



# Akdeniz University

## Computer Engineering Department

CSE206 Computer Organization  
Week03: Computer Systems

Assoc.Prof.Dr. Taner Danişman  
[tdanisman@akdeniz.edu.tr](mailto:tdanisman@akdeniz.edu.tr)

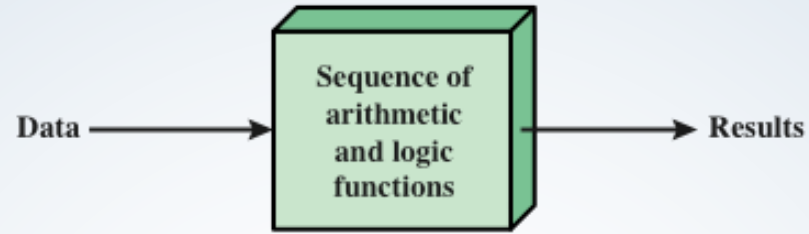
# Course program (Textbook: Stalling 10th Edt.)

Week 1	28-Feb-2023	Introduction	Ch1
Week 2	7-Mar-2023	Computer Evolution	Ch2
Week 3	14-Mar-2023	Computer Systems	Ch3
Week 4	21-Mar-2023	Cache Memory, Direct Cache Mapping	Ch4
Week 5	28-Mar-2023	Associative and Set Associative Mapping	Ch4
Week 6	4-Apr-2023	Internal Memory, External Memory, I/O	Ch5-Ch6-Ch7
Week 7	11-Apr-2023	Number Systems, Computer Arithmetic	Ch9-Ch10
Week 8	18-Apr-2023	Midterm (Expected date, may change)	Ch1-...-Ch10
Week 9	25-Apr-2023	Digital Logic	Ch11
Week 10	2-May-2023	Instruction Sets	Ch12
Week 11	9-May-2023	Addressing Modes	Ch13
Week 12	16-May-2023	Processor Structure and Function	Ch14
Week 13	23-May-2023	RISC, Instruction Level Parallelism	Ch15-Ch16
Week 14	30-May-2023	Assembly Language (TextBook : Assembly Language for x86 Processors)	Kip Irvine
Week 15	6-Jun-2023	Assembly Language (TextBook : Assembly Language for x86 Processors)	Kip Irvine

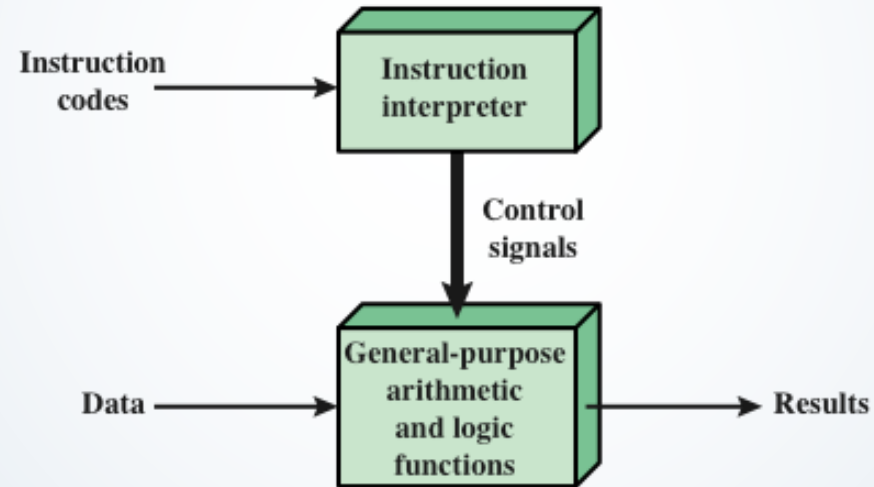
# Computer Components

- ▶ Contemporary computer designs are based on concepts developed by John von Neumann at the Institute for Advanced Studies, Princeton
- ▶ **Referred to as the *von Neumann architecture* and is based on three key concepts:**
  - ▶ Data and instructions are stored in a single read-write memory
  - ▶ The contents of this memory are addressable by location, without regard to the type of data contained there
  - ▶ Execution occurs in a sequential fashion (unless explicitly modified) from one instruction to the next
- ▶ ***Hardwired program***
  - ▶ The result of the process of connecting the various components in the desired configuration

# Hardware and Software Approaches



(a) Programming in hardware



(b) Programming in software

Figure 3.1 Hardware and Software Approaches

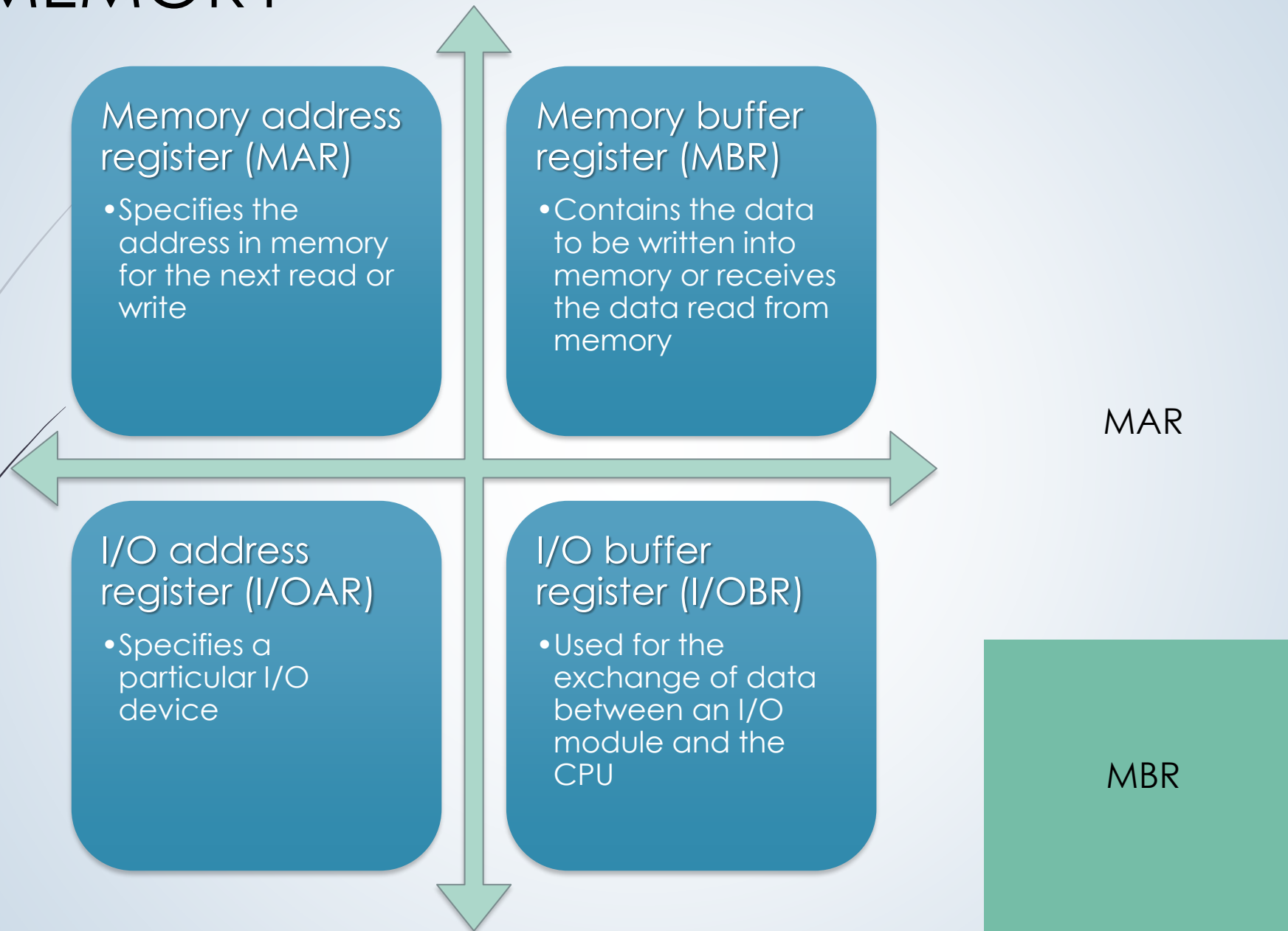
# I/O Components

## Software

- A sequence of codes or instructions
- Part of the hardware interprets each instruction and generates control signals
- Provide a new sequence of codes for each new program instead of rewiring the hardware

