

**F21MP**  
**Masters Project and Dissertation**

**Script execution made simple with**  
**SML/NJ**

**Due on Thursday 04 May 2023**

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## Abstract

“Standard ML (SML) is a safe, modular, strict, functional, polymorphic programming language with compile-time type checking and type inference, garbage collection, exception handling, immutable data types and updatable references, abstract data types, and parametric modules” [2]. Developers of compilers, programmers, and others collaborating with theorem provers have used it to create formal proofs [1]. A free and open-source compiler called SML/NJ offers an SML programming environment. While it does not evaluate, SML/NJ processes the instructions in a manner like the well-known REPL (Read-Eval-Print-Loop) approach. It reads, parses, does type checks, compiles, links, runs, and modifies global environments. It loops back through the full cycle until all the instructions have been carried out. The term "script" in this article refers to a file containing source code which are also executable files. SML/NJ doesn't have the capability to execute a script. *smlnj-script* originally developed by project supervisor Dr Joe Wells, gave the capability to build SML scripts in a file and run like regular operating system files. Additional features of *smlnj-script* include support for Unicode and regular expressions. *smlnj-script* is a slightly modified version of SML/NJ which gives the capability to execute SML programs as a script. *smlnj-script* is written as a layer on top of SML/NJ. This dissertation aims to add and merge all *smlnj-script* features into SML/NJ and eliminate the need to maintain *smlnj-script* as a separate implementation. Students, programmers, and developers will benefit from the ability to write SML source code in a script and have them run like Python or Perl.

## Acknowledgements

I would like to thank my supervisor Dr Joe Wells for his support, guidance and knowledge sharing on this topic.

## List of Abbreviations

SML	Standard ML.
SML/NJ	Standard ML of New Jersey.
CM	Compilation and Library Manager (a component of SML/NJ).
JSON	JavaScript Object Notation.
REPL	Read-Eval-Print Loop (In SML world REPL is used to denote read, parse, type check, compile, link and execute.)
HWU	Heriot Watt University
SME	Subject Matter Expert
CI/CD	Continuous Integration / Continuous Delivery

## Table of Contents

<b>Abstract.....</b>	<b>3</b>
<b>Acknowledgements.....</b>	<b>4</b>
<b>List of Abbreviations .....</b>	<b>5</b>
<b>Table of Contents .....</b>	<b>6</b>
<b>1. Introduction.....</b>	<b>8</b>
1.1 Basic Concepts.....	8
1.2 Aim .....	8
1.3 Objectives .....	9
1.4 Stakeholder .....	9
<b>2. Literature Review .....</b>	<b>10</b>
2.1 Standard ML (SML ‘90 & ‘97).....	10
2.2 Implementations of SML .....	10
2.3 SML/NJ.....	10
2.4 Programming process – Compile, Link and Execute.....	12
2.5 SML/NJ Compilation and Execution.....	13
2.6 <i>smlnj-script</i> .....	14
2.7 What happens when <i>smlnj-script</i> gets loaded? .....	15
2.8 Key features of <i>smlnj-script</i> .....	15
2.9 How to run a script?.....	21
2.10 How to run <i>smlnj-script</i> ? .....	22
2.11 How to work with SML/NJ?.....	22
2.12 What happens when SML/NJ is loaded? .....	23
2.13 What happens when a script is executed?.....	25
2.14 Areas to be investigated in SML/NJ .....	26
<b>3. Requirements Analysis.....</b>	<b>28</b>
3.1 Functional Requirements .....	28
3.2 Non-Functional Requirements .....	29
<b>4. Methodology.....</b>	<b>30</b>
<b>5. Professional, Legal, Ethical, and Social issues .....</b>	<b>31</b>
5.1 Professional Issues .....	31
5.2 Legal Issues.....	31
5.3 Ethical Issues .....	31
5.4 Social Issues.....	31
<b>6. Project Risk Assessment.....</b>	<b>32</b>
<b>7. Design and Implementation.....</b>	<b>33</b>
7.1 Pre-requisite .....	33

7.2	Software and Hardware used .....	33
7.3	Where to be updated in SML/NJ .....	34
7.4	Change details .....	34
7.4.1	File cm-boot.sml .....	34
7.4.2	File boot-env-fn.sml.....	39
7.4.3	Files interact.sig & interact.sml .....	41
7.4.4	File backend.sig .....	43
7.4.5	File backend-fn.sml.....	43
7.4.6	File mutecompiler.sig .....	43
7.4.7	File mutecompiler.sml .....	44
7.4.8	Files INDEX, MAP and core.cm .....	47
7.5	Writing a script .....	48
7.6	Executing a script.....	49
7.7	SML/NJ Implementation and 64-bit Support .....	49
7.8	GitHub Codebase .....	49
7.9	GitHub Test.....	50
7.10	Merging with SML/NJ .....	51
7.11	Issues and Interaction with David MacQueen & John Reppy .....	51
<b>8.</b>	<b><i>Evaluation and Results</i></b> .....	<b>52</b>
8.1	Before the change .....	52
8.2	After the change .....	53
8.2.1	Execute as a script.....	53
8.2.2	Silencing compiler messages .....	55
8.2.3	Un-silencing compiler messages.....	56
8.2.4	Printing silenced compiler messages .....	57
8.2.5	Restoring printing limits .....	61
8.3	Sample Scripts .....	62
8.4	GitHub Test Results .....	64
<b>9.</b>	<b><i>Conclusion and Future Work</i></b> .....	<b>65</b>
9.1	Conclusion .....	65
9.2	Future Work .....	65
9.3	Lessons Learnt .....	65
<b>10.</b>	<b><i>Copyright and Licensing</i></b> .....	<b>67</b>
	<b><i>References</i></b> .....	<b>68</b>
	<b><i>Appendix A</i></b> .....	<b>70</b>
	GitHub Test Screen prints.....	70

# 1. Introduction

## 1.1 Basic Concepts

This paragraph will aid readers in understanding the purpose. Most programming languages break down the execution of a program's source code into stages like compilation, linking, and execution. In contrast to certain other programming languages like Python and Perl, where the programme can be run straight from the source code, the user must conduct or provide instructions for each phase in the former languages. Additionally, in some programming languages, running a programme on the command line requires passing the instructions through a wrapper script. One of the main goals of this project is to execute an SML programme like a script.

## 1.2 Aim

SML is a procedural language that supports higher-order functions and abstraction. Developers of compilers, programmers, and others collaborating with theorem provers have used SML to create formal proofs [1]. Though SML/NJ (a free and open-source compiler) offers an SML programming environment, it doesn't evaluate and processes the instructions in a manner like the well-known REPL (Read-Eval-Print-Loop) approach. It reads, parses, does type checks, compiles, links, runs, and modifies global environments. It loops back through the full cycle until all the instructions have been carried out. SML/NJ doesn't have the capability to execute a script (file containing source code which are also executable files). *smlnj-script* originally developed by project supervisor Dr Joe Wells, gave the capability to build SML scripts in a file and run like regular operating system files. Additional features of *smlnj-scripts* include support for Unicode and regular expressions. *smlnj-script* is a slightly modified version of SML/NJ which gives the capability to execute SML source code as a script. *smlnj-script* is written as a layer on top of SML/NJ. This dissertation aims to add and merge all *smlnj-script* features into SML/NJ and eliminate the need to maintain *smlnj-script* as a separate implementation. Students, programmers, and developers will benefit from the ability to write SML source code in a script and have them run like Python or Perl.

The project's goal is to make it easier for college students to complete their class assignments so they can create and execute SML source code as a script.



Additionally, the project seeks to cooperate with SML/NJ implementors to include these new capabilities and have SML/NJ published in the future with these features.

### 1.3 Objectives

The SML-NJ code will be examined for suitable places where the *smlnj-script* features can be added or merged. Before merging with SML/NJ, *smlnj-script* features will be evaluated for compatibility with the most recent version of SML/NJ and will be modified to meet new standards. SML/NJ implementors will be given information on the new features, and we'll work to have them included in a future version.

### 1.4 Stakeholder

Below are the stakeholders of this work.

Stakeholders	Interests
<i>Primary:</i>	<i>Have direct relationship to the project</i>
Dayanandan Natarajan (myself)	Student working on this project.
Dr Joe Wells	Supervisor and original author of <i>smlnj-script</i> .
SML/NJ Implementors	They are the maintainers of SML/NJ and publish the SML/NJ upgrades.
<i>Secondary:</i>	<i>Have indirect relationship to the project</i>
Students	Beneficiaries of the resultant product.
Programmers	Beneficiaries of the resultant product.
Developers	Beneficiaries of the resultant product.

## 2. Literature Review

### 2.1 Standard ML (SML “90 & “97)

Standard ML “90 is a procedural language that supports higher-order functions and abstraction very strongly because it is highly typed. SML maintains fairly like the  $\lambda$ -calculus in terms of syntax and semantics. The language has undergone a minor change and simplification to become Standard ML “97. A new SML base library was included with the language definition to allow a variety of systems and applications programming. It has a large selection of pre-defined modules that include fundamental types, input/output capabilities, and interfaces for portable operating system interaction [6].

### 2.2 Implementations of SML

There are multiple variants of SML implementations available, from interpreters to byte-code compilers, to incremental compilers, to whole-program compilers. The two main capabilities which differ them are REPL and FFI. An interactive top-level read-eval-print-loop (REPL) capability enables users to enter specific top-level SML declarations, which are subsequently parsed, type verified, executed (either as machine code or as interpreted code), and the results are displayed. FFI (foreign function interface) enables calling of SML from other languages as well as functions in C libraries and other languages [7].

Some of the main implementations are,

- MLton
- SML/NJ
- Poly/ML
- SML#

For our dissertation we will be seeing more of SML/NJ.

### 2.3 SML/NJ

SML/NJ is an open-source compiler and programming environment for SML, which was developed at Bell Labs. It was initially created in collaboration between Princeton University and Bell Laboratories, and later joined by Yale University and

AT&T. The most recent version of SML/NJ implements the Standard ML language's SML '97 revision, including the standard basis library. The SML/NJ source code is provided without charge or warranty. In accordance with the relevant licence and copyright notice, anybody may use, copy, modify, and distribute the software [8].

Except for the runtime system, which is written in C, SML/NJ is primarily written in Standard ML. Numerous big systems have been implemented using it, particularly in the areas of applied logic and verification, programme analysis, and advanced compilers [8].

Some of the key features of the SML/NJ system are below [8],

- a. An interactive top level based on incremental compilation is offered by SML/NJ.
- b. SML/NJ uses Compilation Manager (from Matthias Blume), to simplify the development of large software projects.
- c. A range of general-purpose data structures, algorithms, and utilities are offered by the SML/NJ library (such as finite sets and maps, regular expressions, pretty-printing). The other main SML implementations also use SML/NJ library and are periodically resynchronized. SML/NJ library not necessarily denoted as SML/NJ.
- d. Higher-order functors, OR-patterns (logical "or" operations), first-class continuations, and other helpful features are added to the SML '97 language by SML/NJ.
  - i. Higher-order *functor* is a module that take one or more other modules as a parameter [18][19].
  - ii. First-class continuations are continuations used as values. Continuations can be the value of a variable, an argument for a function call, or the return value of a function. It can be considered as a function since it takes a value and produces a value [17].
- e. A simple quote/unquote mechanism offers support for manipulating "object languages". The command arguments contain the protected characters get passed along [20].

## 2.4 Programming process – Compile, Link and Execute

A program written in any programming language needs to be compiled and linked before execution. Programming languages vary in this process; they can be either compiled or interpreted based on how they are compiled, linked, and executed [15][16].

### Compiled Programming Languages

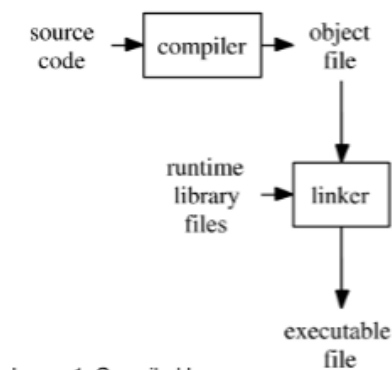


Image 1. Compiled Languages

Image1: Compiled Languages [16].

