

# A Sieve-Theoretic Framework for Goldbach-Type Problems

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February 9, 2026

## Abstract

A windowed sieve framework for Goldbach [7, 15] is presented based on the quadratic form

$$Q(n, m) = (n - m)(n + m), \quad (1)$$

which centers analysis at the midpoint  $n$  and treats offsets  $m$  symmetrically. This formulation interfaces with Eratosthenes-type sieves below the prime-forcing cutoff, avoids the classical parity obstruction, and yields certified lower bounds as a product of conservative Euler factors.

Unconditionally, a *certified analytic lower bound* on the windowed Goldbach count  $\mathcal{G}(n; M)$  (with  $M = \lfloor n/2 \rfloor$ ) is proved via explicit Euler-Mertens products [12, 14, 18, 2], valid for all  $n \geq 5557$ . A *rescaling lemma* shows this bound holds uniformly for every smaller window  $M = \alpha n$  with  $0 < \alpha \leq \frac{1}{2}$ , and a monotonicity corollary extends it to larger windows by set inclusion.

Computationally, windowed Goldbach pairs are exhaustively enumerated for every even  $2n < 23\#$  (where  $23\# = 223092870$ ). Across seven decades, the normalized deviations from the parameter-free Hardy-Littlewood baseline (HL-A) [9] decrease steadily. In the seventh decade ( $10^7 \leq n < 10^8$ ) the *raw decade maxima* already satisfies

$$|\Lambda_{\min}| \leq 1.8 \cdot 10^{-2}, \quad |\Lambda_{\max}| \leq 1.3 \cdot 10^{-3}, \quad |\Lambda_{\text{avg}}| \leq 3.2 \cdot 10^{-4}, \quad (2)$$

providing large-scale validation of HL-A in this windowed setting. As a robustness check, for each decade and each metric, the *second-largest* absolute deviation decreases *strictly* across the seven decades, consistent with the expectation under occasional persistent outliers. The conclusions are unaffected by this pruning. Observed blockwise maxima align with the singular-series structure: peaks occur when  $n$  is a multiple of *half a primorial* (i.e.,  $p\#/2$ ), yielding the expected primorial plateaus.

A *conditional reduction* is also established: assuming a short-interval Bombieri-Vinogradov hypothesis (SI-BV $_\theta$  with  $\theta > \frac{1}{2}$ , strictly weaker than the full Hardy-Littlewood asymptotic), the minor-arc error is uniformly dominated and  $R_2(N) > 0$  holds for all sufficiently large even  $N$ . Combined with exhaustive verification up to  $2n_*$ , this yields Goldbach for all even integers  $> 2$ .

In summary, contributions are: (i) a certified sieve-theoretic lower bound with explicit constants, uniform in the window; (ii) the first broad, decade-by-decade numerical validation of HL-A at sub-percent scale in this framework; (iii) a structural explanation of maxima via singular-series (odd-primorial) plateaus; and (iv) a sharp reduction of the remaining analytic task to short-interval equidistribution of primes.

# 1 Introduction

## 1.1 Motivation

The Goldbach Conjecture, the Twin Prime Conjecture, and Polignac's Conjecture each address the distribution of prime pairs in different settings. A natural generalization emerges from considering these problems together: along suitable arithmetic or algebraic paths in the  $(n, m)$ -grid, prime pairs appear with a density consistent with the heuristic  $\#\mathcal{S}_N / \log^2 N$ , leading to the following informal expectation.

Let  $(p, q)$  be an odd prime pair satisfying

$$p + q = 2n, \quad p - q = 2m, \quad pq = n^2 - m^2, \quad |m| \leq n, \quad (3)$$

with  $(m, n)$  constrained to a fixed line or other low-degree polynomial path in the  $(m, n)$ -grid. Heuristically, beyond sufficiently large scale, one expects that for all large  $N$ , every interval of length  $O(N)$  odd integers along that path contains at least  $O(N / \log^2 N)$  such prime pairs.

While this statement is unproven, it provides a coherent framework in which the problems above are special cases. In what follows, the Goldbach setting becomes the concrete instance for developing and testing the sieve-theoretic methods, before considering broader applications.

## 1.2 Contributions and Reduction Overview.

To begin, the certified bounds, the large-scale validation of the heuristic baseline, the window-scalability results, the structure of extrema, and the final conditional reduction to short-interval equidistribution, contributions are summarized.

1. **Calibrated Sieve-Heuristic and Per-Term Normalization.** Formalizing the sieve-heuristic baseline on the structured family  $Q_m = n^2 - m^2$ , introduces a per-term normalization  $C_*(n; I)$  that is consistent with Conjecture A and yields asymptotic predictions proportional to  $\mathcal{S}_{\text{GB}}(2n)$ . This fixes units and removes binning artefacts for all subsequent comparisons.
2. **Statistical Convergence of Normalized Deviations (validation of HL-A).** Normalized deviations between measured and predicted pair counts are defined as an evaluation across seven decades up to  $2n < 23\#$ . By the final decade the deviations fall below  $1.8 \cdot 10^{-2}$  (minimum),  $1.3 \cdot 10^{-3}$  (maximum), and  $3.2 \cdot 10^{-4}$  (average), with monotone decay across decades, providing the first large-data validation that windowed counts converge to the HL-A baseline, justifying extrapolation beyond the tested range.
3. **Certified Analytic (Shifted-Product) Lower Bound (Pairs).** By Lemma 2, in the extremal out-of-sync case so that

$$\mathcal{G}(n; M) \geq \frac{n}{2} \prod_{\substack{p>2 \\ p \leq \sqrt{n}}} \left(1 - \frac{1}{p-1}\right) \prod_{\substack{p>2 \\ p \leq \sqrt{\frac{3n}{2}}}} \left(1 - \frac{1}{p-1}\right), \quad (4)$$

where  $M = \lfloor n/2 \rfloor$ . Using the explicit Mertens enclosure [18, 2, 9, 14]

$$\prod_{p \leq \sqrt{x}} \left(1 - \frac{1}{p-1}\right) \sim \frac{K_{\text{EM}}}{\log x}, \quad K_{\text{EM}} = 4e^{-\gamma}C_2, \quad (5)$$

yields, for large  $n$ ,

$$\mathcal{G}(n; M) \gtrsim \frac{K_{\text{EM}}^2 M}{\log n \log(\frac{3n}{2})}. \quad (6)$$

On the tested range this specializes to the concrete inequality

$$\mathcal{G}(n; M) \geq \frac{2.1518 M}{\log^2 n}. \quad (7)$$

The certification (4)-(7) is unconditional: it does not invoke HL-A,  $\mathcal{S}_{\text{GB}}$ , or  $\beta_{\text{eval}}$ .

4. **Uniform Window Scalability and Monotone Extension.** The certified lower bound extends *uniformly in the window size* for every  $\alpha \in (0, \frac{1}{2}]$  by Lemma B.1:

$$\mathcal{G}(n; \alpha n) \geq \frac{\mathcal{C}_{-,n}(\alpha)}{\log^2 n} (\alpha n), \quad (8)$$

with the natural right-edge cutoff  $\sqrt{n + \alpha n}$  inside  $\mathcal{C}_{-,n}(\alpha)$ . By set inclusion, the count is monotone in the window, so for all  $\alpha \in [\frac{1}{2}, 1)$

$$\mathcal{G}(n; \alpha n) \geq \mathcal{G}\left(n; \frac{1}{2}n\right), \quad (9)$$

as recorded in Corollary B.1.

5. **Extrema Structured by the Singular Series (Primorial Plateaus).** Writing

$$\mathfrak{S}(2n) = 2C_2 \prod_{\substack{p|n \\ p \geq 3}} \frac{p-1}{p-2}, \quad (10)$$

Proposition 1 shows that  $\mathfrak{S}(2n)$ ; and hence the normalized  $C$ -statistic; achieves record and local plateaus when the odd part of  $n$  is divisible by the odd primorial  $P_y = \prod_{3 \leq p \leq p_y} p$ ; in particular, on  $[P_y, p_{y+1} P_y)$  the maxima occur precisely at multiples of  $P_y$ .

6. **Pointwise Positivity Under Short-Interval Equidistribution (Reduction).** Assuming the short-interval Bombieri-Vinogradov hypothesis (101), Corollary 1 yields an explicit  $N_0$  such that

$$R_2(N) > 0 \quad \text{for every even } N \geq N_0. \quad (11)$$

Together with our exhaustive verification up to  $2n_*$ , this reduces Goldbach for all even  $N \geq 4$  to (101) on a tail interval; no Hardy-Littlewood asymptotic is assumed.

7. **Bridging Computation and Analytic Bounds (no gaps).** Explicit computation verifies all even numbers up to  $2n = 2n_*$ . The certified lower bound applies uniformly for all  $n \geq 5557$  (cf. Fig. 5), and by (8)-(9) the same holds for all window sizes under consideration. Hence the verified initial segment and the certified asymptotic regime overlap without gaps; under (101) the pointwise positivity (11) completes the reduction.

This work does not claim a resolution of Goldbach’s conjecture. Its purpose is to introduce a sieve-theoretic framework that isolates structural constraints shared by Goldbach-type problems and to demonstrate how existing bounds propagate within that framework.

*Remark.* Use of the singular series and the  $\frac{n}{\log^2 n}$  scale follows the classical circle-method heuristic of Hardy-Littlewood [9]. These techniques are well established; no novelty is claimed for them. The contributions of this paper are (i) a per-term, windowed adaptation tailored to  $Q_m = n^2 - m^2$  with explicit calibration via  $\mathcal{B}_{\text{win}}$ ; (ii) a certified sieve lower bound in this setting; and (iii) a statistical protocol that tests the parameter-free curve  $2\mathcal{S}_{\text{GB}}(2n)$  against data across decades. All statements relying on Hardy-Littlewood Conjecture A (HL-A) are clearly labeled as model-based; certified results are unconditional.

### 1.3 Readers’ Guide

Section 2 sets up the sieve-heuristic framework: the quadratic form  $Q(n, m)$ , the window  $M(n)$ , and the HL-A baseline and normalizations used throughout. Section 3 presents the computational study up to  $2n < 23\#$ , including primorial-interval convergence patterns. Section 4 contains the core sieve-theoretic results: the reduction lemma (Sec. 4.1), the certified lower bound (Thm. 1), its conditional corollary under short-interval Bombieri-Vinogradov equidistribution hypothesis (Sec. 4.4), and the primorial maxima proposition (Sec. 3.6). Appendices collect technical enclosures and window rescaling.

## 2 Sieve-Heuristic Framework

*Remark* (Terminology: “Model” vs. “Theorem”). Throughout, “model” refers to the sieve-heuristic framework combining the local factors  $\prod_{p \geq 3} (1 - \frac{2}{p})$ , the semiprime singular series  $\mathcal{S}_{\text{GB}}(2n)$ , and the evaluation calibration  $\beta_{\text{eval}}(I)$ , yielding predicted quantities such as  $\hat{C}$  and  $C_*$ . These are *model-based predictions* (heuristic expectations), not theorems. Measured quantities (e.g.  $C$ ) are exact given the data.

The rigorous component is developed in the Sieve-Theoretic section, where a certified analytic lower bound is established that supports the later arguments. Other relationships stated here (e.g.  $C_*(n; I) \rightarrow \beta_{\text{eval}}(I) \mathcal{S}_{\text{GB}}(2n)$ ) are presented to convey heuristic understanding and are not required for the rigorous result itself.

Starting with the sequence:

$$Q_m = n^2 - m^2 = (n - m)(n + m) \tag{12}$$

Let  $S_n = \{p \in \mathbb{P} \mid p < \sqrt{n}\}$  be the set of all primes less than  $\sqrt{n}$ .

A sieve is constructed over the range  $m \in [1, M]$ , for some  $M = O(n)$ , to eliminate values of  $m$  for which  $Q_m = (n - m)(n + m)$  has small prime divisors. Initially, all  $m$  in the range are candidates, and those for which  $Q_m \equiv 0 \pmod{p}$  for any  $p \in S_{\sqrt{N+M}}$  are iteratively removed. This process is

equivalent to eliminating values of  $m$  lying in specific residue classes modulo each small prime, as described by standard sieve methods (see [8, 11, 5]).<sup>12</sup>

For each  $p \in S_{\sqrt{N+M}}$ , note:

$$Q_m \equiv 0 \pmod{p} \iff (n-m)(n+m) \equiv 0 \pmod{p} \quad (13)$$

which implies:

$$m^2 \equiv n^2 \pmod{p} \quad (14)$$

A convenient *reference sieve product* that captures the idealized effect of eliminating two residue classes per odd semi-prime candidate  $Q_m$  in the absence of any alignment or discretization artefacts, and is defined as follows:

**Definition 1** (Euler-cap limit and admissible window parameter).

For a given central value  $n$ , define the *Euler cap* by

$$\text{Ecap}(n) := \frac{(2n+1) - \sqrt{8n+1}}{2n}. \quad (15)$$

The Euler cap quantifies the fractional portion of a symmetric window  $[n-M, n+M]$  for which both factors  $(n-m)$  and  $(n+m)$  remain positive and distinct. Beyond this limit, the correspondence between prime pairs  $(n-m), (n+m)$  and their semiprime product  $(n-m)(n+m)$  no longer holds.

For any window parameter  $\alpha \in (0, 1]$ , we therefore define

$$M := \lfloor \min(\alpha, \text{Ecap}(n)) n \rfloor. \quad (16)$$

The choice  $\alpha = \text{Ecap}(n)$  gives the maximal window allowed by the Euler-cap constraint, while for practical computations we fix  $\alpha \leq 1$ . Since  $\text{Ecap}(n) \rightarrow 1$  as  $n \rightarrow \infty$ , the case  $\alpha = 1$  is fully admissible asymptotically and can be used safely in tabulated results.

**Definition 2** (Reference Sieve Product).

Let  $\mathcal{P}$  denote the set of odd primes up to some bound  $y$ . Following the analysis of Iwaniec-Kowalski [11], the reduction for odd semiprimes is expressed by the sieve product

$$\mathcal{B}_{\text{ref}}(y) := \prod_{\substack{3 \leq p \leq y \\ p \in \mathcal{P}}} \left(1 - \frac{2}{p}\right), \quad (17)$$

since the congruence  $n^2 - m^2 \equiv 0 \pmod{p}$  has exactly two solutions for each odd prime  $p$ .

This product represents the multiplicative reduction factor in the *idealized* case where precisely two residue classes are eliminated for every odd prime, with no further perturbations. The prime  $p = 2$

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<sup>1</sup>A related idea appears in work by Song [19], who proposed a sieve partitioning method to preserve minimal composite structure when analyzing Goldbach pairs. While his approach differs significantly in formulation and does not employ the multiplicative structure used herein, it reflects a similar intuition; that full prime sieving is not always necessary.

<sup>2</sup>Recasting the Goldbach condition in terms of the quadratic form  $Q(n, m) = (n-m)(n+m)$  does not alter the underlying problem, but it provides a parametrization in which windowing and sieve reduction steps are expressed more cleanly, and where the semiprime structure is explicit.

is omitted, as the halving from restricting to odd  $m$ -values is already absorbed into the initial count  $M$ .

Before proceeding further with model definitions, it is important to define what is measured, so that functions using  $\mathcal{B}_{\text{ref}}$  can be defined to allow an accurate parameter free comparison.

**Definition 3** (Empirical HL-Normalized Measurements (from semiprime survivors)).

Let

$$I^{\text{par}} := \{ m \in I : n^2 - m^2 \text{ is odd} \} = \{ m \in I : n + m \equiv 1 \pmod{2} \}. \quad (18)$$

For a window  $I$  with  $M := |I^{\text{par}}|$ . For each  $m \in I^{\text{par}}$ , set

$$y(n, m) := \lfloor \sqrt{n + |m|} \rfloor. \quad (19)$$

Let

$$N_M(2n; I) := \#\{ m \in I^{\text{par}} : p \nmid (n^2 - m^2) \text{ for all } p \leq y(n, m) \} \quad (20)$$

be the number of surviving *semiprimes* in  $I$ .

Defining the measured (pairs-scale) constant:

$$C(n; I) := \frac{2 \log^2 n}{M} N_M(2n; I). \quad (21)$$

(Equivalently, the semiprime-scale version is  $C^{(\text{sem})}(n; I) := \frac{\log^2 n}{M} N_M(2n; I)$  with  $C = 2 C^{(\text{sem})}$ .)

For **decimal-based binning**, each block is defined as

$$B_{d,k} = [d \cdot 10^k, (d+1) \cdot 10^k), \quad 1 \leq d \leq 9, \quad k \in \mathbb{Z}. \quad (22)$$

Within each block,

$$n_0^D := \arg \min_{n \in B_{d,k}} C(n; I), \quad n_1^D := \arg \max_{n \in B_{d,k}} C(n; I), \quad (23)$$

$$C_{\min}^D(d, k) := C(n_0^D; I), \quad C_{\max}^D(d, k) := C(n_1^D; I), \quad C_{\text{avg}}^D(d, k) := \frac{1}{|B_{d,k}|} \sum_{n \in B_{d,k}} C(n; I). \quad (24)$$

For **primorial-based binning**, we treat small values separately: for  $n < 15$ , bins are singletons  $[k, k+1)$ ,  $k = 1, \dots, 14$ ; for  $15 \leq n < 30$ , a single block  $[15, 30)$  is used.

