text{ s.t. } y \leq x}$

$$(g \circ f)(\pi) = g(f(\pi)) = g(f(3, 4, 19, \dots)) = g(3) = 1$$

$$(g \circ f)(1.10 \dots) = g(f(1, 10, \dots)) = g(1) = 0$$

7.13

$g \circ f$ not surj b/c $\underline{g \text{ is not surj}}$, b/c $\underline{g(x) \geq 0 \forall x \in \mathbf{Z}}$

$g \circ f$ not inj b/c $\underline{f \text{ is not inj}}$

$$(g \circ f)(\boxed{1}) = g(\underline{f(1)}) = g(1) = 0$$

$$(g \circ f)(\boxed{1}) = g(\underline{f(1)}) = g(1) = 0$$

(3) Let $f: \mathbf{Z} \rightarrow P(\mathbf{Z})$ be the function $f(n) = \{n\}$ and let $g: P(\mathbf{Z}) \rightarrow P(\mathbf{Z})$ be the function $g(S) = S \cap \{1\}$.

- (a) What is $g \circ f(0)$?
- (b) What is $g \circ f(1)$?
- (c) Give a formula for $g \circ f$.

$$(g \circ f)(0) = g(f(0)) = g(\{0\}) = \{0\} \cap \{1\} = \emptyset$$

$$(g \circ f)(1) = g(f(1)) = g(\{1\}) = \{1\} \cap \{1\} = \{1\}$$

$$(g \circ f)(n) = g(f(n)) = g(\{n\}) = \{n\} \cap \{1\}$$

$$= \{1\} \quad n = 1$$

$$\} \neq \emptyset \quad n \neq 1$$

X, Y Sets

$$\text{Fun}(X, Y) := \left\{ f : X \rightarrow Y \right\}$$

$$\text{Fun}(A, B) \times \text{Fun}(B, C) \xrightarrow{\circ} \text{Fun}(A, C)$$

$$\cup$$

$$(f, g) \mapsto g \circ f$$

$$\cup$$

$$\mathcal{P}(A) \times \mathcal{P}(A) \xrightarrow{\cup} \mathcal{P}(A)$$

$$\cup$$

$$(S, T) \mapsto S \cap T$$

h o g o f

Lemma: Composition is associative, i.e.,

Let $f: A \rightarrow B$, $g: B \rightarrow C$, $h: C \rightarrow D$ be functions. Then

Then $h \circ (g \circ f) \stackrel{(*)}{=} (h \circ g) \circ f$.

Proof: The domain & codomain agree.

(WTS: $\forall a \in A$, $(h \circ (g \circ f))(a) = ((h \circ g) \circ f)(a)$)

Let $a \in A$. Then $(h \circ (g \circ f))(a) = h((g \circ f)(a)) = h(g(f(a)))$.

Also $((h \circ g) \circ f)(a) = (h \circ g)(f(a)) = h(g(f(a)))$.

These are equal, so $h \circ (g \circ f) = (h \circ g) \circ f$. \square

1. Induction to prove for 4 or more, induction, using \square as base case

$$f \circ (g \circ (h \circ w)) = (f \circ g) \circ (h \circ w) - \square$$

use

$\rightarrow \leftarrow \cdot \nearrow \cdot \searrow \cdot$ to check

(4) Let $f: A \rightarrow B$ and $g: B \rightarrow C$ be functions. Prove or disprove each of the following:

- (a) If f and g are injections, then gf is an injection.
- (b) If f and g are surjections, then gf is a surjection.
- (c) If f and g are bijections, then gf is a bijection.
- (d) If gf is an injection, then f and g are injections.
- (e) If gf is a surjection, then f and g are surjections.
- (f) If gf is a bijection, then f and g are bijections.
- (g) If gf is an injection, then f is an injection.
- (h) If gf is an injection, then g is an injection.
- (HW) If gf is a surjection, then f is a surjection.
- (i) If gf is a surjection, then g is a surjection.
- (j) If gf is a bijection, then f is a bijection.
- (k) If gf is a bijection, then g is a bijection.
- (l) If gf is a bijection, then f is a bijection.
- (m) If gf is an injection and g is a bijection, then f is an injection.