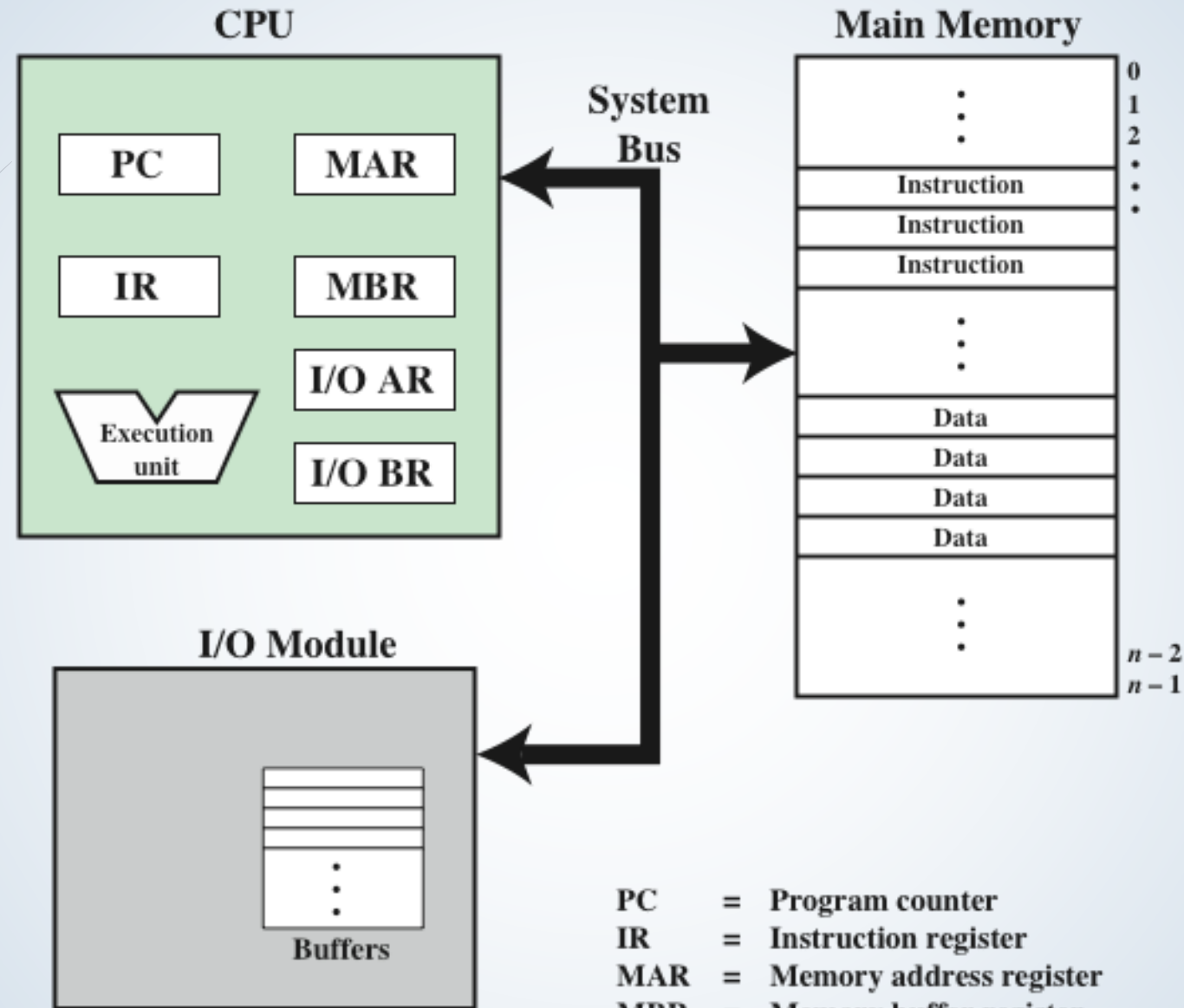
## Major components:

- CPU
  - Instruction interpreter
  - Module of general-purpose arithmetic and logic functions
- I/O Components
  - Input module
    - Contains basic components for accepting data and instructions and converting them into an internal form of signals usable by the system
  - Output module
    - Means of reporting results



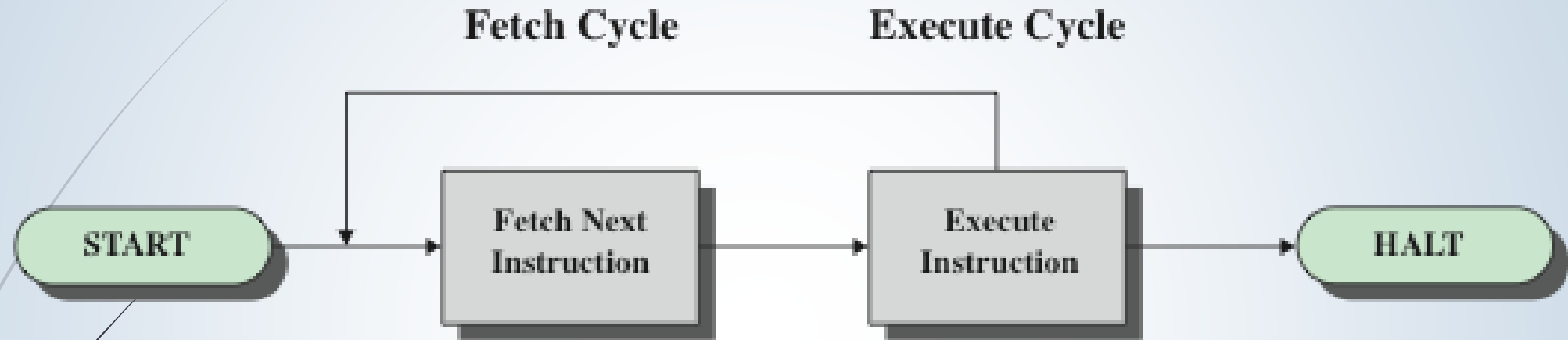


# Computer Components: Top Level View



PC = Program counter  
 IR = Instruction register  
 MAR = Memory address register  
 MBR = Memory buffer register  
 I/O AR = Input/output address register  
 I/O BR = Input/output buffer register