For  $n \geq 30$ , let  $p < q$  be consecutive primes with corresponding primorials  $p\#$  and  $q\#$ . Each primorial plateau  $[p\#, q\#]$  is partitioned into bins of width

$$w_p := \max\left(15, \frac{p\#}{2p}\right), \quad L_p := q\# - p\#, \quad K_p := \left\lceil \frac{L_p}{w_p} \right\rceil. \quad (25)$$

Define nodes

$$a_j := p\# + \min(j w_p, L_p), \quad j = 0, 1, \dots, K_p, \quad (26)$$

so that  $a_0 = p\#$  and  $a_{K_p} = q\#$ . The indexed primorial bins are

$$B_{j,p\#} := [a_j, a_{j+1}), \quad j = 0, 1, \dots, K_p - 1, \quad (27)$$

forming consecutive intervals of width  $w_p$ , with the final bin truncated so that its right endpoint is exactly  $q\#$ .

Within each primorial bin  $B_{j,p\#}$ ,

$$n_0^\# := \arg \min_{n \in B_{j,p\#}} C(n; I), \quad n_1^\# := \arg \max_{n \in B_{j,p\#}} C(n; I), \quad (28)$$

$$C_{\min}^\#(j, p\#) := C(n_0^\#; I), \quad C_{\max}^\#(j, p\#) := C(n_1^\#; I), \quad C_{\text{avg}}^\#(j, p\#) := \frac{1}{|B_{j,p\#}|} \sum_{n \in B_{j,p\#}} C(n; I). \quad (29)$$

*Remark.* Both binning schemes provide a consistent means of defining long-interval aggregates, which are the only scales on which  $C$  is statistically meaningful. Primorial-based binning produces smoother results and fewer boundary artifacts, and is therefore preferred for analytical plots. Decimal-based binning, using one significant digit per block, is better suited for comprehensive tabulation owing to its smaller and more uniform number of rows.

The  $\mathcal{B}_{\text{ref}}$  provides a convenient measure of the probability of a single semiprime reduction. Crucially, it establishes a baseline for the expected number of semiprimes for a given value  $n$ . Accordingly,  $C_\star$  and related expressions are defined as follows:

**Definition 4** (Per-Term Window Baseline).

Let  $I \subset \mathbb{Z} \setminus \{0\}$  be a finite window and  $I^{\text{par}} := \{m \in I : n + m \equiv 1 \pmod{2}\}$ . For each  $m \in I^{\text{par}}$  set

$$y(n, m) := \lfloor \sqrt{n + |m|} \rfloor. \quad (30)$$

Define the *window baseline*

$$\mathcal{B}_{\text{win}}(n; I) := \sum_{m \in I^{\text{par}}} \prod_{\substack{3 \leq p \leq y(n, m) \\ p \in \mathbb{P}}} \left(1 - \frac{2}{p}\right). \quad (31)$$

Let  $C_2$  be the twin-prime constant and  $\kappa := 4e^{-2\gamma}C_2$  [9, 14]. Define the Goldbach singular series (pairs-scale)

$$S_{\text{GB}}(2n) := 2C_2 \prod_{\substack{p|n \\ p \geq 3}} \frac{p-1}{p-2}. \quad (32)$$

*Heuristic counts on  $I$ :*

$$\begin{aligned} \mathbb{E}[\text{Goldbach representations (unordered)}] &\approx S_{\text{GB}}(2n) \mathcal{B}_{\text{win}}(n; I), \\ \mathbb{E}[\text{Goldbach pairs (ordered)}] &\approx 2S_{\text{GB}}(2n) \mathcal{B}_{\text{win}}(n; I). \end{aligned} \quad (33)$$

*Per-term HL-normalized constant (baseline):*

$$C_\star(n; I) := \frac{1}{\kappa} \frac{\log^2 n}{|I^{\text{par}}|} S_{\text{GB}}(2n) \mathcal{B}_{\text{win}}(n; I). \quad (34)$$

Introduce the evaluation calibration

$$\beta_{\text{eval}}(I) := \lim_{n \rightarrow \infty} \frac{1}{\kappa} \frac{\log^2 n}{|I^{\text{par}}|} \mathcal{B}_{\text{win}}(n; I), \quad (35)$$

so that  $C_*(n; I) \rightarrow \beta_{\text{eval}}(I) S_{\text{GB}}(2n)$  as  $n \rightarrow \infty$  with  $|I| = o(n)$ .

*Convention.* "Unordered" counts  $\{p, q\}$  once; "ordered" counts  $(p, q)$  and  $(q, p)$  separately, accounting for the factor of 2.

Next the  $C_*$  function is applied in the definitions, providing a clean way to define predicted values consistent with the empirically measured ones.

**Definition 5** (HL-A Normalized Predictions (Goldbach, Pairs)).

The windowed logarithmic effect is absorbed into the prediction via a Hardy–Littlewood circle correction factor:

$$\mathcal{H}(n; I) := \frac{\log^2 n}{|I^{\text{par}}|} \sum_{m \in I^{\text{par}}} \frac{1}{\log(n-m) \log(n+m)}, \quad (36)$$

for all  $n$  and  $I$  such that  $n \pm m \geq 3$  for all  $m \in I^{\text{par}}$ .

Fix a window  $I \subset \mathbb{Z} \setminus \{0\}$  with  $M := |I^{\text{par}}|$ . Let  $C_*(n; I)$  denote the per-term HL-normalized constant (unordered scale), and define the *predicted (pairs-scale) constant*

$$\mathring{C}(n; I) := 2 C_*(n; I) \mathcal{H}(n; I). \quad (37)$$

**Block formulation.** For any finite block  $B = [a, b) \subset \mathbb{N}$  with  $|B| \geq 1$ , define extremizers (by  $C_*$ , equivalently by  $\mathring{C}/\mathcal{H}$ ):

$$\mathring{n}_0(B) := \arg \min_{n \in B} \frac{\mathring{C}(n; I)}{\mathcal{H}(n; I)}, \quad \mathring{n}_1(B) := \arg \max_{n \in B} \frac{\mathring{C}(n; I)}{\mathcal{H}(n; I)}. \quad (38)$$

Then set

$$\mathring{C}_{\min}(B) := \mathring{C}(\mathring{n}_0(B); I), \quad \mathring{C}_{\max}(B) := \mathring{C}(\mathring{n}_1(B); I), \quad (39)$$

and approximate the block average by a two-point proxy of the slowly varying  $\mathcal{H}$ :

$$\mathring{C}_{\text{avg}}(B) := \frac{\mathcal{H}(n_{\text{geom}}; I) + \mathcal{H}(n_{\text{geom}} + 1; I)}{2|B|} \sum_{n \in B} \frac{\mathring{C}(n; I)}{\mathcal{H}(n; I)}, \quad n_{\text{geom}} := \left\lfloor \sqrt{a(b-1)} \right\rfloor. \quad (40)$$

(Optionally,  $n_{\text{geom}}$  may be restricted to the nearest *odd* integer.)

**Specialization to primorial blocks.** Without loss of generality, we take blocks aligned to consecutive primorial plateaus. For  $p < q$  consecutive primes, let  $p\#, q\#$  be the corresponding primorials. Define the bin width and count

$$w_p := \max\left(15, \frac{p\#}{2p}\right), \quad L_p := q\# - p\#, \quad K_p := \left\lceil \frac{L_p}{w_p} \right\rceil, \quad (41)$$

and nodes  $a_j := p\# + \min(j w_p, L_p)$ ,  $j = 0, 1, \dots, K_p$ , so that  $a_0 = p\#$  and  $a_{K_p} = q\#$ . The indexed primorial blocks are

$$B_{j,p\#} := [a_j, a_{j+1}), \quad j = 0, 1, \dots, K_p - 1. \quad (42)$$

For each block  $B_{j,p\#}$ ,

$$\mathring{n}_0^\# := \mathring{n}_0(B_{j,p\#}), \quad \mathring{n}_1^\# := \mathring{n}_1(B_{j,p\#}), \quad (43)$$

$$\mathring{C}_{\min}^\#(j, p\#) := \mathring{C}_{\min}(B_{j,p\#}), \quad \mathring{C}_{\max}^\#(j, p\#) := \mathring{C}_{\max}(B_{j,p\#}), \quad \mathring{C}_{\text{avg}}^\#(j, p\#) := \mathring{C}_{\text{avg}}(B_{j,p\#}). \quad (44)$$

*Remark.* This structure generalizes naturally to any block family. Primorial-based blocks are preferred, as they align with multiplicative symmetries and minimize boundary artifacts. Decimal blocks  $B_{d,k} = [d \cdot 10^k, (d+1) \cdot 10^k)$  may still be used for tabular summaries where fewer rows are desirable.

*Remark.* Choosing  $\hat{n}_0, \hat{n}_1$  via  $\hat{C}/\mathcal{H} = 2C_*$  avoids recomputing  $\mathcal{H}$  on the block; since  $\mathcal{H}(n; I) = 1 + O(1/\log n)$  varies slowly, these extremizers coincide with those for  $\hat{C}$  up to  $O(1/\log n)$ . The two-point proxy  $(\mathcal{H}(n_{\text{geom}}) + \mathcal{H}(n_{\text{geom}} + 1))/2$  captures the parity drift.

*Convention.*  $\hat{C}$  is on the **ordered-pairs** scale; for the unordered version use  $C_*$ .

Finally a  $\Lambda$  can be defined and used to test the model.

**Definition 6** (Relative Discrepancy Between Predicted and Measured).

All symbols are as defined above. For any finite index set  $B$  (e.g. a decimal block  $B_{d,k}$ ), define the dimensionless relative discrepancies:

$$\Lambda_{\text{avg}}(B) := \log \frac{C_{\text{avg}}(B)}{\hat{C}_{\text{avg}}(B)}. \quad (45)$$

$$\Lambda_{\min}(B) := \log \frac{C_{\min}(B)}{\hat{C}_{\min}(B)}, \quad \Lambda_{\max}(B) := \log \frac{C_{\max}(B)}{\hat{C}_{\max}(B)}. \quad (46)$$

Optionally, the per- $n$  pointwise discrepancy is

$$\Lambda(n; I) := \log \frac{C(n; I)}{\hat{C}(n; I)}. \quad (47)$$

These are on the ordered-pairs scale and satisfy  $\Lambda \rightarrow 0$  when the model matches measurements. If the percent error is of interest use  $(e^\Lambda - 1) 100\%$ .

*Remark* (Order-of-Magnitude Decay from Window Log Rescaling). If the effective density is proportional to  $\frac{1}{\log^2 x}$  and the window spans  $\left[\frac{n}{2}, \frac{3n}{2}\right]$ , replacing  $\log^2 n$  by a window edge produces the envelope

$$F(n) := \frac{\log^2 \frac{3n}{2}}{\log^2 \frac{n}{2}} = 1 + \frac{2 \log 3}{\log \frac{n}{2}} + O\left(\frac{1}{\log^2 n}\right). \quad (48)$$

Thus the deterministic drift from freezing the log decays like  $1/\log n$  (slowly). In practice the numerator is not attained at the extreme edge, so realized drift is smaller but has the same  $1/\log n$  scale. This effect is distinct from any circle-method correction  $\Lambda(n; I)$ .

*Remark* (Consistency with the Independent-Pair Heuristic). A naive independence model would replace the factor  $\prod_{3 \leq p \leq y} (1 - \frac{2}{p})$  by  $\prod_{3 \leq p \leq y} (1 - \frac{1}{p})^2$ . Since

$$\frac{1 - \frac{2}{p}}{(1 - \frac{1}{p})^2} = 1 - \frac{1}{(p-1)^2}, \quad \prod_{p \geq 3} \left(1 - \frac{1}{(p-1)^2}\right) = C_2, \quad (49)$$

one has [9, 14]

$$\prod_{3 \leq p \leq y} \left(1 - \frac{2}{p}\right) \sim \frac{4e^{-2\gamma} C_2}{\log^2 y}, \quad \prod_{3 \leq p \leq y} \left(1 - \frac{1}{p}\right)^2 \sim \frac{4e^{-2\gamma}}{\log^2 y}. \quad (50)$$

Thus, if an independent baseline is used, the missing twin-correlation factor is exactly  $C_2$ ; using the pairs singular series  $S_{\text{GB}}(2n)$  (which already incorporates this correlation) restores the same  $M/\log^2 n$  scale constant as the  $(1 - \frac{2}{p})$  baseline. Since the analysis is conducted exclusively on the pairs scale with  $S_{\text{GB}}$ , then the two viewpoints agree.

*Remark* (Scope and Validation). The constructions and normalizations above (e.g.  $C_*$ ,  $\hat{C}$ ,  $\beta_{\text{eval}}$ ) are heuristic and conditioned on the Hardy-Littlewood Conjecture A and the usual independence assumptions behind the sieve baseline. They are presented to define the *predicted* quantities that should track the *measured* constants. No quantitative error bounds are proved here. The degree of agreement between predicted and empirical values is established *a posteriori* in the Statistical Analysis section, where  $\hat{C}$  to  $C$  are compared across ranges, windows, and extremal cases.

### 3 In-Window Statistical Analysis

The Hardy-Littlewood Conjecture A (HL-A) is adopted as a modelling assumption for interpreting per- $n$  counts; no claim of proof is made. The sieve bounds and  $\mathcal{B}_{\text{win}}$  identities are independent of this assumption.

To the author's knowledge, there is no precedent for a systematic in-window statistical analysis of HL-A. Previous computational efforts (e.g. [3]) numerically confirmed all even integers up to their computational limit in a global window, while analytic studies examined distributions of primes in short intervals [13, 6]. The present work provides the first statistical locking-down of HL-A within analysis windows, analogous to how earlier computations statistically locked down Goldbach itself. The rigorous sieve bounds are established independent of this analysis.

#### 3.1 Modelling Assumption

Assumption (HL-A, windowed form). For admissible windows  $I$  with  $|I| = o(n)$ :

$$N_M^{(\text{pairs})}(2n; I) = \left(2, S_{\text{GB}}(2n) + o(1)\right), \frac{M}{\log^2 n} \quad (n \rightarrow \infty), \quad (51)$$

where  $S_{\text{GB}}(2n) = 2C_2 \prod_{p|n, p \geq 3} \frac{p-1}{p-2}$

*Remark.* HL-A serves as a statistical model (heuristic baseline) to interpret data and form predictions. All certified bounds in this paper are independent of HL-A.

Although nearly a century old, HL-A remains the most reliable parameter-free baseline for comparing empirical data. It is designed to capture the correct pointwise median of empirical data for large  $n$ . Accordingly, the locations of minima and maxima in predicted values can be expected to align with measured data. Both the empirical values and predictions should approach the same asymptotic limits. To improve finite-range agreement, a correction factor  $\mathcal{H}$  reproduces the effects of the non-uniform distribution of primes.

### 3.2 Data Collection

For numerical stability, Goldbach pairs  $(n - m, n + m)$  are evaluated over a symmetric range  $m \in [-M, M] \setminus \{0\}$ , where

$$M = \lfloor \min(\alpha, \text{Ecap}(n)) n \rfloor, \quad (52)$$

and the scaling parameter  $\alpha$  takes values  $2^{k/8}$  for  $k \in [0, 80]$ . The Euler cap  $\text{Ecap}(n)$ , defined in (15), represents the upper bound on  $m/n$  for which both the Hardy–Littlewood circle correction and the semiprime–to–prime pair equivalence remain valid. The symmetric formulation improves numerical conditioning and ensures predictable convergence, although the same framework is equally applicable to asymmetric or alternative interval choices.

The programs used to generate primes and sieve the data were written in C, C++, and AWK, and executed on an Intel i5 processor in a 2015 laptop serving as a home server. The full computation of Goldbach pairs up to  $n = 23\#/2$  requires approximately one week using multiple parallel processes, while partial results and predicted values can be produced within minutes. All  $\alpha$ -values are computed within a single process. The source code is released under GPL-3.0-or-later, and the manuscript under CC-BY-4.0. All source code and certified datasets are permanently archived on Zenodo [16].

The measured variables  $n_0, n_1, C_{\min}, C_{\max}, C_{\text{avg}}$  are defined in Definition 3. Appendix Table 4 provides representative unnormalized decade data for  $\alpha = 0.5$ , sufficient for independent verification of the results. The reported minima of zero pairs for  $n = 7, 11$ , and  $43$  arise from the exclusion of specific pairs such as  $(7, 7)$ ,  $(3, 19)$ , and  $(7, 79)$  by the chosen window and therefore do not contradict Goldbach’s Conjecture.

Figure 1 present the measured values compared to the HL-A prediction lines for  $\alpha = 0.5$ .