bij = inj AND surj

δ^{-1} means

$$\delta(x) = g(\gamma) \Rightarrow x = \gamma$$

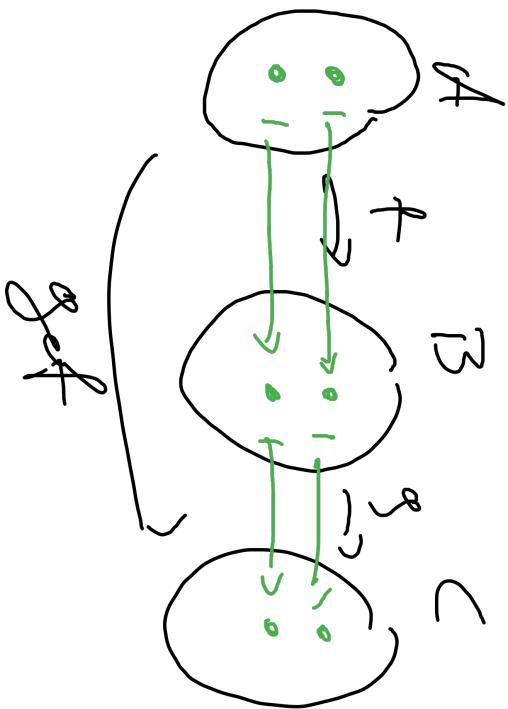
$$\begin{cases} \delta(x) = g(\gamma) \\ \text{Apply w } x = f(a) \\ \gamma = f(b) \end{cases}$$

for proof: Sps f and g are \nearrow . (wtp: $(gof)(a) = (gof)(b) \Rightarrow a=b$)

Let $a, b \in A$. Assume $(gof)(a) = (gof)(b)$. Then

$$g(f(a)) = g(f(b))$$

Since g is \nearrow , $f(a) = f(b)$. Since f is \nearrow , $a = b$. \square



$$(gof)(A) = C$$

It is $\text{im } g \circ f = C$, i.e.

(4) Sps f and g are surj. (wts: $g \circ f$ is surj). $\exists E \forall x \in C, \exists a \in A$ s.t. $(g \circ f)(a) = x$

Let $x \in C$. Since g is surj, $\exists b \in B$ s.t. $g(b) = x$. Since f is surj,

$\exists a \in A$ s.t. $f(a) = b$. Then $(g \circ f)(a) = g(f(a)) = g(b) = x$. \square

(4) Let $f: A \rightarrow B$ and $g: B \rightarrow C$ be functions. Prove or disprove each of the following:

- I
~~(a)~~ If f and g are injections, then gf is an injection.
~~(b)~~ If f and g are surjections, then gf is a surjection.
~~(c)~~ If f and g are bijections, then gf is a bijection.
~~(d)~~ If gf is an injection, then f and g are injections.
~~(e)~~ If gf is a surjection, then f and g are surjections.
~~(f)~~ If gf is a bijection, then f and g are bijections.
~~(g)~~ If gf is an injection, then f is an injection.
~~(h)~~ (HW) If gf is an injection, then g is an injection.
~~(i)~~ (HW) If gf is a surjection, then f is a surjection.
~~(j)~~ (HW) If gf is a surjection, then g is a surjection.
~~(k)~~ If gf is a bijection, then f is a bijection.
~~(l)~~ If gf is a bijection, then g is a bijection.
~~(m)~~ If gf is an injection and g is a bijection, then f is an injection.

$a \text{ and } b \rightarrow c$

$f \circ g$
 $\circ \rightarrow \circ \rightarrow \circ$
Counterexample to d be,
not co ex to g

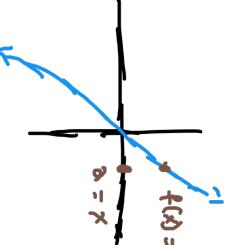
Proof of g: So gf is $\downarrow a$. Let $a, b \in A$.
Suppose $f(a) = f(b)$. Then $g(f(a)) = g(f(b))$,
and since gaf is $\downarrow a$, $a = b$. \square