# Basic Instruction Cycle



**Figure 3.3 Basic Instruction Cycle**



# Fetch Cycle

- At the beginning of each instruction cycle the processor fetches an instruction from memory
- The **program counter** (PC) holds the address of the instruction to be fetched next
- The processor increments the PC after each instruction fetch so that it will fetch the next instruction in sequence
- The fetched instruction is loaded into the **instruction register** (IR)
- The processor interprets the instruction and performs the required action



# Action Categories

## Processor-memory

Data transferred from processor to memory or from memory to processor

## Processor-I/O

Data transferred to or from a peripheral device by transferring between the processor and an I/O module

## Data processing

The processor may perform some arithmetic or logic operation on data

## Control

An instruction may specify that the sequence of execution be altered

# Characteristics of a Hypothetical Machine



(a) Instruction format



(b) Integer format

Program Counter (PC) = Address of instruction  
Instruction Register (IR) = Instruction being executed  
Accumulator (AC) = Temporary storage

(c) Internal CPU registers

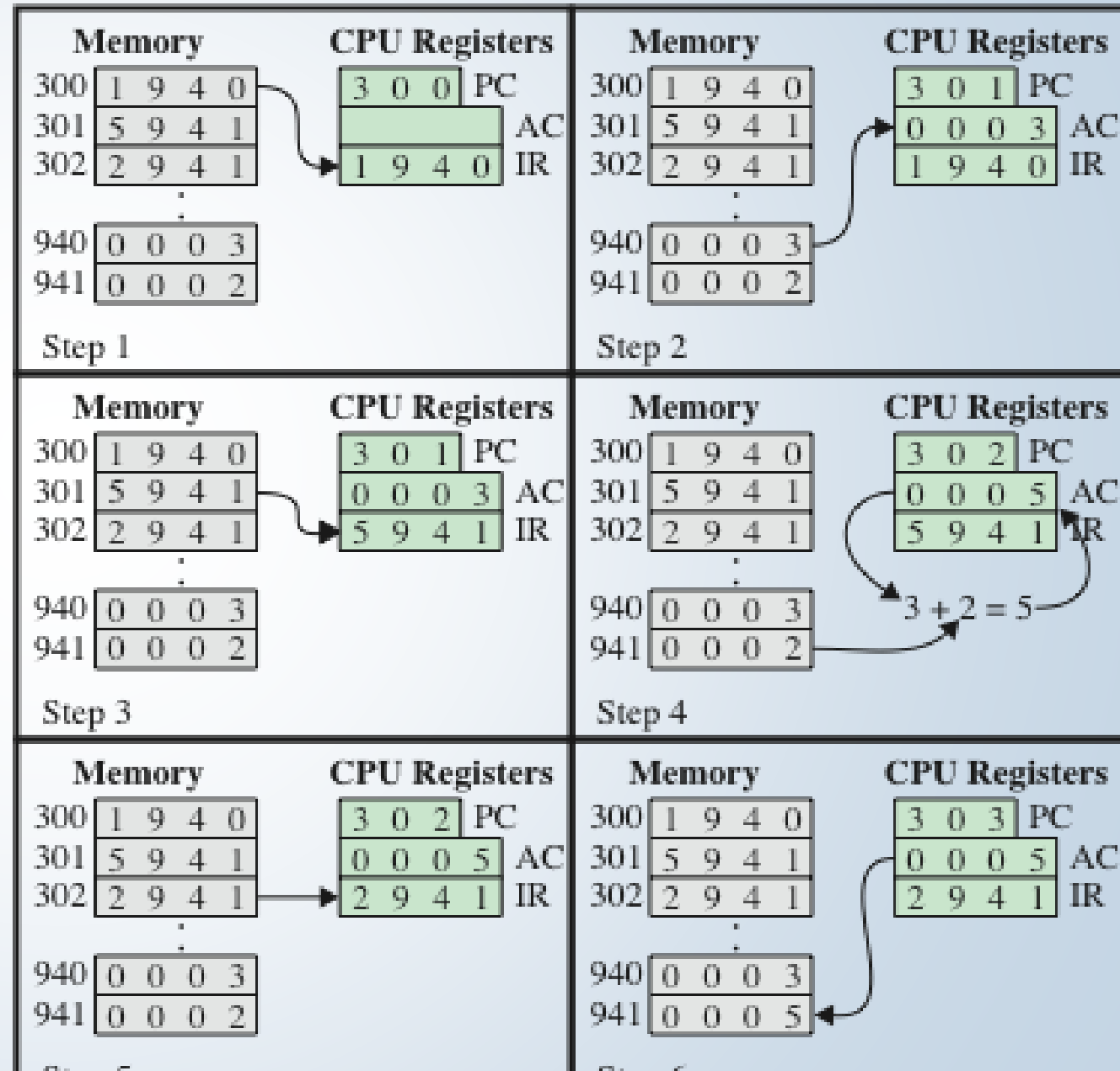
0001 = Load AC from Memory  
0010 = Store AC to Memory  
0101 = Add to AC from Memory