### Interpreted Programming Languages

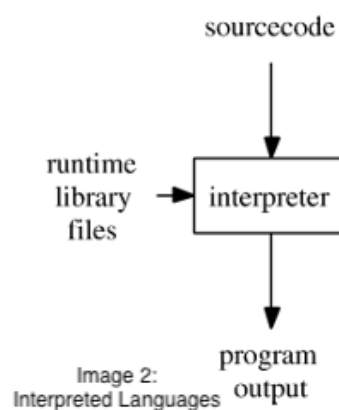


Image 2:  
Interpreted Languages

Image2: Interpreted Languages (ex: Python/Perl) [16].

Programming language like Java doesn't fit either in compiled or interpreted language, they have their own compiler which converts the source code into byte code which can be run on any machine using their own interpreter.

## 2.5 SML/NJ Compilation and Execution

When SML/NJ is started over the command line, it gets loaded and prompts to enter a top-level declaration through keyboard. When a top-level declaration is given to the prompt with a termination character semicolon “;”, SML/NJ understands the end of declaration and starts to process the declaration. “SML/NJ performs the following actions: elaborate (performs type checking and static analyses), compile (obtain executable machine code), execute, and finally print the result of this declaration” [22].

Rather than typing the program in top-level declarations it can be put together in a single file and loaded into SML/NJ system. SML/NJ processes the instructions using a method like the well-known REPL (Read-Eval-Print-Loop) methodology, though it does not evaluate. It reads, parses, does type checks, compiles, links, runs, and modifies global environments. It repeats the entire cycle in loops until all the instructions are processed. Once all the declarations in the program have been compiled and executed. Here comes an optional handy feature dumping, which creates a completely pre-compiled version of the program that can start-up faster.

At any given point of time, the snapshot of application’s memory called as heap dump contains information in binary format, they contain information like objects in memory, their values, their size, and reference to other objects [23]. SML/NJ gives the capability to ship a program heap; upon loading a file in to SML/NJ compiler, it gets parsed, compiled, and executed. Dumping saves the system heap into a file. This file, which is a virtual memory image and not an executable, can be exported [22]. The exported heap can be executed by feeding the argument of the heap file to SML over a shell script.

Compilation and Library Manager (CM) is used for managing the compilation of SML programs and its required libraries. SML programs can be organized into separate modules with its own source file with CM. It also allows each module to be compiled independently, and compiled code to be linked together to form a complete program. Each module is compiled and linked in the proper order thanks to CM's management of inter module dependencies. This helps to prevent problems with circular dependencies and guarantees that the programme is successfully compiled and linked.

A library manager is another feature of the CM system that enables the management and usage of SML libraries in SML programmes. An assortment of SML modules that can be utilised in various SML programmes make up a library. With a separate source file for each module, libraries can be provided using the library manager in the same way that modules are. The library manager also handles the dependencies between modules in the library, ensuring that they are compiled and linked in the correct order. When a library is used in an SML program, the library manager automatically handles the loading and linking of the required library modules.

Overall, the CM system provides a powerful and flexible way to manage the compilation and linking of SML programs and libraries. By organizing programs and libraries into separate modules and managing the dependencies between them, CM helps to ensure that SML programs are correct, reliable, and maintainable [24]. CM gives the capability to load a specific file from the libraries which are precompiled. It can be loaded and added into the memory image whenever required.

To execute a SML program, the end-user needs to have two files, a program, and a wrapper script. The SML program will be run using the instructions in the wrapper script. Sometimes a long command line can be used as an alternative to a wrapper script. For example, to run a program “helloworld.sml” there should be wrapper script “helloworld.sh” available with instructions as below,

```
#!/bin/sh
sml /full/path/of/directories/to/helloworld.sml
```

## 2.6 *smlnj-script*

*smlnj-script* program is an extension to SML/NJ implementation. In SML/NJ, reading, parsing, type checking, compiling, linking, execution, modify global environments and looping happens in separate phases and the program is executed through instructions in a wrapper script. Unix has a special support for script; with a special instruction (program name preceded by a sharp exclamation) inside the script, the path that was initially used when attempting to start the script is passed to the interpreter programme that is given by the loader as an argument, allowing the programme to use the file as input data [14]. *smlnj-script* is an extension to SML/NJ and designed to perform all the operations in a single phase through a single file like

Python or Perl. It also includes run time machinery, with additional utility functions defined and a specialized start-up routine. Here are *smlnj-script* capabilities,

- a. All the features of the SML language (as implemented by SML/NJ).
- b. A portion of SML Basis Library that is provided by the SML/NJ implementation.
- c. SML/NJ library.
- d. New functions added as part of SmlnjScriptUtils *structure*.

## 2.7 What happens when *smlnj-script* gets loaded?

Whenever a user tries to run a script or load the *smlnj-script* directly, the SML/NJ gets preloaded, then it dumps the entire virtual memory image (of everything that's loaded in SML/NJ at that moment) into a heap file, converts the file into assembly code, gets assembled by an assembler into a binary file which can be executed at a later stage. Upon execution it recreates a memory image and starts it up again as *smlnj-script*.

## 2.8 Key features of *smlnj-script*

The *smlnj-script* has the following key features,

- a. Execute as a script:

The very first and most important feature of *smlnj-script* is the ability to run the SML source code as a script. It has the capability to recognise the script as a SML source code and process the instructions in the file by passing them through REPL.

- b. Silencing the Compiler:

Running a script through *smlnj-script* goes thru this redevelopment loop, cycle of read, parse, type check, compile, links, execute and modify global environment. The compiler produces bunch of output messages printed for each operation which is going to confuse the user especially when there is a compilation error. `silenceCompiler` function gives the capability to turns off these compiler messages like print of the name, type, and value of all

identifiers that get added to the top-level environment, autoloading and compilation error messages. This completely silences the compiler and hides type error messages. There is no option to turn off the messages while keeping the error messages [10]. `silenceCompiler` function intercept the compiler messages, accumulates these messages, and stores them in a variable to be printed to the user at a later stage when there is a compilation error or program crashes.

```
val silenceCompiler : unit -> unit =  
  SmlnjScriptUtils.silenceCompiler
```

c. Overloading `toString` function:

SML/NJ has the limitation on overloading `toString` operator especially with complex types like constructors with arguments. To print out a list of integers or real numbers or float data types a custom-built function is needed. The `toString` overloading function needs further improvement and advancements to support string functions.

The `toString` (and `%`) overloaded operator:

```
val % : 'a -> string = ...  
val toString : 'a -> string = ...
```

The `%` and `toString` operators are both overloaded operators to convert many types into strings. They are equivalent to `Int.toString`, `Bool.toString`, `Real.toString`, etc. There is a room for improvement to add support for more types, provided the type is “atomic” with `SmlnjScriptUtils.extendToString` [10].

d. The `SmlnjScriptUtils` structure:

The `SmlnjScriptUtils` *structure*, denoted by the short name “U”, is a collection of functions with its own features.



```
structure SmlnjScriptUtils = ...
structure U = SmlnjScriptUtils
```

Some of those functions are,

- i. The `raisePrintingLimitsToMax` function:

```
val raisePrintingLimitsToMax : unit -> unit = ...
```

The function is very useful while debugging and the default limits cause too much of the data structures to be truncated when printed by the compiler [10].

- ii. The `interact` and `continue` functions:

```
val interact : unit -> unit = ...
val continue : unit -> unit = ...
```

The function `interact` will interrupt the script and transfers control to the user who is at the keyboard, and function `continue` resumes execution by returning from a call to `interact` [10].

- iii. The `extendToString` function

```
val extendToString : string -> unit = ...
```

The function `extendToString` extends the top-level overloaded operators `toString` and `%` with an additional case. Supply the name of a function of type `T -> string` where `T` is an “atomic” type. The type `T` is atomic when it is a type name without arguments and there is no definition available in scope that allows the compiler to deduce it is equal to a non-atomic type [10].

iv. The `evalString` and `useString` functions:

```
val evalString : string -> unit = ...  
val useString : string -> unit = ...
```

Both `evalString` and `useString` enter into a loop of parsing the string to find a prefix that is a top-level SML form, type checking and compiling this form, and executing the resulting machine code, and then repeating with the rest of the string. The difference is that `evalString` throws any new top-level definitions away when it is done and returns (only their side-effects persist) while `useString` extends the top-level environment with the new definitions when it returns. In both cases new top-level definitions are available for remaining forms processed during the call to `evalString` or `useString` [10].

v. The `StringSet` and `StringMap` *structures*:

```
structure StringSet : ORD_SET = ...  
structure StringMap : ORD_MAP = ...
```

The *structures* `StringSet` and `StringMap` provide implementations of sets and finite maps (finite maps are sometimes called “dictionaries” or “associative arrays”) for the type “string”. The operations available are defined in the `ORD\_SET` and `ORD\_MAP` *signatures*, which are part of the SML/NJ library [10].

vi. The `q` and `qq` quoting functions:

```
val q = SmlnjScriptUtils.q  
val qq = SmlnjScriptUtils.qq
```

The functions `q` and `qq` are abbreviations for things in `SmlnjScriptUtils`, giving you a compact syntax for writing

strings with stuff spliced into a string template. These operators are intended to be used with SML/NJ's quasiquote syntax to make quoting operators [10]. The quasiquote syntax is also supported by Moscow ML and the ML Kit, but is not supported by some other major SML implementations like MLton and Poly/ML.

To illustrate, here are some expressions that evaluate to true:

```
let val x = 7 in q`The value of x is ^(% x)!` =  
"The value of x is 7!" end  
let val y = "Jack" in qq`My name is ^y.\n` = "My  
name is Jack.\n" end
```

In the above code, the use of “% x” is a use of toString feature. Both “q” and “%” are two independent features, combined together they give the capability to print a string in middle of a string or include value of a variable. The difference between q and qq is that qq interprets SML-style escape syntax in the literal text while q does not. For example, this evaluates to true:

```
q`This is not a newline: \n` = "This is not a  
newline: \n"
```

The names q and qq are loosely inspired by Perl's q and qq operators [10].

e. Dumping and Creating executable:

SML/NJ gives the ability to dump and load back contents of virtual memory in the form of heap image into a file. When a SML is started again (later), by passing a special command line instructions through another script, the contents of the memory can be discarded and loaded with the contents of the heap image. *smlnj-script* uses SML/NJ's built-in utility functions *exportML* to save the heap image (current state of SML/NJ) in a file [22], *exportFN* to save the current state in the form of a function (function upon restarting takes in shell command-line arguments) in a file [22], and

“heap2exec” to convert heap image into an executable which can be run faster. “heap2exec” wraps the binary SML/NJ runtime image and converts heap image into one executable. It makes use of built-in utility “heap2asm” to build a stand-alone executable from the image [25], which can be used at a later stage.

f. CM Autoloading:

Some of the libraries which come with SML/NJ as part of Compilation and Library Manager are not loaded initially. It needs some expertise to load these libraries which is unusual for an end user. *smlnj-script* loads such libraries and make them available to the user, one of the common is regular expression library. This eliminates the need for end user’s knowledge on CM.autoload.

g. Parse, Eval & Format:

Parse, Eval & Format gives the capability to make use of compilers ability to pretty print data which it has for implementing the redevelopment loop. In the last phase of the loop the results are printed out, SML/NJ has built-in ability to print out results of many types though they are not ordinarily available to the end users and programmers. At the end of each iteration loop this stores the value of the expressions which need to be printed out in a variable. It ignores the irrelevant part and keeps only the relevant part. This helps to convert the data of many types to strings.

h. Parser Combinators:

Parser Combinators gives the capability to operate on streams of arbitrary types rather than just streams of chars. This allows to implement Unicode support using the utf8 encoding.

i. UTF8 processing:

This module provides a structure with types and utility functions to work with Unicode characters written in utf8. It also provides structure with definitions which are needed during bootstrap and normal use.

This gives the ability to create stream of Unicode characters from a stream of bytes (old style ASCII characters).

j. Process function:

*smlnj-script* gives a value-added interface on top of the subprocess handling available in SML/NJ.

k. XML Wrapping:

This module is derived from SML XML library and work of Tom Murphy VII, which is again a wrapper of “fxp” [27]. It is used to parse XML files.

l. Utility functions:

*smlnj-script* have utility functions to handle basic data operations and load required libraries. Structures String and CharVector have good overlap, some members in CharVector don't have no equivalents in String and the same goes for Substring and CharVectorSlice. The function is customized to give the capability of string repetitions, concatenations, comparison, and conversion operations. One of the capabilities is to auto load JSON library, so that end user need not to load the JSON library.