Week 13: Inverse Functions

I def: f^{-1} "undoes" f

$$f(x) = x^3$$

$$f^{-1}(x) = x^{1/3}$$



$$f^{-1}(8) = 2$$

$$f(3) = 27$$

$$f^{-1}(27) = 3$$

"the x s.t. $f(x) = 26$ " = $f^{-1}(26) = 26^{1/3} = \dots?$

To solve the eqn $f(x) = 26$ for x

$f(x) = x^3 = y$ & solve for x

$$x = y^{1/3}$$

"Works", but isn't the domain

(*) $(f^{-1} \circ f)(x) = \underline{f^{-1}(f(x))} = f^{-1}(x^3) = (x^3)^{1/3} = x = \text{id}_{\mathbb{R}}(x)$

$$f^{-1} \circ f = \text{id}_{\mathbb{R}}$$

$\text{id}_A: A \rightarrow A$

$$x \longmapsto x$$

"the" identity fn

$$\text{id}_A(x) = x, \forall x \in A$$

Prop: Let $f: A \rightarrow B$ be any fn. Then

$$(i) f \circ \text{id}_A = f \quad A \xrightarrow{\text{id}_A} A \xrightarrow{f} B$$

$$(ii) \text{id}_B \circ f = f \quad A \xrightarrow{f} B \xrightarrow{\text{id}_B} B$$

Proof: i) $f \circ \text{id}_A$ and f have the same domain & codomain.

Let $a \in A$. Then $(f \circ \text{id}_A)(a) = f(\text{id}_A(a)) = f(a)$.

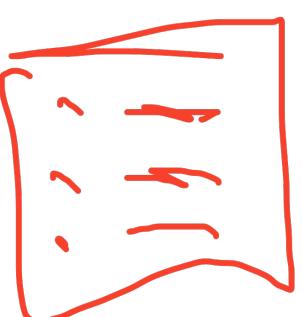
ii) is similar.

Defn: We say that a f'm $f: A \rightarrow B$ is invertible if

$$\exists g: B \rightarrow A \text{ s.t. } \begin{aligned} f \circ g &= \text{id}_B \\ g \circ f &= \text{id}_A \end{aligned}$$

When such a g exists, we call g an inverse of f and sometimes write $g = f^{-1}$.

Warnings: (i) $f^{-1} \neq \frac{1}{f}$



(ii) not every f has an inverse!

$$A = B = \mathbb{R}$$

$$f: A \rightarrow A$$

$$x \mapsto x^3$$

$$g: x \mapsto x^{1/3}$$

$$(g \circ f)(x) = g(f(x)) = g(x^3) = (x^3)^{1/3} = x = \text{id}_{\mathbb{R}}(x)$$

$$(f \circ g)'(x) = f(g(x)) = f(x^{1/3}) = x^{1/3} = x = \text{id}_{\mathbb{R}}(x)$$

Example:

What is the inverse g of f ?

$$B \xrightarrow{g} A$$

$$A \xrightarrow{f} B$$

$$f \circ g = \text{id}_B$$

$$g \circ f = \text{id}_A$$

$$\begin{matrix} 4 & 0 & 1 & 7 & 0 \\ \swarrow & \searrow & \downarrow & \nearrow & \nwarrow \\ 5 & 0 & 1 & 2 & 0 \\ \searrow & \swarrow & \downarrow & \nearrow & \nwarrow \\ 6 & 0 & 1 & 3 & 0 \end{matrix}$$

$$(g \circ f)(1) = g(f(1)) = g(5) = 1$$

$$(g \circ f)(z) = g(f(z)) = z$$

$$g^{-1}(y) = z$$

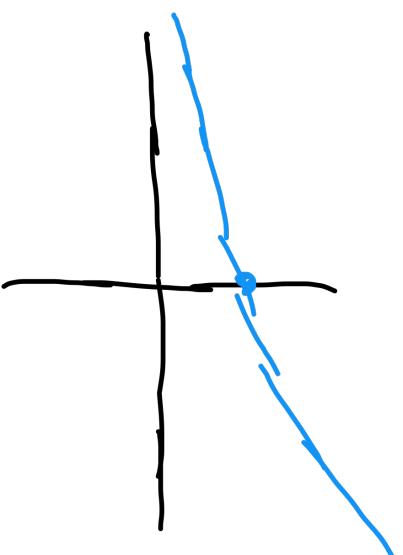
" $g(x)$ is the solution to $f(x) = x$ "

$$g(x) = z \Rightarrow x = f(z)$$

Example:

$$\mathbb{R} \xrightarrow{e} \mathbb{R}_{>0}$$

$$x \mapsto e^x$$



This has an inverse \ln

$$\mathbb{R}_{>0} \xrightarrow{\ln} \mathbb{R}$$

$$x \mapsto \ln x$$

$$\ln \circ e = \text{id}_{\mathbb{R}}$$

$$e \circ \ln = \text{id}_{\mathbb{R}_{>0}}$$

$$e^{\pi} = 1$$

$\ln 1 = "the solution to $e^x = 1"$$

$$e^\pi = e^{\pi} \dots$$

$$\ln(e^\pi) = \pi$$

$$\ln(e^x) = x$$

$$e^{\ln x} = x$$

$$e^x = y$$

$$\ln y = x$$

$$\mathbb{R} \xrightarrow{f} \mathbb{R}$$

$$x \mapsto x^5 + 4x$$

$$f'(x) = 5x^4 + 4 \geq 4$$

$\Rightarrow f$ is inj

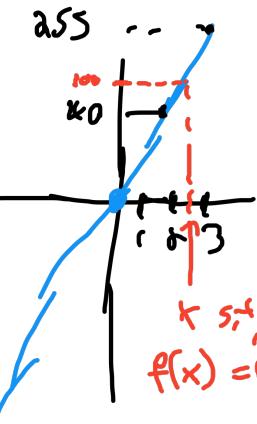
$\exists v \in \mathbb{R} \Rightarrow f(v) = 100$

HS way

Solve $x^5 + 4x = y$ for x

-- can't do this

nicely



FACT f has an inverse g .

$$f(0) = 0 \quad g(0) = 0$$

$$f(1) = 5 \quad g(5) = 1$$

$$f(x) = y \Leftrightarrow g(y) = x$$

$$f(2) = 40 \quad g(40) = 2$$

$$f(x) = -5 \Leftrightarrow g(-5) = x$$

$$\Leftrightarrow$$

$$x^5 + 4x = -5 \Leftrightarrow x = -1$$

$$\text{Thus } g(-5) = -1$$

$$g(100) = x \Leftrightarrow f(x) = 100$$

$$\Leftrightarrow x^5 + 4x = 100$$

$$\Leftrightarrow x^5 + 4x - 100 = 0$$

By FvT, has a sol, b/c

$$f(2) = 40$$

$$f(3) = 3^5 + 4 \cdot 3 = 253$$

$\Rightarrow \exists x \in [2,3] \text{ s.t. } f(x) = 100$

$$f^{-1}(100) = \text{the } x \in [2,3] \dots$$

Sometimes f has no inverse

$$\mathbb{R} \xrightarrow{f} \mathbb{R}$$

$$x \mapsto x^2$$

① What is $f^{-1}(-1)$?

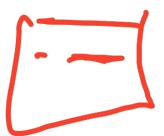
There is ~~no~~ input s.t. $f(x) = -1$

Math hack: replace codomain of image

② What is $f^{-1}(4)$?

$$f(a) = 4$$

$$f(-a) = 4$$



Problem: 13

$$f^{-1}(4) = a \text{ or } -a ?$$

Need inj to have an inverse

What about $x^{1/2}$? ↴ not def for $x < 0$

↳ ambiguous (+ & - root)

THM: $f: A \rightarrow B$ has an inverse $\Leftrightarrow f$ is bijective.

Proof: " \Rightarrow " Assume f has an inverse g .

(Inj , $w\exists s \forall a b \in A, f(a) = f(b) \Rightarrow a = b$)

Let $a, b \in A$. Suppose $f(a) = f(b)$. Then, $g(f(a)) = g(f(b))$.

Since $g = f^{-1}$, $g \circ f = \text{id}$, so $a = b$. ($(g \circ f)(a) = g(f(a)) = \text{id}(a) = a$)

$(S, w\forall s \forall b \in B, \exists q \in A \text{ s.t. } f(q) = b)$

Let $b \in B$. Let $a = g(b)$. Then $f(a) = f(g(b)) = (f \circ g)(b) = \text{id}_B(b) = b$.

" \Leftarrow " Assume f is bijection. Let's define $g: B \rightarrow A$ as follows.

