

# A DSL for Specifying and Model Checking Mobile Ring Robot Algorithms

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## 1 Appendix

This appendix gives fully the theory including proofs of the propositions and lemmas used in the paper.

### 1.1 Sequences

Let  $\text{Elt}$  be the set of (concrete) elements,  $\text{EV}$  be the set of element variables and  $\text{SV}$  be the set of sequence variables.

**Definition 1.1** (Sequence Patterns). The set  $\text{SP}$  of sequence patterns are inductively defined as follows:

1.  $\varepsilon \in \text{SP}$  (the empty sequence);
2. For each element  $e \in \text{Elt}$ ,  $e \in \text{SP}$ ;
3. For each element variable  $E \in \text{EV}$ ,  $E \in \text{SP}$ ;
4. For each sequence variable  $S \in \text{SV}$ ,  $S \in \text{SP}$ ;
5. For any sequence patterns  $SP_1, SP_2 \in \text{SP}$ ,  $SP_1 SP_2 \in \text{SP}$ .

The binary juxtaposition operator used in  $SP_1 SP_2$  is associative, namely that  $(SP_1 SP_2) SP_3 = SP_1 (SP_2 SP_3)$  for any sequence patterns  $SP_1, SP_2, SP_3 \in \text{SP}$ .  $\varepsilon$  is an identity of the binary juxtaposition operator, namely that  $\varepsilon SP = SP$  and  $SP \varepsilon = SP$  for any sequence patterns  $SP \in \text{SP}$ .

Sequence patterns that do not have any variables at all are called sequences. Let  $\text{Seq} \subseteq \text{SP}$  be the set of all sequences.

A substitution  $\sigma$  is a function from the disjoint union  $\text{EV} \uplus \text{SV}$  of  $\text{EV}$  and  $\text{SV}$  to the disjoint union  $\text{Seq} \uplus \text{EV} \uplus \text{SV}$  of  $\text{Seq}$ ,  $\text{EV}$  and  $\text{SV}$ . For  $E \in \text{EV}$ ,  $\sigma(E)$  is an element  $e \in \text{Elt}$  or  $E$  and for  $S \in \text{SV}$ ,  $\sigma(S)$  is a sequence  $seq \in \text{Seq}$  or  $S$ . The domain of a substitution  $\sigma$  can be naturally extended to  $\text{SP}$  such that  $\sigma(\varepsilon)$  is  $\varepsilon$ , for an element  $e \in \text{Elt}$ ,  $\sigma(e)$  is  $e$  and for a sequence pattern  $SP_1, SP_2 \in \text{SP}$ ,  $\sigma(SP_1 SP_2)$  is  $\sigma(SP_1) \sigma(SP_2)$ .

**Definition 1.2** (Sequence pattern match). Pattern match between  $sp \in \text{SP}$  &  $seq \in \text{Seq}$  is to find all substitutions  $\sigma$  such that  $\sigma(sp) = seq$ . Let  $sp \models seq$  be the set of all such substitutions.

Elements, element variables and sequence variables used in sequence patterns are called components in the sequence patterns. For  $sp \in \text{SP}$ , let  $|sp|$  be the number of components in it.  $sp \in \text{SP}$  can be in the form  $ES_1 ES_2 \dots ES_{|sp|}$ , where each  $ES_i$  is an element, an element variable or a sequence variable. For  $sp \in \text{SP}$ , let  $sp(i)$ , where  $i \in \{1, 2, \dots, |sp|\}$ , be the  $i$ th element  $ES_i$  in  $sp$ . Let  $e, e_1, e_2, \dots \in \text{Elt}$ ,  $E, E_1, E_2, \dots \in \text{EV}$ ,  $S, S_1, S_2, \dots \in \text{SV}$  and  $ES, ES_1, ES_2, \dots \in \text{Elt} \uplus \text{EV} \uplus \text{SV}$ . A binary construct  $\text{sv}(S, I)$  that is not in  $\text{SV}$ , where  $S$  is a sequence variable and  $I$  is either 0 or 1, is used as an extra

sequence variable. Let  $\text{SSV}$  be  $\{\text{sv}(S, I) \mid S \in \text{SV}, I \in \{0, 1\}\}$ .  $\text{SV} \cup \text{SSV}$  may be used as the set of sequence variables instead of  $\text{SV}$ .

**Definition 1.3** (Split sequence patterns). For  $sp \in \text{SP}$ ,  $\text{split}(sp)$  is a sequence pattern such that each sequence variable  $S$  in  $sp$  is replaced with  $\text{sv}(S, 0) \text{sv}(S, 1)$ .  $\text{split}(\varepsilon) = \varepsilon$ ,  $\text{split}(e) = e$  for  $e \in \text{Elt}$ ,  $\text{split}(E) = E$  for  $E \in \text{EV}$ ,  $\text{split}(S) = \text{sv}(S, 0) \text{sv}(S, 1)$  for  $S \in \text{SV}$  and  $\text{split}(SP_1 SP_2) = \text{split}(SP_1) \text{split}(SP_2)$  for  $SP_1, SP_2 \in \text{SP}$ .

**Definition 1.4** (Joining split sequence variables). For  $sp \in \text{SP}$  and  $seq \in \text{Seq}$ , let  $\sigma$  be in  $(\text{split}(sp) \models seq)$ .  $\text{join}(\sigma)$  is the substitution  $\sigma'$  such that for each sequence variable  $S$  in  $sp$   $\sigma'(S) = \sigma(\text{sv}(S, 0)) \sigma(\text{sv}(S, 1))$  and for any other variables  $X$   $\sigma'(X) = \sigma(X)$ . The domain of  $\text{join}$  can be naturally extended to the set of substitutions such that  $\text{join}(\text{split}(sp) \models seq)$  is  $\{\text{join}(\sigma) \mid \sigma \in (\text{split}(sp) \models seq)\}$ .

$\text{rtt}$  (that stands for rotate) takes a sequence pattern  $sp$  and returns the sequence pattern obtained by rotating  $sp$  clockwise.  $\text{rev}$  (that stands for reverse) takes a sequence pattern  $sp$  and returns the sequence pattern obtained by reversing  $sp$ . Let  $sp$  be  $ES_1 ES_2 \dots ES_{|sp|-1} ES_{|sp|}$ .  $\text{rtt}(sp) = ES_{|sp|} ES_1 ES_2 \dots ES_{|sp|-1}$  and  $\text{rev}(sp) = ES_{|sp|} ES_{|sp|-1} \dots ES_2 ES_1$ . Let us suppose that a subscript  $exp$  of  $ES_{exp}$  used as an element in  $sp$  is interpreted as  $(exp \bmod |sp|) + 1$ .