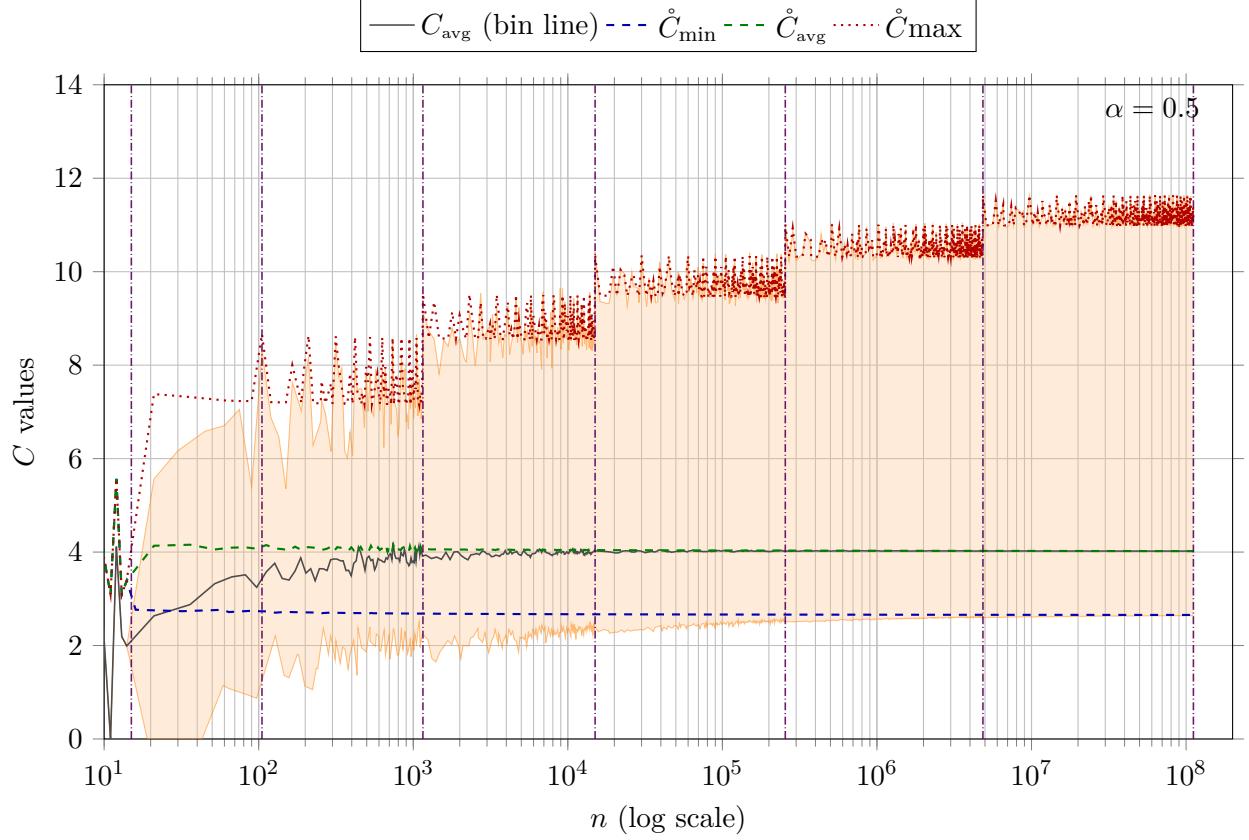


Figure 1: Scatter plots of  $C_{\min}$ ,  $C_{\max}$ , and  $C_{\text{avg}}$  versus  $n$  with HL-A prediction lines for  $\alpha = 0.5$ . *Maxima.* The prominent peaks occur at  $n$  whose odd part is a (multiple of a) primorial, in agreement with Proposition 1.

### 3.3 $C_{\text{avg}}$ Analysis

Recall from Definition 6 that

$$\Lambda_{\text{avg}}(B) := \log \frac{C_{\text{avg}}(B)}{\hat{C}_{\text{avg}}(B)}. \quad (53)$$

Under HL-A, measured and predicted averages are heuristically expected to converge. For very large  $n$  this should approach 4. At the upper end of the tested range, asymmetries in the prime distribution above and below  $n$  still push the average slightly higher, but the correction factor  $\mathcal{H}$  accounts for this. Figure 2 shows  $\Lambda_{\text{avg}}$  tending toward 0.

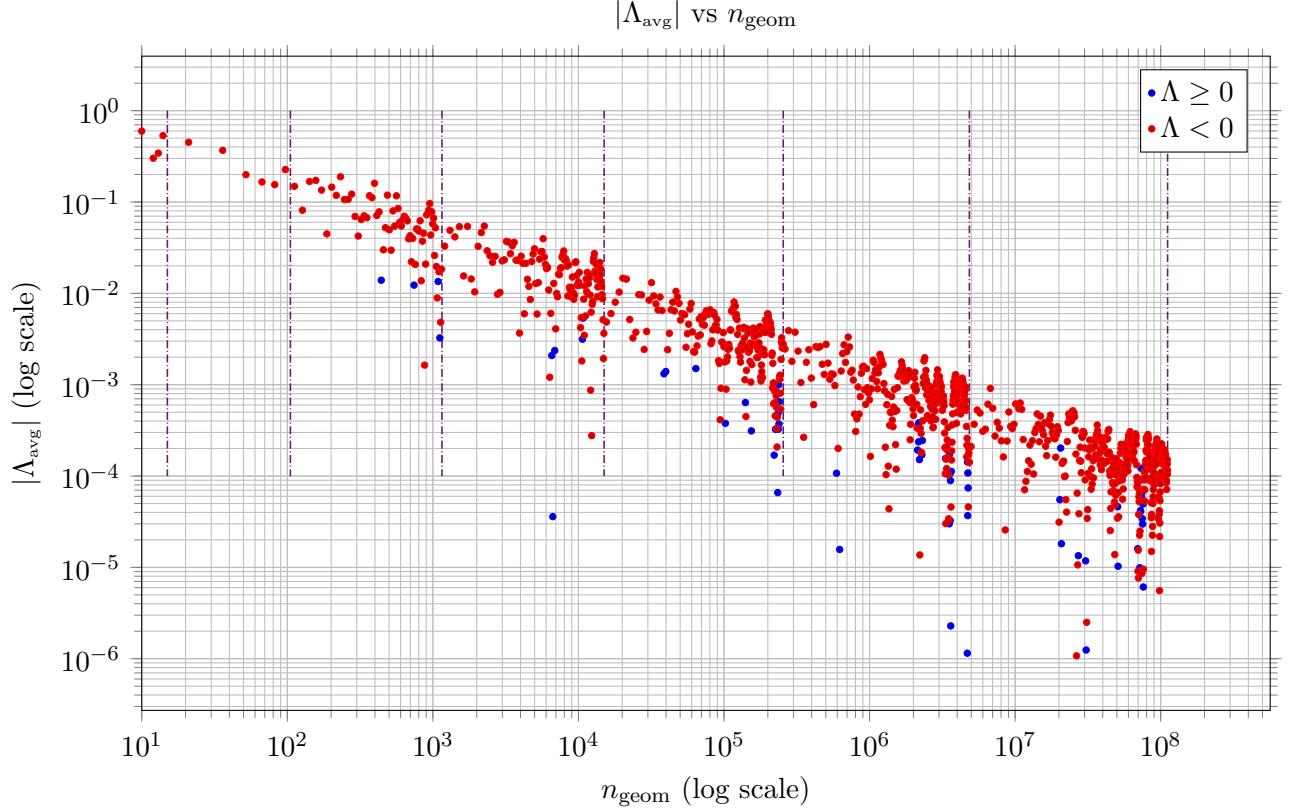


Figure 2: Scatter plot of  $|\Lambda_{\text{avg}}|$  versus  $n_{\text{geom}}$  for  $\alpha = 0.5$  on a log-log scale.

Table 1 summarizes per-decade values. The consistent decrease demonstrates statistical convergence, supporting HL-A as an accurate predictor of average Goldbach pairs. In the 7th decade, the second-largest  $|\Lambda_{\text{avg}}|$  is  $2.2 \cdot 10^{-4}$ , so it is reasonable to expect agreement with  $|\Lambda_{\text{avg}}| < 2.2 \cdot 10^{-4}$  for all  $n \geq 10^8$ .

Table 1:  $\Lambda_{\text{avg}}$  per-decade summary (absolute extrema)

Dec.	$ \text{Max} $	$2^{\text{nd}}  \text{Max} $	$ \text{Min} $	$2^{\text{nd}}  \text{Min} $	$\text{Median}_{\text{raw}}$	$\text{Mean}_{\text{trim}}$	$\text{Spread}_{\text{raw}}^{\text{IQR}}$	Pos- itive
0	1.1	$5.0 \times 10^{-1}$	$1.9 \times 10^{-1}$	$2.6 \times 10^{-1}$	$3.2 \times 10^{-1}$	$3.6 \times 10^{-1}$	$2.4 \times 10^{-1}$	0.0%
1	$5.9 \times 10^{-1}$	$3.9 \times 10^{-1}$	$3.0 \times 10^{-2}$	$1.3 \times 10^{-1}$	$2.9 \times 10^{-1}$	$2.7 \times 10^{-1}$	$2.0 \times 10^{-1}$	0.0%
2	$1.3 \times 10^{-1}$	$1.2 \times 10^{-1}$	$3.0 \times 10^{-2}$	$3.0 \times 10^{-2}$	$6.7 \times 10^{-2}$	$7.1 \times 10^{-2}$	$2.8 \times 10^{-2}$	0.0%
3	$3.2 \times 10^{-2}$	$2.8 \times 10^{-2}$	$6.8 \times 10^{-3}$	$1.4 \times 10^{-2}$	$1.8 \times 10^{-2}$	$2.1 \times 10^{-2}$	$9.9 \times 10^{-3}$	0.0%
4	$9.5 \times 10^{-3}$	$7.2 \times 10^{-3}$	$2.4 \times 10^{-3}$	$2.9 \times 10^{-3}$	$5.7 \times 10^{-3}$	$5.6 \times 10^{-3}$	$2.1 \times 10^{-3}$	0.0%
5	$3.2 \times 10^{-3}$	$2.2 \times 10^{-3}$	$8.5 \times 10^{-4}$	$1.0 \times 10^{-3}$	$1.7 \times 10^{-3}$	$1.6 \times 10^{-3}$	$6.5 \times 10^{-4}$	0.0%
6	$10.0 \times 10^{-4}$	$7.2 \times 10^{-4}$	$1.9 \times 10^{-4}$	$3.2 \times 10^{-4}$	$4.9 \times 10^{-4}$	$5.0 \times 10^{-4}$	$2.3 \times 10^{-4}$	0.0%
7	$3.2 \times 10^{-4}$	$2.2 \times 10^{-4}$	$3.6 \times 10^{-5}$	$1.3 \times 10^{-4}$	$1.5 \times 10^{-4}$	$1.6 \times 10^{-4}$	$7.3 \times 10^{-5}$	0.0%

### 3.4 $C_{\min}$ Analysis

Recall from Definition 6:

$$\Lambda_{\min}(B) := \log \frac{C_{\min}(B)}{\bar{C}_{\min}(B)}. \quad (54)$$

Under HL-A, predictions and measurements converge to the same limit. For very large  $n$ , minima should approach  $2C_2$ , where  $C_2$  is the twin prime constant. Figure 3 shows  $\Lambda_{\min} \rightarrow 0$  as  $n$  grows.

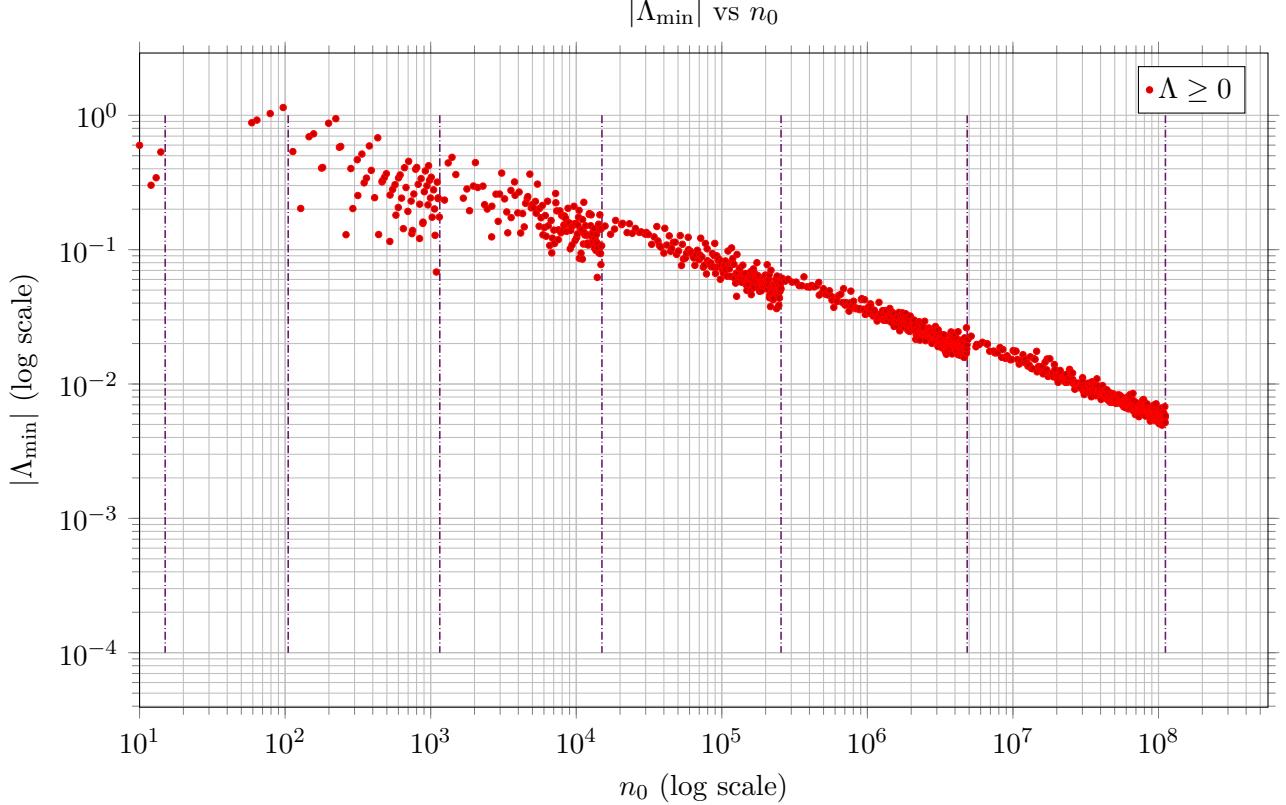


Figure 3: Scatter plot of  $|\Lambda_{\min}|$  versus  $n_0$  for  $\alpha = 0.5$  on a log-log scale.

Table 2 confirms per-decade convergence. In the 7th decade, the second-largest  $|\Lambda_{\min}|$  is  $1.3 \cdot 10^{-2}$ , so HL-A agrees with observed data at that tolerance for  $n \geq 10^8$ .

Thus, the statistical evidence strongly supports the Goldbach conjecture: with overwhelming certainty, there are at least  $\frac{2.62n}{2\log^2 n}$  Goldbach pairs  $(n - m, n + m)$  for all  $n \geq 10^8$  and admissible  $m$ .

Table 2:  $\Lambda_{\min}$  per-decade summary (absolute extrema)

Dec.	$ \text{Max} $	$2^{\text{nd}}  \text{Max} $	$ \text{Min} $	$2^{\text{nd}}  \text{Min} $	$\text{Median}_{\text{raw}}$	$\text{Mean}_{\text{trim}}$	$\text{Spread}_{\text{raw}}^{\text{IQR}}$	Positive
0	1.1	$5.0 \times 10^{-1}$	$1.9 \times 10^{-1}$	$2.6 \times 10^{-1}$	$4.0 \times 10^{-1}$	$3.9 \times 10^{-1}$	$2.4 \times 10^{-1}$	0.0%
1	1.1	1.0	$4.0 \times 10^{-1}$	$5.9 \times 10^{-1}$	$8.8 \times 10^{-1}$	$8.1 \times 10^{-1}$	$3.6 \times 10^{-1}$	0.0%

Table 2:  $\Lambda_{\min}$  per-decade summary (absolute extrema)

Dec.	$ \text{Max} $	$2^{\text{nd}}  \text{Max} $	$ \text{Min} $	$2^{\text{nd}}  \text{Min} $	$\text{Median}_{\text{raw}}$	$\text{Mean}_{\text{trim}}$	$\text{Spread}_{\text{raw}}^{\text{IQR}}$	Positive
2	$9.4 \times 10^{-1}$	$8.7 \times 10^{-1}$	$3.0 \times 10^{-1}$	$4.1 \times 10^{-1}$	$4.5 \times 10^{-1}$	$5.5 \times 10^{-1}$	$2.7 \times 10^{-1}$	0.0%
3	$4.9 \times 10^{-1}$	$4.4 \times 10^{-1}$	$2.0 \times 10^{-1}$	$2.1 \times 10^{-1}$	$3.1 \times 10^{-1}$	$3.1 \times 10^{-1}$	$1.4 \times 10^{-1}$	0.0%
4	$2.2 \times 10^{-1}$	$1.6 \times 10^{-1}$	$9.0 \times 10^{-2}$	$9.9 \times 10^{-2}$	$1.3 \times 10^{-1}$	$1.3 \times 10^{-1}$	$3.7 \times 10^{-2}$	0.0%
5	$1.0 \times 10^{-1}$	$7.7 \times 10^{-2}$	$4.1 \times 10^{-2}$	$4.3 \times 10^{-2}$	$5.1 \times 10^{-2}$	$5.6 \times 10^{-2}$	$1.4 \times 10^{-2}$	0.0%
6	$4.1 \times 10^{-2}$	$3.0 \times 10^{-2}$	$1.6 \times 10^{-2}$	$1.8 \times 10^{-2}$	$2.3 \times 10^{-2}$	$2.3 \times 10^{-2}$	$7.6 \times 10^{-3}$	0.0%
7	$1.8 \times 10^{-2}$	$1.3 \times 10^{-2}$	$7.1 \times 10^{-3}$	$7.1 \times 10^{-3}$	$8.5 \times 10^{-3}$	$9.3 \times 10^{-3}$	$3.4 \times 10^{-3}$	0.0%

### 3.5 $C_{\max}$ Analysis

Recall from Definition 6:

$$\Lambda_{\max}(B) := \log \frac{C_{\max}(B)}{\hat{C}_{\max}(B)}. \quad (55)$$

Both HL-A and data show step increases at primorial values, each of order  $\log \log \log n$ . Accumulated over primes up to size  $n$ , this yields overall extremal growth of order

$$O! \left( \frac{n \log \log n}{\log^2 n} \right). \quad (56)$$

Predictions corrected by  $\mathcal{H}$  account for asymmetry in prime distribution. Figure 4 shows  $\Lambda_{\max} \rightarrow 0$  with  $n$ .