Let $b \in B$. Since f is surj, $\exists q \in A \text{ s.t. } f(q) = b$. Since f is inj,
there is only one such q . Define $g(b) = q$. Then

$$(g \circ f)(a) = g(f(a)) = a.$$

$$(f \circ g)(b) = f(g(b)) = b.$$

to help us prove

(1) USE THM \hookrightarrow to help w/ counterexamples

(2) Given f , to find f^{-1} , "solve" for y

$$\begin{aligned} f(y) &= \text{target} \\ g(f) &= \text{input for } g \end{aligned}$$

(3') If you have a guess for y ,

verify your guess by plugging y into

(ii) x^2 not bi \Rightarrow not invertible

$$f(x) = x^2 + x^3 + x^4$$

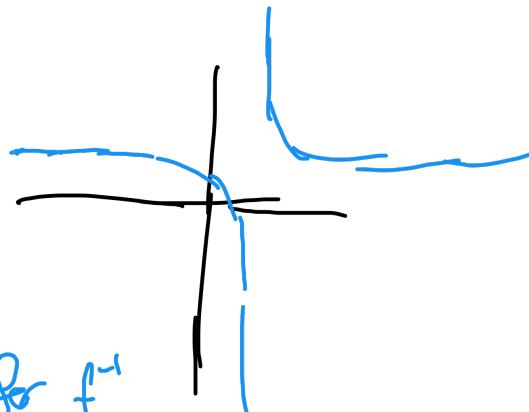
vs inv. by Thm

$$b_1 > b_2,$$

(Still need to explain why $f(3)$ & $f(4)$)

Ex: $\mathbb{R} - \{-1\} \xrightarrow{f} \mathbb{R} - \{1\}$

$$x \mapsto \frac{x+1}{x-1}$$



How to find a formula for f^{-1}

$$f(x) = y \Leftrightarrow x = g(y)$$

" $g(y)$ is the x s.t. $f(x) = y$ "

$$\frac{x+1}{x-1} = y \quad \frac{2}{x-1} = y-1 \Rightarrow \frac{x-1}{2} = \frac{1}{y-1}$$

||

$$\frac{x-1+\lambda}{x-1} = 1 + \frac{\lambda}{x-1} \stackrel{?}{=} y \Rightarrow x-1 = \frac{\lambda}{y-1} \Rightarrow x = \boxed{1 + \frac{\lambda}{y-1}} = g(y)$$

FE $(f \circ g)(x) \stackrel{?}{=} x$

$$f\left(\frac{1+\lambda}{x-1}\right) = \frac{\left(\frac{1+\lambda}{x-1}\right) + 1}{\left(\frac{1+\lambda}{x-1}\right) - 1} \stackrel{MAW}{=} x$$

E-xample:

$$\mathbb{R}^2 \xrightarrow{f} \mathbb{R}^2$$

$$\begin{pmatrix} x \\ x+y \\ x-y \end{pmatrix} \xrightarrow{g} \begin{pmatrix} x \\ x \\ x \end{pmatrix}$$

$$\mathbb{R}^2 \xrightarrow{g} \mathbb{R}^2$$

$$\begin{pmatrix} x \\ x+y \\ x-y \end{pmatrix} \xrightarrow{f} \begin{pmatrix} x \\ x \\ x \end{pmatrix}$$

$$(g \circ f)(x) = g(f(x)) =$$

$$g\left(\begin{pmatrix} x \\ x+y \\ x-y \end{pmatrix} \right) =$$

$$\begin{pmatrix} x \\ (x+y)+(x-y) \\ x+y+x-y \end{pmatrix} =$$

$$\begin{pmatrix} x \\ 2x \\ 2x \end{pmatrix} =$$

kom
us
Gof

Week 14: Relations (4.2)

Informally: a "relation" is a way to compare "things"

Example: $S = \mathbb{R}$, \geq is a relation

$\forall a, b \in S$, " $a \geq b$ " is either true or false

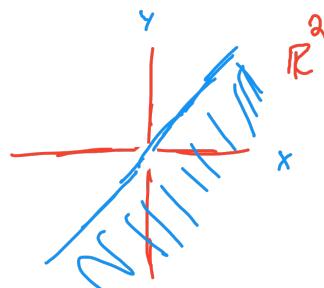
Defn: let S be a set. A relation on S is a subset $R \subseteq S \times S$.

Usually use \sim
If $(a, b) \in R$, we say that " a is related to b "
and write $[a \sim b]$ (or $a \underset{R}{\sim} b$).

