

1. Let G be a graph with adjacency matrix

$$A = \begin{pmatrix} 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \end{pmatrix}$$

Use the fact that the ij -th entry of A^k gives the number of walks (defined in problem 2, below) of length k from vertex i to vertex j in G to argue whether or not G is connected. Drawing a picture of the graph is not sufficient.

I wrote following MATLAB codes for this question:

```
% @Author: Baihan Lin, @Date: Oct 2016

% initiate adjacency matrix A.
A = [0 0 1 0 0 0 0 0 1;
      0 0 0 1 1 0 0 0 0;
      1 0 0 0 0 0 1 0 0;
      0 1 0 0 1 0 1 0 0;
      0 1 0 1 0 1 0 0 0;
      0 0 0 0 1 0 0 1 1;
      0 0 1 1 0 0 0 0 0;
      0 0 0 0 0 1 0 0 0;
      1 0 0 0 0 1 0 0 0];

% find minimum steps for every vertex can connect to each other (complete).
n = whenComplete(A, 1);

n % minimum steps to reach a complete graph (if n = -1, not connected).
An = A^n % the final complete adjacency graph in n steps.

-----
% to find out in how many steps can each vertex be connected to another one.
function n = whenComplete(A, n)
% A is the input adjacency matrix
% n is the starting steps (set to -1 if it is not connected at all)

if n ~= -1
    % only in connected graphs, we find minimum steps towards complete.
    if prod(sum(A)) == 0 % if one column of 0, not connected, set n to -1.
        n = -1;
    else
        if prod(prod(A^n)) == 0 % not complete graph yet within n step.
            n = whenComplete(A, n+1);
        end
    end
end
end
end
```

Here is the output on MATLAB console:

n =

6

An =

21	15	1	11	10	24	17	1	1
15	41	10	43	43	35	27	9	11
1	10	21	25	17	9	3	7	17
11	43	25	65	46	42	19	10	17
10	43	17	46	73	21	34	25	33
24	35	9	42	21	54	17	2	3
17	27	3	19	34	17	29	8	9
1	9	7	10	25	2	8	12	17
1	11	17	17	33	3	9	17	28

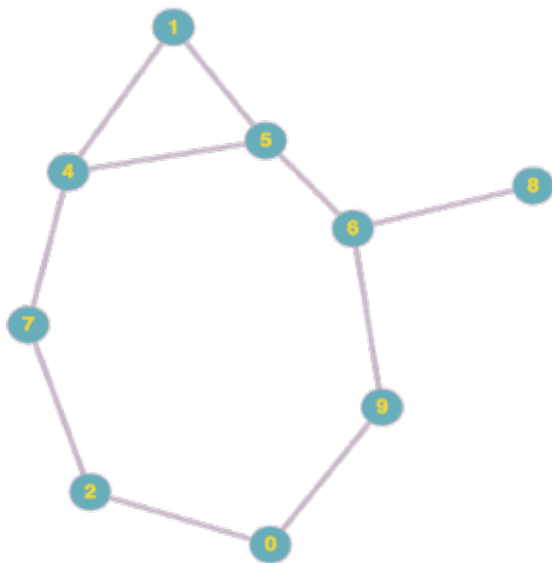


Figure 1. Graph based on adjacency matrix

As shown in Figure 1, based on the adjacency matrix, this graph is connected.

But we can further prove this by definition of this adjacency matrix: A^k gives the number of walks of length k from vertex i to vertex j in graph G .

In another word, if G is connected, there must be a number k of walks to satisfy the path from any vertex i to any vertex j in G .

If so, it means from every vertex, it is possible to travel to another vertex in certain steps, which suggests the connectivity of the graph.

Therefore, I wrote a program to find the minimum steps it takes, to let every vertex in this graph to travel towards any other vertices (more details in comments).

I found that when $n = 6$, every vertex has at least one possible ways to travel to any other vertex in 6 steps. Since A^6 has no zero entries, we know there is a path (of length 6) between every pair of vertices.

Thus, this proves that the graph is connected.

2. Let a and b be positive integers. We say that a divides b iff $b = ak$ for some integer k .

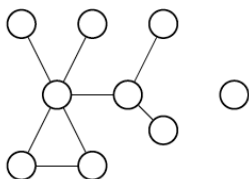
So 2 divides 10, and 3 does not divide 8.

Let $V = \{2, 3, 4, \dots, 22\} = \{j \in \mathbb{Z} : 2 \leq j \leq 22\}$.

Define a graph H with V as its vertex set and edge set E defined by $(v_1, v_2) \in E$ iff $v_1 \neq v_2$ and v_1 divides v_2 or v_2 divides v_1 . So $(2, 6)$ is an edge in H ; $(3, 4)$ is not.