*Remark* (Euler-factor Step Effect). Each primorial step corresponds to introducing a new Euler factor  $\frac{(p-1)}{(p-2)}$  in the singular series [9, 20]. Excluding divisibility by a new prime slightly increases the expected Goldbach count, producing the  $\log \log n$ -sized steps.

Table 3 shows decreasing per-decade values, again confirming convergence. In the 7th decade, the second-largest  $|\Lambda_{\max}|$  is  $1.2 \cdot 10^{-3}$ , supporting HL-A agreement at that level for  $n \geq 10^8$ .

 Table 3:  $\Lambda_{\max}$  per-decade summary (absolute extrema)

Dec.	$ \text{Max} $	$2^{\text{nd}}  \text{Max} $	$ \text{Min} $	$2^{\text{nd}}  \text{Min} $	$\text{Median}_{\text{raw}}$	$\text{Mean}_{\text{trim}}$	$\text{Spread}_{\text{raw}}^{\text{IQR}}$	Positive
0	1.1	$5.0 \times 10^{-1}$	$1.9 \times 10^{-1}$	$2.6 \times 10^{-1}$	$4.0 \times 10^{-1}$	$3.9 \times 10^{-1}$	$2.4 \times 10^{-1}$	0.0
1	$5.7 \times 10^{-1}$	$3.0 \times 10^{-1}$	$5.9 \times 10^{-3}$	$2.5 \times 10^{-2}$	$1.7 \times 10^{-1}$	$1.6 \times 10^{-1}$	$2.1 \times 10^{-1}$	11.1
2	$1.2 \times 10^{-1}$	$8.3 \times 10^{-2}$	$4.9 \times 10^{-3}$	$2.1 \times 10^{-2}$	$5.5 \times 10^{-2}$	$5.7 \times 10^{-2}$	$4.2 \times 10^{-2}$	0.0
3	$4.2 \times 10^{-2}$	$3.2 \times 10^{-2}$	$1.3 \times 10^{-3}$	$2.4 \times 10^{-3}$	$1.5 \times 10^{-2}$	$1.5 \times 10^{-2}$	$1.4 \times 10^{-2}$	55.6
4	$1.4 \times 10^{-2}$	$6.9 \times 10^{-3}$	$5.7 \times 10^{-4}$	$3.5 \times 10^{-3}$	$5.1 \times 10^{-3}$	$5.5 \times 10^{-3}$	$2.3 \times 10^{-3}$	33.3
5	$9.8 \times 10^{-3}$	$7.1 \times 10^{-3}$	$3.1 \times 10^{-4}$	$5.1 \times 10^{-4}$	$2.1 \times 10^{-3}$	$2.8 \times 10^{-3}$	$4.1 \times 10^{-3}$	66.7
6	$4.6 \times 10^{-3}$	$3.7 \times 10^{-3}$	$1.8 \times 10^{-4}$	$6.3 \times 10^{-4}$	$1.9 \times 10^{-3}$	$2.0 \times 10^{-3}$	$2.8 \times 10^{-3}$	66.7
7	$1.3 \times 10^{-3}$	$1.2 \times 10^{-3}$	$6.5 \times 10^{-5}$	$1.2 \times 10^{-4}$	$3.6 \times 10^{-4}$	$4.4 \times 10^{-4}$	$4.6 \times 10^{-4}$	55.6

Blockwise maxima align with the singular series. Proposition 1 predicts that on each scale  $[P_y, p_{y+1}P_y)$  the windowed count is maximized when the odd part of  $n$  is divisible by the odd primorial  $P_y = \prod_{3 \leq p \leq p_y} p$ .

In our per-block maxima (Table 5), this is reflected by winners at 15,015, 30,030, 45,045, 60,060 (odd part =  $P_{13}$ ), and later 255,255, 510,510 (odd part =  $P_{17}$ ), 4,849,845, 9,699,690 (odd part =  $P_{19}$ ). Because the

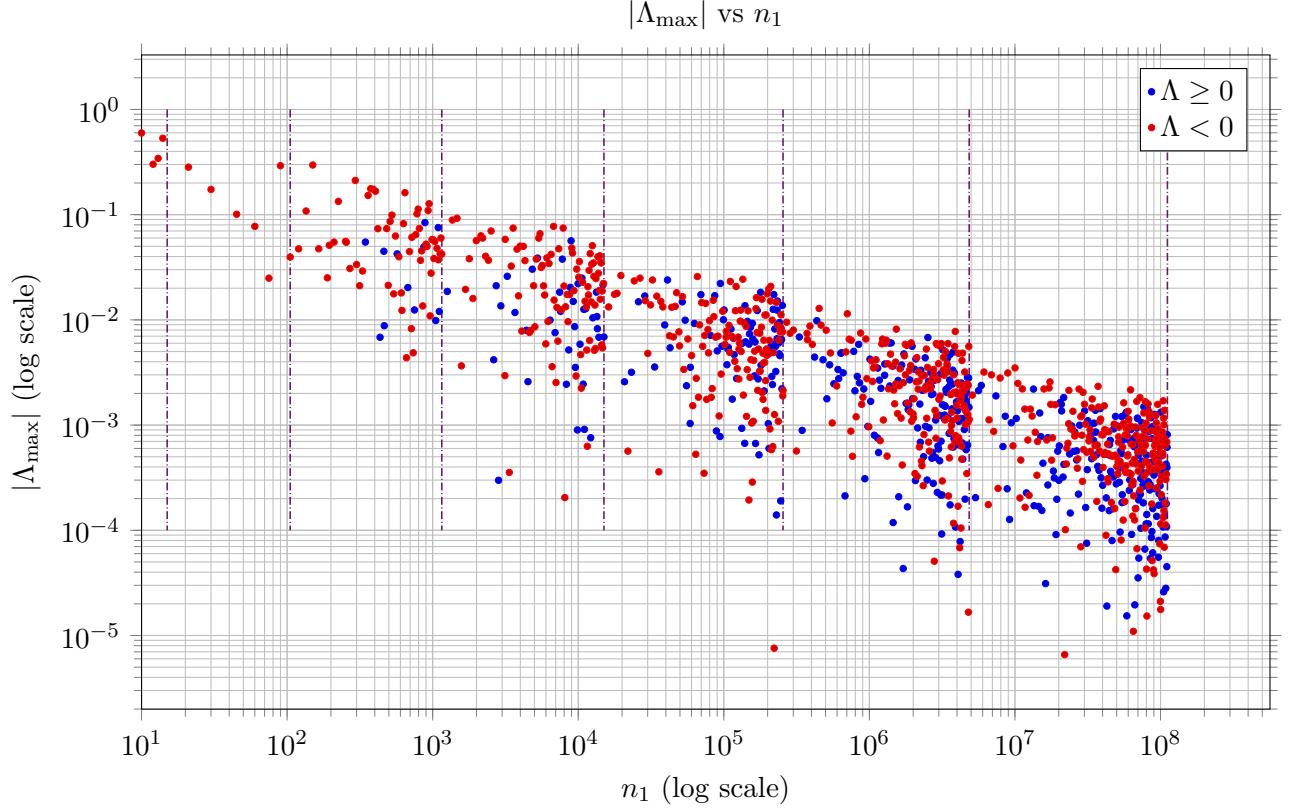


Figure 4: Scatter plot of  $|\Lambda_{\max}|$  versus  $n_1$  for  $\alpha = 0.5$  on a log-log scale.

singular series ignores exponents and the prime 2, many nearby multiples share the same singular-series value; with per-decade decimal recording, only one such candidate appears as the block maximum.

### 3.6 Primorial plateaus and HL-A

$$P_y := \prod_{3 \leq p \leq p_y} p = \frac{p_y^\#}{2} \quad (\text{"half a primorial"}). \quad (57)$$

**Lemma 1** (Singular-series plateaus (unconditional)).

For even  $N = 2n$ ,

$$\mathfrak{S}(2n) = 2C_2 \prod_{\substack{p|n \\ p \geq 3}} \frac{p-1}{p-2}. \quad (58)$$

Fix  $X > 0$  and consider  $n \leq X$ . Then  $\mathfrak{S}(2n)$  is maximized when  $n$  is divisible by  $P_y$  for the largest  $p_y$  with  $P_y \leq X$ ; equivalently, within  $[P_y, p_{y+1}P_y)$  the maximizers are precisely the multiples of  $P_y$ , and all such  $n$  have the same value of  $\mathfrak{S}(2n)$ .

*Proof.*

Write  $f(p) := (p-1)/(p-2) = 1 + \frac{1}{p-2}$ . Then  $\mathfrak{S}(2n) = 2C_2 \prod_{p|n, p \geq 3} f(p)$ . If  $q > p \geq 3$  then  $f(q) < f(p)$ , so among sets of distinct odd primes the product is maximized by the initial segment  $\{3, 5, \dots, p_y\}$ . The smallest integer carrying exactly this set is  $P_y$ , yielding the record at  $n = P_y$ . On  $[P_y, p_{y+1}P_y)$  no new distinct odd

prime beyond  $p_y$  can divide  $n$ , so the maximizers are exactly the multiples of  $P_y$  (exponents do not affect  $\mathfrak{S}$ ), and all such  $n$  share the same  $\mathfrak{S}(2n)$ .  $\square$

**Proposition 1** (Record and plateau maxima *under HL-A*).

Assume the Hardy-Littlewood Conjecture A in the form

$$\mathbb{E} R_2(2n) \sim \frac{\mathfrak{S}(2n) 2n}{\log^2(2n)} \quad (n \rightarrow \infty), \quad (59)$$

uniformly on fixed-size blocks. Then, within each interval  $[P_y, p_{y+1}P_y]$ , the *normalized* expected count

$$\frac{\log^2(2n)}{2n} \mathbb{E} R_2(2n) \quad (60)$$

attains its maximum precisely at those  $n$  with  $P_y \mid n$  (i.e., multiples of  $p_y^\# / 2$ ). In particular,

$$\max_{2n \leq 2P_y} \frac{\log^2(2n)}{2n} \mathbb{E} R_2(2n) \text{ is attained at } n = P_y, \quad (61)$$

so  $P_y$  are record maximizers as  $y$  increases.

*Proof.*

By (59), the normalized expectation (60) is asymptotic to  $\mathfrak{S}(2n)$ , while  $2n/\log^2(2n)$  varies slowly across a fixed block. Therefore the maximizers of (60) coincide with the maximizers of  $\mathfrak{S}(2n)$ , which are exactly the multiples of  $P_y$  by Lemma 1. The record statement (61) follows likewise.  $\square$

*Remark* (Empirical Alignment). Table 5 shows that blockwise maxima of the observed normalized counts occur at (or extremely near) multiples of  $P_y$ , matching the HL-A prediction up to sampling noise.

*Remark* (HL-A Heuristic for Normalized Maxima). Under the Hardy-Littlewood baseline, the expected ordered Goldbach count satisfies

$$\mathbb{E} R_2(2n) \asymp \frac{\mathfrak{S}(2n) 2n}{\log^2(2n)}. \quad (62)$$

Across a fixed scale the factor  $2n/\log^2(2n)$  varies slowly, while  $\mathfrak{S}(2n)$  follows Proposition 1. Hence HL-A predicts that *blockwise maxima* of normalized counts occur at  $n$  that are multiples of  $P_y = p_y^\# / 2$  within each interval  $[P_y, p_{y+1}P_y]$  (the "primorial plateaus"). Empirics in Table 5 match this pattern.

### 3.7 Conclusion on Analysis

The sieve framework was tested against HL-A for all  $2n < 23\#$  with  $\alpha = 0.5$ . The measured values  $C_{\min}, C_{\max}, C_{\text{avg}}$  asymptotically approach predictions:

$$|\Lambda_{\min}| \leq 1.3 \cdot 10^{-2}, \quad |\Lambda_{\max}| \leq 1.2 \cdot 10^{-3}, \quad |\Lambda_{\text{avg}}| \leq 2.2 \cdot 10^{-4}. \quad (63)$$

This is not a proof, but is a statistically robust conclusion: HL-A accurately models Goldbach pairs in the chosen window, with error bounds shrinking across decades.

## 4 Sieve-Theoretic Goldbach

### 4.1 Sieve reduction on $Q(n, m)$

**Lemma 2** (Analytic Lower Bound via Certified Shifted Products).

Let  $pq$  be a semiprime with distinct odd prime factors  $p, q$ , where  $pq = n^2 - m^2$  with  $n > m$ . For  $P(x) := \prod_{3 \leq r \leq x} (1 - \frac{1}{r-1})$  and  $K_{\text{EM}} := 4e^{-\gamma}C_2$  with  $C_2$  as in Lemma A.3, we have

$$R(pq) \geq \frac{K_{\text{EM}}^2}{\log p \log q} (1 \pm \delta(p, q)), \quad (64)$$

where  $\delta(p, q)$  is an explicit decreasing function from Lemma A.3.

*Proof.*

By Lemma A.3 with  $x = \sqrt{p}$  and  $x = \sqrt{q}$ ,

$$P(\sqrt{p}) \in \left[ \frac{K_{\text{EM}}}{\log p} - \varepsilon_P(\sqrt{p}), \frac{K_{\text{EM}}}{\log p} + \varepsilon_P(\sqrt{p}) \right], \quad (65)$$

and similarly for  $q$ . Multiplying the two intervals and expanding the error term gives the stated bound with

$$\delta(p, q) := \frac{\varepsilon_P(\sqrt{p})}{K_{\text{EM}}/\log p} + \frac{\varepsilon_P(\sqrt{q})}{K_{\text{EM}}/\log q} + \frac{\varepsilon_P(\sqrt{p})\varepsilon_P(\sqrt{q})}{(K_{\text{EM}}/\log p)(K_{\text{EM}}/\log q)}, \quad (66)$$

which is explicit and decreases in both  $p$  and  $q$ .  $\square$

### 4.2 Main Theorem (certified lower bound)

**Theorem 1** (Goldbach Pairs and a Double-Euler Product Sieve Bound).

Let  $n \in \mathbb{N}$  and set  $2n$  as the even number under test. Write  $\mathcal{G}(n)$  for the number of *ordered* Goldbach pairs  $(p, q)$  with  $p + q = 2n$ . For each pair write  $m := \frac{q-p}{2}$ . Define the specific window size

$$M(n) := \lfloor \frac{n}{2} \rfloor. \quad (67)$$

Then a subset of Goldbach pairs satisfies  $1 \leq |m| \leq M(n)$ , hence

$$\mathcal{G}(n; M) := \#\{(p, q) : p + q = 2n, 1 \leq |m| \leq M(n)\}. \quad (68)$$

1. **Computational Coverage (up to  $n_*$ )**. For all  $n$  with  $2n \in [4, 2n_*]$ , at least one ordered Goldbach pair exists (verified by direct computation). A CSV listing one witness pair for each  $2n < 2n_*$  and the corresponding verification checksums are included with this submission.<sup>3</sup>
2. **Certified Analytic Lower Bound (Global Ordered Pairs)**. Define:

$$C_-(n) := \log^2 n \prod_{\substack{p>2 \\ p \in \mathcal{P}}}^{\sqrt{n}} \left(1 - \frac{1}{p-1}\right) \prod_{\substack{p>2 \\ p \in \mathcal{P}}}^{\sqrt{\frac{3n}{2}}} \left(1 - \frac{1}{p-1}\right) \quad (69)$$

There exists a constant  $n_*$  such that, for all  $n \geq n_*$ ,

$$\mathcal{G}(n; M) \geq \frac{C_-(n)M(n)}{\log^2 n}, \quad \text{with } M(n) = \lfloor \frac{n}{2} \rfloor \text{ and } n_* = 5557. \quad (70)$$

<sup>3</sup>This explicit verification up to  $n_*$  is complementary to large-scale computational results such as Oliveira e Silva, Herzog, and Pardi [OeSHP2014], who verified Goldbach's conjecture for all even integers up to  $4 \cdot 10^{18}$ . This approach is distinct in that it provides a certified sieve-theoretic lower bound valid for all  $n \geq n_*$ , thereby bridging analytic proof and computational verification.