Example: $S = \mathbb{R}$

$R \subseteq \mathbb{R} \times \mathbb{R}$ given by

$$R \stackrel{\text{def}}{=} \{(a, b) \in \mathbb{R} \times \mathbb{R} \mid a \geq b\}$$



Example: $S = \{0, 1, 2\}$

$$R \subseteq S \times S$$

$$R \stackrel{\text{def}}{=} \{(0, 1), (0, 2), (1, 2)\}$$

$0 \sim 1$ true b/c $(0, 1) \in R$

$1 \sim 2$ true

$2 \sim 0$ false b/c $(2, 0) \notin R$

"Equivalence" Relations

Axiomatize the notion of "the same" or "more or less the same"

There are many properties & A's which don't depend on



Size or orientation.



but still the same

not literally the same

Example:



$$\frac{2}{3} = \frac{4}{6}$$

Defn: let S be a set.. Let R be a relation on S

We say that R is an equivalence relation if

$$(R) \forall a \in S, a \sim a$$

• Reflexive

$$(S) \forall a, b \in S, a \sim b \Rightarrow b \sim a$$

• Symmetric

$$(T) \forall a, b, c \in S, a \sim b \wedge b \sim c \Rightarrow a \sim c$$

• transitive

Let $a \in S$. We define the equivalence class of a to be

$$[a] \stackrel{\text{def}}{=} \{b \in S \mid b \sim a\}$$

Warning:

NOTE: if $b \in [a]$ then $[b] = [a]$

↑
Set

↑
Sets

Ref interval
notation.

$$a \in [a] \quad b \in a \sim a$$

new
notation

$$[a] \text{ vs } \{a\}$$

↗

↑ the set whose only elt is a

More
elts

Examples: $S = \mathbb{Z}_-$, $a \sim b$ if $a \mid a-b$

$$2 \sim 4$$
 b/c $2-4 = -2$ and $2 \mid -2$

$$1 \sim 3$$

$$1-3 = -2$$

$$1 \not\sim 2$$

$$1-2 = -1$$

$$2 \sim 1$$

$(2,4) \in R$
 $(2,3) \notin R$

i.e., $a \sim b$ if a and b have the same remainder when you divide by 2

i.e., $a \sim b$ if they have the same parity

Claim: This is an equiv. relation.

Pf: (R) Let $a \in \mathbb{Z}$. (WTP: $a \sim a$, i.e., $2|a-a$)

Since $a-a=0$, $2|a-a$, so $a \sim a$.

(S) Let $a, b \in \mathbb{Z}$. Suppose $a \sim b$. Then $2|a-b$,

(WTP: $b \sim a$, i.e., $2|b-a$). Since $b-a = -(a-b)$,

$2|b-a$, so $b \sim a$.

(T) Let $a, b, c \in \mathbb{Z}$. Suppose $a \sim b$ and $b \sim c$. Then $2|a-b$ and $2|b-c$.

(WTS: $a \sim c$, i.e., $2|a-c$) Then $2|(a-b)+(b-c)$, so $2|a-c$. Thus $a \sim c$. \square

①: What are the equiv. classes?

$$\begin{aligned} [0] &= \{a \in \mathbb{Z} \text{ s.t. } a \sim 0\} & a \sim 0 &\iff 2|a-0 \\ &= \{a \in \mathbb{Z} \text{ s.t. } 2|a\} & &\iff 2|a \\ &= 2\mathbb{Z} = \mathbb{E} \end{aligned}$$

$$\begin{aligned} [1] &= \{a \in \mathbb{Z} \text{ s.t. } a \sim 1\} & a \sim 1 &\iff 2|a-1 \\ &= \{a \in \mathbb{Z} \text{ s.t. } a \text{ is odd}\} & &\iff a-1 \text{ is even} \\ &= 2\mathbb{Z} + 1 \text{ or } \mathbb{O} & &\iff a \text{ is odd} \end{aligned}$$

NOTE: $[0] \cup [1] = \mathbb{Z}$ AND $[0] \cap [1] = \emptyset$ "partition"

$$\begin{aligned} [2] &= \{a \in \mathbb{Z} \text{ s.t. } a \sim 2\} & a \sim 2 &\iff 2|a-2 \\ &= \mathbb{E} & &\iff 2|a \end{aligned}$$

$$[0] = [2]$$

$$[0] \neq [1]$$

$$0 \in [0] \text{ but } 0 \notin [1]$$

$S = \mathbb{R}$ $\forall x, y$ if $x < y$

NOT AN E.R. b/c

Not (R) or (S). (IS (T))

Pf: Let $a=0$. Then $a < 0$ is false, so $a \neq 0$.

