

# Data Structures and Algorithms

## Binary Search Tree and Heap

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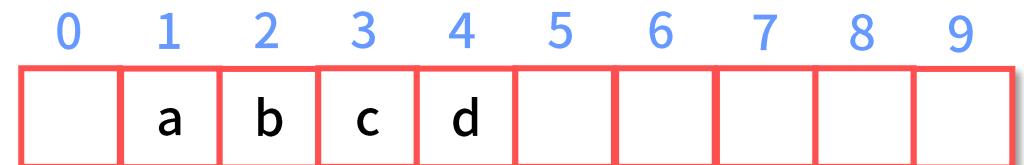
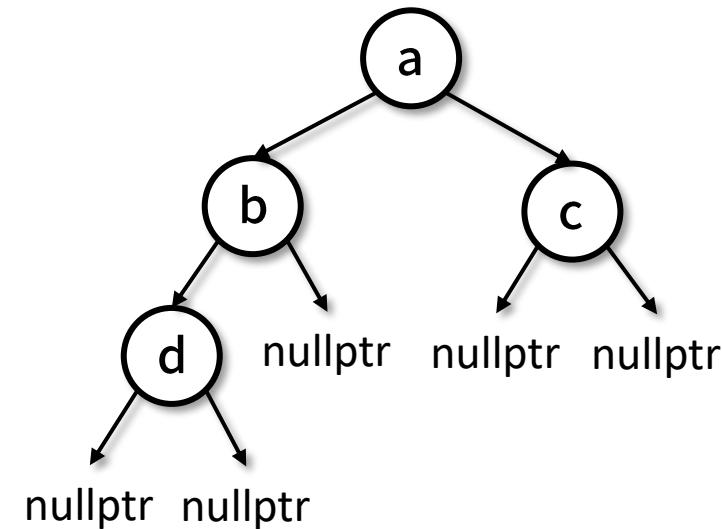
# Overview

- Data Structures:
  - Linked-List, Array
  - Stack, Queue
  - Union Find, Hash Table
  - **Binary Search Tree, Heap**
  - Graph
- Algorithms
  - Big-O Notation
  - Sorting
  - Graph Algorithms
  - Dynamic Programming

# Binary Tree

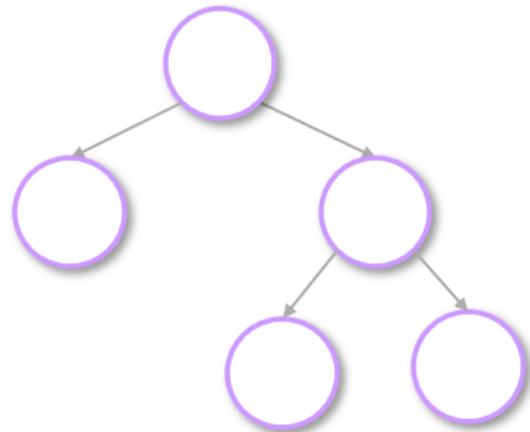
# Binary Tree (1)

- Each parent node can have at most 2 children node.
- The top node is called the **root node**.
- The node without children is called the **leaf node**.
- Can be implemented using linked-list or array.

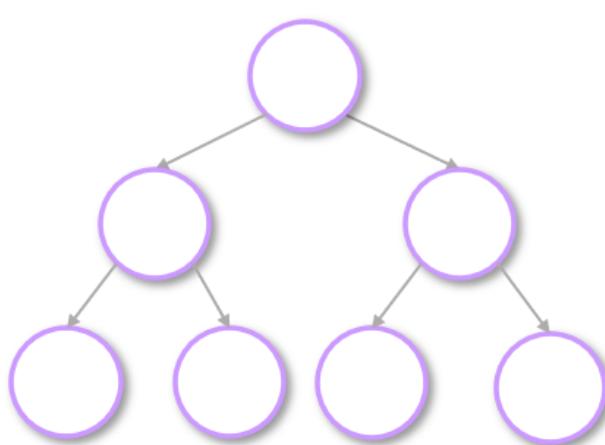


# Binary Tree (2)

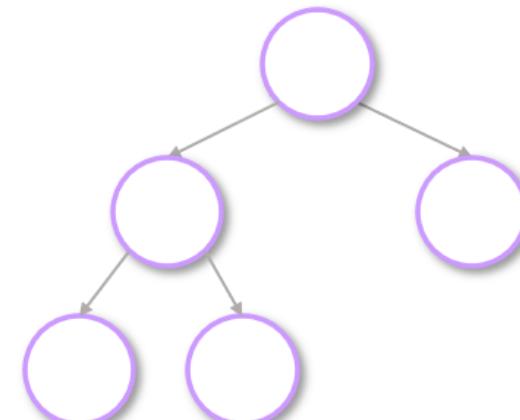
Full



Perfect



Complete

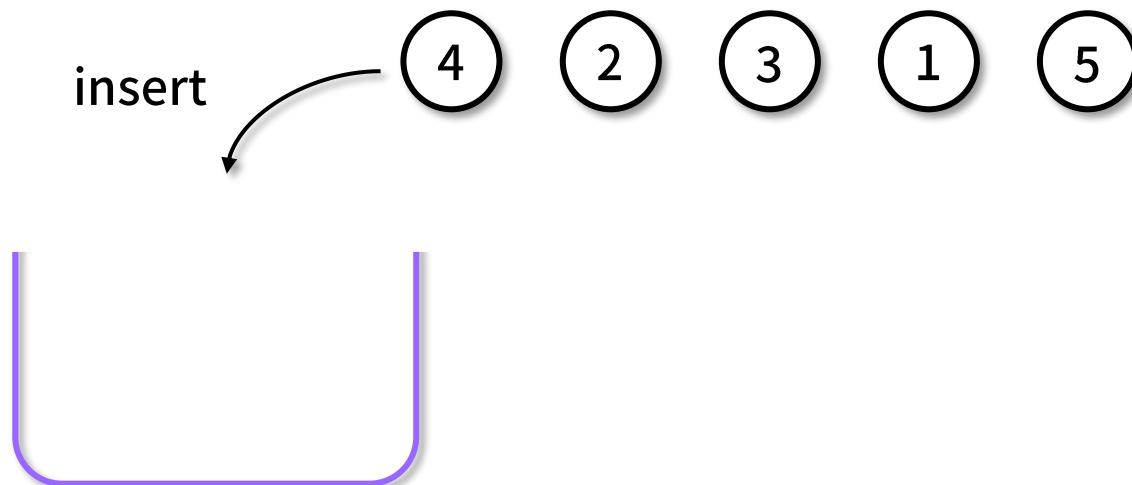


- all nodes have 0 or 2 children
- all interior nodes have 2 children
- all leaves on the same level
- all interior nodes (exclude last layer) is perfect
- all leaves are pushed to the left

# Heap

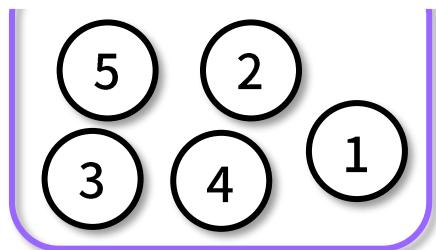
# Minimum Heap (1)

- Always pop the smallest element.



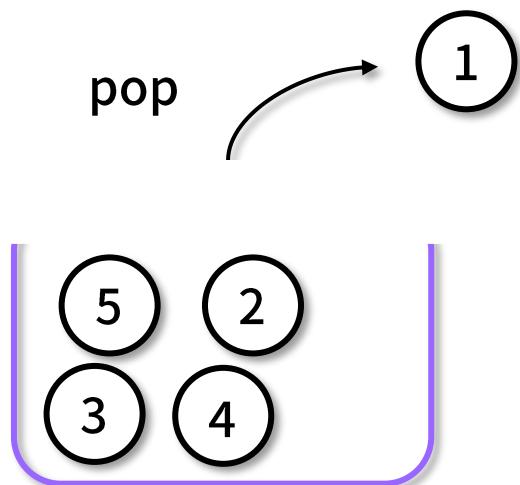
# Minimum Heap (2)

- Always pop the smallest element.



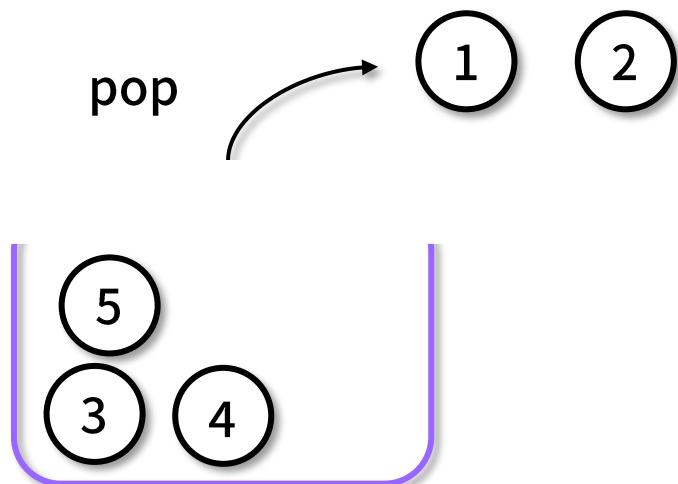
# Minimum Heap (2)

- Always pop the smallest element.



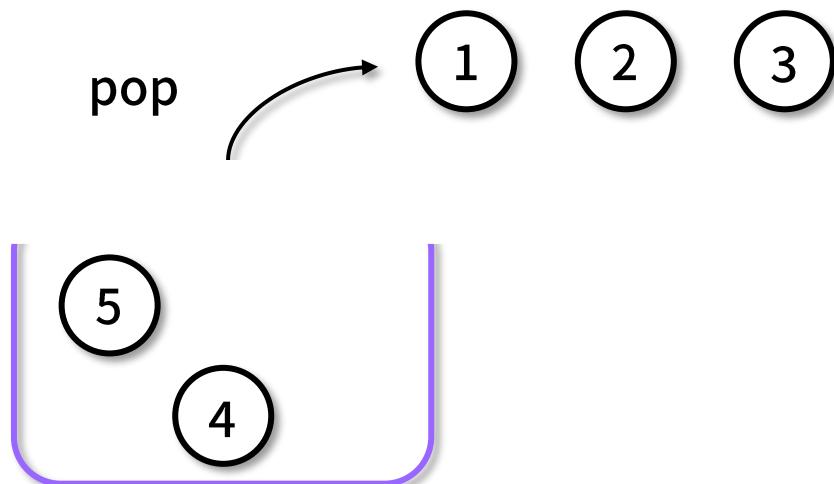
# Minimum Heap (3)

- Always pop the smallest element.



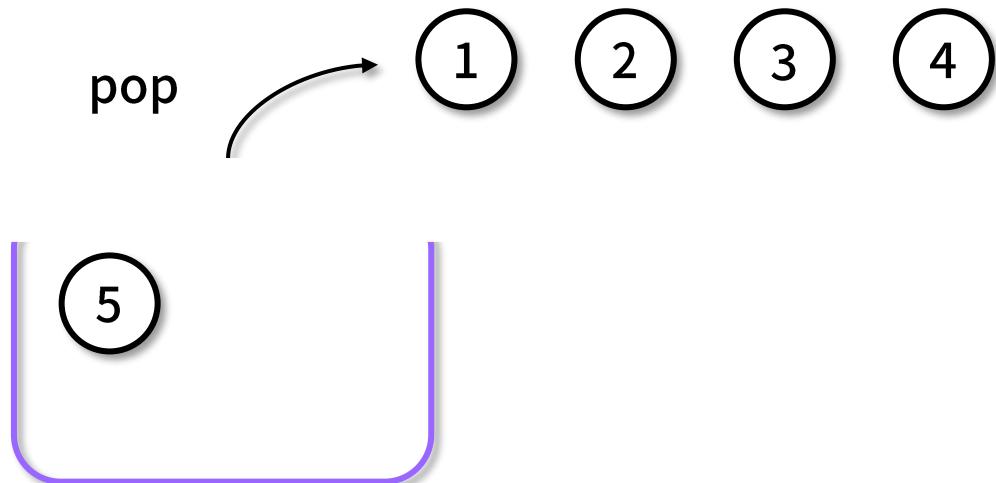
# Minimum Heap (4)

- Always pop the smallest element.



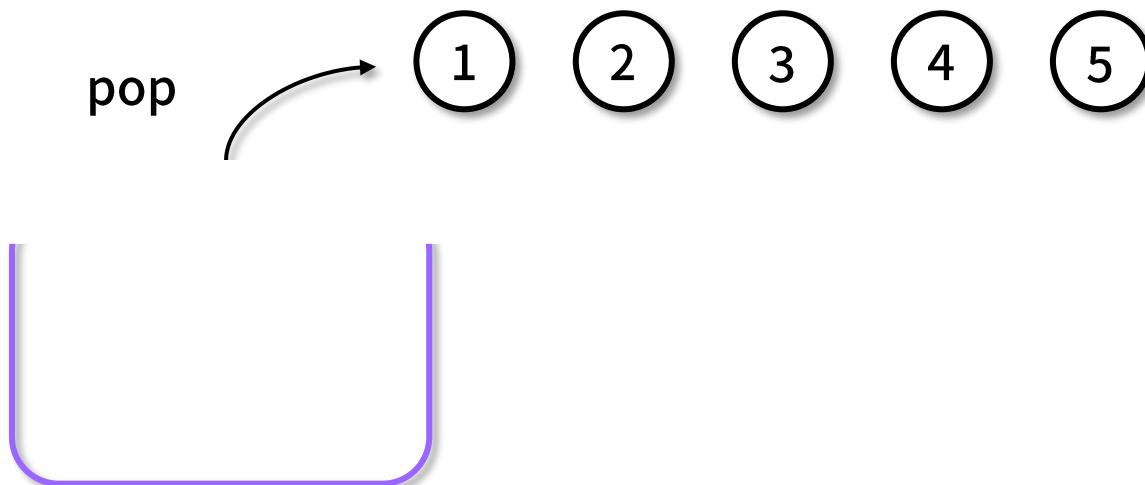
# Minimum Heap (5)

- Always pop the smallest element.



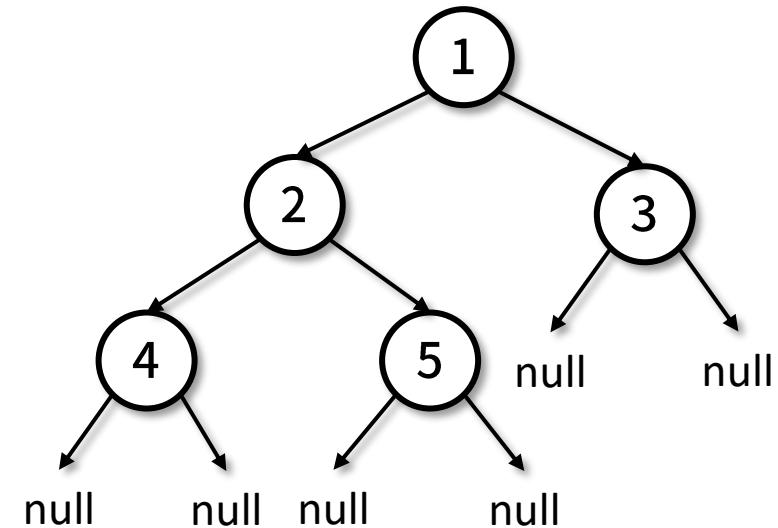
# Minimum Heap (6)

- Always pop the smallest element.



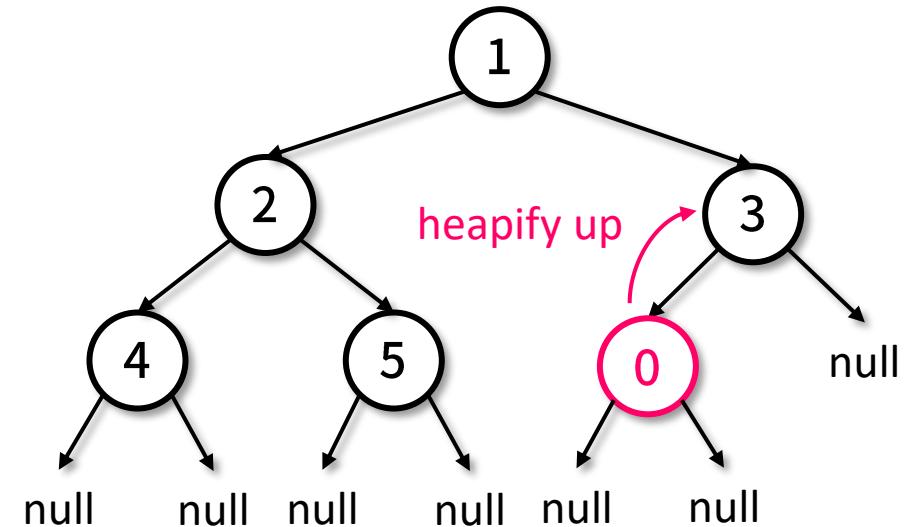
# Minimum Heap (7)

- Use a binary tree to implement the heap.
- The key of the parent node cannot be larger than the key of the children node.
- Need additional operation for *insert* and *pop* to maintain the structure to ensure the first rule is invariant.



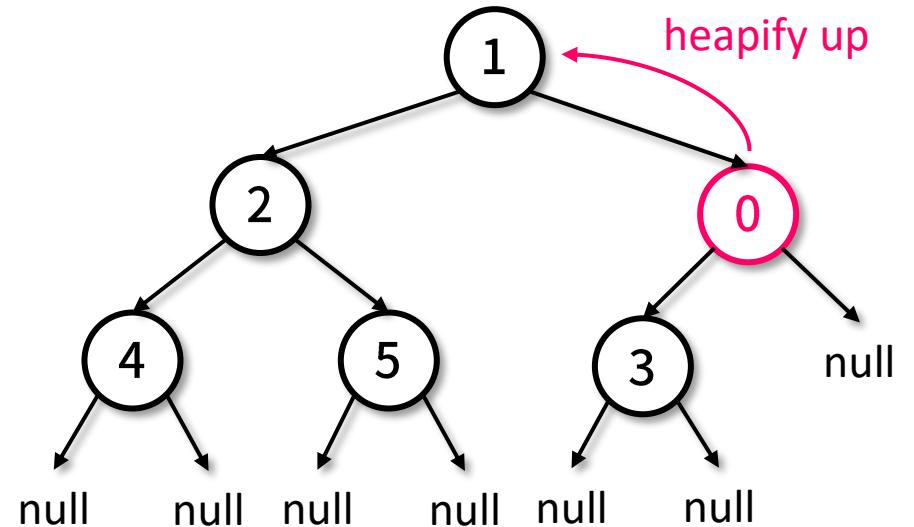
# Minimum Heap Insert(1)

1. Insert to the next empty leaf node  
(e.g. insert node with key 0).
2. Swap the inserted node up to the tree  
until smaller item is encountered  
**(heapify up)**.



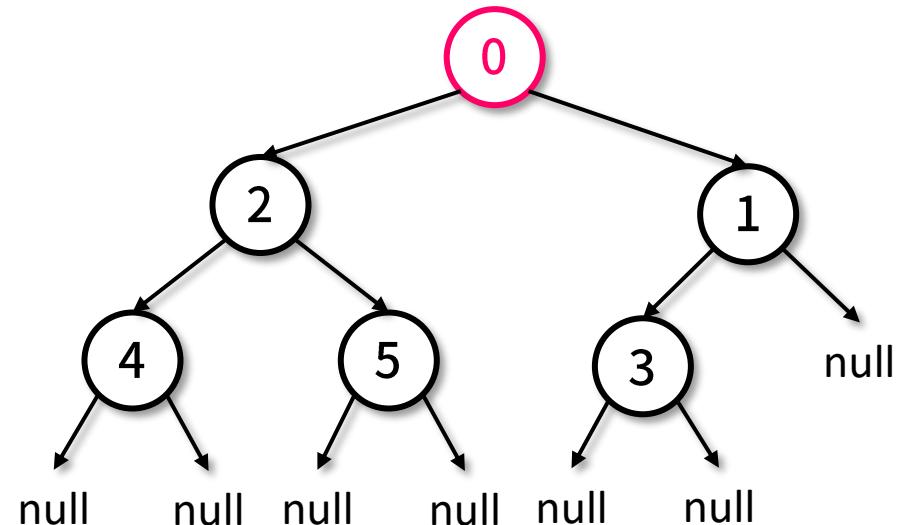
# Minimum Heap Insert(2)

1. Insert to the next empty leaf node  
(e.g. insert node with key 0).
2. Swap the inserted node up to the tree  
until smaller item is encountered  
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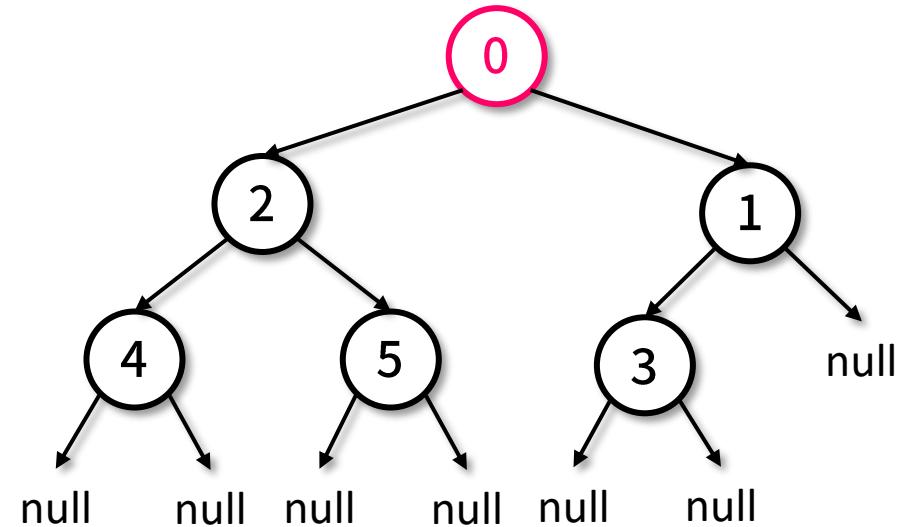
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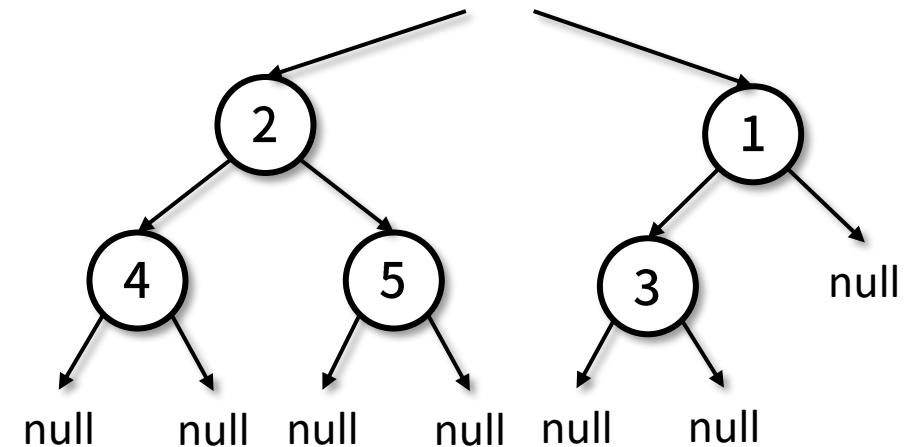
# Minimum Heap Pop (1)

- The value to be return is the value of the root node.



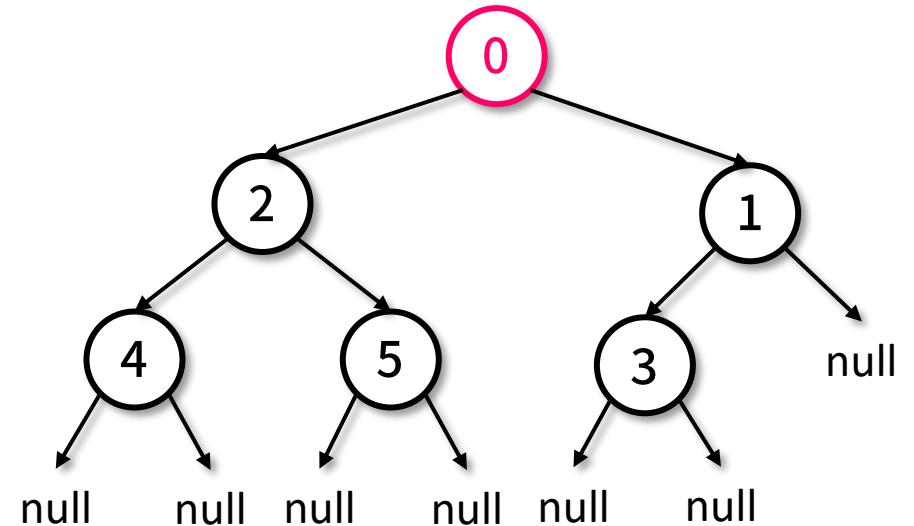
# Minimum Heap Pop (2)

- The value to be return is the value of the root node.
- Simply removing the node will break the binary tree.



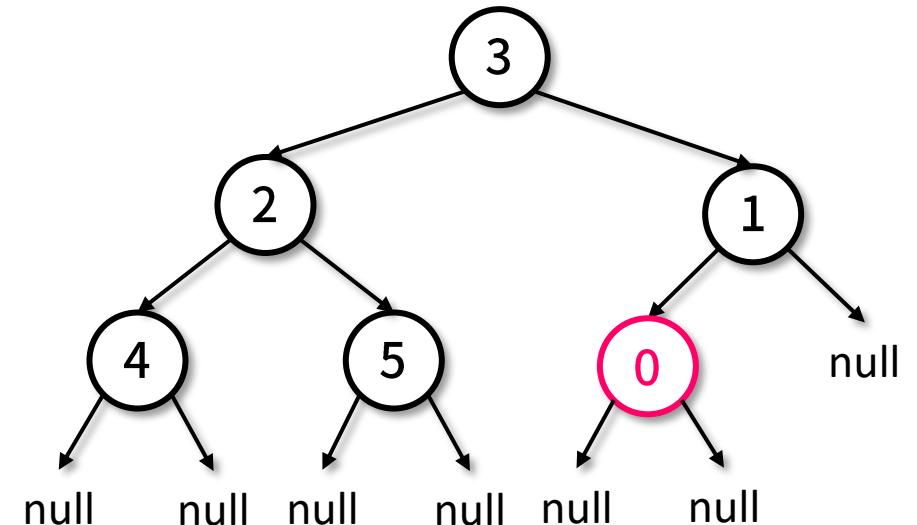
# Minimum Heap Pop (3)

- The value to be return is the value of the root node.
- Simply removing the node will break the binary tree.
- Instead, we swap with the last node (relatively large) and pop it.



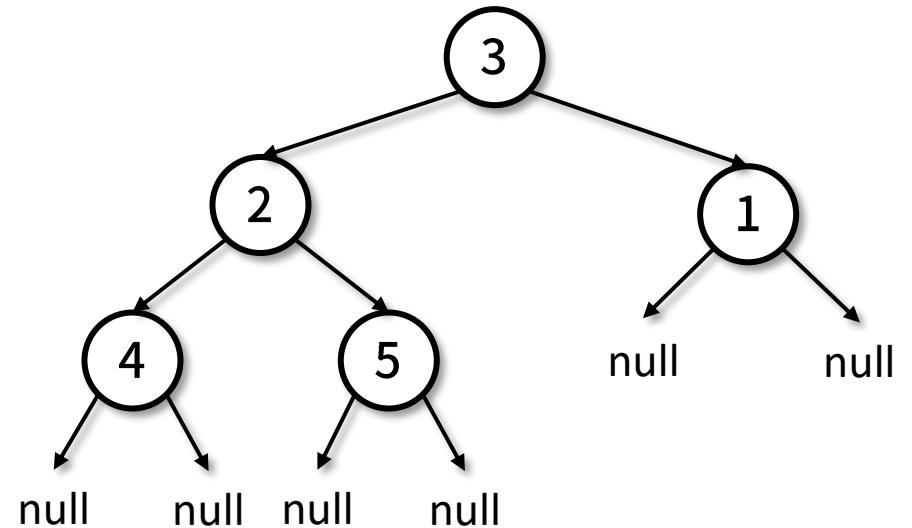
# Minimum Heap Pop (4)

- The value to be return is the value of the root node.
- Simply removing the node will break the binary tree.
- Instead, we swap with the last node (relatively large) and pop it.



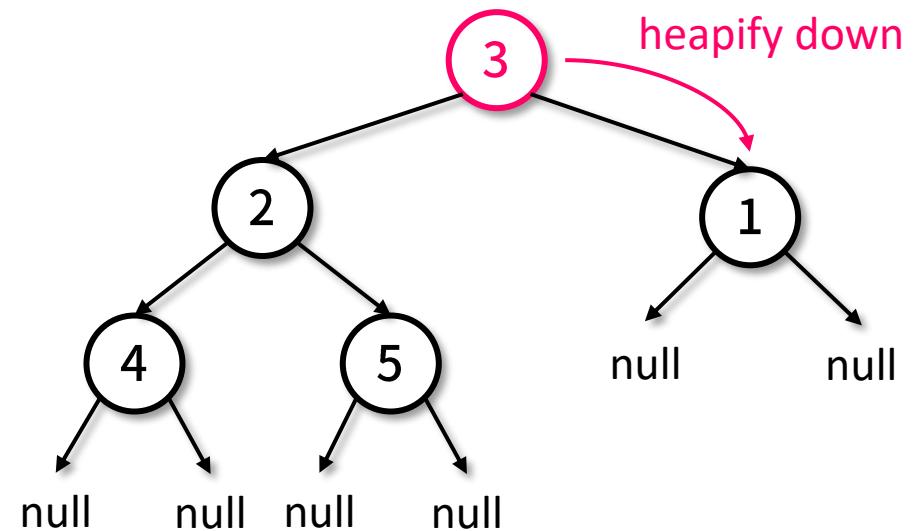
# Minimum Heap Pop (5)

- The value to be return is the value of the root node.
- Simply removing the node will break the binary tree.
- Instead, we swap with the last node (relatively large) and pop it.
- Then, we swap the root node down to the tree until larger node is encountered or reach the leaf (**heapify down**).



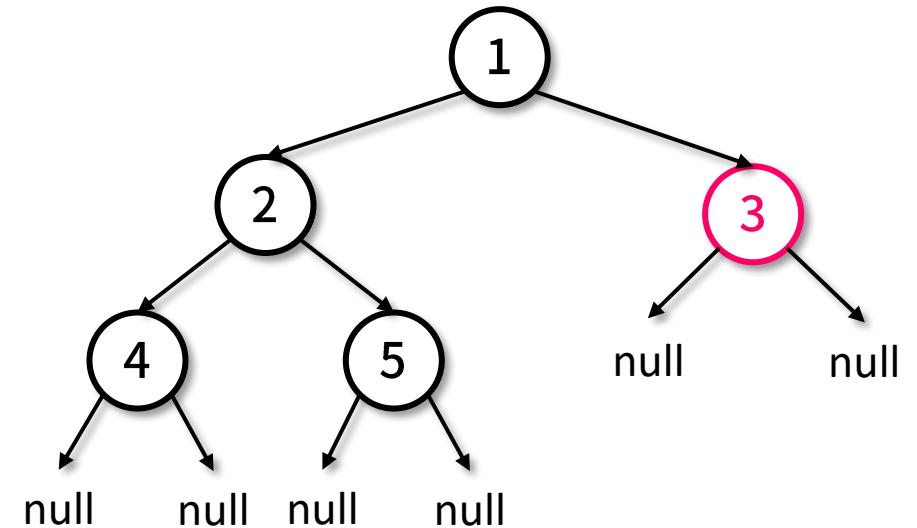
# Minimum Heap Pop (6)

- The value to be return is the value of the root node.
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- Instead, we swap with the last node (relatively large) and pop it.
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# Minimum Heap Pop (6)

- The value to be return is the value of the root node.
- Simply removing the node will break the binary tree.
- Instead, we swap with the last node (relatively large) and pop it.
- Then, we swap the root node down to the tree until larger node is encountered or reach the leaf (**heapify down**).



# Minimum Heap (8)

## Initialization

```
class MinHeap(Iterable):
    def __init__(self):
        self.data: Array = Array(1)
        self.n = 0

    def __len__(self) -> int: ...

    def __iter__(self) -> None: ...

    def insert(self, x: ComparableT) -> Iterator: ...
        insert an item

    def peek(self) -> ComparableT: ...
        check what is the minimum value in the Heap

    def pop(self) -> ComparableT: ...
        return and remove the minimum value in the Heap

    def heapify_up(self, index: int) -> None: ...
        reorganize the Heap after insertion

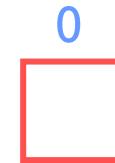
    def heapify_down(self, index: int) -> None: ...
        reorganize the Heap after deletion
```

## Command

```
heap = MinHeap()
```

## Visualization

n = 0



data

# Minimum Heap Insert(1)

## Insertion

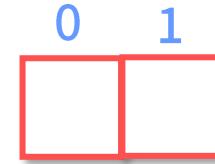
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

```
heap.insert(4)
```

## Visualization

$n = 1$



## Tree Visualization

# Minimum Heap Insert (2)

## Insertion

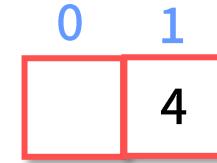
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
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    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
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    parent_index = index // 2  
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        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

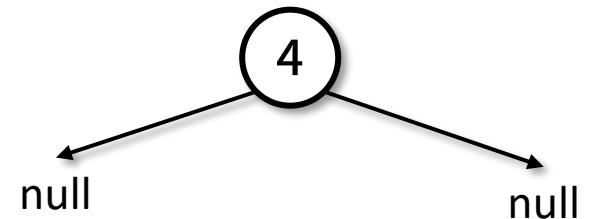
```
heap.insert(4)
```

## Visualization

$n = 1$



## Tree Visualization



# Minimum Heap Insert (3)

## Insertion

```
def insert(self, x: ComparableT) -> None:
    if self.n == len(self.data) - 1:
        self.data.resize(2 * len(self.data))

    self.n += 1
    self.data[self.n] = x
    self.heapify_up(self.n)
    return

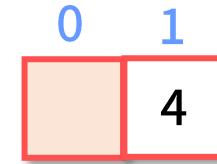
def heapify_up(self, index: int) -> None:
    parent_index = index // 2
    while index > 1 and self.data[parent_index] > self.data[index]:
        self.data.exchange(index, parent_index)
        index = parent_index
        parent_index = index // 2
    return
```

## Command

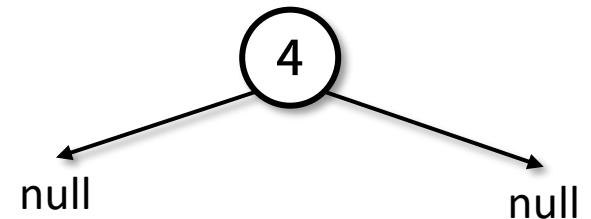
```
heap.insert(4)
```

## Visualization

$n = 1$



## Tree Visualization



# Minimum Heap Insert (4)

## Insertion

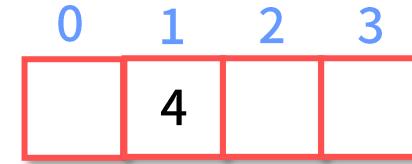
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

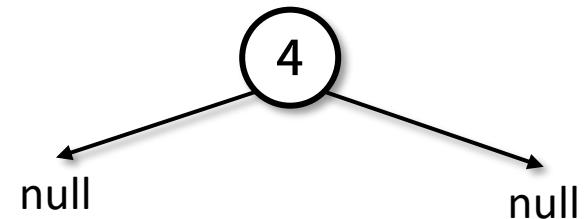
```
heap.insert(2)
```

## Visualization

$n = 1$



## Tree Visualization



# Minimum Heap Insert (5)

## Insertion

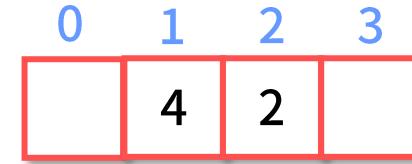
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
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    self.n += 1  
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    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

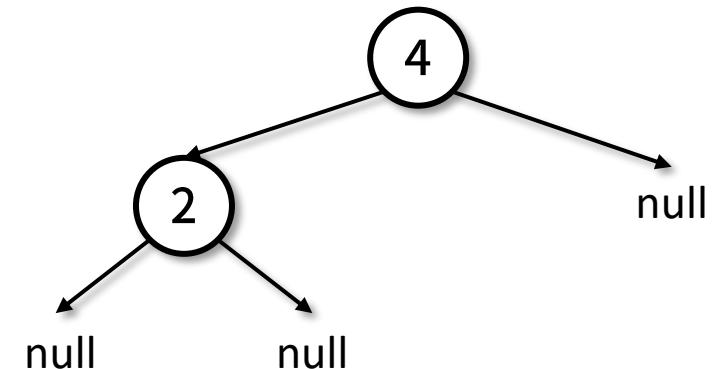
```
heap.insert(2)
```

## Visualization

$n = 2$



## Tree Visualization



# Minimum Heap Insert(6)

## Insertion

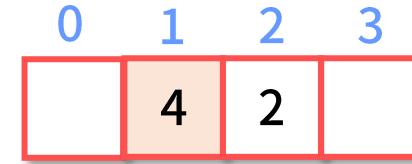
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

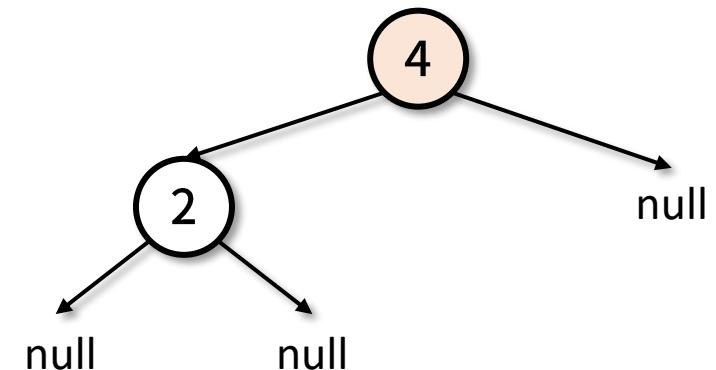
```
heap.insert(4)
```

## Visualization

$n = 2$



## Tree Visualization



# Minimum Heap Insert (7)

## Insertion

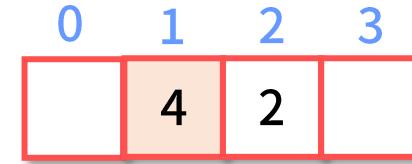
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

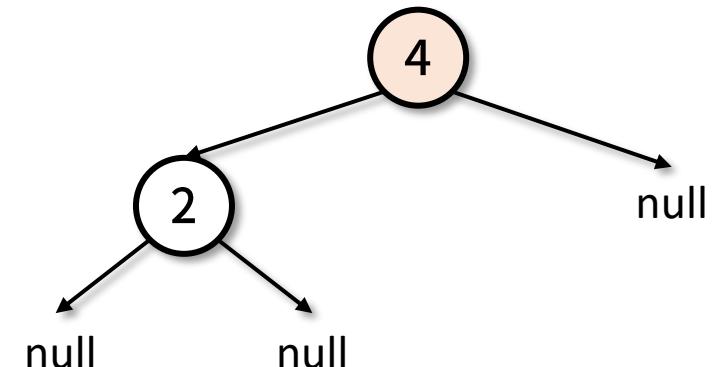
```
heap.insert(4)
```

## Visualization

$n = 2$



## Tree Visualization



# Minimum Heap Insert (8)

## Insertion

```
def insert(self, x: ComparableT) -> None:
    if self.n == len(self.data) - 1:
        self.data.resize(2 * len(self.data))

    self.n += 1
    self.data[self.n] = x
    self.heapify_up(self.n)
    return

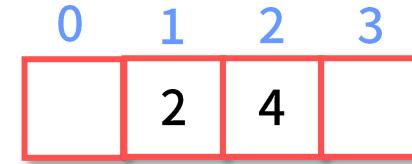
def heapify_up(self, index: int) -> None:
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        self.data.exchange(index, parent_index)
        index = parent_index
        parent_index = index // 2
    return
```

## Command

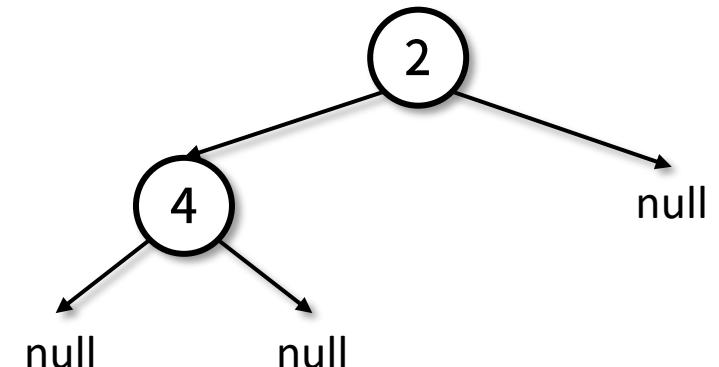
```
heap.insert(4)
```

## Visualization

$n = 2$



## Tree Visualization



# Minimum Heap Insert (9)

## Insertion

```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

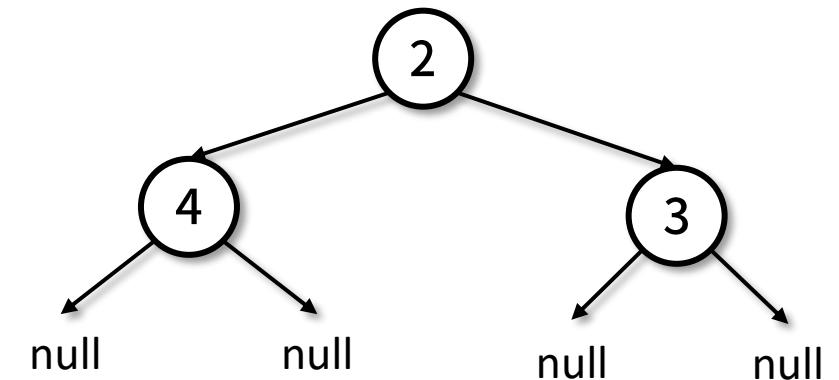
```
heap.insert(3)  
heap.insert(1)
```

## Visualization

$n = 3$

0	1	2	3
	2	4	3

## Tree Visualization



# Minimum Heap Insert(10)

## Insertion

```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

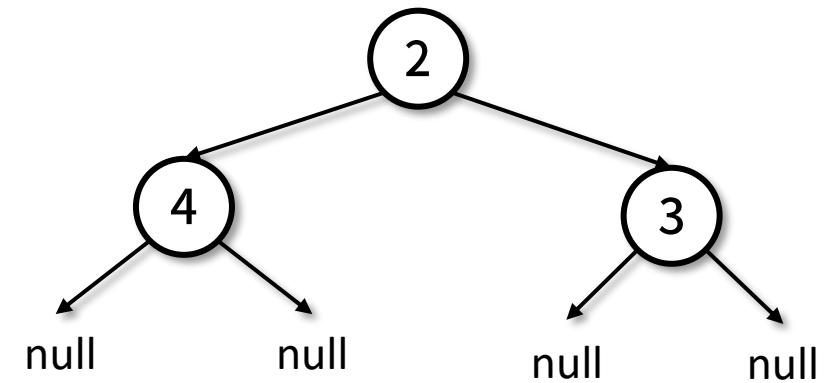
```
heap.insert(3)  
heap.insert(1)
```

## Visualization

$n = 3$

0	1	2	3	4	5	6	7
		2	4	3			

## Tree Visualization



# Minimum Heap Insert(11)

## Insertion

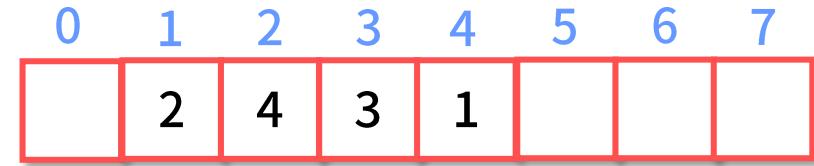
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

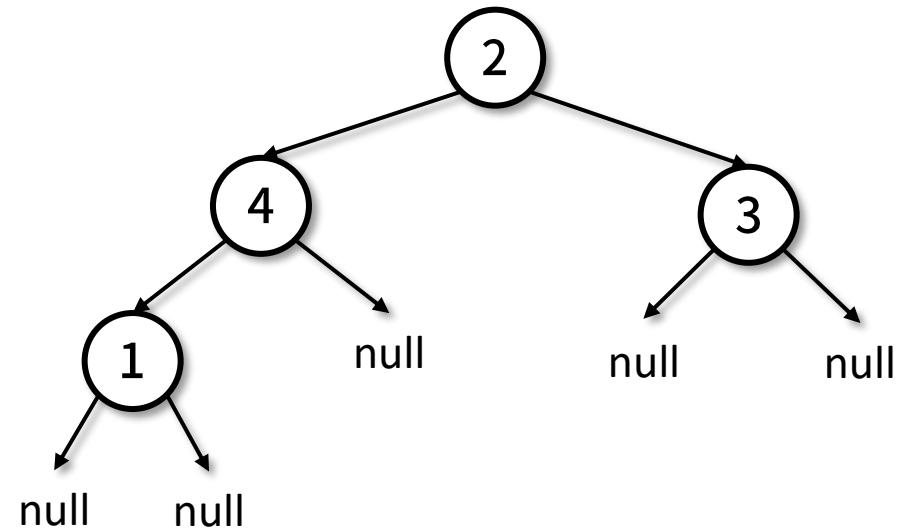
```
heap.insert(3)  
heap.insert(1)
```

## Visualization

$n = 4$



## Tree Visualization



# Minimum Heap Insert(12)

## Insertion

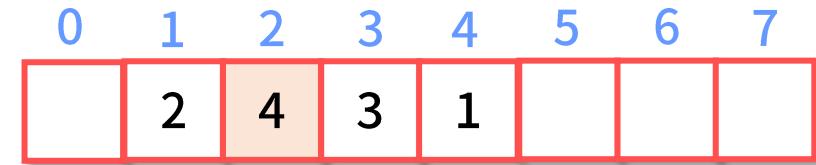
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