(d) Partial list of opcodes

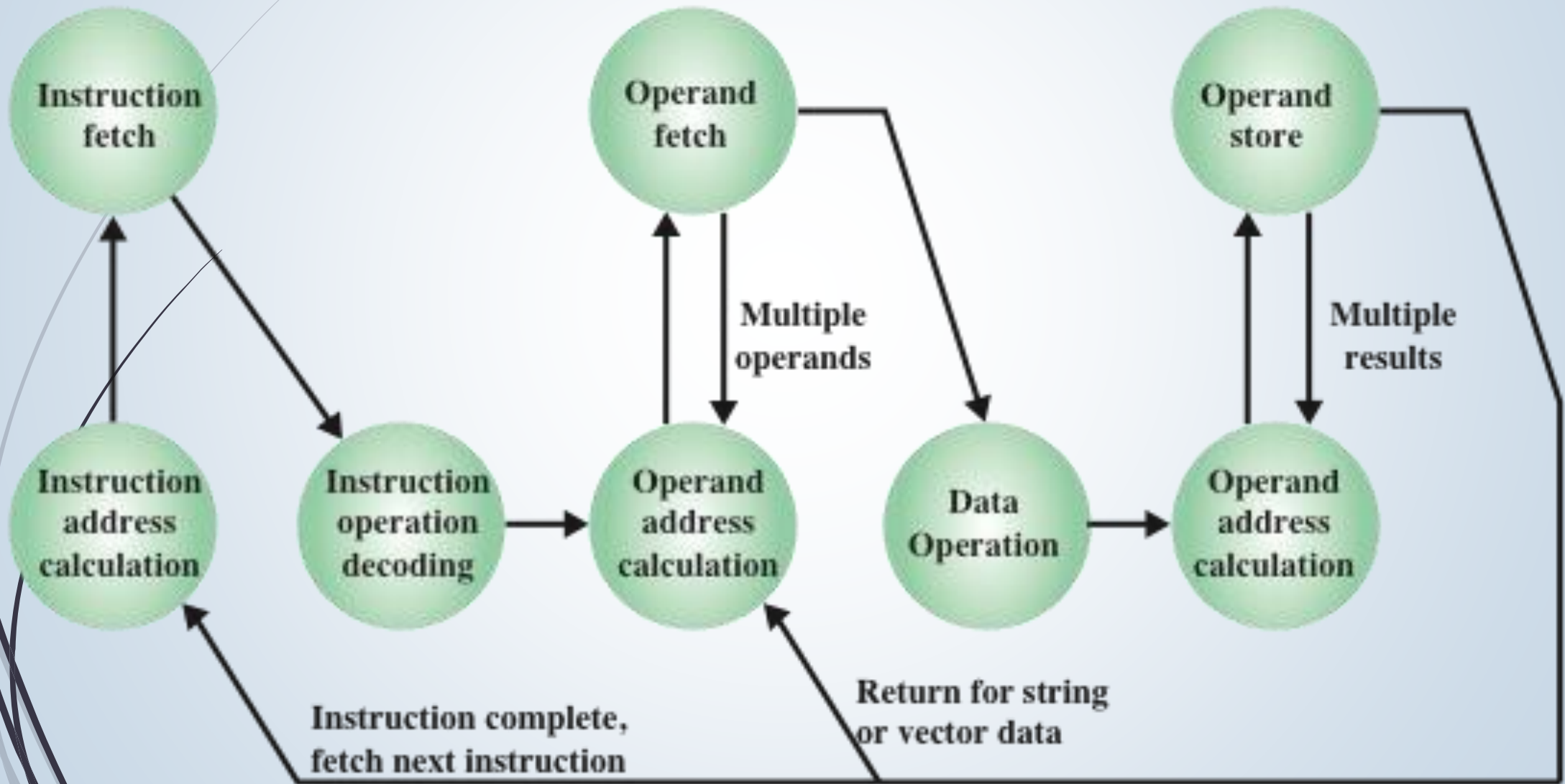
# Example of Program Execution

**Figure 3.5 Example of Program Execution**  
(contents of memory and registers in hexadecimal)

0001 = Load AC from Memory  
0010 = Store AC to Memory  
0101 = Add to AC from Memory



# Instruction Cycle State Diagram

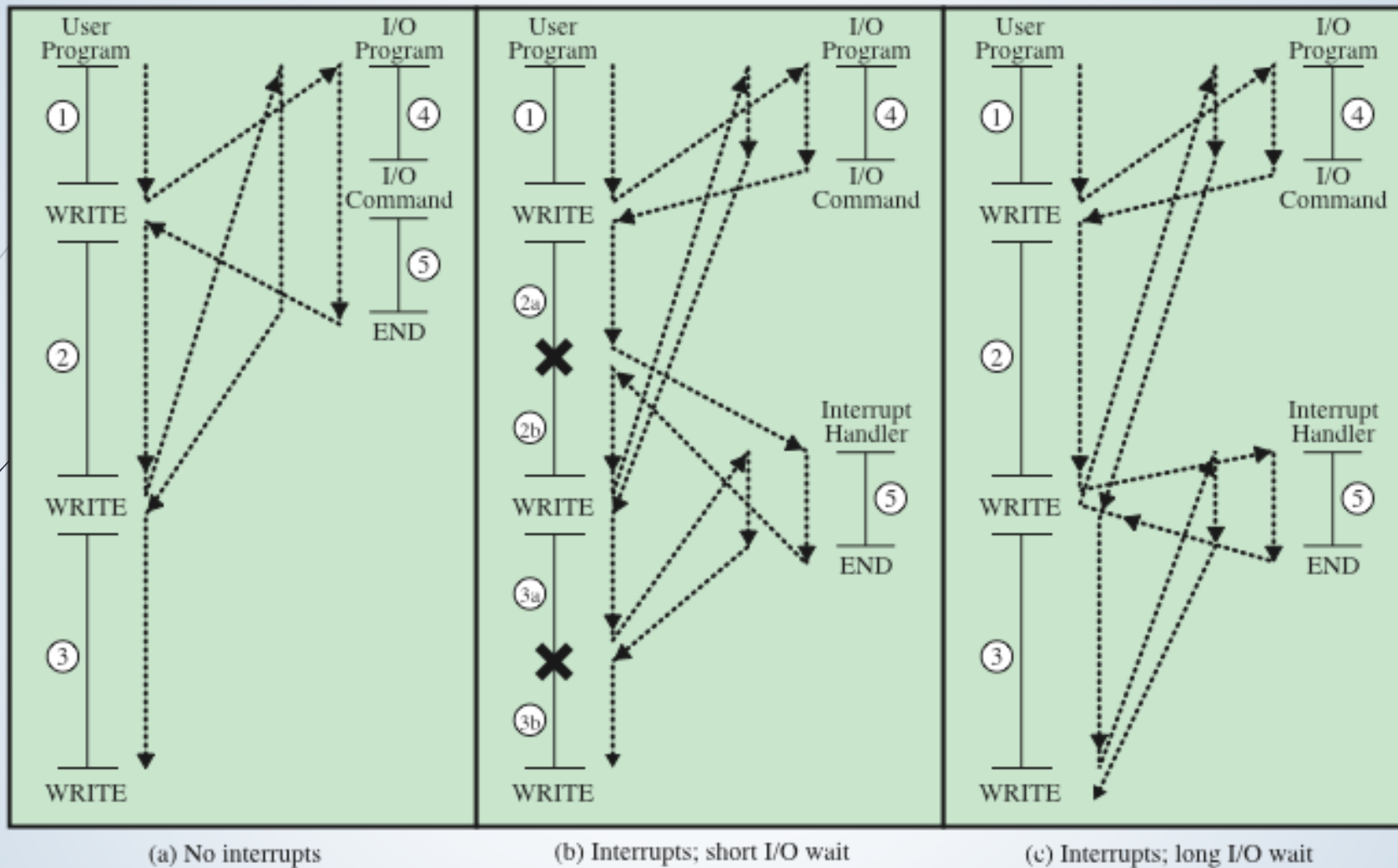


# Classes of Interrupts

<b>Program</b>	Generated by some condition that occurs as a result of an instruction execution, such as arithmetic overflow, division by zero, attempt to execute an illegal machine instruction, or reference outside a user's allowed memory space.
<b>Timer</b>	Generated by a timer within the processor. This allows the operating system to perform certain functions on a regular basis.
<b>I/O</b>	Generated by an I/O controller, to signal normal completion of an operation, request service from the processor, or to signal a variety of error conditions.
<b>Hardware failure</b>	Generated by a failure such as power failure or memory parity error.