## 2.9 How to run a script?

In UNIX, to execute a file or script, the script should start with a hashbang (“#!”) character followed by the path to the interpreter and any optional argument as follows,

```
#!PATH-TO-INTERPRETER OPTIONAL-ARGUMENT
```

and followed by instructions in the subsequent lines. The script also needs execute permission. Section 2.13 will explain in detail on what happens when a script is executed.

## 2.10 How to run *smlnj-script*?

*smlnj-script* file also follows the same rule as mentioned in 2.9. Let's see a sample script "HelloWorld" that runs *smlnj-script*, the script will start with the below statement followed by instructions in the subsequent lines.

```
#!/usr/bin/env smlnj-script P
```

When the above script is executed as below,

```
./HelloWorkd
```

the operating system will run the "/usr/bin/env" program with the arguments "smlnj-script" and P.

The Interpreter can be invoked directly passing it the name of the script as its first argument. For example, to run a script "HelloWorld" with three arguments arg1, arg2 and arg3, the interpreter can be invoked as below,

```
$ smlnj-script HelloWorld arg1 arg2 arg3
```

The script "HelloWorld" doesn't need to have "execute" permission if it's going to be passed as an argument to interpreter.

## 2.11 How to work with SML/NJ?

Users can invoke SML/NJ by typing the keyword "SML" over command prompt or terminal in Linux, which will load SML in interaction mode,

```
$ sml
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 16
18:42:54 2023]
-
```

Users can break the interactive mode and come out of SML/NJ by issuing the keyboard shortcut "Ctrl+D" or by issuing the SML command (OS.Process.exit OS.Process.success).

```
$ sml
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 16
18:42:54 2023]
- ^D
$
```

Users can type in the instructions line by line with the delimiter semicolon (;) at the end of line. The semicolon denotes the end of line and instructs the compiler to process the line thru REPL.

```
$ sml
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 16
18:42:54 2023]
- val x = "Hello World\n";
val x = "Hello World\n" : string
- print x;
Hello World
val it = () : unit
-
```

If an instruction is not delimited by a semicolon, the compiler will assume the instructions are continuing in next line and wait for a semicolon to mark the end of the instruction and submit it for REPL.

```
- val x = "Hello World\n"
= ;
val x = "Hello World\n" : string
- print
= x
= ;
Hello World
val it = () : unit
-
```

## 2.12 What happens when SML/NJ is loaded?

We now know that when a user uses SML/NJ, an interactive system loads and becomes available to the user, allowing the user to type in the instructions line by line for processing. Let's look a little deeper into what happens after the SML/NJ is called and before the user gets the interactive terminal.

SML/NJ is based on the SML language as well as C. The primary operating system-interaction routines are written in C, and SML is used to build additional functions on top of them. If we carefully examine the SML/NJ installation log, we can

see that all C programmes will be first built and loaded to produce a run file, and then all SML code will be parsed, compiled, and loaded to produce a heap image. The heap image is passed as an argument to the executable run file consisting of C binaries in order to load the SML.

Logs from terminal when SML is invoked,

```
$ sml
CMD=`basename "$0"`
basename "$0"
++ basename /usr/local/smlnj/bin/sml
+ CMD=sml
...
...
exec "$RUN" @SMLcmdname="$0" "$HEAP" $ALLOC "$@"
+ exec /usr/local/smlnj/bin/.run/run.x86-darwin
@SMLcmdname=/usr/local/smlnj/bin/sml
@SMLload=/usr/local/smlnj/bin/.heap/sml
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 16
18:42:54 2023]
-
```

The boot process is initiated by C executables, which installs the default basic functions and builds a boot environment, establishing a pervasive environment for the interactive system. If any command line parameters are present when the boot environment is loaded, they are checked and validated against the pre-defined values before the appropriate action is taken. Basic operations like arithmetic, string manipulation, symbol infix binding, and overloaded bindings that make use of overloaded declarations are tied to the pervasive environment.

Only the structures that contain the fundamental functions are loaded when the interactive system is made available to the user; the others are not. To reduce memory usage and improve efficiency, not all the structures are loaded. The pervasive environment's auto-loading procedure aids in loading the required libraries. The top-level environment is made available and a chosen few crucial libraries are set up to load automatically as needed.



## 2.13 What happens when a script is executed?

In this section, we'll talk about what happens when a script is run from a Unix machine's command prompt or terminal. When a script is run, the shell interpreter checks to see if it has the necessary execute rights, and if it does, it moves on to choose the appropriate interpreter. The interpreter continues to read and carry out the file's instructions after being identified.

Example:

```
$ ./samplescript
or
$ ./samplescript.sh
```

If the file does not already have the execute permission, it can be granted using the “chmod” command at the command line or terminal.

Example:

```
$ chmod +x samplescript
$ ./samplescript
or
$ chmod +x samplescript.sh
$ ./samplescript.sh
```

According to section 2.9, Unix scripts begin with the character sequence “#!”, often known as a hashbang, sharp-exclamation, or shebang. When a file starting with a hashbang is executed over operating system command prompt, the program loader parses the first line as an interpreter directive [26]. The program loader will execute the specified interpreter along with the path of the script passed as an argument to the interpreter.

For example, if a script starts as below,

```
#!/usr/bin/env perl
```

The script will be executed using the `env` command with the user's `$PATH` environment variable for the specified interpreter, in this case its Perl. Any additional arguments specified in the script first line will also be passed along to the interpreter.

For example, if a script starts as below,

```
#!/usr/bin/env -Ssml -arg1
```

The script will be executed using the `env` command with the user's `$PATH` environment variable for the specified interpreter. The `-S` option is used to pass the interpreter as an argument to the `env` command. The “sml” argument following the `-S` option specifies the interpreter to use. In this case, the interpreter being used to execute the script is SML. Any argument following the interpreter will be passed on as additional argument to the interpreter. In this case, the argument “arg1” is passed to the SML interpreter. If there are more additional arguments, they all will be passed along followed by end the path of the script.

## 2.14 Areas to be investigated in SML/NJ

As we know from the main objective of this dissertation, SML doesn't have the feature where a script with SML source code can't be executed. In other words, even if the SML interpreter is supplied as an argument to `env`, a script with SML source code cannot be run. The script is logically executed by the operating system, but it fails because SML is unable to accept the arguments that are supplied by the programme loader when a script is executed. Technically SML receives the argument but fails to recognise and decide what to be done with this argument. If the argument is not part of the pre-defined list.

For example, SML can take arguments like below,

```
$ SML -h
$ SML -H
$ SML -S
$ SML sample.sml
```

For example, the below scripts will fail,

```
#!/usr/bin/env -Ssml  
(or)  
#!/usr/bin/env -Ssml --script
```

Though SML interpreter is identified and passed with argument “--script”, the argument will fail the interpreter since the interpreter won’t recognise the argument. This is one of the causes for the development of *smlnj-script* as a layer on top of SML/NJ, which made it simple to add the ability to run an SML source file as a script. It also had additional features, which are described in section 2.6.

This is the most important and key change that need to be incorporated in SML/NJ to accept or recognise the new argument “—script” and able to process the instructions in the script. Technically this change will be interrupting the start-up or loading process of SML/NJ whereas other features planned are mostly interrupting the evaluation part. Implementation section will cover all the change information in detail.

### 3. Requirements Analysis

#### 3.1 Functional Requirements

The following are the functional requirements originally planned to be analyzed, built, tested, and delivered.

No.	Requirement Description	MoSCoW	Priority
1	Identify the right module and merge <i>smlnj-script</i> 's execute as a script capability into SML/NJ.	M	High
2	Identify the right module and merge <i>smlnj-script</i> 's SilenceCompiler function into SML/NJ.	M	High
3	Identify the right module and merge <i>smlnj-script</i> 's toString overloading capability into SML/NJ.	M	High
4	Identify the right module and merge <i>smlnj-script</i> 's CM Autoloading capability into SML/NJ.	M	High
5	Identify the right module and merge <i>smlnj-script</i> 's Parse, Eval & Format functions into SML/NJ.	M	High
6	Identify the right module and merge <i>smlnj-script</i> 's Parser Combinators capability into SML/NJ.	M	Medium
7	Identify the right module and merge <i>smlnj-script</i> 's Utf8 processing capability into SML/NJ.	M	Medium
8	Identify the right module and merge <i>smlnj-script</i> 's Process functions into SML/NJ.	S/C	Low
9	Identify the right module and merge <i>smlnj-script</i> 's XML Wrapping capability into SML/NJ.	S/C	Low
10	Identify the right module and merge <i>smlnj-script</i> 's Utility functions into SML/NJ.	S/C	Low

### 3.2 Non-Functional Requirements

The following non-functional requirements are planned to be analyzed, built, tested, and delivered.

No.	Requirement Description	MoSCoW	Priority
1	Usability – End users, programmers, developers, and students should be able to use this final product to run the scripts with minimal effort.	M	High
2	Stability – The final product should be stable enough to run the scripts without getting crashed.	M	High
3	Reliability – The final product should be capable to producing accurate results.	M	High

## 4. Methodology

This dissertation adopted the modern software development practice Continuous Integration and Continuous Development (CI/CD), where incremental code changes are reliably and regularly made. Since we are updating a language itself, we anticipated lot of trial and errors, going back and forth, and exploring new ideas, so we decided to adopt CI/CD. This would help us to manage the change and the flexibility to add or drop an idea and a change.

Code updates that are merged into the repository are made reliable by automated build-and-test procedures that are sparked by CI. As part of the delivery process, the change is quickly and seamlessly delivered. As an extension to Continuous Integration, Continuous delivery (CD) is the automated delivery of completed code to the testing and production environments. The CI/CD pipeline, as used in the software industry, is the automation that enables developers to quickly and reliably deliver the incremental code changes from their development environment to production environment.

The CI/CD helps to automate the testing process, whenever a change is made, the code is automatically moved to a test environment and gets executed or tested. It also gives access to a unique platform to manage all the software artefacts and to keep track of each change. This in turn helps to identify and fix errors at the earliest. The next phase after continuous delivery is continuous deployment. There are numerous production deployments because every change that passes the automated tests is automatically moved into production.

The project's goal is to introduce new features to SML/NJ implementors so they can be included in a later release. In the event of a redistribution, the initiative will assist licensing agreements between HWU and SML/NJ implementors. We anticipate that HWU will either unilaterally license the project code under the same conditions as SML/NJ or will assign its copyright interest in *smlnj-script* to "The Fellowship of SML/NJ," with either course of action being discussed with SML/NJ maintainers before moving forward.

## 5. Professional, Legal, Ethical, and Social issues

### 5.1 Professional Issues

- This project conforms to the terms and conditions of the SML/NJ licence and use of all software and technologies provided by HWU in line with their terms of licence and terms and conditions.
- If SML/NJ implementors agree to the change, this project will be included back into SML/NJ with the appropriate licencing approval from Dr Joe Wells and HWU for utilising *smlnj-script* features.
- Any open-source software from a third party that is utilised will receive due acknowledgement.

### 5.2 Legal Issues

- This project adheres to HWU's licencing policies, and any source code protected by intellectual property rights is reused.
- If an agreement between HWU and SML/NJ implementors is necessary to add project features to SML/NJ, it will be arranged. If the agreement is unsuccessful or if either party is unwilling, the project shall remain the property of the author and HWU.
- For any exploitation and dissemination of third-party content, formal approval from the contributors will be obtained.

### 5.3 Ethical Issues

- No human survey has been done or will be done; this project is a technical implementation on top of an existing open-source tool; it does not involve any sensitive personal data of an individual. Only source code and author information are gathered for this project.

### 5.4 Social Issues

- This project stores source code and compiler information, as well as personally identifiable information like author names, which are both preserved and cited in the project papers.

## 6. Project Risk Assessment

The following risks have been identified, impact analysed, and mitigated.