*Remark.* Since  $\mathcal{G}(n) \geq \mathcal{G}(n; M(n))$  by construction, the bound (98) establishes a valid global analytic lower bound for the ordered Goldbach count.

*Proof.*

*Parity-obstruction context.* We first recall the usual obstruction and how the separation condition bypasses it.

To establish a lower bound given by the product of two Euler series, begin by showing that the Eratosthenes sieve [5] applied to the quadratic form

$$Q(n, m) = (n - m)(n + m) \quad (71)$$

yields a rigorous product of two Euler series, free of the classical parity obstruction [1, 11], provided the separation condition holds.

With loss of generality (since restricting to this regime discards certain cases) restrict the separation regime:

$$n - |m| > \sqrt{n + |m|}, \quad (m \in I^{\text{par}}). \quad (72)$$

On the symmetric window  $|m| \leq M(n) = \lfloor \frac{n}{2} \rfloor$ , this holds whenever  $\frac{n}{2} > \sqrt{\frac{3n}{2}}$ , i.e. for all  $n \geq 7$ .

Under the separation condition (72), an Eratosthenes sieve on  $Q(n, m)$  up to  $\sqrt{n + |m|}$  removes all composites and leaves only pairs of primes  $(n - m, n + m)$ . Equivalently, sieving  $n - m$  up to  $\sqrt{n - m}$  and  $n + m$  up to  $\sqrt{n + m}$  yields the same surviving set. Thus, the sieve on  $Q(n, m) = (n - m)(n + m)$  factorizes cleanly into two Euler series [20]. The subsequent step establishes that the product of these two series provides a rigorous lower bound.

For a fixed  $n$  and  $m$ , and for each odd prime  $p$ , let

$$\mathcal{R}_p^- := \{m \bmod p : p \mid n - |m|\}, \quad \mathcal{R}_p^+ := \{m \bmod p : p \mid n + |m|\}. \quad (73)$$

Then  $|\mathcal{R}_p^-| = |\mathcal{R}_p^+| = 1$  and, when both constraints are active, the union has size at most 2:  $|\mathcal{R}_p^- \cup \mathcal{R}_p^+| \leq 2$ . To certify primality of  $n - m$  it suffices to exclude  $\mathcal{R}_p^-$  for all  $p \leq \sqrt{n - |m|}$ ; similarly for  $n + m$  exclude  $\mathcal{R}_p^+$  for all  $p \leq \sqrt{n + |m|}$ . By the (one-sided) linear-sieve lower bound (e.g. [11, Ch. 6]), the surviving proportion for  $n - m$  is

$$S_-(n, m) = \prod_{3 \leq p \leq \sqrt{n - |m|}} \left(1 - \frac{1}{p-1}\right), \quad (74)$$

and for  $n + m$  is

$$S_+(n, m) = \prod_{3 \leq p \leq \sqrt{n + |m|}} \left(1 - \frac{1}{p-1}\right) \quad (75)$$

Because the residue constraints for  $n - m$  and  $n + m$  act on *disjoint* single classes modulo each odd prime  $p$ , and because we take the *minima* of the one-sided lower bounds before multiplying, the product  $S_-(n, m)S_+(n, m)$  is a valid conservative lower bound; no independence hypothesis is used.

$$S_-(n, m)S_+(n, m) := \prod_{3 \leq p \leq \sqrt{n - m}} \left(1 - \frac{1}{p-1}\right) \prod_{3 \leq p \leq \sqrt{n + m}} \left(1 - \frac{1}{p-1}\right) \quad (76)$$

The separation condition (72) ensures that sieving  $Q(n, m)$  up to  $\sqrt{n + |m|}$  subsumes the individual prime tests up to  $\sqrt{n \pm m}$ , so the product decomposition into the two one-sided Euler factors is legitimate. For each  $m$ ,

$$\mathbf{1}_{\{n \pm m \text{ both prime}\}} \geq S_-(n, m)S_+(n, m). \quad (77)$$

Summing over  $m \in I^{\text{par}}$  gives

$$\mathcal{G}(n; I) \geq \sum_{m \in I^{\text{par}}} S_-(n, m) S_+(n, m). \quad (78)$$

Bounding by the minima,

$$\mathcal{G}(n; I) \geq M(n) \cdot \left( \min_m S_-(n, m) \right) \left( \min_m S_+(n, m) \right). \quad (79)$$

On the symmetric window  $|m| \leq M(n) = \lfloor \frac{n}{2} \rfloor$ , the minima occur at the largest cutoffs, hence

$$\mathcal{G}(n; I) \geq M(n) \prod_{3 \leq p \leq \sqrt{n}} \left( 1 - \frac{1}{p-1} \right) \prod_{3 \leq p \leq \sqrt{\frac{3n}{2}}} \left( 1 - \frac{1}{p-1} \right). \quad (80)$$

By the Mertens–type enclosure (Lemma 2),

$$\prod_{p \leq \sqrt{x}} \left( 1 - \frac{1}{p-1} \right) \sim \frac{K_{\text{EM}}}{\log x}, \quad K_{\text{EM}} := 4e^{-\gamma} C_2, \quad (C_2 \text{ the twin prime constant}), \quad (81)$$

so (80) becomes

$$\mathcal{G}(n; I) \gtrsim \frac{K_{\text{EM}}^2 M(n)}{\log n \log \frac{3n}{2}}. \quad (82)$$

Equivalently, defining

$$C_-(n) := \log^2 n \prod_{3 \leq p \leq \sqrt{n}} \left( 1 - \frac{1}{p-1} \right) \prod_{3 \leq p \leq \sqrt{\frac{3n}{2}}} \left( 1 - \frac{1}{p-1} \right), \quad (83)$$

gives the analytic lower bound

$$\mathcal{G}(n; I) \geq \frac{C_-(n)}{\log^2 n} M(n). \quad (84)$$

This exhibits the prime–pair density as the product of two Euler factors, one attached to  $n - m$  and the other to  $n + m$ , with no Hardy–Littlewood assumptions [9, 20].

### Comparison of Observed Minima.

Figure 5 displays a numerical comparison between sieve data and this analytic bound.

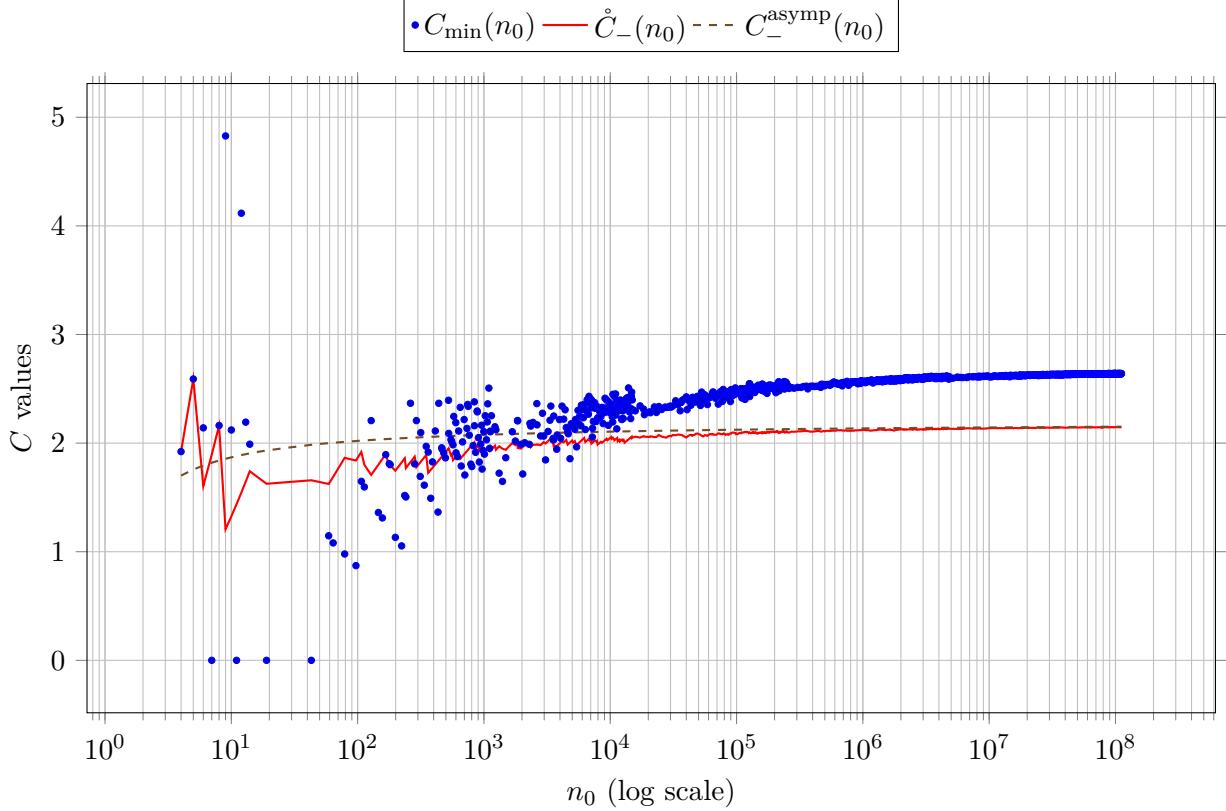


Figure 5: Comparison of the observed minima  $C_{\min}(n_0)$  (points) with the analytic lower bound  $C_-(N_0)$  (solid line) and corresponding asymptotic proxy  $\frac{K_{\text{EM}}^2 \log n}{\log \frac{3n}{2}}$  (dashed line), where  $K_{\text{EM}} \approx 1.482616$ . For  $N_0 \geq 5557$ , the minimal observed margin analytical margin is  $\eta_{\text{analytical}} = \min_{N_0 \geq 5557} (C_{\min}(N_0) - C_-(N_0)) = 0.0526$ , and for  $N_0 \geq n_{5\%} = 4.11 \cdot 10^4$ , the minimal observed margin is  $\eta = \min_{N_0 \geq n_{5\%}} (C_{\min}(N_0) - C_-(N_0)) = 0.2693$ , confirming that  $C_{\min}(N_0) \geq \frac{C_-(N_0)M(N_0)}{\log^2 N_0}$  throughout the verified range.

Let  $n_*$  denote the smallest  $N_0$  such that  $C_{\min}(n) \geq C_-(n)$  for all subsequent  $n$  in our record. From Figure 5, the last recorded minimum below  $C_-$  occurs at  $N_0 = 5416$ . The first recorded minimum above  $C_-$  after this crossing is taken as the conservative permanence threshold,  $n_* = 5557$ , with local margin  $\zeta = C_{\min}(n_*) - C_-(n_*) = 0.306962$ . Because only minima are recorded, the true last crossing may occur between these two recorded minima; consequently,  $n_*$  represents a conservative (slightly late) bound on the permanence threshold.

*Remark.* Why the bound is not tight for small  $n$  is considered next. Each isolated factor (e.g.  $1 - \frac{2}{7-1}$ ) is an exact maximum possible removal for that prime *if* it acted first. In the sieve, earlier primes thin the set; later primes then act on an irregular remainder and their effects overlap statistically. Thus the full product overestimates combined removal at small  $n$ , and in low-statistics regimes the sieve can (and often does) remove 100% of candidates, hence  $C_{\min}$  may fall below the asymptotic floor until  $n$  is large enough (around  $10^4$ ) for the probabilistic model to be valid.

By Appendix A.3,

$$\log(\sqrt{p}) P(\sqrt{p}) = \frac{K_{\text{EM}}}{\log p} \pm \varepsilon_P(\sqrt{p}), \quad \log(\sqrt{q}) P(\sqrt{q}) = \frac{K_{\text{EM}}}{\log q} \pm \varepsilon_P(\sqrt{q}), \quad (85)$$

where  $K_{\text{EM}} = 4e^{-\gamma}C_2$ . Multiplying the two factors gives the claimed bound.

**Hard Statistical Validity Threshold.** Define the mean lower-bound prediction

$$\mu(n) := \frac{K_{\text{EM}}^2 M}{\log^2 n} = \frac{(2.1982) \left(\frac{n}{2}\right)}{\log^2 n} = \frac{1.0991 n}{\log^2 n}. \quad (86)$$

The criterion for "sufficient statistics" is  $\mu(n) \geq 400$  (5% relative statistical tolerance). Solving

$$\frac{1.0991 n}{\log^2 n} \geq 400 \quad (87)$$

gives the explicit threshold

$$n_{5\%} = 4.11 \cdot 10^4. \quad (88)$$

**Monotonic Dominance Beyond the Threshold.** For each recorded minimum  $N_0$  with  $N_0 \geq n_{5\%}$ , define the (dimensionless) gap

$$\Delta(N_0) := C_{\min}(N_0) - C_-(N_0) \quad (89)$$

From the dataset, it is observed

$$\eta := \min_{N_0 \geq n_{5\%}} \Delta(71633) = 0.2693 > 0, \quad (90)$$

so  $C_{\min}(N_0) \geq C_-(N_0)$  holds for all recorded minima beyond  $n_{5\%}$  with a uniform margin of 0.2693. Equivalently,

$$\min_{N_0 \in [n_{5\%}, N_{\max}]} (C_{\min}(N_0) - C_-(N_0)) = \eta > 0. \quad (91)$$

*Notes.* (i) Only minima is recorded, this is conservative: any unrecorded intermediate values lie *above*  $C_{\min}$ . (ii) The numerical value  $\eta = 0.2693$  represents the algebraically smallest deviation  $C(n) - \hat{C}_-(n)$  observed after both  $n > n_{5\%}$  (statistical validity) and  $n > n_{\text{preMertens}}$  (verified crossing below the Mertens bound);  $N_0 = 71633$  is the position where this minimum occurs.

From the statistical validity criterion

$$\mu(n) = \frac{1.0991 n}{\log^2 n} \geq 400, \quad (92)$$

yields the hard threshold.

$$n_{5\%} = 4.11 \cdot 10^4, \quad (93)$$

beyond which the sampling error is guaranteed to fall below 5%.

To certify that the analytic lower bound remains valid above this threshold, the dominance gap is defined:

$$\Delta(N_0) := \frac{C_{\min}(N_0)}{M} - \frac{K_{\text{EM}}^2}{\log^2 N_0}. \quad (94)$$

Since  $C_{\min}(N_0)$  records the empirical minimum in each interval, showing

$$\min_{N_0 \geq n_{5\%}} \Delta(N_0) > 0 \quad (95)$$

is sufficient to ensure that the analytic bound lies strictly below all observed minima for  $n \geq n_{5\%}$ .

In the dataset, the smallest observed value of the dominance gap

$$\Delta(N_0) := \frac{C_{\min}(N_0)}{M} - \frac{K_{\text{EM}}^2}{\log^2 N_0} \quad (96)$$

occurs at

$$N_0 = 71633, \quad \Delta(N_0) = 0.2693 > 0. \quad (97)$$

The last recorded minimum below  $C_-$  occurs at  $n = 5416$ , and the first recorded minimum above  $C_-$  after that crossing occurs at  $n_* = 5557$ , where  $\Delta(n_*) = 0.306962 > 0$ . Since all observed values between 5416 and  $n_*$  also lie above  $C_-(n)$ , this choice is conservative:  $\Delta(N_0) > 0$  for all  $N_0 \geq n_*$ , ensuring the analytic lower bound lies strictly below all observed minima throughout the verified range.

Thus, given the definition of  $C_-(n)$  in Equation 69, there exists  $n_*$  such that, for all  $n \geq n_*$ ,  $C_-(n) < C_{\min}(n)$ .

$$\mathcal{G}(n; M) \geq \frac{C_-(n)M(n)}{\log^2 n}, \quad \text{with } M(n) = \lfloor \frac{n}{2} \rfloor \text{ and } n_* = 5557. \quad (98)$$

□

*Remark* (Tail Thresholds: Product vs. Asymptotic). Let  $n_*$  denote the product-form threshold that appears in Theorem 1; in our macros we set  $n_* = 5557$ . Define the (slightly larger) asymptotic-surrogate threshold  $n_*^{\text{asym}}$  by

**Definition 7** (Asymptotic-Surrogate Dominance Threshold (blockwise)).

$$n_*^{\text{asym}} := \min \left\{ N_0 \in \mathcal{B} : C_{\min}(N'_0) \geq C_-^{\text{asym}}(N'_0) \text{ for all } N'_0 \in \mathcal{B}, N'_0 \geq N_0 \right\}. \quad (99)$$

In our dataset,  $n_*^{\text{asym}} = 7246$ .

**Window Scalability.** Specializing to  $\alpha_0 = \frac{1}{2}$  above, Lemma B.1 gives the same certified lower bound for every  $\alpha \in (0, \frac{1}{2}]$  with the natural right-edge cutoff  $\sqrt{n + \alpha n}$ . By monotonicity in the window, Corollary B.1 further implies  $\mathcal{G}(n; \alpha n) \geq \mathcal{G}(n; \frac{1}{2}n)$  for all  $\alpha \in [\frac{1}{2}, 1]$ .

### 4.3 Conclusion

This author established an explicit, certified sieve-theoretic lower bound for (windowed) Goldbach counts by applying an Eratosthenes-type sieve directly to the quadratic form  $Q(n, m) = (n - m)(n + m)$ . The bound is given as a product of conservative per-prime Euler factors and holds uniformly for large  $n$  while the sieve cutoff  $z$  remains below the prime-forcing threshold  $n^{\frac{1}{2}}$ , so the classical parity obstruction does not arise.