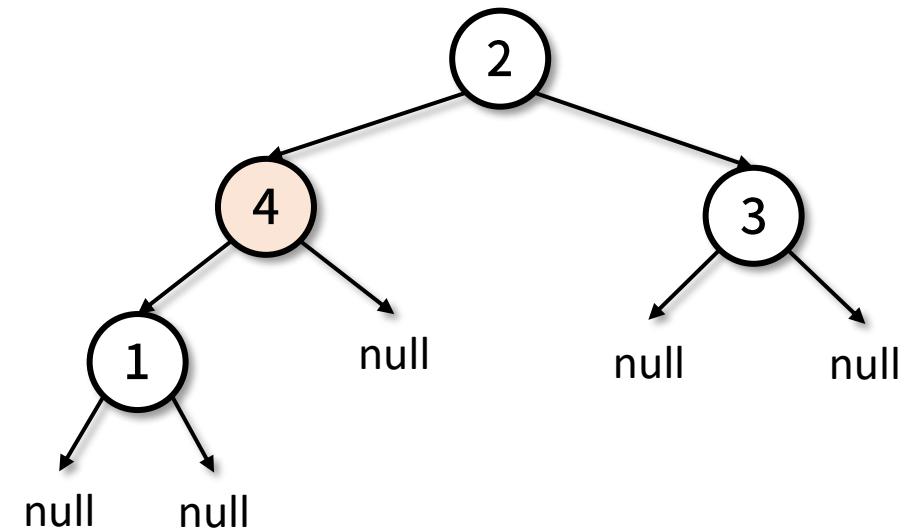
```
heap.insert(3)  
heap.insert(1)
```

## Visualization

$n = 4$



## Tree Visualization



# Minimum Heap Insert (13)

## Insertion

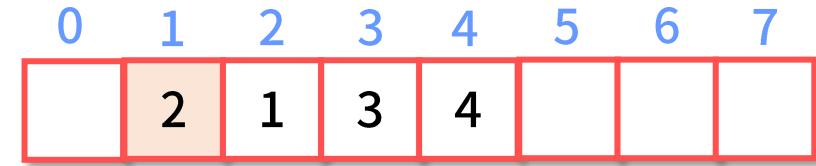
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

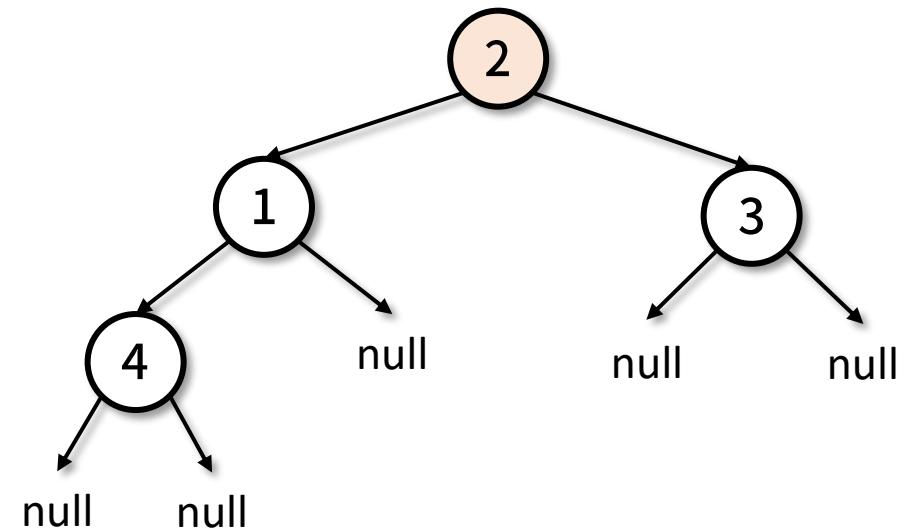
```
heap.insert(3)  
heap.insert(1)
```

## Visualization

$n = 4$



## Tree Visualization



# Minimum Heap Insert (14)

## Insertion

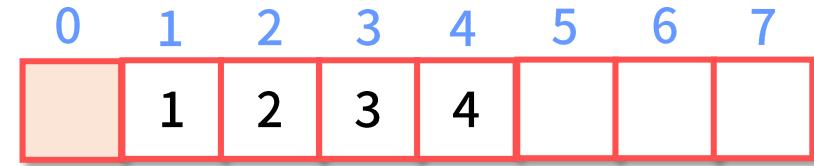
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

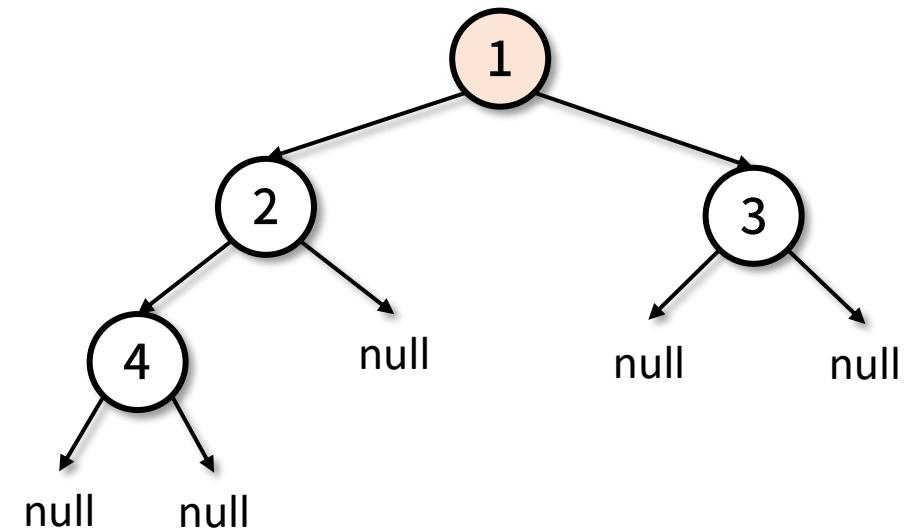
```
heap.insert(3)  
heap.insert(1)
```

## Visualization

$n = 4$



## Tree Visualization



# Minimum Heap Insert (15)

## Insertion

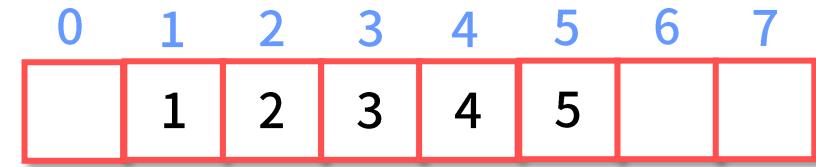
```
def insert(self, x: ComparableT) -> None:  
    if self.n == len(self.data) - 1:  
        self.data.resize(2 * len(self.data))  
  
    self.n += 1  
    self.data[self.n] = x  
    self.heapify_up(self.n)  
    return  
  
def heapify_up(self, index: int) -> None:  
    parent_index = index // 2  
    while index > 1 and self.data[parent_index] > self.data[index]:  
        self.data.exchange(index, parent_index)  
        index = parent_index  
        parent_index = index // 2  
    return
```

## Command

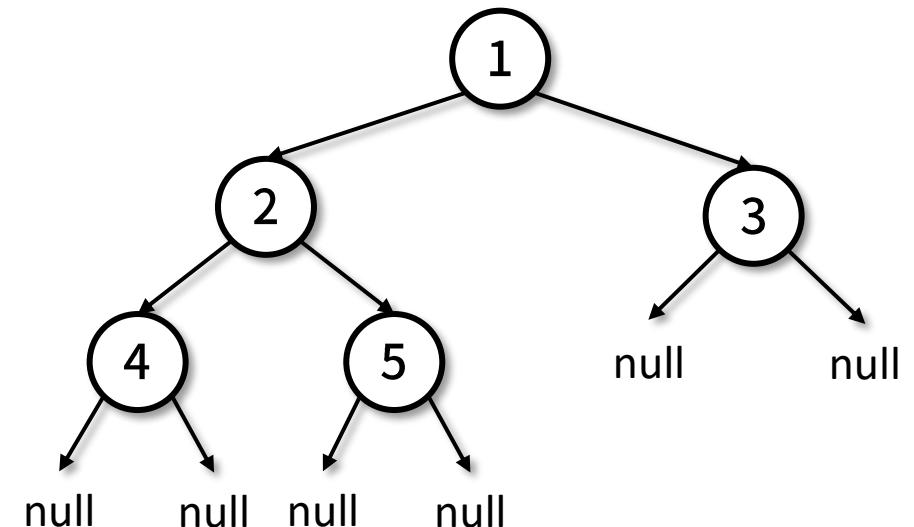
```
heap.insert(5)
```

## Visualization

$n = 5$



## Tree Visualization



# Minimum Heap Pop (1)

## Pop

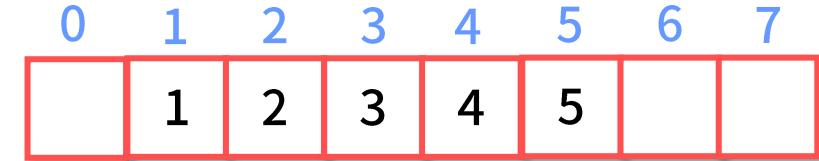
```
def pop(self) -> ComparableT:  
    result = self.data[1]  
    self.data.exchange(1, self.n)  
    self.n -= 1  
    self.heapify_down(1)  
    self.data[self.n + 1] = None  
    if self.n > 0 and self.n <= (len(self.data) - 1) / 4:  
        self.data.resize(int(0.5 * len(self.data)))  
    return result
```

## Command

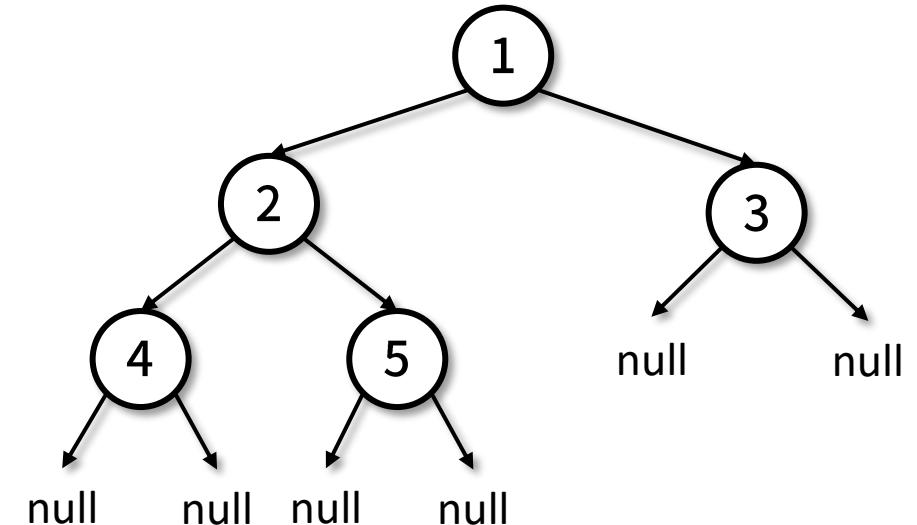
```
heap.pop()
```

## Visualization

$n = 5 \quad \text{result} = 1$



## Tree Visualization



# Minimum Heap Pop (2)

## Pop

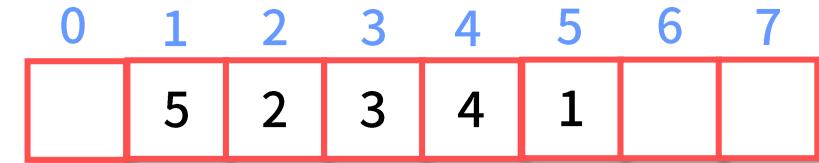
```
def pop(self) -> ComparableT:  
    result = self.data[1]  
    self.data.exchange(1, self.n)  
    self.n -= 1  
    self.heapify_down(1)  
    self.data[self.n + 1] = None  
    if self.n > 0 and self.n <= (len(self.data) - 1) / 4:  
        self.data.resize(int(0.5 * len(self.data)))  
    return result
```

## Command

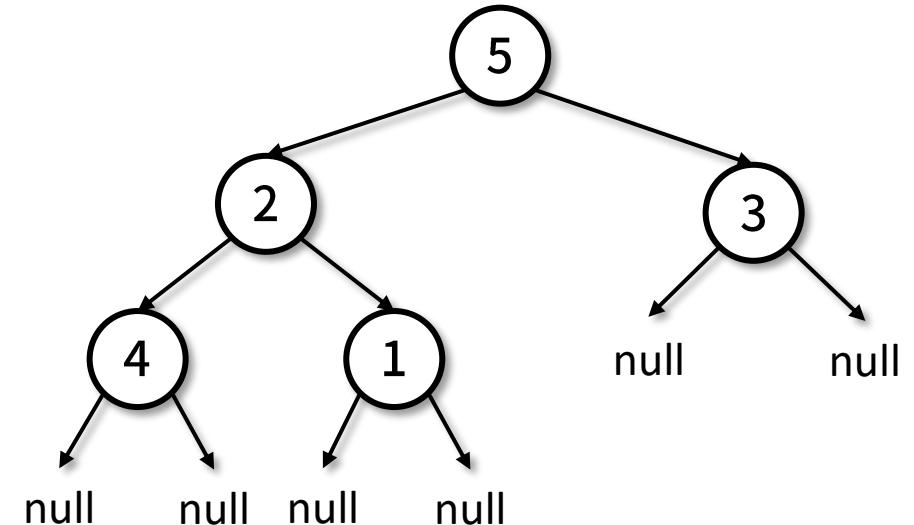
```
heap.pop()
```

## Visualization

$n = 5 \quad \text{result} = 1$



## Tree Visualization



# Minimum Heap Pop (3)

## Pop

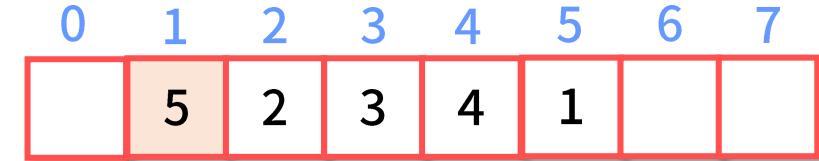
```
def pop(self) -> ComparableT:  
    result = self.data[1]  
    self.data.exchange(1, self.n)  
    self.n -= 1  
    self.heapify_down(1)  
    self.data[self.n + 1] = None  
    if self.n > 0 and self.n <= (len(self.data) - 1) / 4:  
        self.data.resize(int(0.5 * len(self.data)))  
    return result
```

## Command

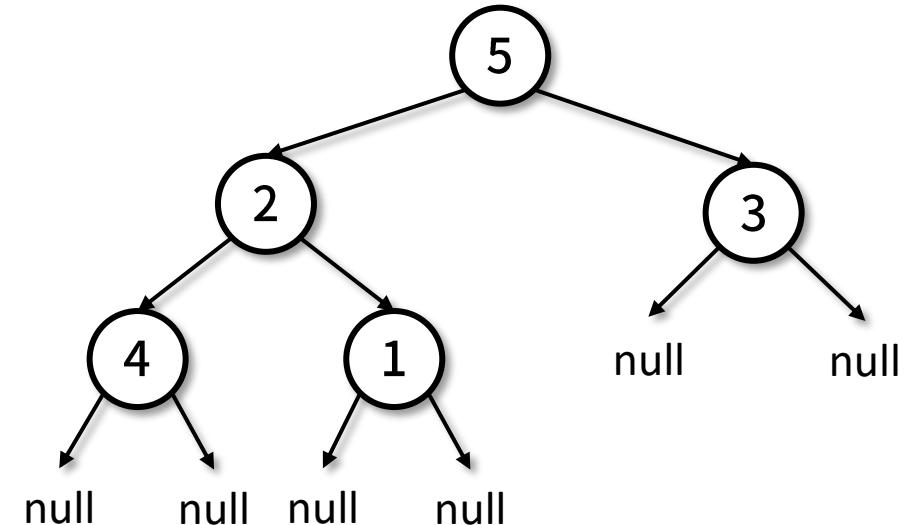
```
heap.pop()
```

## Visualization

$n = 4 \quad \text{result} = 1$



## Tree Visualization



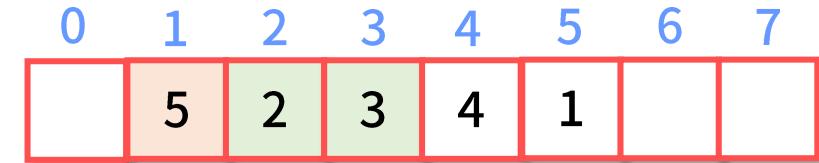
# Minimum Heap Pop (4)

## Pop

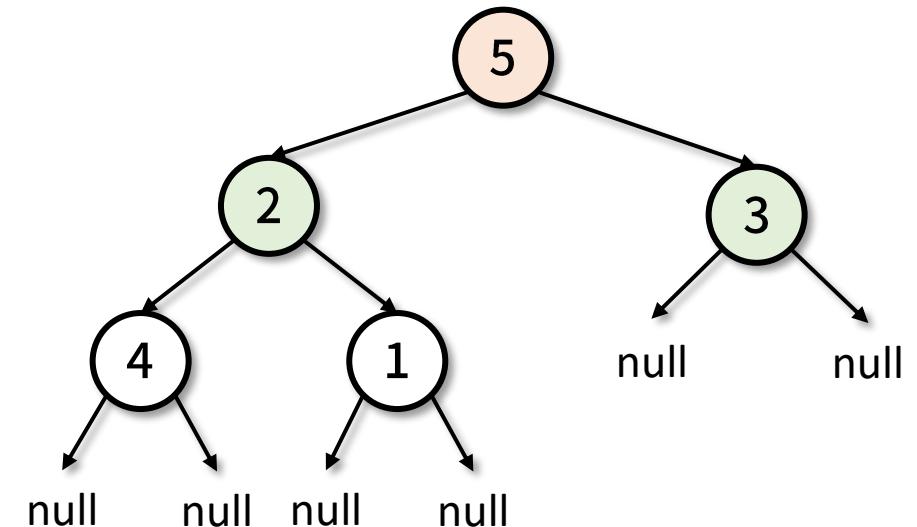
```
def heapify_down(self, index: int) -> None:
    while 2 * index <= self.n:
        child_a_index = index * 2
        child_b_index = index * 2 + 1
        if child_b_index > self.n:
            exchange_to = child_a_index
        elif self.data[child_b_index] < self.data[child_a_index]:
            exchange_to = child_b_index
        else:
            exchange_to = child_a_index
        if self.data[index] <= self.data[exchange_to]:
            break
        self.data.exchange(index, exchange_to)
        index = exchange_to
    return
```

## Visualization

$n = 4$  result = 1



## Tree Visualization



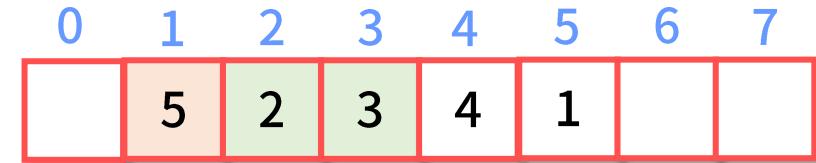
# Minimum Heap Pop (5)

## Pop

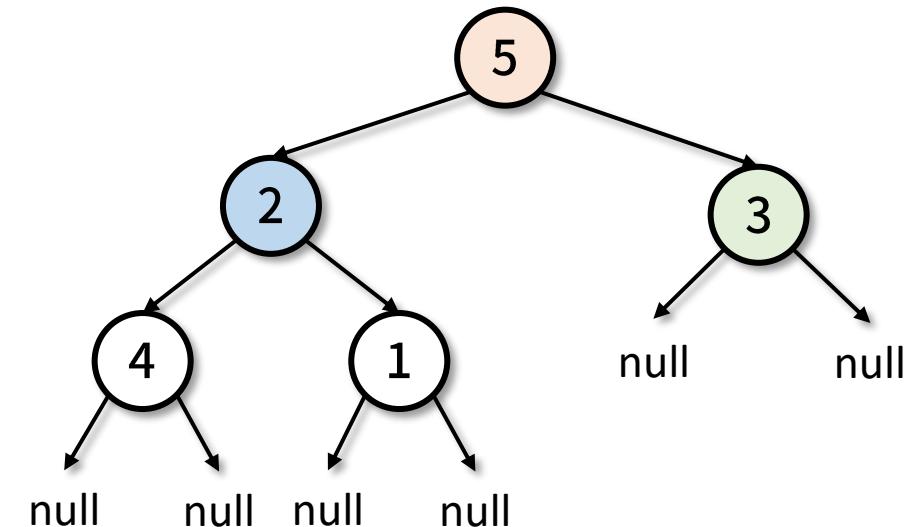
```
def heapify_down(self, index: int) -> None:  
    while 2 * index <= self.n:  
        child_a_index = index * 2  
        child_b_index = index * 2 + 1  
        if child_b_index > self.n:  
            exchange_to = child_a_index  
        elif self.data[child_b_index] < self.data[child_a_index]:  
            exchange_to = child_b_index  
        else:  
            exchange_to = child_a_index  
        if self.data[index] <= self.data[exchange_to]:  
            break  
        self.data.exchange(index, exchange_to)  
        index = exchange_to  
    return
```

## Visualization

$n = 4 \quad \text{result} = 1$



## Tree Visualization



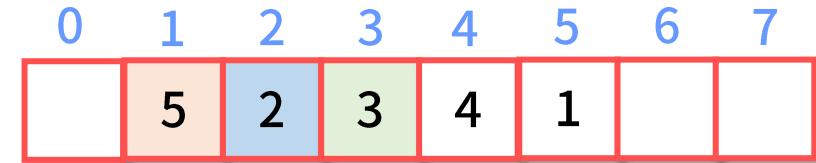
# Minimum Heap Pop (6)

## Pop

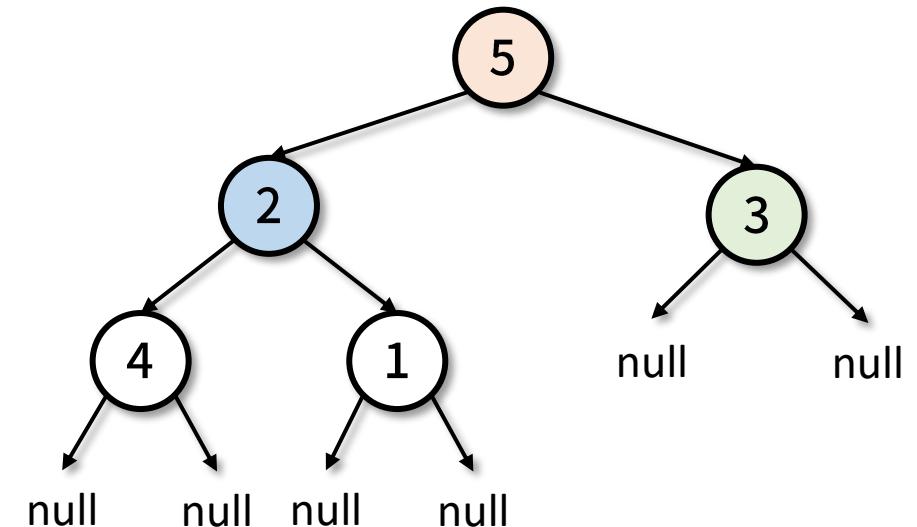
```
def heapify_down(self, index: int) -> None:  
    while 2 * index <= self.n:  
        child_a_index = index * 2  
        child_b_index = index * 2 + 1  
        if child_b_index > self.n:  
            exchange_to = child_a_index  
        elif self.data[child_b_index] < self.data[child_a_index]:  
            exchange_to = child_b_index  
        else:  
            exchange_to = child_a_index  
        if self.data[index] <= self.data[exchange_to]:  
            break  
        self.data.exchange(index, exchange_to)  
        index = exchange_to  
  
    return
```

## Visualization

$n = 4$  result = 1



## Tree Visualization



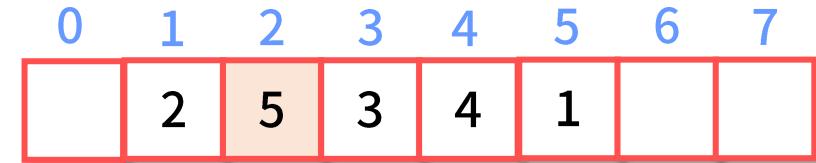
# Minimum Heap Pop (7)

## Pop

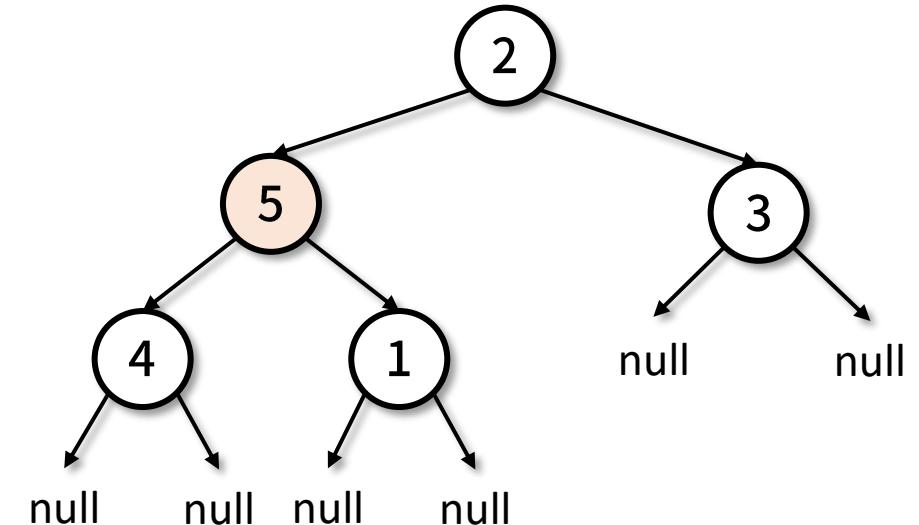
```
def heapify_down(self, index: int) -> None:  
    while 2 * index <= self.n:  
        child_a_index = index * 2  
        child_b_index = index * 2 + 1  
        if child_b_index > self.n:  
            exchange_to = child_a_index  
        elif self.data[child_b_index] < self.data[child_a_index]:  
            exchange_to = child_b_index  
        else:  
            exchange_to = child_a_index  
        if self.data[index] <= self.data[exchange_to]:  
            break  
        self.data.exchange(index, exchange_to)  
        index = exchange_to  
  
    return
```

## Visualization

$n = 4$  result = 1



## Tree Visualization



# Minimum Heap Pop (7)

## Pop

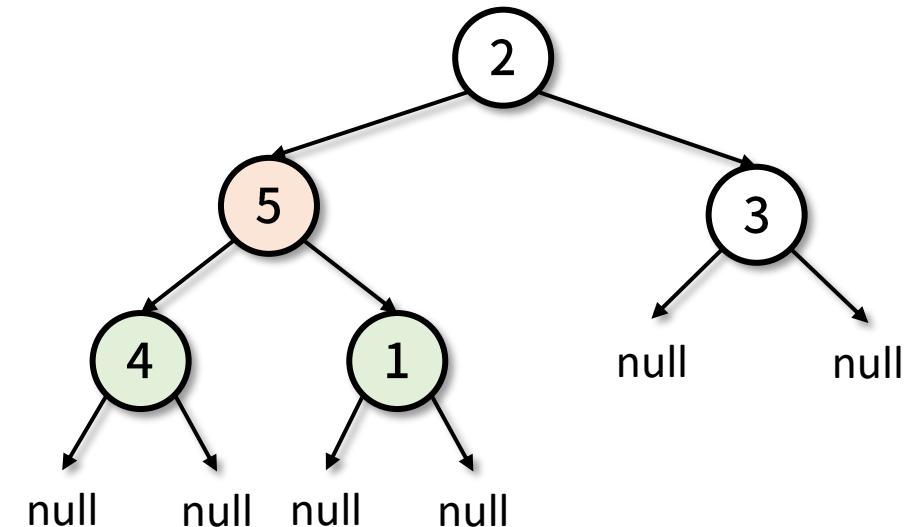
```
def heapify_down(self, index: int) -> None:
    while 2 * index <= self.n:
        child_a_index = index * 2
        child_b_index = index * 2 + 1
        if child_b_index > self.n:
            exchange_to = child_a_index
        elif self.data[child_b_index] < self.data[child_a_index]:
            exchange_to = child_b_index
        else:
            exchange_to = child_a_index
        if self.data[index] <= self.data[exchange_to]:
            break
        self.data.exchange(index, exchange_to)
        index = exchange_to
    return
```

## Visualization

$n = 4$  result = 1



## Tree Visualization



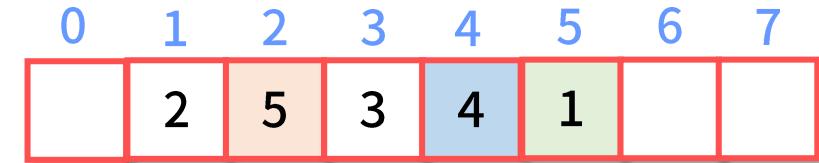
# Minimum Heap Pop (8)

## Pop

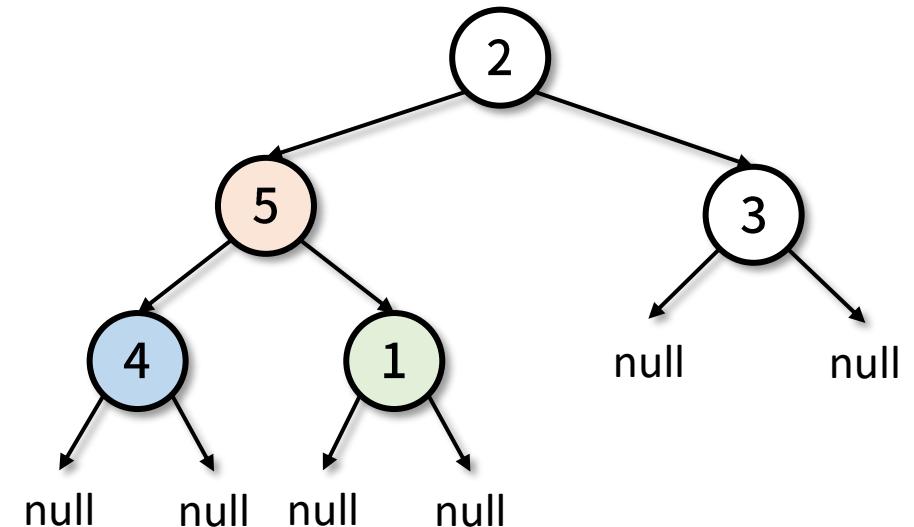
```
def heapify_down(self, index: int) -> None:  
    while 2 * index <= self.n:  
        child_a_index = index * 2  
        child_b_index = index * 2 + 1  
        if child_b_index > self.n:  
            exchange_to = child_a_index  
        elif self.data[child_b_index] < self.data[child_a_index]:  
            exchange_to = child_b_index  
        else:  
            exchange_to = child_a_index  
        if self.data[index] <= self.data[exchange_to]:  
            break  
        self.data.exchange(index, exchange_to)  
        index = exchange_to  
    return
```

## Visualization

$n = 4 \quad \text{result} = 1$



## Tree Visualization



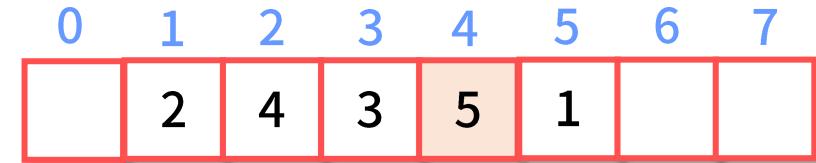
# Minimum Heap Pop (9)

## Pop

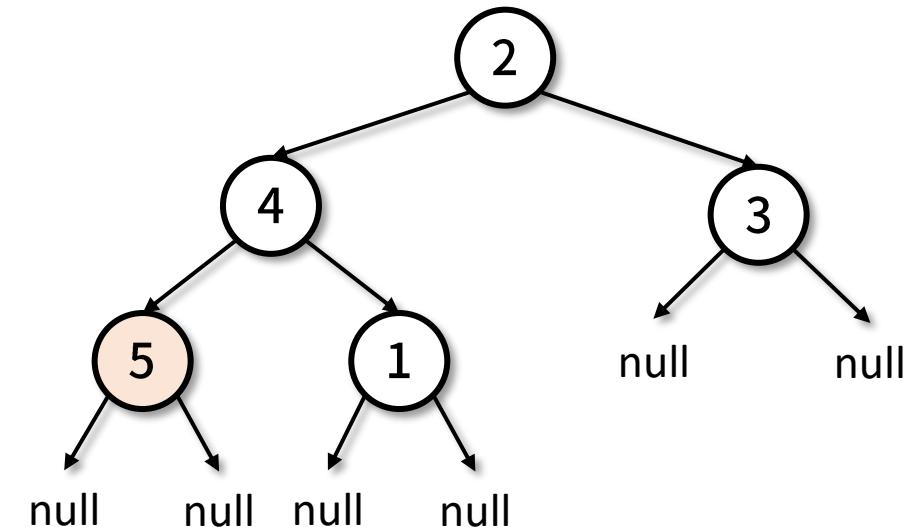
```
def heapify_down(self, index: int) -> None:  
    while 2 * index <= self.n:  
        child_a_index = index * 2  
        child_b_index = index * 2 + 1  
        if child_b_index > self.n:  
            exchange_to = child_a_index  
        elif self.data[child_b_index] < self.data[child_a_index]:  
            exchange_to = child_b_index  
        else:  
            exchange_to = child_a_index  
        if self.data[index] <= self.data[exchange_to]:  
            break  
        self.data.exchange(index, exchange_to)  
        index = exchange_to  
  
    return
```

## Visualization

$n = 4 \quad \text{result} = 1$



## Tree Visualization



# Minimum Heap Pop (10)

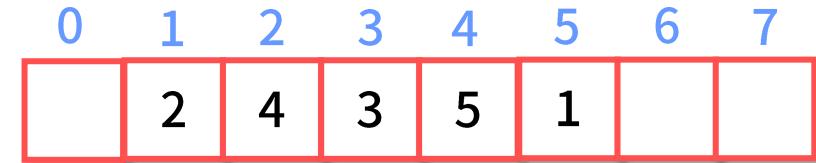
## Pop

```
def heapify_down(self, index: int) -> None:  
    while 2 * index <= self.n:  
        child_a_index = index * 2  
        child_b_index = index * 2 + 1  
        if child_b_index > self.n:  
            exchange_to = child_a_index  
        elif self.data[child_b_index] < self.data[child_a_index]:  
            exchange_to = child_b_index  
        else:  
            exchange_to = child_a_index  
        if self.data[index] <= self.data[exchange_to]:  
            break  
        self.data.exchange(index, exchange_to)  
        index = exchange_to  
    return
```

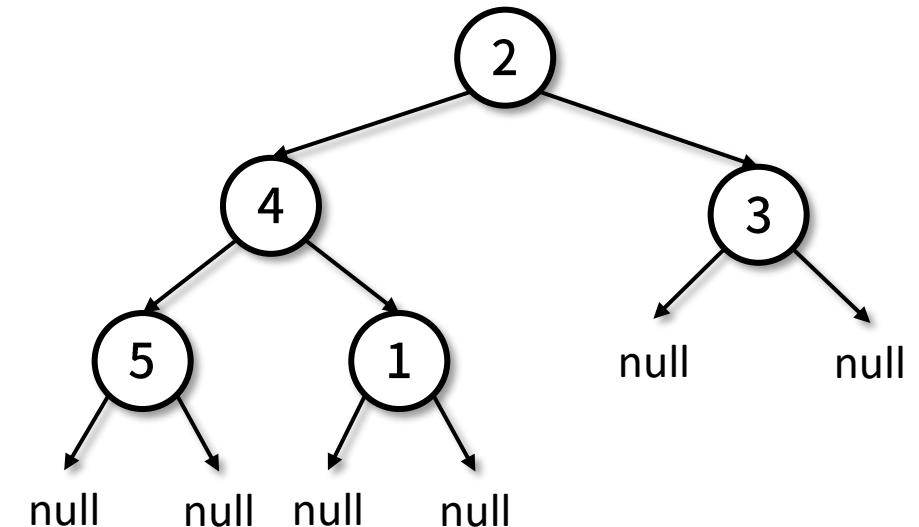
can we change the second if statement to the opposite?

## Visualization

$n = 4$  result = 1



## Tree Visualization



# Minimum Heap Pop (11)

## Pop

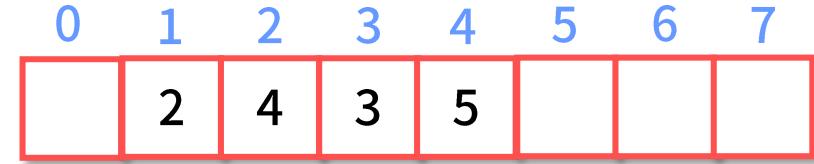
```
def pop(self) -> ComparableT:  
    result = self.data[1]  
    self.data.exchange(1, self.n)  
    self.n -= 1  
    self.heapify_down(1)  
    self.data[self.n + 1] = None  
    if self.n > 0 and self.n <= (len(self.data) - 1) / 4:  
        self.data.resize(int(0.5 * len(self.data)))  
    return result
```

## Command

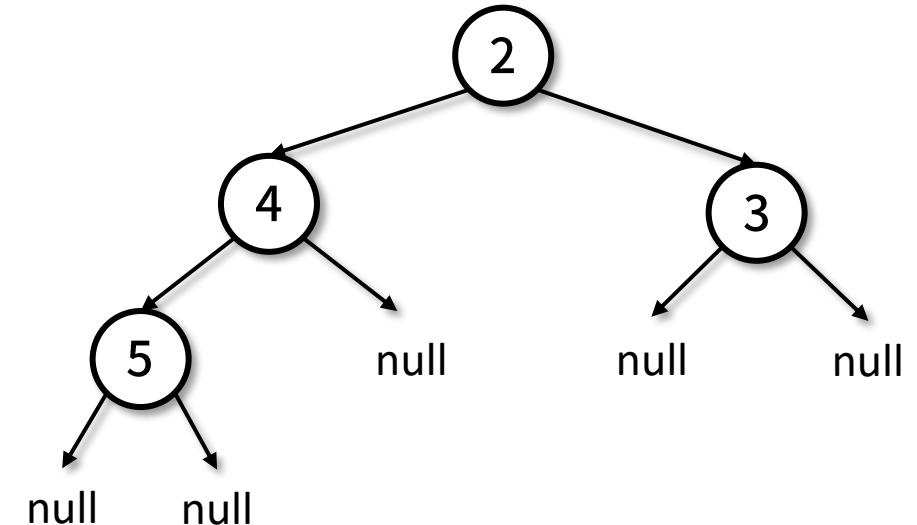
```
heap.pop()
```

## Visualization

$n = 4 \quad \text{result} = 1$



## Tree Visualization



# Heap Summary

## Worst Time Complexity

	Heap	Sorted Array
Insert	$O(\log N)^*$	$O(N)$
Pop	$O(\log N)^*$	$O(1)^*$

\*amortized time complexity

# Minimum Heap Practices

- Maximum Product of Two Elements in an Array (Leetcode Problem 1464)
- Kth Largest Element in an Array (Leetcode Problem 215)

# Binary Search Tree

# Binary Search

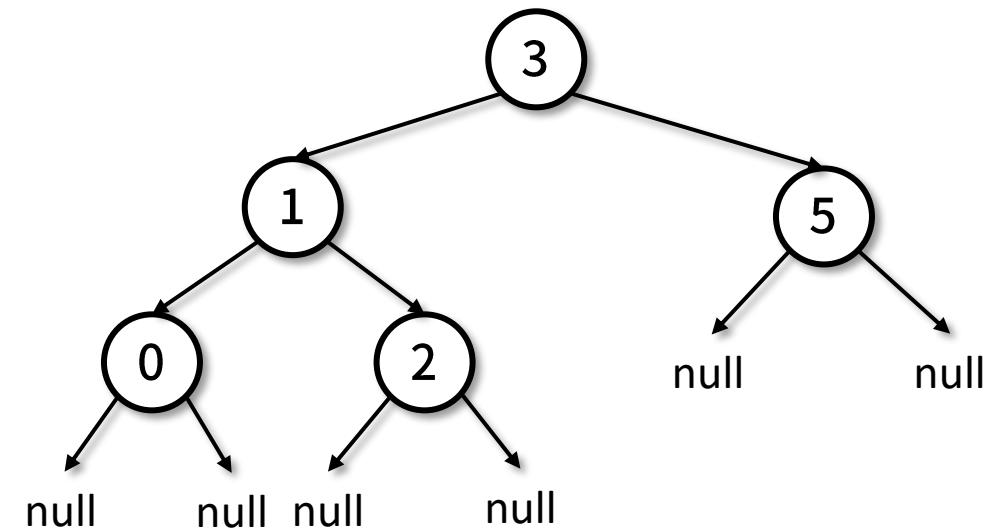
## Function

```
def binary_search(sorted_array, item, start: int, end: int):
    if end < start:
        return -1
    middle = (start + end) // 2
    if sorted_array[middle] == item:
        return middle
    elif sorted_array[middle] > item:
        return binary_search(sorted_array, item, start, middle-1)
    elif sorted_array[middle] < item:
        return binary_search(sorted_array, item, middle+1, end)
```

- Given a sorted array, we could search an item in logarithm time complexity.
- However, insert new element into a sorted array takes linear time.
- Binary search tree is a data structure that uses the idea of binary search to ensure logarithm time complexity for *search* and *insertion*.
- It also support arbitrary comparable key.

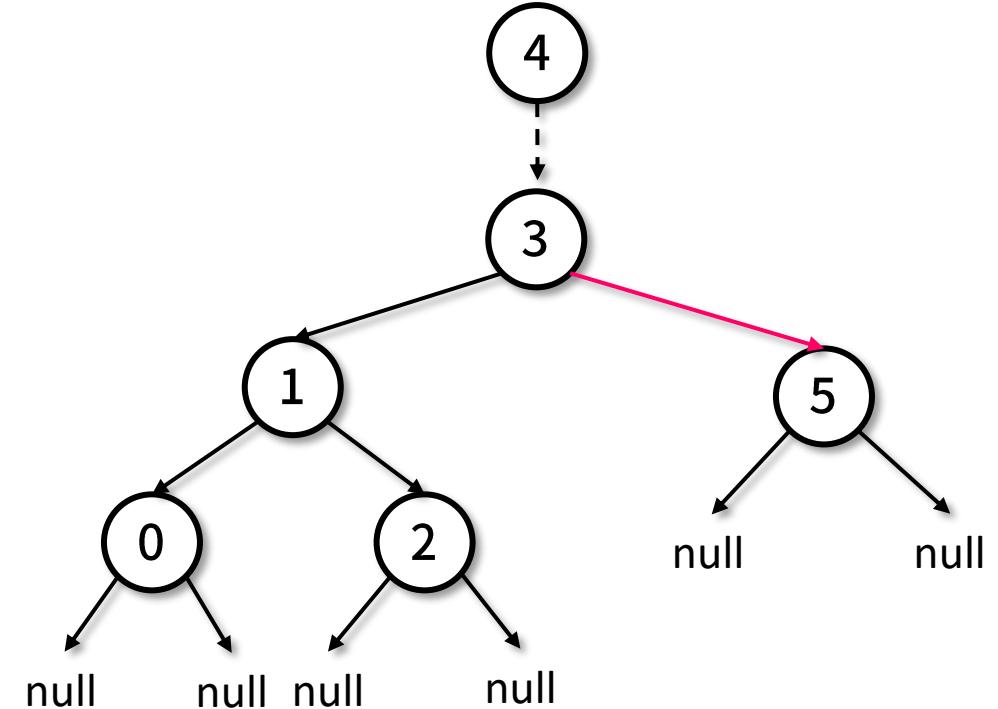
# Binary Search Tree (1)

- The key of the parent node cannot be larger than any of the key node in the right subtree.
- The key of the parent node cannot be smaller than any of the key node in the left subtree.
- Need additional operation for *insert* and *delete* to maintain the structure to ensure the first rule is invariant.



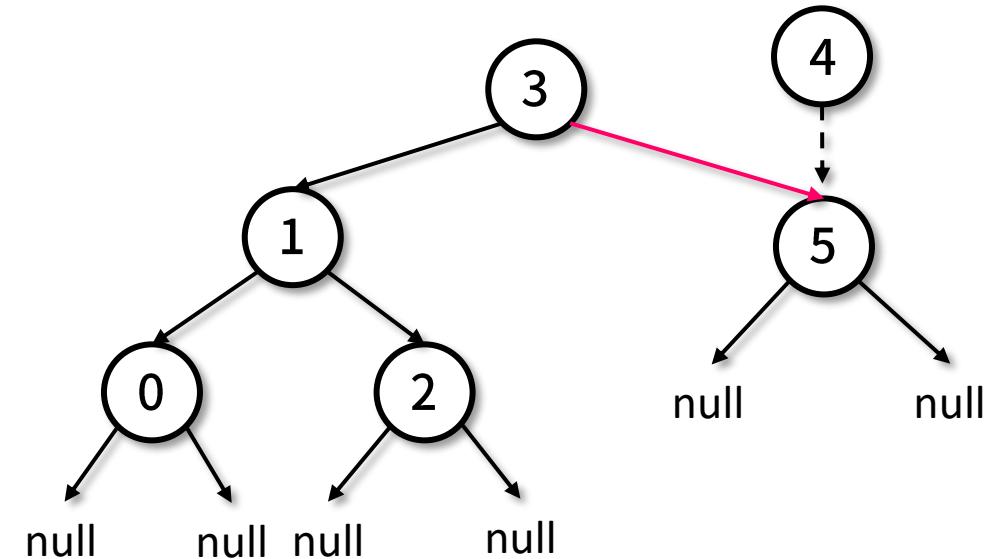
# Binary Search Tree Insert (1)

1. From the root, if the inserted node with key (e.g. 4) is larger than the root, check the right subtree.
2. If the inserted node with key is smaller than the root, check the left subtree.
3. Iterate until the same key is found, or reach the.



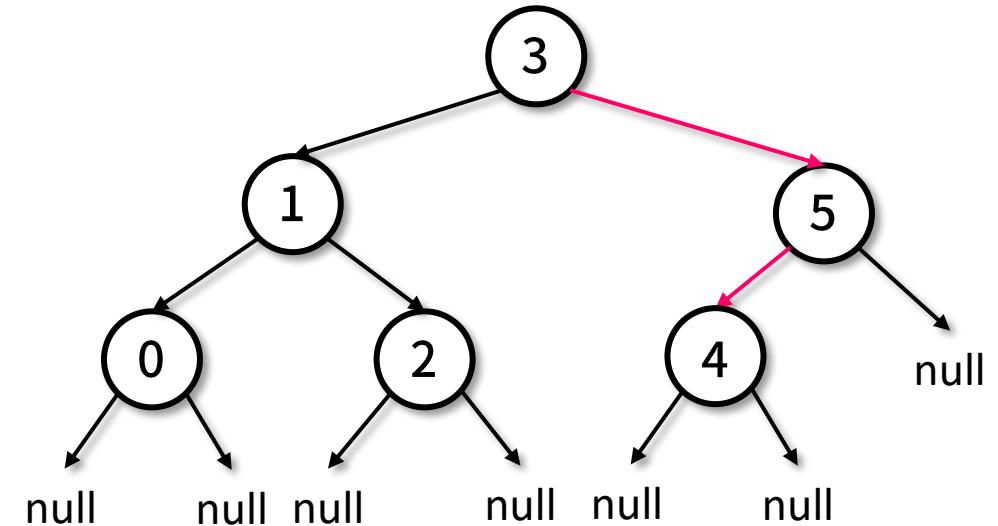
# Binary Search Tree Insert (2)

1. From the root, if the inserted node with key (e.g. 4) is larger than the root, check the right subtree.
2. If the inserted node with key is smaller than the root, check the left subtree.
3. Iterate until the same key is found, or reach the.