<b>Risk</b>	<b>Likelihood</b>	<b>Impact</b>	<b>Mitigation</b>
Author's medical emergency.	Medium	High	Work with supervisor and University for extension or resubmission.
Supervisor's medical emergency.	Medium	High	Work with University and personal tutor for additional support.
Loss of project artifacts.	Low	Medium	Periodic backup will be taken in GitHub and University locations.
SML/NJ releases a new update.	Low	Low	Features that are planned to merge with SML/NJ will be modified to adapt to latest update from SML/NJ.
Unavailability of technical information	High	High	Reach out to SML/NJ implementors for details.
Project gets delayed and overrun planned timeframe.	Medium	Medium	High priority requirements will take precedence.
Software not compatible with Author's hardware.	Low	Low	University resources will be utilized.



## 7. Design and Implementation

### 7.1 Pre-requisite

Dr Joe Wells came up with the concept for *smlnj-script* a few years ago. It was built with SML/NJ version 110.78, which at the time only supported 32-bit, and it was tested with that version. SML/NJ got numerous updates over time and started supporting 64-bit. The most recent stable versions are 110.99.3, 2021.1, and 2022.1. Before we wanted to start on the dissertation, we tried to recompile *smlnj-script* with the latest versions of SML/NJ in a macOS machine. The exercise was originally unsuccessful; therefore, we need to update the original *smlnj-script* build file to reflect the most recent XML libraries and source files (which it was using from the SML/NJ library) before it will work again. To support SML source and heap2asm modules, the SML configuration file must also be changed. We were able to verify that *smlnj-script* is supported by both the 32-bit and 64-bit versions of SML 110.99.3 after a successful installation of SML and recompilation of the script. Not all of the features, however, were successful. For example, the *smlnj-script* feature that was intended to facilitate string overloading using the special character "%" did not work.

### 7.2 Software and Hardware used

SML/NJ version 110.99.3 (32 bit) was utilised with macOS 10.13.16 on an Intel-based MacBook Pro. Even though SML's most recent versions, 2021.1 and 2022.2 were available their support is available for 64-bit only, and in 64-bit versions we ran into a problem with the structure Backend, which SML can detect but cannot load along with the related structures and functions in interaction mode. We didn't have any such problem with 32-bit and got confirmed with SML implementors that this was a known problem with SML 110.99.3 (64-bit). Since 32-bit support ended with 110.99.3, we ultimately chose to use the 32-bit version of SML.

The standard SML implementation is not accompanied by source code. SML requires manual installation, to make change in source code, recompile and reinstall. In manual installation, upon downloading the config.tgz files from SML/NJ portal, the targets file in the configuration folder needs to be modified to request for source code before installation. This will make it possible to update the source code, rebuild, and reinstall SML.

### 7.3 Where to be updated in SML/NJ

Following files in the SML/NJ library were amended to include the *smlnj-script* capabilities in to SML/NJ,

- `smlnj/base/cm/main/cm-boot.sml`
- `smlnj/base/system/smlnj/internal/boot-env-fn.sml`
- `smlnj/base/compiler/TopLevel/interact/interact.sig`
- `smlnj/base/compiler/TopLevel/interact/interact.sml`
- `smlnj/base/compiler/TopLevel/backend/backend-fn.sml`
- `smlnj/base/compiler/TopLevel/backend/backend.sig`
- `smlnj/base/compiler/TopLevel/interact/mutecompiler.sig` (New component)
- `smlnj/base/compiler/TopLevel/interact/mutecompiler.sml` (New component)
- `smlnj/base/compiler/INDEX`
- `smlnj/base/compiler/MAP`
- `smlnj/base/compiler/core.cm`

### 7.4 Change details

The enhancements made to SML/NJ using Dayanandan Natarajan's own research and features taken from the *smlnj-script* created by Dr Joe Wells are explained in the section that follows.

#### 7.4.1 File `cm-boot.sml`

Change made in this file is the result of author's own research and development. File *cm-boot.sml* is the very first file that is loaded when the CM system is invoked. It contains the basic definitions and functions that are needed to initialize the system. Some of the key structures and functions are defined in this file. Structure CM is defined with main functions like CM.make, CM.recompile, and CM.sources to manage the compilation process. Functor LinkCM (which is linked with Structure Backend) is defined with key functions and one such is to process the command-line arguments passed by user through functions *args* and *carg*.

Based on the command-line arguments passed, appropriate functions are invoked. From the command prompt or terminal along with sml interpreter if a file is passed as an argument, it passed thru function *processFile* for further processing. In

*processFile*, the file extension is evaluated for and put through another function *useFile* for extensions “sml”, “sig” and “fun”, and function *make* for extension “cm”.

In addition to that it also defines the basic error handling and exception handling functions that are used throughout the CM system. These functions are used to handle errors and exceptions that occur during the compilation process in a more efficient way. File *cm-boot.sml* plays a vital role in initializing the CM system and providing the basic functionality that is needed for managing the compilation of SML programs.

The following changes were made in file *cm-boot.sml* to recognise the new command line parameter passed from script.

- a) *init* function is the point of entry for all the functionalities defined in file *cm-boot.sml*. It's been called with other key functions imported from other structures as arguments. A new function *useScriptFile* added in Structure *Interact* has been added as an additional argument to *init* function. Purpose and usage of function *useScriptFile* will be detailed in the upcoming sections.

```
fun init (bootdir, de, er, useStream, useScriptFile, useFile,
errorwrap, icm) = let ...
```

- b) Function *procCmdLine* is one of the key functions in file *cm-boot.sml* which process the command-line instructions and arguments. A new function *processFileScript* is added as a sub function to *procCmdLine*. Function *processFileScript* takes the script file name passed as argument, performs a validation through function *checkSharpbang*, consumes the contents of first line through function *eatuntilnewline* and calls the function *useScriptFile* to process the script. Function *useScriptFile* is called with file name and its content in the form of stream as arguments.

Source code:

```
fun processFileScript (fname) = let
  val stream = TextIO.openIn fname
  val isscript = checkSharpbang stream
in
```

```

        if (isscript) = false
        then      ( Say.say [ "!* Script file doesn't start
                        with #!. \n" ] )
        else      ( useScriptFile (fname, stream) )
    end
end

```

- c) Function *checkSharpbang* takes file stream as an argument, which is called with the script file contents as a stream from the calling function. The function is used to verify whether the first line of the stream starts with “#!”. In that case the function returns a Boolean response “true” and if the stream doesn’t start with “#!” then the function returns the Boolean response “false”.

Source code:

```

fun checkSharpbang (instream : TextIO.instream): bool = let
    val c = TextIO.input1 instream
in
    case c of
        SOME #"#" => (
            case TextIO.lookahead instream of
                SOME #"!" => eatuntilnewline
                _ => false
            )
        | SOME c => false
        | NONE => false
    )
    | SOME c => false
    | NONE => false
end

```

- d) The contents of the first line of the script up until the new line character “\n” are eaten since the REPL won’t be able to handle the first line if the file that is passed as an argument to the function *processFileScript* is confirmed to be a script by function *checkSharpbang*. SML won’t be able to grasp the instructions because the first line of the script, which begins with “#!” is more for the operating system to understand. A few other strategies were tested until the first line consumption solution was found. These included either remarking the first line using the single-line comment “(\*)” or commenting the entire line using “(\* \*)”. In both methods, the partial file must either be moved to a new file, or the original content must be

changed before it can be sent for evaluation. It is chosen to utilise the SML feature to consume the already-accessed characters to get around this. The script file stream is supplied as an argument when the script function `eatuntilnewline` is used to consume the first line of the script.

- e) Function *eatuntilnewline* will consume all the characters in the script first line until a new line “\n” character is encountered. Once the new line character is reached the function returns Boolean value “true” to the calling function, which at that point will have the file stream stripped with first line, passes the file stream for further processing.

Source code:

```
fun eatuntilnewline (instream : TextIO.instream): bool = let
  val c = TextIO.input1 instream
in
  case TextIO.lookahead instream of
    SOME #"\n" => true
  | SOME c => eatuntilnewline instream
  | NONE => false
end
```

- f) SML can take in three type of arguments “rtsargs” (Runtime system arguments), “options” and “files”. “rtsargs” arguments are used to load specific information from SML Basic Library. “options” are used either to set CM variables or to invoke help or to change any settings. “files” are used either to load instructions to the top level environment through (.sml/.sig/.fun) files or load custom libraries through .cm files.

Function *args* is one of the key functions in functor *LinkCM* which validates the command line arguments. When a valid argument is passed from the command prompt or terminal, appropriate task is performed by calling the respective functions. Only “options” and “files” arguments are handled by function *args*. “rtsargs” arguments are handled separate in a different functor which is discussed later in this document. The arguments will be looked for whether a standalone argument or the following argument in the list is a supplementary argument to previous argument. Standalone arguments referred

here are the ones which are used to perform a specific task without the need for next argument, for example passing “-h” or “-H” for help in command prompt. Consider another example “-type f” where the first argument needs the second argument to be present, this is what I referred here as supplementary argument. If the arguments are not in the predefined list in function *args* or standalone argument, the control is transferred to function *carg*. Only the first two characters of the argument is passed to function *carg*, the assumption here is only the standalone arguments are passed. Function *carg* does almost the same as function *args*, validates the argument passed and appropriate task is performed by calling its corresponding functions.

Function *args* is updated to recognise the new argument “--script” which is passed by the program loader from the script. Whenever “—script” is passed as an argument, a new function *nextargscript* is called.

Source code:

```
fun args ([], _) = ()
  | args ("-a" :: _, _) = nextarg autoloader
  | args ("-m" :: _, _) = nextarg make'
  | args (["-H"], _) = (help NONE; quit ())
  | args ("-H" :: _ :: _, mk) = (help NONE; nextarg mk)
  | args (["-S"], _) = (showcur NONE; quit ())
  | args ("-S" :: _ :: _, mk) = (showcur NONE; nextarg mk)
  | args (["-E"], _) = (show_envvars NONE; quit ())
  | args ("-E" :: _ :: _, mk) = (show_envvars NONE; nextarg
mk)

  | args ("--script" :: _, _) = (nextargscript ()) (* line
added by DAYA *)
  | args ("@CMbuild" :: rest, _) = mlbuild rest
  | args (["@CMredump", heapfile], _) = redump_heap heapfile
  | args (f :: rest, mk) =
(carg (String.substring (f, 0, 2)
  handle General.Subscript => "",
  f, mk, List.null rest);
  nextarg mk)
and nextarg mk =
let val l = SMLofNJ.getArgs ()
in SMLofNJ.shiftArgs (); args (l, mk)
end
```

- g) A new function *nextargscript* is defined to look for the remaining values in the arguments list, retrieve the head of the remaining argument's list (which is expected to be the file name of the script) and pass the file name as an argument to another new function *processFileScript*.

The general assumption here is, first line in the script starting with “#!” doesn't have any argument passed after “—script”. In case of multiple argument to be passed, “—script” will be last one in the list. The reason behind is, in the script when option “-S” is used, it passes the interpreter as an argument to the “env” followed by all other entries as arguments to the interpreter and finally the name of the script passed as the last argument to the interpreter. This means, the name of the script will be in the argument list immediately after entry “--script”. When “—script” is passed from script, it gets picked by function *args* and function *nextargscript* is called to retrieve the script file name and call the function *processFileScript* with the file name. Upon completion of function *processFileScript*, the process gets terminated and returns control to the command prompt or terminal.

Source code:

```
and nextargscript () =  
    let val l = SMLofNJ.getArgs ()  
    in SMLofNJ.shiftArgs (); processFileScript (hd l); quit ()  
    end
```

#### 7.4.2 File *boot-env-fn.sml*

Change made in this file is the result of author's own research and development. File *boot-env-fn.sml* hosts the signature *BOOTENV* and functor *BootEnvF*, they are the key components in creating the bootstrap environment. Key functions that are required to initiate the pervasive environment are declared here for importing. Function *bootArgs* plays a key role in the bootstrap environment process, it handles the runtime system arguments passed from command line. These runtime arguments starting with “@SML\*” are validated here and appropriate function is performed. Functions that are passed as arguments to the *init* function in file *cm-boot.sml* are declared here with their arguments data type and return data type.

In functor *BootEnvF*, function *cminit* declaration is amended to include the argument and return data type of newly added function *useScriptFile*.