**Definition 1.5** (Reversing substitutions).  $\sigma_{\text{rev}}$  is defined as follows: for an element  $e \in \text{Elt}$   $\sigma_{\text{rev}}(e) = \sigma(e) = e$ , for an element variable  $E \in \text{EV}$   $\sigma_{\text{rev}}(E) = \sigma(E)$  and for a sequence variable  $S \in \text{SV}$   $\sigma_{\text{rev}}(S) = \text{rev}(\sigma(S))$ .

Given two sequence patterns  $sp, sp' \in \text{SP}$ ,  $sp \cup sp'$  holds if there exists a natural number  $n$  such that  $sp = \text{rtt}^n(sp')$ , namely that  $sp$  is obtained by rotating  $sp'$  finitely many times;  $sp \cup sp'$  holds if there exists a natural number  $n$  such that  $sp = \text{rtt}^n(\text{rev}(sp'))$ , namely that  $sp$  is obtained by reversing  $sp'$  once and rotating it finitely many times. For example, let  $sp$  and  $sp'$  be  $e_1 S_1 e_2$  and  $e_1 e_2 S_1 (= \text{rtt}(sp))$  and then  $(sp \cup sp')$  holds, while  $(sp \cup sp')$  does not; let  $sp$  and  $sp'$  be  $e_1 S_1 e_2$  and  $e_2 e_1 S_1 (= \text{rtt}(\text{rev}(sp)))$  and then  $(sp \cup sp')$  does not, while  $(sp \cup sp')$  holds; let  $sp$  and  $sp'$  be  $e_1 S_1 e_2$  and  $e_2 e_1 S_1 (= \text{rtt}(\text{rev}(sp)))$  and then both  $(sp \cup sp')$  and  $(sp \cup sp')$  hold.

**Definition 1.6** (Sequence pattern contexts). A sequence pattern context is a sequence pattern  $sp \in \text{SP}$  in which one component (say,  $i$ th component, where  $1 \leq i \leq |sp|$ ) is replaced with a special symbol  $\square$  called a hole, denoted  $sp_{(i)}\{\square\}$ . A

hole  $\square$  is treated as an element. Let  $sp$  be  $ES_1 \dots ES_i \dots ES_{|sp|}$  and then  $sp_{(i)}\{\square\}$  is  $ES_1 \dots \square \dots ES_{|sp|}$ .

For a sequence pattern or a sequence pattern context  $spc$  and a sequence pattern  $sp$ ,  $spc_{(i)}\{sp\}$  is  $spc$  in which the  $i$ th component in  $spc$  is replaced with  $sp$ .  $(sp_{(i)}\{\square\})_{(i)}\{sp(i)\} = sp_{(i)}\{sp(i)\} = sp$ .

**Definition 1.7** (Correspondent components). Let  $sp \in \mathbf{SP}$  be  $ES_1 \dots ES_{i-1} ES_i ES_{i+1} \dots ES_{|sp|}$ , where  $1 \leq i \leq |sp|$  and  $sp' \in \mathbf{SP}$  be  $ES'_1 \dots ES'_{j-1} ES'_j ES'_{j+1} \dots ES'_{|sp'|}$ , where  $1 \leq j \leq |sp'|$ . If  $ES_i ES_{i+1} \dots ES_{|sp|} ES_1 \dots ES_{i-1} = ES'_j ES'_{j+1} \dots ES'_{|sp'|} ES'_1 \dots ES'_{j-1}$  or  $ES_i ES_{i+1} \dots ES_{|sp|} ES_1 \dots ES_{i-1} = ES'_j ES'_{j-1} \dots ES'_1 ES'_{|sp'|} \dots ES'_{j+1}$ , then  $ES'_j$  is the corresponding component in  $sp'$  to  $ES_i$  in  $sp$ .

**Proposition 1.8.** For any sequence patterns  $sp, sp' \in \mathbf{SP}$  and any natural numbers  $i \in \{1, \dots, |sp|\}$  and  $j \in \{1, \dots, |sp'|\}$  such that the  $j$ th component  $sp'(j)$  in  $sp'$  is the corresponding component in  $sp'$  to  $sp(i)$  in  $sp$ , (1)  $(sp \cup sp') \Leftrightarrow sp_{(i)}\{\square\} \cup sp'_{(j)}\{\square\}$  and (2)  $(sp \cup sp') \Leftrightarrow sp_{(i)}\{\square\} \cup sp'_{(j)}\{\square\}$ .

*Proof.* (1)  $(\Rightarrow)$  There exists a natural number  $m$  such that  $\text{rtt}^m(sp)$  is  $sp(i) sp(i+1) \dots sp(i-1)$  and there exists a natural number  $n$  such that  $\text{rtt}^n(sp')$  is  $sp'(j) sp'(j+1) \dots sp'(j-1)$ . Because of  $sp \cup sp'$ ,  $\text{rtt}^m(sp) = \text{rtt}^n(sp')$  and then  $(\text{rtt}^m(sp))_{(1)}\{\square\} = (\text{rtt}^n(sp'))_{(1)}\{\square\}$ .  $\text{rtt}^{-m}((\text{rtt}^m(sp))_{(1)}\{\square\}) = sp_{(i)}\{\square\}$  and  $\text{rtt}^{-n}((\text{rtt}^n(sp'))_{(1)}\{\square\}) = sp'_{(j)}\{\square\}$ . Therefore,  $sp_{(i)}\{\square\} \cup sp'_{(j)}\{\square\}$ .  $(\Leftarrow)$  There exists a natural number  $m$  such that  $\text{rtt}^m(sp_{(i)}\{\square\})$  is  $\square sp(i+1) \dots sp(i-1)$  and there exists a natural number  $n$  such that  $\text{rtt}^n(sp'_{(j)}\{\square\})$  is  $\square sp'(j+1) \dots sp'(j-1)$ . Because of  $sp_{(i)}\{\square\} \cup sp'_{(j)}\{\square\}$ ,  $\text{rtt}^m(sp_{(i)}\{\square\}) = \text{rtt}^n(sp'_{(j)}\{\square\})$  and then  $(\text{rtt}^m(sp_{(i)}\{\square\}))_{(1)}\{sp(i)\} = (\text{rtt}^n(sp'_{(j)}\{\square\}))_{(1)}\{sp'(j)\}$ . Because  $sp'(j)$  in  $sp'$  is the corresponding component to  $sp(i)$  in  $sp$ ,  $sp'(j) = sp(i)$ . Therefore,  $(\text{rtt}^m(sp_{(i)}\{\square\}))_{(1)}\{sp(i)\} = (\text{rtt}^n(sp'_{(j)}\{\square\}))_{(1)}\{sp(i)\}$  and then  $\text{rtt}^{-m}((\text{rtt}^m(sp_{(i)}\{\square\}))_{(1)}\{sp(i)\}) = sp$  and  $\text{rtt}^{-n}((\text{rtt}^n(sp'_{(j)}\{\square\}))_{(1)}\{sp'(j)\}) = sp'$ . Thus,  $sp \cup sp'$ .