# Binary Search Tree Insert (3)

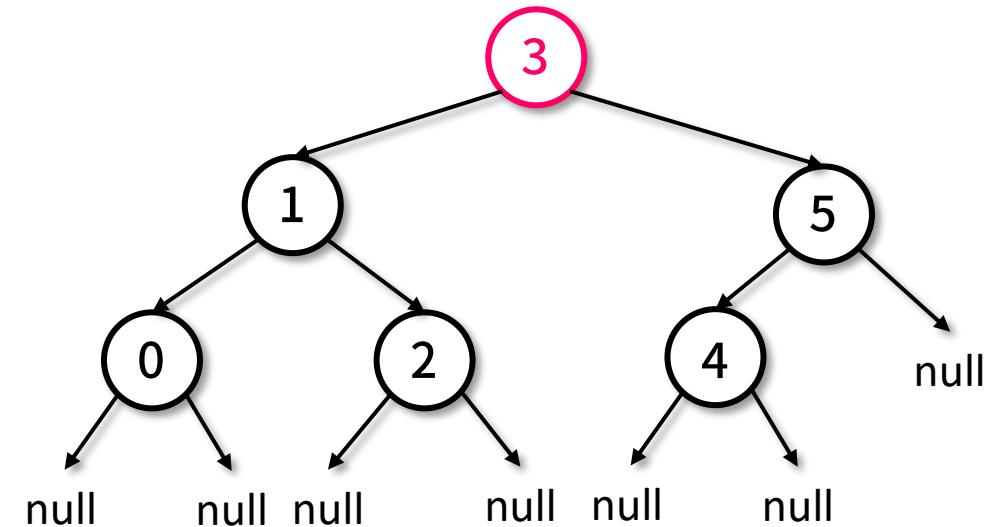
1. From the root, if the inserted node with key (e.g. 4) is larger than the root, check the right subtree.
2. If the inserted node with key is smaller than the root, check the left subtree.
3. Iterate until the same key is found, or reach the.



Binary Search Tree

# Deletion (1)

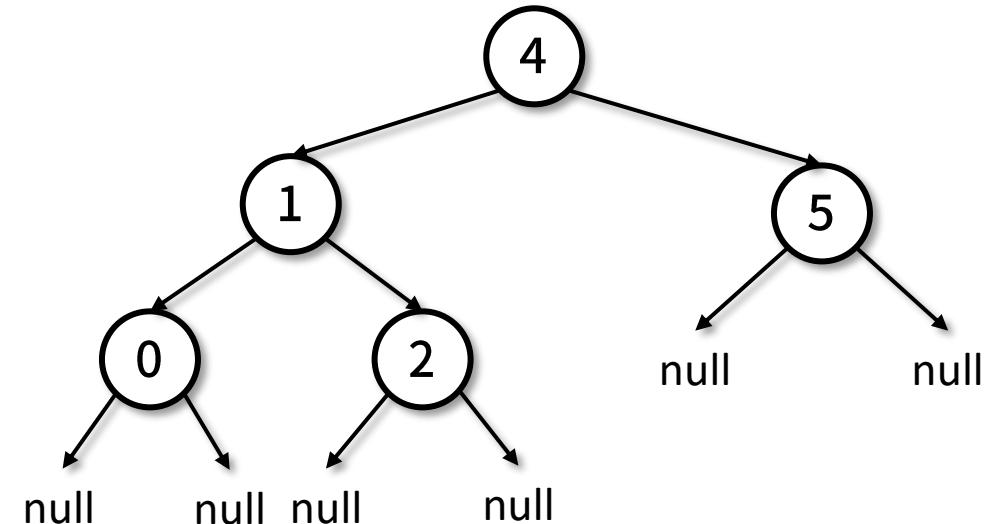
1. Use the same approach to reach we want to delete (e.g. node with key 3)



## Binary Search Tree

# Deletion (2)

1. Use the same approach to reach we want to delete (e.g. node with key 3)
2. Replace the removed node with the smallest node in the right subtree.



# Binary Search Tree (2)

## Initialization

```
class BinarySearchTree(MutableMapping):
    def __init__(self):
        self.n: int = 0
        self.root: BinarySearchTreeNode = None

    def __setitem__(self, key: ComparableT, value: Any) -> None:
        insert or update item
    def __getitem__(self, key: ComparableT) -> Any: ...
        get access to the item with given key
    def __delitem__(self, key: ComparableT) -> None: ...
        delete an item with given key
    def insert(...):
        helper function for insertion
    def delete(...):
        helper function for deletion
    def min_of_subtree(...):
        return the smallest node in the subtree
    def __iter__(self) -> BinarySearchTreeIterator: ...
    def __len__(self) -> int: ...
```

```
class BinarySearchTreeNode:
    def __init__(self, key: ComparableT, value: Any):
        self.key = key
        self.value = value
        self.left: BinarySearchTreeNode = None
        self.right: BinarySearchTreeNode = None
```

## Command

```
bst = BinarySearchTree()
```

## Visualization

$n = 0$

# Binary Search Tree Set Item (1)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
    return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

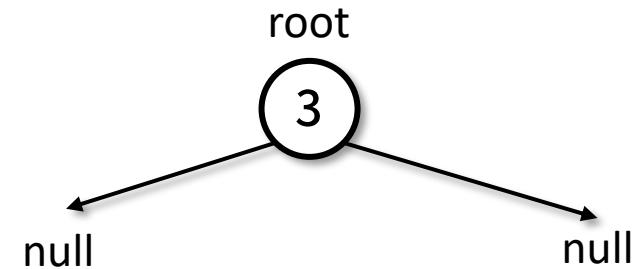
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[3] = '3'
```

## Visualization

n = 1



# Binary Search Tree Set Item (2)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
    return
    else:
        self.root = self.insert(self.root, key, value)
    return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

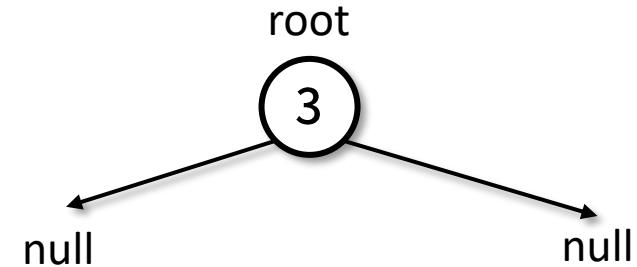
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[1] = '1'
```

## Visualization

n = 1



# Binary Search Tree Set Item (3)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return

def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

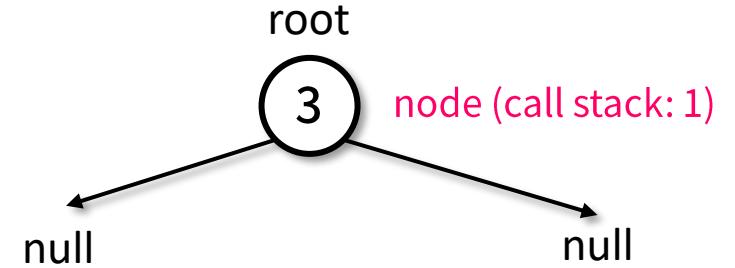
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[1] = '1'
```

## Visualization

n = 1



# Binary Search Tree

# Set Item (4)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)
```

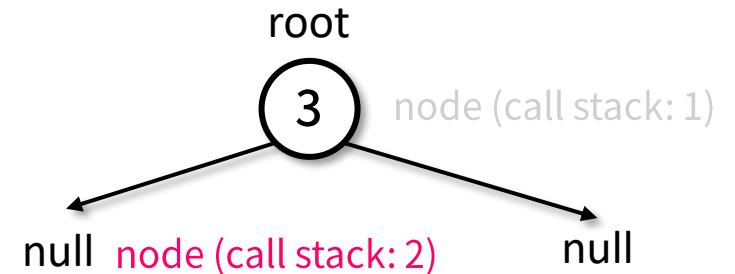
```
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[1] = '1'
```

## Visualization

n = 2



# Binary Search Tree Set Item (5)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

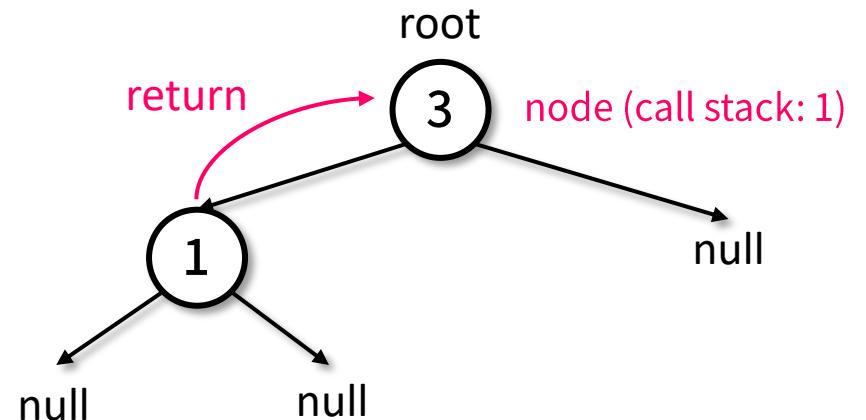
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[1] = '1'
```

## Visualization

n = 2



# Binary Search Tree Set Item (6)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

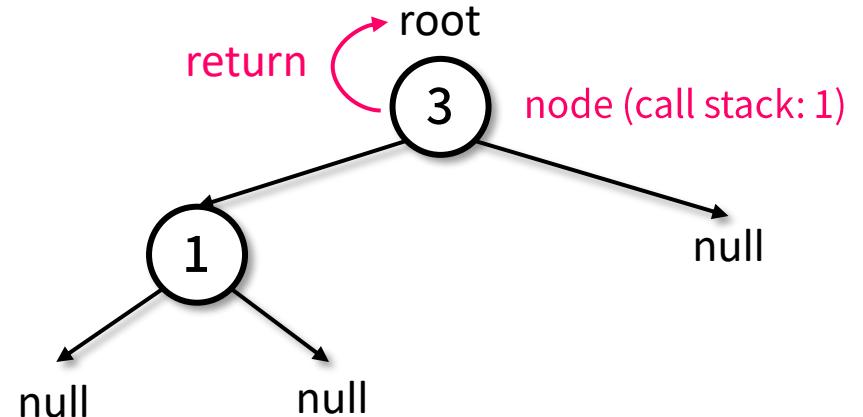
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[1] = '1'
```

## Visualization

n = 2



# Binary Search Tree Set Item (7)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
    return
    else:
        self.root = self.insert(self.root, key, value)
    return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

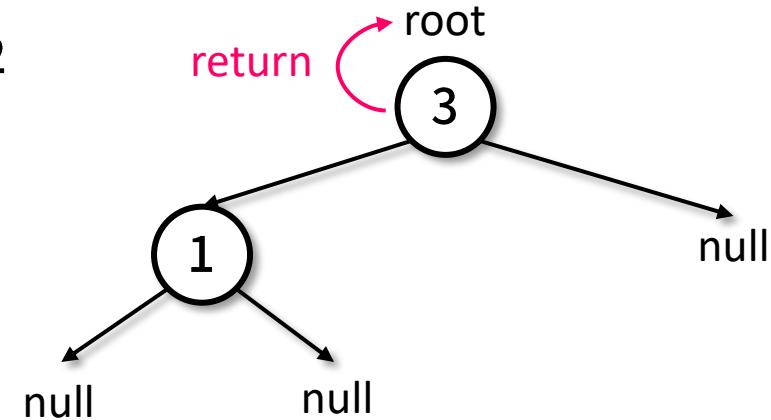
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[1] = '1'
```

## Visualization

n = 2



# Binary Search Tree Set Item (8)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
    return
    else:
        self.root = self.insert(self.root, key, value)
    return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

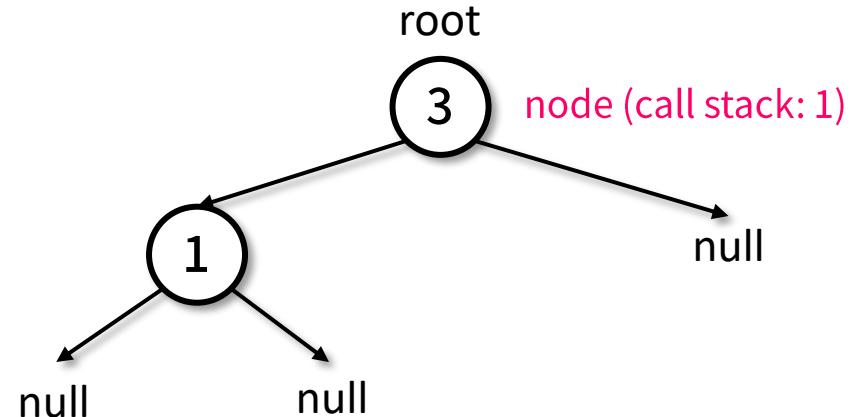
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[2] = '2'
```

## Visualization

n = 2



# Binary Search Tree Set Item (9)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

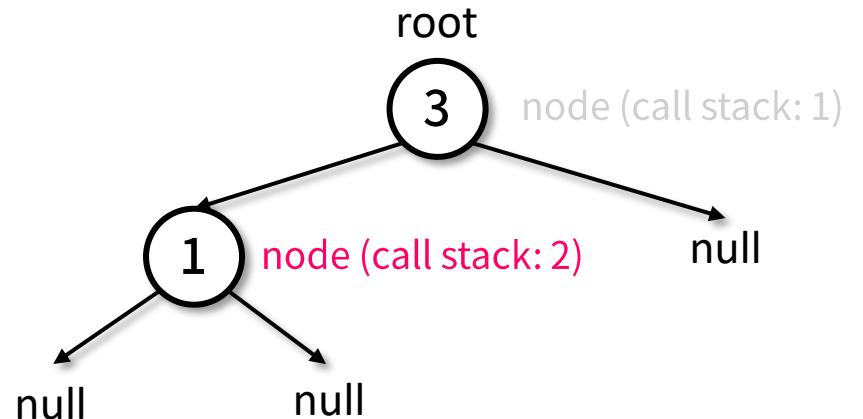
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[2] = '2'
```

## Visualization

$n = 2$



# Binary Search Tree Set Item (10)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

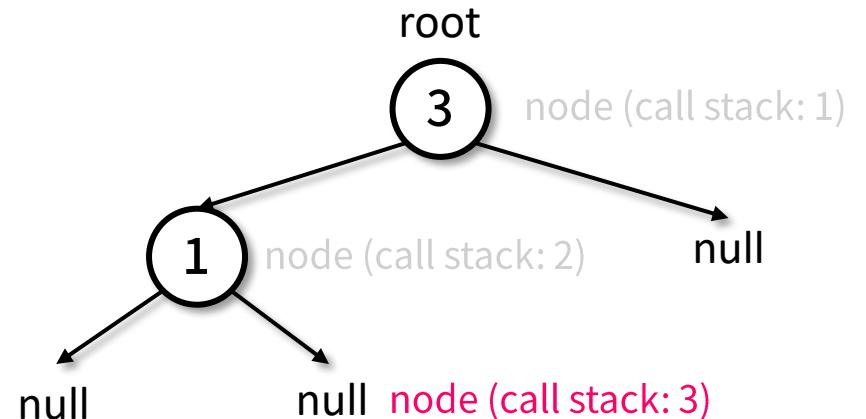
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[2] = '2'
```

## Visualization

n = 3



# Binary Search Tree Set Item (11)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

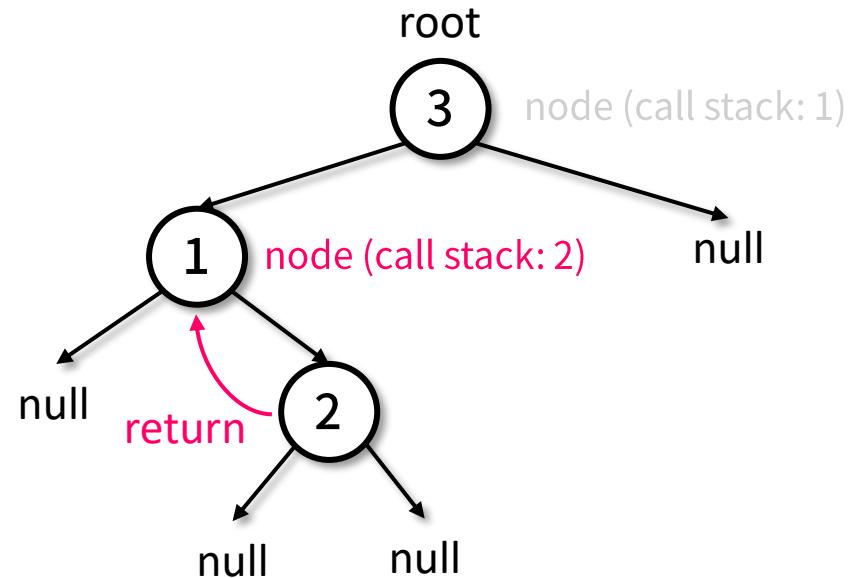
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[2] = '2'
```

## Visualization

n = 3



# Binary Search Tree Set Item (12)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

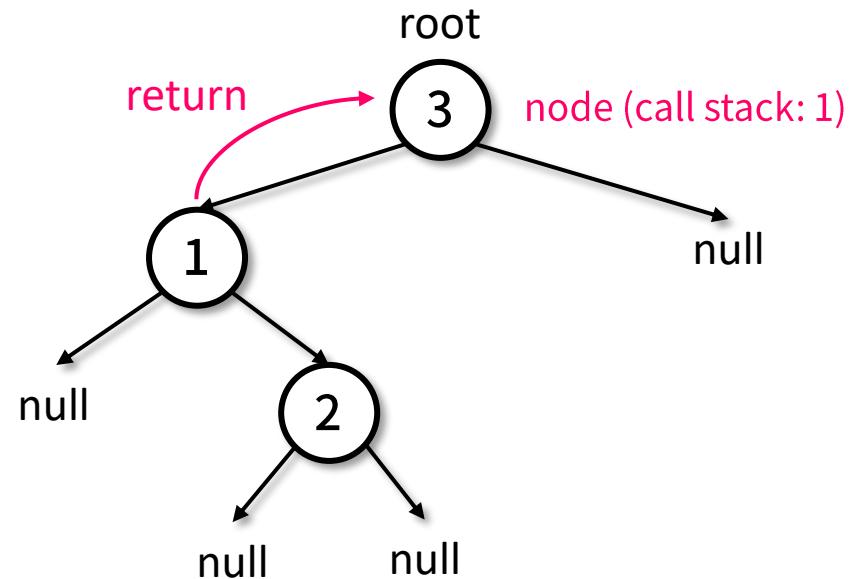
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[2] = '2'
```

## Visualization

n = 3



# Binary Search Tree Set Item (13)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
        return
    else:
        self.root = self.insert(self.root, key, value)
        return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

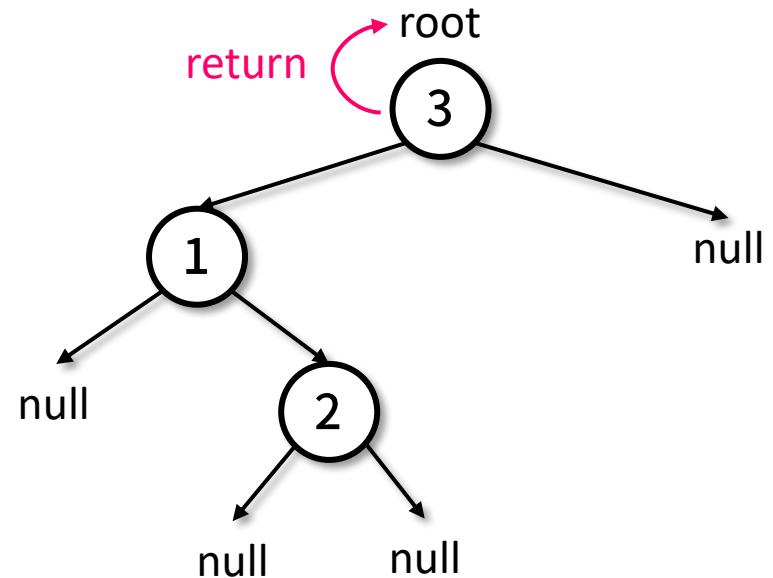
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[2] = '2'
```

## Visualization

n = 3



# Binary Search Tree

# Set Item (14)

## Set Item

```
def __setitem__(self, key: ComparableT, value: Any) -> None:
    if self.n == 0:
        self.root = BinarySearchTreeNode(key, value)
        self.n += 1
    return
    else:
        self.root = self.insert(self.root, key, value)
    return
def insert(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT,
    value: Any) -> BinarySearchTreeNode:
    if node is None:
        self.n += 1
        return BinarySearchTreeNode(key, value)

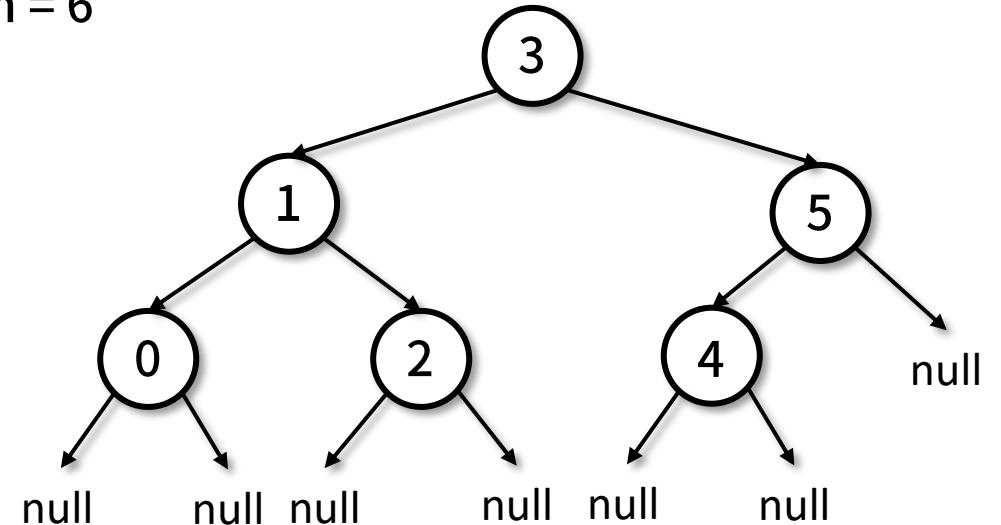
    if key < node.key:
        node.left = self.insert(node.left, key, value)
    elif key > node.key:
        node.right = self.insert(node.right, key, value)
    else:
        node.value = value
    return node
```

## Command

```
bst[0] = '0'
bst[5] = '5'
bst[4] = '4'
```

## Visualization

$n = 6$



Binary Search Tree

# Delete Item (1)

## Delete Item

```
def _delitem_(self, key: ComparableT) -> None:  
    self.delete(self.root, key)  
    self.n -= 1  
    return
```

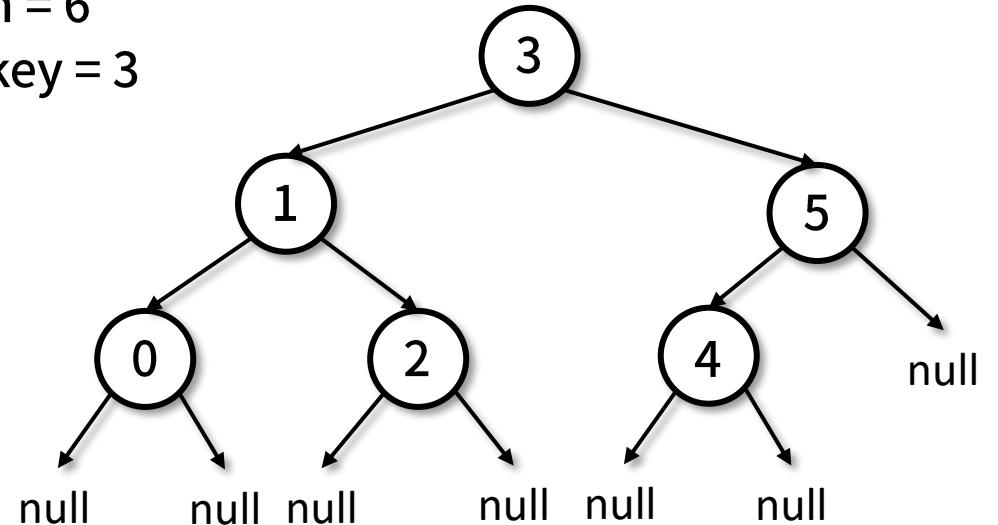
## Command

```
del bst[3]
```

## Visualization

$n = 6$

key = 3



# Binary Search Tree Delete Item (2)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

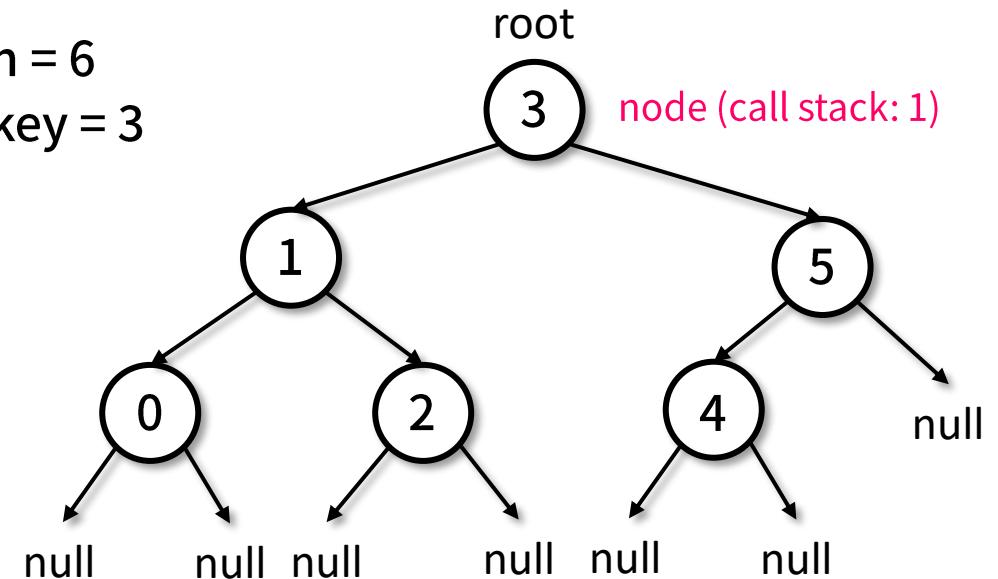
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3$



# Binary Search Tree Delete Item (3)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

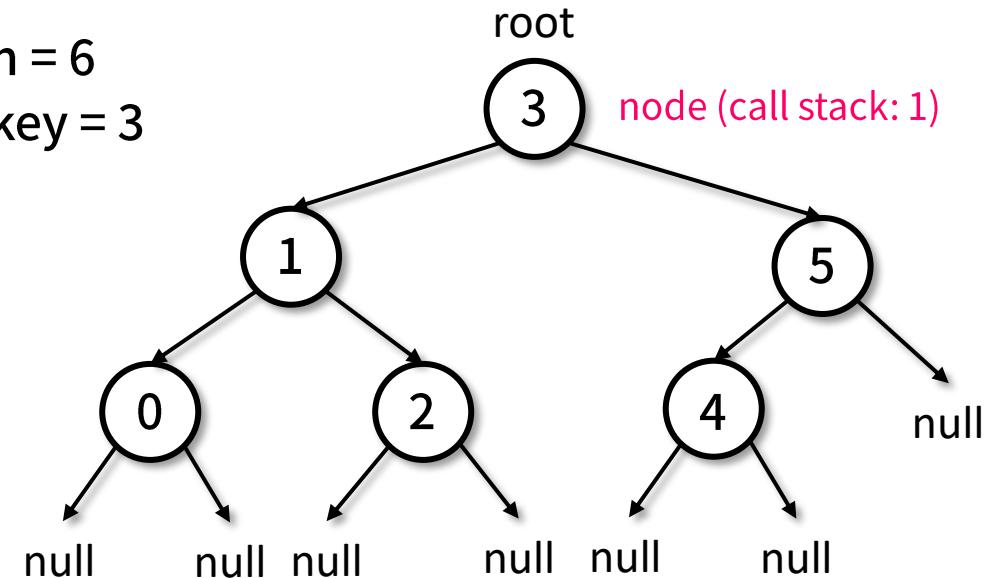
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3$



# Binary Search Tree Delete Item (4)

## Delete Item

```
def min_of_subtree(  
    self,  
    node: BinarySearchTreeNode) -> BinarySearchTreeNode:  
    while node.left is not None:  
        node = node.left  
    return node
```

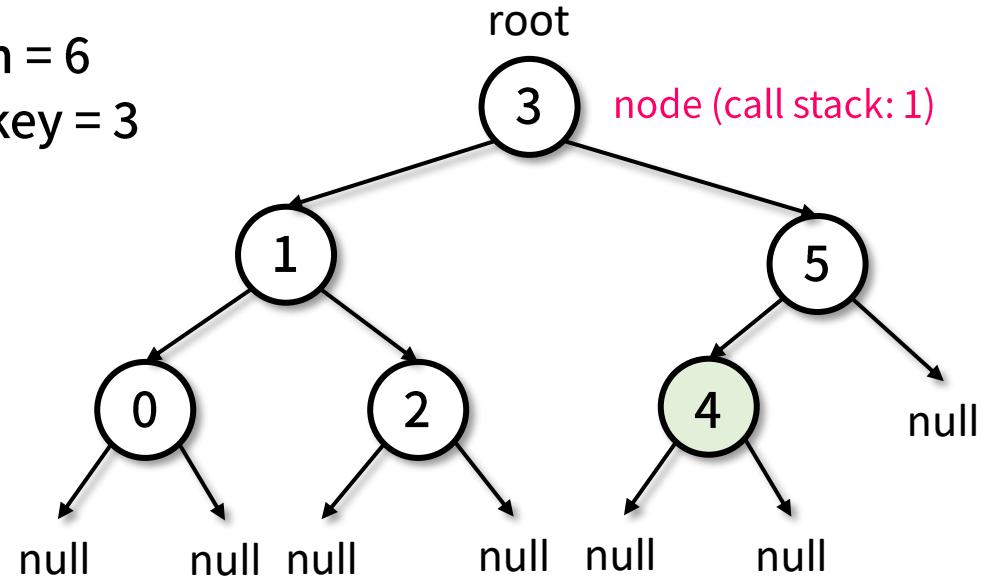
## Command

```
del bst[3]
```

## Visualization

$n = 6$

key = 3



# Binary Search Tree Delete Item (5)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

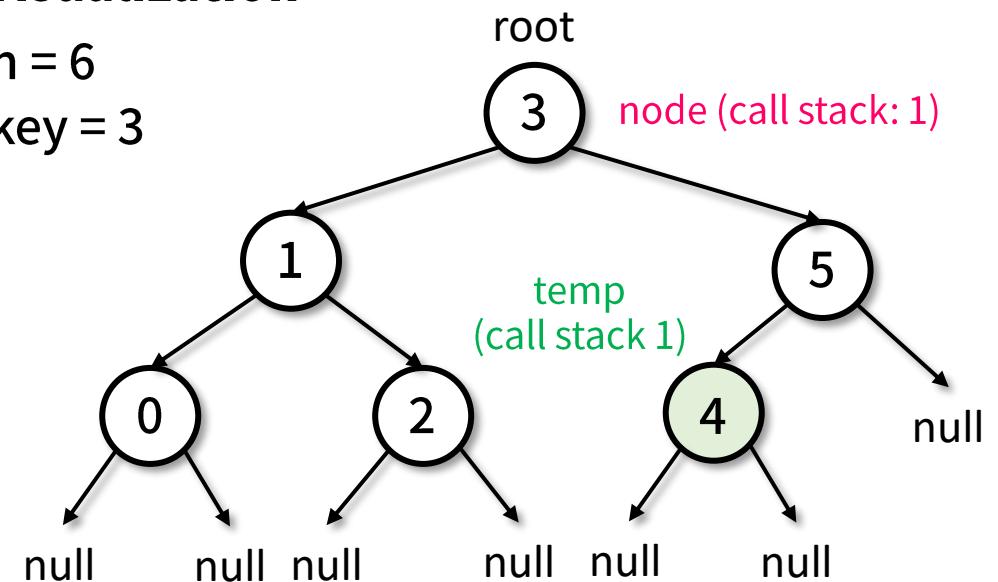
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3$



# Binary Search Tree Delete Item (6)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

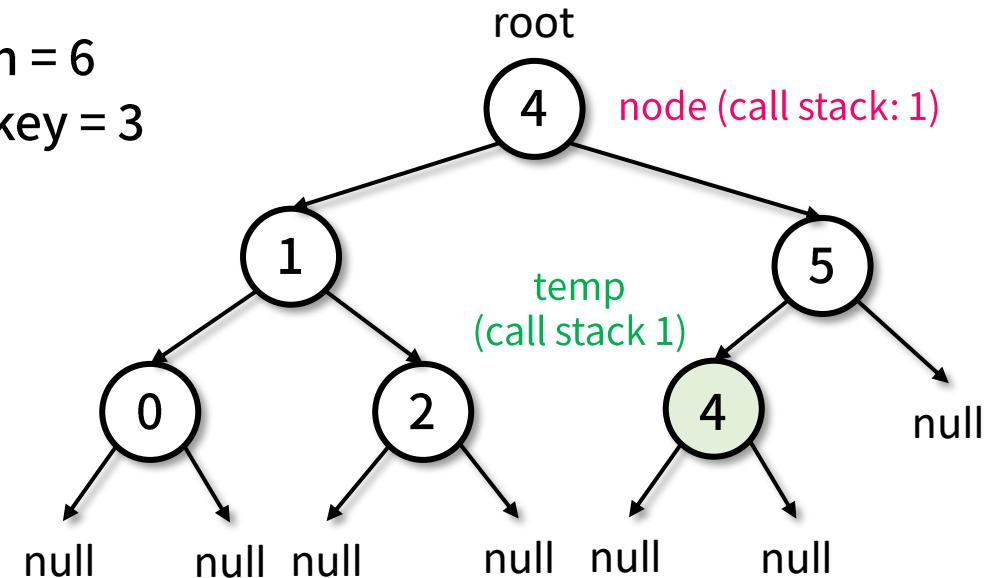
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3$



# Binary Search Tree Delete Item (7)

# Delete Item

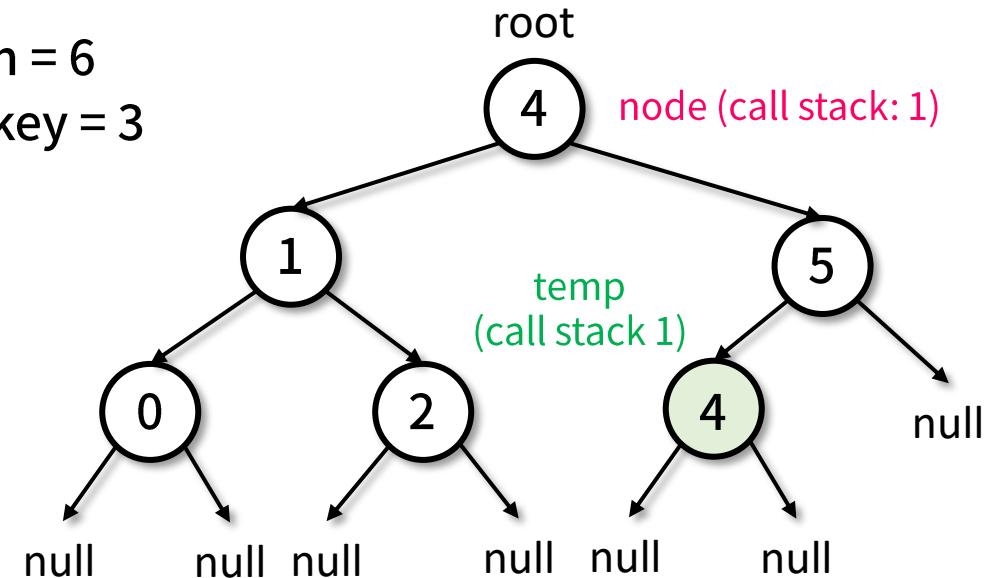
```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')
    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, tem
    return node
```

# Command

```
del bst[3]
```

# Visualization

$$n = 6$$
  
$$\text{key} = 3$$



# Binary Search Tree Delete Item (8)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

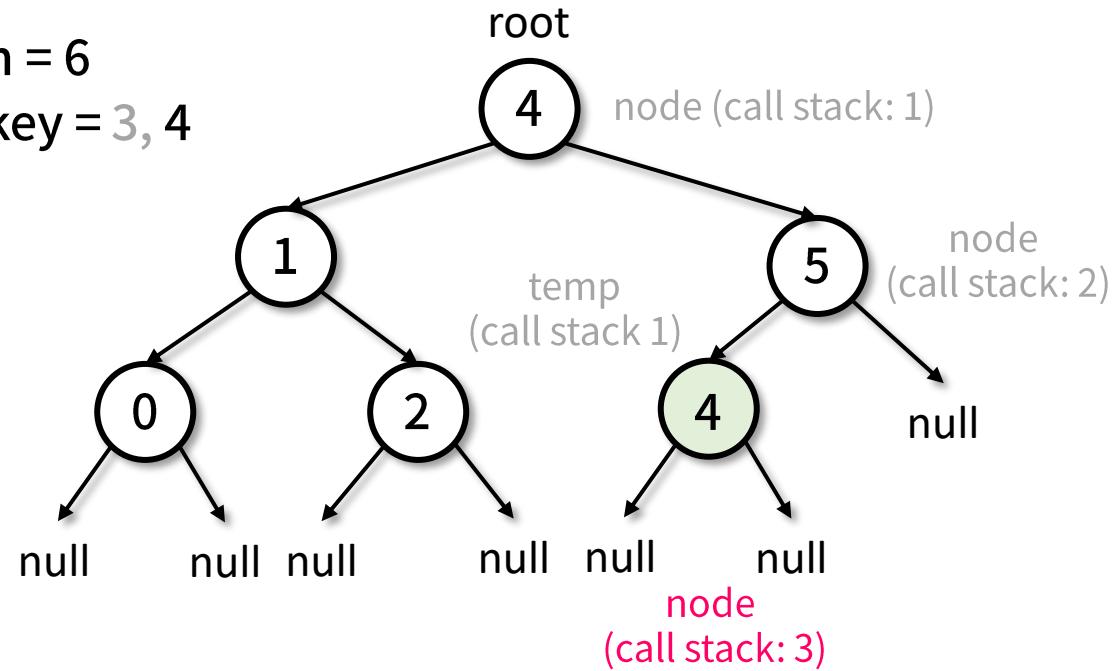
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3, 4$



# Binary Search Tree Delete Item (9)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

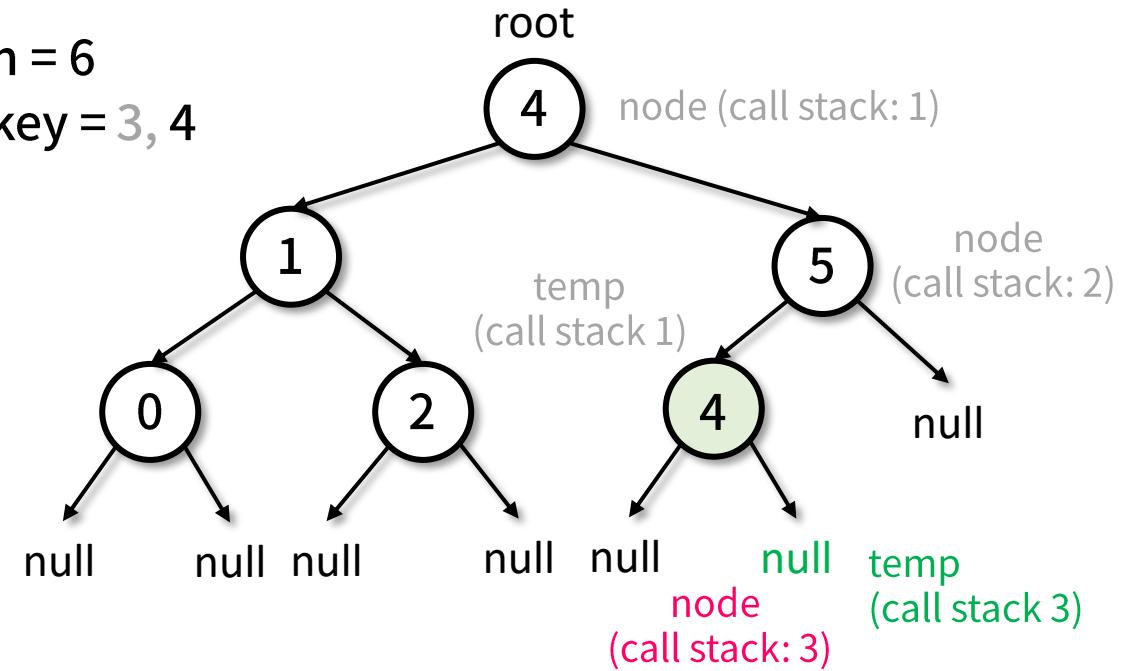
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3, 4$



# Binary Search Tree Delete Item (10)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

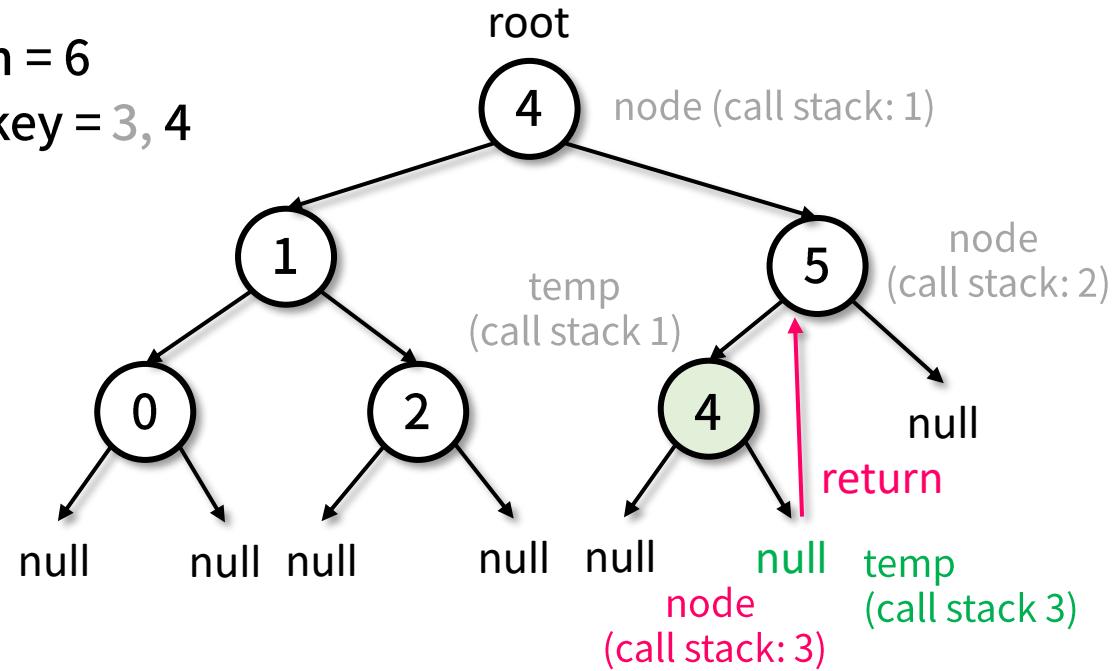
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3, 4$



# Binary Search Tree Delete Item (11)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

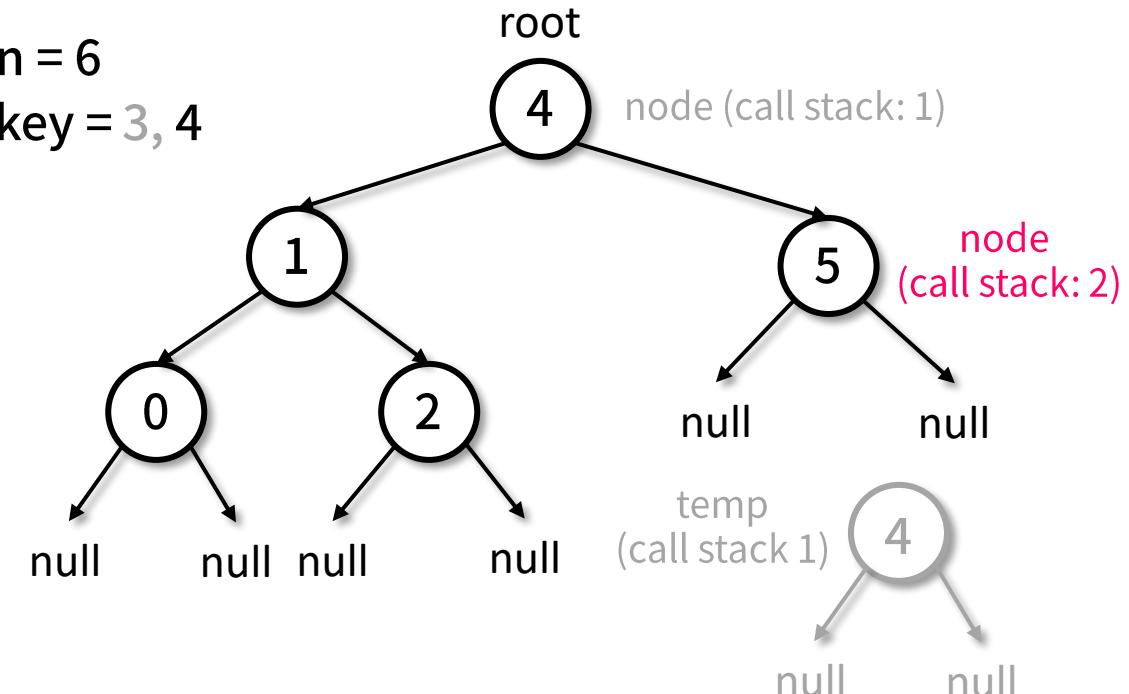
    return node
```

