



國立臺灣大學

National  
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University

# Nachos Project Assignment 3

## Memory Management / CPU part 2

EE 5173 Operating System  
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# Project 3 – part1

## Memory Management

# What will you learn from this project?

1. Memory Swap management
2. Swap replacement algorithm
3. Demand paging

# Why we need swap management

1. It expands the amount of memory a process may use.
2. A significant number of the pages referenced by a process early in its life may only be used for initialization and then never used again

## Disadvantages:

1. Performance decrease (Disk is slow.)

# Swap Management

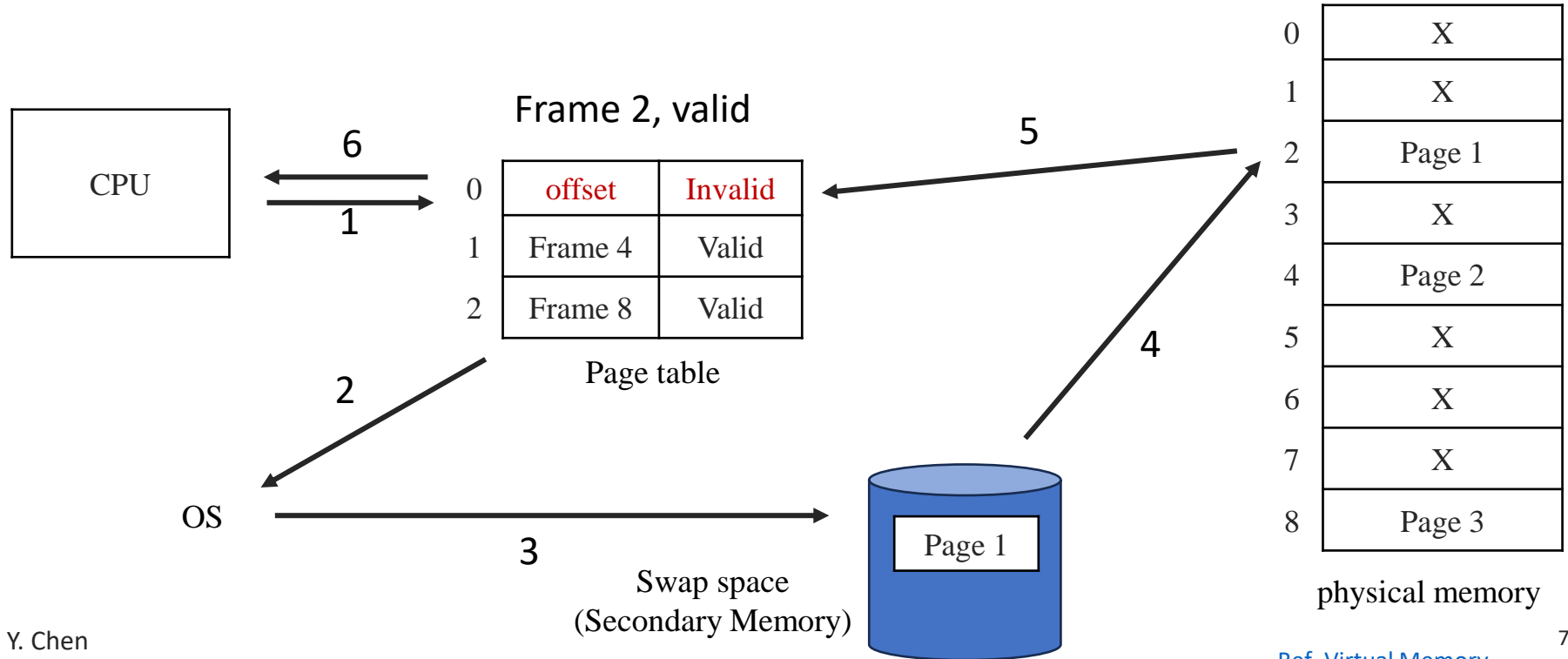
Goal:

- Successfully execute the following test cases, concurrently.
  - `/test/matmalt.c`
  - `/test/sort.c`
- Both cases use require lots of memory
  - Larger than the physical memory

# Demand Paging

- Only load pages into memory when necessary. (Actually, we loads some pages into memory to avoid lots of page fault).
- Strictly follow above definition which is called pure demand paging.
- You can decide how many preload pages.
- When page not in memory (present bit off), trap to OS via page fault.

