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**RV COLLEGE OF ENGINEERING®**  
**(An Autonomous Institution Affiliated to VTU)**  
**V Semester B. E. Examinations April/May -2024**  
**Computer Science and Engineering**  
**THEORY OF COMPUTATION**

Time: 03 Hours

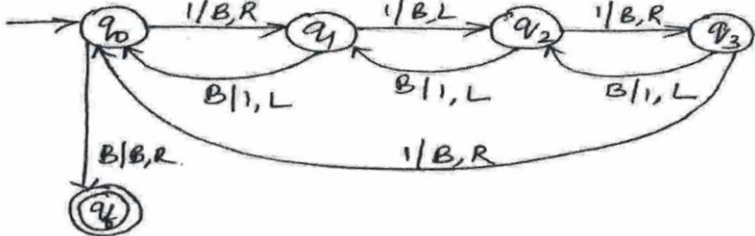
Maximum Marks: 100

Instructions to candidates:

1. Answer all questions from Part A. Part A questions should be answered in first three pages of the answer book only.
2. Answer FIVE full questions from Part B. In Part B question number 2 is compulsory. Answer any one full question from 3 and 4, 5 and 6, 7 and 8, 9 and 10.

**PART-A**

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1	1.1	If the DFA $A1 = (\{q_0, q_1, q_2\}, \Sigma, \delta, \{q_0, q_1\}, q_0)$ and DFA $A2 = (\{p_0, p_1, p_2, p_3\}, \Sigma, \delta, \{p_3\}, p_0)$ , then DFA of $(L(A1) \cap L(A2))^R$ has _____ final state(s).	01	1	3
	1.2	Let $x$ be a positive integer represented using unary notation. Identify the function $f(x)$ computed by Turing machine shown in Fig 1.2. Note $x \in 1^*$ .			
		 <p style="text-align: center;">Fig 1.2</p>	01	3	4
	1.3	Consider the machine $M = (Q = \{q_0, q_1, q_2, q_3, q_4\}, \Sigma = \{a_1, a_2, a_3\}, \Gamma = \{a_1, a_2, a_3\}, \delta, q_0, B, F = \{q_4\})$ . A universal Turing machine Mu is an automaton that, given as input the description of any Turing machine $M$ and a string $w$ , can simulate the computation of $M$ on $w$ . For the transition $\delta(q_1, a_1) = (q_2, a_3, R)$ , the encoding is _____.	01	2	4
	1.4	Construct the grammar or the given regular expression $a(a + b)^* b$ .	02	2	3
	1.5	Consider the Turing machine $M = (Q = \{q_0, q_1\}, \Sigma = \{a, b\}, \Gamma = \{A, B\}, \delta, q_0, B, F = \{q\})$ with one of the transition ( $\delta$ ) defined by $\delta(q, a, b) = (q, A, B, L)$ . Then the Turing machine is a _____.	01	1	1
	1.6	Identify the nullable variable and eliminate all $\lambda$ -productions from $S \rightarrow AaB \mid aaB, A \rightarrow \lambda, B \rightarrow bbA \mid \lambda$ .	02	3	3
	1.7	Define GNF. Convert the grammar $S \rightarrow abSb \mid aa$ into GNF.	02	2	3
	1.8	State any one decidable property of context free language.	01	1	1
	1.9	Consider the regular expression $ab^*(a + b)$ . Construct the finite automata for this regular expression.	02	2	2
	1.10	Show that the following grammar is ambiguous: $S \rightarrow AB \mid aaB, A \rightarrow a \mid Aa, B \rightarrow b$	02	3	3
	1.11	Distinguish between recursive and recursively enumerable language.	02	2	2