## Command

```
del bst[3]
```

## Visualization

$n = 6$   
 $\text{key} = 3, 4$



# Binary Search Tree Delete Item (12)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

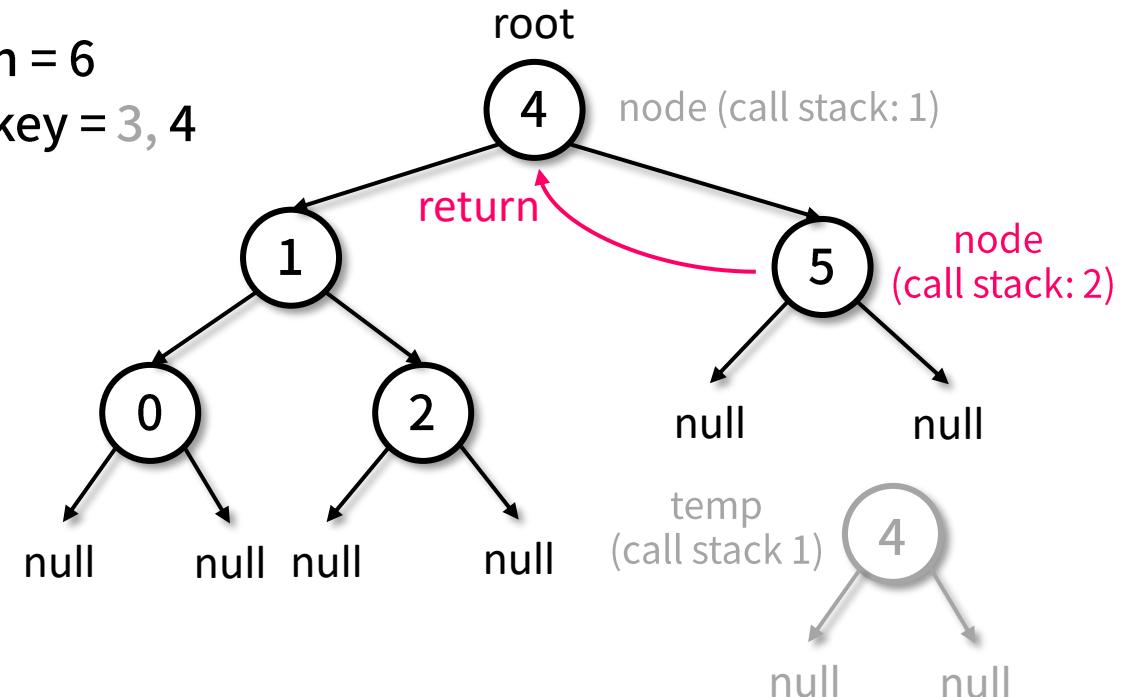
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3, 4$



# Binary Search Tree Delete Item (13)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

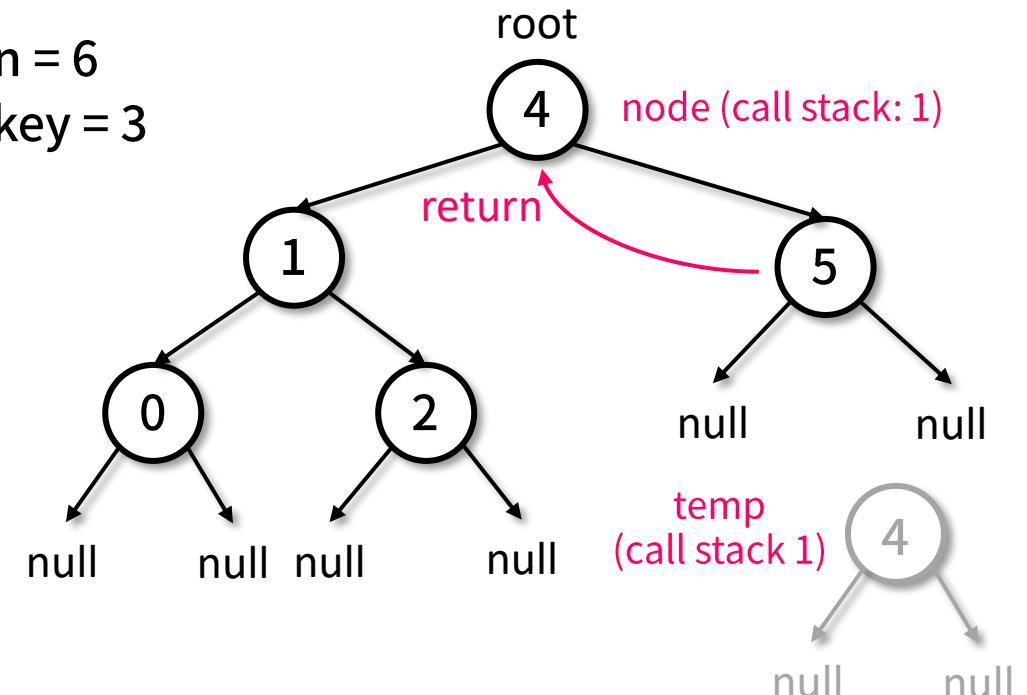
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$\text{key} = 3$



# Binary Search Tree Delete Item (14)

## Delete Item

```
def delete(
    self,
    node: BinarySearchTreeNode,
    key: ComparableT) -> BinarySearchTreeNode:
    if node is None:
        raise KeyError(f'{key}')

    if key < node.key:
        node.left = self.delete(node.left, key)
    elif key > node.key:
        node.right = self.delete(node.right, key)
    else:
        if node.left is None:
            temp = node.right
            node = None
            return temp
        elif node.right is None:
            node = node.left
            node = None
            return temp
        else:
            temp = self.min_of_subtree(node.right)
            node.key = temp.key
            node.value = temp.value
            node.right = self.delete(node.right, temp.key)

    return node
```

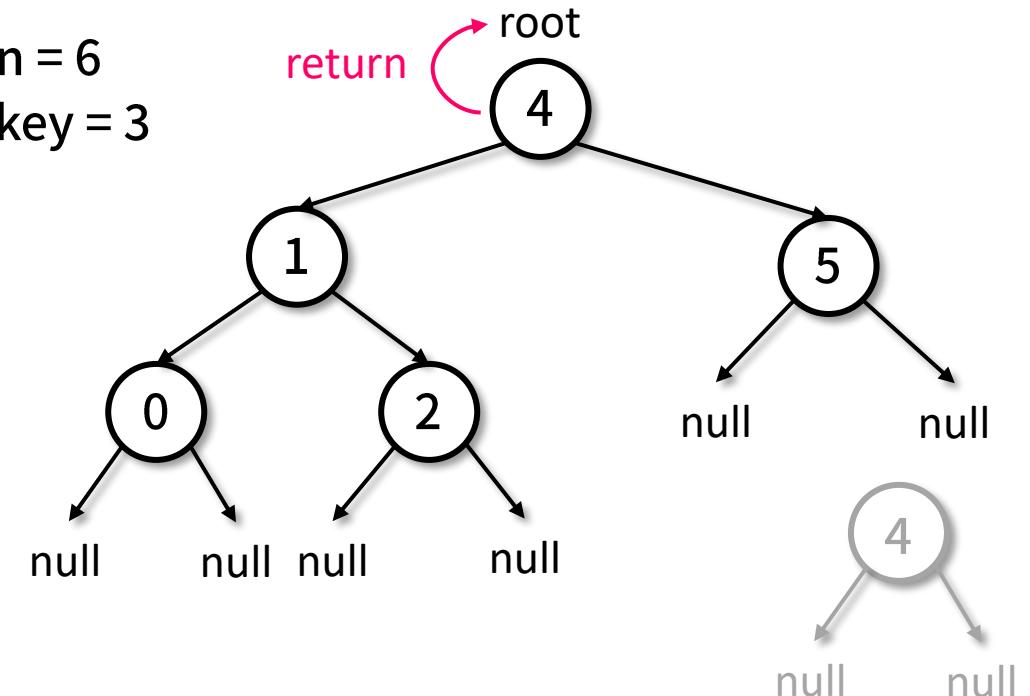
## Command

```
del bst[3]
```

## Visualization

$n = 6$

$key = 3$



## Binary Search Tree

# Delete Item (15)

### Delete Item

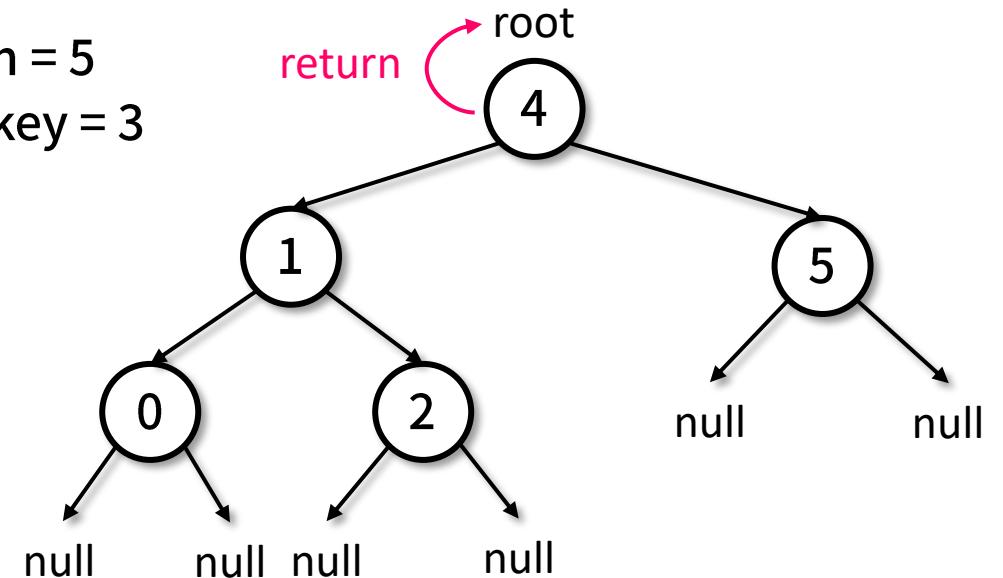
```
def __delitem__(self, key: ComparableT) -> None:  
    self.delete(self.root, key)  
    self.n -= 1  
    return
```

### Command

```
del bst[3]
```

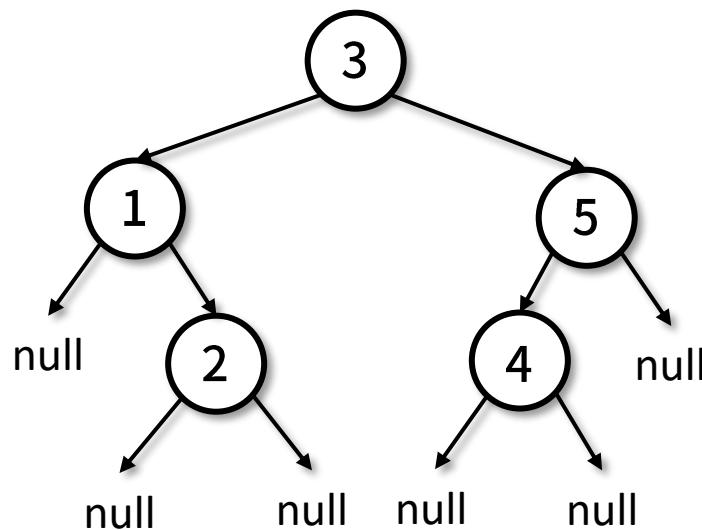
### Visualization

n = 5  
key = 3



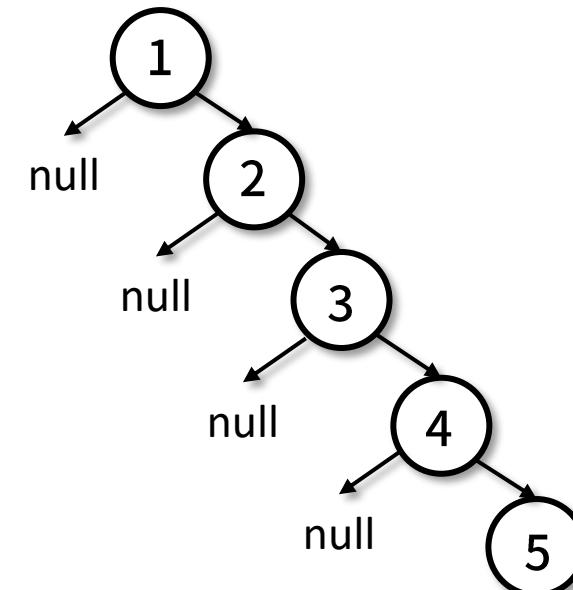
# Problem of Binary Search Tree

- The tree structure depends on the order of insertion.



insertion: 3, 5, 4, 1, 2

average search time:  $2.2T$



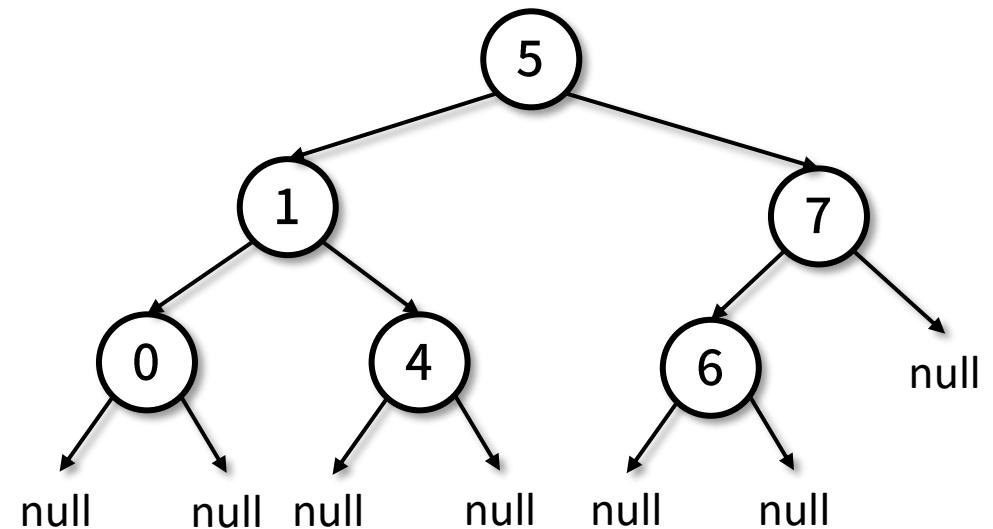
insertion: 1, 2, 3, 4, 5

average search time:  $3T$

# Adelson-Velsky Landis Tree (AVL Tree)

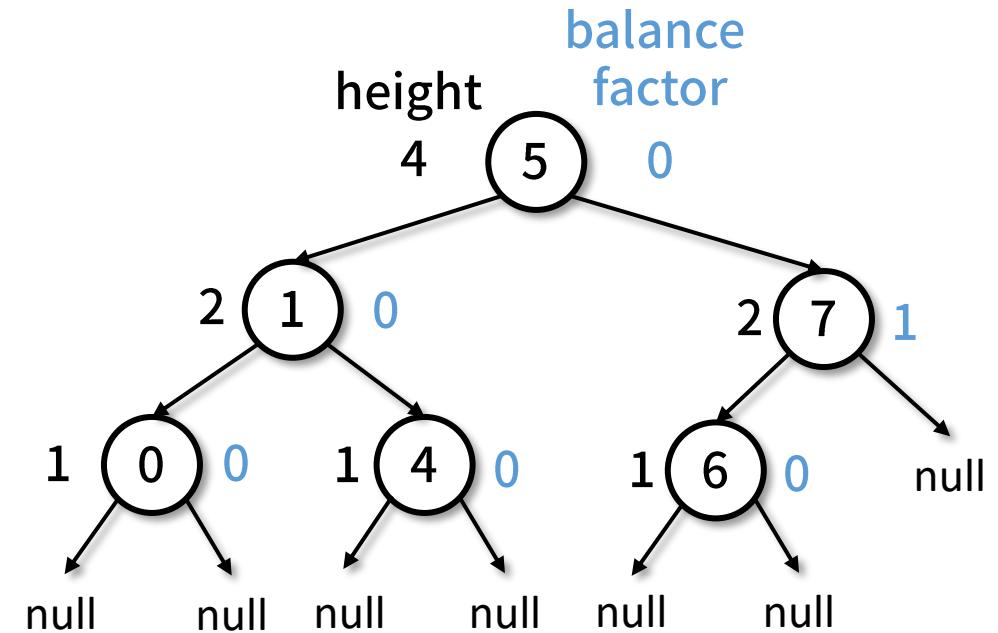
# AVL Tree (1)

- The AVL tree is a **balanced** binary search tree.
- Any subtree in the AVL tree is also an AVL tree.
- Need additional operations for *insert* and *delete* to maintain the structure to ensure the rules are invariant.



# AVL Tree (2)

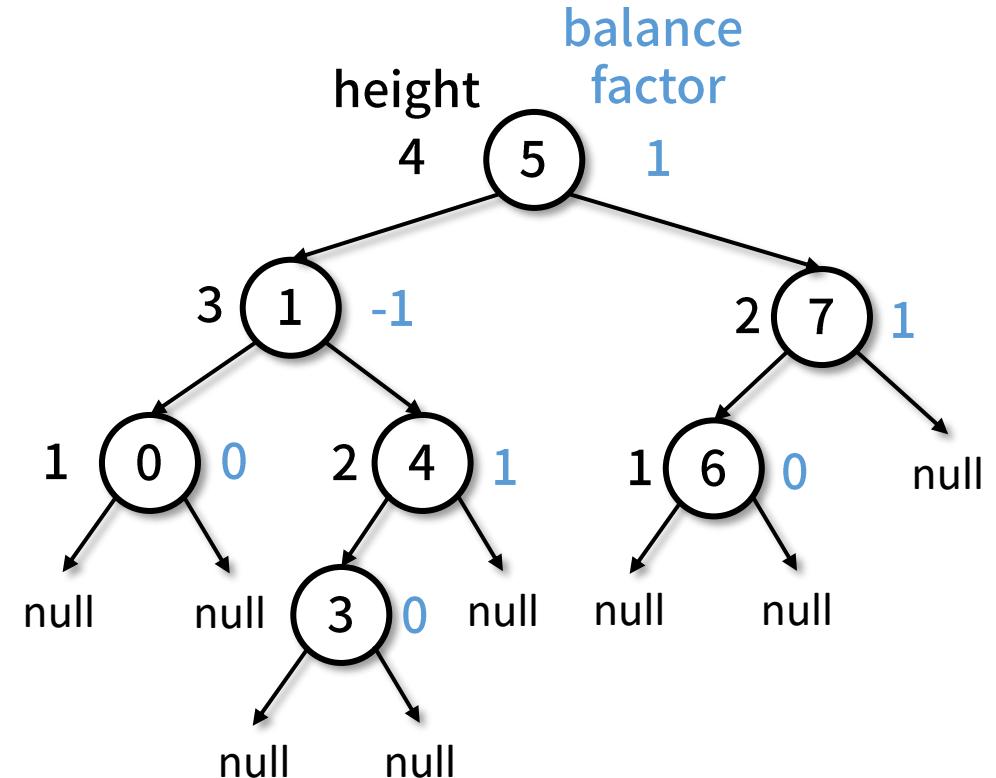
- Define the balance factor as:
  - $height(left\ subtree) - height(right\ subtree)$
- **Balance Tree**: the balance factor of every node inside is either -1, 0 or 1.



AVL Tree

# Insertion (1)

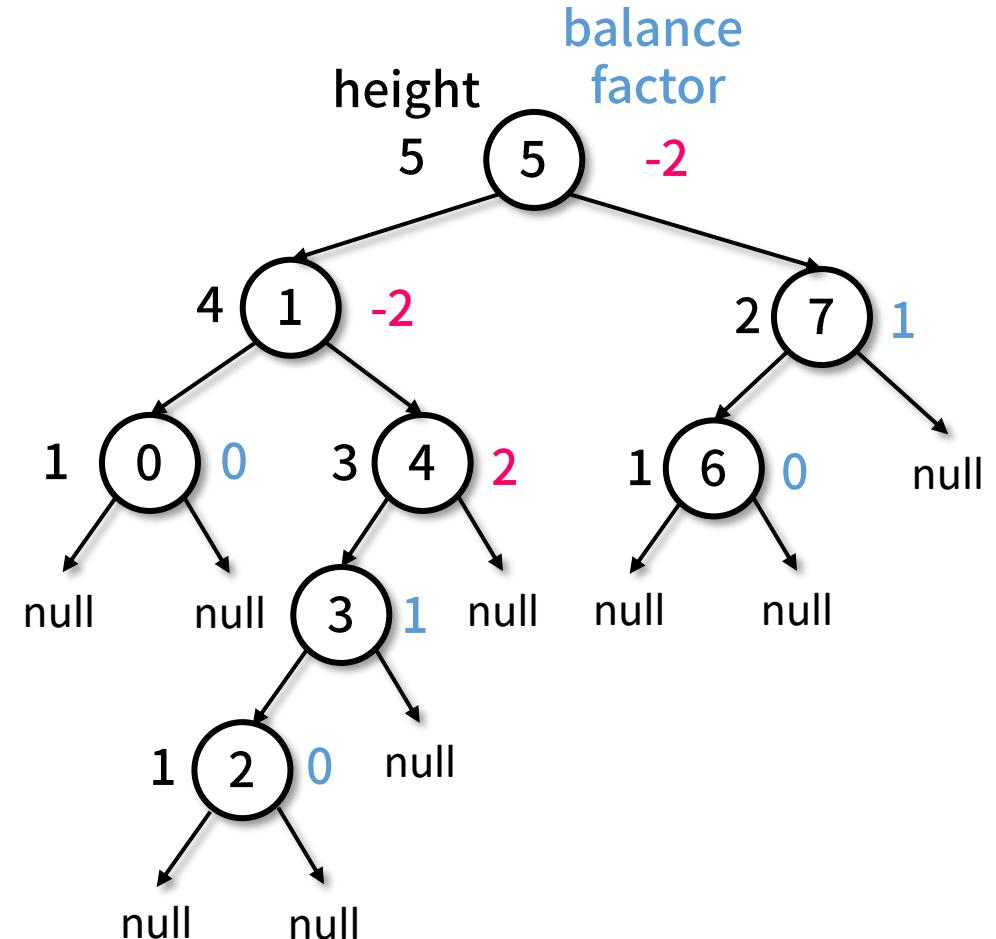
- Define the balance factor as:
  - $height(left\ subtree) - height(right\ subtree)$
- Balance Tree: the balance factor of every node inside is either -1, 0 or 1.
- Let us try insert a node with key 3.



AVL Tree

# Insertion (2)

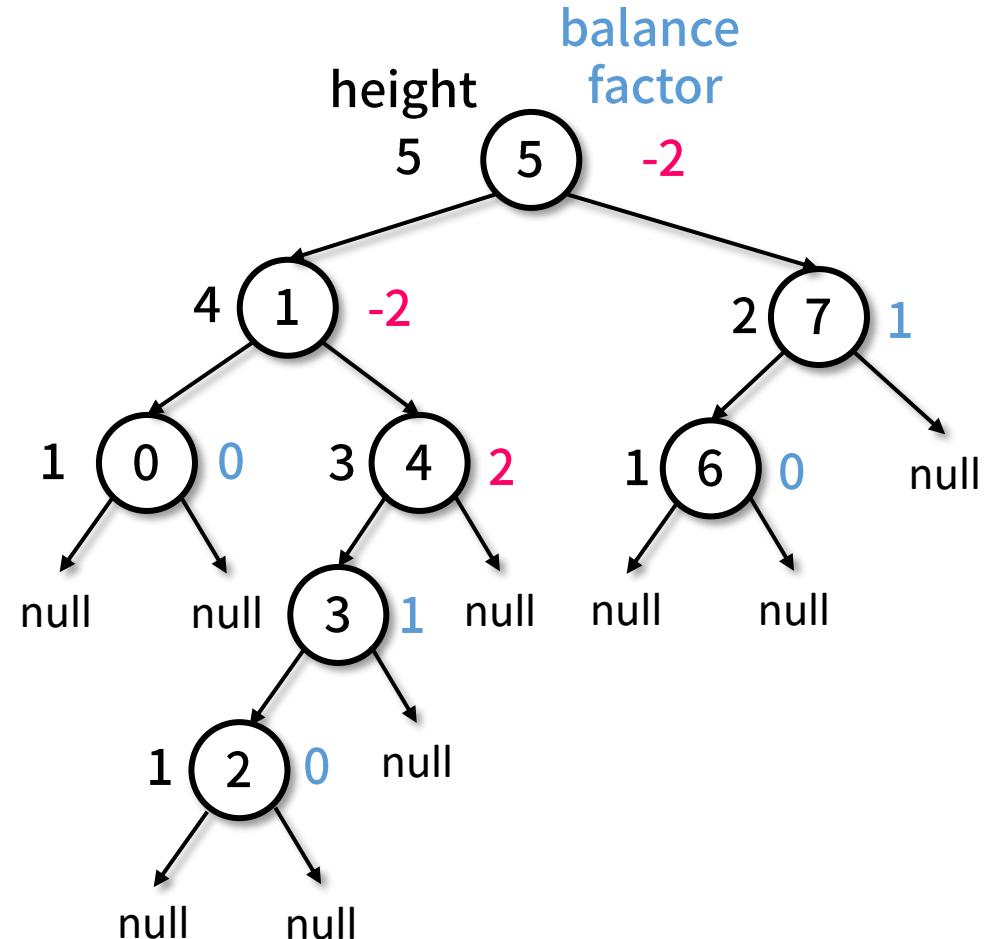
- Define the balance factor as:
  - $height(left\ subtree) - height(right\ subtree)$
- Balance Tree: the balance factor of every node inside is either -1, 0 or 1.
- Let us try insert a node with key 3.
- Let us try insert a node with key 2. The balance is broken.



AVL Tree

# Insertion (3)

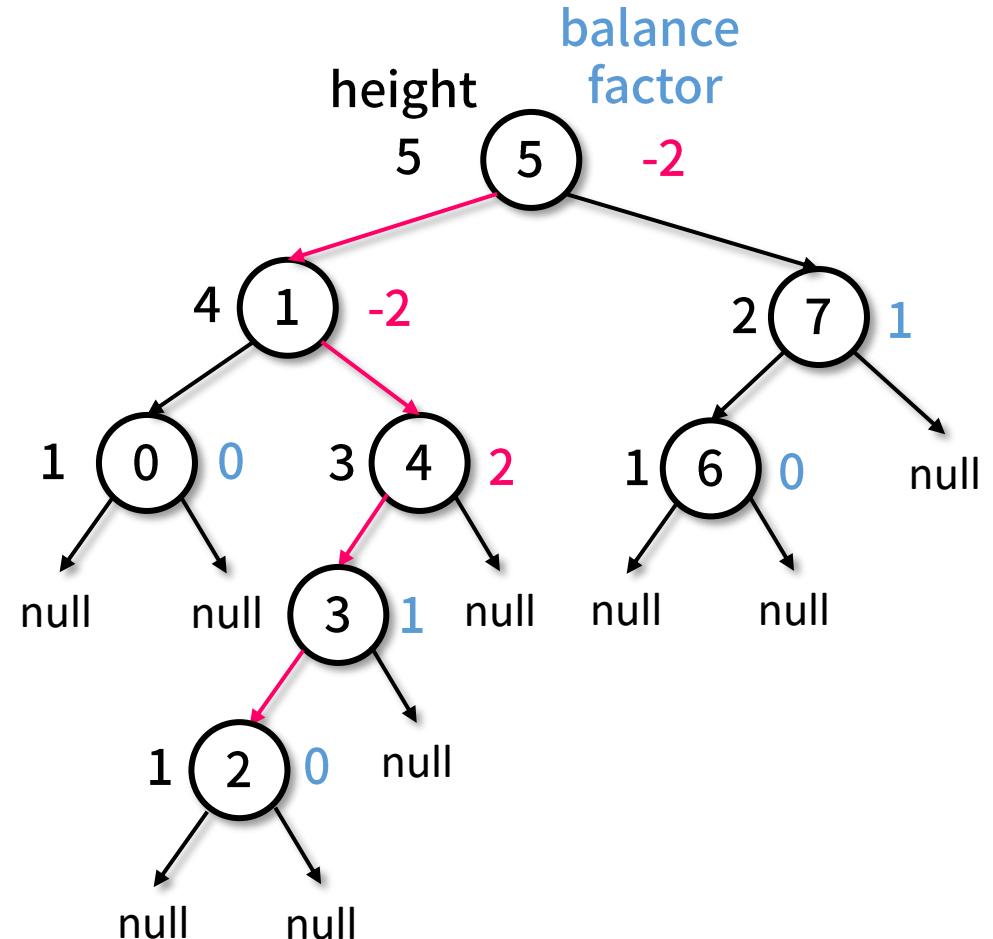
- Define the balance factor as:
  - $height(left\ subtree) - height(right\ subtree)$
- Balance Tree: the balance factor of every node inside is either -1, 0 or 1.
- In what situation the balance will be broken?



AVL Tree

# Insertion (4)

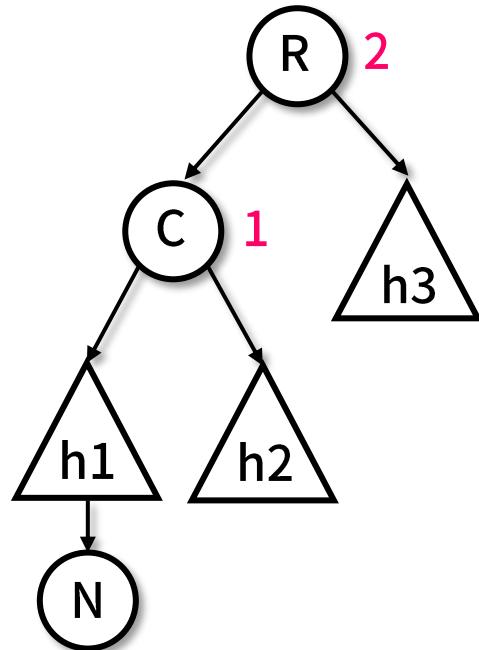
- Define the balance factor as:
  - $height(left\ subtree) - height(right\ subtree)$
- Balance Tree: the balance factor of every node inside is either -1, 0 or 1.
- In what situation the balance will be broken?
- What can we observe from the example?



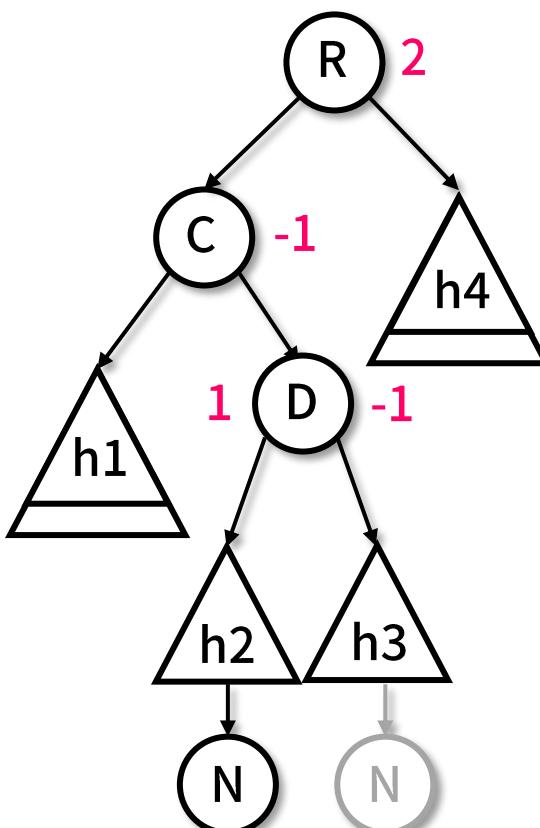
AVL Tree

# Insertion (5)

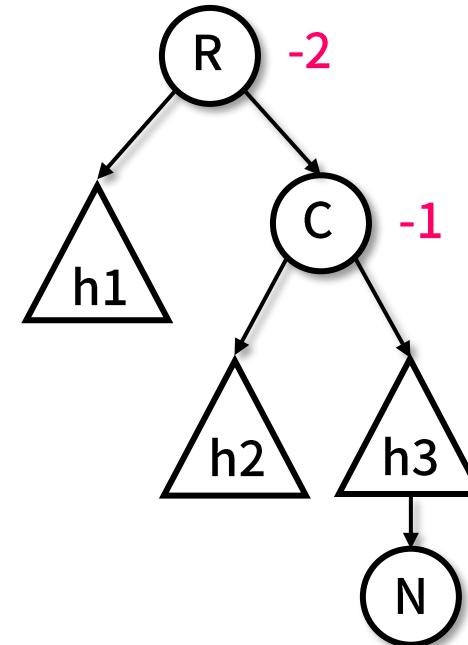
**Left Left**



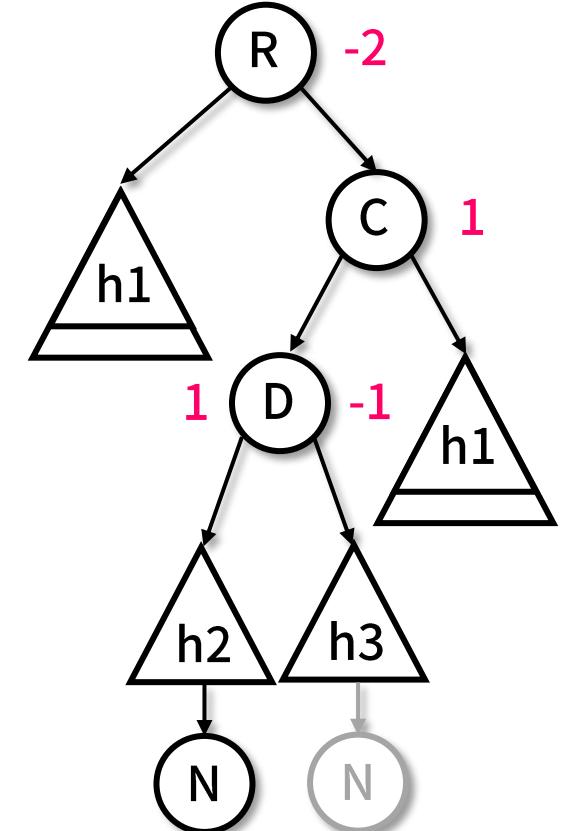
**Left Right**



**Right Right**



**Right Left**



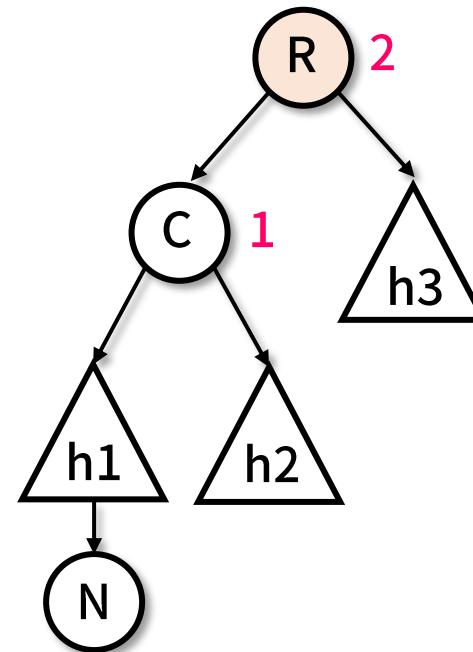
AVL Tree

# Right Rotation

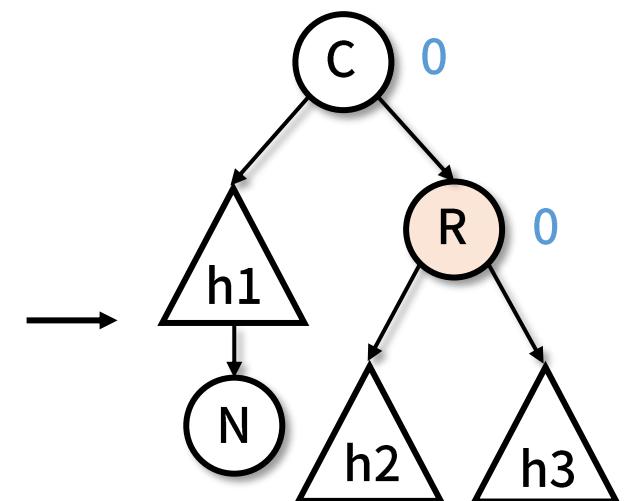
## Right Rotation

```
def right_rotate(self, node: AVLTreeNode) -> AVLTreeNode:  
  
    x = node.left  
    y = x.right  
  
    x.right = node  
    node.left = y  
  
    node.height = 1 + max(  
        self.get_node_height(node.left),  
        self.get_node_height(node.right))  
    x.height = 1 + max(  
        self.get_node_height(x.left),  
        self.get_node_height(x.right))  
  
    return x
```

## Left Left



## Balance



the height remains the same after correction

AVL Tree

# Left Rotation

## Left Rotation

```
def left_rotate(self, node: AVLTreeNode) -> AVLTreeNode:

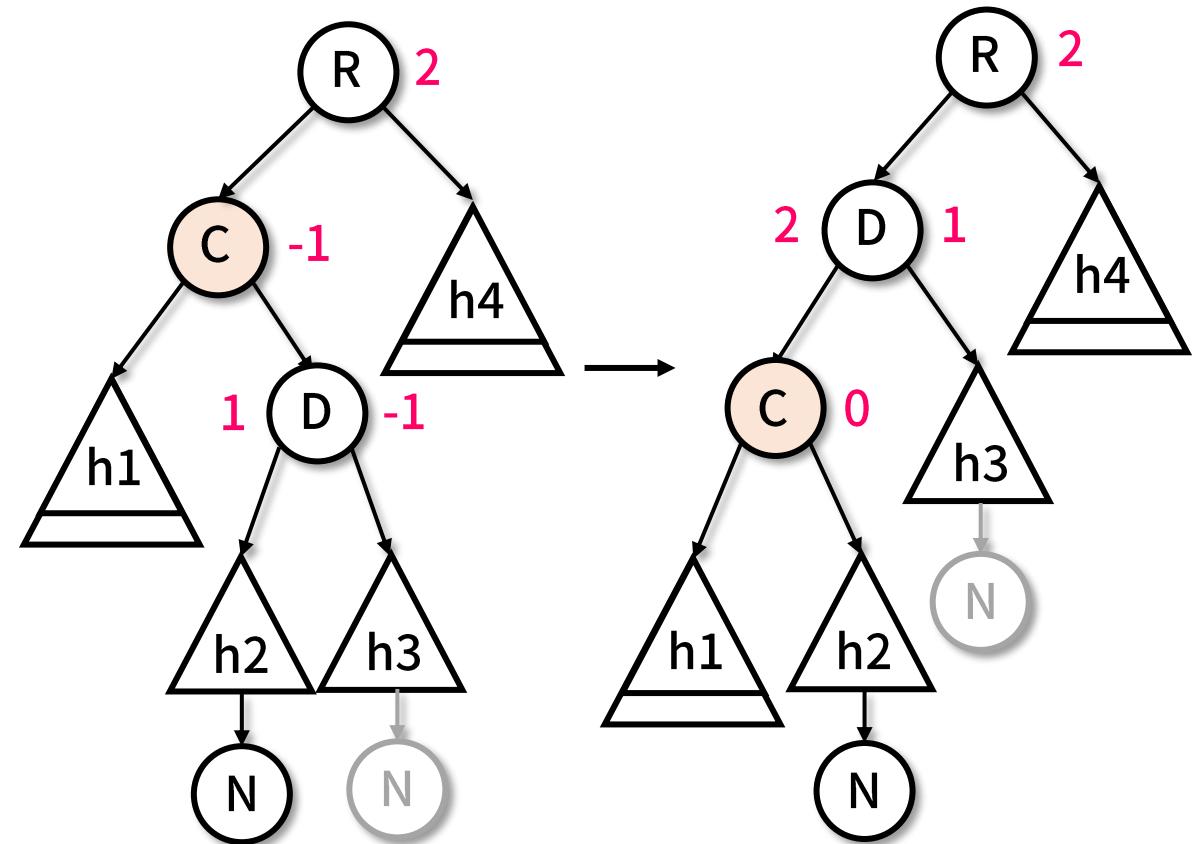
    x = node.right
    y = x.left

    x.left = node
    node.right = y

    node.height = 1 + max(
        self.get_node_height(node.left),
        self.get_node_height(node.right))
    x.height = 1 + max(
        self.get_node_height(x.left),
        self.get_node_height(x.right))

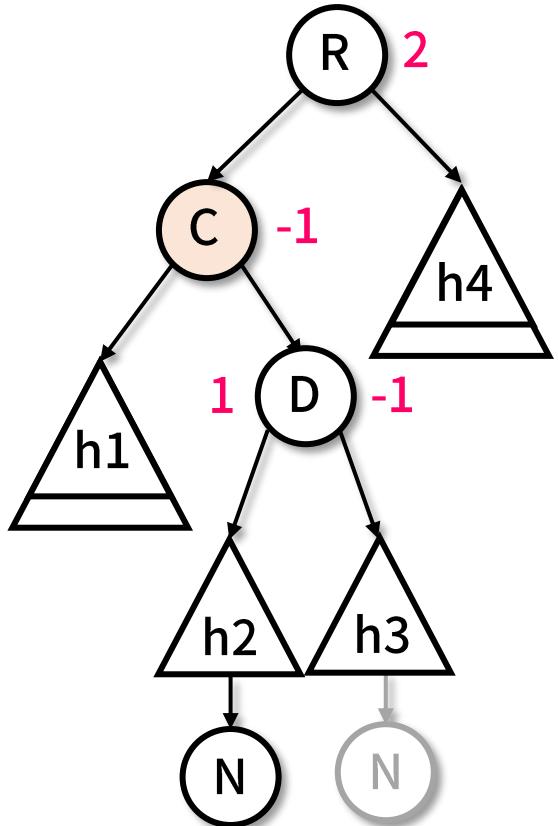
    return x
```

## Left Right

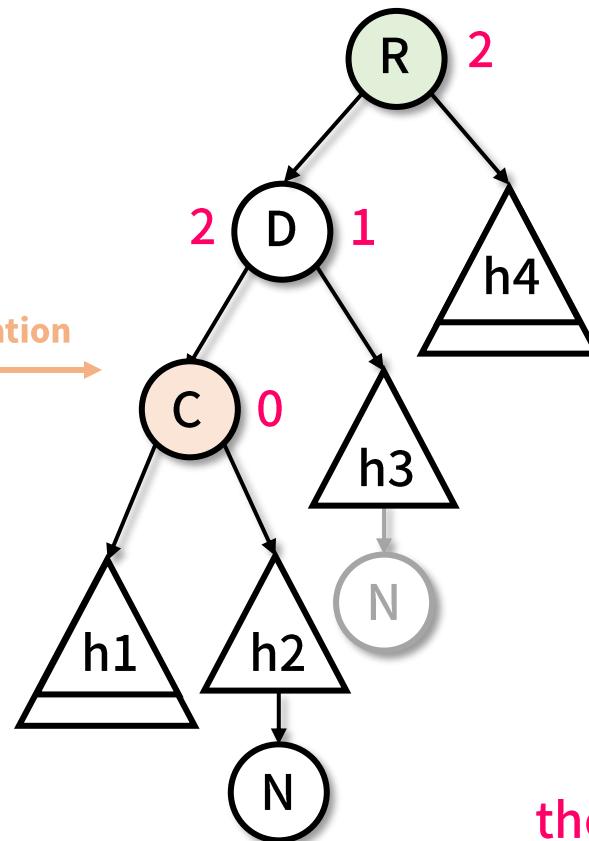


# AVL Tree Rotation

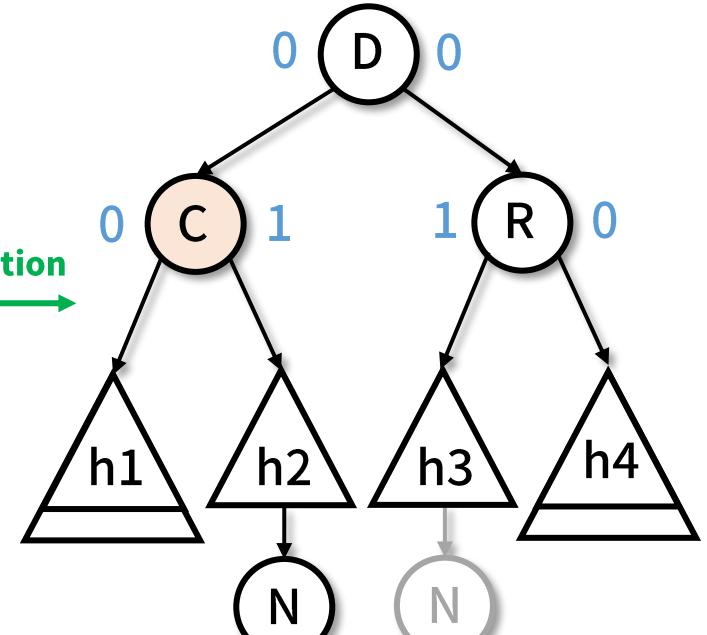
**Left Right**



Left Rotation



**Balance**



the height remains the same after correction

AVL Tree

# Insertion Correction

## Insertion Correction

```
if balance_factor > 1 and key < node.left.key:  
    return self.right_rotate(node)  
  
if balance_factor < -1 and key > node.right.key:  
    return self.left_rotate(node)  
  
if balance_factor > 1 and key > node.left.key:  
    node.left = self.left_rotate(node.left)  
    return self.right_rotate(node)  
  
if balance_factor < -1 and key < node.right.key:  
    node.right = self.right_rotate(node.right)  
    return self.left_rotate(node)
```

**Left Left**



**Right Right**



**Left Right**



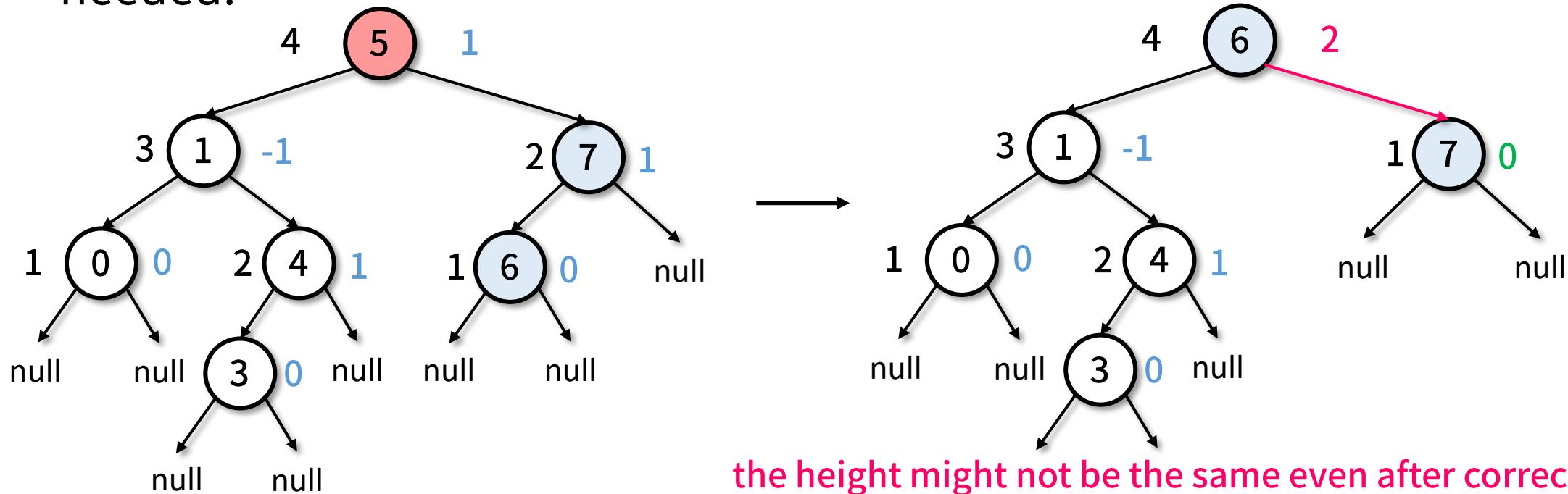
**Right Left**



AVL Tree

# Deletion (1)

- Perform the ordinary binary search tree deletion.
- Check all the nodes iterated in the process, correct the structure if needed.

