

Mathematics 2, part 2, lecture 3

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Extra optimization topics

- ▶ Nelder-Mead method
- ▶ Local search, case study
- ▶ Interior point methods, pushing the boundaries

Nelder - Mead (simplex) method

- heuristic method performing search in n -dim space
- samples $n+1$ points (forming a simplex) and produces a new sample of $n+1$ points (sometimes replacing a single point)
- zero-order method
- **can** crawl out of local extrema
- a.k.a. amoeba method

What is a simplex?

Simplex (n -simplex) is a convex hull of $n+1$ affinely independent points.

$x_0, x_1, x_2, \dots, x_n$ are affinely independent

How to construct a simplex with equal edge lengths?

$$x_0, x_0 + (1, 0, \dots), x_0 + (0, 1, 0, \dots), \dots, x_0 + (0, \dots, 0, 1)$$

Does this work?

Nelder-Mead notation

$f \leftarrow$ cost function

minimization

$x_0, x_1, \dots, x_n \leftarrow$ sample points

$y_0 \leq y_1 \leq \dots \leq y_n \leftarrow$ f -values

$$y_i = f(x_i)$$

$x_0, x_{n-1}, x_n \leftarrow$ best, lousy, worst points

How to produce next $n+1$ sample points?

① replace x_n with a better point x_{NEXT} , keep x_0, \dots, x_{n-1}

OR

② only keep x_0 and replace all the rest

Nelder-read operations

Reflect

Expand

Contract (outside)
(inside)

Shrink

Nelder-Mead termination criteria

- points x_0, x_1, \dots, x_n sufficiently close
- values y_0, y_1, \dots, y_n sufficiently close
- combination of above

Local search

- heuristic method for solving (hard) optimization problems

$\Omega \leftarrow$ search space / set of feasible solutions

$f \leftarrow$ cost function minimization

$x \in \Omega \leftarrow$ feasible solution

$N(x) \subseteq \Omega \leftarrow$ neighborhood of x

$x_0, x_1, x_2, x_3, \dots, x_i, \dots$

$\forall i \quad f(x_{i+1}) < f(x_i)$

$\forall i \quad x_{i+1} \in N(x_i)$

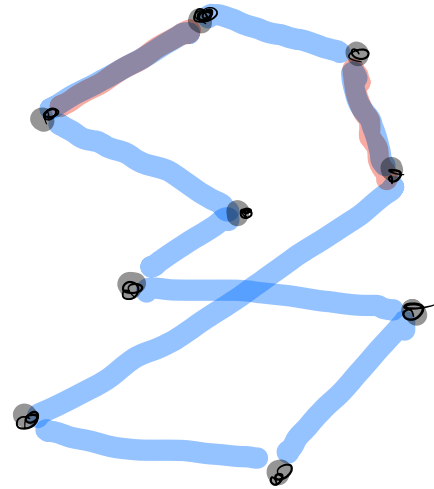
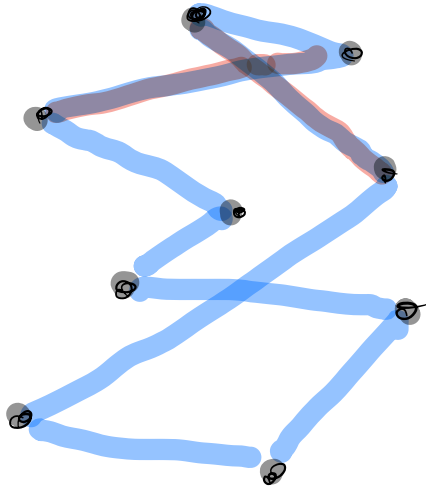
Local search - example

Travelling salesman problem TSP (geometric version)

input: p_1, p_2, \dots, p_n points in \mathbb{R}^2

output: shortest closed walk containing all above points
(it should be a Hamiltonian cycle)

Neighborhood



also called a 2-opt move

Proposition

If TSTour T is 2-optimal
then T is without crossings.

Local search can get stuck. What now?

- simulated annealing
allow bad steps if temperature is high
- tabu search
do not reverse recent steps
- jump moves
allow bad steps sometimes

Interior point methods

- not limited to LP, can be used in other convex opt. problems

Quadratic prog.

$$\min x^T Q x + c^T x$$

$$Ax = b$$

$Q \succeq 0$ \leftarrow Q PSD matrix

Semidefinite prog.

$$\min c^T x$$

$$Ax = b$$

$$x \succeq 0$$

$X \succeq 0$ \nearrow X PSD matrix

Conic opt.

$$\min c^T x$$

$$Ax = b$$

$$x \in K$$

\nearrow cone