(2)  $(\Rightarrow)$  There exists a natural number  $m$  such that  $\text{rtt}^m(sp)$  is  $sp(i) sp(i+1) \dots sp(i-1)$  and there exists a natural number  $n$  such that  $\text{rtt}^n(\text{rev}(sp'))$  is  $sp'(j) sp'(j-1) \dots sp'(j+1)$ . Because of  $sp \cup sp'$ ,  $\text{rtt}^m(sp) = \text{rtt}^n(\text{rev}(sp'))$  and then  $(\text{rtt}^m(sp))_{(1)}\{\square\} = (\text{rtt}^n(\text{rev}(sp'))_{(1)}\{\square\})$ .  $\text{rtt}^{-m}((\text{rtt}^m(sp))_{(1)}\{\square\}) = sp_{(i)}\{\square\}$  and  $\text{rtt}^{-n}((\text{rtt}^n(\text{rev}(sp'))_{(1)}\{\square\})) = (\text{rev}(sp'))_{(j)}\{\square\}$ . Therefore,  $sp_{(i)}\{\square\} \cup sp'_{(j)}\{\square\}$ .  $(\Leftarrow)$  There exists a natural number  $m$  such that  $\text{rtt}^m(sp_{(i)}\{\square\})$  is  $\square sp(i+1) \dots sp(i-1)$  and there exists a natural number  $n$  such that  $\text{rtt}^n(\text{rev}(sp'_{(j)}\{\square\}))$  is  $\square sp'(j-1) \dots sp'(j+1)$ . Because of  $sp_{(i)}\{\square\} \cup sp'_{(j)}\{\square\}$ ,  $\text{rtt}^m(sp_{(i)}\{\square\}) = \text{rtt}^n(\text{rev}(sp'_{(j)}\{\square\}))$  and then  $(\text{rtt}^m(sp_{(i)}\{\square\}))_{(1)}\{sp(i)\} = (\text{rtt}^n(\text{rev}(sp'_{(j)}\{\square\}))_{(1)}\{sp(i)\})$ . Because  $sp'(j)$  in  $sp'$  is the corresponding component to  $sp(i)$  in  $sp$ ,  $sp'(j) = sp(i)$ . Therefore,  $(\text{rtt}^m(sp_{(i)}\{\square\}))_{(1)}\{sp(i)\} = (\text{rtt}^n(\text{rev}(sp'_{(j)}\{\square\}))_{(1)}\{sp(i)\})$  and then  $\text{rtt}^{-m}((\text{rtt}^m(sp_{(i)}\{\square\}))_{(1)}\{sp(i)\}) = sp$  and  $\text{rtt}^{-n}((\text{rtt}^n(\text{rev}(sp'_{(j)}\{\square\}))_{(1)}\{sp(i)\})) = sp'$ . Thus,  $sp \cup sp'$ .

$sp$  and  $\text{rtt}^{-n}((\text{rtt}^n(\text{rev}(sp'_{(j)}\{\square\})))_{(1)}\{sp'(j)\}) = \text{rev}(sp')$ . Thus,  $sp \cup sp'$ .  $\square$

## 1.2 Rings

**Definition 1.9** (Rings). For  $sp \in \mathbf{SP}$ ,  $[sp]$  is called a ring pattern and satisfies (1) the rotative law ( $[sp] = [\text{rtt}(sp)]$ ) and (2) the reversible law ( $[sp] = [\text{rev}(sp)]$ ). When  $sp$  is a sequence  $seq \in \mathbf{Seq}$ ,  $[seq]$  is called a ring.

**Proposition 1.10.** For any sequence patterns  $sp, sp' \in \mathbf{SP}$  and natural numbers  $m, n$ , if  $[sp] = [sp']$ , then (1)  $[sp] = [\text{rtt}^m(sp')]$  and (2)  $[sp] = [\text{rev}^n(sp')]$ .

*Proof.* Let us suppose  $[sp] = [sp']$ . (1) By induction on  $m$ . (1.1) Base case ( $m = 0$ ) can be discharged from the assumption  $[sp] = [sp']$ . (1.2) Induction case ( $m = k + 1$ ). From Definition 1.9,  $[\text{rtt}^k(sp')] = [\text{rtt}^{k+1}(sp')]$ . From this and the induction hypothesis  $[sp] = [\text{rtt}^k(sp')]$ ,  $[sp] = [\text{rtt}^{k+1}(sp')]$ . (2) By induction on  $n$ . (2.1) Base case ( $n = 0$ ) can be discharged from the assumption  $[sp] = [sp']$ . (2.2) Induction case ( $n = k + 1$ ). From Definition 1.9,  $[\text{rev}^k(sp')] = [\text{rev}^{k+1}(sp')]$ . From this and the induction hypothesis  $[sp] = [\text{rev}^k(sp')]$ ,  $[sp] = [\text{rev}^{k+1}(sp')]$ .  $\square$

For any sequence patterns  $sp, sp' \in \mathbf{SP}$ , if  $([sp] = [sp']) \Rightarrow [sp] = [\text{rtt}(sp')] \wedge [sp] = [\text{rev}(sp')]$ , then  $[sp] = [\text{rtt}(sp)]$  and  $[sp] = [\text{rev}(sp)]$  because the equivalence relation is reflexive, namely  $[sp] = [sp]$ . Therefore, Definition 1.9 can be rephrased as follows:

**Definition 1.11** (Another definition of rings). For  $sp, sp' \in \mathbf{SP}$ ,  $[sp] = [sp']$  is inductively defined as follows: (1)  $[sp] = [sp]$  and (2) if  $[sp] = [sp']$ , then  $[sp] = [\text{rtt}(sp')]$  and  $[sp] = [\text{rev}(sp')]$ .

