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# Part 1: Multi-Programming

```
test1 (source code and execution result)
執行結果為 9~6 依序輸出
#include "syscall.h"
main()

{
    int     n;
    for (n=9;n>5;n--)
        PrintInt(n);
}
```

```
r12631055@r12631055-VirtualBox:~/Nachos/nachos-4.0/code/userprog$ ./nachos -e ..
/test/test1
Total threads number is 1
Thread ../test/test1 is executing.
Print integer:9
Print integer:7
Print integer:6
return value:0
No threads ready or runnable, and no pending interrupts.
Assuming the program completed.
Machine halting!

Ticks: total 200, idle 66, system 40, user 94
Disk I/O: reads 0, writes 0
Console I/O: reads 0, writes 0
Paging: faults 0
Network I/O: packets received 0, sent 0
r12631055@r12631055-VirtualBox:~/Nachos/nachos-4.0/code/userprog$
```

test2 (source code and execution result)

```
r12631055@r12631055-VirtualBox:~/Nachos/nachos-4.0/code/userprog$ ./nachos -e ..
/test/test2
Total threads number is 1
Thread ../test/test2 is executing.
Print integer:20
Print integer:21
Print integer:22
Print integer:23
Print integer:24
Print integer:25
return value:0
No threads ready or runnable, and no pending interrupts.
Assuming the program completed.
Machine halting!

Ticks: total 200, idle 32, system 40, user 128
Disk I/O: reads 0, writes 0
Console I/O: reads 0, writes 0
Paging: faults 0
Network I/O: packets received 0, sent 0
r12631055@r12631055-VirtualBox:~/Nachos/nachos-4.0/code/userprog$
```

executing test1 and test2 simultaneously, the result seems weird. The output is not consistent with source code.

同時執行兩程式,結果並未按照原先方式輸出

```
r12631055@r12631055-VirtualBox:~/Nachos/nachos-4.0/code/userprog$ ./nachos -e ..
/test/test1 -e ../test/test2
Total threads number is 2
Thread ../test/test1 is executing.
Thread ../test/test2 is executing.
Print integer:9
Print integer:8
Print integer:7
Print integer:20
Print integer:21
Print integer:22
Print integer:23
Print integer:24
Print integer:6
Print integer:7
Print integer:8
Print integer:9
Print integer:10
Print integer:12
Print integer:13
Print integer:14
Print integer:15
Print integer:16
Print integer:16
Print integer:17
Print integer:18
Print integer:19
Print integer:20
Print integer:17
Print integer:18
```

可以觀察到執行到interger:6時,程式開始出錯,開始遞增數列,粗估推測會導致程式執行結果而與預期不同的原因可能是:Nachos系統的預設並沒有對多個程式做記憶體管理,會將全部的實體記憶體分頁都分配出去,因此當多個執行緒在執行時,兩個程式會被分到同一個page,導致執行時出現錯誤。雖然有開啟虛擬記憶體,但是基本上沒有作為,所以當多程式同時執行時就會重疊到其他程式正在使用的page,然後發生錯誤,為此我們要修改addrspace,使程式的虛擬記憶體映射到沒有人使用的實體記憶體,而不是互相糾纏。

要讓程式的虛擬記憶體能map到實體記憶體上,首先第一步先修改addrspace.h 在addrspace.h檔中,多增加一行變數

static bool usedPhysicalPage[NumPhysPages],用來記錄page的使用情形,確認當下有哪些page是可以使用的。

```
class AddrSpace {
 public:
   AddrSpace();
                                        // Create an address space.
    ~AddrSpace();
                                        // De-allocate an address space
    void Execute(char *fileName);
                                        // Run the the program
                                        // stored in the file "executable"
   void SaveState();
                                        // Save/restore address space-specific
    void RestoreState();
                                                on a context switch
   static bool usedPhysicalPage[NumPhysPages];
                                        // Assume linear page table translation
    TranslationEntry *pageTable;
                                        // for now!
   unsigned int numPages;
                                        // Number of pages in the virtual
                                        // address space
   bool Load(char *fileName);
                                        // Load the program into memory
                                        // return false if not found
    void InitRegisters();
                                        // Initialize user-level CPU registers,
                                        // before jumping to user code
```

然後在addrspace.cc檔中初始化宣告usedPhysicalPage:

在addrspace.cc中的AddrSpace::AddrSapce()函式,將函式中的所有程式做註解,因為進行多執行緒時,只需足夠大小的記憶體即可,無須分配實際記憶體的大小,之後在一併把映射記憶體這件事挪到 AddrSpace::Load 中