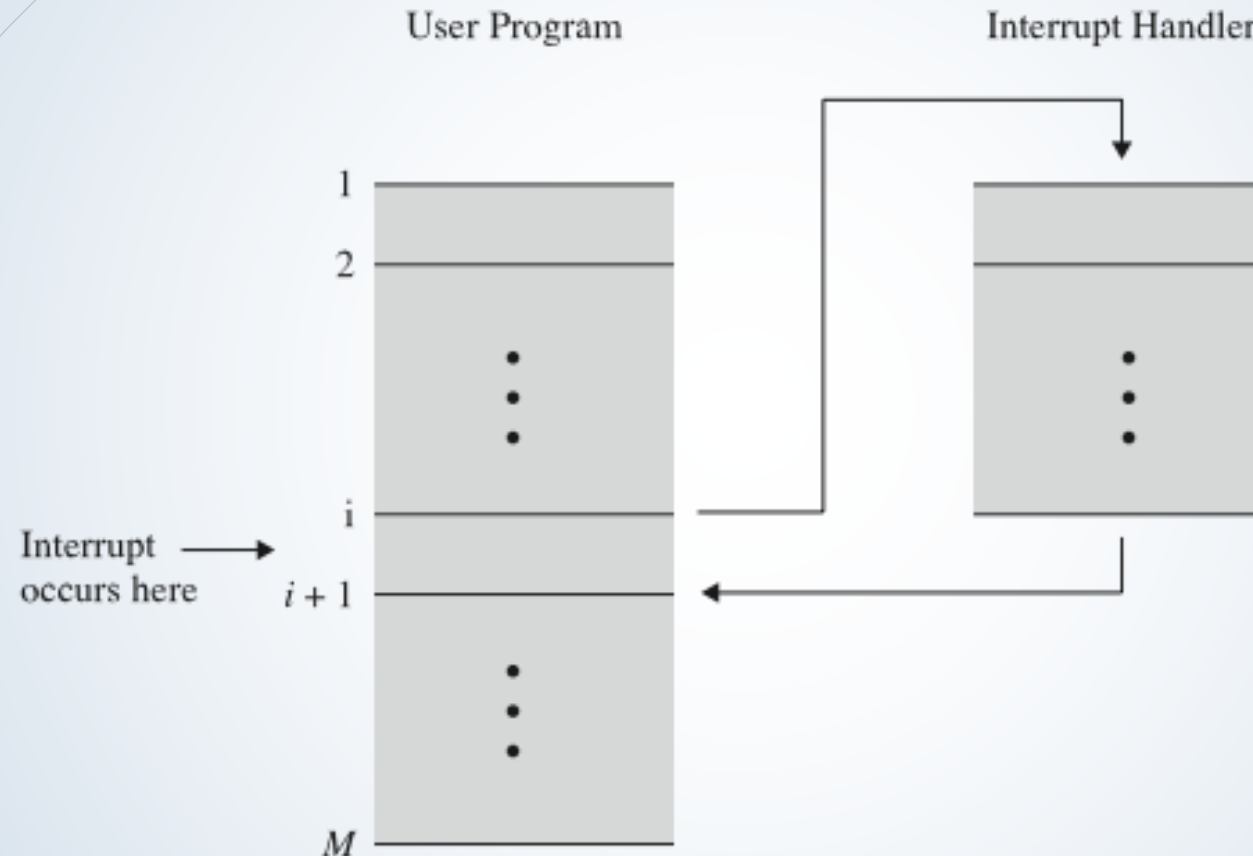


# Program Flow Control



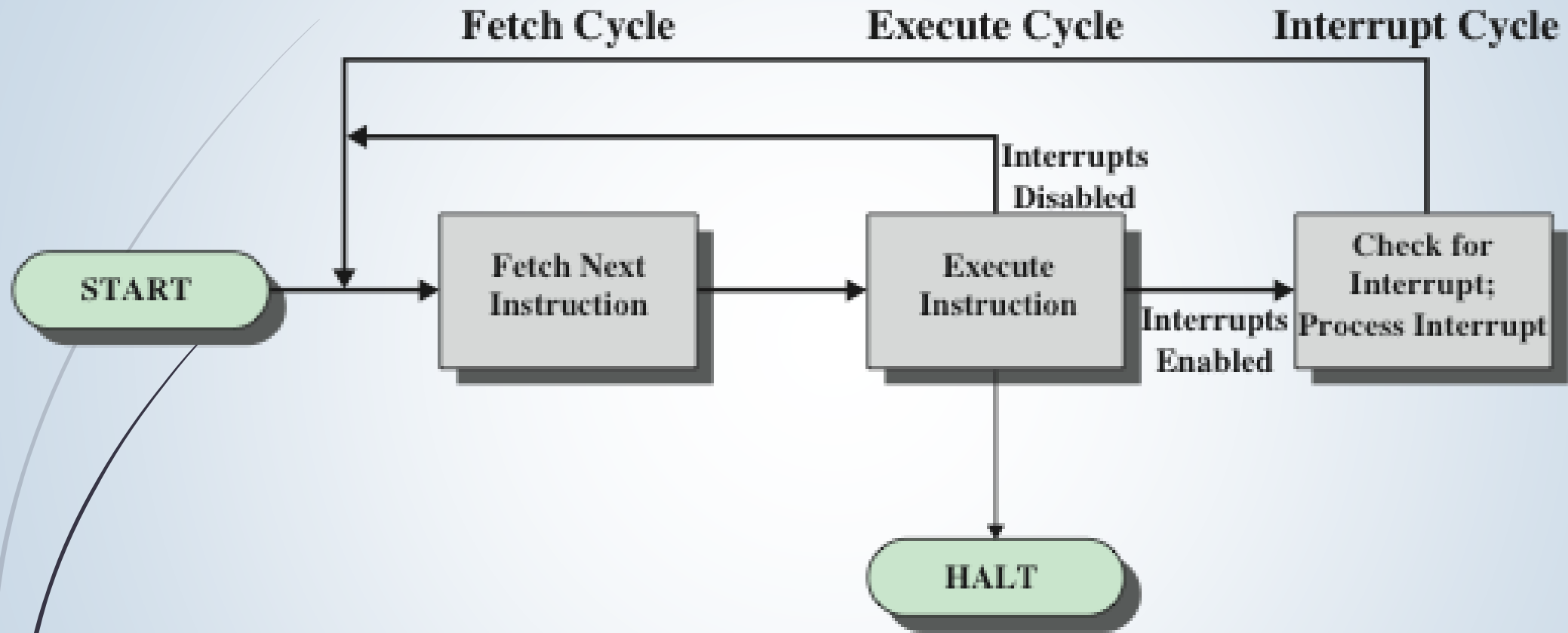
**Figure 3.7 Program Flow of Control Without