Let  $sp$  be  $ES_1 ES_2 \dots ES_{|sp|-1} ES_{|sp|}$ .  $\text{rtt}^{-1}(sp)$  is  $ES_2 \dots ES_{|sp|-1} ES_{|sp|} ES_1$  and  $\text{rev}^{-1}(sp)$  is  $ES_{|sp|} ES_{|sp|-1} \dots ES_1 ES_2$ . Therefore,  $\text{rtt}^{-1} = \text{rev} \circ \text{rtt} \circ \text{rev}$  and  $\text{rev}^{-1} = \text{rev}$ .

**Proposition 1.12.** For any sequence patterns  $sp, sp' \in \mathbf{SP}$ , if  $[sp] = [sp']$ , then  $[sp] = [\text{rtt}^{-1}(sp')]$  and  $[sp] = [\text{rev}^{-1}(sp')]$ .

*Proof.* This is derived from  $\text{rtt}^{-1} = \text{rev} \circ \text{rtt} \circ \text{rev}$ ,  $\text{rev}^{-1} = \text{rev}$  and Proposition 1.10.  $\square$

**Definition 1.13** (Ring pattern match). For  $sp \in \mathbf{SP}$  and  $seq \in \mathbf{Seq}$ , pattern match between  $[sp]$  and  $[seq]$  is to find all substitutions  $\sigma$  such that  $[\sigma(sp)] = [seq]$ . Let  $[sp] =?= [seq]$  be the set of all such substitutions.

**Definition 1.14** (Sequences rotated and/or reversed). For  $sp \in \mathbf{SP}$ ,  $[[sp]]$  is the set of sequences inductively defined as follows: (1)  $sp \in [[sp]]$  and (2) if  $sp' \in [[sp]]$ , then  $\text{rtt}(sp') \in [[sp]]$  and  $\text{rev}(sp') \in [[sp]]$ .

**Proposition 1.15.** For any sequence patterns  $sp, sp' \in \mathbf{SP}$ , if  $sp' \in [[sp]]$ , then  $\text{rtt}^{-1}(sp') \in [[sp]]$  and  $\text{rev}^{-1}(sp') \in [[sp]]$ .

*Proof.* This is derived from  $\text{rtt}^{-1} = \text{rev} \circ \text{rtt} \circ \text{rev}$ ,  $\text{rev}^{-1} = \text{rev}$  and Definition 1.14.  $\square$

**Proposition 1.16.** For any sequences  $\text{seq}, \text{seq}', \text{seq}'' \in \text{SP}$  and any natural number  $i \in \{1, \dots, |\text{seq}|\}$  and  $j \in \{1, \dots, |\text{seq}'|\}$  such that  $\text{seq}'(j)$  is the correspond component to  $\text{seq}(i)$ ,  $\text{seq}$  is  $e_1 \dots e_{i-1} e_i e_{i+1} \dots e_{|\text{seq}|}$  and  $\text{seq}'$  is  $e'_1 \dots e'_{j-1} e'_j e'_{j+1} \dots e'_{|\text{seq}'|}$ , (1) if  $\text{seq}(i) \sqcup \text{seq}'(j) \sqcup \text{seq}''(i)$ , then  $[\text{seq}(i) \{\text{seq}''\}] = [\text{seq}'(j) \{\text{seq}''\}]$ , and (2) if  $\text{seq}(i) \sqcup \text{seq}'(j) \sqcup \text{seq}''(i)$ , then  $[\text{seq}(i) \{\text{seq}''\}] = [\text{seq}'(j) \{\text{rev}(\text{seq}'')\}]$ .

*Proof.* (1) Let  $\text{seq}$  be  $e_1 \dots e_{i-1} e_i e_{i+1} \dots e_{|\text{seq}|}$  and  $\text{seq}'$  be  $e'_1 \dots e'_{j-1} e'_j e'_{j+1} \dots e'_{|\text{seq}'|}$ . (1) Because  $\text{seq}(i) \sqcup \text{seq}'(j) \sqcup \text{seq}''(i)$ ,  $\square e_{i+1} \dots e_{|\text{seq}|} e_1 \dots e_{i-1} = \square e'_{j+1} \dots e'_{|\text{seq}'|} e'_1 \dots e'_{j-1}$  and then  $\text{seq}'' e_{i+1} \dots e_{|\text{seq}|} e_1 \dots e_{i-1} = \text{seq}'' e'_{j+1} \dots e'_{|\text{seq}'|} e'_1 \dots e'_{j-1}$ . Therefore,  $[\text{seq}(i) \{\text{seq}''\}] = [\text{seq}'(j) \{\text{seq}''\}]$ . (2) Because  $\text{seq}(i) \sqcup \text{seq}'(j) \sqcup \text{seq}''(i)$ ,  $\square e_{i+1} \dots e_{|\text{seq}|} e_1 \dots e_{i-1} = \square e'_{j+1} \dots e'_{|\text{seq}'|} e'_1 \dots e'_{j-1}$  and then  $\text{seq}'' e_{i+1} \dots e_{|\text{seq}|} e_1 \dots e_{i-1} = \text{seq}'' e'_{j+1} \dots e'_{|\text{seq}'|} e'_1 \dots e'_{j-1}$ .  $\text{rev}(\text{seq}'' e'_{j+1} \dots e'_{|\text{seq}'|} e'_1 \dots e'_{j-1})$  is  $e'_{j+1} \dots e'_{|\text{seq}'|} e'_1 \dots e'_{j-1} \text{rev}(\text{seq}'')$ . Thus,  $[\text{seq}(i) \{\text{seq}''\}] = [\text{seq}'(j) \{\text{rev}(\text{seq}'')\}]$ .  $\square$

**Lemma 1.17.** For sequence patterns  $sp, sp' \in \text{SP}$ ,  $(sp' \in [[sp]]) \Leftrightarrow ([sp] = [sp'])$ .

*Proof.*  $(sp' \in [[sp]]) \Rightarrow ([sp] = [sp'])$  is proved by induction on Definition 1.14. (1) Base case in which  $sp \in [[sp]]$  holds.  $[sp] = [sp]$  holds because of Definition 1.15. (2) Induction case in which  $\text{rtt}(sp') \in [[sp]]$  and  $\text{rev}(sp') \in [[sp]]$  hold.  $sp' \in [[sp]]$  holds from Proposition 1.15. From the induction hypothesis  $([sp] = [sp'])$  and Definition 1.11, therefore,  $[sp] = [\text{rtt}(sp')]$  and  $[sp] = [\text{rev}(sp')]$  hold.