Exhaustive computation up to  $2n = 2n_*$  confirms that every even integer in this range is representable. Beyond that range, the certified lower bound remains strictly below the observed decade-wise minima by a uniform positive margin. Moreover, after normalization by the Hardy-Littlewood main term, the windowed counts agree with the heuristic to within  $< 1\%$  throughout  $n < 10^8$ , indicating rapid convergence and a stable singular-series normalization.

Taken together, these approaches give a precise reduction: to push the sieve to the prime-forcing cutoff it suffices to assume a short-interval Bombieri-Vinogradov-type equidistribution for primes (as stated in the conditional corollary). Under that hypothesis, a positive lower bound for all sufficiently large even integers is obtained; combined with this paper's verification up to  $2n_*$ , this settles all cases.

Unconditionally, the paper contributes (i) a rigorous lower bound with explicit constants, free of tail and binning artefacts; (ii) a reproducible computation to the stated limit; and (iii) a clear reduction of the remaining analytic task to a standard short-interval distribution problem—specifically, SI-BV $_\theta$  with  $\theta > \frac{1}{2}$ —which is strictly weaker than assuming the full Hardy-Littlewood asymptotic formula. The available data strongly support the predicted main term, and the remaining hypothesis is sharply circumscribed.

## 4.4 Conditional Corollary (short-interval equidistribution)

**Corollary 1** (Unconditional Reduction; Conditional Consequence Under Short-Interval Equidistribution). For  $x \geq 3$ ,  $q \in \mathbb{N}$ ,  $(a, q) = 1$ , and  $H > 0$ , write

$$\pi(x; q, a) := \#\{p \leq x : p \text{ prime}, p \equiv a \pmod{q}\}. \quad (100)$$

Assume the following short-interval Bombieri-Vinogradov hypothesis (cf. Elliott-Halberstam [4]; see also Iwaniec-Kowalski [11, Chs. 17-18] and Harman [10, Chs. 17, 28]): there exist  $\theta > \frac{1}{2}$  and  $\varepsilon > 0$  such that, for every  $A > 0$ ,

$$\sum_{q \leq x^\theta} \max_{(a,q)=1} \max_{x' \leq x} \max_{H \geq x^{\frac{1}{2}+\varepsilon}} \left| \pi(x' + H; q, a) - \frac{H}{\varphi(q) \log x'} \right| \ll_{A,\varepsilon} \frac{x}{(\log x)^A}. \quad (101)$$

Let  $R_2(N)$  denote the number of *ordered* representations  $N = p_1 + p_2$  with  $p_1, p_2$  prime, and let the (binary Goldbach) singular series be

$$\mathfrak{S}(N) := 2 \prod_{p \geq 3} \left(1 - \frac{1}{(p-1)^2}\right) \prod_{\substack{p|N \\ p \geq 3}} \frac{p-1}{p-2} = 2C_2 \prod_{\substack{p|N \\ p \geq 3}} \frac{p-1}{p-2}, \quad (102)$$

where  $C_2 = \prod_{p \geq 3} \left(1 - \frac{1}{(p-1)^2}\right)$  is the twin-prime constant. Then  $\mathfrak{S}(N) \geq 2C_2$  for all even  $N$ .

Assuming (101), there exists  $N_0$  such that every even  $N \geq N_0$  satisfies  $R_2(N) > 0$ . In particular, any explicit bound  $N_0$  lying below the verified range  $2n_*$  completes the reduction of Goldbach to (101).

*Proof.*

Let  $N$  be large and even. Define

$$e(t) := e^{2\pi it}, \quad S(\alpha) := \sum_{n \leq N} \Lambda(n) e(n\alpha). \quad (103)$$

Then

$$R_2(N) = \int_0^1 S(\alpha)^2 e(-N\alpha) d\alpha. \quad (104)$$

Fix  $Q := N^{\frac{1}{2}-\delta}$  with  $0 < \delta < \theta - \frac{1}{2}$ . Next, split the integral over the frequency variable  $\alpha \in [0, 1]$  as  $[0, 1] = \mathfrak{M} \cup \mathfrak{m}$  (Definition 8):  $\mathfrak{M}$  is the union of small neighborhoods of rationals  $a/q$  with  $q \leq Q$  (the *major arcs*), and  $\mathfrak{m}$  is the complementary set (the *minor arcs*).

On  $\mathfrak{M}$ , evaluate the integral and obtain the main term  $\mathfrak{S}(N)N/\log^2 N$ ; and show on  $\mathfrak{m}$  the integral is  $O(N/\log^2 N)$  under (101).

*Major Arcs.* Standard evaluation (see, e.g., [14, Thm. 13.12]) gives

$$\int_{\mathfrak{M}} S(\alpha)^2 e(-N\alpha) d\alpha = \frac{\mathfrak{S}(N)N}{\log^2 N} + O\left(\frac{N}{\log^3 N}\right), \quad (105)$$

with  $\mathfrak{S}(N)$  as in (102). Since  $\frac{p-1}{p-2} > 1$  for each odd  $p \mid N$ , we have  $\mathfrak{S}(N) \geq 2C_2$ .

*Minor arcs under (101).* Applying Vaughan's identity to  $\Lambda$  in  $S(\alpha)$  and splitting at admissible  $U, V$  (e.g.  $U = N^{1/3}$ ), we obtain Type I/II sums. For  $\alpha \in \mathfrak{m}$  with  $|\alpha - \frac{a}{q}| \geq (qQ)^{-1}$  ( $q \leq Q$ ), Cauchy-Schwarz and the large sieve bound these by mean-square discrepancies of primes in progressions over short intervals of length  $H \asymp N^{\frac{1}{2}+\varepsilon}$ . The short-interval hypothesis (101) (with  $\theta > \frac{1}{2}$ ) then yields, for every  $A > 0$ ,

$$\int_{\mathfrak{m}} S(\alpha)^2 e(-N\alpha) d\alpha \ll_{A,\varepsilon} \frac{N}{(\log N)^A}. \quad (106)$$

(See the dispersion/large-sieve treatment in [11, Chs. 17-18] or [10, Chs. 17, 28]; the short-interval input replaces the classical BV step.)

Combining (104), (105), and (106),

$$R_2(N) \geq \frac{(2C_2)N}{\log^2 N} - K_m \frac{N}{(\log N)^A} \quad (107)$$

for some  $K_m = K_m(A, \varepsilon)$ . Choosing  $A \geq 3$  and  $N_0$  so that  $\frac{2C_2}{\log^2 N_0} > \frac{K_m}{(\log N_0)^A}$  gives  $R_2(N) > 0$  for all even  $N \geq N_0$ . The exhaustive computation up to  $2n_*$  covers the remaining  $N < N_0$ .  $\square$

*Remark* (Weaker Sufficient Inputs for the Reduction). The corollary is proved with a minor-arc bound of the form

$$\int_m S(\alpha)^2 e(-N\alpha) d\alpha \ll_{A,\varepsilon} \frac{N}{(\log N)^A} \quad \text{for some } A > 2, \quad (108)$$

with  $S(\alpha)$  as in (103). The full hypothesis (101) is a convenient sufficient condition for (108), but it is not necessary. Any of the following implies (108) and hence the corollary:

- (i) *Short-Interval BDH/ $L^2$ -type Estimate.* There exist  $\delta, \varepsilon > 0$  such that, for every  $A > 0$ ,

$$\sum_{q \leq N^{\frac{1}{2}+\delta}} \sum_{\substack{a \bmod q \\ (a,q)=1}} \max_{H \geq N^{\frac{1}{2}+\varepsilon}} \int_N^{2N} \left| \pi(x+H; q, a) - \frac{H}{\varphi(q) \log x} \right|^2 dx \ll_{A,\varepsilon} \frac{N^2}{(\log N)^A}. \quad (109)$$

Via Vaughan's identity, Cauchy-Schwarz and the large sieve, this delivers (108).

- (ii) *Almost-Everywhere Short-Interval Equidistribution.* For some  $\delta, \varepsilon > 0$  and every  $A > 0$ , all but  $O(N/(\log N)^A)$  starting points  $x \in [N, 2N]$  satisfy

$$\max_{q \leq N^{\frac{1}{2}+\delta}} \max_{(a,q)=1} \max_{H \geq N^{\frac{1}{2}+\varepsilon}} \left| \pi(x+H; q, a) - \frac{H}{\varphi(q) \log x} \right| \ll \frac{H}{(\log N)^A}. \quad (110)$$

This yields (108) after integrating over  $x$  and summing dyadically.

- (iii) *Any stronger hypothesis implying (i) or (ii)* (e.g. a GEH/EH-type statement in short intervals, or BV in short intervals for a rich class of moduli together with a standard dispersion argument).

Thus the reduction is robust: it requires only that the minor-arc contribution be smaller than the major-arc main term by a fixed power of  $\log N$ . The full SI-BV $_\theta$  statement (101) is one natural way to guarantee this, but strictly weaker inputs suffice.

*Remark* (Scope, Logical Independence, and Status of the Reduction). All bounds stated as theorems in this paper are *unconditional*. In particular, the certified windowed lower bound (Theorem 1; cf. (70) with the product defined in (69)) is proved via a one-sided sieve on  $n \pm m$  with explicit Euler-product factors and a finite edge term; its validity is independent of any circle-method or distributional hypothesis. The accompanying computations are exhaustive on the stated range.

Separately, Corollary 1 is an *unconditional reduction*: it proves the implication

$$(101) \implies \text{Goldbach for all sufficiently large even } N, \quad (111)$$

without further assumptions. The corollary is "conditional" only in the sense that the antecedent (101) is not established here.

Finally, this work does not by itself yield an unconditional proof that every sufficiently large even  $N$  is a Goldbach number. The classical parity barrier prevents pushing a lower-bound sieve to the prime-forcing threshold  $z \asymp \sqrt{N}$  with a uniform positive constant. Thus a full resolution requires additional short-interval *equidistribution* input of the SI-BV type; this author's contribution is to isolate this precise reduction while providing a certified sieve bound and comprehensive data that are logically independent of it.

*Remark* (How the Reduction is Used). Assume the minor-arc input (108) with some  $A > 2$  (e.g. under SI-BV $_{\theta}$ ).

(i) *Positivity.* Since  $\mathfrak{S}(N) \geq S_0 = 2C_2$  and

$$R_2(N) = \frac{\mathfrak{S}(N)N}{\log^2 N} + O\left(\frac{N}{(\log N)^A}\right), \quad (112)$$

there exists  $N_0$  with  $R_2(N) > 0$  for all even  $N \geq N_0$ .

(ii) *Tail From the Certified Product-Form Bound.* From Theorem 1, for  $n \geq n_*$  and a fixed window  $M(n) \asymp n$  (e.g.  $M(n) = \lfloor n/2 \rfloor$ ),

$$\mathcal{G}(n; M) \geq \frac{C_-(n)}{\log^2 n} M(n), \quad C_-(n) = \log^2 n \prod_{3 \leq p \leq \sqrt{n}} \left(1 - \frac{1}{p-1}\right) \prod_{3 \leq p \leq \sqrt{n+M(n)}} \left(1 - \frac{1}{p-1}\right). \quad (113)$$

Set  $\kappa := \inf_{n \geq n_*} C_-(n) > 0$  (the recorded positive margin). Then for all  $N \geq 2n_*$ ,

$$R_2(N) \geq \frac{c_{\text{prod}} N}{\log^2 N}, \quad c_{\text{prod}} = c(S_0, \kappa) > 0, \quad (114)$$

by a dyadic decomposition and the comparison  $M(n) \asymp n$ .

*Note (Asymptotic Surrogate).* If, instead of (113), work with the asymptotic surrogate  $C_-^{\text{asym}}(n)$  (replacing the products by their Mertens asymptotics involving  $K_{\text{EM}}$ ) is chosen, then a slightly larger threshold  $n_*^{\text{asym}} \geq n_*$  so that  $C_-^{\text{asym}}(n) \leq C_-(n)$  for all  $n \geq n_*^{\text{asym}}$  should be used. The same conclusion then holds for all  $N \geq 2n_*^{\text{asym}}$  with a (possibly smaller) constant  $c_{\text{asym}} > 0$ .

**Definition 8** (Exponential Sum and Major/Minor Arcs).

Set  $e(t) := e^{2\pi it}$  and

$$S(\alpha) := \sum_{n \leq N} \Lambda(n) e(n\alpha), \quad \alpha \in [0, 1]. \quad (115)$$

Fix  $Q := N^{1/2-\delta}$  with  $0 < \delta < \theta - \frac{1}{2}$ . For each reduced fraction  $a/q$  with  $1 \leq q \leq Q$  and  $(a, q) = 1$ , define the major arc

$$\mathfrak{M}(q, a) := \left\{ \alpha \in [0, 1] : \left| \alpha - \frac{a}{q} \right| \leq \frac{1}{2qQ} \right\}. \quad (116)$$

Let  $\mathfrak{M} := \bigcup_{1 \leq q \leq Q} \bigcup_{(a, q)=1} \mathfrak{M}(q, a)$  and  $\mathfrak{m} := [0, 1] \setminus \mathfrak{M}$ .

## A Certified Enclosures for Euler Products

### A.1 Mertens Product Enclosure

**Lemma A.1** (Explicit Mertens Enclosure [18, 2]).

There exists  $x_0$  (e.g.  $x_0 = 5557$ ) such that for all  $x \geq x_0$ ,

$$e^{-\gamma} \frac{1}{\log x} \left(1 - \frac{1}{20 \log^3 x} - \frac{316}{\log^4 x}\right) \leq M(x) \leq e^{-\gamma} \frac{1}{\log x} \left(1 + \frac{1}{20 \log^3 x} + \frac{3}{16 \log^4 x} + \frac{1.02}{(x-1) \log x}\right), \quad (\text{A.1})$$

hence

$$|\log x \cdot M(x) - e^{-\gamma}| \leq e^{-\gamma} E_M(x), \quad (\text{A.2})$$

with

$$E_M(x) := \frac{1}{20 \log^3 x} + \max \left\{ \frac{316}{\log^4 x}, \frac{3}{16 \log^4 x} + \frac{1.02}{(x-1) \log x} \right\}. \quad (\text{A.3})$$

## A.2 $C_2$ Tail Bound

**Lemma A.2** (Certified Tail for  $C_2$ ).  
For all  $x \geq 3$ ,

$$0 \leq 1 - \frac{C_2(x)}{C_2} \leq T(x), \quad T(x) := \frac{1}{x-1} + \frac{1}{3(x-1)^3}, \quad (\text{A.4})$$

so  $|C_2(x) - C_2| \leq T(x)$ .

Combining the lemmas,

$$|\log x \cdot P(x) - C_-^{(1)}| \leq e^{-\gamma} C_2 E_M(x) + e^{-\gamma} T(x) = \varepsilon_P(x), \quad (\text{A.5})$$

which is explicit and strictly decreasing for  $x \geq x_0$ .

## A.3 Shifted Product Enclosure

**Lemma A.3** (Certified Enclosure for the Shifted Product).  
Define

$$P(x) := \prod_{\substack{3 \leq p \leq x \\ p \text{ prime}}} \left(1 - \frac{1}{p-1}\right). \quad (\text{A.6})$$

There exists  $x_0$  such that for all  $x \geq x_0$ ,

$$|\log x \cdot P(x) - C_-^{(1)}| \leq \varepsilon_P(x), \quad (\text{A.7})$$

where  $C_-^{(1)} = e^{-\gamma} C_2$  and

$$\varepsilon_P(x) := C_2 E_M(x) + e^{-\gamma} T(x) + E_M(x) T(x). \quad (\text{A.8})$$

Here  $E_M(x)$  and  $T(x)$  are explicit, strictly decreasing functions given in (A.11) and (A.12) below.

*Proof.*

For  $p \geq 3$ ,

$$1 - \frac{1}{p-1} = \left(1 - \frac{1}{p}\right) \left(1 - \frac{1}{(p-1)^2}\right), \quad (\text{A.9})$$

so with

$$M(x) := \prod_{p \leq x} \left(1 - \frac{1}{p}\right), \quad C_2(x) := \prod_{\substack{3 \leq p \leq x \\ p \text{ prime}}} \left(1 - \frac{1}{(p-1)^2}\right), \quad (\text{A.10})$$

we have  $P(x) = M(x) C_2(x)$ .

For Mertens' product we use the explicit enclosure

$$|\log x \cdot M(x) - e^{-\gamma}| \leq E_M(x), \quad x \geq x_0, \quad (\text{A.11})$$

with  $E_M(x)$  strictly decreasing to 0. For the twin-prime factor we use the monotone tail bound

$$0 < C_2 - C_2(x) \leq T(x), \quad T(x) := \sum_{p>x} \frac{1}{(p-1)^2}, \quad (\text{A.12})$$

which is strictly decreasing in  $x$  and satisfies  $T(x) \leq \sum_{n>x} \frac{1}{(n-1)^2} \leq \frac{1}{x-1}$ .