Recall that the edge $(2, 6)$ is the same as the edge $(6, 2)$.

- (a) The *degree sequence* of a graph G is the sequence of the degrees of all vertices in G . For example, the graph below has degree sequence 5, 3, 2, 2, 1, 1, 1, 1, 0 (in decreasing order).



Give the degree sequence of H in decreasing order.

Based on the definition of E , we can enumerate out all combinations of (v_1, v_2) :

For vertex 2, we have 10 edges:

$(2, 4), (2, 6), (2, 8), (2, 10), (2, 12), (2, 14), (2, 16), (2, 18), (2, 20), (2, 22)$
in another word, vertex 2 has degree 10

For vertex 3, we have 6 edges:

$(3, 6), (3, 9), (3, 12), (3, 15), (3, 18), (3, 21)$
in another word, vertex 3 has degree 6

For vertex 4, we have 5 edges:

$(2, 4), (4, 8), (4, 12), (4, 16), (4, 20)$
in another word, vertex 4 has degree 5

For vertex 5, we have 3 edges:

$(5, 10), (5, 15), (5, 20)$
in another word, vertex 5 has degree 3

For vertex 6, we have 4 edges:

$(2, 6), (3, 6), (6, 12), (6, 18)$
in another word, vertex 6 has degree 4

For vertex 7, we have 2 edges:

$(7, 14), (7, 21)$
in another word, vertex 7 has degree 2

For vertex 8, we have 3 edges:

$(2, 8), (4, 8), (8, 16)$
in another word, vertex 8 has degree 3

For vertex 9, we have 2 edges:

$(3, 9), (9, 18)$
in another word, vertex 9 has degree 2

For vertex 10, we have 2 edges:

(5, 10), (10, 20)

in another word, vertex 10 has degree 2

For vertex 11, we have 1 edges:

(11, 22)

in another word, vertex 11 has degree 1

For vertex 12, we have 4 edges:

(2, 12), (3, 12), (4, 12), (6, 12)

in another word, vertex 12 has degree 4

For vertex 13, we have 0 edges:

in another word, vertex 13 has degree 0

For vertex 14, we have 2 edges:

(2, 14), (7, 14)

in another word, vertex 14 has degree 2

For vertex 15, we have 2 edges:

(3, 15), (5, 15)

in another word, vertex 15 has degree 2

For vertex 16, we have 3 edges:

(2, 16), (4, 16), (8, 16)

in another word, vertex 16 has degree 3

For vertex 17, we have 0 edges:

in another word, vertex 17 has degree 0

For vertex 18, we have 4 edges:

(2, 18), (3, 18), (6, 18), (9, 18)

in another word, vertex 18 has degree 4

For vertex 19, we have 0 edges:

in another word, vertex 19 has degree 0

For vertex 20, we have 4 edges:

(2, 20), (4, 20), (5, 20), (10, 20)

in another word, vertex 20 has degree 4

For vertex 21, we have 3 edges:

(3, 21), (7, 21)

in another word, vertex 21 has degree 3

For vertex 22, we have 2 edges:

(2, 22), (11, 22)

in another word, vertex 22 has degree 2

As these vertices are vertices of graph H , the degree sequence of H is:

10, 6, 5, 4, 4, 4, 4, 3, 3, 3, 3, 2, 2, 2, 2, 2, 1, 0, 0, 0 (in decreasing order)

- (b) A *path* from vertex u to vertex v in a graph G is an alternating sequence of vertices and edges

$$u = u_0, e_1, u_1, e_2, \dots, u_{n-1}, e_n, u_n = v$$

where $e_i = (u_{i-1}, u_i)$ and none of the vertices are repeated.

(A *walk* is the same as a path except that we allow repetition of edges and vertices.)

The *length* of a path is the number of edges in the path.

The *distance* from vertex u to vertex v in a graph G is the length of the shortest path from u to v .

A graph G is connected if, for all pairs of vertices u and v in G , there is a path from u to v .

Remove the vertices that have degree zero from H (along with their incident edges), to get the *subgraph* H' . Is H' connected? Support your claim fully.

What two vertices in H' are farthest apart? Support your claim fully.

From (a), we find that vertices 13, 17 and 19 have degree 0. Thus we remove them from H , H' consists of vertices 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 14, 15, 16, 18, 20, 21, and 22.