在解構子AddrSpace::~AddrSapce()函式中,用for迴圈把所有有用到的Page都改回 false,釋放資源以供之後的程式使用。

在AddrSpace::Load()函式中,把原本於AddrSpace::AddrSapce()註解掉的程式移至此並稍作修改。透過while迴圈查看每個physical page是否已被使用,若已被使用(值為true)則看下一頁,直到找到值為false的page,將其設為true並將資料存至該頁。

```
numPages = divRoundUp(size, PageSize);
//
          cout << "number of pages of " << fileName<< " is "<<numPages<<endl;</pre>
     pageTable = new TranslationEntry[NumPhysPages];
     for (unsigned int i = 0; i < NumPhysPages; i++) {
   pageTable[i].virtualPage = i; // for now, virt page # = phys page #
   while(j<NumPhysPages && AddrSpace::usedPhysicalPage[j]==true){</pre>
          AddrSpace::usedPhysicalPage[j]=TRUE;
          PageTable[i].physicalPage = j;
          //pageTable[i].physicalPage = i;
          pageTable[i].physicalPage = 0;
pageTable[i].valid = TRUE;
II
          pageTable[i].valid = FALSE;
//
          pageTable[i].use = FALSE;
          pageTable[i].dirty = FALSE;
          pageTable[i].readOnly = FALSE;
     size = numPages * PageSize;
```

最後是在AddrSpace::Load()函式中,因為要找出map後的位置,因此更改程式碼中兩個executable->ReadAt的部分,使virtualAddr先除以PageSize,以此得知使用的是第幾個page,再透過pageTable找到所對應實體頁的頁數,並乘上每個page的大小得到該頁的physical memory,接著再加上virtualAddr除以PageSize的餘數值,以得到page內的offset,如此即得到相對應的physical address。

```
// then, copy in the code and data segments into memory
         if (noffH.code.size > 0) {
        DEBUG(dbgAddr, "Initializing code segment.");
DEBUG(dbgAddr, noffH.code.virtualAddr << ", "
                                                             << noffH.code.size);
                  executable->ReadAt(
                  &(kernel->machine->mainMemory[pageTable[noffH.code.virtualAddr/
| PageSize].physicalPage*PageSize+(noffH.code.virtualAddr%PageSize
noffH.code.size, noffH.code.inFileAddr);
         if (noffH.initData.size > 0) {
        DEBUG(dbgAddr, "Initializing data segment.");
DEBUG(dbgAddr, noffH.initData.virtualAddr <</pre>
noffH.initData.size);
        executable->ReadAt(
                  &(kernel->machine->mainMemory[pageTable[noffH.code.virtualAddr/
PageSize].physicalPage*PageSize+(noffH.code.virtualAddr%PageSize)]),
noffH.code.size, noffH.code.inFileAddr);
    delete executable;
                                              // close file
    return TRUE;
                                              // success
```

修改之後,再同時執行test1與test2,即可發現程式已能正常執行。

```
Total threads number is 2
Thread ../test/test1 is executing.
Thread ../test/test2 is executing.
Print integer:9
Print integer:8
Print integer:7
Print integer:20
Print integer:21
Print integer:22
Print integer:23
Print integer:24
Print integer:6
return value:0
Print integer:25
return value:0
No threads ready or runnable, and no pending interrupts.
Assuming the program completed.
Machine halting!
Ticks: total 300, idle 8, system 70, user 222
Disk I/O: reads 0, writes 0
Console I/O: reads 0, writes 0
Paging: faults 0
Network I/O: packets received 0, sent 0
```

心得:光是在建置環境,研究文章跟trace code就花了不少時間。

## Part2: System call tracing

#### 使用gdb來進行System call tracing

```
(gdb) info b
Num
       Type
                      Disp Enb Address
                                         What
       breakpoint
                             0x0804d3a5 in Interrupt::Halt()
                      keep y
                                         at ../machine/interrupt.cc:236
       breakpoint already hit 1 time
(gdb) list 236
       //--
231
               Shut down Nachos cleanly, printing out performance statistics.
232
233
234
       void
       Interrupt::Halt()
235
236
           cout << "Machine halting!\n\n";</pre>
237
238
239
       }
240
```

觀察內部的system code

Show the current state of call stack.

以下分別來討論SC\_Halt各個功能:

### SC Halt

- Machine/mipssims.cc
- 模擬 MIPS P2/3000 processor

```
Machine:Run()

// 模擬執行 program in thread

// 切換到 UserMode

Machine::OneInstruction()

// 定義執行一個 instruction 所需的參數與步驟

// 例如:一個 byte 大小, rs, rt, rd 等 register

// 判斷 opCode 並執行對應的 instruction
```

- Machine/machine.cc
- 判斷 host 的 位元組順序 big-endian 或 little-endian

```
Machine::RaiseException()
// 當 interrupt 或是 exception 發生時,紀錄 program 當前位置
// 設定 User Mode -> Kernel Mode
// 呼叫 Exception Handler
// 設定 Kernel Mode -> User Mode
```

• userprog/exception.cc

```
ExceptionHandler()
// entry point
// 定義 exception 發生時所對應的指令
// syscall (interrupt) 或 exception 發生時會呼叫 (Halt 為其中一種 syscall)
// 後續呼叫 SysHalt()
```

• machine/interrupt.cc

```
Interrupt::Halt()
// Show halt 的訊息
// delete kernel <- 這個動作?
```

### Reference

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