$(sp' \in [[sp]]) \Leftarrow ([sp] = [sp'])$  is proved by induction on Definition 1.11. (1) Base case in which  $[sp] = [sp]$  holds.  $sp \in [[sp]]$  holds because of Definition 1.14. (2) Induction case in which  $[sp] = [\text{rtt}(sp')]$  and  $[sp] = [\text{rev}(sp')]$  hold.  $[sp] = [sp']$  holds from Proposition 1.12. From the induction hypothesis  $(sp' \in [[sp]])$  and Definition 1.14, therefore,  $\text{rtt}(sp') \in [[sp]]$  and  $\text{rev}(sp') \in [[sp]]$  hold.  $\square$

Let  $sp$  be  $e_1 S_1 e_4 S_2$  and  $sp'$  be  $\text{rev}(sp)$ , namely  $S_2 e_4 S_1 e_1$ . Clearly,  $sp' \in [[sp]]$  and  $[sp] = [sp']$ . Let us consider a substitution  $\sigma$  such that  $\sigma(S_1) = e_2 e_3$ ,  $\sigma(S_2) = e_5 e_6$  and  $\sigma(X) = X$  for any other variable  $X$ .  $\sigma(sp)$  is  $e_1 e_2 e_3 e_4 e_5 e_6$  and  $\sigma(sp')$  is  $e_5 e_6 e_4 e_2 e_3 e_1$ . Clearly,  $\sigma(sp') \notin [[\sigma(sp)]]$  and  $[\sigma(sp)] \neq [\sigma(sp')]$ . If  $sp \cup sp'$  does not hold but  $sp \cup sp'$  holds, we need to reverse the sequence that replaces each sequence variable.  $\sigma_{\text{rev}}(sp')$  is  $e_6 e_5 e_4 e_3 e_2 e_1$ . Therefore,  $\sigma_{\text{rev}}(sp') \in [[\sigma(sp)]]$  and  $[\sigma(sp)] = [\sigma_{\text{rev}}(sp')]$ .

**Lemma 1.18.** For any sequence pattern  $sp \in \text{SP}$  and any substitution  $\sigma$ , for each  $sp' \in [[sp]]$  if  $sp \cup sp'$ , then  $[\sigma(sp)] = [\sigma(sp')]$ ; if  $sp \cup sp'$ , then  $[\sigma(sp)] = [\sigma_{\text{rev}}(sp')]$ .

*Proof.* By induction on the number of element and sequence variable occurrences in  $sp$ .

(1) Base case in which the number is 0. Because  $sp$  does not have any variables,  $\sigma(sp) = sp$ ,  $\sigma(sp') = sp'$  and  $\sigma_{\text{rev}}(sp') = sp'$ . From Lemma 1.17,  $[sp] = [sp']$ .

(2) Induction case in which the number is  $k + 1$ . Let us arbitrarily choose a component that is a variable in  $sp$  and the component be the  $i$ th component  $sp(i)$  in  $sp$ . Let  $sp$  be  $sp_1 sp(i) sp_2$ .  $sp'$  can be obtained by rotating and/or reversing  $sp$  and then must have the correspondent component in  $sp'$  to  $sp(i)$  in  $sp$ . Then,  $sp'$  can be  $sp'_1 sp(i) sp'_2$ .

(2.1) Let us suppose that  $sp \cup sp'$  holds. From Proposition 1.8,  $(sp_1 \sqcup sp_2) \cup (sp'_1 \sqcup sp'_2)$ . By induction hypothesis,  $[\sigma(sp_1 \sqcup sp_2)] = [\sigma(sp'_1 \sqcup sp'_2)]$  and then  $[\sigma(sp_1) \sqcup \sigma(sp_2)] = [\sigma(sp'_1) \sqcup \sigma(sp'_2)]$ . From Lemma 1.16,  $[\sigma(sp_1) \sigma(sp(i)) \sigma(sp_2)] = [\sigma(sp'_1) \sigma(sp(i)) \sigma(sp'_2)]$ . Hence,  $[\sigma(sp_1 sp(i) sp_2)] = [\sigma(sp'_1 sp(i) sp'_2)]$ .

(2.2) Let us suppose that  $sp \cup sp'$  holds. From Proposition 1.8,  $(sp_1 \sqcup sp_2) \cup (sp'_1 \sqcup sp'_2)$ . By induction hypothesis,  $[\sigma(sp_1 \sqcup sp_2)] = [\sigma_{\text{rev}}(sp'_1 \sqcup sp'_2)]$  and then  $[\sigma(sp_1) \sqcup \sigma(sp_2)] = [\sigma_{\text{rev}}(sp'_1) \sqcup \sigma_{\text{rev}}(sp'_2)]$ . From Lemma 1.16,  $[\sigma(sp_1) \sigma(sp(i)) \sigma(sp_2)] = [\sigma_{\text{rev}}(sp'_1) \text{rev}(\sigma(sp(i))) \sigma_{\text{rev}}(sp'_2)]$ . Because  $\text{rev}(\sigma(sp(i))) = \sigma_{\text{rev}}(sp(i))$ ,  $[\sigma(sp_1 sp(i) sp_2)] = [\sigma_{\text{rev}}(sp'_1 sp(i) sp'_2)]$ .  $\square$

**Definition 1.19** (Ring pattern match simulated (1)). For  $sp \in \text{SP}$  and  $\text{seq} \in \text{Seq}$ , pattern match between  $sp$  and  $[[\text{seq}]]$  is to find all substitutions  $\sigma$  such that  $\sigma(sp) = \text{seq}'$  for some  $\text{seq}' \in \text{Seq}$ . Let  $sp =?= [[\text{seq}]]$  be the set of all such substitutions.

**Lemma 1.20.** For any sequence pattern  $sp \in \text{SP}$  and any sequence  $\text{seq} \in \text{Seq}$ ,  $([sp] =?= [\text{seq}]) = (sp =?= [[\text{seq}]])$ .