Source code:

```
functor BootEnvF (datatype envrequest = AUTOLOAD | BARE
  val architecture: string
  val cminit : string * DynamicEnv.env * envrequest
    * (TextIO.instream -> unit) (* useStream *)
    * (string * TextIO.instream -> unit) (*
useScriptFile *)
    * (string -> unit) (* useFile *)
    * ((string -> unit) -> (string -> unit))
      (* errorwrap *)
    * ({ manageImport:
      Ast.dec * EnvRef.envref -> unit,
    managePrint:
      Symbol.symbol * EnvRef.envref -> unit,
    getPending : unit -> Symbol.symbol list }
-> unit)
    -> (unit -> unit) option
  val cmbmake: string * bool -> unit) :> BOOTENV =
struct

fun useFile f = if Backend.Interact.useFile f
  then ()
  else OS.Process.exit OS.Process.failure
in
  U.pStruct := U.NILrde;
  cminit (bootdir, de, er,
    Backend.Interact.useStream,
    Backend.Interact.useScriptFile, (* added as part
of Execute as a script change *)
    errorwrap false useFile,
    errorwrap true,
    Backend.Interact.installCompManagers)
end
```



### 7.4.3 Files `interact.sig` & `interact.sml`

Change made in this file is the result of author's own research and development. The idea of calling function *mcdummyfn* by passing it as an instruction; to be evaluated through another function *evalstream* is given by Dr Joe Wells. One of the key structures in SML is Backend, it encapsulates the low-level details of the compilation and interpretation process. It allows the compiler and interpreter to operate at the higher level of abstraction. Interact is a sub structure within Backend which provides functions to perform interactive input and output operations over a command prompt or terminal environment. Some of the key functions of Interact are *useFile* (which takes SML source in a file and add those declarations in the file to the top-level environment) *useStream* and *evalStream* (which takes SML source in the form of stream and add those declarations in the stream to the top-level environment) and *withErrorHandling* (which handles the error). A new function *useScriptFile* is added to Interact structure, which takes the file name and its content as a stream and process the stream by passing it to function *evalStream* hosted in structure *EvalLoop*.

- a) New function declaration is added to `interact.sig`,

Source code:

```
val useStream : TextIO.instream -> unit
val useScriptFile : string * TextIO.instream -> unit (*
Added by DAYA *)
val evalStream : TextIO.instream *
Environment.environment -> Environment.environment
```

- b) New function definition for *useScriptFile* is added to `interact.sml`.

Function *useScriptFile* takes in script file and stream as arguments and pass the same to function *evalStream* to process the stream over REPL. From the new structure *Mutecompiler*, functions *silenceCompiler*, *unsilenceCompiler* and *printStashedCompilerPutput* are called-in here to perform specific tasks which are part of the key requirements. We will discuss more in detail on these functions later in this document. In short, function *silenceCompiler* is used to save the compiler output to a reference cell, function *unsilenceCompiler* will revert the saved compiler output

from reference cell to compiler and function *printStashedCompilerPutput* will print the saved or stashed output to the terminal for user.

Before the file stream is passed to function *evalStream* for processing, the compiler messages are muted, a dummy function is called (which does nothing) and the compiler messages are unmuted. This task is very much needed to make the script output clear of unnecessary messages presented to the user. This will suppress the auto-loading messages for structure *Mutecompiler* in the script output which is totally unnecessary and could be avoided for pretty printing.

Source code:

```
fun useScriptFile (fname, stream) = (  
    Mutecompiler.silenceCompiler () ;  
    EvalLoop.evalStream ("<instream>",  
    (TextIO.openString "Backend.Mutecompiler.mcdummyfn () ;")  
    ) ;  
    Mutecompiler.unsilenceCompiler () ;  
  
    (EvalLoop.evalStream (fname, stream))  
    handle exn => (  
        Mutecompiler.printStashedCompilerOutput () ;  
        Mutecompiler.unsilenceCompiler () ;  
        EvalLoop.uncaughtExnMessage exn  
    )  
    )  
    )
```

In case of error in the script file, the compiler will throw the error message along with file name and line number information. And it prints the details of the exception thrown to the terminal. But this is not the case if the user has decided to mute the compiler by calling function *silenceCompiler*. When the compiler messages are muted, the error messages are stashed and unavailable to be printed to the terminal. These messages need to be retrieved from stash and printed to the user; and the compiler need to be unmuted to print the exception details. Function *printStashedCompilerPutput* is called to print the stashed output to user and function *unsilenceCompiler* is called to unmute the compiler before

throwing the uncaught exception, so that the exception details are printed to the terminal.

#### 7.4.4 File backend.sig

Change made in this file is the result of author's own research and development. A new structure *Mutecompiler* is declared within signature BACKEND,

Source code:

```
signature BACKEND = sig
  structure Profile : PROFILE
  structure Compile : COMPILE
  structure Interact : INTERACT
  structure Mutecompiler : MUTECOMPILER
  structure Machine : MACHINE
  val architecture: string
  val abi_variant: string option
end
```

#### 7.4.5 File backend-fn.sml

Change made in this file is the result of author's own research and development. New structure *Mutecompiler* is defined within functor *BackendFn*,

Source code:

```
structure Mutecompiler = Mutecompiler
```

#### 7.4.6 File mutecompiler.sig

Change made in this file is the result of author's own research and development. Though the idea of this signature file is from the author, the declaration of few variables and functions is a necessity which in turn is derived from *smlnj-script* developed by Dr Joe Wells. The details of these variables and functions are described in the next section. New signature *MUTECOMPILER* is defined with all the global variables and functions that are part of Structure *Mutecompiler*.

Source code:

```
signature MUTECOMPILER =
```

```

sig
  val printlineLimit : int ref
  val compilerMuted : bool ref
  val isNewline : char -> bool
  val push : "a list ref -> "a -> unit
  val installPrintingLimitSettings : int list -> unit
  val saveControlPrintOut : unit -> unit
  val stashCompilerOutput : string -> unit
  val savePrintingLimits : unit -> unit
  val lowerPrintingLimitsToMin : unit -> unit
  val restoreControlPrintOut : unit -> unit
  val restorePrintingLimits : unit -> unit
  val outputFlush : TextIO.outstream -> TextIO.vector -
> unit
  val silenceCompiler : unit -> unit
  val unsilenceCompiler : unit -> unit
  val printStashedCompilerOutput : unit -> unit
  val mcdummyfn : unit -> unit
end (* signature MUTECOMPILER *)

```

#### 7.4.7 File `mutecompiler.sml`

Change made in this file is the result of author's own research, development and capabilities derived from *smlnj-script* which is developed by Dr Joe Wells. Though most of the functions are derived from *smlnj-script*, the idea to have a new structure, placement of this structure in the master structure *Backend* under *Interact*, design and built of this structure is author's work. New structure *Mutecompiler* is defined with global variables, core functions and sub-functions which are required to perform the tasks like muting the compiler, unmuting the compiler, printing the stashed compiler messages back to the terminal and increase the print limits.

- a. Function *silenceCompiler* and the functions called in it are derived from *smlnj-script* developed from Dr Joe Wells, and with a minor change by the author to assign value to a Boolean variable *compilerMuted* which is later used by other function. Function *silenceCompiler* is defined to mute the compiler messages. The compiler messages are suppressed by saving all the "Control.Print.out" values to a reference cell. It also saves the current printing limits in a ref cell and then set them all to zero.

### Source code:

```
fun silenceCompiler () =
  (compilerMuted := true;
   saveControlPrintOut ();
   Control.Print.out := { flush = fn () => (), say
= stashCompilerOutput };
   savePrintingLimits ();
   lowerPrintingLimitsToMin ());

fun saveControlPrintOut () =
  if isSome (! savedControlPrintOut)
  then ()
  else savedControlPrintOut := SOME (!
Control.Print.out);

fun stashCompilerOutput string
  = case String.fields isNewline string
  of nil => raise (Fail "impossible ")
  | [chunk] => push compilerOutputCurrentLine chunk
  | chunk :: lines
  => (if chunk <> "" then push
compilerOutputCurrentLine chunk else ();
      push compilerOutputPreviousFullLines
        (String.concat (rev (!
compilerOutputCurrentLine))));
      let val (last :: others) = rev lines
      in app (push
compilerOutputPreviousFullLines)
        (rev others);
        compilerOutputCurrentLine
        := (if last <> "" then [last] else [])
      end);

fun savePrintingLimits () =
  if isSome (! savedPrintingLimitSettings)
  then ()
  else savedPrintingLimitSettings := SOME (map !
printingLimitRefs);

fun lowerPrintingLimitsToMin () =
  List.app (fn r => r := 0) printingLimitRefs;
```

- b. Function *unsilenceCompiler* and the functions called in it are derived from *smlnj-script* developed from Dr Joe Wells, and with a minor change by the author to assign value to a Boolean variable *compilerMuted* which is later used by other function. Function *unsilenceCompiler* unmutes the compiler messages by restoring the printing limits and value of “Control.Print.out”.

Source code:

```
fun unsilenceCompiler () = (compilerMuted := false;
                           restoreControlPrintOut ();
                           restorePrintingLimits ());

fun restoreControlPrintOut () =
  case ! savedControlPrintOut of
    NONE => ()
  | SOME value => (savedControlPrintOut :=
NONE;
                  Control.Print.out :=
value);

fun restorePrintingLimits () =
  case ! savedPrintingLimitSettings of
    NONE => ()
  | SOME settings =>
(savedPrintingLimitSettings := NONE;
 installPrintingLimitSettings settings);
```

- c. Function *printStashedCompilerOutput* and the functions called in it are derived from *smlnj-script* developed from Dr Joe Wells with few changes by the author. Author’s own research and development include, using variable *compilerMuted* to decide to print only when the compiler is silenced and a global variable *printlineLimit* which gives the capability to increase or decrease the printing line limits. Function *printStashedCompilerOutput* is defined to print the saved or stashed compiler output back to the terminal. Whenever the compiler messages are suppressed through *silenceCompiler* function and an error occurs, these error messages are stashed and need to be printed back to the terminal. By

invoking this function, the stashed messages can be printed back to the terminal.

- d. Function *dummyfn* is an idea put forward by Dr Joe Wells during the development of this dissertation. Function *dummyfn* is created (which does nothing) and is called to preload the *Mutecompiler* structure before the script is passed to *evalStream* function. This is to suppress the structure auto-loading logs in the script results. If this dummy function is not called, then the user can see the auto-loading logs in the script result when one of the *Mutecompiler* function or variable is accessed. This is more of a cosmetic feature, to make sure the script results don't have unnecessary logs in it.
- e. Variable *compilerMuted* is introduced by author to know whether the compiler messages are silenced.
- f. Variable *printlineLimit* is introduced by author to increase or decrease the number of lines is retrieved from the stashed output before they are printed to the terminal. By default, this limit is set to 5, but the user can change this limit if they want to see more details from the stashed output.

The limit can be changed as below,

```
Backend.Mutecompiler.printlineLimit := 10;
```

#### 7.4.8 Files INDEX, MAP and core.cm

Change made in these files are the result of author's own research and development. Files *INDEX*, *MAP* and *core.cm* are updated with definitions for signature *MUTECOMPILER* and structure *Mutecompiler*, this will let CM know there is a new signature and structure defined and to be made available for usage.

- a. In file *INDEX* the definitions are added to reflect the new signature and structure as below,

```
MUTECOMPILER
  TopLevel/interact/mutecompiler.sig
```

```
Mutecompiler
  TopLevel/interact/mutecompiler.sml
```

- b. In file *MAP* the definitions are added to reflect the new signature and structure as below,

```
interact/
  envref.sml
    supports top-level environment management
    defs: ENVREF, EnvRef : ENVREF
  evalloop.sig,sml
    top-level read-eval-print loop
    defs: EVALLOOP, EvalLoopF: TOP_COMPILE => EVALLOOP
  interact.sig,sml
    creating top-level loops
    defs: INTERACT, Interact: EVALLOOP => INTERACT
  mutecompiler.sig,sml
    allow compiler silencing
    defs: MUTECOMPILER, Mutecompiler
```

- c. In file *core.cm* the definitions are added to reflect the new signature and structure as below,

```
TopLevel/interact/mutecompiler.sig
TopLevel/interact/mutecompiler.sml
```

## 7.5 Writing a script

The script should start with “#!” in the first line followed by the environment location, command-line parameters “-Ssml” and “—script”, a new line and then followed by the SML code or program.