1.12	Find the $\epsilon$ - <i>CLOSURE</i> (8). Refer Fig 1.12.			
1.13	Define <i>DPDA</i> .	01	1	2
		02	2	3

### PART-B

2	a	Show that the language $L = \{awa : w \in \{a,b\}^*\}$ is regular using deterministic automata.	03	2	3
	b	Convert the following <i>NFA</i> to <i>DFA</i> and minimize the converted <i>DFA</i> .			
		Fig 2b	08	3	2
	c	Discuss the algebraic properties of regular expressions.	05	2	1
3	a	Define <i>CFG</i> . Construct <i>CFG</i> for the following languages.			
		i) $L = \{a^i b^j c^k : i \neq j\}$			
		ii) $L = \{a^i b^j c^k : j = i + k\}$	04	2	2
	b	Define left linear grammar. Find all left linear grammar for the language accepted by the finite automata in Fig 3b.			
		Fig 3b	07	3	3
	c	Transform the grammar with productions into Chomsky normal form: $S \rightarrow abAB$ , $A \rightarrow bAB   \lambda$ , $B \rightarrow BAa   A   \lambda$	05	3	4
		<b>OR</b>			
4	a	Construct a finite automaton that accepts the language generated by the right linear grammar.			
		$S \rightarrow cA   baS$ , $A \rightarrow bB   aC   \epsilon$ , $B \rightarrow aA   bbC   \epsilon$ , $C \rightarrow abA   baC$	04	2	3
	b	Let $M_1$ and $M_2$ be the FA's as shown in Fig 4b, accepting languages $L_1$ and $L_2$ respectively. Draw FA's accepting the following languages.			
		i) $L_1 \cup L_2$			
		ii) $L_2 - L_1$			
		iii) $L_1^R$			
		iv) $\overline{L_1}$			
		Fig 4b	08	3	2

c		Eliminate the useless productions form the productions listed below and justify your answer. $S \rightarrow c aA B C, A \rightarrow aB \lambda, B \rightarrow Aa, C \rightarrow cCD, D \rightarrow ddd$	04	3	3
5	a	State and prove pumping lemma for <i>CFL's</i> . Show that the language $L = \{a^n b^n c^n : n \geq 0\}$ is not context free.	08	2	2
	b	Construct an <i>NPDA</i> that accepts the language $L = \{a^n b^m : n \geq 0, n \neq m\}$ . Give the graphical representation for <i>PDA</i> obtained. Show the moves made by the <i>PDA</i> for the string <i>aabbbb</i> .	08	3	3
		<b>OR</b>			
6	a	Show that <i>CFL's</i> are closed under union and concatenation.	04	2	1
	b	Give an algorithm to convert <i>PDA</i> to <i>CFG</i> . Convert the following Push Down Automata transitions to <i>CFG</i> . $\delta(q_0, a, z) = \{(q_0, A_z)\}$ $\delta(q_0, a, A) = \{(q_0, A)\}$ $\delta(q_0, b, A) = \{(q_1, \lambda)\}$ $\delta(q_0, \lambda, z) = \{(q_2, \lambda)\}$	08	3	2
	c	Discuss and elaborate on the languages accepted by <i>PDA</i> .	04	3	1
7	a	Define Turing machine. Let <i>x</i> and <i>y</i> are two positive integers represented using unary notation. Design a Turing machine that computes the function. i) $f(x, y) = x + y$ where $W(x), W(y) \in 1^+$ and $q \circ W(x) \circ W(y) \vdash qfW(x + y)0$ ii) $f(x) = x \bmod 4, x \in 1^+$	07	4	4
	b	Write a note on the following: i) Semi-infinite Tape Turing machine ii) Multi-stack Turing machine iii) Offline Turing machine.	09	2	1
		<b>OR</b>			
8	a	Design Turing machine to compute the function $f(w) = w^R$ where $w \in \{a, b\}^*$ . Using instantaneous descriptions, show operation on $w = abb$ .	06	3	4
	b	Define multitape <i>TM</i> . Explain how multitape <i>TM</i> can be simulated using single tape <i>TM</i> .	06	2	1
	c	If $L_1$ and $L_2$ are recursively enumerable languages over $\Sigma$ , then prove that $L_1 \cap L_2$ and $L_1 \cup L_2$ are recursively enumerable.	04	1	2
9	a	Define post correspondence problem. Identify whether it is possible to solve the <i>PCP</i> given below. i) $A = \{1, 10111, 10\}$ and $B = \{111, 10, 0\}$ ii) $A = \{100, 101, 110\}$ and $B = \{10, 01, 1010\}$	08	3	4
	b	With the help of a neat diagram, discuss the relationship between the families of languages.	08	1	2
		<b>OR</b>			
10	a	Define context sensitive grammar. Design <i>CSG</i> for the language $L = \{a^n b^n c^n   n \geq 1\}$ . Derive the string $a^3 b^3 c^3$ from the constructed <i>CSG</i> .	08	3	2
	b	Discuss the following: i) Halting problem of a <i>TM</i> ii) Unrestricted grammar.	08	2	1