and With Interrupts**

# Transfer of Control via Interrupts



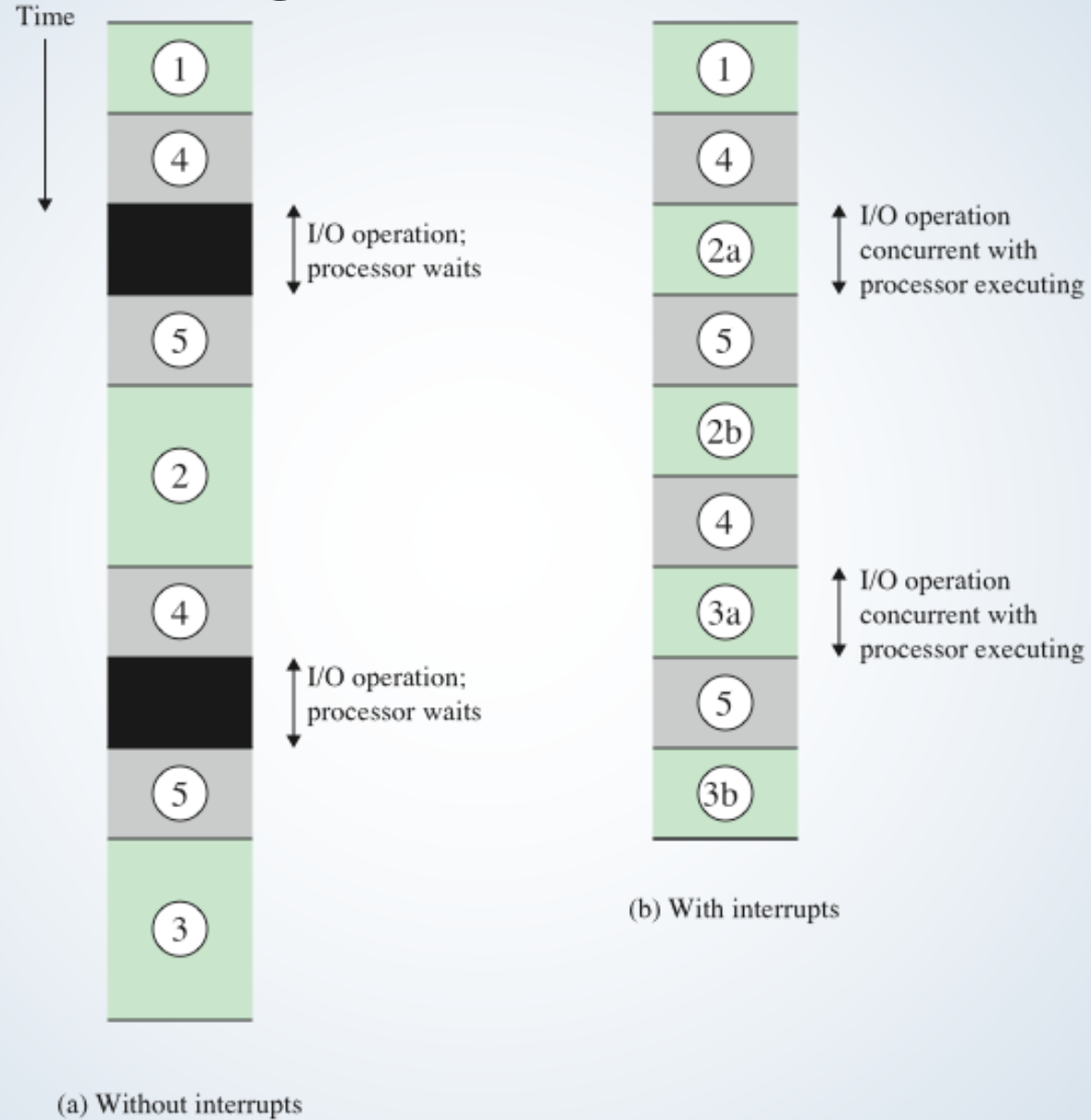
**Figure 3.8** Transfer of Control via Interrupts