Write  $C_2(x) = C_2 - \delta(x)$  with  $0 \leq \delta(x) \leq T(x)$ . Then

$$\log x \cdot P(x) = (\log x \cdot M(x)) C_2(x) = (e^{-\gamma} \pm E_M(x)) (C_2 - \delta(x)). \quad (\text{A.13})$$

Expanding and bounding the error terms gives

$$\left| \log x \cdot P(x) - e^{-\gamma} C_2 \right| \leq C_2 E_M(x) + e^{-\gamma} \delta(x) + E_M(x) \delta(x) \leq C_2 E_M(x) + e^{-\gamma} T(x) + E_M(x) T(x), \quad (\text{A.14})$$

which is (A.7)-(A.8). The monotonicity of  $E_M, T$  makes  $\varepsilon_P$  strictly decreasing as well.  $\square$

#### A.4 Application to the Lower Bound Product

Let

$$F_1(n) = \log n \cdot P(\sqrt{n}), \quad F_2(n) = \log\left(\frac{3n}{2}\right) \cdot P\left(\sqrt{\frac{3n}{2}}\right), \quad \widehat{C}_- := (e^{-\gamma} C_2)^2. \quad (\text{A.15})$$

Let  $x_1 = \sqrt{n}$  and  $x_2 = \sqrt{\frac{3n}{2}}$ . By Appendix A.3, for  $i = 1, 2$ ,

$$\left| F_i(n) - \widehat{C}_-^{(1)} \right| \leq \varepsilon_P(x_i), \quad (\text{A.16})$$

and hence

$$\left| F_1(n)F_2(n) - \widehat{C}_- \right| \leq \widehat{C}_- (\varepsilon_P(x_1) + \varepsilon_P(x_2)) + \varepsilon_P(x_1)\varepsilon_P(x_2) := \varepsilon(n), \quad (\text{A.17})$$

with  $\varepsilon(n)$  explicit and strictly decreasing in  $n$ .

## B Window rescaling

**Lemma B.1** (Window Rescaling Without Re-Certification).

Fix  $\alpha_0 \in (0, 1)$  and suppose the certified lower bound

$$\mathcal{G}(n; \alpha_0 n) \geq \frac{\mathcal{C}_{-, \alpha_0}(n)}{\log^2 n} (\alpha_0 n) \quad (\text{B.1})$$

holds for all sufficiently large  $n$ , where

$$\mathcal{C}_{-, \alpha}(n) := \log^2 n \prod_{3 \leq p \leq \sqrt{n}} \left(1 - \frac{1}{p-1}\right) \prod_{3 \leq p \leq \sqrt{n+\alpha n}} \left(1 - \frac{1}{p-1}\right). \quad (\text{B.2})$$

Then for every  $\alpha \in (0, \alpha_0]$  and all sufficiently large  $n$ ,

$$\mathcal{G}(n; \alpha n) \geq \frac{\mathcal{C}_{-, \alpha}(n)}{\log^2 n} (\alpha n) .$$

(B.3)

*Proof.*

Summing the one-sided lower bounds over  $|m| \leq \alpha n$  proceeds exactly as in the  $\alpha_0$  case. Shrinking the window reduces the number of offsets linearly by  $\alpha/\alpha_0$ , while

$$\sqrt{n + \alpha n} \leq \sqrt{n + \alpha_0 n} \quad (\text{B.4})$$

tightens the right-edge cutoff in the second Euler product, which can only *increase* the conservative product in (B.2). Hence the same certification yields (B.3) for all  $\alpha \leq \alpha_0$ .  $\square$

*Remark* (Checking Validity for Rescaling). While Lemma B.1 allows rescaling from  $\alpha_0$  to any  $\alpha \leq \alpha_0$ , the Mertens certification threshold  $n_{\text{Mertens}}$  depends on  $\alpha$ . Table 8 provides empirical Mertens positions and convergence statistics for 82 different  $\alpha$  values. Before applying the rescaling lemma, readers should consult this table to verify that the target  $\alpha$  has sufficient data support (i.e.,  $n_{\text{Mertens}}$  is well-separated from the dataset boundary) to justify the rescaling without requiring new certification computations. In practice, so long as  $n \geq n_{\text{Mertens}}$  for the chosen  $\alpha$ , the rescaling should yield a normalized lower bound of approximately  $\gtrsim 2$ .

**Corollary B.1** (Monotone Extension to Larger Windows).

Under the hypotheses of Lemma B.1, for every  $\alpha \in [\alpha_0, 1)$  and all sufficiently large  $n$ ,

$$\boxed{\mathcal{G}(n; \alpha n) \geq \mathcal{G}(n; \alpha_0 n) \geq \frac{\mathcal{C}_{-,n}(\alpha_0)}{\log^2 n} (\alpha_0 n).} \quad (\text{B.5})$$

*Proof.*

Monotonicity in the window is immediate from the set inclusion

$$\{ |m| \leq \alpha_0 n \} \subseteq \{ |m| \leq \alpha n \}, \quad (\text{B.6})$$

which implies  $\mathcal{G}(n; \alpha n) \geq \mathcal{G}(n; \alpha_0 n)$ . The second inequality in (B.5) is exactly (B.1).  $\square$

*Remark* (Uniform-in- $\alpha$  Certification). Because the one-sided sieve factors  $S_{\pm}(n, m)$  are pointwise in  $m$ , the same argument that proves (B.1) works verbatim for each fixed  $\alpha \in (0, 1)$ ; in particular, for all sufficiently large  $n$ ,

$$\mathcal{G}(n; \alpha n) \geq \frac{\mathcal{C}_{-,n}(\alpha)}{\log^2 n} (\alpha n). \quad (\text{B.7})$$

No re-tuning of sieve weights is required; only the right-edge cutoff  $\sqrt{n + \alpha n}$  in  $\mathcal{C}_{-,n}(\alpha)$  changes.

## C Decadal Statistics for Goldbach Pair Distribution

Table 4: Per-decade Statistics for Goldbach Pair Counts for  $|m| \in [1, \lfloor \frac{n}{2} \rfloor]$

Dec.	Min At	Min	Max At	Max	$n_{\text{geom}}$	$\langle \text{Count} \rangle$
0	4	2	4	2	4	2.0
0	5	2	5	2	5	2.0
0	6	2	6	2	7	2.0
0	7	0	7	0	7	0.0
0	8	2	8	2	9	2.0
0	9	4	9	4	9	4.0
1	11	0	12	4	15	2.2
1	22	2	21	6	25	3.2
1	31	2	30	8	35	4.2
1	43	0	45	10	45	4.2
1	53	2	57	10	55	5.8
1	61	2	60	12	65	6.0
1	79	2	75	14	75	7.8
1	82	4	81	10	85	7.0
1	97	2	90	12	95	7.8
2	107	4	195	26	141	10.6
2	223	4	210	30	245	14.7
2	302	8	315	40	347	19.1
2	433	8	495	50	447	23.0
2	508	14	570	56	547	26.1
2	601	14	660	62	649	29.8
2	706	14	735	72	749	33.7
2	802	16	840	76	849	36.6
2	919	18	975	78	949	38.3
3	1 009	20	1 995	148	1 415	54.9
3	2 029	30	2 730	208	2 449	80.4
3	3 076	44	3 990	250	3 465	103.7
3	4 051	60	4 830	310	4 473	126.3
3	5 416	72	5 775	358	5 477	146.6
3	6 353	88	6 930	424	6 481	169.5
3	7 219	94	7 770	442	7 483	187.0
3	8 116	112	8 925	520	8 485	206.4
3	9 014	124	9 975	544	9 487	225.9
4	10 462	134	19 635	990	14 143	323.9
4	20 023	234	28 665	1 312	24 495	488.5
4	30 332	332	39 270	1 790	34 641	641.1
4	40 597	416	49 665	2 050	44 721	785.9
4	51 826	516	58 905	2 476	54 773	926.6
4	60 413	604	69 615	2 826	64 807	1 064.8
4	71 633	676	78 540	3 108	74 833	1 194.1
4	80 441	786	87 780	3 374	84 853	1 324.8
4	91 958	860	98 175	3 708	94 869	1 455.4
5	101 467	948	195 195	6 716	141 421	2 117.9
5	204 928	1 688	285 285	9 808	244 949	3 252.3
5	300 739	2 396	390 390	12 048	346 411	4 319.0
5	401 509	3 044	495 495	14 828	447 213	5 340.3
5	500 417	3 742	570 570	17 786	547 723	6 334.5

(continued)

Dec.	Min At	Min	Max At	Max	$n_{\text{geom}}$	$\langle \text{Count} \rangle$
5	603 182	4 352	690 690	20 546	648 075	7 298.4
5	700 268	4 948	765 765	22 942	748 331	8 241.7
5	804 191	5 550	855 855	25 114	848 529	9 177.1
5	909 037	6 154	990 990	26 788	948 683	10 089.6
6	1 004 449	6 742	1 996 995	51 734	1 414 213	14 890.7
6	2 012 212	12 360	2 984 520	71 382	2 449 489	23 157.9
6	3 004 042	17 494	3 993 990	94 150	3 464 101	31 002.9
6	4 015 034	22 544	4 849 845	118 980	4 472 135	38 562.7
6	5 001 482	27 418	5 870 865	139 510	5 477 225	45 926.9
6	6 002 812	32 242	6 891 885	152 328	6 480 741	53 114.6
6	7 010 638	36 882	7 912 905	177 818	7 483 315	60 199.5
6	8 007 488	41 544	8 843 835	195 128	8 485 281	67 166.4
6	9 001 429	46 072	9 699 690	217 942	9 486 833	74 015.4
7	10 030 684	50 364	19 399 380	400 846	14 142 135	110 283.3
7	20 007 184	93 132	29 099 070	572 870	24 494 897	173 140.1
7	30 032 203	133 266	38 798 760	738 184	34 641 017	233 156.3
7	40 002 659	172 084	48 498 450	900 422	44 721 359	291 303.5
7	50 008 249	209 830	58 198 140	1 060 096	54 772 255	348 071.9
7	60 010 597	246 670	67 897 830	1 213 536	64 807 407	403 718.9
7	70 017 487	282 866	77 597 520	1 367 996	74 833 147	458 571.4
7	80 015 692	318 898	87 297 210	1 518 344	84 852 813	512 553.2
7	90 020 452	353 874	99 804 705	1 692 366	94 868 329	565 927.0

Table 5: Normalized by  $\frac{\log^2 n}{M}$  Per-Decade Statistics for Goldbach Pair Counts for  $|m| \in [1, \lfloor \frac{n}{2} \rfloor]$

<b>Dec.</b>	$n_0$	$C_{\min}(n_0)$	$n_1$	$C_{\max}(n_1)$	$n_{\text{geom}}$	$C_{\text{avg}}$
0	4	1.9218	4	1.9218	4	1.92181
0	5	2.5903	5	2.5903	5	2.59029
0	6	2.1403	6	2.1403	6	2.14027
0	7	0.0000	7	0.0000	7	0.00000
0	8	2.1620	8	2.1620	8	2.16204
0	9	4.8278	9	4.8278	9	4.82780
1	11	0.0000	15	4.1906	15	2.22523
1	28	1.5862	21	5.5615	25	2.76778
1	37	1.4487	30	6.1697	35	3.11072
1	43	0.0000	45	6.5867	45	2.74658
1	59	1.1466	57	5.8380	55	3.44494
1	64	1.0810	60	6.7055	65	3.26267
1	79	0.9791	75	7.0532	75	3.92285
1	89	1.8316	81	4.8278	85	3.28736
1	97	0.8720	90	5.3995	95	3.43676
2	199	1.1321	105	8.3305	141	3.58341
2	223	1.0536	210	8.1690	245	3.60234
2	379	1.4922	315	8.4311	347	3.75489
2	433	1.3650	420	7.9919	447	3.81991
2	569	1.9839	570	7.9121	547	3.78829
2	661	1.7890	660	7.9189	649	3.84952
2	706	1.7065	735	8.5455	749	3.94429
2	802	1.7842	840	8.2041	849	3.91997
2	967	1.7610	975	7.5867	949	3.79061
3	1402	1.6476	1155	9.1356	1415	3.93033
3	2029	1.7158	2730	9.5391	2449	3.93819
3	3076	1.8453	3465	9.2051	3465	3.94779
3	4801	1.8562	4620	9.4320	4473	3.97557
3	5416	1.9651	5775	9.3025	5477	3.95664
3	6353	2.1246	6930	9.5702	6481	4.02160
3	7219	2.0559	7770	9.1297	7483	3.97249
3	8777	2.1795	8925	9.6435	8485	3.97681
3	9649	2.1637	9240	9.6375	9487	3.99115
4	11272	2.1315	15015	10.4223	14143	4.00416
4	20816	2.2799	21945	10.0363	24495	4.01074
4	35792	2.2977	30030	10.2932	34641	4.01184
4	40597	2.3078	45045	10.2676	44721	4.01124
4	51826	2.3466	58905	10.1422	54773	4.01506
4	67904	2.4136	60060	10.2886	64807	4.02457
4	71633	2.3588	75075	10.1865	74833	4.01279
4	89459	2.3832	87780	9.9601	84853	4.01645
4	92357	2.4345	90090	10.3847	94869	4.02541
5	116728	2.4025	150150	10.4044	141421	4.02084
5	204928	2.4642	255255	10.9888	244949	4.02288
5	366794	2.4992	345345	10.8231	346411	4.02349
5	463549	2.5131	435435	10.8082	447213	4.02272
5	548461	2.5320	510510	11.0269	547723	4.02481
5	686398	2.5271	690690	10.7554	648075	4.02369
5	770558	2.5323	765765	10.9991	748331	4.02222

(continued)

Dec.	$n_0$	$C_{\min}(n_0)$	$n_1$	$C_{\max}(n_1)$	$n_{\text{geom}}$	$C_{\text{avg}}$
5	804 191	2.552 0	855 855	10.950 6	848 529	4.025 35
5	915 961	2.547 1	930 930	10.674 7	948 683	4.024 42
6	1 201 553	2.553 5	1 276 275	11.043 5	1 414 213	4.023 67
6	2 053 553	2.579 8	2 042 040	11.036 4	2 449 489	4.023 69
6	3 004 042	2.591 1	3 573 570	11.047 5	3 464 101	4.023 94
6	4 792 159	2.588 5	4 849 845	11.628 0	4 472 135	4.023 15
6	5 167 067	2.597 6	5 870 865	11.544 5	5 477 225	4.023 42
6	6 175 451	2.603 3	6 561 555	11.429 8	6 480 741	4.022 32
6	7 376 626	2.610 5	7 402 395	11.421 2	7 483 315	4.023 27
6	8 143 934	2.607 6	8 273 265	11.322 4	8 485 281	4.023 60
6	9 121 549	2.613 9	9 699 690	11.630 4	9 486 833	4.022 61
7	10 030 684	2.609 8	14 549 535	11.638 0	14 142 135	4.022 24
7	24 496 594	2.621 7	29 099 070	11.629 7	24 494 897	4.021 85
7	30 099 763	2.626 0	38 798 760	11.618 7	34 641 017	4.021 57
7	41 344 276	2.629 5	48 498 450	11.629 2	44 721 359	4.021 29
7	53 699 671	2.633 0	58 198 140	11.645 8	54 772 255	4.021 12
7	66 759 878	2.632 3	67 897 830	11.624 9	64 807 407	4.020 61
7	78 822 322	2.634 3	77 597 520	11.636 9	74 833 147	4.021 11
7	82 476 448	2.635 8	82 447 365	11.630 5	84 852 813	4.020 56
7	96 281 998	2.635 6	96 996 900	11.629 5	94 868 329	4.020 44

*Remark.* Primorials consistently correspond to maxima. Many unnormalized binned maxima have occurred at values equal to  $19\#$  or its multiples, and many of the normalized maxima align with these values as well. In contrast, the minima are more likely to occur at values that are either prime or semiprime.