Example script named “sample”,

```
-----beginning of the script-----
#!/usr/bin/env -Ssml -script
; (*--SML--*)
val () = print "Hello World\n";
-----end of the script-----
```



## 7.6 Executing a script

The script “sample” can be executed from terminal or command prompt as a regular OS script as below (provided it is given execute permission),

```
$ ./sample
```

## 7.7 SML/NJ Implementation and 64-bit Support

As mentioned in section 7.1 and 7.2 the change was initially started with 32-bit version of SML/NJ 110.99.3. As part of our change, we are preloading structure *Mutecompiler* through a dummy function to avoid auto-loading logs in the output. This in-turn preloads structure Backend and solved the issue reported in 7.2. We were able to test our changes successfully in SML/NJ 110.99.3 64-bit version. There is no separate code base for 32-bit and 64-bit, the code base is same and only installation instruction for SML/NJ has a minor variation for 32-bit, which is as per the standard implementation instructions.

SML/NJ can be implemented in one of two methods described below,

1. Installation file config.tgz downloaded directly from SML/NJ portal (<https://smlnj.org/dist/working/110.99.3/index.html>) and follow the command line instructions in the same portal. The installation should include source file and re-compiling support, and this is achieved by enabling request for source in “targets” file (part of installation files in “config” directory).
2. Installation from GitHub repository “smlnj/legacy - <https://github.com/smlnj/legacy>”. The installation and re-compile instructions are available in the same portal.

SML/NJ package installation is available which doesn’t have the support for source code and recompiling, it is recommended to follow one of two installation procedures recommended above.

## 7.8 GitHub Codebase

Codebase for this dissertation is forked from SML/NJ’s GitHub repository “legacy - <https://github.com/smlnj/legacy>“. The code is forked into author’s GitHub repository <https://github.com/dn2007hw/legacy>.

The changes are available in two branches as below,

- 1) Branch “dn2007hw-patch-1-execute-as-a-script” include change for execute as a script functionality. (<https://github.com/dn2007hw/legacy/tree/dn2007hw-patch-1-execute-as-a-script>)
- 2) Branch “dn2007hw-patch-2-with-mutecompiler” include changes for execute as a script capability and Mutecompiler functions.  
(<https://github.com/dn2007hw/legacy/tree/dn2007hw-patch-2-with-mutecompiler>)

## 7.9 GitHub Test

A test workflow is also created and available to test these changes over GitHub test environment. Two separate workflows were created to test in macOS and Ubuntu operating systems. The workflows are available in the following path,

1. MacOS -  
<https://github.com/dn2007hw/legacy/blob/main/.github/workflows/FullTestMacOS.yml>
2. Ubuntu -  
<https://github.com/dn2007hw/legacy/blob/main/.github/workflows/FullTestUbuntu.yml>

These workflows are configured to run manually on user request. By default, these workflow take the changes from branch “dn2007hw-patch-2-with-mutecompiler” and a complete end-2-end testing will be performed. These workflows will first checkout the code from the specified branch and install SML. Upon completion of install, SML recompile from source procedure will follow. On completion of recompilation and reinstallation, a test script is prepared and run on the terminal.

The workflows can be reconfigured to test the changes from a different branch or execute a test script from the repository rather than dynamically creating one.

## 7.10 Merging with SML/NJ

Two separate pull requests have been raised with SML/NJ to merge the changes into SML/NJ legacy repository and they are under review by the SML/NJ implementors David MacQueen and John H. Reppy from University of Chicago. The pull requests are available as below,

1. #276 for “Execute as a script” from branch “dn2007hw-patch-1-execute-as-a-script”, <https://github.com/smlnj/legacy/pull/276>.
2. #275 for “Execute as a script with Mutecompiler” from branch “dn2007hw-patch-2-with-mutecompiler“, <https://github.com/smlnj/legacy/pull/275>.

## 7.11 Issues and Interaction with David MacQueen & John Reppy

Along the development of this change, we raised few issue requests with SML/NJ team in GitHub to understand few concepts and behaviour, here are the details.

- Unable to load structure Backend in interactive mode of SML/NJ 110.99.3 64 bit, SML was recognising the structure but fails to load the substructures and functions declared in it. SML SME David MacQueen pointed out about the similar requests that were made in the past and a solution was provided, but that didn't serve our need. Details available at <https://github.com/smlnj/legacy/issues/268>.
- We were attempting to make a custom-built function defined in a structure be available in top-level (and also without being associated with any structure). We had an interesting conversation with David MacQueen on the basics of CM, top-level and pervasive environment. Details available at <https://github.com/smlnj/legacy/issues/273>.
- In SML/NJ V100.99.3 32-bit, the recompile from source procedure was failing at the very step, in ‘CMB Make’. There was a change in 2023 January by SML SME John Reppy to change random number generator module. The 32-bit components were failing, but not with 64-bit components. As a result of this issue, SML SME John Reppy fixed the issue, and we were able to continue with recompilation in 32-bit version. Details available at <https://github.com/smlnj/legacy/issues/277>.

## 8. Evaluation and Results

### 8.1 Before the change

As we saw in section 2.13, SML doesn't have the capability to allow a SML source file to be executed like a script. Let's see an example, a simple script "sample" as below,

```
$ cat sample
#!/usr/bin/env -Ssml
>(* SML code starts here *)
val x = "Hello World x\n";
val () = print x;
$
```

When the above script is executed with the SML logs enabled, it failed with the below error messages,

```
$ ./sample
.....
exec "$RUN" @SMLcmdname="$0" "$HEAP" $ALLOC "$@"
+ exec /usr/local/smlnj/bin/.run/run.amd64-darwin
@SMLcmdname=/usr/local/smlnj/bin/sml @SMLload=./sample
/usr/local/smlnj/bin/sml: Fatal error -- incorrect byte order
in heap image
$
```

Though the program loader recognised the interpreter as SML, the SML was not able to load its heap image due to incorrect arguments passed on to the interpreter. If we add an additional argument "--script" to the same script,

```
$ cat sample
#!/usr/bin/env -Ssml --script
>(* SML code starts here *)
val x = "Hello World x\n";
val () = print x;
$
```

The program loader recognised the interpreter as SML, the SML got loaded with heap image and the remaining arguments were passed on to the interpreter. Due to unavailability of the feature, SML was unable to recognise the argument "--script" and process the file; it continued to interactive mode.

```
$ ./sample
...
exec "$RUN" @SMLcmdname="$0" "$HEAP" $ALLOC "$@"
```

```

+ exec /usr/local/smlnj/bin/.run/run.amd64-darwin
@SMLcmdname=/usr/local/smlnj/bin/sml
@SMLload=/usr/local/smlnj/bin/.heap/sml --script ./sample
Standard ML of New Jersey (64-bit) v110.99.3 [built: Thu Jul 28
00:35:16 2022]
!* unable to process "--script" (unknown extension "<none>")
!* unable to process "./sample" (unknown extension "<none>")
- ^D
$

```

## 8.2 After the change

### 8.2.1 Execute as a script.

Post implementation of the changes, we were able to execute SML source code as scripts. Let's consider the same example,

```

$ cat sample
#!/usr/bin/env -Ssml --script
>(* SML code starts here *)
val x = "Hello World x\n";
val () = print x;
$

```

Now that we have implemented the capability, we were able to execute the source code as a script,

```

$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x
$

```

By enabling the logs to verify the local variables, we can see the program loader able to identify the SML interpreter and pass on the arguments to the interpreter.

```

$ ./sample
...
...
exec "$RUN" @SMLcmdname="$0" "$HEAP" $ALLOC "$@"
+ exec /Users/ndaya/withmc32/bin/.run/run.x86-darwin
@SMLcmdname=/Users/ndaya/withmc32/bin/sml
@SMLload=/Users/ndaya/withmc32/bin/.heap/sml --script ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x
$

```

Below is the log when SML is loaded directly from the command prompt and the arguments passed to the interpreter.

```

exec "$RUN" @SMLcmdname="$0" "$HEAP" $ALLOC "$@"
+ exec /Users/ndaya/withmc32/bin/.run/run.x86-darwin
@SMLcmdname=/Users/ndaya/withmc32/bin/sml
@SMLload=/Users/ndaya/withmc32/bin/.heap/sml
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]

```

Below is the log when a SML source code “sample” is executed as a script, we can see the interpreter is loaded with default parameters @SMLcmdname and @SMLload, followed by newly introduced argument “--script” and the script file name “sample”.

```

exec "$RUN" @SMLcmdname="$0" "$HEAP" $ALLOC "$@"
+ exec /Users/ndaya/withmc32/bin/.run/run.x86-darwin
@SMLcmdname=/Users/ndaya/withmc32/bin/sml
@SMLload=/Users/ndaya/withmc32/bin/.heap/sml --script ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]

```

The newly introduced changes as described in section 7.3.1, 7.3.2 and 7.3.3 have recognised the script, loaded the SML and executed the instructions in the script file.

### 8.2.2 Silencing compiler messages

We were also able to successfully verify the capability we have introduced to silence the compiler messages. This is achieved by introducing a new structure *Mutecompiler* and custom-built functions within it. The features include, silencing/muting compiler messages, un-silencing/unmuting compiler messages, printing silenced messages and increase the print limits while printing silenced messages.

Let's consider the same example,

```
$ cat sample
#!/usr/bin/env -Ssml --script
>(* SML code starts here *)
val x = "Hello World x\n";
val () = print x;
$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x
$
```

In the above result, the “val x = "Hello World x\n" : string” is the compiler message printed as part of the variable declaration. And “Hello World x” is the output of the print instruction. One of the goals for our dissertation is to introduce the capability to hide/silence/mute these compiler messages. The changes we have introduced as described in sections 7.3.3 thru 7.3.8 have enabled this capability. By calling the appropriate functions from structure *Mutecompiler* we can achieve the desired results.

Let's consider the same example, now call the `silenceCompiler` function to mute compiler messages,

```
$ cat sample
#!/usr/bin/env -Ssml --script
val _ = Backend.Mutecompiler.silenceCompiler ();
val x = "Hello World x\n";
val () = print x;
```

```

$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
Hello World x
$

```

We can see the compiler message is muted but not the output. The muting takes effect only for the compiler messages that are generated from the point where the function is called.

Let's consider another example and result,

```

$ cat sample
#!/usr/bin/env -Ssml --script
val x = "Hello World x\n";
val () = print x;
val _ = Backend.Mutecompiler.silenceCompiler ();
val y = "Hello World y\n";
val () = print y;
$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x
Hello World y
$

```

In the above results we can see the compiler message generated as part of variable declaration before invoking the `silenceCompiler` function has been printed to the output. The compiler message for the variable declaration after invoking the `silenceCompiler` function has been suppressed or muted.

### 8.2.3 Un-silencing compiler messages

We were able to verify the capability to un-mute the compiler messages by invoking the function `unsilenceCompiler`. This is only required if the compiler messages are muted earlier. This feature gives us the capability to suppress the compiler messages for a selective part of the source code.



Let's consider the example and result,

```
$ cat sample
#!/usr/bin/env -Ssml --script
val x = "Hello World x\n";
val () = print x;
val _ = Backend.Mutecompiler.silenceCompiler ();
val y = "Hello World y\n";
val () = print y;
val _ = Backend.Mutecompiler.unsilenceCompiler ();
val z = "Hello World z\n";
val () = print z;
$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x
Hello World y
val z = "Hello World x\n" : string
Hello World z
$
```

From the results above we can see the compiler messages are suppressed for declaration of variable “y”. The calling of function `silenceCompiler` before variable declaration and then calling function `unsilenceCompiler` after the variable declaration for “y” have suppressed the compiler messages. The compiler messages for declarations before and after the functions were printed.

#### 8.2.4 Printing silenced compiler messages

Silencing the compiler message helps to generate a pretty out but it also suppresses the error messages. In case there is an error in the instructions passed from script file, the details of the error won't be printed out to the user.

Let's consider the example with a forced error in the variable declaration,

```
$ cat sample
#!/usr/bin/env -Ssml --script
```

```

val x == "Hello World x\n";
val () = print x;
$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
./sample:2.27-3.4 Error: syntax error: deleting SEMICOLON VAL

uncaught exception Compile [Compile: "syntax error"]
  raised at: ../compiler/Parse/main/smlfile.sml:19.24-19.46
             ../compiler/TopLevel/interact/evalloop.sml:45.54
             ../compiler/TopLevel/interact/evalloop.sml:306.20-
306.23
$

```

In the above example, we can see the error details called out with the file name (./sample) and specific line number (./sample:2.27-3.4) in the file where the error have occurred.