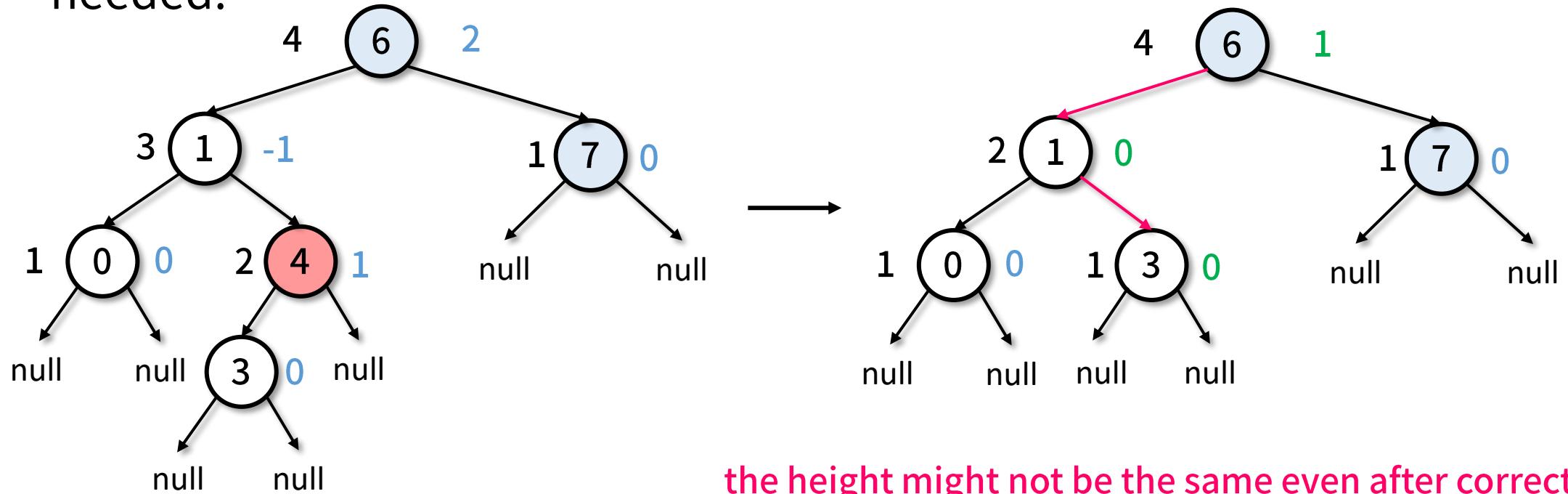


the height might not be the same even after correction,  
need to check through all node in the iteration.

AVL Tree

# Deletion (2)

- Perform the ordinary binary search tree deletion.
- Check all the nodes iterated in the process, correct the structure if needed.



AVL Tree

# Deletion Correction

## Deletion Correction

```
if ((balance_factor > 1) and  
    (self.get_balance_factor(node.left) >= 0)):  
    return self.right_rotate(node)  
  
if ((balance_factor < -1) and  
    (self.get_balance_factor(node.right) <= 0)):  
    return self.left_rotate(node)  
  
if ((balance_factor > 1) and  
    (self.get_balance_factor(node.left) < 0)):  
    node.left = self.left_rotate(node.left)  
    return self.right_rotate(node)  
  
if ((balance_factor < -1) and  
    (self.get_balance_factor(node.right) > 0)):  
    node.right = self.right_rotate(node.right)  
    return self.left_rotate(node)
```

**Left Left**



**Right Right**



**Left Right**



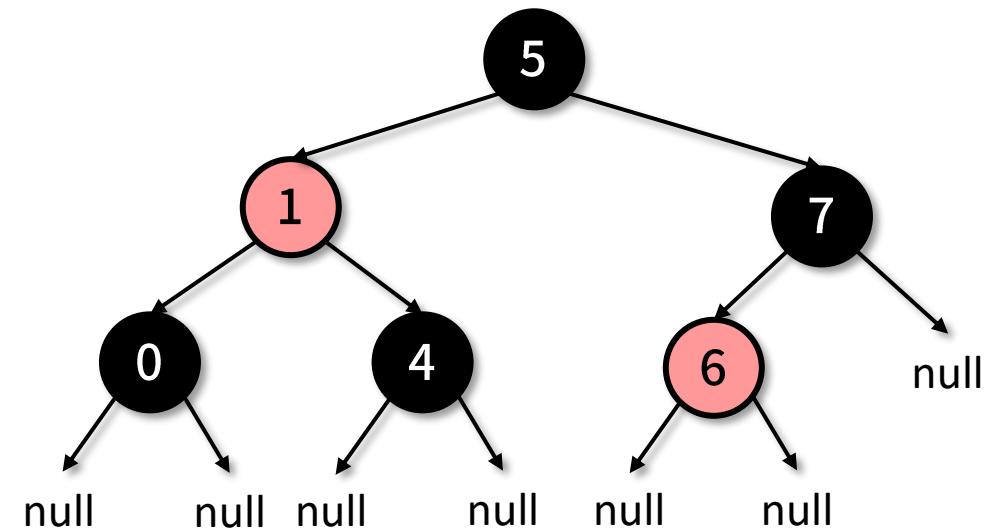
**Right Left**



# Red-Black Tree

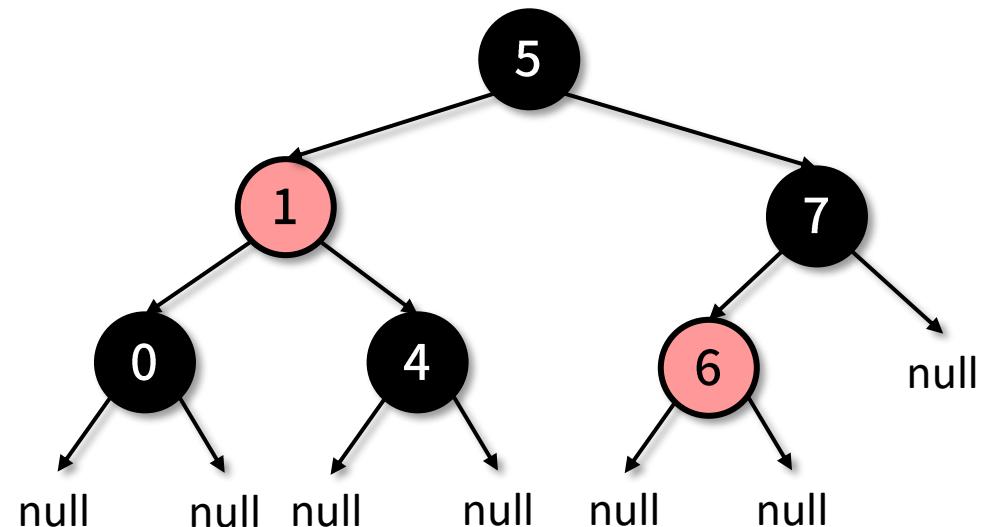
# Red-Black Tree (1)

- Red-black tree rules:
  - Node can be either black or red.
  - Root node is always black.
  - Null node is considered as black.
  - The number of black node from the root to the leaf is the same.
  - Red node can only have black children.
  - Black node can have black or red children.
- Need additional operations for *insert* and *delete* to maintain the structure to ensure the rules are invariant.



# Red-Black Tree (2)

- Red-black tree properties:
  - Any subtree with a black root is also a red-black tree.
  - The maximum height difference between the two subtree is
$$2 \times \min(\text{height}_{\text{left}}, \text{height}_{\text{right}})$$
  - Not strictly balanced binary search tree.

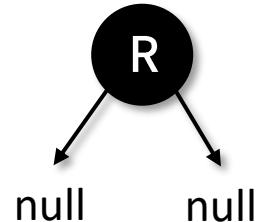


Red-Black Tree

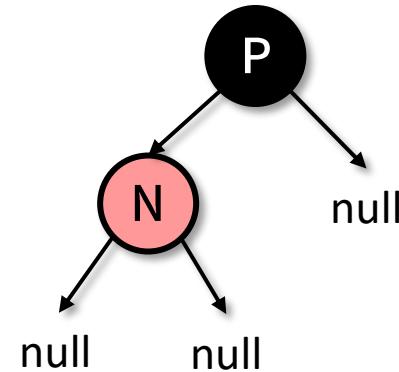
# Insertion (1)

- We always add a red node using ordinary insertion of binary search tree, and then correct the tree if rules are broken.

**case 1:  
empty tree**



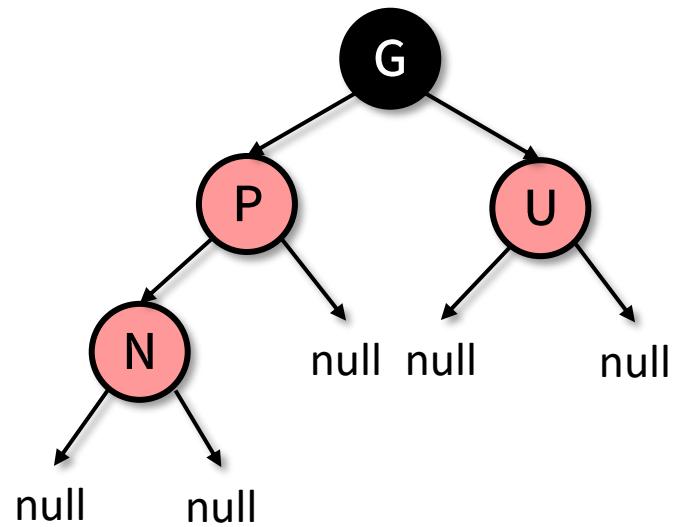
**case 2:  
add to a black parent**



the number of black nodes to the leaf node remains the same

# Red-Black Tree Insertion (3)

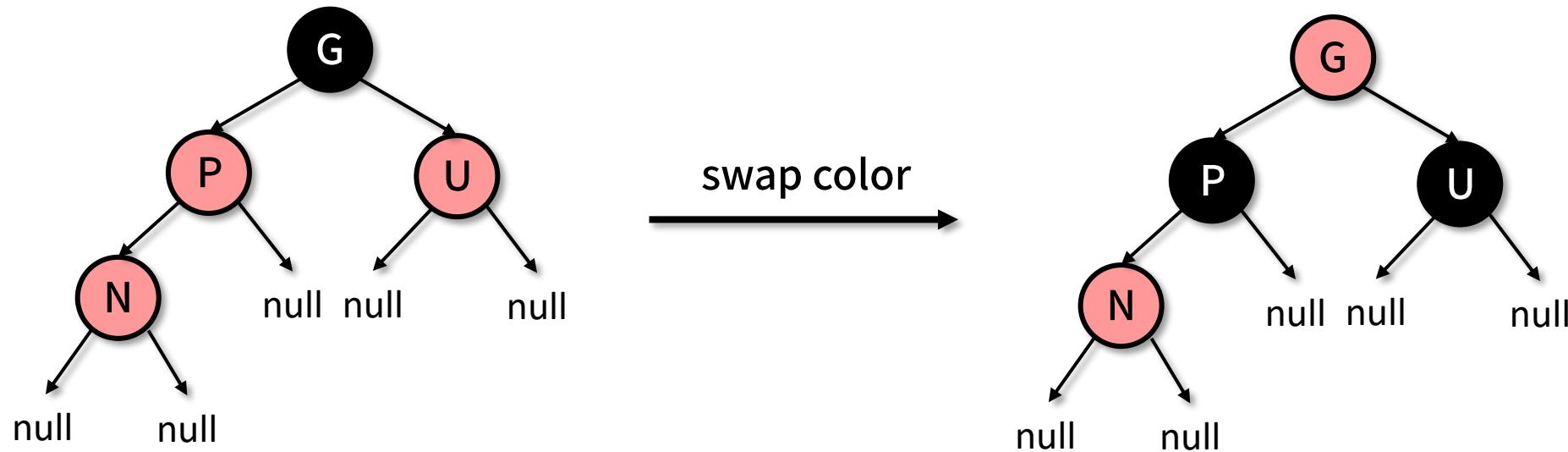
**case 3:**  
**add to a red parent with a red uncle**



# Red-Black Tree Insertion (4)

**case 3:**  
**add to a red parent with a red uncle**

iterate to parent of grandparent  
to check the rules

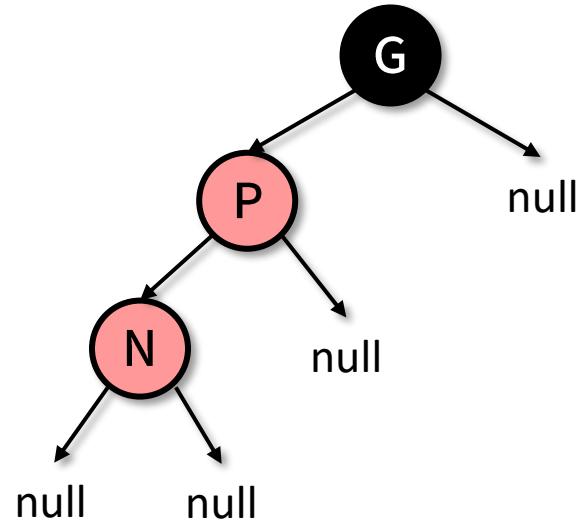


the number of black nodes to the leaf node remains the same

# Red-Black Tree Insertion (5)

**case 4: left left**

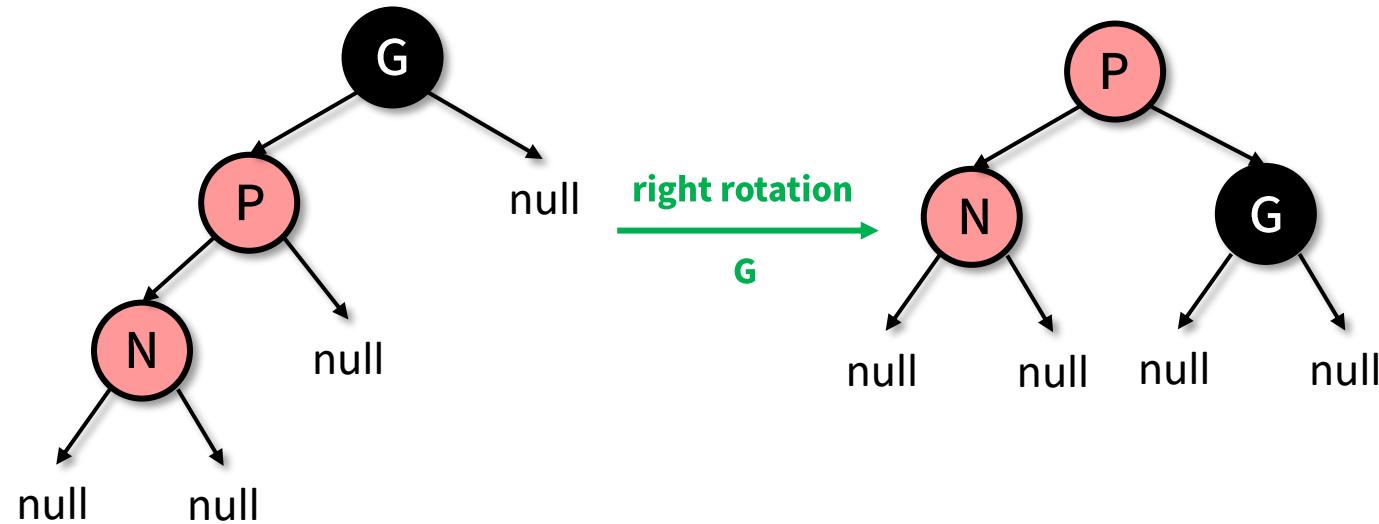
1. add to a red parent without uncle
2. new node is a left child of the parent
3. parent is a left child of the grandparent



# Red-Black Tree Insertion (6)

**case 4: left left**

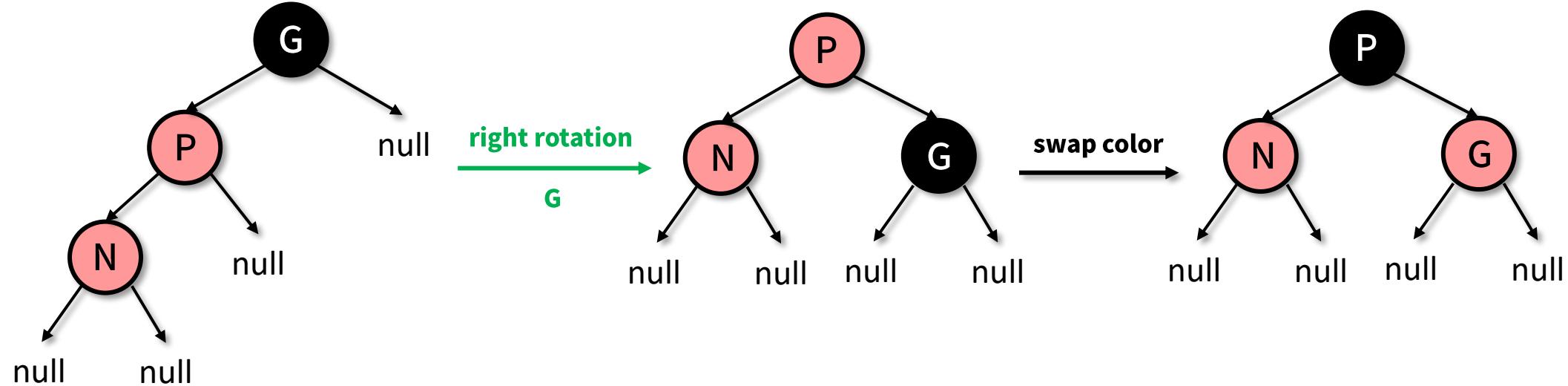
1. add to a red parent without uncle
2. new node is a left child of the parent
3. parent is a left child of the grandparent



# Red-Black Tree Insertion (7)

**case 4: left left**

1. add to a red parent without uncle
2. new node is a left child of the parent
3. parent is a left child of the grandparent

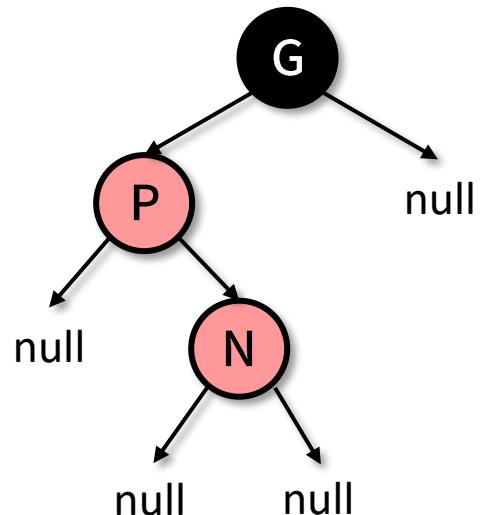


the number of black nodes to the leaf node remains the same

# Red-Black Tree Insertion (8)

**case 5: left right**

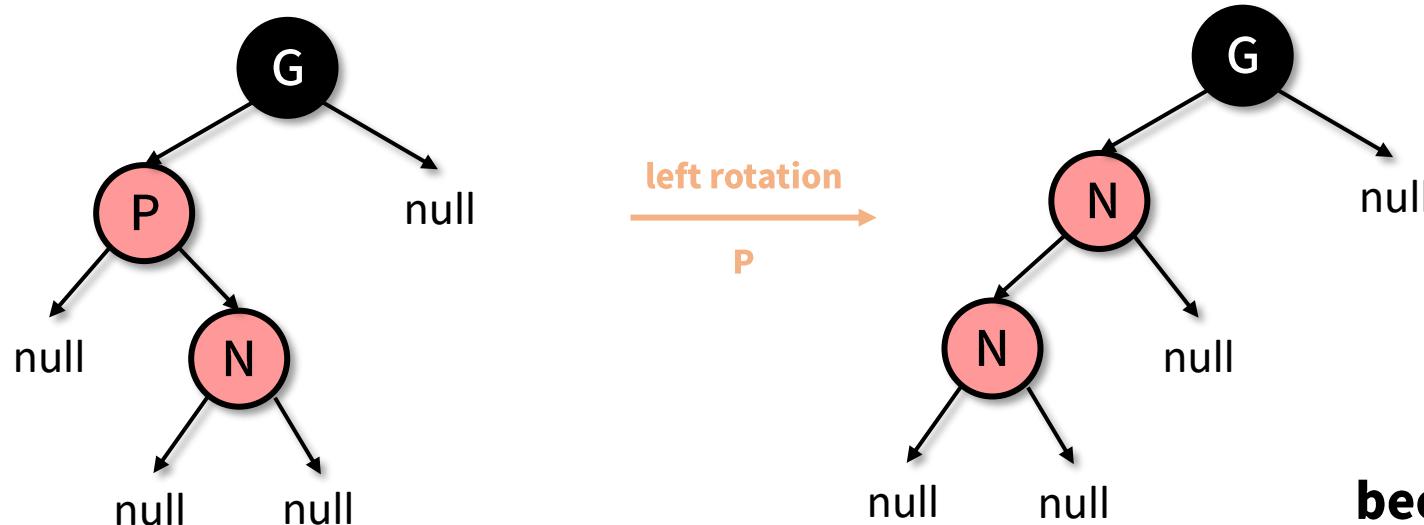
1. add to a red parent without uncle
2. new node is a right child of the parent
3. parent is a left child of the grandparent



# Red-Black Tree Insertion (9)

**case 5: left right**

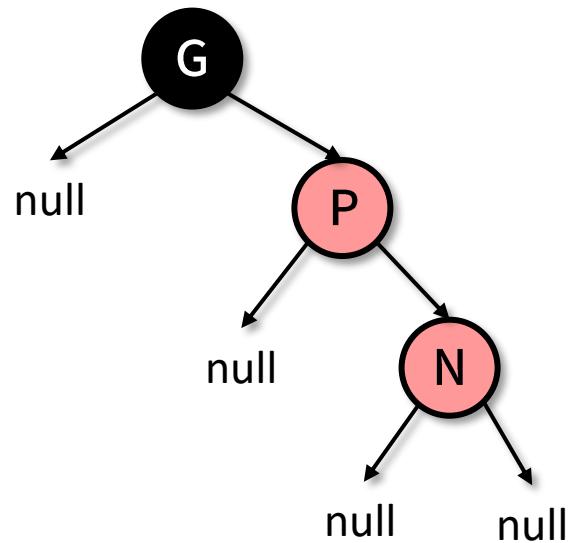
1. add to a red parent without uncle
2. new node is a right child of the parent
3. parent is a left child of the grandparent



Red-Black Tree  
**Insertion (10)**

**case 6: right right**

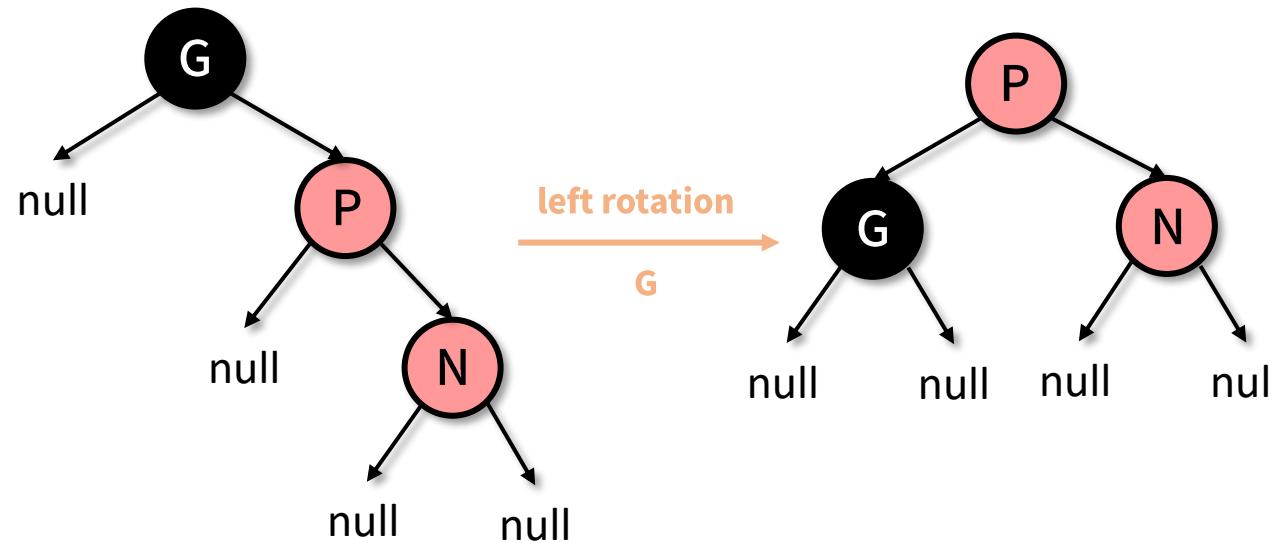
1. add to a red parent without uncle
2. new node is a right child of the parent
3. parent is a right child of the grandparent



# Red-Black Tree Insertion (11)

**case 6: right right**

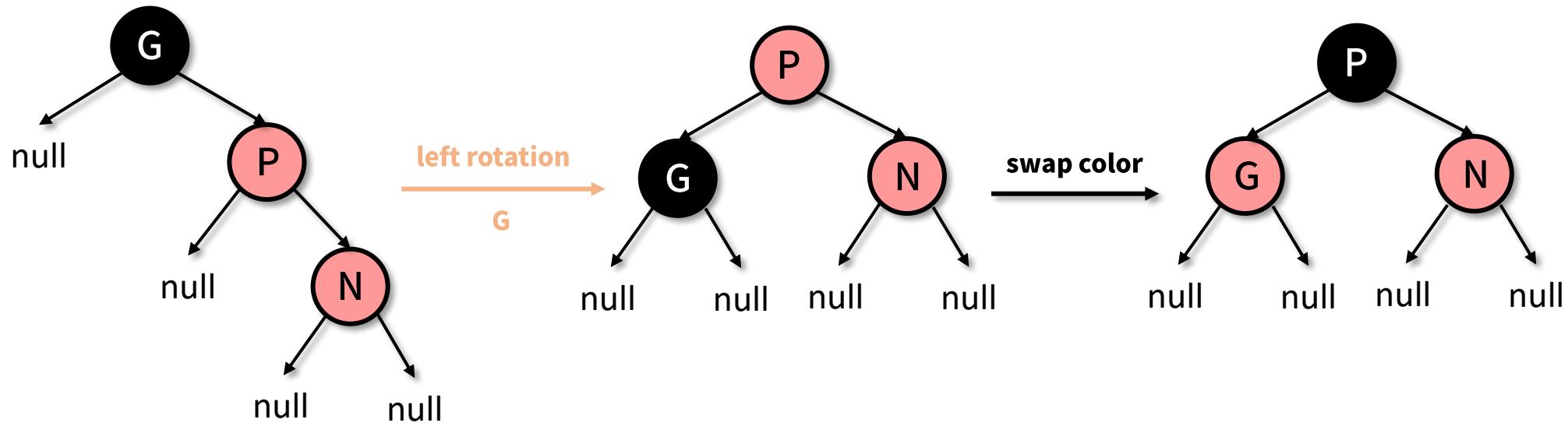
1. add to a red parent without uncle
2. new node is a right child of the parent
3. parent is a right child of the grandparent



# Red-Black Tree Insertion (12)

**case 6: right right**

1. add to a red parent without uncle
2. new node is a right child of the parent
3. parent is a right child of the grandparent

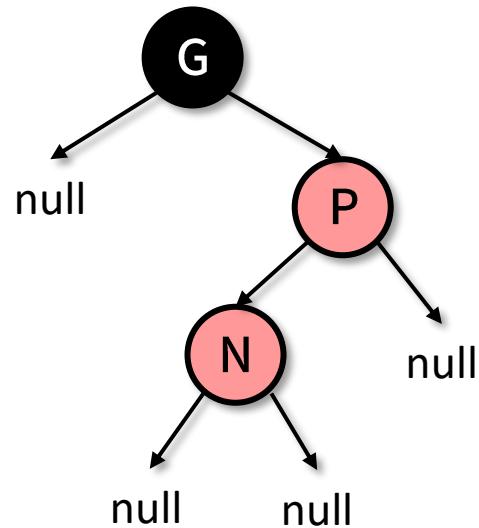


the number of black nodes to the leaf node remains the same

Red-Black Tree  
**Insertion (13)**

**case 7: right left**

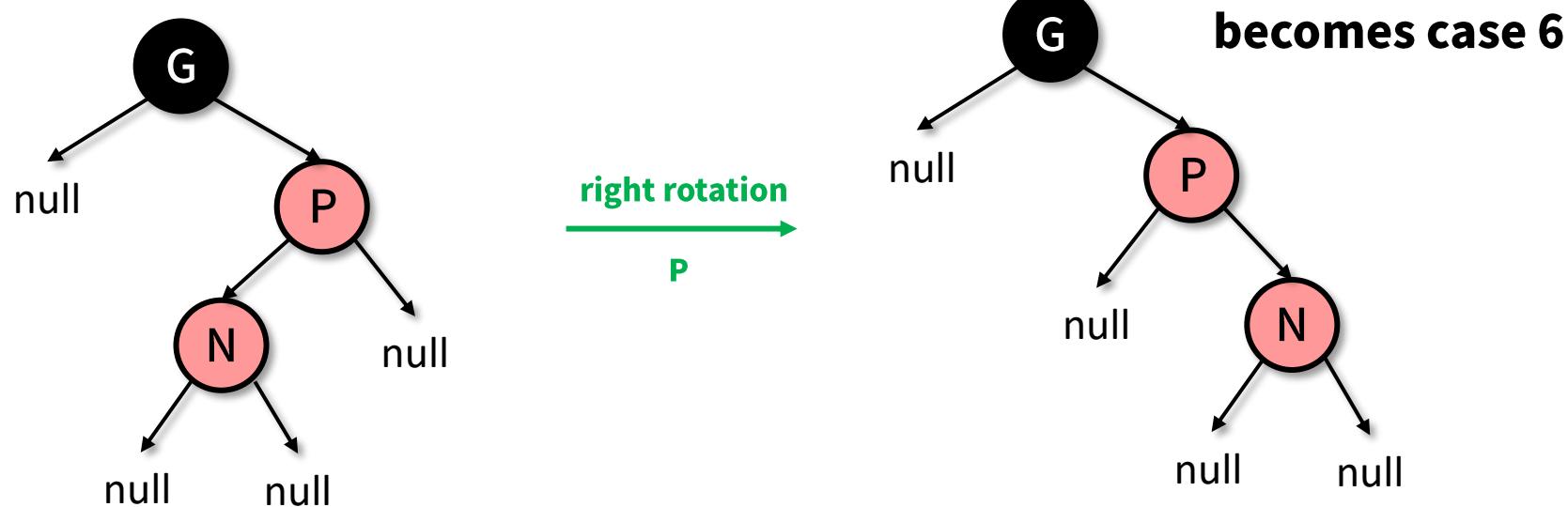
1. add to a red parent without uncle
2. new node is a left child of the parent
3. parent is a right child of the grandparent



# Red-Black Tree Insertion (14)

**case 7: right left**

1. add to a red parent without uncle
2. new node is a left child of the parent
3. parent is a right child of the grandparent

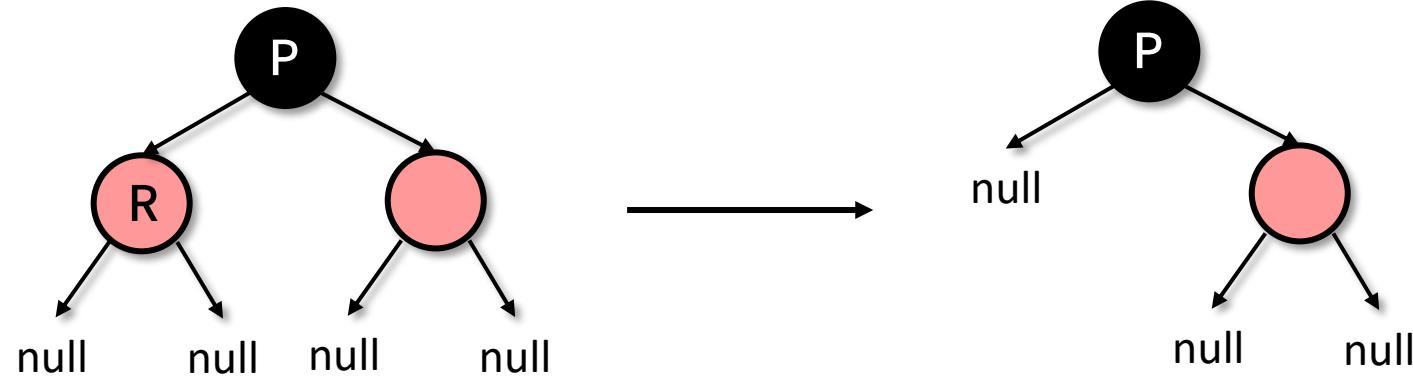


Red-Black Tree

# Deletion (1)

- Need custom operations for different cases, not using the ordinary deletion in binary search tree.

**case 1: remove any red leaf node**

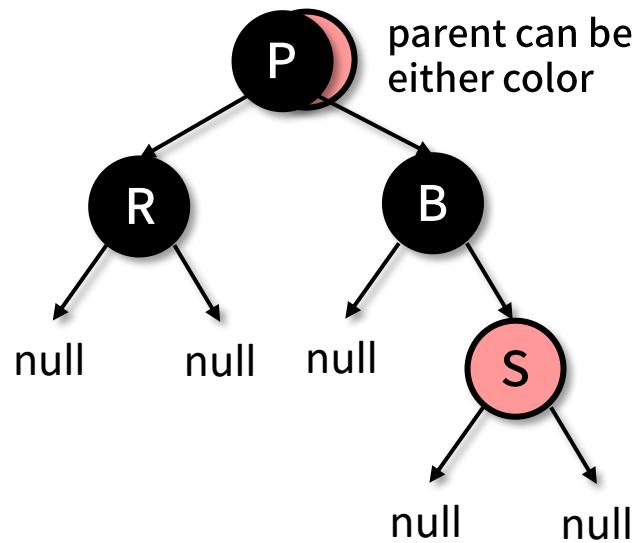


no need for correction

# Red-Black Tree Deletion (2)

**case 2:**

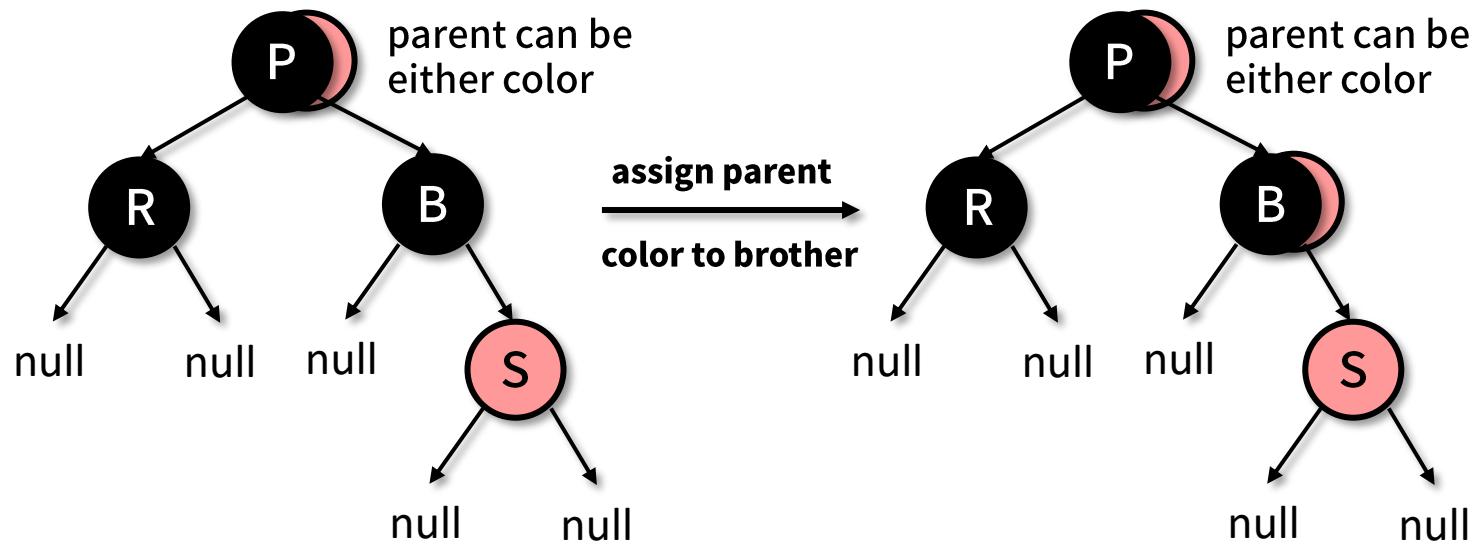
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red right child



# Red-Black Tree Deletion (3)

**case 2:**

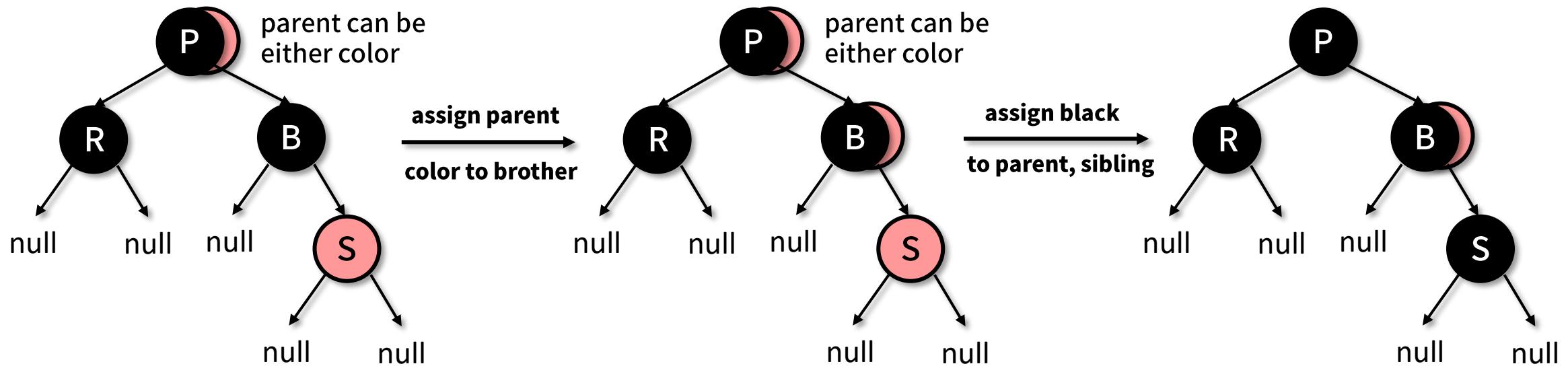
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red right child



# Red-Black Tree Deletion (4)

**case 2:**

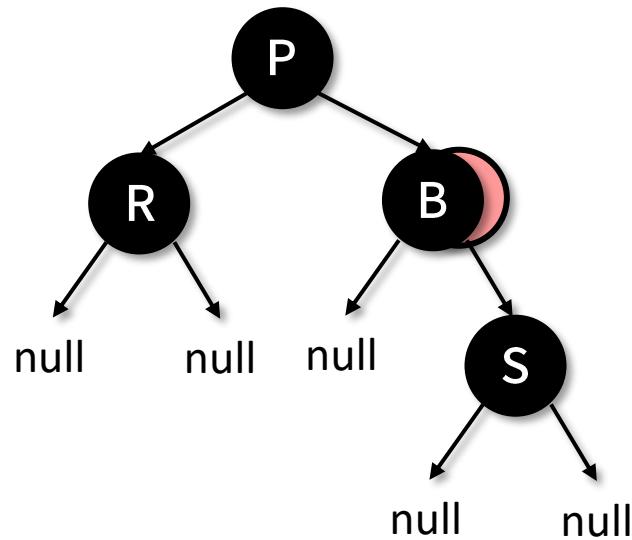
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red right child



# Red-Black Tree Deletion (5)

**case 2:**

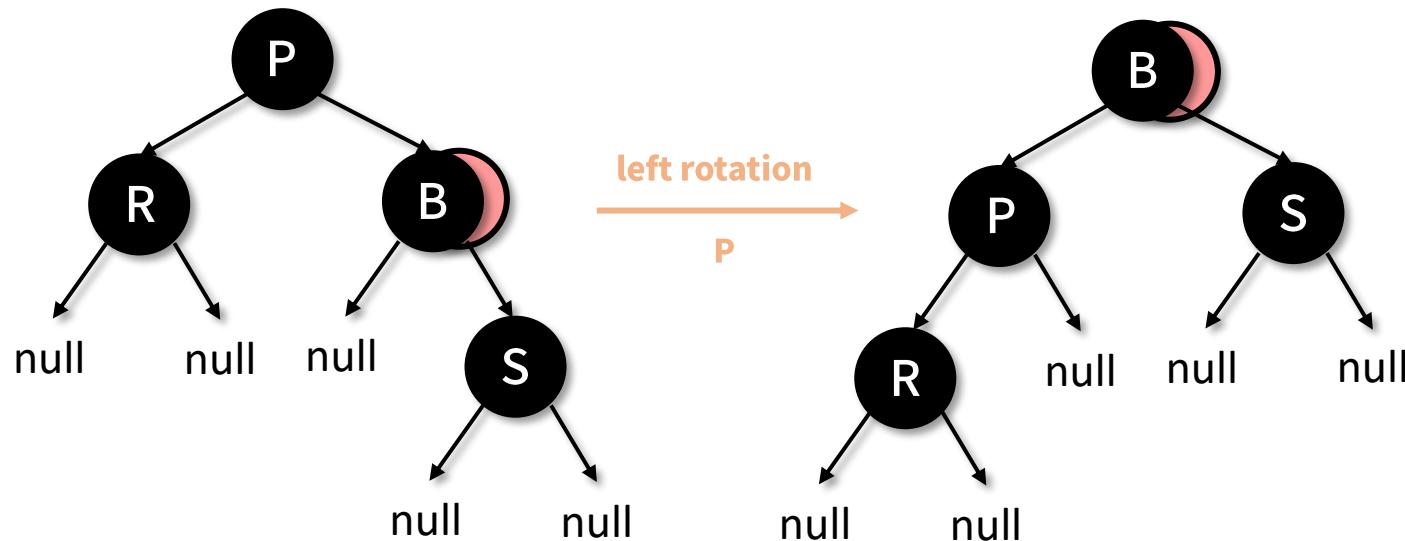
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red right child



# Red-Black Tree Deletion (6)

**case 2:**

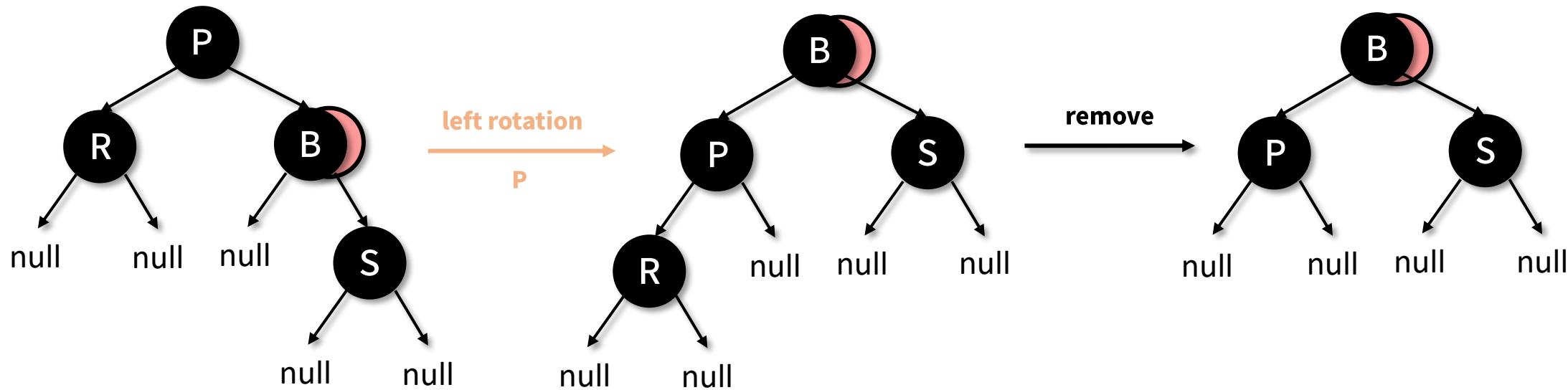
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red right child



# Red-Black Tree Deletion (7)

**case 2:**

1. remove black leaf node
2. the brother is a black node
3. the brother has only one red right child

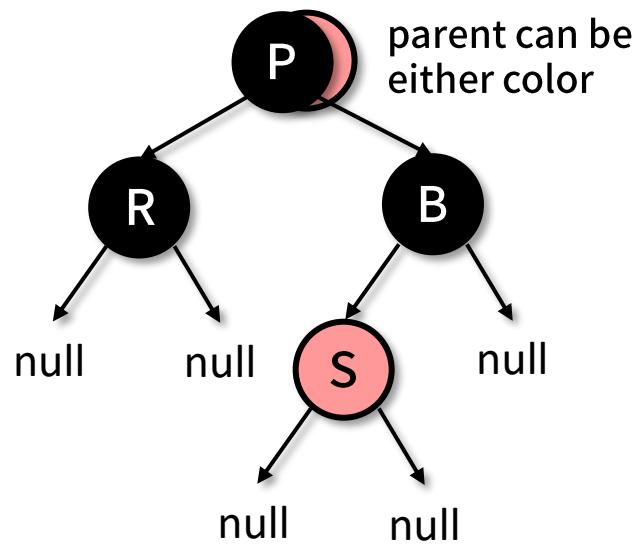


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (8)