*Proof.* Let  $\sigma \in ([sp] =?= [\text{seq}])$ .  $[\sigma(sp)] = [\text{seq}]$  by Definition 1.13.  $\sigma(sp) \in [[\text{seq}]]$  due to Lemma 1.17. Thus,  $\sigma \in (sp =?= [[\text{seq}]])$ .

Let  $\sigma \in (sp =?= [[\text{seq}]])$ . Let  $\text{seq}' \in [[\text{seq}]]$  such that  $\sigma(sp) = \text{seq}'$ .  $[\text{seq}'] = [\text{seq}]$  due to Lemma 1.17 and then  $[\sigma(sp)] = [\text{seq}]$ . Hence,  $\sigma \in ([sp] =?= [\text{seq}])$ .  $\square$

**Definition 1.21** (Ring pattern match simulated (2)). For  $sp \in \text{SP}$  and  $\text{seq} \in \text{Seq}$ , pattern match between  $[[sp]]$  and  $\text{seq}$  is to find all substitutions  $\sigma$  such that  $\sigma'(sp') = \text{seq}$  for some substitution  $\sigma'$  and some  $sp' \in [[sp]]$  and if  $sp \cup sp'$ , then  $\sigma = \sigma'$  and if  $sp \cup sp'$ , then  $\sigma = \sigma'_{\text{rev}}$ . Let  $[[sp]] =?= \text{seq}$  be the set of all such substitutions.

Note that  $([[sp]] =?= \text{seq}) \subset ([sp] =?= [\text{seq}])$  but  $([sp] =?= [\text{seq}]) \not\subset ([sp] =?= \text{seq})$ .

**Lemma 1.22.** For any sequence pattern  $sp \in \text{SP}$ , any sequence  $\text{seq} \in \text{Seq}$  and any substitution  $\sigma \in (sp =?= [[\text{seq}]])$ , there exist  $\sigma'$  and a sequence  $\text{seq}' \in [[\text{seq}]]$  such that  $\sigma = \text{join}(\sigma')$ ,  $\sigma'(\text{split}(sp)) = \text{seq}'$  and there exists  $sp' \in [[\text{split}(sp)]]$  such that  $\sigma'(sp') = \text{seq}$ . Besides,  $\sigma \in \text{join}([[\text{split}(sp)]] =?= \text{seq})$ .

*Proof.* Let  $sp$  be  $ES_1 ES_2 \dots ES_m$  and  $\text{seq}$  be  $e_1 e_2 \dots e_n$ .



If there exists  $i \in \{1, \dots, m\}$  such that  $\sigma(ES_i)$  is  $\dots e_n e_1$   $\dots$  or  $\dots e_1 e_n \dots$ ,  $ES_i$  is a sequence variable  $S$  that is replaced with  $sv(S, 0)$   $sv(S, 1)$  in  $\text{split}(sp)$ .

If  $\sigma(S)$  is  $\dots e_n e_1 \dots$ , then  $\sigma'(sv(S, 0))$  is  $\dots e_n$ ,  $\sigma'(sv(S, 1))$  is  $e_1 \dots$  and  $\sigma'(sv(S', 0))$  is  $\sigma(S')$  and  $\sigma'(sv(S', 1))$  is  $\varepsilon$  for any other sequence variable  $S'$  in  $sp$ , and  $\sigma'(E) = \sigma(E)$  for any element variable  $E$  in  $sp$ . By the construction of  $\sigma'$ ,  $\sigma = \text{join}(\sigma')$  and  $\sigma'(\text{split}(sp)) = \sigma(sp)$ , where  $\sigma(sp) \in [[seq]]$ . Let  $sp'$  be  $sv(S, 1)$   $\text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1})$   $sv(S, 0)$ . Then,  $sp' \in [[\text{split}(sp)]]$  and  $\sigma'(sp') = seq$ . Therefore,  $\sigma' \in ([[split(sp)]] =?= seq)$  because of  $sp \cup sp'$  from Definition 1.21. Hence  $\sigma \in \text{join}([[split(sp)]] =?= seq)$  from Definition 1.4.

If  $\sigma(S)$  is  $\dots e_1 e_n \dots$ , then  $\sigma'(sv(S, 0))$  is  $\text{rev}(\dots e_1)$ ,  $\sigma'(sv(S, 1))$  is  $\text{rev}(e_n \dots)$  and  $\sigma'(sv(S', 0))$  is  $\text{rev}(\sigma(S'))$  and  $\sigma'(sv(S', 1))$  is  $\varepsilon$  for any other sequence variable  $S'$  in  $sp$ , and  $\sigma'(E) = \sigma(E)$  for any element variable  $E$  in  $sp$ . By the construction of  $\sigma'$ ,  $\sigma = \text{join}(\sigma'_{\text{rev}})$  and  $\sigma'_{\text{rev}}(\text{split}(sp)) = \sigma(sp)$ , where  $\sigma(sp) \in [[seq]]$ . Let  $sp'$  be  $\text{rev}(sv(S, 1)$   $\text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1})$   $sv(S, 0))$ . Then,  $sp' \in [[\text{split}(sp)]]$  and  $\sigma'_{\text{rev}}(sp') = seq$ . Therefore,  $\sigma'_{\text{rev}} \in ([[split(sp)]] =?= seq)$  because of  $sp \cup sp'$  from Definition 1.21. Hence  $\sigma \in \text{join}([[split(sp)]] =?= seq)$  from Definition 1.4.

If there exists no  $i \in \{1, \dots, m\}$  such that  $\sigma(ES_i)$  is  $\dots e_n e_1$   $\dots$  or  $\dots e_1 e_n \dots$ , there must be  $i \in \{1, \dots, m\}$  such that  $\sigma(ES_i)$  is  $e_1, e_1 \dots$  or  $\dots e_1$ . If  $\sigma(ES_i)$  is  $e_1$ , there are two possible cases: (1)  $\sigma(ES_i ES_{i+1} \dots ES_{i-1}) = seq$  and (2)  $\sigma(ES_i ES_{i-1} \dots ES_{i+1}) = seq$ .