Thus $<$ is not reflexive.

Let $a=0$ and $b=1$. Then $a < 1$, so $a \sim 1$, but $1 \not\sim 0$, so $1 \neq 0$.

(T) $a \sim b \wedge b \sim c \Rightarrow a \sim c$.

$S = \mathbb{R}$, $x, y \in \mathbb{R}$, $x \sim y$ if $x - y \in \mathbb{Q}$

$$\pi \sim \pi + 1 \quad \pi - (\pi + 1) = -1 \in \mathbb{Q}$$

$$0 \neq \bar{\pi} \text{ b/c } \bar{\pi} - 0 = \bar{\pi} \notin \mathbb{Q}$$

Claim: this is an EP.

Pf: (P) Let $a \in \mathbb{R}$. Then $a - a = 0 \in \mathbb{Q}$. Thus $a \sim a$.

(S) Let $a, b \in \mathbb{R}$. Suppose $a \sim b$, i.e., $a - b \in \mathbb{Q}$.

Then $b - a \in \mathbb{Q}$, so $b \sim a$.

(T) Let $a, b, c \in \mathbb{R}$. Suppose $a \sim b$ and $b \sim c$, i.e., $a - b \in \mathbb{Q}$ and $b - c \in \mathbb{Q}$.

Adding gives $a - c = (a - b) + (b - c) \in \mathbb{Q}$, thus $a \sim c$. \blacksquare

A, B sets. Define $A \sim B$ if

\exists a bijection $f: A \rightarrow B$.

Thm $f: B \text{ bi} \iff f \text{ has an inverse}$

$$\{1\} \sim \{2\} \text{ via } f(1) = 2$$

$$\{1, 3\} \not\sim \{1, 2\}$$

Claim: This \sim is an E.R.

Pf: (R) Let A be a set. (wtp: $A \sim A$, i.e., \exists bijection $A \xrightarrow{f} A$)

The $f = \text{id}_A$ is a bijection from A to A , because the inverse of id_A is id_A , i.e., $\text{id}_A \circ \text{id}_A = \text{id}_A$.

(S) Let A, B be sets. Suppose \exists a bijection $f: A \rightarrow B$.

By the Thm, $\exists g: B \rightarrow A$ s.t. $g = f^{-1}$. Since g is invertible, g is a bijection.

(T) Let A, B, C be sets. Suppose $\exists f: A \rightarrow B$ and $g: B \rightarrow C$ s.t. f and g are bijective. Then $g \circ f: A \rightarrow C$ is bijective (bc we proved in L1d).

Thus $A \sim C$. \blacksquare

A, B sets, $A \sim B$ if

$\exists f: A \rightarrow B$ s.t.

f is surj.

\sim is R and (T) (Same pf)

But: not (S).

$$A = \{1, 3\}, B = \{2, 3\}$$

There are 2 fns from $A \rightarrow B$

$$\begin{cases} \exists \text{ surj } B \rightarrow A \\ \text{No surj } A \rightarrow B \end{cases}$$

$$f(1) = 2$$

$$g(1) = 3$$

Neither is a surjection. \square

$$S = R \quad x \sim y \quad \text{if} \quad x=1 \text{ or } y=1$$

$$\begin{array}{ll} 1 \sim 1 & 2 \not\sim 3 \\ 1 \sim 2 & 2 \not\sim 3 \\ 2 \sim 1 & . \end{array}$$

TRUE

$$\neg(\neg R) \Leftrightarrow \neg\neg R$$

(S) is true.

Pf: Let $a, b \in R$. Sps $a \sim b$, i.e., $a=1$ or $b=1$.

(wtk: $b \sim a$, i.e., $b=1$ or $a=1$) Since "or" is commutative, $b=1$ or $a=1$. Thus $b \sim a$.

$$\left. \begin{array}{l} \neg R \Leftrightarrow R \wedge S \wedge T \\ \neg\neg R \Leftrightarrow \neg R \vee \neg S \wedge \neg T \end{array} \right\}$$

$$(T) \stackrel{\vee}{\sim} a=1, b=1, c=3$$

then $a \sim b \sim c$, but $a \not\sim c$