

<https://www.chegg.com/homework-help/questions-and-answers/given-array-jobs-every-job-deadline-associated-profit-job-finished-deadline-also-given-eve-q67106293>

Given An Array Of Jobs Where Every Job Has A Deadline And Associated Profit If The Job Is Finished Before The Deadline. It Is Also Given That Every Job Takes A Single Unit Of Time, So The Minimum Possible

Given an array of jobs where every job has a deadline and associated profit if the job is finished before the deadline. It is also given that every job takes a single unit of time, so the minimum possible deadline for any job is 1. The job sequencing problem is to maximize total profit if only one job can be scheduled at a time. Below is an example output for the optimal solution to this problem.

Input: Five Jobs with the following deadlines and profits

<i>JobID</i>	<i>Deadline</i>	<i>Profit</i>
<i>a</i>	<i>2</i>	<i>100</i>
<i>b</i>	<i>1</i>	<i>19</i>
<i>c</i>	<i>2</i>	<i>27</i>
<i>d</i>	<i>1</i>	<i>25</i>
<i>e</i>	<i>3</i>	<i>15</i>

Output: Following is the maximum profit sequence of jobs c, a, e

- (a) Describe, in words, a greedy algorithm for the job sequencing problem.
(b) Prove that your greedy choice is optimal (e.g. proof by contradiction).

Greedy algorithm for the job sequencing problem:-

Step-01:

Sort all the given jobs in decreasing order of their profit.

Step-02:

Check the value of maximum deadline.

Draw a Gantt chart where maximum time on Gantt chart is the value of maximum deadline.

Step-03:

Pick up the jobs one by one.

Put the job on Gantt chart as far as possible from 0 ensuring that the job gets completed before its deadline.

2.) We can prove it by taking

Two components:

1.) Optimal substructure

2.) Greedy Choice Property: There exists an optimal solution that is consistent with the greedy choice made in the first step of the algorithm.

Greedy Choice Property

1. Let S_k be nonempty subproblem containing the set of activities that finish after activity a_k .

2. Let a_m be an activity in S_k with the earliest finish time.

3. Then a_m is included in some maximum-size subset of mutually compatible activities of S_k .

Proof

- Let A_k be a maximum-size subset of mutually compatible activities in S_k ,
- let a_j be the activity in A_k with the earliest finish time.
- If $a_j = a_m$, we are done, since we have shown that a_m is in some maximum-size subset of mutually compatible activities of S_k .
- If $a_j \neq a_m$, let the set $A_0 k = A_k - \{a_j\} \cup \{a_m\}$
- The activities in $A_0 k$ are disjoint, because – the activities in A_k are disjoint, – a_j is the first activity in A_k to finish, – $f_m \leq f_j$.
- Since $|A_0 k| = |A_k|$, we conclude that $A_0 k$ is a maximum-size subset of mutually compatible activities of S_k , and it includes a_m .