Let's consider an example where the compiler messages are silenced and an error occurs,

```

$ cat sample
#!/usr/bin/env -Ssml --script
val x = "Hello World x\n";
val () = print x;
val _ = Backend.Mutecompiler.silenceCompiler ();
val y == "Hello World y\n";
val () = print y;
$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x

uncaught exception Compile [Compile: "syntax error"]
  raised at: ../compiler/Parse/main/smlfile.sml:19.24-19.46
             ../compiler/TopLevel/interact/evalloop.sml:45.54
             ../compiler/TopLevel/interact/evalloop.sml:306.20-
306.23
$

```

In the above result we can see the output doesn't have any info on what the error is and where the error is occurring, this is because we have silenced the compiler messages. This is where the new function "printStashedCompilerOutput" comes in to get these error messages printed. The users don't need to call the function or invoke in scripts, this is handled automatically. The above example is only for illustration.

Let's consider an example where the compiler messages are silenced and an error occurs, now the function "printStashedCompilerOutput" is implemented.

```
$ cat sample
#!/usr/bin/env -Ssml --script
val x = "Hello World x\n";
val () = print x;
val _ = Backend.Mutecompiler.silenceCompiler ();
val y == "Hello World y\n";
val () = print y;
$ ./sample
Standard ML of New Jersey (32-bit) vl10.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x

_____

_____
The last 5 lines 30 through 34 of suppressed compiler messages
are:
[library $smlnj/MLRISC/IA32.cm is stable]
[library $SMLNJ-MLRISC/IA32.cm is stable]
[autoloading done]
val it = # : unit
./sample:5.27-6.4 Error: syntax error: deleting SEMICOLON VAL
_____End of suppressed compiler
messages._____

uncaught exception Compile [Compile: "syntax error"]
  raised at: ../compiler/Parse/main/smlfile.sml:19.24-19.46
             ../compiler/TopLevel/interact/evalloop.sml:45.54
             ../compiler/TopLevel/interact/evalloop.sml:306.20-
306.23
```

\$

In the above result we can see the error messages printed to the output. The compiler messages that got stashed are retrieved and printed to the output by the new function. By default, only the last 5 lines of the stashed messages are retrieved and printed. This line limit can be increased or decreased based on user needs.

Let's consider this example with an error, compiler messages silenced, and print line limit increased to 10.

```
$ cat sample
#!/usr/bin/env -Ssml --script
Backend.Mutecompiler.printlnLimit := 10;
val x = "Hello World x\n";
val () = print x;
val _ = Backend.Mutecompiler.silenceCompiler ();
val y = "Hello World y\n";
val () = print 3;
$ ./sample
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x

-----

The last 10 lines 30 through 39 of suppressed compiler messages
are:
[library $smlnj/MLRISC/IA32.cm is stable]
[library $SMLNJ-MLRISC/IA32.cm is stable]
[autoloading done]
val it = # : unit
val y = # : string
./sample:7.5-7.17 Error: operator and operand do not agree
[overload - bad instantiation]
  operator domain: string
  operand:         `Z[INT]
  in expression:
    <exp>
```

```

_____End of suppressed compiler
messages._____

uncaught exception Error
  raised at: ../compiler/TopLevel/interact/evalloop.sml:69.13-
69.21
          ../compiler/TopLevel/interact/evalloop.sml:45.54
          ../compiler/TopLevel/interact/evalloop.sml:306.20-
306.23$
$

```

From the result above, we can see the last 10 lines of the stashed compiler messages are printed.

### 8.2.5 Restoring printing limits

From the examples in the previous section, you may notice the compiler message for variable declarations will differ from the compiler messages that are stashed. This is because we have reset the printing limits to avoid saving all the variable information in the memory.

```

val x = "Hello World x\n" : string
(vs)
val y = # : string

```

This information can be restored by pre-setting the printing limits. Let's see the sample example with printing limits restored.

```

$ cat sample
#!/usr/bin/env -Ssml --script
Backend.Mutecompiler.printlnLimit := 10;
val x = "Hello World x\n";
val () = print x;
val _ = Backend.Mutecompiler.silenceCompiler ();
val _ = Backend.Mutecompiler.restorePrintingLimits ();
val y = "Hello World y\n";
val () = print 3;
$ ./sample

```

```
Standard ML of New Jersey (32-bit) vl10.99.3 [built: Sun Apr 23
20:39:28 2023]
val x = "Hello World x\n" : string
Hello World x
```

---

The last 10 lines 30 through 39 of suppressed compiler messages are:

```
[library $smlnj/MLRISC/IA32.cm is stable]
[library $SMLNJ-MLRISC/IA32.cm is stable]
[autoloading done]
val it = # : unit
val y = "Hello World y\n" : string
./sample:7.5-7.17 Error: operator and operand do not agree
[overload - bad instantiation]
  operator domain: string
  operand:         "Z[INT]
  in expression:
    <exp>
```

```
_____End of suppressed compiler
messages._____
```

```
uncaught exception Error
  raised at: ../compiler/TopLevel/interact/evalloop.sml:69.13-
69.21
          ../compiler/TopLevel/interact/evalloop.sml:45.54
          ../compiler/TopLevel/interact/evalloop.sml:306.20-
306.23$
$
```

In the results above you can see the compiler message for variable declaration of “y” saved in stash without any change.

### 8.3 Sample Scripts

Here are some sample scripts and their results.

Example #1:

Script:

```
$ cat sample1
#!/usr/bin/env -Ssml --script
```

```

; (*--SML--*)
fun cube x = x * x * x ;
fun square x = x * x ;
val () = print (Int.toString(square(9)) );
val () = print "\n";
val () = print (Int.toString(cube(3)) );
val () = print "\n";
print (Int.toString(cube(10)) );
$

```

**Result:**

```

$ ./sample1
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun
Apr 23 20:39:28 2023]
val cube = fn : int -> int
val square = fn : int -> int
[autoloading]
[autoloading done]
81
27
1000val it = () : unit
$

```

**Example #2:**

**Script:**

```

$ cat sample2
#!/usr/bin/env -Ssml --script
; (*--SML--*)
val _ = Backend.Mutecompiler.silenceCompiler ();
fun cube x = x * x * x ;
fun square x = x * x ;
val () = print (Int.toString(square(9)) );
val () = print "\n";
val () = print (Int.toString(cube(3)) );
val () = print "\n";
print (Int.toString(cube(10)) );
$

```

**Result:**

```

$ ./sample2
Standard ML of New Jersey (32-bit) v110.99.3 [built: Sun
Apr 23 20:39:28 2023]

```

81  
27  
1000\$

#### 8.4 GitHub Test Results

We were able to successfully run the test workflows set up for Ubuntu and MacOS environments. Screen prints are available in Appendix. For testing, we dynamically created a test script file in the environment thru YAML scripts and the test script is executed inside the environment. This can also be done by having a test script in the repository itself and changing the script to give execute permissions and execute.



## 9. Conclusion and Future Work

### 9.1 Conclusion

This dissertation intended to make SML programming simpler for developers, programmers, and students such that they can write SML source code in a file and execute it as a script. And to integrate *smlnj-script* capabilities into SML/NJ and do away with the requirement to maintain a separate *smlnj-script* implementation.

By analysing the SML/NJ installation procedures, scripts, and logs, we were able to understand how the SML/NJ is built on both C and SML language codebase. Every time SML is loaded, the C executables load the SML heap image which is built as part of the installation process and takes in command line arguments to perform desired functions. This was the entry point of our dissertation. We intercepted the modules which received the command line arguments and identified the right place to add our changes so that the script execution will be recognised, and the instructions are passed through evaluation loop.

We were able to implement one of the key *smlnj-script* capability, silencing the compiler, which is to suppress the compiler output messages. We introduced a new structure *Mutecompiler* with built-in functions which gave the capability to silence and un-silence compiler messages, and to restore the suppressed messages for printing. The results discussed in section 8 is a proof that this dissertation achieved the core objective and a high-level requirement.

### 9.2 Future Work

Though we implemented the main objective and few key features in this dissertation, there are few other nice to have features available in *smlnj-script* (described in section 2.8) that can be merged on to SML/NJ. It needs a considerable amount of time to analyse the SML/NJ code base, for example string overloading is like complete overhaul of string handling and due diligence is needed to analyse every usage of it in SML. The same goes to other features in *smlnj-script*. With the support and guidance from subject matter experts, appropriate modules can be identified, changes built and merged back into SML/NJ.

### 9.3 Lessons Learnt

In this section, we'll discuss a few of the things we learned while conducting our research and writing this dissertation. The author had no prior knowledge of SML and

*smlnj-scripts* and just a fundamental understanding of UNIX and shell scripts prior to writing this dissertation and research paper.

To begin with, the author must grasp the fundamentals of the Unix and Linux operating systems, as well as what shell scripts are, how they operate, how to build scripts, what permissions they require, what happens when they are executed, and how to send inputs from the script to the interpreter. To use to search the SML libraries for study, the author must fully master Unix utilities and command line arguments such "Find," "GREP," and "XARGS." It was essential to comprehend and troubleshoot SML installation scripts as well as scripts that do recompile, build, and reinstall to learn how to execute scripts with logs enabled using the "set -vx" options.

The author has to become familiar with the effects of installing an SML. SML can be installed manually, through package installation, or through home brew. Updated source code, recompilation, and reinstallation all require source code support, which is only available during manual installation. The core layer of SML, which consists of C executables and SML programmes that are built on top of C binaries, must be thoroughly understood by the author before they can understand the installation process. To fully understand the Compilation Library Manager of SML, extensive research was required.

The chance to study CI/CD, actions, and workflows in GitHub was opened up by this dissertation because the author has very little experience using GitHub. This made it easier to put up an entirely automated test that could be run on several operating systems. Early on in the dissertation, there was a development restriction that limited the test to macOS 10.13.6; however, now that the full test configuration has been completed in GitHub, the change can be tested in any acceptable environment. Utilising GitHub made it simple to communicate the change with SML/NJ implementors and provided a straightforward method for merging the change into SML/NJ.

## 10. Copyright and Licensing

I, Dayanandan Natarajan, the research and document's author, grant permission under open-source licencing agreements for my changes to be incorporated into SML/NJ by its implementors.

A separate licensing needs to be agreed with Dr Joe Wells and Heriot Watt University for the features derived from *smlnj-script*.

Section 7.4 of this document has called out the changes that are author's original idea, research and development, and details of the changes that are derived from *smlnj-script*.