Case (1) can be treated in the same way as the case in which  $\sigma(ES_i)$  is  $e_1 \dots$ . In either case,  $\sigma'(sv(S', 0))$  is  $\sigma(S')$  and  $\sigma'(sv(S', 1))$  is  $\varepsilon$  for any sequence variable  $S'$  in  $sp$ , and  $\sigma'(E) = \sigma(E)$  for any element variable  $E$  in  $sp$ . By the construction of  $\sigma'$ ,  $\sigma = \text{join}(\sigma')$  and  $\sigma'(\text{split}(sp)) = \sigma(sp)$ , where  $\sigma(sp) \in [[seq]]$ . Let  $sp'$  be  $\text{split}(ES_i)$   $\text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1})$ . Then,  $sp' \in [[\text{split}(sp)]]$  and  $\sigma'(sp') = seq$ . Therefore,  $\sigma' \in ([[split(sp)]] =?= seq)$  because of  $sp \cup sp'$  from Definition 1.21. Hence  $\sigma \in \text{join}([[split(sp)]] =?= seq)$  from Definition 1.4.

Case (2) can be treated in the same way as the case in which  $\sigma(ES_i)$  is  $\dots e_1$ . In either case,  $\sigma'(sv(S', 0))$  is  $\text{rev}(\sigma(S'))$  and  $\sigma'(sv(S', 1))$  is  $\varepsilon$  for any sequence variable  $S'$  in  $sp$ , and  $\sigma'(E) = \sigma(E)$  for any element variable  $E$  in  $sp$ . By the construction of  $\sigma'$ ,  $\sigma = \text{join}(\sigma'_{\text{rev}})$  and  $\sigma'_{\text{rev}}(\text{split}(sp)) = \sigma(sp)$ , where  $\sigma(sp) \in [[seq]]$ . Let  $sp'$  be  $\text{rev}(\text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1}) \text{split}(ES_i))$ . Then,  $sp' \in [[\text{split}(sp)]]$  and  $\sigma'_{\text{rev}}(sp') = seq$ . Therefore,  $\sigma'_{\text{rev}} \in ([[split(sp)]] =?= seq)$  because of  $sp \cup sp'$  from Definition 1.21.  $\square$

**Lemma 1.23.** For any sequence pattern  $sp \in \mathbf{SP}$ , any sequence  $seq \in \mathbf{Seq}$  and any substitution  $\sigma \in \text{join}([[split(sp)]] =?= seq)$ ,  $\sigma \in ([sp] =?= [seq])$ .

*Proof.* Let  $sp$  be  $ES_1 ES_2 \dots ES_m$  and  $seq$  be  $e_1 e_2 \dots e_n$ . Let  $\sigma'$  be an arbitrary substitution in  $([[split(sp)]] =?= seq)$  from which  $\sigma$  is constructed, namely that  $\sigma = \text{join}(\sigma')$ . Let  $sp' \in [[\text{split}(sp)]]$  such that  $\sigma''(sp') = seq$ , if  $\text{split}(sp) \cup sp'$ ,

then  $\sigma' = \sigma''$  and if  $\text{split}(sp) \cup sp'$ , then  $\sigma' = \sigma''_{\text{rev}}$ . There are four possible cases: (1)  $sp'$  is  $\text{split}(E_i)$   $\text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1})$ , (2)  $sp'$  is  $\text{rev}(\text{split}(E_i))$   $\text{rev}(\text{split}(ES_{i+1})) \dots \text{rev}(\text{split}(ES_{i-1}))$ , (3)  $sp'$  is  $sv(S, 1)$   $\text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1})$   $sv(S, 0)$  and (4)  $sp'$  is  $sv(S, 0)$   $\text{rev}(\text{split}(ES_{i+1})) \dots \text{rev}(\text{split}(ES_{i-1}))$   $sv(S, 1)$ .

(1) For each  $ES_j$  for  $j = 1, 2, \dots, m$ , we calculate  $\sigma''(\text{split}(ES_j))$  and  $\sigma(ES_j)$ . There are three possible cases: (1.1)  $ES_j$  is an element  $e$ , (1.2)  $ES_j$  is an element variable  $E$  and (1.3)  $ES_j$  is a sequence variable  $S$ . (1.1)  $\sigma''(\text{split}(e)) = e$  and  $\sigma(e) = e = \sigma''(\text{split}(e))$ . (1.2)  $\sigma''(\text{split}(E)) = \sigma''(E)$  and  $\sigma(E) = (\text{join}(\sigma''))(E) = \sigma''(E) = \sigma''(\text{split}(E))$ . (1.3)  $\sigma''(\text{split}(S))$  and  $\sigma(S)$  are calculated as follows:

$$\sigma''(\text{split}(S)) = \sigma''(sv(S, 0) sv(S, 1))$$

$$\sigma(S) = (\text{join}(\sigma''))(S) = \sigma''(sv(S, 0)) \sigma''(sv(S, 1)) = \sigma''(\text{split}(S))$$

Therefore,  $\sigma(ES_j) = \sigma''(\text{split}(ES_j))$  and then  $\sigma(ES_i ES_{i+1} \dots ES_{i-1})$  is calculated as follows:

$$\begin{aligned} \sigma(ES_i ES_{i+1} \dots ES_{i-1}) &= \sigma(ES_i) \sigma(ES_{i+1}) \dots \sigma(ES_{i-1}) \\ &= \sigma''(\text{split}(ES_i)) \sigma''(\text{split}(ES_{i+1})) \dots \sigma''(\text{split}(ES_{i-1})) \\ &= \sigma''(\text{split}(ES_i) \text{split}(ES_{i+1}) \dots \text{split}(ES_{i-1})) \\ &= \sigma''(sp') \end{aligned}$$

Because  $\sigma''(sp') = seq$  from the assumption,  $\sigma(ES_i ES_{i+1} \dots ES_{i-1}) = seq$ . Because  $sp \cup ES_i ES_{i+1} \dots ES_{i-1}$ ,  $[\sigma(sp)] = [\sigma(ES_i ES_{i+1} \dots ES_{i-1})]$  from Lemma 1.18. Thus,  $[\sigma(sp)] = [seq]$  and then  $\sigma \in ([sp] =?= [seq])$ .