The adjacency matrix for H' is

	2	3	4	5	6	7	8	9	10	11	12	14	15	16	18	20	21	22
2	0	0	1	0	1	0	1	0	1	0	1	1	0	1	1	1	0	1
3	0	0	0	0	1	0	0	1	0	0	1	0	1	0	1	0	1	0
4	1	0	0	0	0	0	1	0	0	0	1	0	0	1	0	1	0	0
5	0	0	0	0	0	0	0	0	1	0	0	0	1	0	0	1	0	0
6	1	1	0	0	0	0	0	0	0	0	1	0	0	0	1	0	0	0
7	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0	1	0
8	1	0	1	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0
9	0	1	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0
10	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	1	0	0
11	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1
12	1	1	1	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0
14	1	0	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0
15	0	1	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0
16	1	0	1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0
18	1	1	0	0	1	0	0	1	0	0	0	0	0	0	0	0	0	0
20	1	0	1	1	0	0	0	0	1	0	0	0	0	0	0	0	0	0
21	0	1	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0
22	1	0	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	0

So as question 2, I wrote following MATLAB codes to approach this question:

```

-----
% @Author: Baihan Lin, @Date: Oct 2016

Hp = [0 0 1 0 1 0 1 0 1 0 1 1 0 1 1 1 0 1;
      0 0 0 0 1 0 0 1 0 0 1 0 1 0 1 0 1 0;
      1 0 0 0 0 0 1 0 0 0 1 0 0 1 0 1 0 0;
      0 0 0 0 0 0 0 0 1 0 0 0 1 0 0 1 0 0;
      1 1 0 0 0 0 0 0 0 0 1 0 0 0 1 0 0 0;
      0 0 0 0 0 0 0 0 0 0 0 1 0 0 0 0 1 0;
      1 0 1 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0;
      0 1 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0 0;
      0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 1 0 0;
      0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1;
      1 1 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0;
      1 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0;
      0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0;
      1 0 1 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0;
      1 1 0 0 1 0 0 1 0 0 0 0 0 0 0 0 0 0;
      1 0 1 1 0 0 0 0 1 0 0 0 0 0 0 0 0 0;
      0 1 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0;
      1 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0];

% find minimum steps for every vertex can connect to each other (complete).
nH = whenComplete(Hp, 1);

nH % minimum steps to reach a complete graph (if n = -1, not connected).
Hn = Hp^nH % the final complete adjacency graph in n steps.

-----
% to find out in how many steps can each vertex be connected to another one.
function n = whenComplete(A, n)
% A is the input adjacency matrix
% n is the starting steps (set to -1 if it is not connected at all)

if n ~= -1
    % only in connected graphs, we find minimum steps towards complete.
    if prod(sum(A)) == 0 % if one column of 0, not connected, set n to -1.
        n = -1;
    else
        if prod(prod(A^n)) == 0 % not complete graph yet within n step.
            n = whenComplete(A, n+1);
        end
    end
end
end
end
-----

```

Here is the output on MATLAB console:

nH =

6

Hn =

2281	1148	1515	586	1173	220	981	450	767	141	1192	419	274	981	967	1028	204	413
1127	768	694	282	664	140	416	298	330	63	615	193	157	416	575	432	138	187
1478	691	1065	375	807	119	711	287	574	82	837	332	201	711	673	740	150	327
384	205	261	136	181	38	152	65	144	22	177	62	45	152	145	180	24	62
1168	668	829	265	798	98	582	338	476	56	811	314	245	582	708	619	210	299
219	141	120	56	98	37	65	43	47	15	87	21	15	65	78	66	12	21
963	417	729	227	580	64	551	200	434	48	660	276	169	550	520	585	127	271
447	299	291	94	337	43	200	165	164	22	326	114	113	200	301	216	104	104
307	132	219	101	151	24	138	54	133	18	158	59	51	138	121	166	27	58
141	64	83	36	56	15	49	22	35	12	55	15	9	49	41	50	6	15
1172	616	858	261	809	86	660	326	528	54	897	357	266	660	757	714	220	342
419	195	349	94	316	21	278	115	232	15	359	165	109	278	298	301	92	156
241	148	180	54	223	14	147	108	137	8	241	94	102	147	215	179	84	85
963	417	729	227	580	64	550	200	434	48	660	276	169	551	520	585	127	271
962	578	695	211	709	78	522	302	427	41	759	297	236	522	673	564	204	282
824	357	646	214	537	48	512	188	430	36	632	269	188	512	499	580	132	263
202	138	153	34	209	12	126	104	109	6	219	91	88	126	203	141	86	78
413	189	344	94	301	21	273	105	227	15	344	156	100	273	283	295	79	153

If H' is connected, there must be a number k of walks to satisfy the path from any vertex i to any vertex j in H' . If so, it means from every vertex, it is possible to travel to another vertex in certain steps, which suggests the connectivity of the graph.

Therefore, I wrote a program to find the minimum steps it takes, to let every vertex in this graph to travel towards any other vertices (more details in comments).