**case 3:**

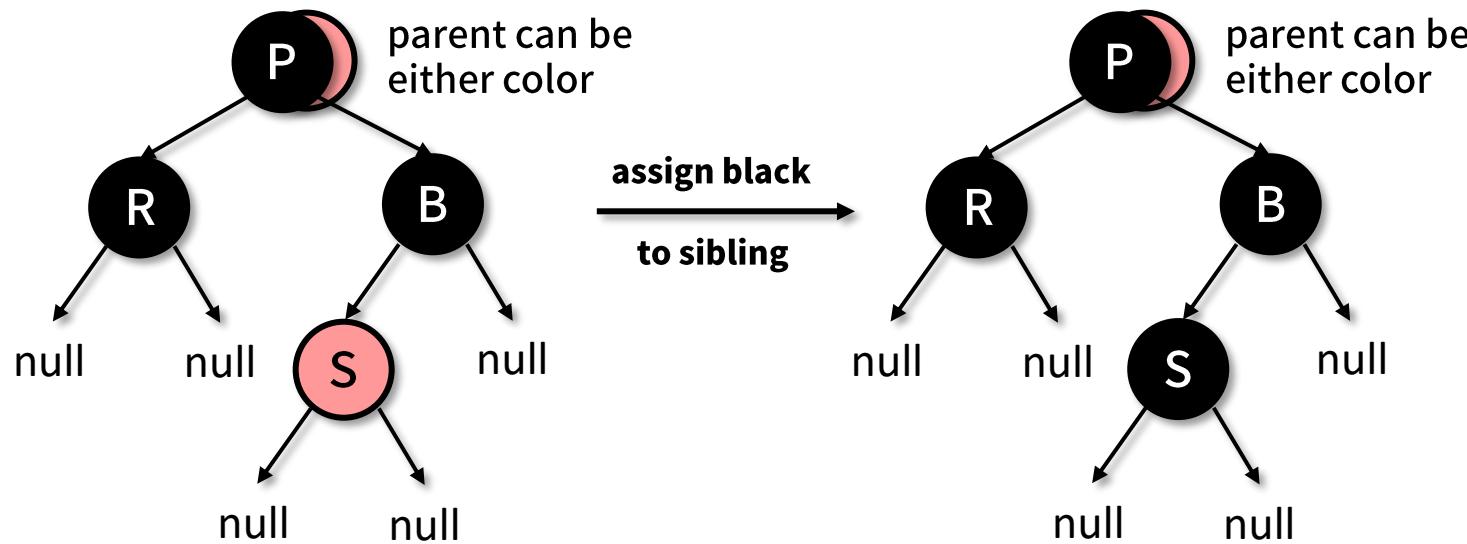
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red left child



# Red-Black Tree Deletion (9)

**case 3:**

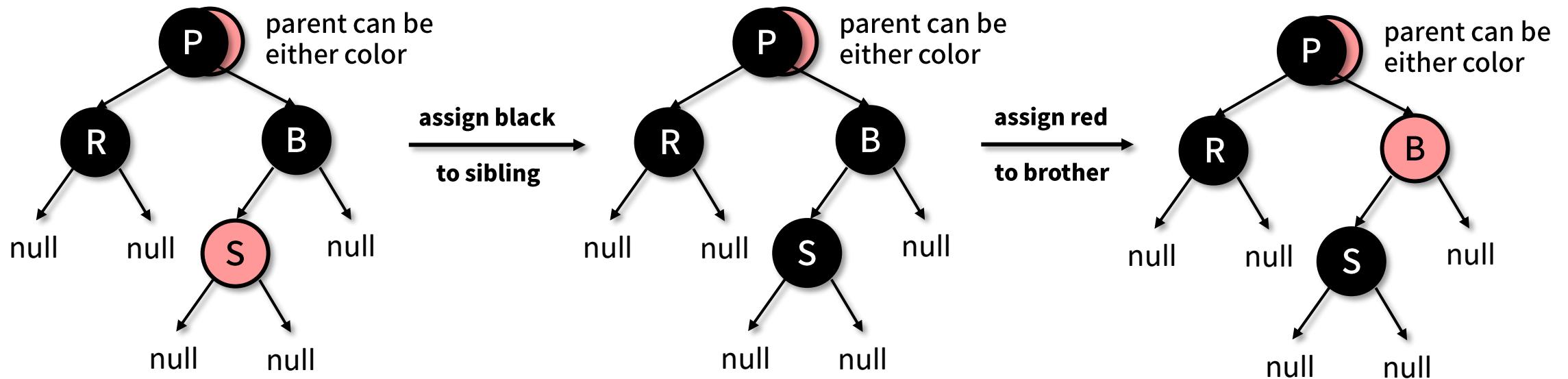
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red left child



# Red-Black Tree Deletion (10)

**case 3:**

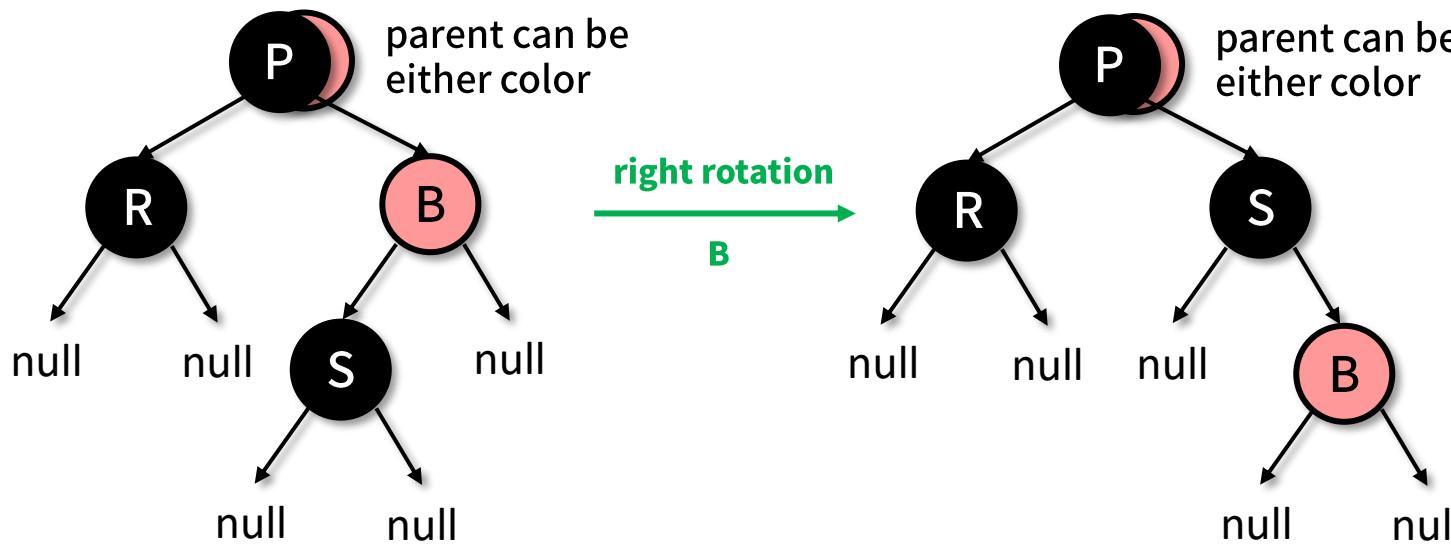
1. remove black leaf node
2. the brother is a black node
3. the brother has only one red left child



# Red-Black Tree Deletion (11)

**case 3:**

1. remove black leaf node
2. the brother is a black node
3. the brother has only one red left child

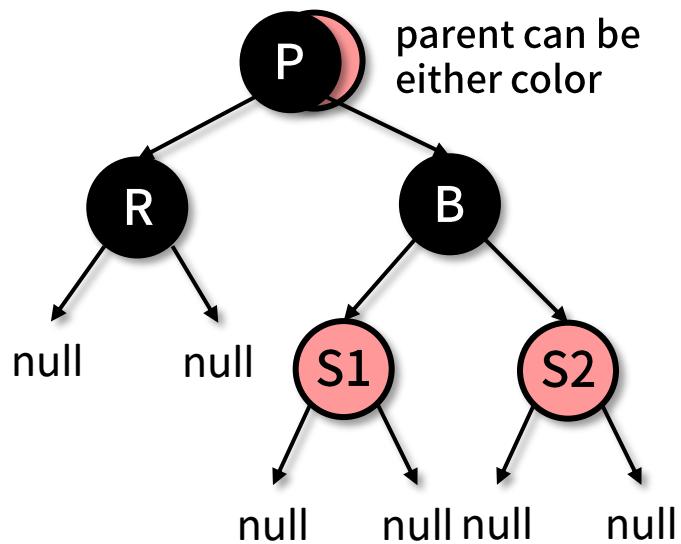


**becomes case 2**

# Red-Black Tree Deletion (12)

**case 4:**

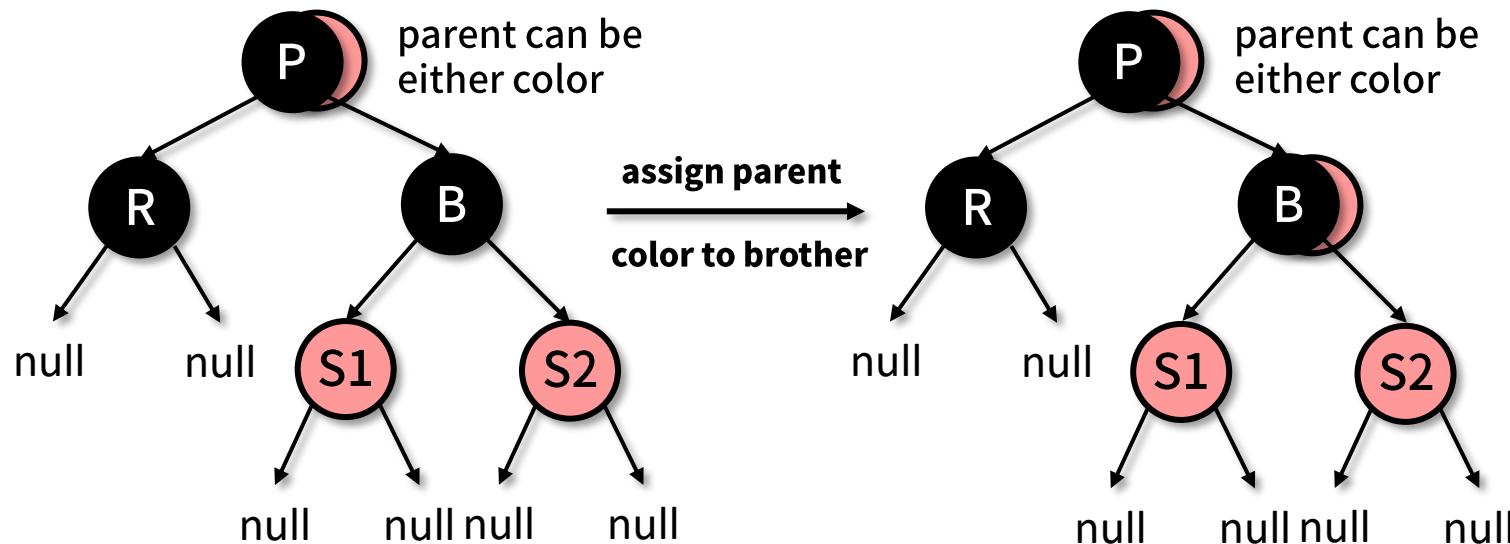
1. remove black leaf node
2. the brother is a black node
3. the brother has two red children



# Red-Black Tree Deletion (13)

**case 4:**

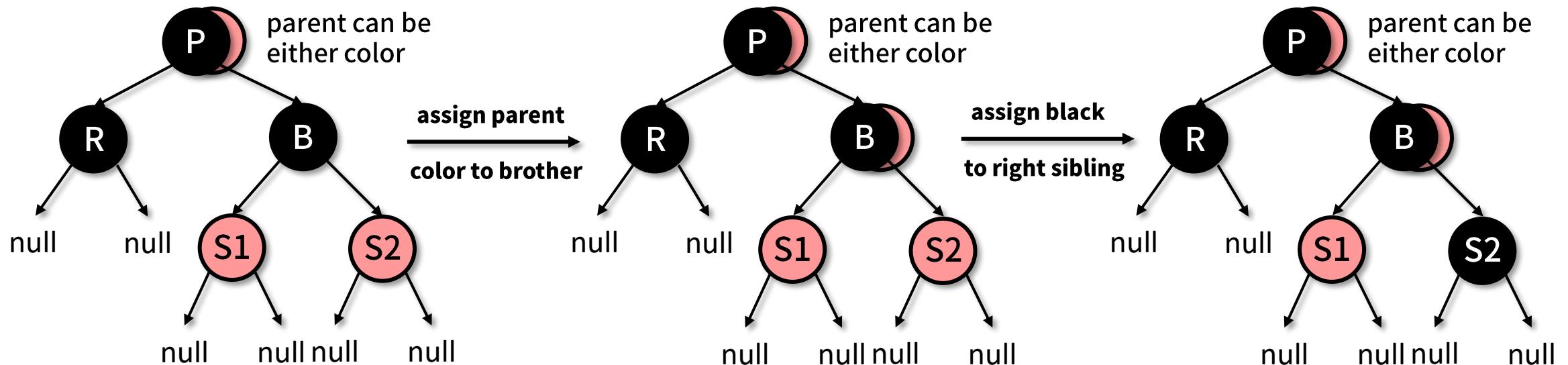
1. remove black leaf node
2. the brother is a black node
3. the brother has two red children



# Red-Black Tree Deletion (14)

**case 4:**

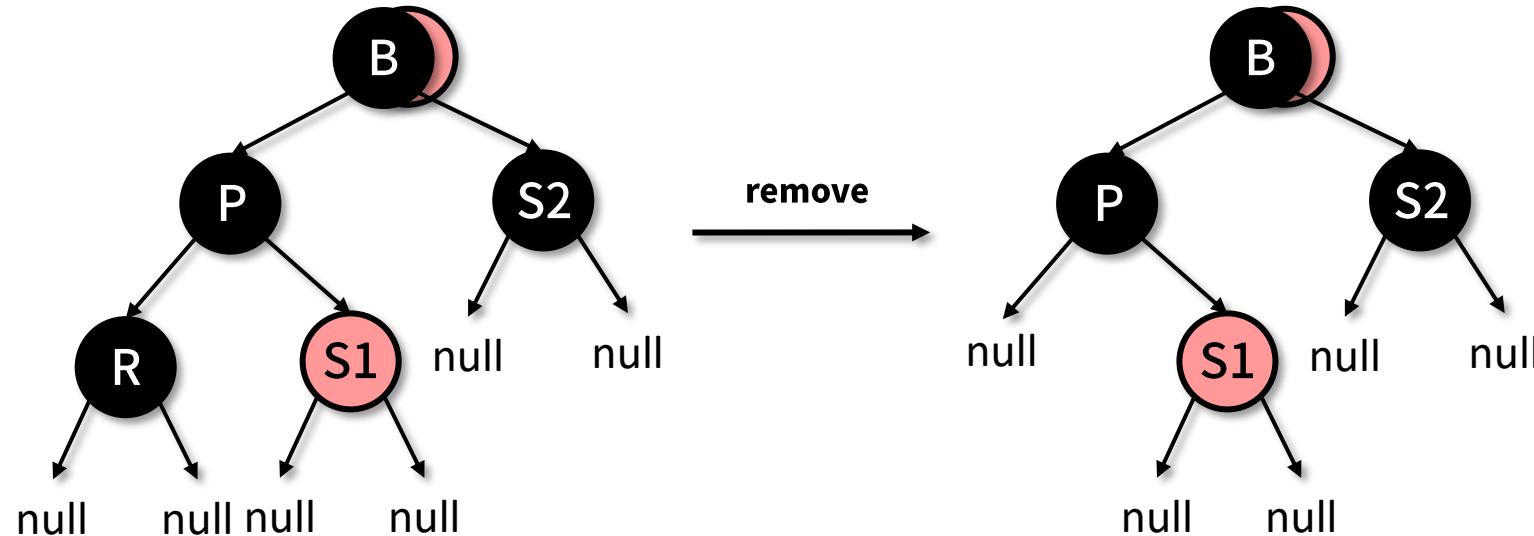
1. remove black leaf node
2. the brother is a black node
3. the brother has two red children



# Red-Black Tree Deletion (15)

**case 4:**

1. remove black leaf node
2. the brother is a black node
3. the brother has two red children

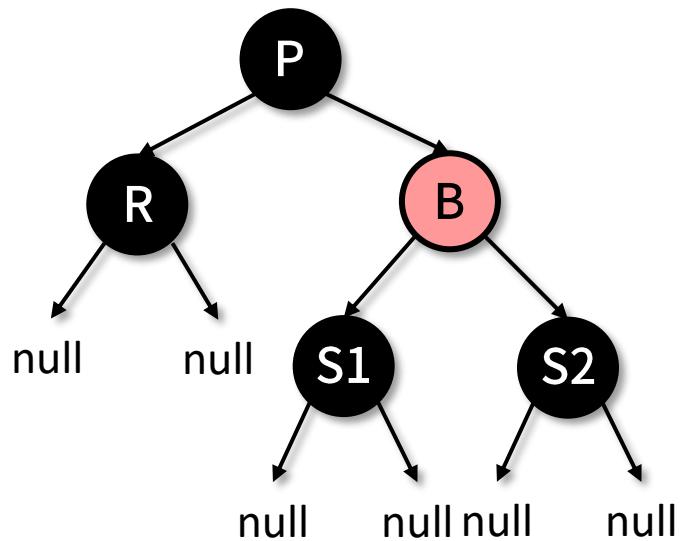


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (16)

**case 5:**

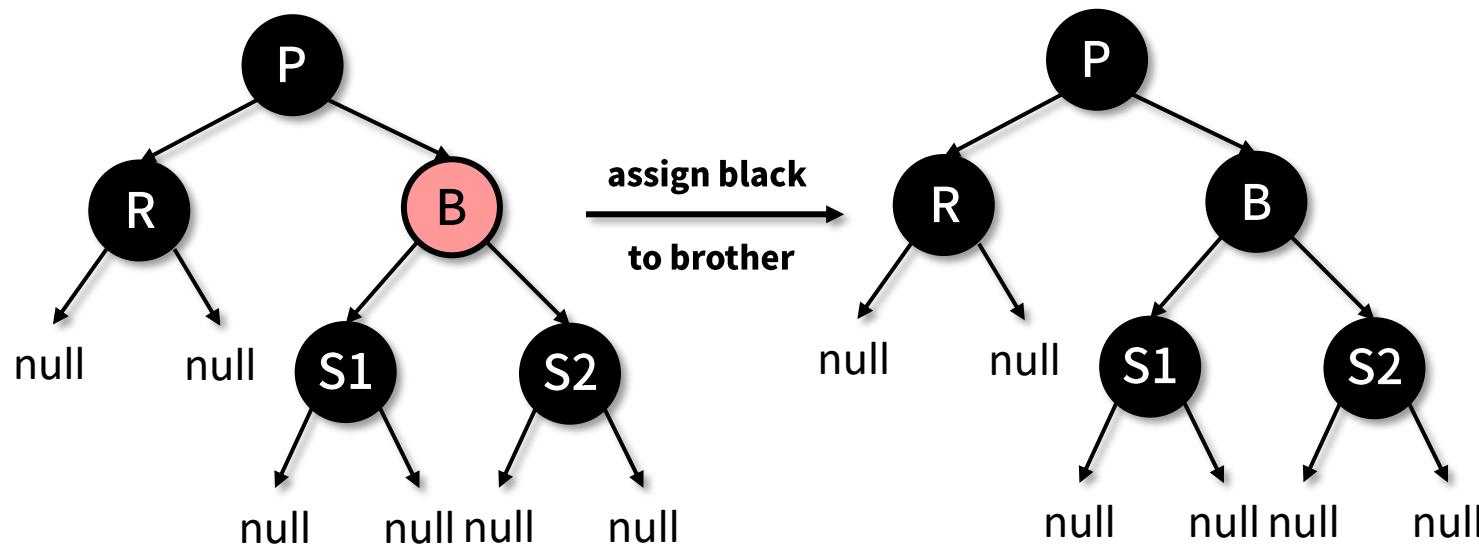
1. remove black leaf node
2. the brother is a red node
3. the brother has two red children



# Red-Black Tree Deletion (17)

**case 5:**

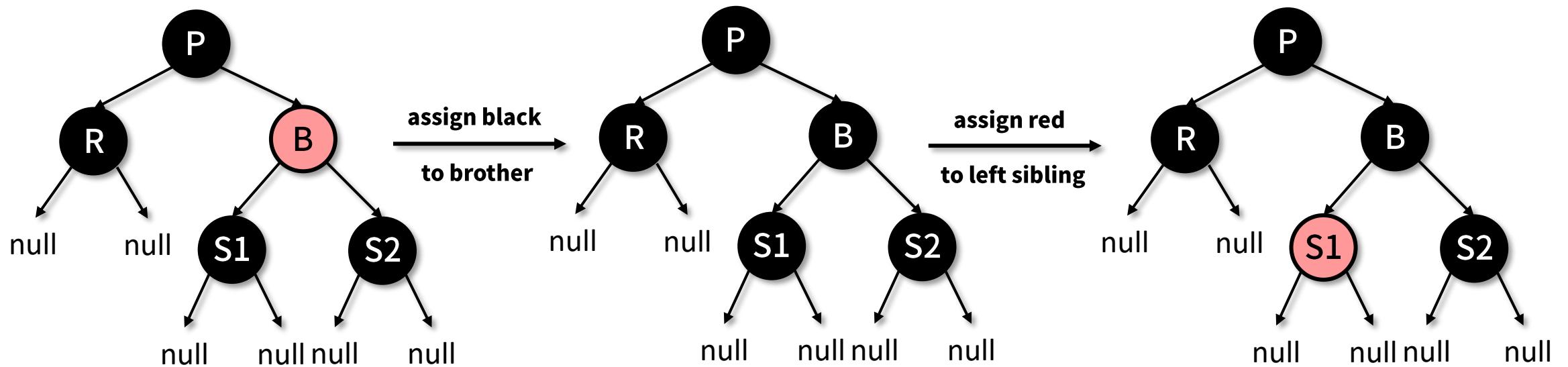
1. remove black leaf node
2. the brother is a red node
3. the brother has two red children



# Red-Black Tree Deletion (18)

**case 5:**

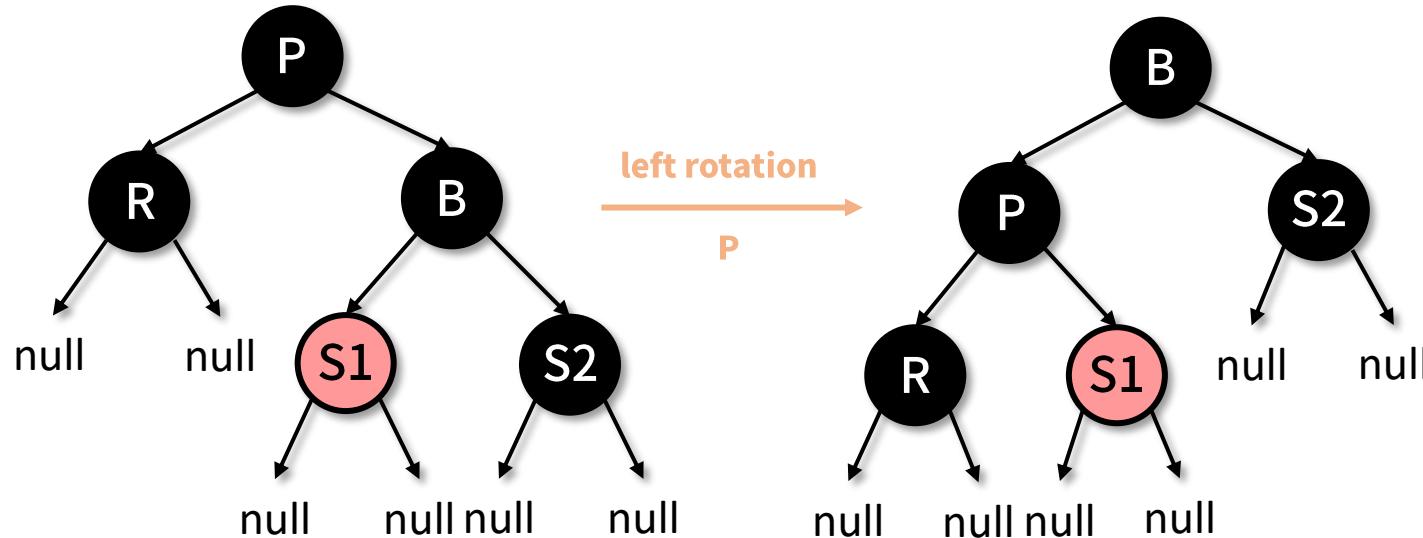
1. remove black leaf node
2. the brother is a red node
3. the brother has two red children



# Red-Black Tree Deletion (19)

**case 5:**

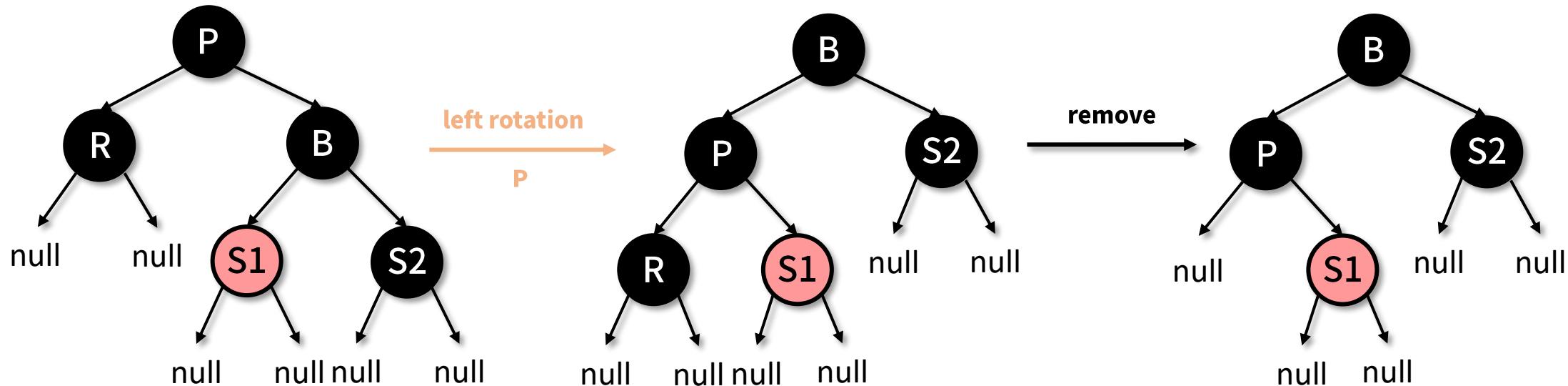
1. remove black leaf node
2. the brother is a red node
3. the brother has two red children



# Red-Black Tree Deletion (20)

**case 5:**

1. remove black leaf node
2. the brother is a red node
3. the brother has two red children

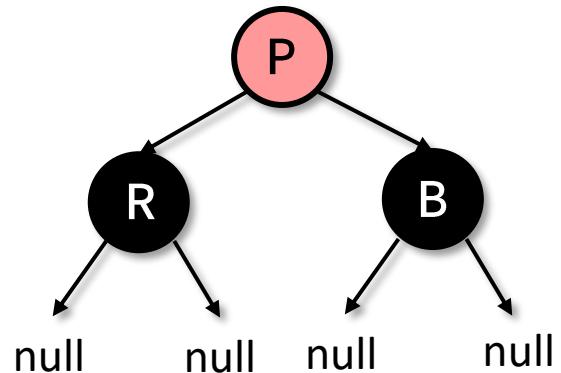


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (21)

**case 6:**

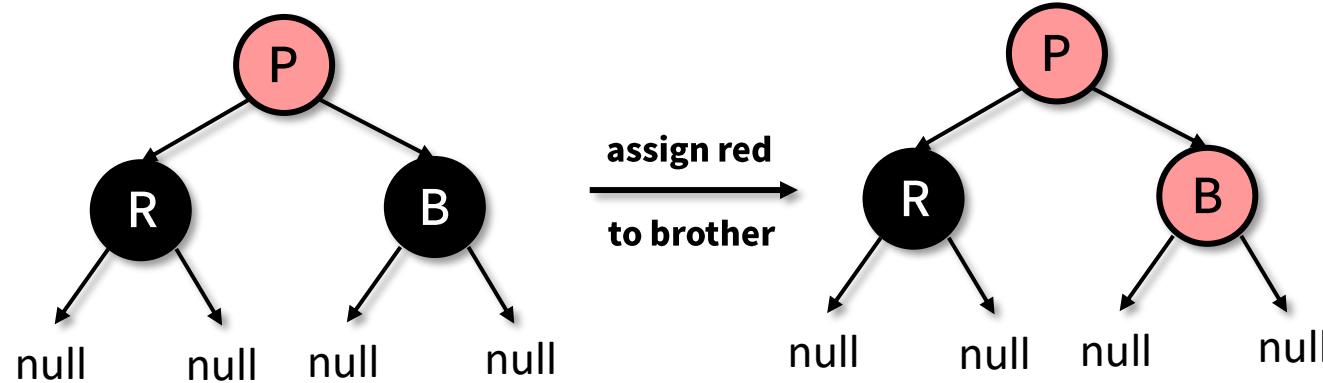
1. remove black leaf node
2. the parent is red
3. the brother is a black node
4. the brother has no children



# Red-Black Tree Deletion (22)

**case 6:**

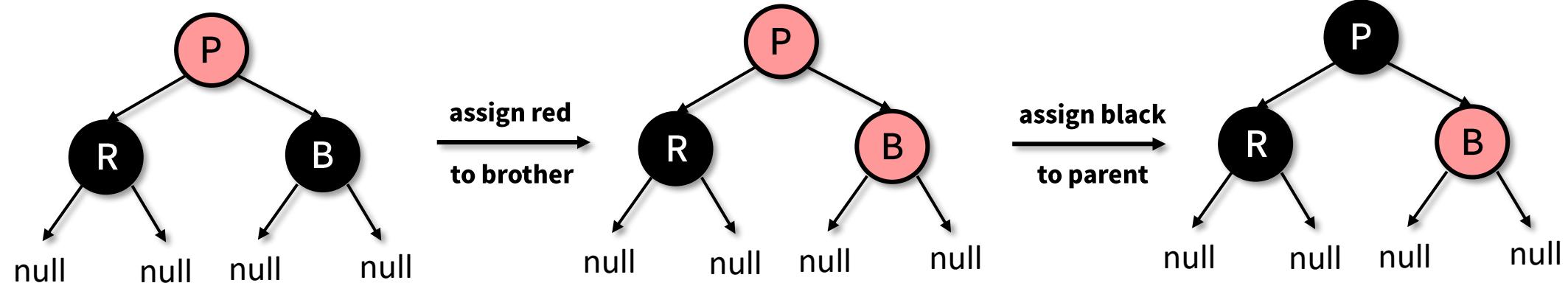
1. remove black leaf node
2. the parent is red
3. the brother is a black node
4. the brother has no children



# Red-Black Tree Deletion (23)

**case 6:**

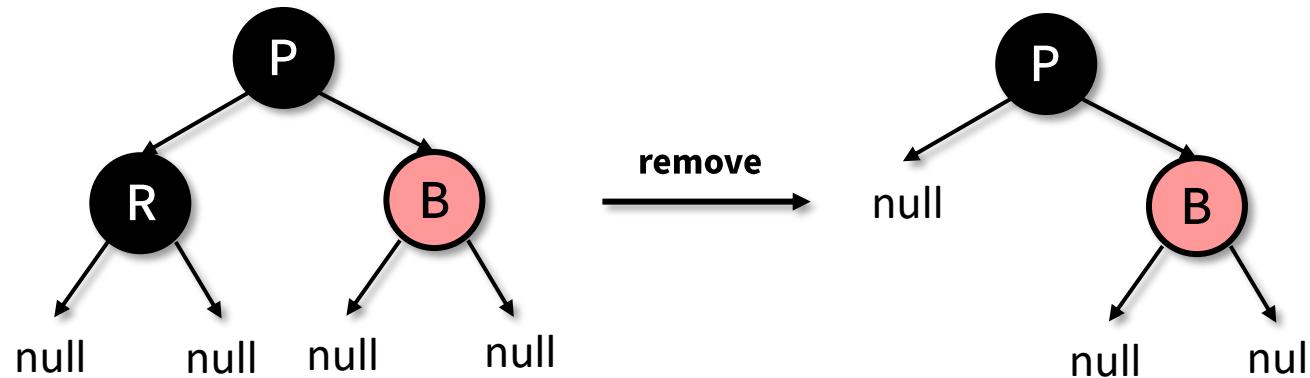
1. remove black leaf node
2. the parent is red
3. the brother is a black node
4. the brother has no children



# Red-Black Tree Deletion (24)

**case 6:**

1. remove black leaf node
2. the parent is red
3. the brother is a black node
4. the brother has no children

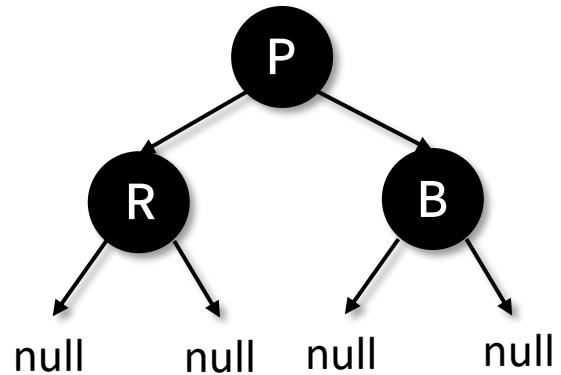


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (25)

**case 7:**

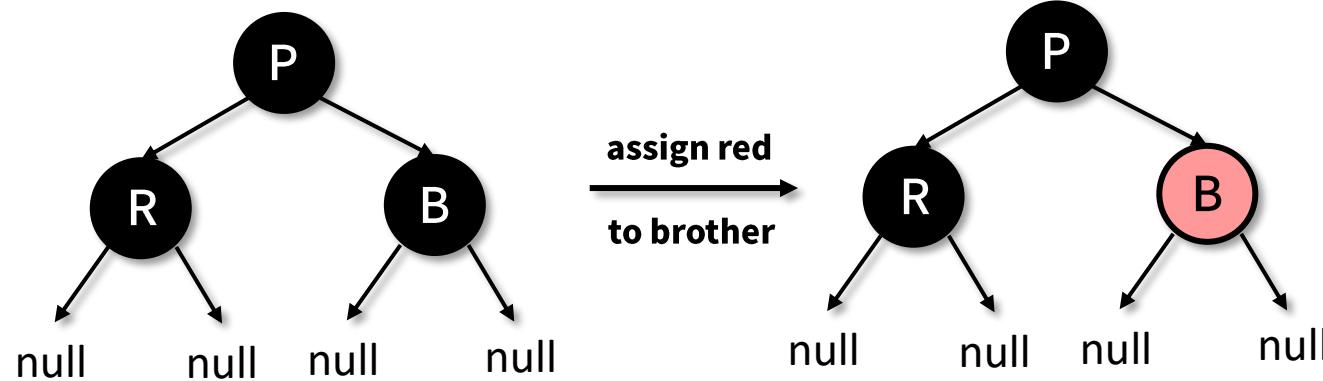
1. remove black leaf node
2. the parent is black
3. the brother is a black node
4. the brother has no children



# Red-Black Tree Deletion (26)

**case 7:**

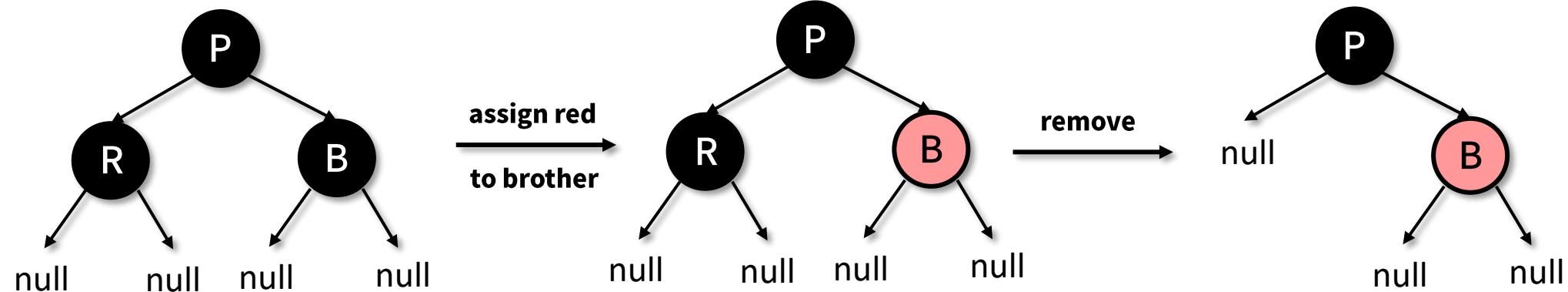
1. remove black leaf node
2. the parent is black
3. the brother is a black node
4. the brother has no children



# Red-Black Tree Deletion (27)

**case 7:**

1. remove black leaf node
2. the parent is black
3. the brother is a black node
4. the brother has no children

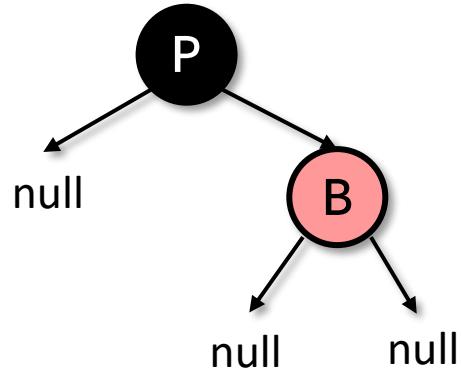


the number of black nodes to the leaf node decrease by one.  
need to further correct for it.

# Red-Black Tree Deletion (28)

**case 7.1:**

**1. parent is root**

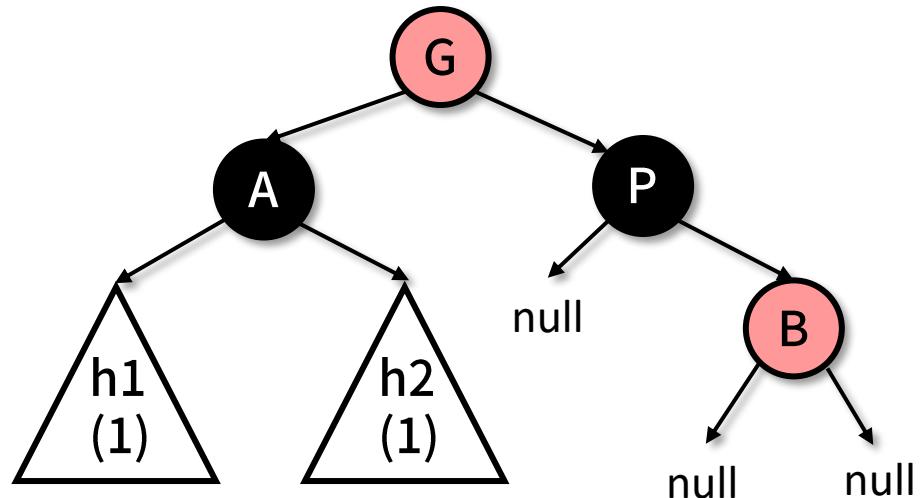


no need for correction

# Red-Black Tree Deletion (29)

**case 7.2:**

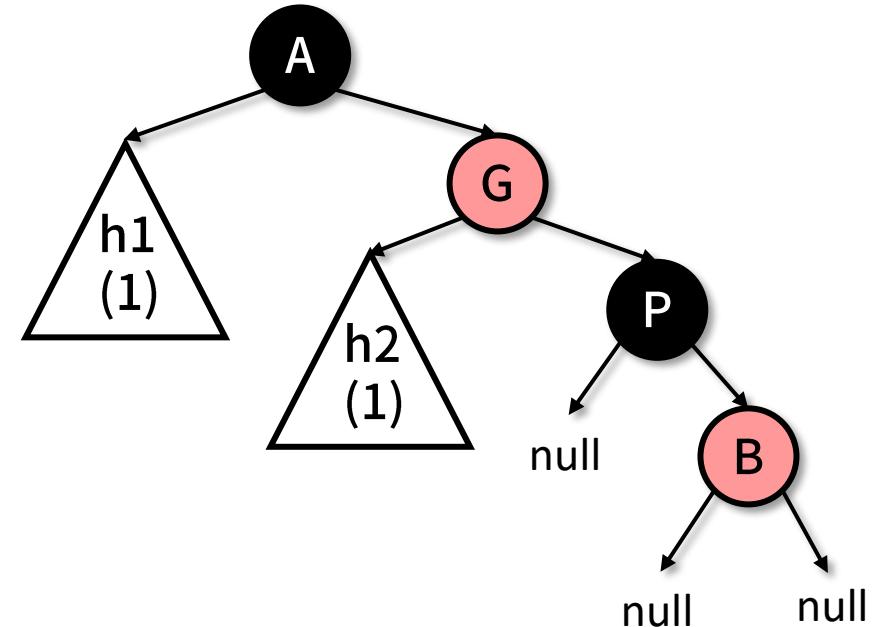
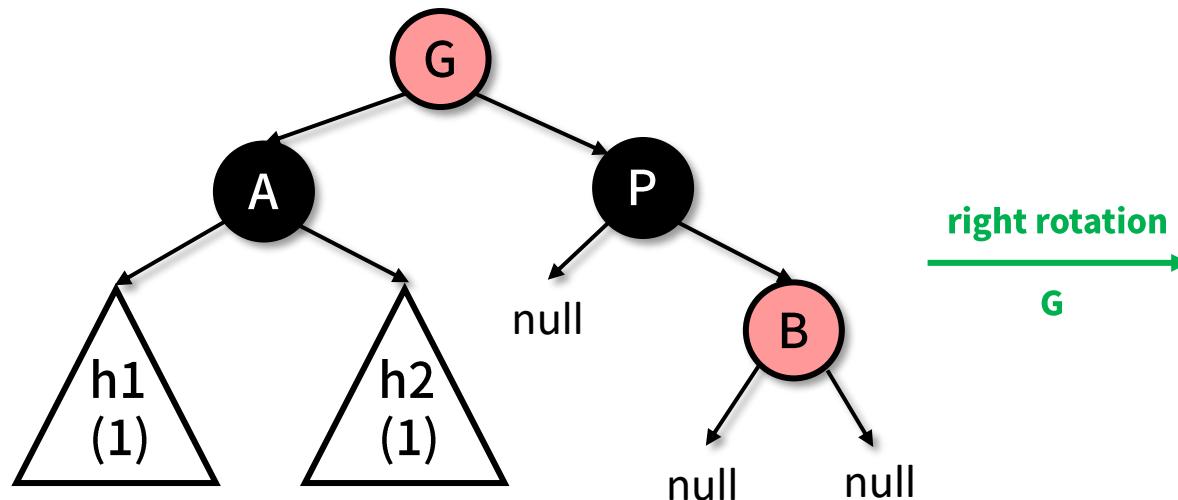
**1. grandparent is red (number of black: 2)**



# Red-Black Tree Deletion (30)

case 7.2:

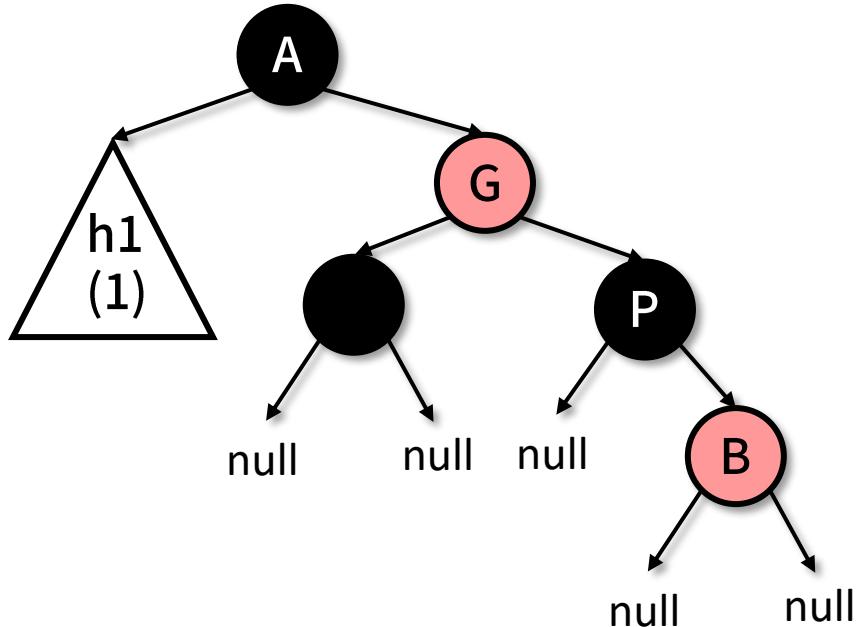
1. grandparent is red (number of black: 2)



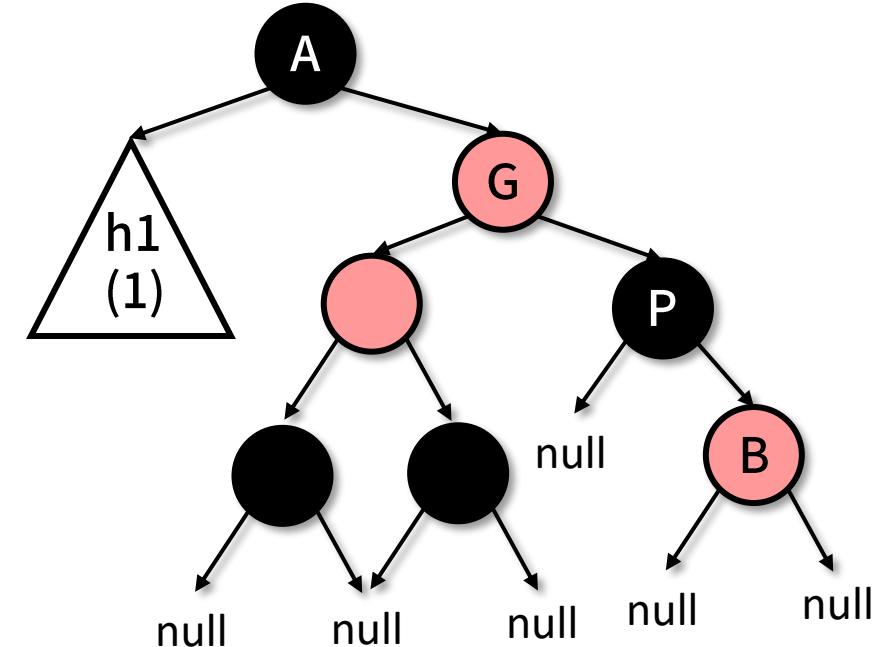
# Red-Black Tree Deletion (31)