# Instruction Cycle With Interrupts

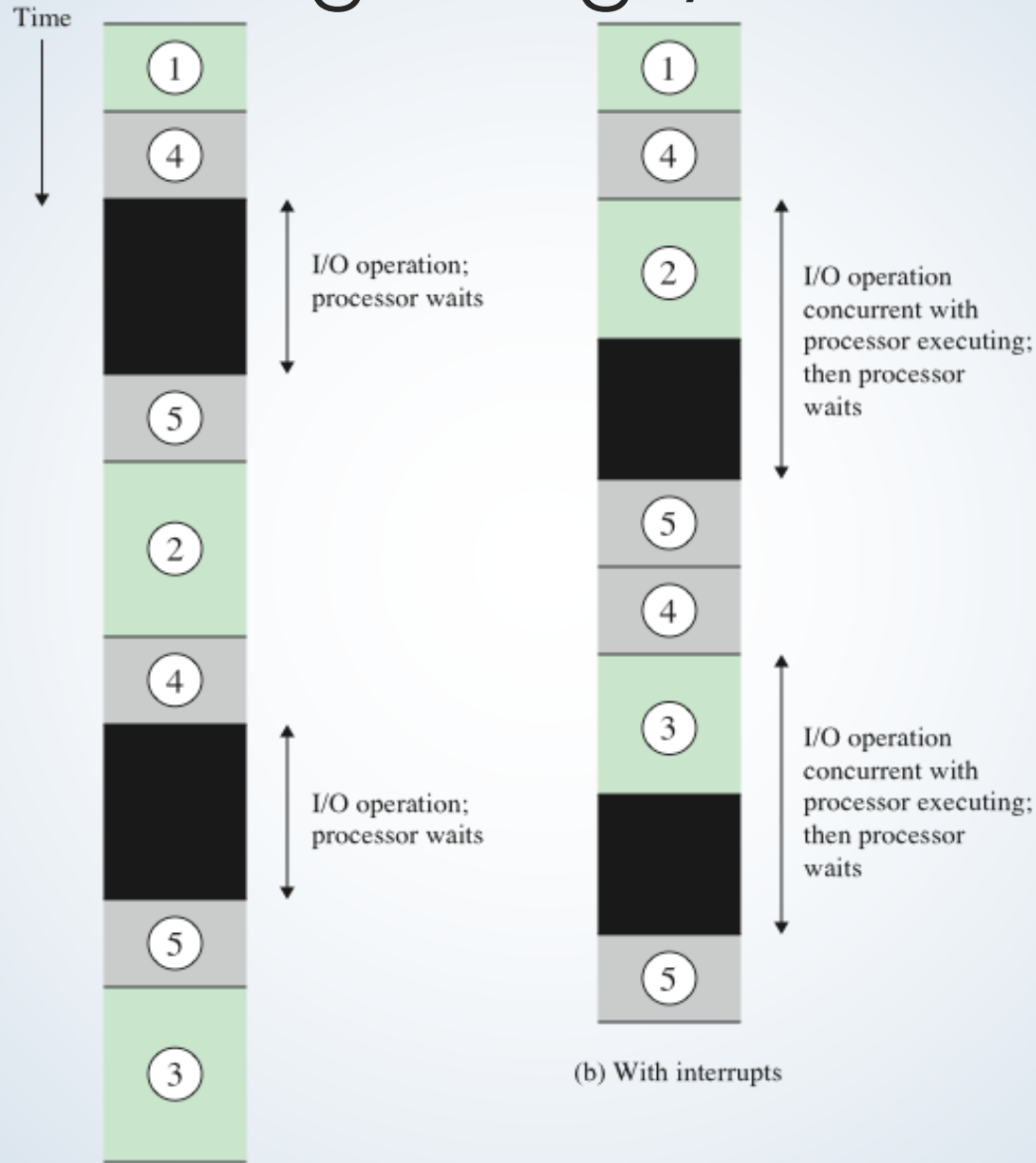


**Figure 3.9 Instruction Cycle with Interrupts**

# Program Timing: Short I/O Wait



# Program Timing: Long I/O Wait



(a) Without interrupts

(b) With interrupts

# Instruction Cycle State Diagram With Interrupts

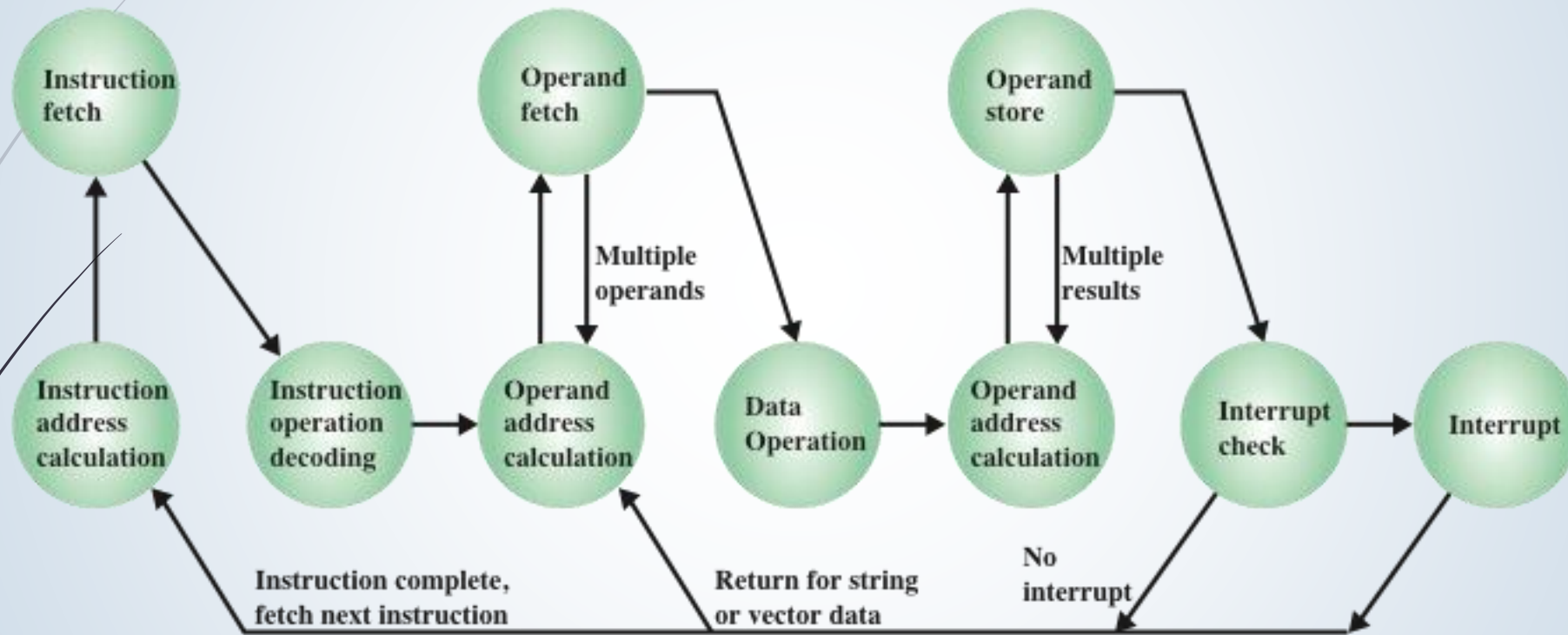
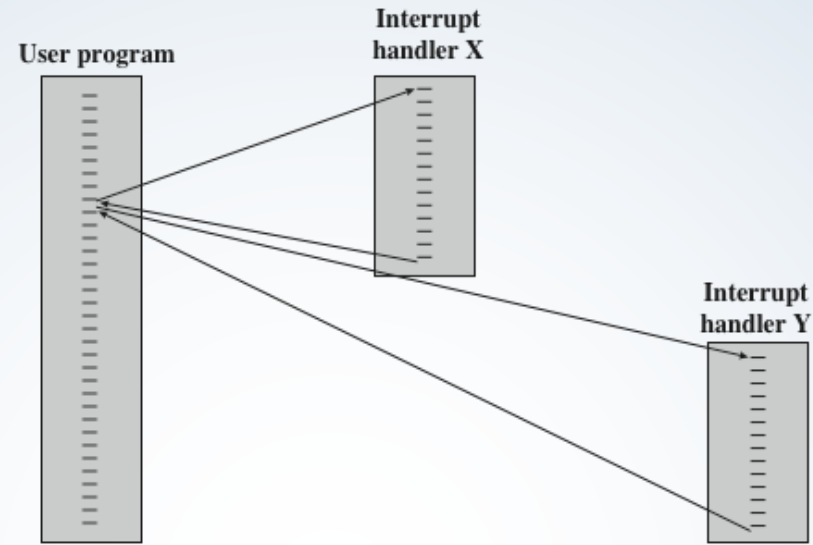


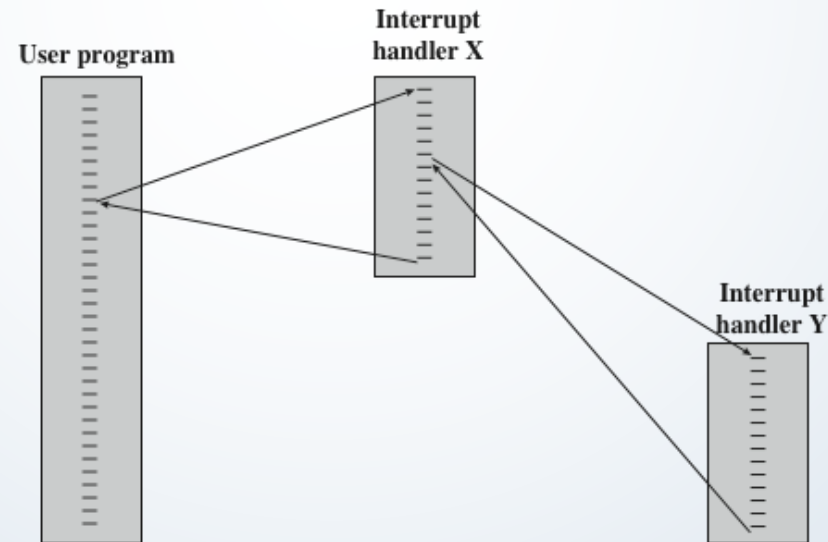
Figure 3.12 Instruction Cycle State Diagram, With Interrupts