Table 6: Normalized by  $\frac{\log^2 n}{M}$  Per-Decade HL-A Predictions for Goldbach Pair Counts for  $|m| \in [1, \lfloor \frac{n}{2} \rfloor]$

Dec.	$\dot{n}_0$	$\dot{C}_{\min}(n_0)$	$\dot{n}_1$	$\dot{C}_{\max}(n_1)$	$n_{\text{geom}}$	$\dot{C}_{\text{avg}}(n_{\text{geom}})$
0	4	2.870 1	4	2.870 1	4	2.640 65
0	5	4.266 1	5	4.266 1	5	4.266 09
0	6	6.218 9	6	6.218 9	6	6.218 87
0	7	3.393 0	7	3.393 0	7	3.393 04
0	8	2.814 6	8	2.814 6	8	2.814 64
0	9	5.820 4	9	5.820 4	9	5.820 37
1	16	2.765 1	15	7.380 0	15	4.018 77
1	29	2.855 7	21	6.664 3	25	4.020 54
1	32	2.734 6	30	7.340 7	35	4.330 99
1	47	2.783 0	45	7.286 7	45	4.052 32
1	59	2.764 6	51	5.803 6	55	3.923 09
1	64	2.715 4	60	7.244 9	65	4.344 50
1	79	2.745 9	75	7.231 4	75	4.043 42
1	89	2.743 2	84	6.503 4	85	3.910 75
1	97	2.738 6	90	7.231 9	95	4.320 55
2	128	2.702 5	105	8.666 9	141	4.080 98
2	256	2.693 3	210	8.629 8	245	4.064 24
2	397	2.696 1	315	8.611 2	347	4.097 36
2	499	2.692 0	420	8.602 2	447	4.057 44
2	512	2.686 4	525	8.597 1	547	4.051 35
2	691	2.687 9	630	8.592 1	649	4.082 69
2	797	2.686 5	735	8.587 2	749	4.065 58
2	887	2.685 2	840	8.584 1	849	4.041 22
2	997	2.684 2	945	8.582 0	949	4.079 48
3	1 024	2.681 1	1 155	9.530 1	1 415	4.057 33
3	2 048	2.676 9	2 310	9.516 0	2 449	4.050 96
3	3 989	2.674 3	3 465	9.508 7	3 465	4.053 81
3	4 096	2.673 5	4 620	9.504 0	4 473	4.048 63
3	5 987	2.672 3	5 775	9.500 6	5 477	4.045 53
3	6 997	2.671 7	6 930	9.498 1	6 481	4.048 93
3	7 993	2.671 1	7 140	9.117 7	7 483	4.045 50
3	8 192	2.670 7	8 085	9.496 0	8 485	4.043 51
3	9 973	2.670 2	9 240	9.494 2	9 487	4.047 20
4	16 384	2.668 3	15 015	10.350 8	14 143	4.042 46
4	29 989	2.666 6	21 945	10.042 0	24 495	4.039 79
4	32 768	2.666 3	30 030	10.342 8	34 641	4.038 77
4	49 999	2.665 2	45 045	10.338 6	44 721	4.037 48
4	59 999	2.664 8	58 905	10.106 3	54 773	4.036 63
4	65 536	2.664 5	60 060	10.335 8	64 807	4.036 30
4	79 999	2.664 1	75 075	10.333 8	74 833	4.035 60
4	89 989	2.663 8	87 780	10.028 5	84 853	4.035 12
4	99 991	2.663 6	90 090	10.332 2	94 869	4.035 00
5	131 072	2.663 0	105 105	10.330 9	141 421	4.033 55
5	262 144	2.661 6	255 255	11.012 3	244 949	4.031 93
5	399 989	2.660 9	345 345	10.813 4	346 411	4.031 02
5	499 979	2.660 5	435 435	10.702 6	447 213	4.030 33
5	524 288	2.660 4	510 510	11.007 3	547 723	4.029 82
5	699 967	2.660 0	690 690	10.808 8	648 075	4.029 45
5	799 999	2.659 8	765 765	11.004 7	748 331	4.029 09

(continued)

<b>Dec.</b>	$\dot{n}_0$	$\dot{C}_{\min}(n_0)$	$\dot{n}_1$	$\dot{C}_{\max}(n_1)$	$n_{\text{geom}}$	$\dot{C}_{\text{avg}}(n_{\text{geom}})$
5	899 981	2.659 6	855 855	10.923 0	848 529	4.028 79
5	999 983	2.659 4	930 930	10.671 4	948 683	4.028 58
6	1 048 576	2.659 3	1 021 020	11.002 9	1 414 213	4.027 68
6	2 097 152	2.658 4	2 042 040	10.998 9	2 449 489	4.026 57
6	3 999 971	2.657 6	3 063 060	10.996 7	3 464 101	4.025 91
6	4 194 304	2.657 5	4 849 845	11.641 2	4 472 135	4.025 45
6	5 999 993	2.657 1	5 870 865	11.517 1	5 477 225	4.025 09
6	6 999 997	2.656 9	6 561 555	11.431 8	6 480 741	4.024 81
6	7 999 993	2.656 8	7 402 395	11.399 6	7 483 315	4.024 56
6	8 388 608	2.656 7	8 273 265	11.315 3	8 485 281	4.024 35
6	9 999 991	2.656 6	9 699 690	11.637 8	9 486 833	4.024 18
7	16 777 216	2.656 0	14 549 535	11.636 0	14 142 135	4.023 55
7	29 999 999	2.655 5	24 249 225	11.633 8	24 494 897	4.022 74
7	33 554 432	2.655 4	33 948 915	11.632 5	34 641 017	4.022 25
7	49 999 991	2.655 0	43 648 605	11.631 5	44 721 359	4.021 91
7	59 999 999	2.654 9	53 348 295	11.630 7	54 772 255	4.021 65
7	67 108 864	2.654 8	63 047 985	11.630 1	64 807 407	4.021 43
7	79 999 987	2.654 6	72 747 675	11.629 6	74 833 147	4.021 25
7	89 999 999	2.654 5	82 447 365	11.629 1	84 852 813	4.021 10
7	99 999 989	2.654 4	92 147 055	11.628 7	94 868 329	4.020 96

Table 7:  $\Lambda$  Calculations for Euler Product Series Products

Dec.	$n_0$	$C_{\min}$	$C_-$	$C_{\min} - C_-$	$C_-^{\text{asymp}}$	$C_{\min} - C_-^{\text{asymp}}$
0	4	1.922	0.961	0.961	1.701	0.221
0	5	2.590	1.295	1.295	1.756	0.835
0	6	2.140	1.605	0.535	1.793	0.348
0	7	0.000	1.893	-1.893	1.819	-1.819
0	8	2.162	2.162	0.000	1.840	0.323
0	9	4.828	1.207	3.621	1.856	2.972
1	11	0.000	1.438	-1.438	1.880	-1.880
1	28	1.586	1.561	0.025	1.960	-0.374
1	37	1.449	1.528	-0.079	1.976	-0.528
1	43	0.000	1.658	-1.658	1.984	-1.984
1	59	1.147	1.624	-0.477	1.999	-0.853
1	64	1.081	1.689	-0.608	2.003	-0.922
1	79	0.979	1.865	-0.885	2.012	-1.032
1	89	1.832	1.771	0.061	2.016	-0.184
1	97	0.872	1.839	-0.967	2.019	-1.147
2	199	1.132	1.746	-0.614	2.042	-0.910
2	223	1.054	1.822	-0.768	2.045	-0.991
2	379	1.492	1.754	-0.261	2.058	-0.565
2	433	1.365	1.833	-0.468	2.061	-0.696
2	569	1.984	1.843	0.141	2.066	-0.082
2	661	1.789	1.866	-0.077	2.069	-0.280
2	706	1.707	1.904	-0.198	2.070	-0.364
2	802	1.784	1.979	-0.195	2.073	-0.288
2	967	1.761	1.895	-0.134	2.076	-0.315
3	1 402	1.648	1.948	-0.301	2.082	-0.434
3	2 029	1.716	1.966	-0.250	2.087	-0.371
3	3 076	1.845	1.996	-0.151	2.093	-0.247
3	4 801	1.856	2.006	-0.150	2.098	-0.242
3	5 416	1.965	1.983	-0.018	2.099	-0.134
3	6 353	2.125	2.010	0.115	2.101	0.024
3	7 219	2.056	2.003	0.053	2.102	-0.046
3	8 777	2.180	2.012	0.168	2.104	0.075
3	9 649	2.164	2.033	0.131	2.105	0.059
4	11 272	2.132	2.045	0.087	2.107	0.025
4	20 816	2.280	2.063	0.217	2.112	0.168
4	35 792	2.298	2.077	0.220	2.116	0.181
4	40 597	2.308	2.058	0.250	2.117	0.191
4	51 826	2.347	2.076	0.271	2.119	0.228
4	67 904	2.414	2.078	0.336	2.121	0.293
4	71 633	2.359	2.090	0.269	2.121	0.238
4	89 459	2.383	2.090	0.293	2.123	0.261
4	92 357	2.435	2.096	0.338	2.123	0.312
5	116 728	2.403	2.104	0.299	2.124	0.278
5	204 928	2.464	2.106	0.358	2.128	0.337
5	366 794	2.499	2.110	0.389	2.131	0.368
5	463 549	2.513	2.106	0.407	2.132	0.381
5	548 461	2.532	2.112	0.420	2.133	0.399
5	686 398	2.527	2.115	0.412	2.134	0.393
5	770 558	2.532	2.114	0.418	2.134	0.398
5	804 191	2.552	2.112	0.440	2.135	0.418
5	915 961	2.547	2.118	0.429	2.135	0.412

(continued)

<b>Dec.</b>	$n_0$	$C_{\min}(n_0)$	$C_-(n_0)$	$C_{\min}(n_0) - C_-(n_0)$	$C_-^{\text{asymp}}(n_0)$	$C_{\min}(n_0) - C_-^{\text{asymp}}(n_0)$
6	1 201 553	2.554	2.117	0.437	2.136	0.417
6	2 053 553	2.580	2.125	0.455	2.139	0.441
6	3 004 042	2.591	2.126	0.465	2.140	0.451
6	4 792 159	2.589	2.134	0.455	2.142	0.447
6	5 167 067	2.598	2.131	0.466	2.142	0.456
6	6 175 451	2.603	2.131	0.473	2.143	0.461
6	7 376 626	2.611	2.134	0.477	2.143	0.467
6	8 143 934	2.608	2.134	0.473	2.144	0.464
6	9 121 549	2.614	2.134	0.480	2.144	0.470
7	10 030 684	2.610	2.137	0.473	2.144	0.466
7	24 496 594	2.622	2.142	0.480	2.147	0.475
7	30 099 763	2.626	2.141	0.485	2.148	0.478
7	41 344 276	2.630	2.142	0.487	2.149	0.481
7	53 699 671	2.633	2.144	0.489	2.149	0.484
7	66 759 878	2.632	2.146	0.487	2.150	0.483
7	78 822 322	2.634	2.145	0.489	2.150	0.484
7	82 476 448	2.636	2.146	0.490	2.150	0.486
7	96 281 998	2.636	2.146	0.489	2.151	0.485

## D CPS Summary for Multiple Alpha Values

Table 8 presents the complete CPS (Certified Pair Summary) analysis for 82 different alpha values ranging from  $\alpha = 0.0009765625$  to  $\alpha = 0.5$ , computed over the range  $4 \leq 2n < 23\#$ .

For each alpha value, the table shows:

- **Alpha:** The window parameter  $\alpha$  where  $M = \alpha n$
- **PreMertens, Mertens:** Positions where Mertens-type bounds are validated
- **DeltaMertens:** Normalized deviation at the Mertens position
- **n\_5percent:** The 5% threshold position in the tested range
- **NzeroStat, EtaStat:** Statistical measures of minimum location and deviation
- **MertensAsymp, DeltaMertensAsymp:** Asymptotic analogs of the Mertens measures
- **NzeroStatAsymp, EtaStatAsymp:** Asymptotic statistical measures

**Certification Status:** Only  $\alpha = 0.5$  has been fully certified in this work. The first five rows ( $\alpha < 0.0015$ ) exhibit edge effects where the minimum  $C(n)$  occurs close to the dataset boundary, providing insufficient data to reliably verify the Mertens condition. Values with  $\alpha \geq 0.0078125$  may be suitable for certification but require individual inspection of convergence behavior. All data is provided for transparency and to enable independent verification by other researchers.

Table 8: CPS summary statistics for multiple alpha values.

$\alpha$	PreMertens	Mertens	$\Delta$ Mertens	$n_{5\%}$	$N_0$ Stat	$\eta$ Stat
0.0009765625	111512498	0	0.000000	59742208	111514064	0.007676
0.00106494895768	111469532	0	0.000000	54188513	111512498	0.015869
0.00116133507324	109372352	109613477	0.058396	49147745	111469532	0.018778
0.00126644487759	105556148	105764111	0.045734	44572805	105604172	0.028492
0.001381067932	99297428	99791746	0.096924	40420894	103662451	0.006795
0.00150606525919	84700843	84937997	0.042446	36653127	93011768	0.013970
0.00164237581104	78296887	78405373	0.029156	33234172	80607004	0.019601
0.00179102352188	69815183	70038676	0.043155	30131933	75636646	0.016557
0.001953125	73484561	73620566	0.127773	27317248	74097862	0.005516
0.00212989791536	74097862	74444324	0.080115	24763629	74149787	0.058726
0.00232267014649	46229168	46528574	0.057570	22447010	74097862	0.002601
0.00253288975518	41926246	42189997	0.105026	20345931	61389073	0.007697
0.00276213586401	35480609	35870878	0.068037	18439716	39833536	0.013621
0.00301213051837	35480609	35849614	0.055440	16710764	38236591	0.035569
0.00328475162208	35480609	35873053	0.083271	15143002	39813646	0.033458
0.00358204704377	25265521	25276486	0.043753	13720926	27332504	0.000846
0.00390625	23699191	23775286	0.043396	12431616	24333362	0.003985
0.00425979583072	23699191	23879929	0.042563	11262277	25139893	0.028715
0.00464534029298	22187332	22340671	0.017289	10202267	22340671	0.017289
0.00506577951036	20173091	20459074	0.042176	9241026	22187332	0.013895
0.00552427172802	17193173	17390482	0.103087	8369792	20173091	0.017955
0.00602426103675	14077352	14378131	0.063070	7580017	14577527	0.010423
0.00656950324417	14077352	14450707	0.055862	6864142	14638991	0.024813

0.00716409408754	13067672	13481324	0.066309	6215301	17173202	0.050072
0.0078125	9057157	9417181	0.021194	5627264	9159812	0.011142
0.00851959166145	12973039	13050508	0.090449	5094376	13009069	0.077090
0.00929068058596	7238339	7638497	0.016485	4611503	12973039	0.006079
0.0101315590207	6384029	6651937	0.030187	4173988	7216597	0.008751
0.011048543456	5719144	6068224	0.021493	3777603	5744512	0.015045
0.0120485220735	5506681	5771728	0.026958	3418511	6384029	0.026452
0.0131390064883	4540013	4553048	0.072624	3093309	5051327	0.008551
0.0143281881751	4347407	4362716	0.037246	2798680	4774582	0.011659
0.015625	4343501	4362716	0.065480	2531904	5042162	0.019217
0.0170391833229	3936271	3952873	0.013783	2290310	4290386	0.010622
0.0185813611719	3936271	3952873	0.057429	2071485	4328171	0.045007
0.0202631180414	2451667	2468068	0.065506	1873404	3008938	0.018217
0.0220970869121	2451667	2468068	0.093749	1694070	2737006	0.053774
0.024097044147	2324848	2330203	0.054640	1531724	2368547	0.037871
0.0262780129767	1875931	1882459	0.078942	1384770	2324848	0.007527
0.0286563763501	1726924	1742443	0.036204	1251764	1742443	0.036204
0.03125	1665278	1681271	0.085831	1131392	1726924	0.045435
0.0340783666458	1665278	1681271	0.127754	1022497	2006657	0.056839
0.0371627223438	1665278	1681271	0.135617	923937	2620586	0.072703
0.0405262360828	944362	951124	0.025422	834793	1385326	0.003184
0.0441941738242	1385326	1400299	0.089043	754150	1542473	0.081842
0.048194088294	733504	736279	0.071743	681184	1037518	0.036352
0.0525560259534	611264	627271	0.056101	615211	1037518	0.028871
0.0573127527003	594974	606104	0.087853	555549	648502	0.054999
0.0625	594974	606104	0.106550	501616	638167	0.062160
0.0681567332916	476416	480659	0.053695	452839	594974	0.015143
0.0743254446877	424313	435443	0.096208	408743	476416	0.000840
0.0810524721657	266374	281291	0.005420	368897	594974	0.020242
0.0883883476483	284938	294824	0.109189	332872	594974	0.047356
0.096388176588	284938	286733	0.097151	300328	306707	0.080601
0.105112051907	266374	284938	0.080503	270921	286733	0.076978
0.114625505401	266374	283982	0.104699	244353	286733	0.096563
0.125	119563	120418	0.053110	220352	266374	0.003820
0.136313466583	119563	120418	0.087315	198682	266374	0.056946
0.148650889375	109981	111119	0.048076	179105	266374	0.094326
0.162104944331	109981	111637	0.146522	161433	266374	0.124711
0.176776695297	92381	92489	0.131961	145472	266374	0.092065
0.192776353176	92381	92843	0.146649	131069	201904	0.158435
0.210224103813	58543	59027	0.084778	118070	201904	0.148279
0.229251010801	28921	30749	0.112918	106338	108193	0.109940
0.25	25052	25633	0.059169	95756	108193	0.118388
0.272626933166	20947	22772	0.074661	86206	201904	0.161273
0.297301778751	12574	12671	0.169945	77595	79312	0.202401
0.324209888663	23201	24256	0.168985	69829	82919	0.200254
0.353553390593	23201	24256	0.231317	62826	64984	0.213057
0.385552706352	6046	6184	0.129128	56512	79256	0.233419
0.420448207627	4801	4909	0.154197	50822	65029	0.266569
0.458502021602	4801	4933	0.219921	45693	53624	0.264122
0.5	5416	5557	0.306962	41072	71633	0.269257
0.545253866333	1966	2029	0.027049	36910	51826	0.249996
0.594603557501	1669	1733	0.083101	33160	51826	0.309681
0.648419777326	1109	1114	0.393096	29784	30908	0.318380
0.707106781187	841	859	0.061293	26745	29177	0.317629

0.771105412704	556	571	0.231394	24008	29177	0.383976
0.840896415254	556	577	0.378010	21545	29177	0.415774
0.917004043205	199	214	0.394117	19330	35059	0.466658
1	199	214	0.394117	17571	47638	0.544827

## E Reproducibility

All source code, certification tools, and datasets used in this work are permanently archived on Zenodo [16]. The repository includes build scripts, certification outputs, and checksums to ensure bitwise reproducibility of all results.

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