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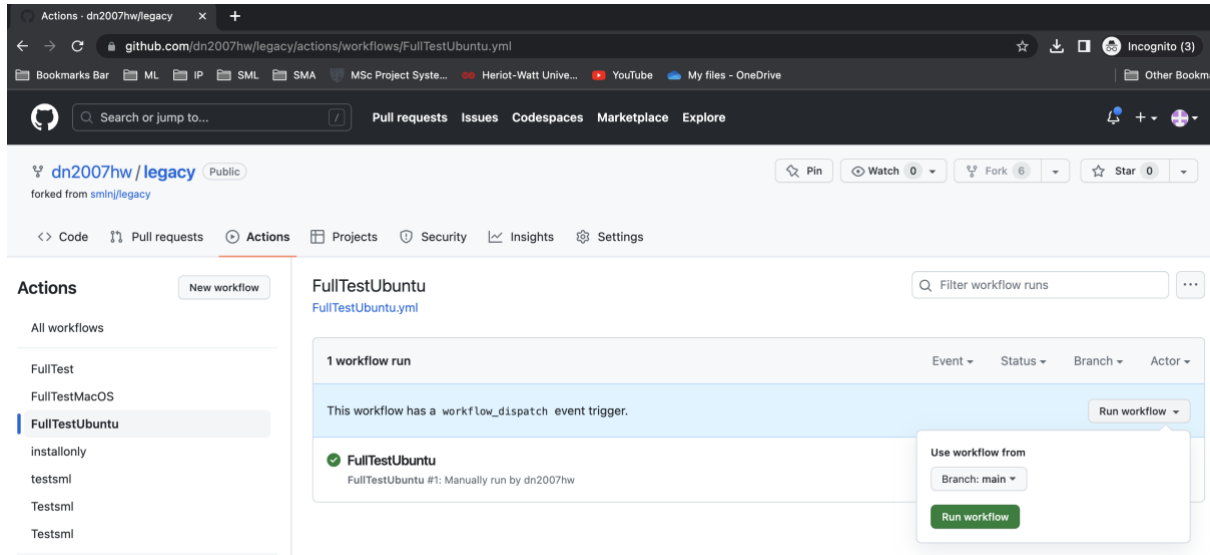
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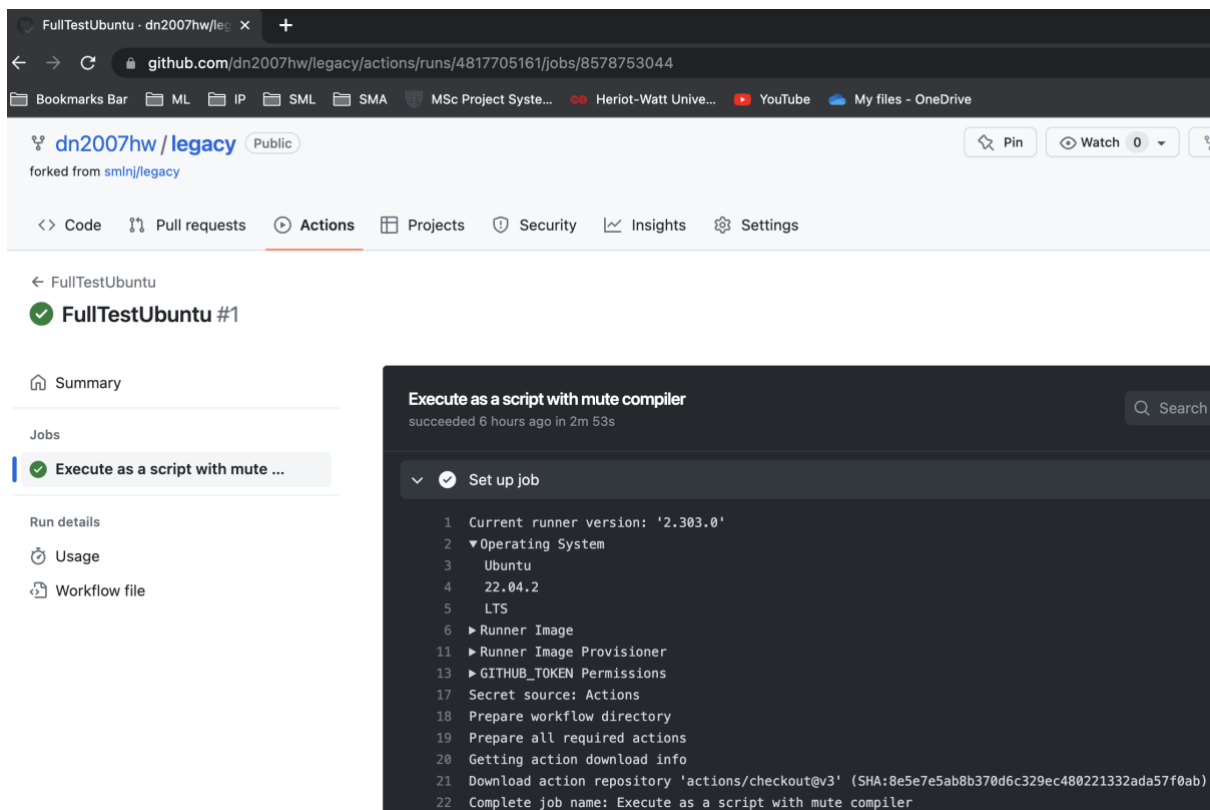
## Appendix A

### GitHub Test Screen prints

#### Running a workflow to test in Ubuntu Linux



#### Results from Ubuntu Linux



FullTestUbuntu · dn2007hw/leg
+

github.com/dn2007hw/legacy/actions/runs/4817705161/jobs/8578753044

Bookmarks Bar
ML
IP
SML
SMA
MSc Project Syste...
Heriot-Watt Unive...
YouTube
My files - OneDrive

### Summary

#### Jobs

Execute as a script with mute ...

#### Run details

Usage
Workflow file

## Execute as a script with mute compiler

succeeded 6 hours ago in 2m 53s

>
Install SML (ubuntu)

>
Recompile - CMB Make (ubuntu)

>
Recompile - makeml (ubuntu)

>
Recompile - installml (ubuntu)

>
Re-Install SML (ubuntu)

Create test script

```

1  Run cat > test-script <<EOF
2  cat > test-script <<EOF
3  #!/usr/bin/env -Ssml --script
4  val x = "Hello World x\n";
5  val () = print x;
6  val _ = Backend.Mutecompiler.silenceCompiler ();
7  val y = "Hello World y\n";
8  val () = print y;
9  val () = print y;
10 EOF
11 shell: /usr/bin/bash -e {0}

```

Execute test script

```

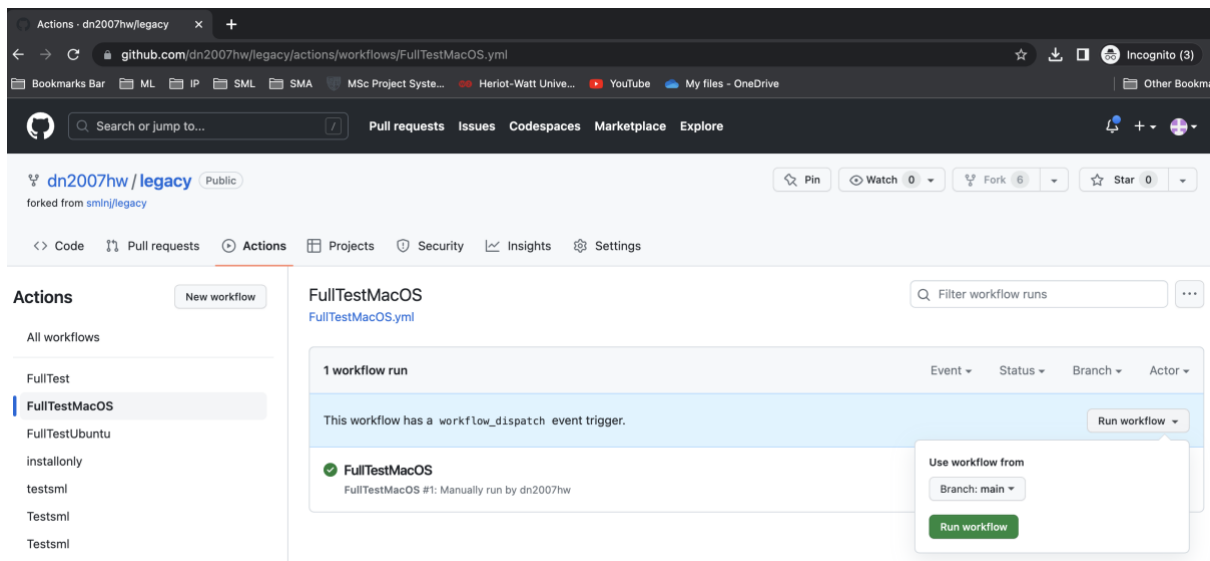
1  Run chmod a+x test-script
2  chmod a+x test-script
3  export PATH="/home/runner/work/legacy/legacy/bin:$PATH"
4  ./test-script
5  shell: /usr/bin/bash -e {0}
6  Standard ML of New Jersey (64-bit) v110.99.3 [built: Thu Apr 27 08:35:14 2023]
7  val x = "Hello World x\n" : string
8  Hello World x
9  Hello World y
10 Hello World y

```

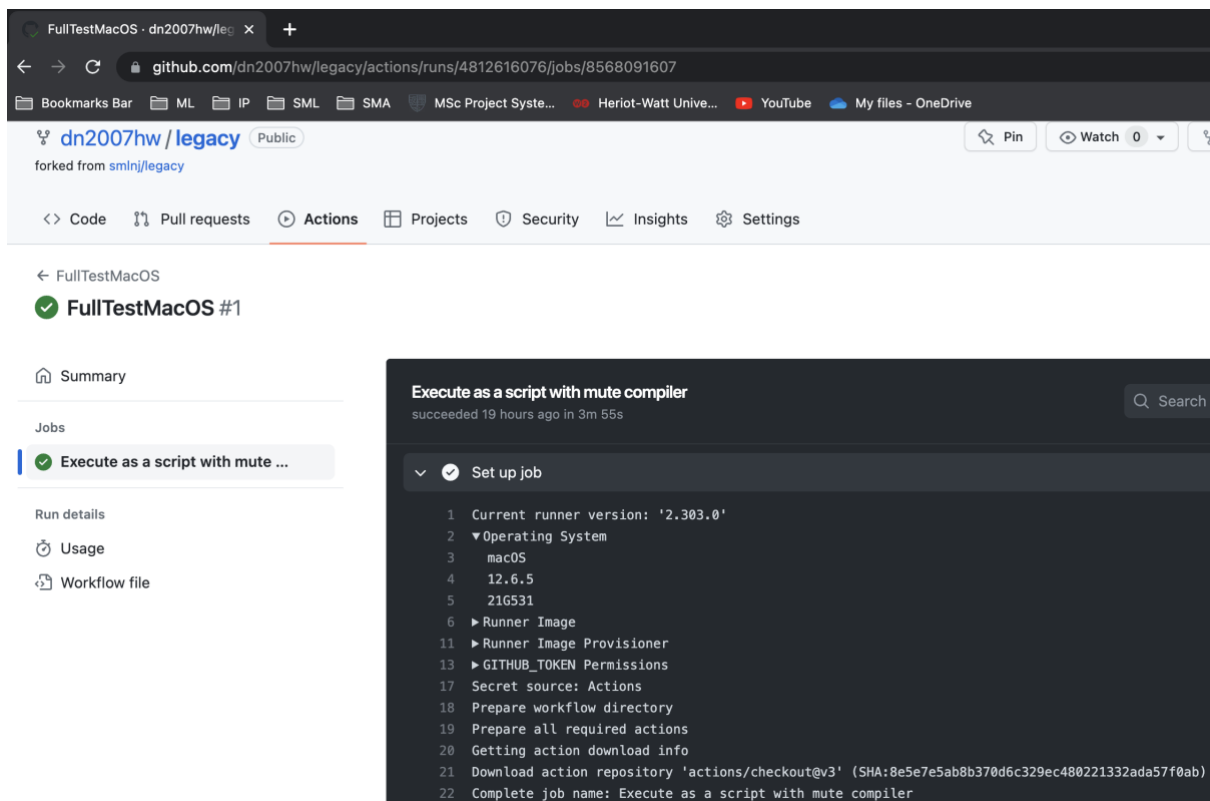
>
Post Run actions/checkout@v3

>
Complete job

## Running a workflow to test in MacOS



## Results from MacOS





FullTestMacOS · dn2007hw/leg x +
github.com/dn2007hw/legacy/actions/runs/4812616076/jobs/8568091607
Bookmarks Bar ML IP SML SMA MSc Project Syste... Heriot-Watt Unive... YouTube My files - OneDrive

### Summary

Jobs

- Execute as a script with mute ...

Run details

Usage

Workflow file

### Execute as a script with mute compiler

succeeded 19 hours ago in 3m 55s

- Run actions/checkout@v3
- Install SML (macos)
- Recompile - CMB Make (macos)
- Recompile - makeml (macos)
- Recompile - installml (macos)
- Re-Install SML (macos)
- Create test script

```

1 Run cat > test-script <<EOF
2 cat > test-script <<EOF
3 #!/usr/bin/env -Ssml --script
4 val x = "Hello World x\n";
5 val () = print x;
6 val _ = Backend.Mutecompiler.silenceCompiler ();
7 val y = "Hello World y\n";
8 val () = print y;
9 val () = print y;
10 EOF
11 shell: /bin/bash -e {}

```
- Execute test script

```

1 Run chmod a+x test-script
2 chmod a+x test-script
3 export PATH="/Users/runner/work/legacy/legacy/bin:$PATH"
4 ./test-script
5 shell: /bin/bash -e {}
6 Standard ML of New Jersey (64-bit) v110.99.3 [built: Wed Apr 26 19:33:58 2023]
7 val x = "Hello World x\n" : string
8 Hello World x
9 Hello World y
10 Hello World y

```
- Post Run actions/checkout@v3
- Complete job