I found that when $n = 6$, every vertex has at least one possible ways to travel to any other vertex in 6 steps. Since H^6 has no zero entries, we know there is a path (of length 6) between every pair of vertices in this graph H' .

Thus, this proves that the graph H' is connected.

For the next question, I wrote another MATLAB function to find the farthest vertices.

```
% find minimum distance between every two vertices.
Hd = minD(Hp, nH)
```

Here in part (b), the code also shares part (a)'s other MATLAB code as inputs of this function:

```
% to minimum distance between every two vertices (min path length).
```

```
function d = minD(H, n)
% H is the input adjacency matrix
% n is the minimum steps to make the graph complete.

fin = [H ~= 0]; % finished connecting pairs
d = 1*[H ~= 0]; % the vertex pairs with distance of 1.
for k = 2:n
    d = d + k.*(H^k ~= 0) - fin; % add
    fin = [H^k ~= 0]; % update the already connected pairs
end
```

Here is the output on MATLAB console:

```
Hn =
```

2	2	1	2	1	3	1	2	1	3	1	2	3	1	1	1	3	2
2	2	2	3	1	3	3	1	3	4	1	3	2	3	1	3	2	3
1	2	2	2	2	3	1	3	2	3	1	2	3	1	2	1	3	2
2	3	2	2	3	4	3	3	1	4	3	3	2	3	3	1	4	3
1	1	2	3	2	3	2	2	2	3	1	2	2	2	1	2	2	2
3	3	3	4	3	4	3	3	3	5	3	3	4	3	3	3	3	4
1	3	1	3	2	3	2	3	2	3	2	2	4	1	2	2	4	2
2	1	3	3	2	3	3	2	3	4	2	3	2	3	1	3	2	3
2	3	2	1	3	4	3	4	2	4	3	3	2	3	3	1	4	3
3	4	3	4	3	5	3	4	3	4	3	4	5	3	3	3	5	3
1	1	1	3	1	3	2	2	2	3	2	2	2	2	2	2	2	2
2	3	2	3	2	3	2	3	2	4	2	3	4	2	2	2	3	3
3	2	3	2	2	4	4	2	2	5	2	4	3	4	2	2	3	4
1	3	1	3	2	3	1	3	2	3	2	2	4	2	2	2	4	2
1	1	2	3	1	3	2	1	2	3	2	2	2	2	2	2	2	2
1	3	1	1	2	3	2	3	1	3	2	2	2	2	2	2	4	2
3	2	3	4	2	3	4	2	4	5	2	3	3	4	2	4	3	4
2	3	2	3	2	4	2	3	2	3	2	3	4	2	2	2	4	3

As shown in this matrix, which shows the minimum distance of path between each two vertices, we can see that the maximum distance is 5 (next page).

	2	3	4	5	6	7	8	9	10	11	12	14	15	16	18	20	21	22
2	2	2	1	2	1	3	1	2	1	3	1	2	3	1	1	1	3	2
3	2	2	2	3	1	3	3	1	3	4	1	3	2	3	1	3	2	3
4	1	2	2	2	2	3	1	3	2	3	1	2	3	1	2	1	3	2
5	2	3	2	2	3	4	3	3	1	4	3	3	2	3	3	1	4	3
6	1	1	2	3	2	3	2	2	2	3	1	2	2	2	1	2	2	2
7	3	3	3	4	3	4	3	3	3	5	3	3	4	3	3	3	3	4
8	1	3	1	3	2	3	2	3	2	3	2	2	4	1	2	2	4	2
9	2	1	3	3	2	3	3	2	3	4	2	3	2	3	1	3	2	3
10	2	3	2	1	3	4	3	4	2	4	3	3	2	3	3	1	4	3
11	3	4	3	4	3	5	3	4	3	4	3	4	5	3	3	3	5	3
12	1	1	1	3	1	3	2	2	2	3	2	2	2	2	2	2	2	2
14	2	3	2	3	2	3	2	3	2	4	2	3	4	2	2	2	3	3
15	3	2	3	2	2	4	4	2	2	5	2	4	3	4	2	2	3	4
16	1	3	1	3	2	3	1	3	2	3	2	2	4	2	2	2	4	2
18	1	1	2	3	1	3	2	1	2	3	2	2	2	2	2	2	2	2
20	1	3	1	1	2	3	2	3	1	3	2	2	2	2	2	2	4	2
21	3	2	3	4	2	3	4	2	4	5	2	3	3	4	2	4	3	4
22	2	3	2	3	2	4	2	3	2	3	2	3	4	2	2	2	4	3

And the farthest paths (or vertices pairs) are:

(7 \leftrightarrow 11)

(11 \leftrightarrow 15)

(11 \leftrightarrow 21)

They all needs at least 5-degree distance apart in H' .