# Steps in handling a page fault



# Page fault

- When valid bit set to fault, this will throw a PageFaultException and trap to Kernel

```
213     } else if (!pageTable[vpn].valid) {  
214         DEBUG(dbgAddr, "Invalid virtual page # " << virtAddr);  
215         return PageFaultException;  
216     }
```

/machine/kernel/translate.cc

- We can catch the exception in **ExceptionHandler** and handle the page fault.
- Add some codes in `/userprog/exception.cc` and create other files if needed.
- BadVAddrReg** will store the VA which is not in memory.



# Page fault

- You can create a Virtual Memory Manger class to handle page fault.
- Note: Load 1 page
- If memory is enough , get a free frame and put the page in.
- If memory is full, page-replacement algorithm to select victim frame to swap out.
  - Design your own page replacement algorithm to get the frame
  - Specify the method in your report
  - More algorithm more socre

# Swap Space – Create Disk

- swap = new **SynchDisk** in userkernel.cc
- Use the disk as swap space.
- Access the virtual memory in the disk by...
  - **kernel->swap->WriteSector**
  - **kernel->swap->ReadSector**
- See “**synchdisk.cc**” and other files in **/filesys** for details!
  - Read the header in those files first

# SynchDisk

- SynchDisk object resides above the raw disk (Nachos)
- It **blocks** the calling thread until after the corresponding I/O complete interrupt takes place. (synchronous)
- Nachos is multi-thread and disk cannot service more than one request at a time.
- It provides mutual exclusion, so that multiple threads can safely call the SynchDisk routines concurrently

# SynchDisk

SynchDisk provides the following operations

- `SynchDisk(char *name):`

Constructor takes the name of Unix file that holds the disk's contents.

- `ReadSector(int sectorNumber, char *data)`
- `WriteSector(int sectorNumber, char *data)`

Invoke the underlying `Disk::ReadSector/ WriteSector` and then wait for interrupt

- `Callback():`

Disk interrupt handler is called when an I/O operation completes

# Project Tasks

- Disk – swap space
- Record some necessary information to manage memory swapping
  - PageTable
  - Swap replacement information
  - Memory usage
  - ...
- Demand paging , page fault handler and page replacement algorithm (FIFO, LRU...)

# Some files that might be useful

- For the disk usage details, see:
  - `/filesys/synchdisk.h`
  - `/filesys/synchdisk.cc`
  - Other files in `/filesys(Optional)`
- For the swap space initialization:
  - `/userprog/userkernel.h`
  - `/userprog/userkernel.cc`
- For the page fault handling:
  - `/userprog/exception.h`
  - `/userprog/exception.cc`

# Some files that might be useful

- For the loading of pages
  - `userprog/addrspace.h`
  - `userprog/addrspace.cc`
- Header file (.h) and comments are helpful
- Based on your implementation, there might be different files that your need to see and modify.

# Project 3 – part2

## CPU scheduling 2



# CPU scheduling 2

- Please following the instruction in the assignment
- Assignment will be released on NTU cool.
- TA will explain next week.

# Report

## Part 1 (50%)

- What is your motivation, problem analysis and plan ?
- Explain the details of code snippet you added or modified.
- Experiment result and some discussion, observation.
- What problem you face and how to tackle it?

## Part 2 (50%)

\* If your code are more different than reference, more score

# Policy

- Please save as [Student ID]\_project3\_1.pdf, [Student ID]\_project3\_2.pdf
  - E.g. f10921a18\_project3\_1.pdf, f10921a18\_project3\_2.pdf
- Upload to NTU cool. **DDL: 2023/12/28**
- Penalty: decrease 5% per day
- **No plagiarism**

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# Reference

- [Nachos SynchDisk](#)
- [Linux Swap Management](#)



Thanks for your attention !

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