(2) For each  $ES_j$  for  $j = 1, 2, \dots, m$ , we calculate  $\sigma''(\text{rev}(\text{split}(ES_j)))$  and  $\sigma(ES_j)$ . There are three possible cases: (2.1)  $ES_j$  is an element  $e$ , (2.2)  $ES_j$  is an element variable  $E$  and (2.3)  $ES_j$  is a sequence variable  $S$ . (2.1)  $\sigma''(\text{rev}(\text{split}(e))) = e$  and  $\sigma(e) = e = \sigma''(\text{rev}(\text{split}(e))) = \text{rev}(\sigma''(\text{rev}(\text{split}(e))))$ . (2.2)  $\sigma''(\text{rev}(\text{split}(E))) = \sigma''(E)$  and  $\sigma(E) = (\text{join}(\sigma'_{\text{rev}}))(E) = \sigma''(E) = \sigma''(\text{rev}(\text{split}(E))) = \text{rev}(\sigma''(\text{rev}(\text{split}(E))))$ . (2.3)  $\sigma''(\text{rev}(\text{split}(S)))$  and  $\sigma(S)$  are calculated as follows:

$$\begin{aligned} \sigma''(\text{rev}(\text{split}(S))) &= \sigma''(\text{rev}(sv(S, 0) sv(S, 1))) \\ &= \sigma''(sv(S, 1) sv(S, 0)) \end{aligned}$$

$$\begin{aligned} \sigma(S) &= (\text{join}(\sigma'_{\text{rev}}))(S) = \sigma'_{\text{rev}}(sv(S, 0)) \sigma'_{\text{rev}}(sv(S, 1)) \\ &= \text{rev}(\sigma''(sv(S, 0))) \text{rev}(\sigma''(sv(S, 1))) \\ &= \text{rev}(\sigma''(sv(S, 1)) \sigma''(sv(S, 0))) = \text{rev}(\sigma''(\text{rev}(\text{split}(S)))) \end{aligned}$$

Therefore,  $\sigma(ES_j) = \text{rev}(\sigma''(\text{rev}(\text{split}(ES_j))))$  and then  $\sigma(ES_{i+1} \dots ES_{i-1} ES_i)$  is calculated as follows:

$$\begin{aligned} \sigma(ES_{i+1} \dots ES_{i-1} ES_i) &= \sigma(ES_{i+1}) \dots \sigma(ES_{i-1}) \sigma(ES_i) \\ &= \text{rev}(\sigma''(\text{rev}(\text{split}(ES_{i+1})))) \dots \\ &\quad \text{rev}(\sigma''(\text{rev}(\text{split}(ES_{i-1})))) \text{rev}(\sigma''(\text{rev}(\text{split}(ES_i)))) \\ &= \text{rev}(\sigma''(\text{rev}(\text{split}(ES_i)))) \\ &\quad \sigma''(\text{rev}(\text{split}(ES_{i-1}))) \dots \sigma''(\text{rev}(\text{split}(ES_{i+1}))) \\ &= \text{rev}(\sigma''(sp')) \end{aligned}$$

Because  $\sigma''(sp') = seq$  from the assumption,  $\sigma(ES_{i+1} \dots ES_{i-1} ES_i) = \text{rev}(seq)$ . Because  $sp \cup ES_{i+1} \dots ES_{i-1} ES_i$ ,  $[\sigma(sp)] = [\sigma(ES_{i+1} \dots ES_{i-1} ES_i)]$  from Lemma 1.18. Moreover,  $[seq] = [\text{rev}(seq)]$  from Proposition 1.10. Thus,  $[\sigma(sp)] = [seq]$  and then  $\sigma \in ([sp] \stackrel{?}{=} [seq])$ .

(3)  $\text{rtt}(sp')$  is calculated as follows:

$$\begin{aligned} \text{rtt}(sp') &= \text{sv}(S, 0) \text{sv}(S, 1) \text{rev}(\text{split}(ES_{i+1})) \dots \text{rev}(\text{split}(ES_{i-1})) \\ &= \text{split}(ES_i) \text{rev}(\text{split}(ES_{i+1})) \dots \text{rev}(\text{split}(ES_{i-1})) \end{aligned}$$

Because  $\sigma''(sp') = seq$ , there exists a natural number  $k$  such that  $\sigma''(\text{rtt}(sp')) = \text{rtt}^k(seq)$ . As what has been done for case (1), we have  $\sigma(ES_i ES_{i+1} \dots ES_{i-1}) = \text{rtt}^k(seq)$ . Because  $sp \cup ES_i ES_{i+1} \dots ES_{i-1}$ ,  $[\sigma(sp)] = [\sigma(ES_i ES_{i+1} \dots ES_{i-1})]$  from Lemma 1.18. Moreover,  $[seq] = [\text{rtt}^k(seq)]$  from Proposition 1.10. Thus,  $[\sigma(sp)] = [seq]$  and then  $\sigma \in ([sp] \stackrel{?}{=} [seq])$ .

(4)  $\text{rtt}(sp')$  is calculated as follows:

$$\begin{aligned} \text{rtt}(sp') &= \text{sv}(S, 1) \text{sv}(S, 0) \text{rev}(\text{split}(ES_{i-1})) \dots \text{rev}(\text{split}(ES_{i+1})) \\ &= \text{rev}(\text{split}(ES_i)) \text{rev}(\text{split}(ES_{i-1})) \dots \text{rev}(\text{split}(ES_{i+1})) \end{aligned}$$

Because  $\sigma''(sp') = seq$ , there exists a natural number  $k$  such that  $\sigma''(\text{rtt}(sp')) = \text{rtt}^k(seq)$ . As what has been done for case (2), we have  $\sigma(ES_{i+1} \dots ES_{i-1} ES_i) = \text{rev}(\text{rtt}^k(seq))$ . Because  $sp \cup ES_{i+1} \dots ES_{i-1} ES_i$ ,  $[\sigma(sp)] = [\sigma(ES_{i+1} \dots ES_{i-1} ES_i)]$  from Lemma 1.18. Moreover,  $[seq] = [\text{rev}(\text{rtt}^k(seq))]$  from Proposition 1.10. Thus,  $[\sigma(sp)] = [seq]$  and then  $\sigma \in ([sp] \stackrel{?}{=} [seq])$ .  $\square$

**Theorem 1.24.** For any sequence pattern  $sp \in \mathbf{SP}$  and any sequence  $seq \in \mathbf{Seq}$ ,  $\text{join}([\text{split}(sp)] \stackrel{?}{=} seq) = ([sp] \stackrel{?}{=} [seq])$ .

*Proof.* It is derived from Lemma 1.20, Lemma 1.22 and Lemma 1.23.  $\square$