**case 7.2:**

**1. grandparent is red (number of black: 2)**



no need for correction

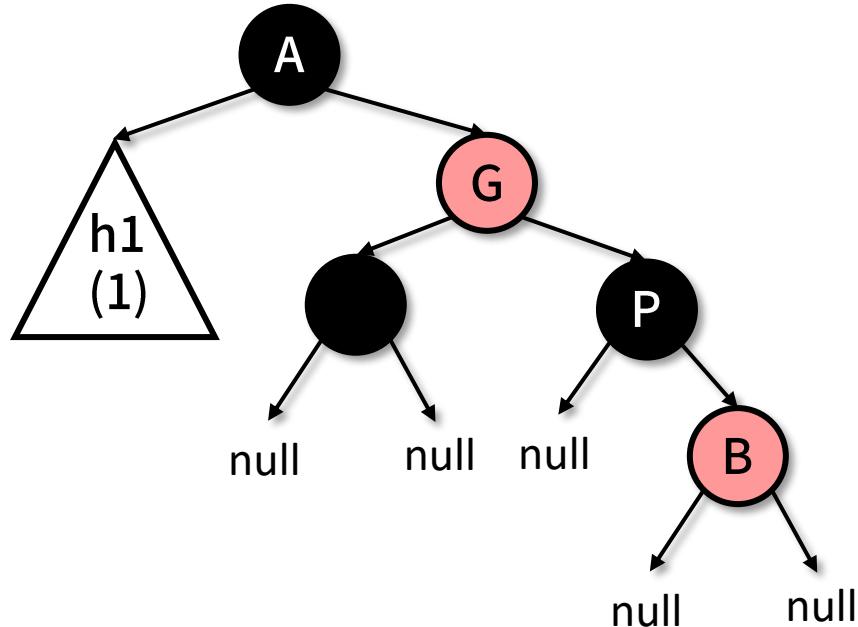


need to swap the color

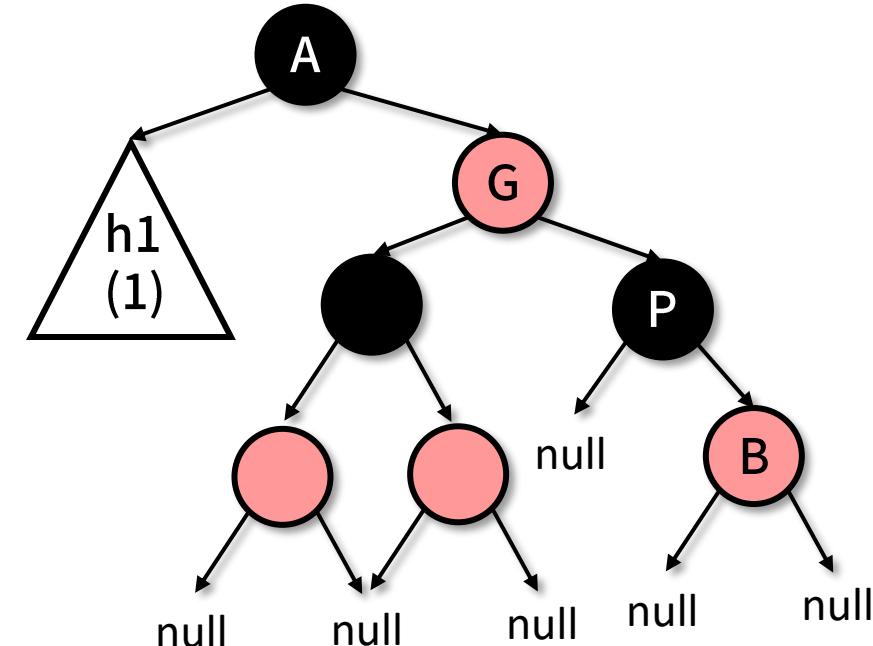
# Red-Black Tree Deletion (32)

case 7.2:

1. grandparent is red (number of black: 2)



no need for correction



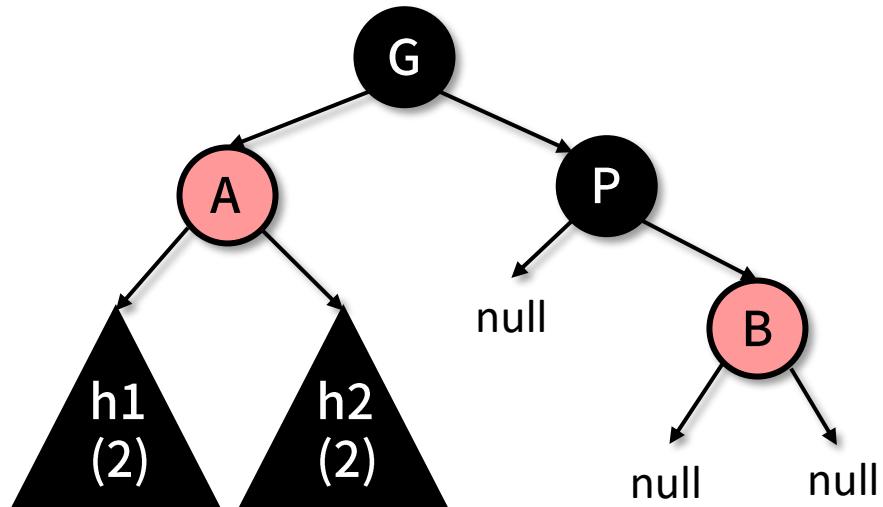
need to swap the color

the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (33)

**case 7.3:**

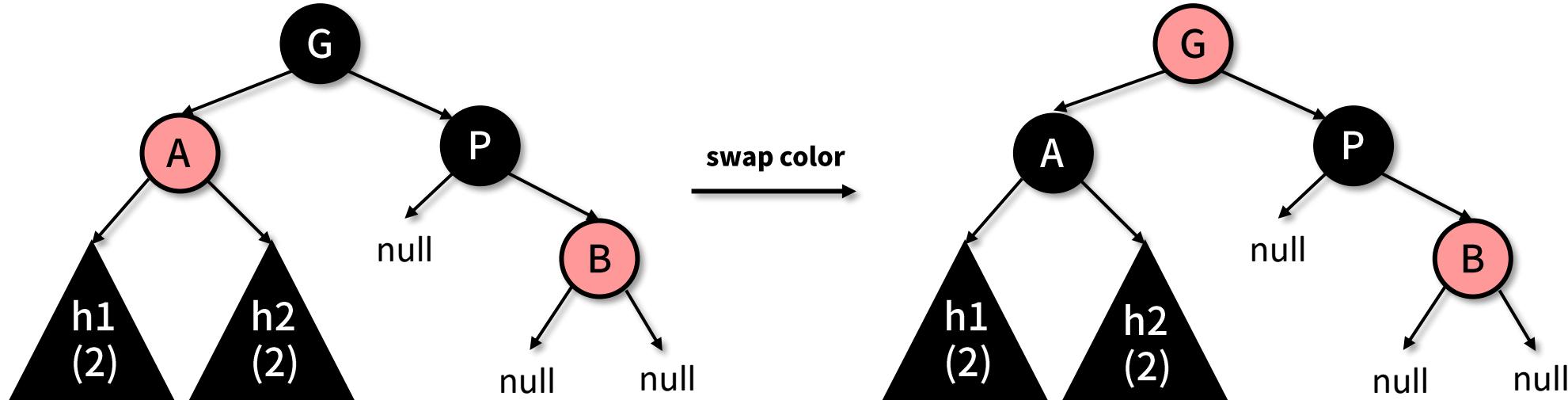
1. **grandparent is black (number of black: 3)**
2. **brother of parent is red**



# Red-Black Tree Deletion (34)

**case 7.3:**

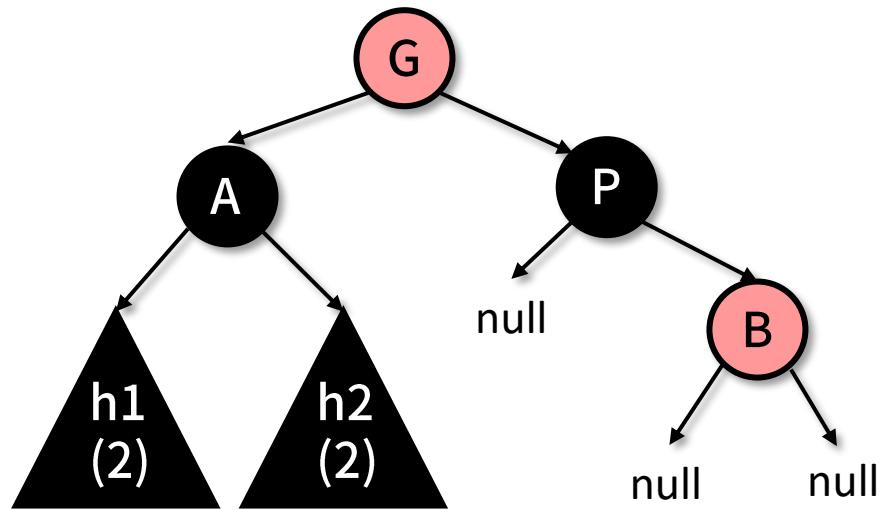
1. **grandparent is black (number of black: 3)**
2. **brother of parent is red**



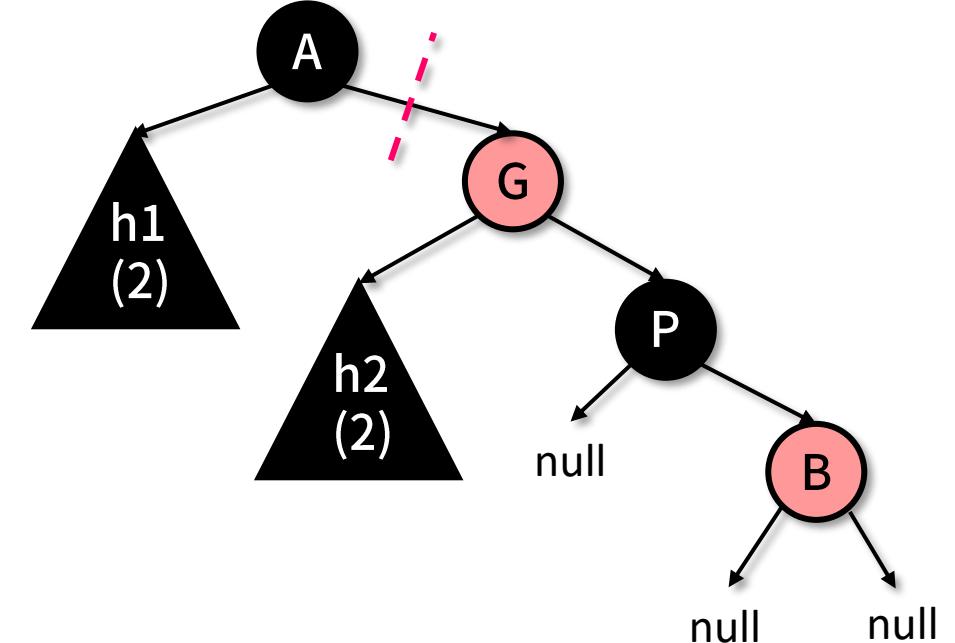
# Red-Black Tree Deletion (35)

**case 7.3:**

1. **grandparent is black (number of black: 3)**
2. **brother of parent is red**



right rotation  
G

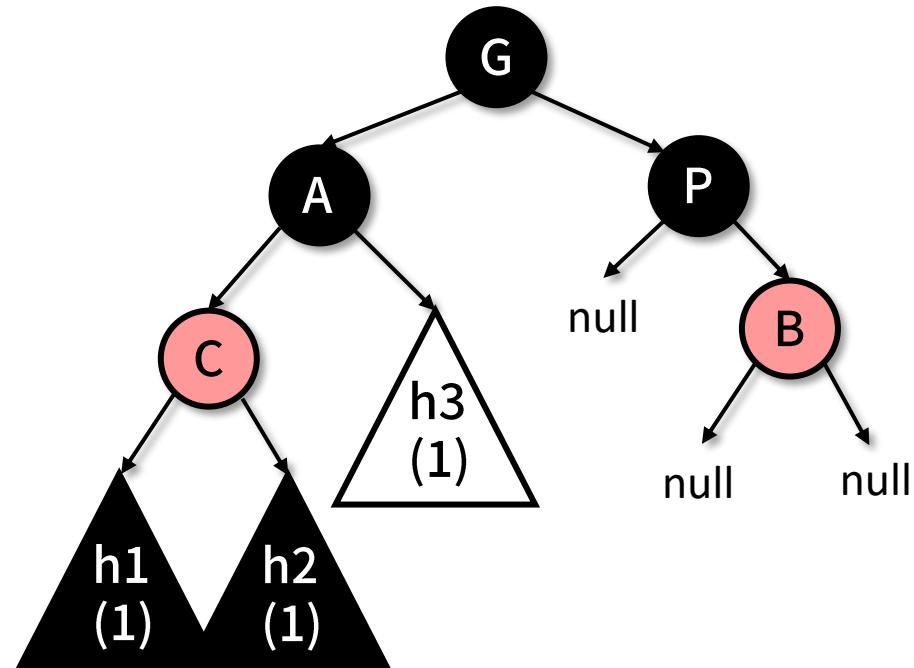


becomes case 7.2

# Red-Black Tree Deletion (36)

**case 7.4:**

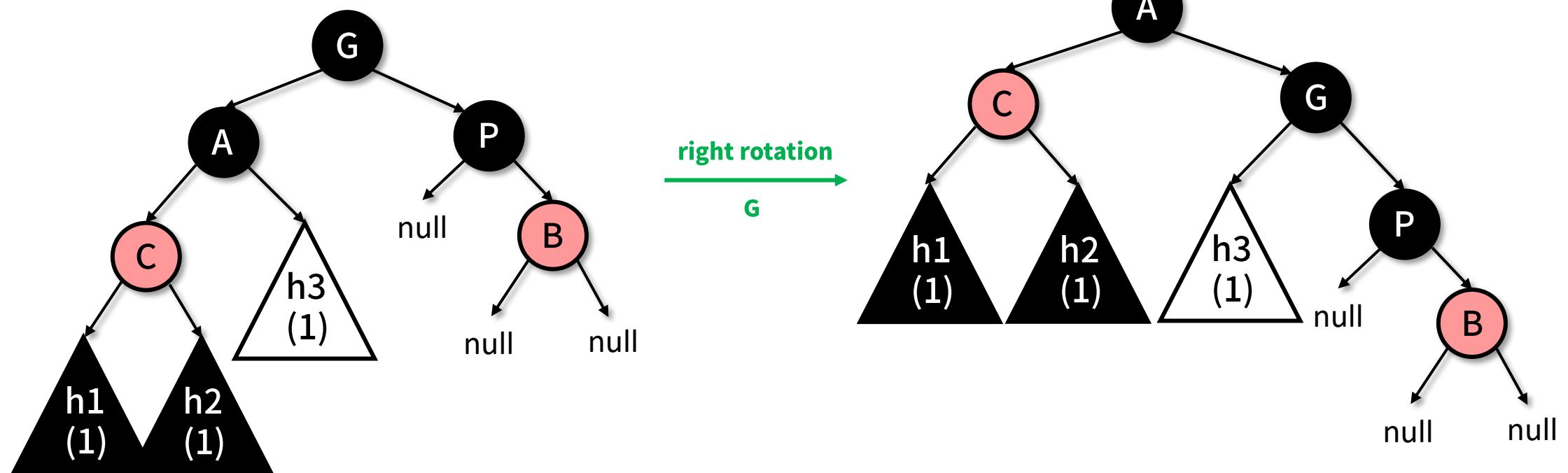
1. **grandparent is black (number of black: 3)**
2. **brother of parent is black**
3. **left child of brother of parent is red.**



# Red-Black Tree Deletion (37)

**case 7.4:**

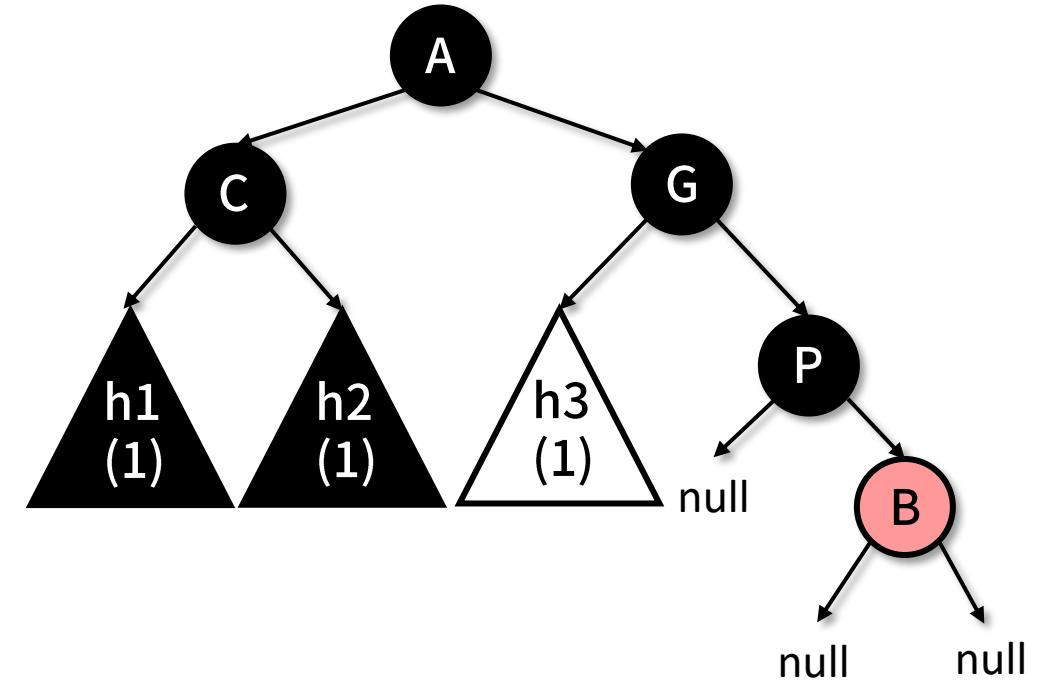
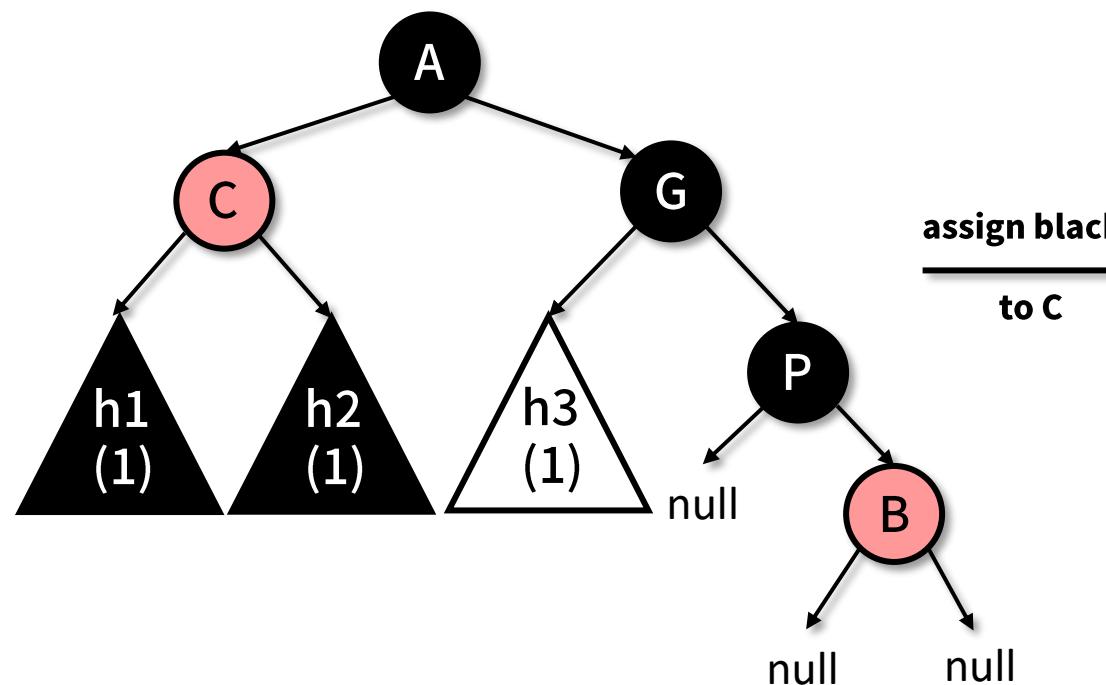
1. **grandparent is black (number of black: 3)**
2. **brother of parent is black**
3. **left child of brother of parent is red.**



# Red-Black Tree Deletion (38)

**case 7.4:**

1. **grandparent is black (number of black: 3)**
2. **brother of parent is black**
3. **left child of brother of parent is red.**

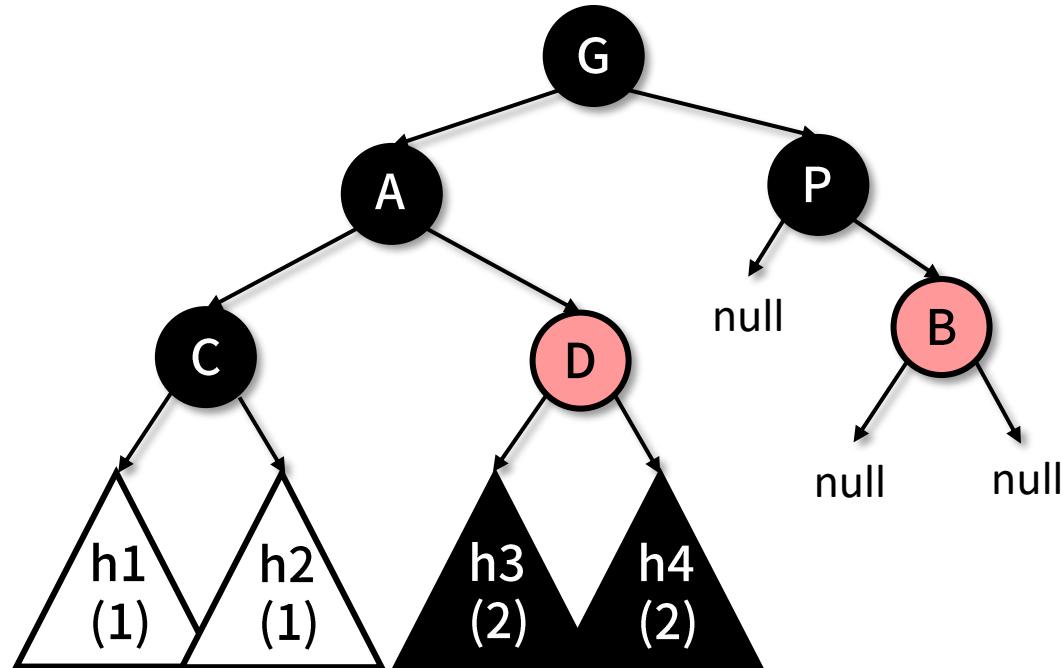


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (39)

**case 7.5:**

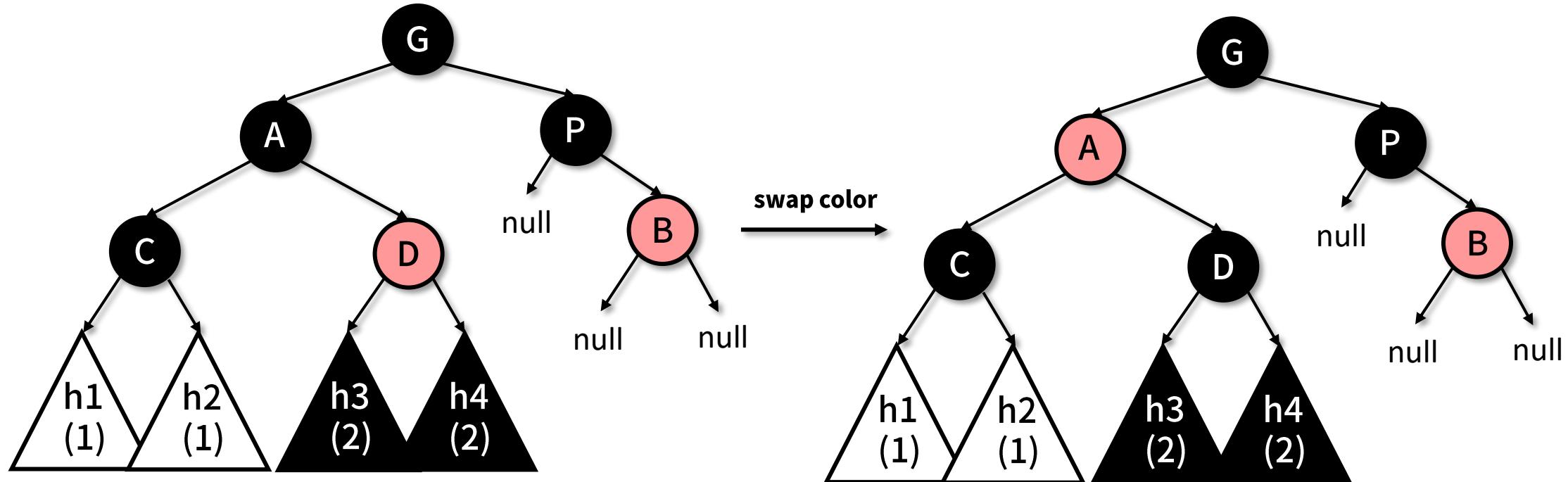
- 1. grandparent is black (number of black: 3)
- 2. brother of parent is black
- 3. left child of brother of parent is black.
- 4. right child of brother of parent is red



# Red-Black Tree Deletion (40)

**case 7.5:**

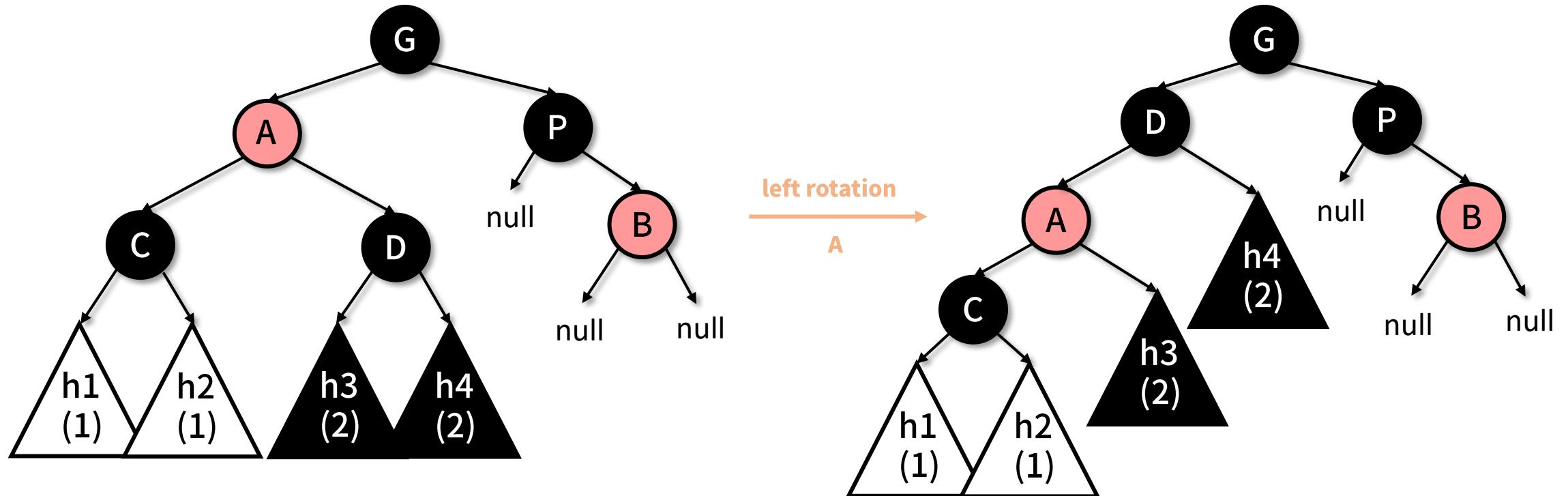
1. grandparent is black (number of black: 3)
2. brother of parent is black
3. left child of brother of parent is black.
4. right child of brother of parent is red



# Red-Black Tree Deletion (41)

**case 7.5:**

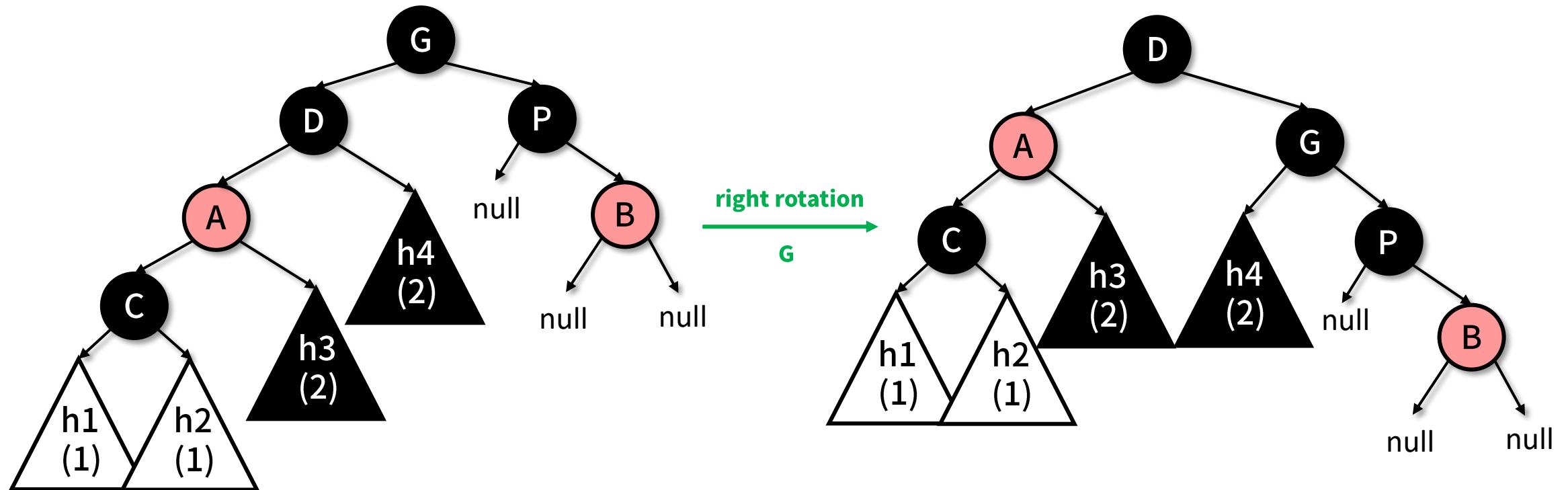
1. grandparent is black (number of black: 3)
2. brother of parent is black
3. left child of brother of parent is black.
4. right child of brother of parent is red



# Red-Black Tree Deletion (42)

**case 7.5:**

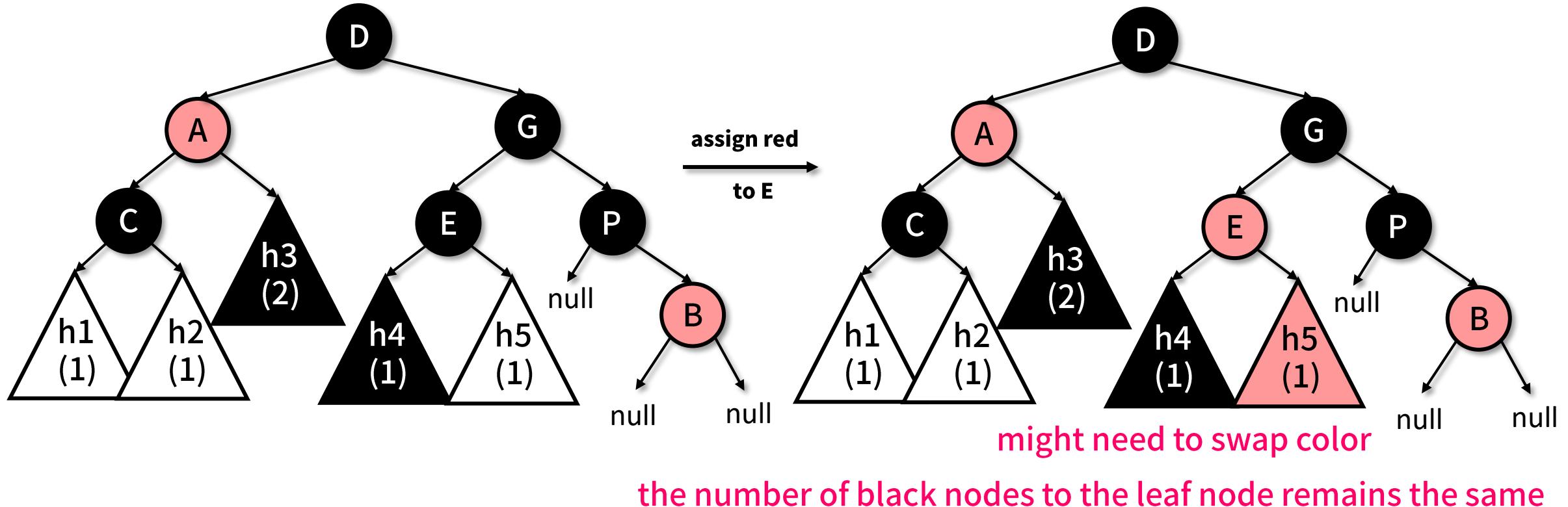
1. grandparent is black (number of black: 3)
2. brother of parent is black
3. left child of brother of parent is black.
4. right child of brother of parent is red



# Red-Black Tree Deletion (43)

**case 7.5:**

1. grandparent is black (number of black: 3)
2. brother of parent is black
3. left child of brother of parent is black.
4. right child of brother of parent is red

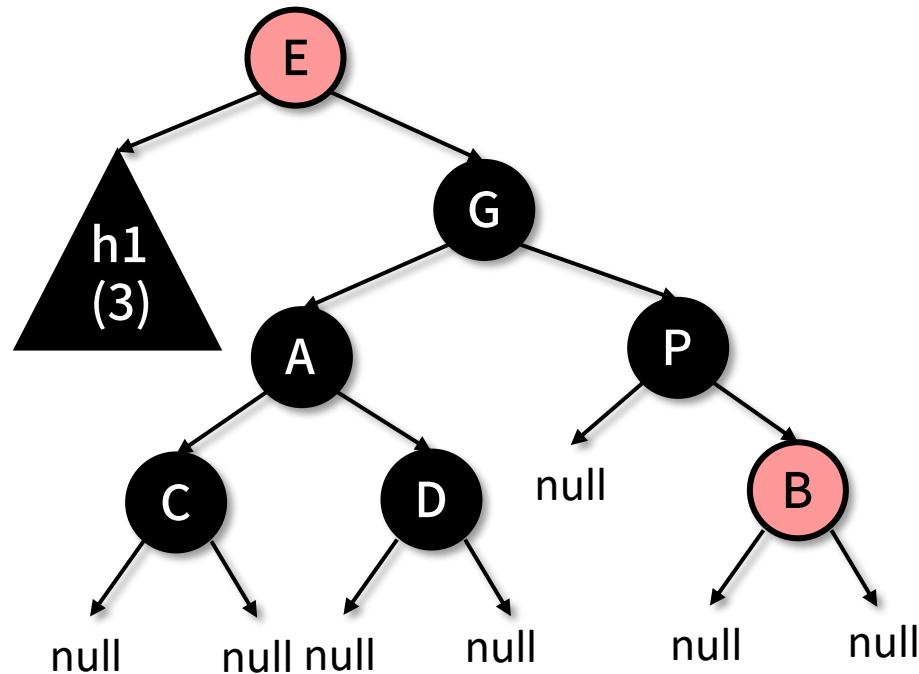


Red-Black Tree

# Deletion (44)

**case 7.6:**

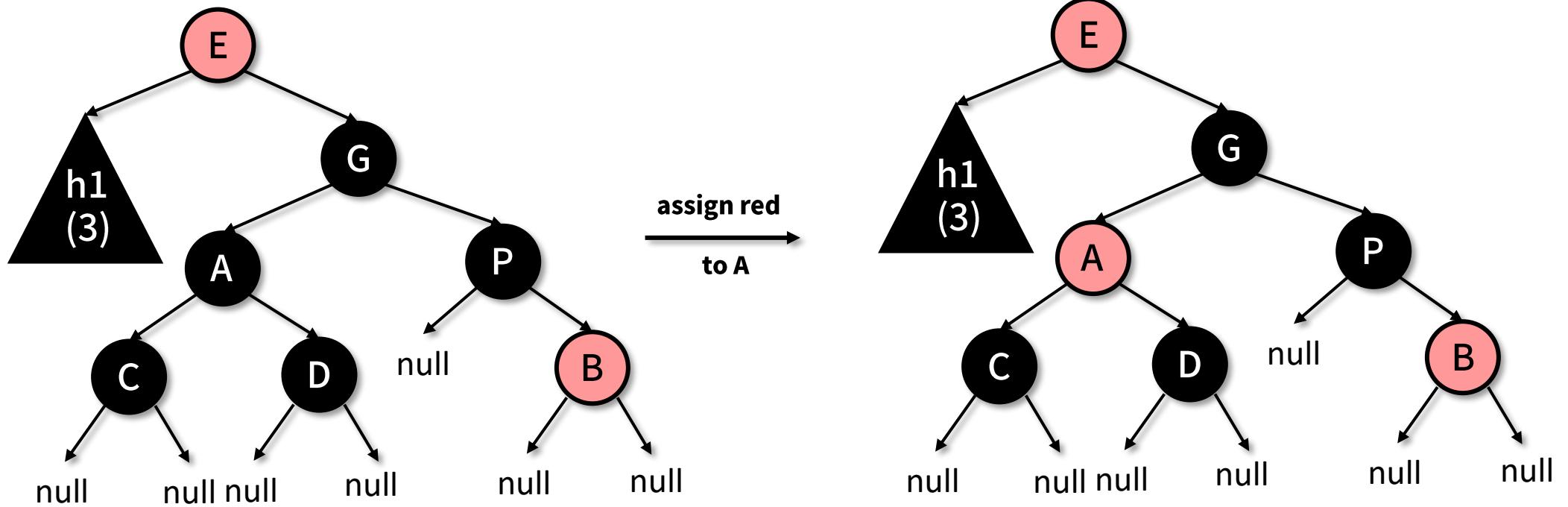
- 1. grandparent is black (number of black: 3)
- 2. brother of parent is black
- 3. children of brother of parent is black
- 4. grand-grand parent is red



# Red-Black Tree Deletion (45)

**case 7.6:**

1. grandparent is black (number of black: 3)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is red

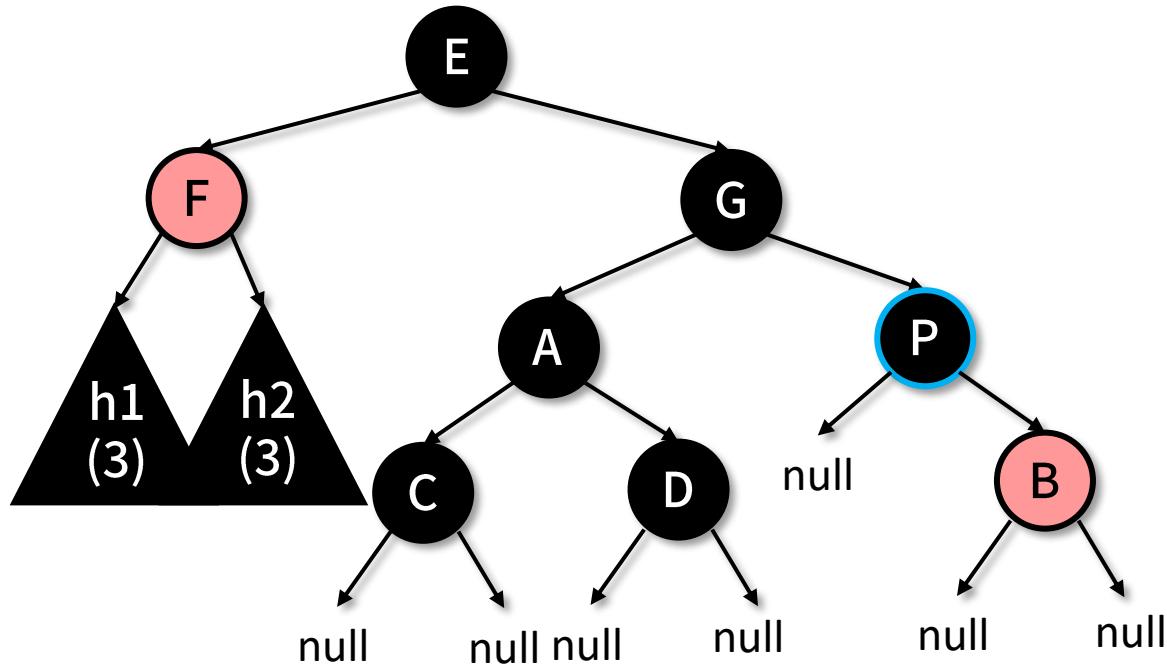


**becomes case 7.2**

# Red-Black Tree Deletion (46)

**case 7.6:**

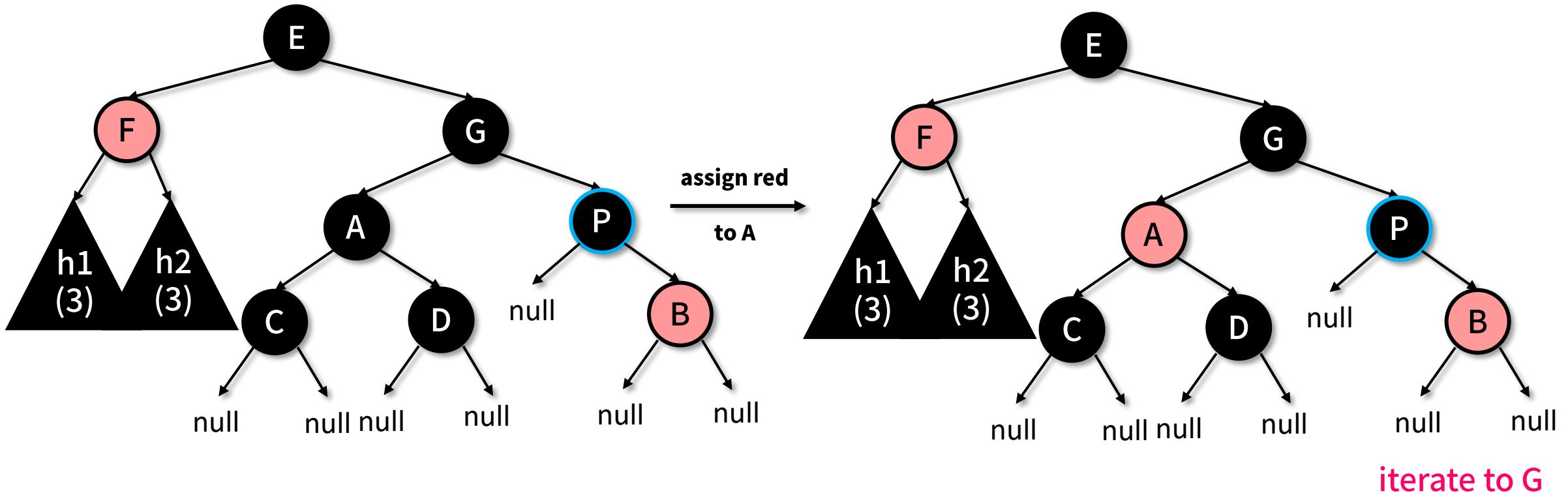
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is red



# Red-Black Tree Deletion (47)

**case 7.6:**

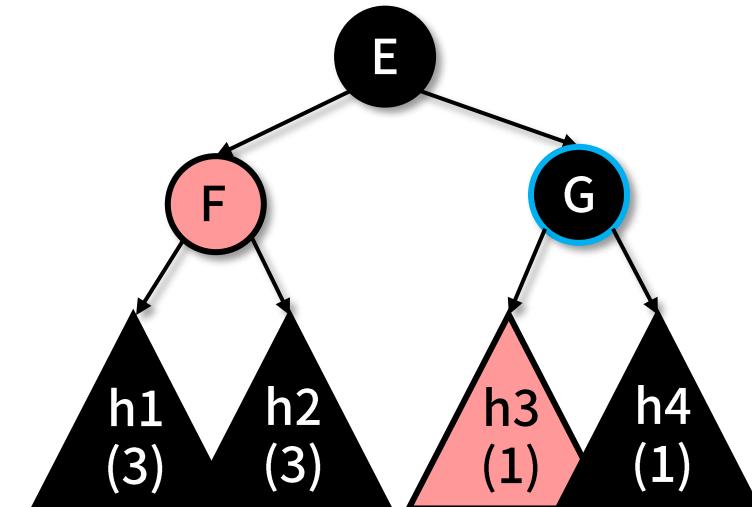
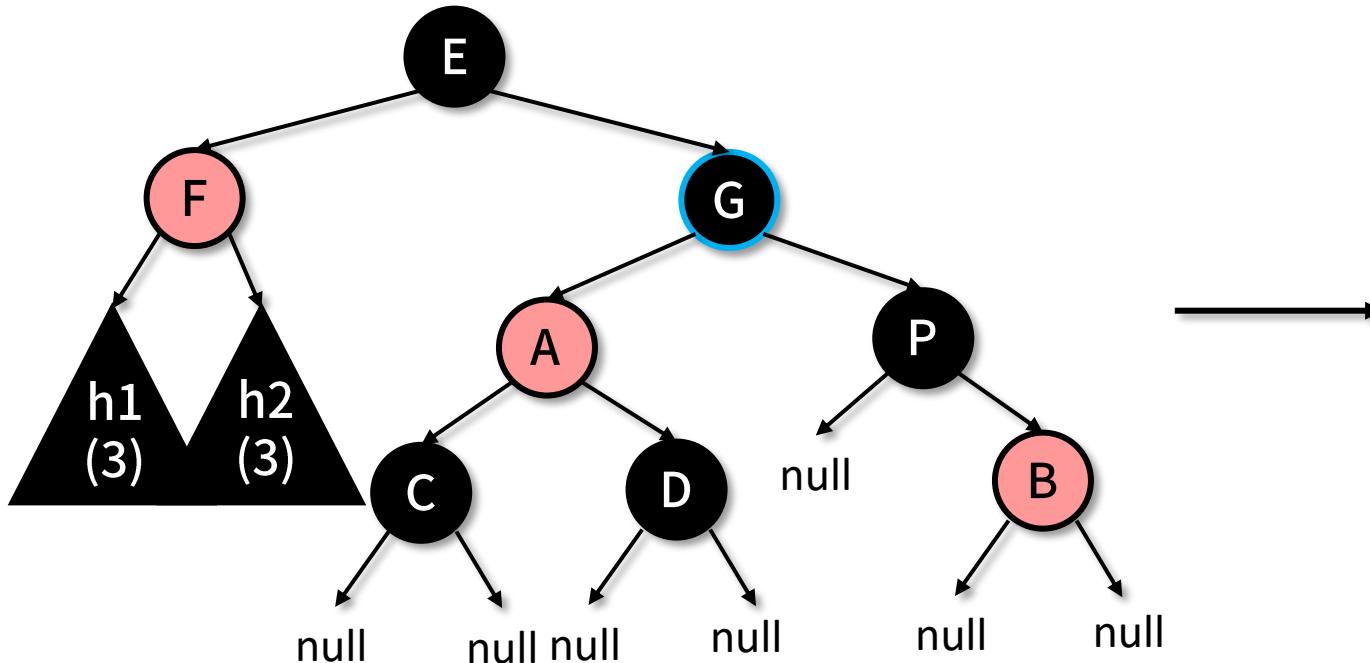
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is red



# Red-Black Tree Deletion (48)

**case 7.6:**

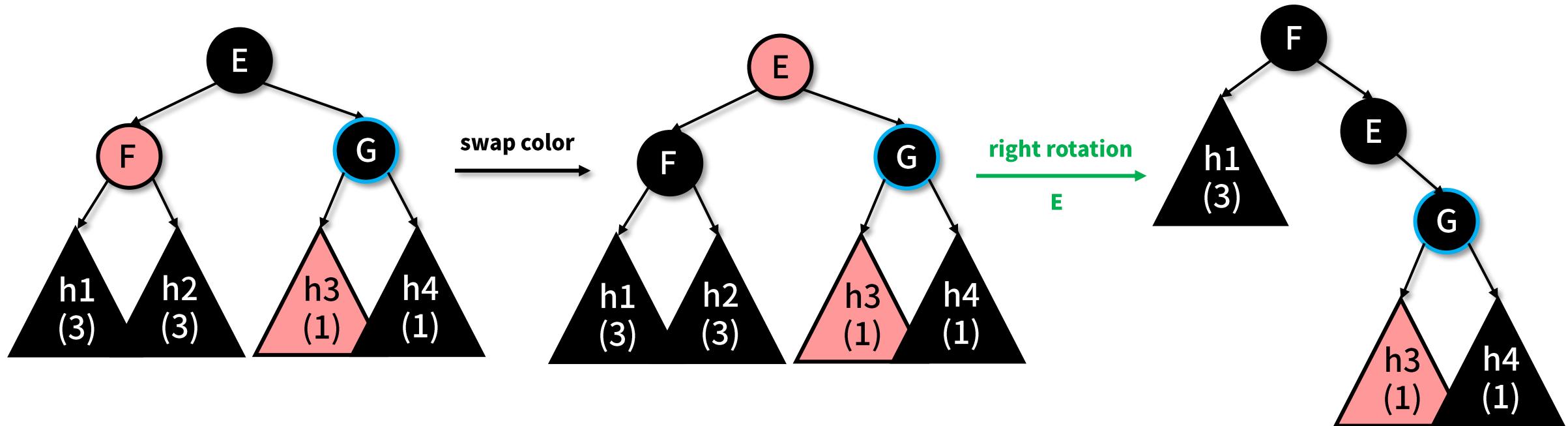
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is red



# Red-Black Tree Deletion (49)

**case 7.6:**

1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is red

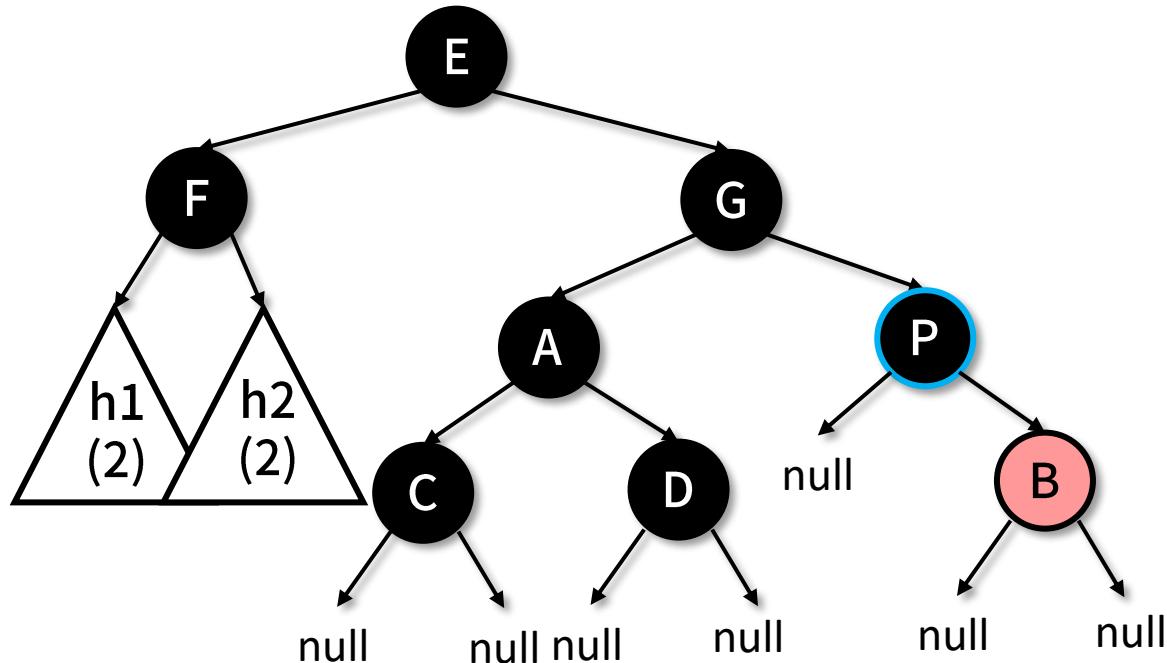


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (50)

**case 7.7:**

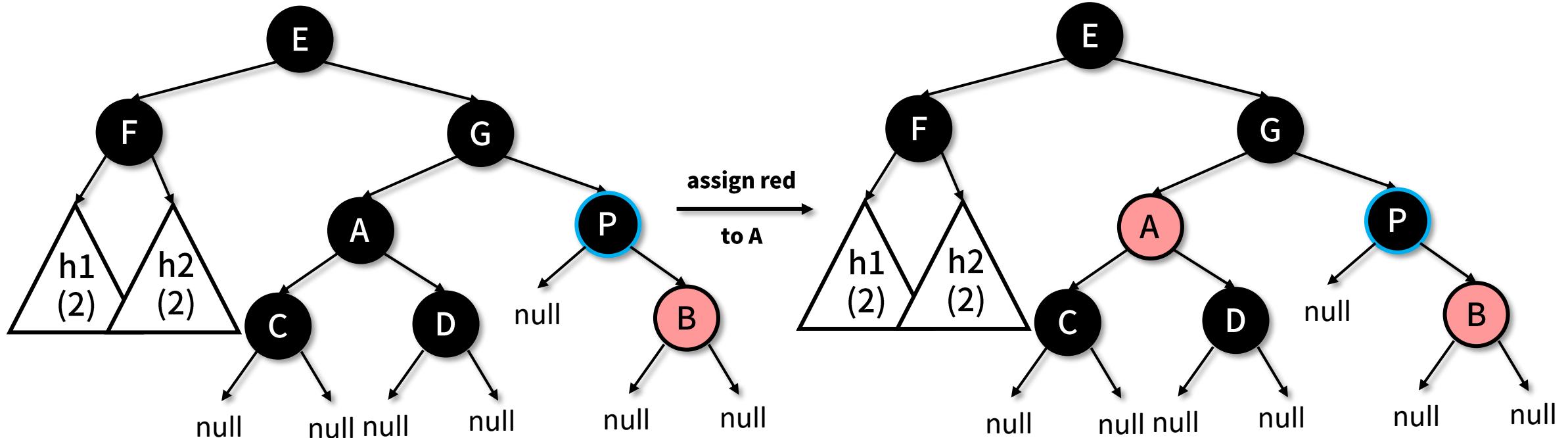
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black



# Red-Black Tree Deletion (51)

**case 7.7:**

1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black



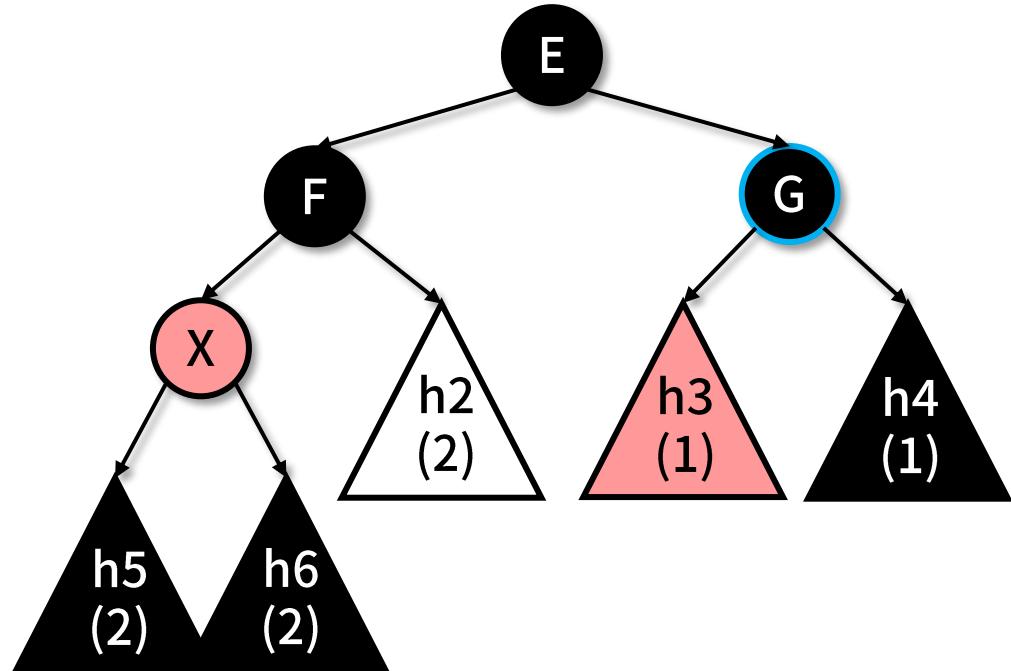
iterate to G

Red-Black Tree

# Deletion (52)

**case 7.7.1:**

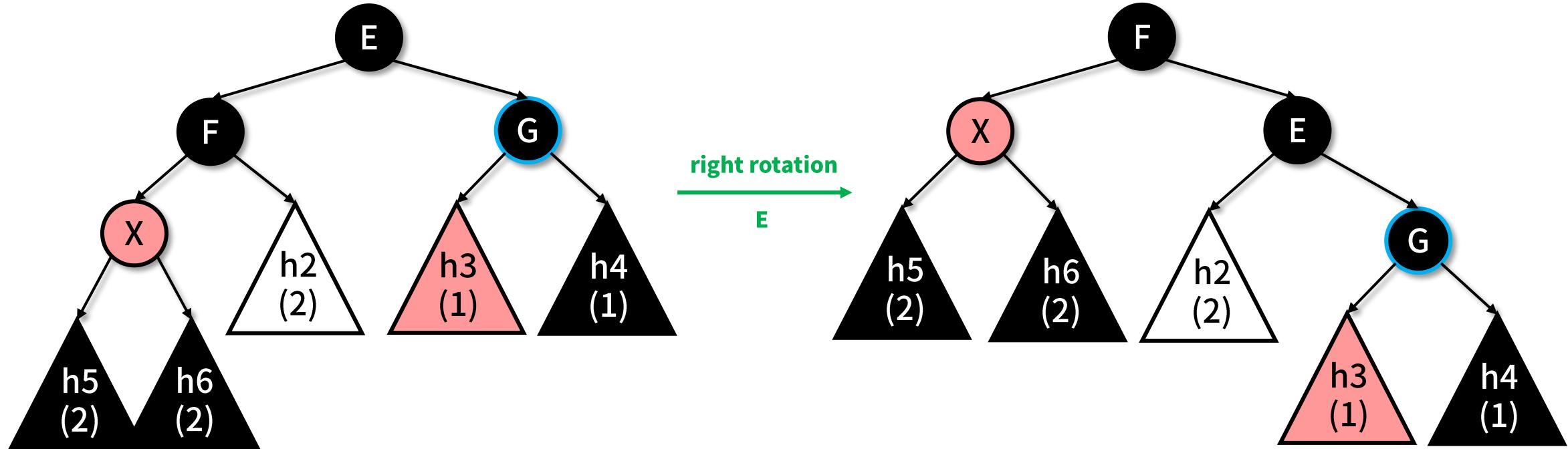
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black
6. left child of brother of grand-parent is red



# Red-Black Tree Deletion (53)

**case 7.7.1:**

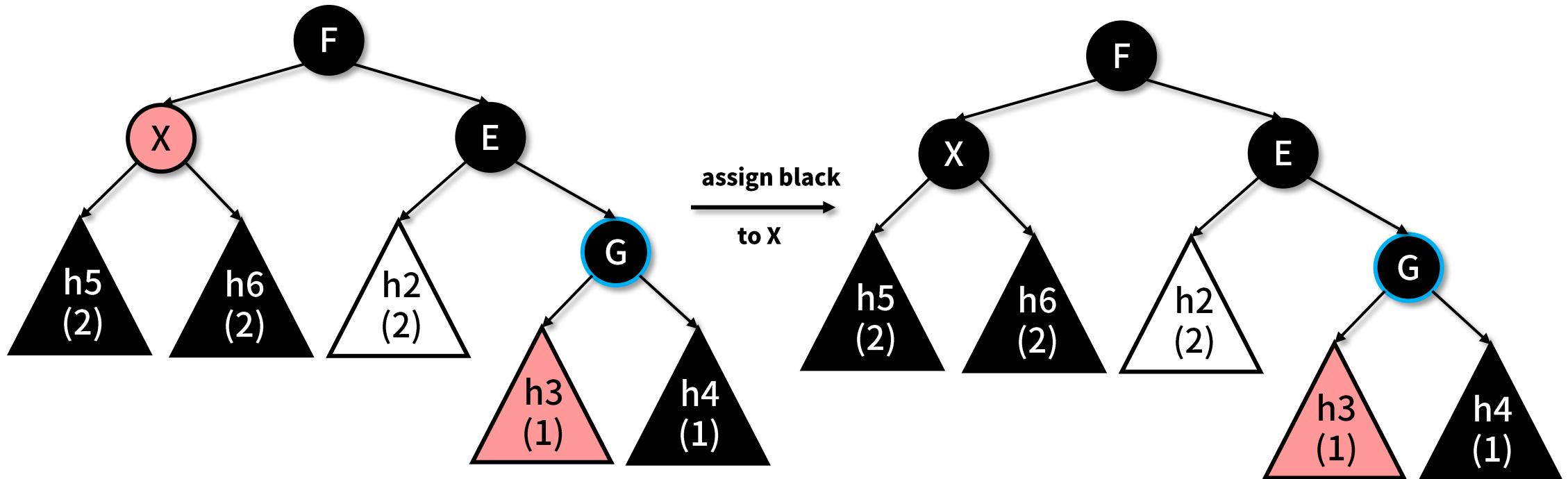
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black
6. left child of brother of grand-parent is red



# Red-Black Tree Deletion (54)

**case 7.7.1:**

- 1. grandparent is black (number of black: 4)
- 2. brother of parent is black
- 3. children of brother of parent is black
- 4. grand-grand parent is black
- 5. left child of grand-grand parent is black
- 6. left child of brother of grand-parent is red

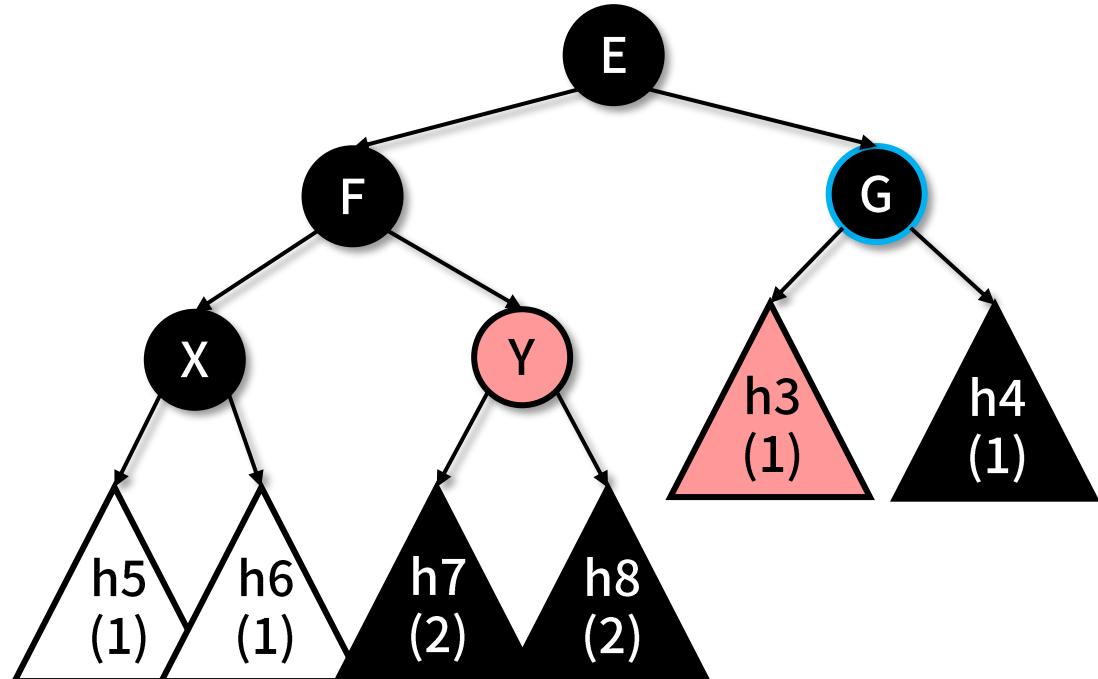


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (55)

**case 7.7.2:**

1. **grandparent is black (number of black: 4)**
2. **brother of parent is black**
3. **children of brother of parent is black**
4. **grand-grand parent is black**
5. **left child of grand-grand parent is black**
6. **left child of brother of grand-parent is black**
7. **right child of brother of grand-parent is red**

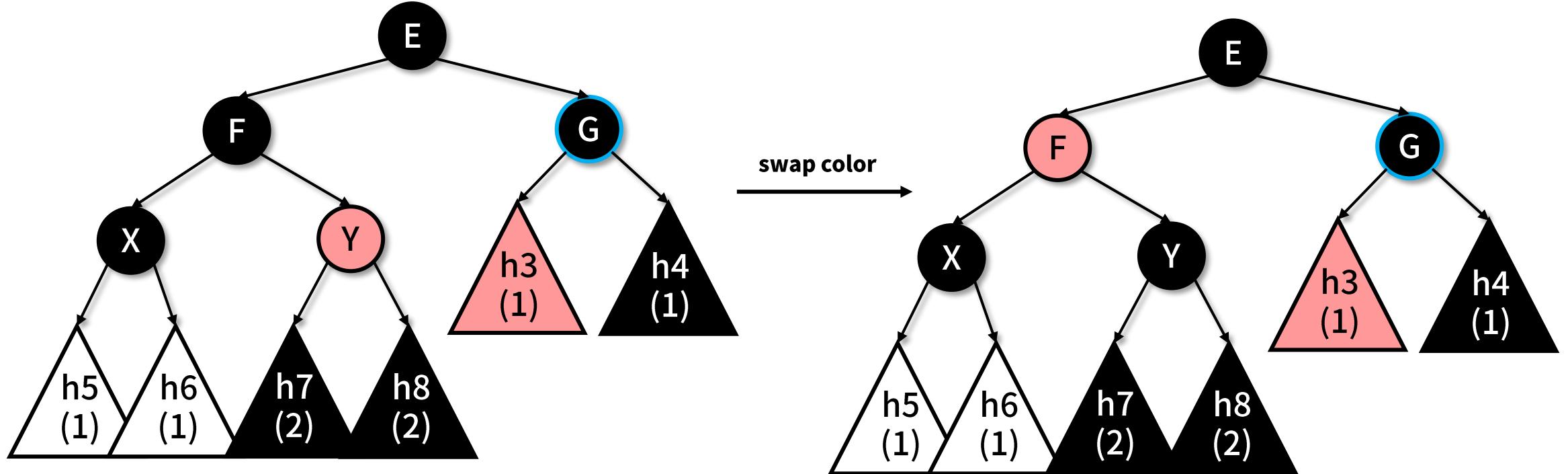


# Red-Black Tree Deletion (56)

**case 7.7.2:**

1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black

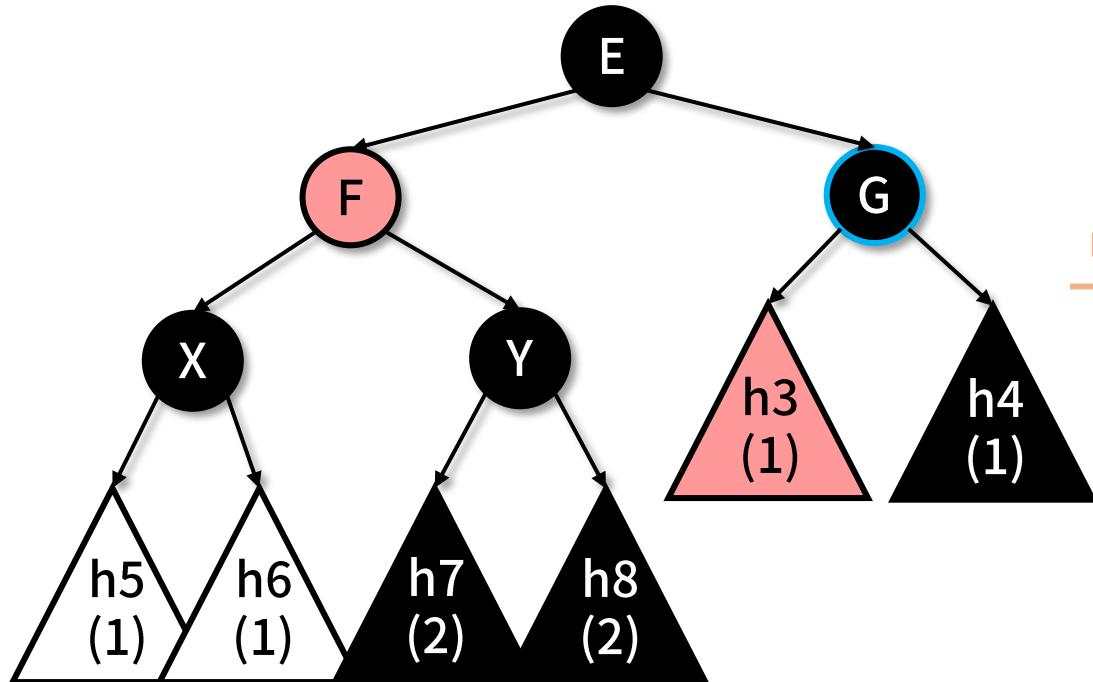
4. grand-grand parent is black
5. left child of grand-grand parent is black
6. left child of brother of grand-parent is black
7. right child of brother of grand-parent is red



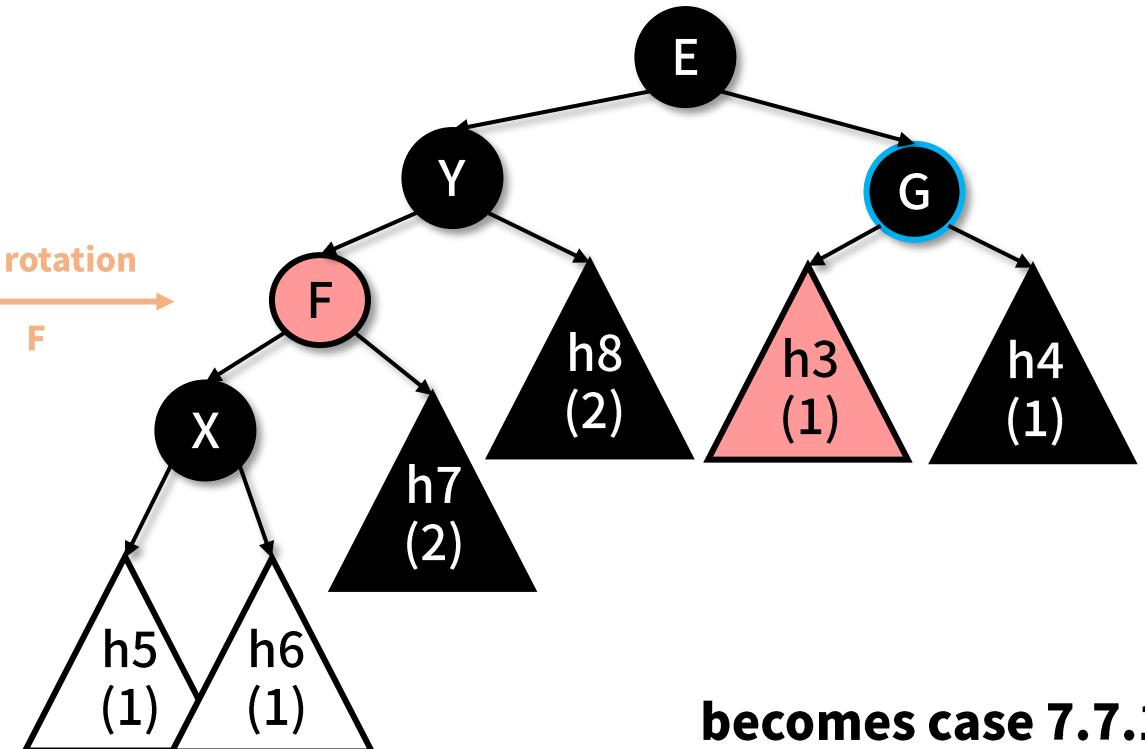
# Red-Black Tree Deletion (57)

**case 7.7.2:**

1. grandparent is black (number of black: 4)
  2. brother of parent is black
  3. children of brother of parent is black
4. grand-grand parent is black
  5. left child of grand-grand parent is black
  6. left child of brother of grand-parent is black
  7. right child of brother of grand-parent is red



Left rotation  
F



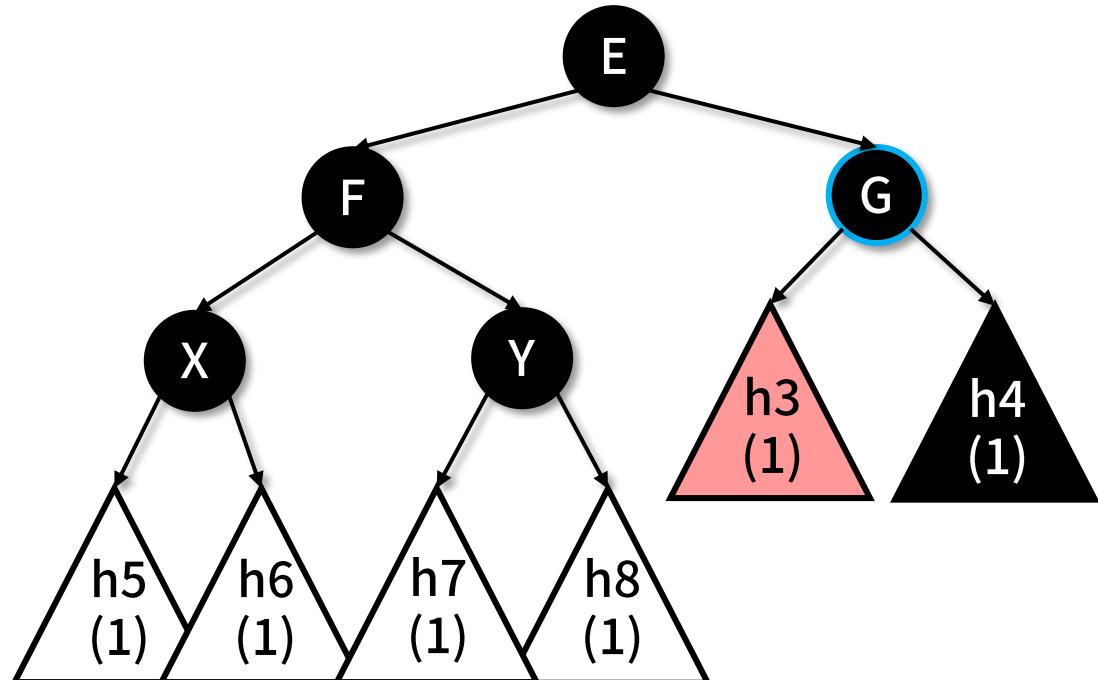
**becomes case 7.7.1**

Red-Black Tree

# Deletion (58)

**case 7.7.3:**

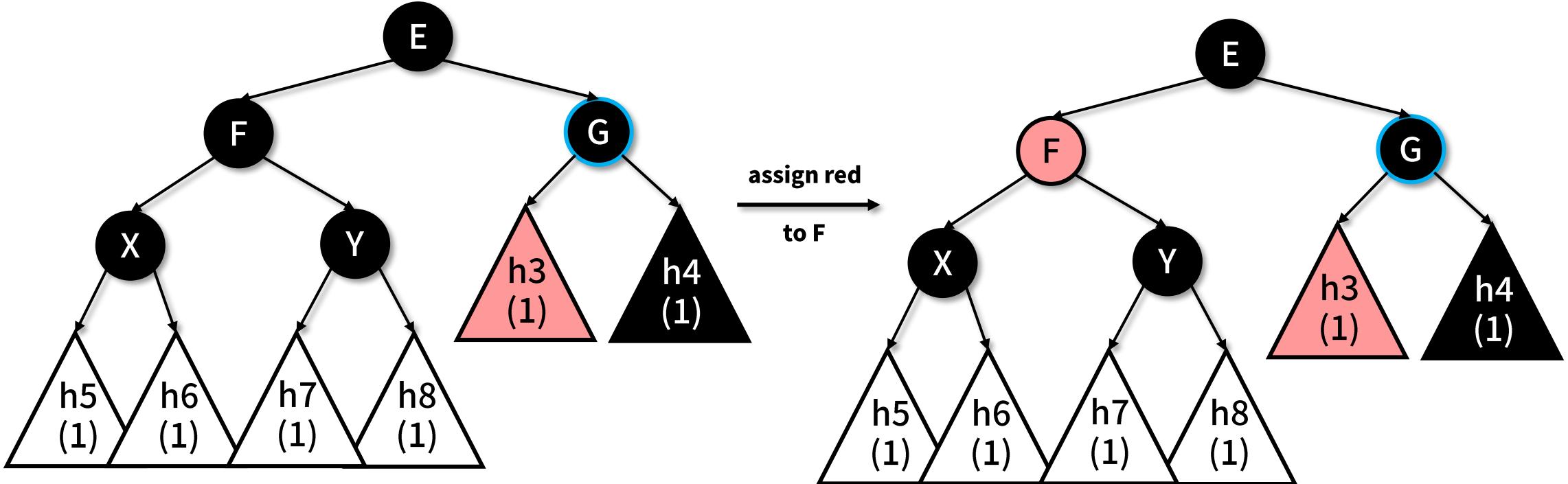
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black
6. children of brother of grand-parent are black



# Red-Black Tree Deletion (59)

**case 7.7.3:**

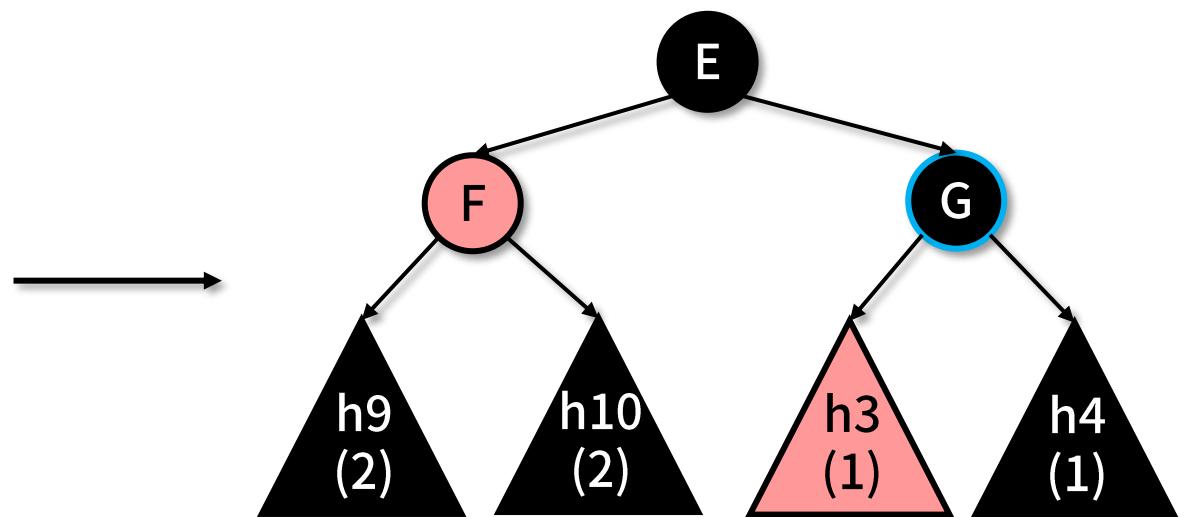
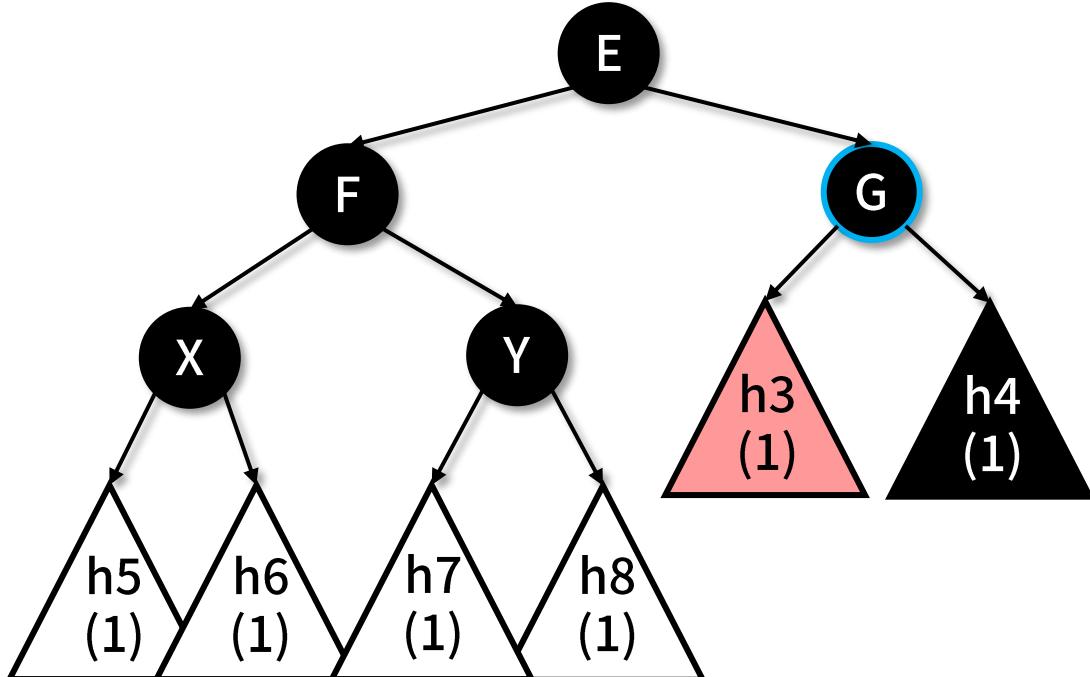
1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black
6. children of brother of grand-parent are black



# Red-Black Tree Deletion (60)

**case 7.7.3:**

1. grandparent is black (number of black: 4)
2. brother of parent is black
3. children of brother of parent is black
4. grand-grand parent is black
5. left child of grand-grand parent is black
6. children of brother of grand-parent are black



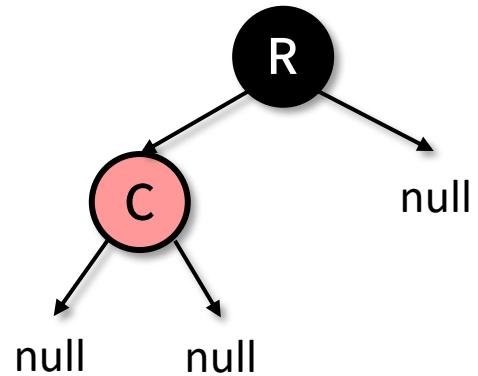
becomes 7.7  
continue iteration until reaching  
root node, 7.6, 7.7.1 or 7.7.2.

Red-Black Tree

# Deletion (61)

**case 8:**

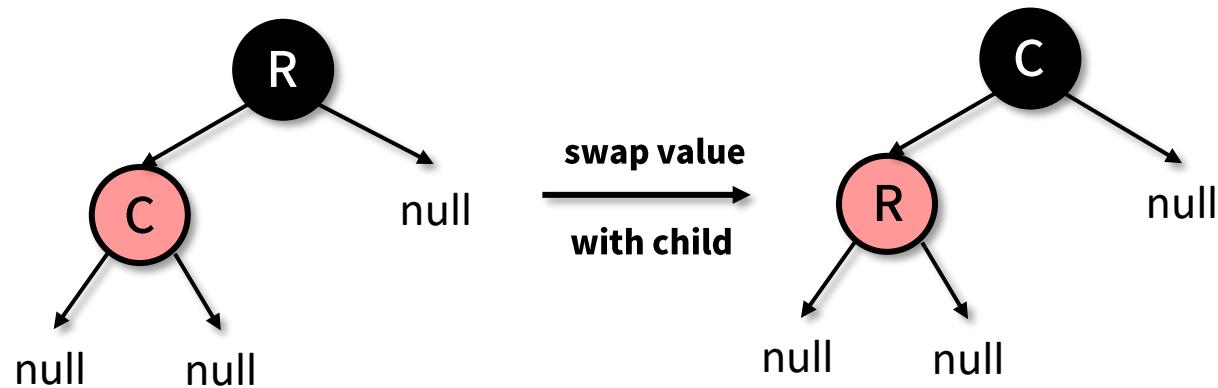
- 1. remove black node with one red child leaf node**



# Red-Black Tree Deletion (62)

**case 8:**

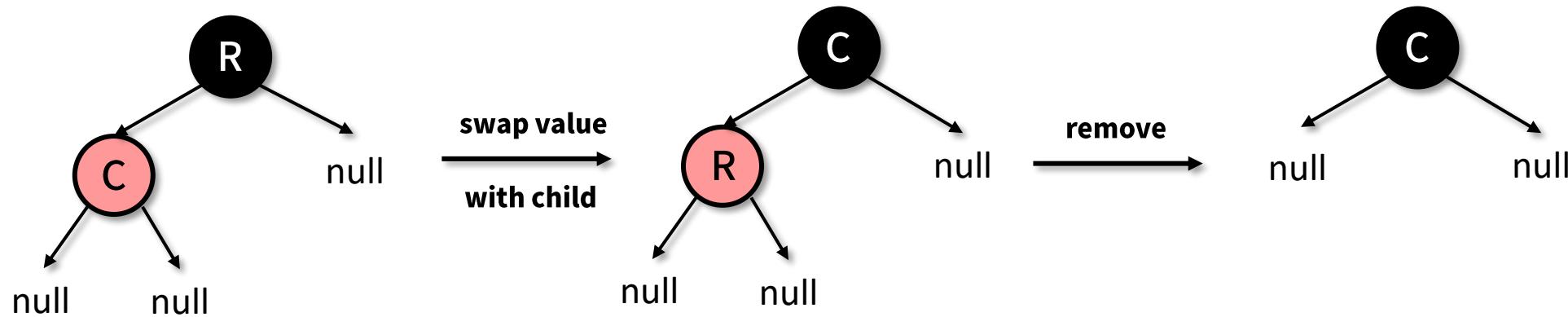
**1. remove black node with one red child leaf node**



# Red-Black Tree Deletion (63)

**case 8:**

**1. remove black node with one red child leaf node**

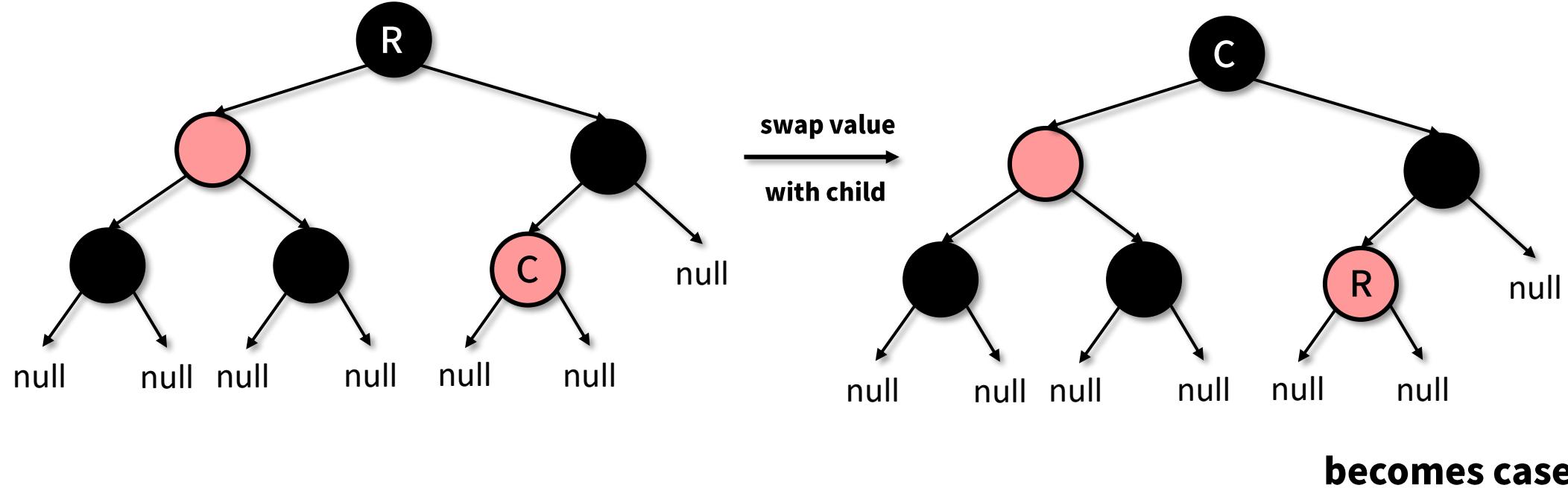


the number of black nodes to the leaf node remains the same

# Red-Black Tree Deletion (64)

**case 9:**

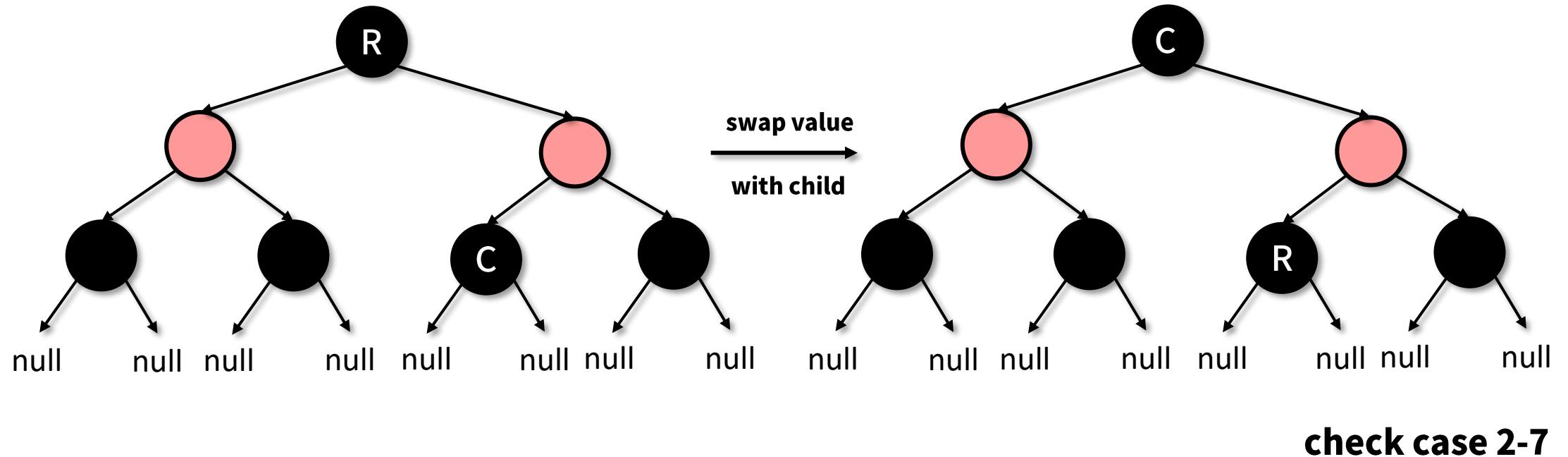
1. remove black node with two children
2. the minimum node in right subtree is red



# Red-Black Tree Deletion (65)

**case 9:**

1. remove black node with two children
2. the minimum node in right subtree is black



# Binary Search Tree Summary

## Worst Time Complexity

	Binary Search Tree	AVL Tree	Red-Black Tree
Search	$O(N)$	$O(\log N)$	$O(\log N)$ *
Insert	$O(N)$	$O(\log N)$ (2 rotations)	$O(\log N)$ (2 rotations)
Delete	$O(N)$	$O(\log N)$ (multiple)	$O(\log N)$ (3 rotations)

\* 2 times longer than AVL Tree

# Binary Search Tree Summary

## Average Time Complexity

	Binary Search Tree	AVL Tree	Red-Black Tree
Search	$O(\log N)$	$O(\log N)$	$O(\log N)$
Insert	$O(\log N)$	$O(\log N)$	$O(\log N)$
Delete	$O(\log N)$	$O(\log N)$	$O(\log N)$

# Mutable Mapping Summary

## Average Time Complexity

	Array	Hash Table	Binary Search Tree
Key	integer	hashable object	comparable object
Ranged search	supported	not supported	supported
Search	$O(1)$	$O(1)$	$O(\log N)$
Insert	$O(N)$	$O(1)$	$O(\log N)$
Delete	$O(N)$	$O(1)$	$O(\log N)$

Binary Search Tree

# Practices

- Range Sum of BST (Leetcode Problem 938)
- Unique Binary Search Trees (Leetcode Problem 96)