# Transfer of Control with Multiple Interrupts



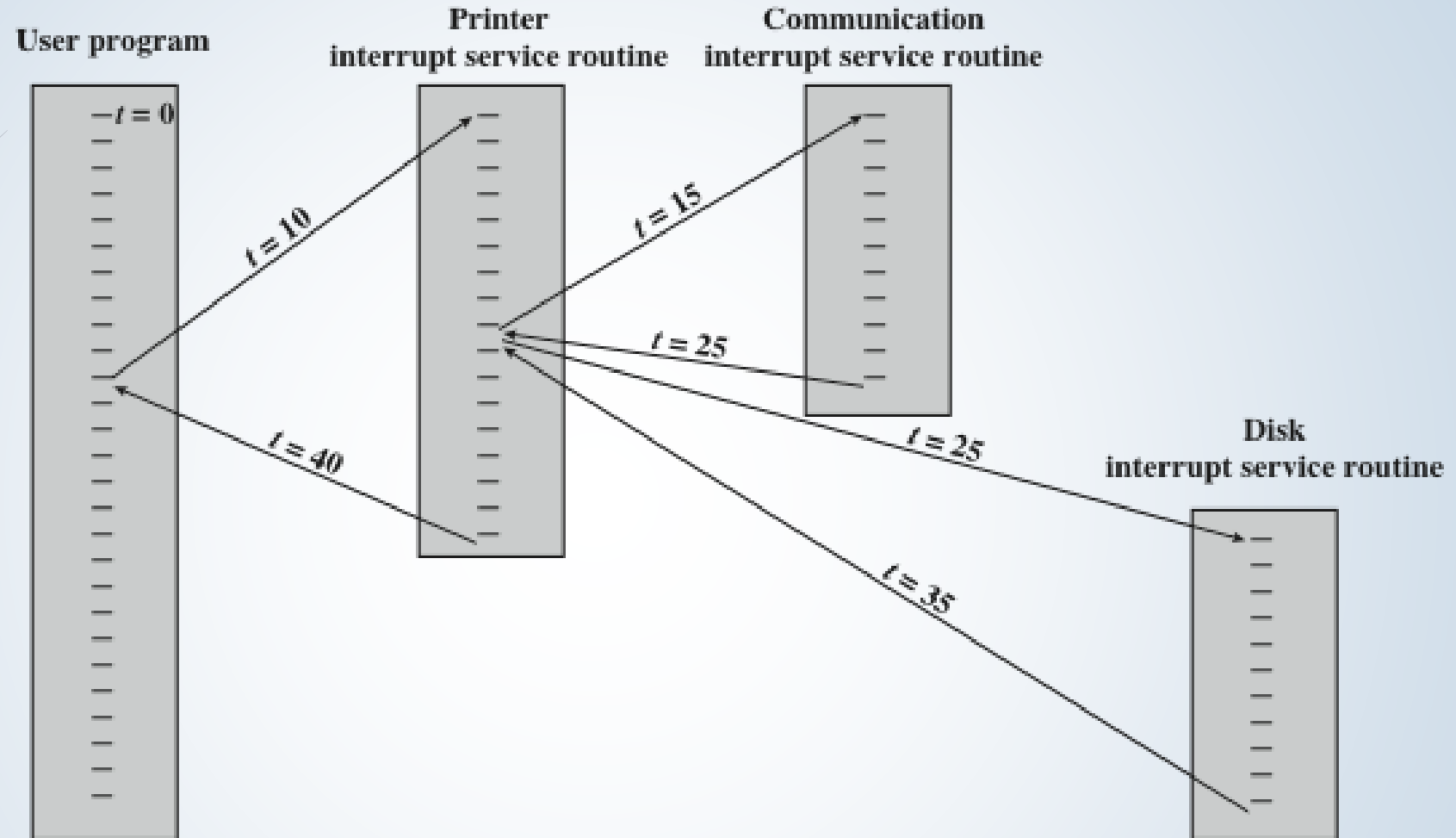
(a) Sequential interrupt processing



(b) Nested interrupt processing

Figure 3.13 Transfer of Control with Multiple Interrupts

# Time Sequence of Multiple Interrupts



**Priorities: Communication > Disk > Printer**

**Interrupts:  $t=10$  Printer,  $t=15$  Communication,  $t=20$  Disk**

**Figure 3.14 Example Time Sequence of Multiple Interrupts**

# I/O Function

- I/O module can exchange data directly with the processor
- Processor can read data from or write data to an I/O module
  - Processor identifies a specific device that is controlled by a particular I/O module
  - I/O instructions rather than memory referencing instructions
- In some cases it is desirable to allow I/O exchanges to occur directly with memory
  - The processor grants to an I/O module the authority to read from or write to memory so that the I/O memory transfer can occur without tying up the processor
  - The I/O module issues read or write commands to memory relieving the processor of responsibility for the exchange
  - This operation is known as direct memory access (DMA)

# Computer Modules

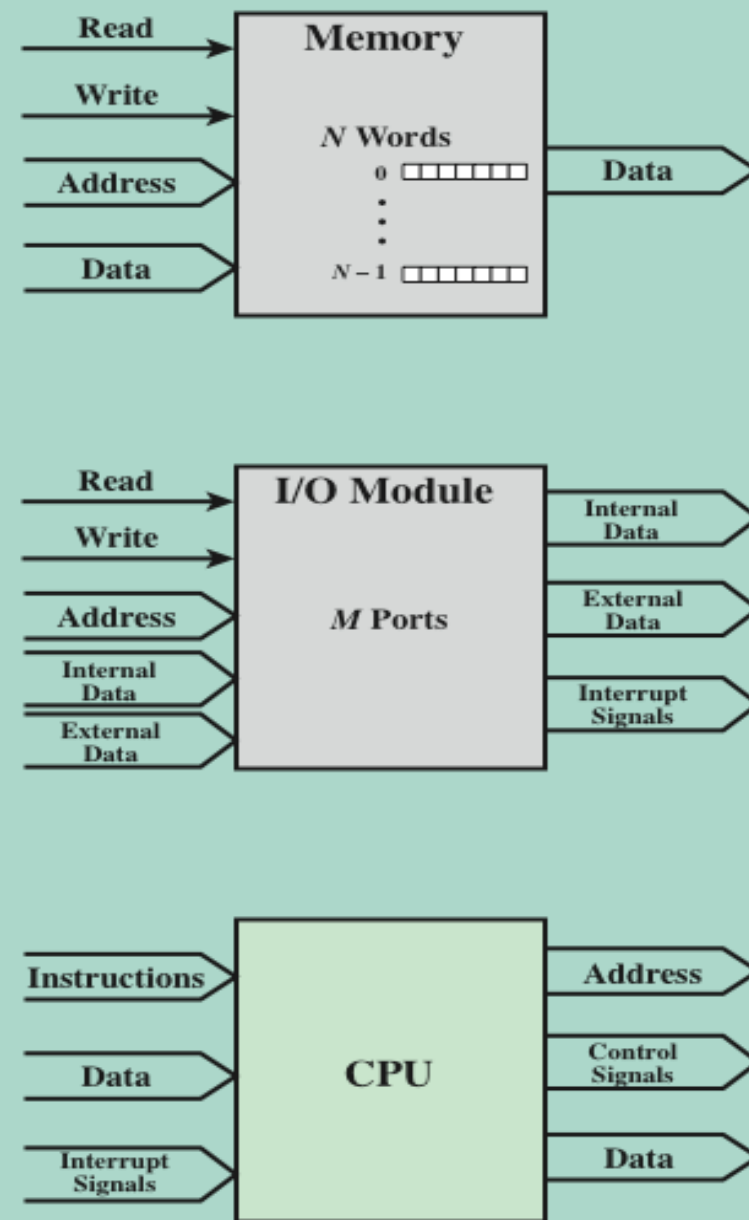


Figure 3.15 Computer Modules

The interconnection structure must support the following types of transfers:

Memory  
to  
processor

Processor  
reads an  
instruction or  
a unit of  
data from  
memory

Processor  
to  
memory

Processor  
writes a unit  
of data to  
memory

I/O to  
processor

Processor  
reads data  
from an I/O  
device via  
an I/O  
module

Processor  
to I/O

Processor  
sends data  
to the I/O  
device

I/O to or  
from  
memory

An I/O  
module is  
allowed to  
exchange  
data directly  
with memory  
without  
going  
through the  
processor  
using direct  
memory  
access

# B u s I n t e r c o n n e c t i o n

A communication pathway connecting two or more devices

- Key characteristic is that it is a shared transmission medium



Signals transmitted by any one device are available for reception by all other devices attached to the bus

- If two devices transmit during the same time period their signals will overlap and become garbled

Typically consists of multiple communication lines

- Each line is capable of transmitting signals representing binary 1 and binary 0

Computer systems contain a number of different buses that provide pathways between components at various levels of the computer system hierarchy



*System bus*

- A bus that connects major computer components (processor, memory, I/O)

The most common computer interconnection structures are based on the use of one or more system buses



# Data Bus

- Data lines that provide a path for moving data among system modules
- May consist of 32, 64, 128, or more separate lines
- The number of lines is referred to as the *width* of the data bus
- The number of lines determines how many bits can be transferred at a time
- The width of the data bus is a key factor in determining overall system performance



# Address Bus

- Used to designate the source or destination of the data on the data bus
  - If the processor wishes to read a word of data from memory it puts the address of the desired word on the address lines
- Width determines the maximum possible memory capacity of the system
- Also used to address I/O ports
  - The higher order bits are used to select a particular module on the bus and the lower order bits select a memory location or I/O port within the module

# Control Bus

- Used to control the access and the use of the data and address lines
- Because the data and address lines are shared by all components there must be a means of controlling their use
- Control signals transmit both command and timing information among system modules
- Timing signals indicate the validity of data and address information
- Command signals specify operations to be performed

# Bus Interconnection Scheme

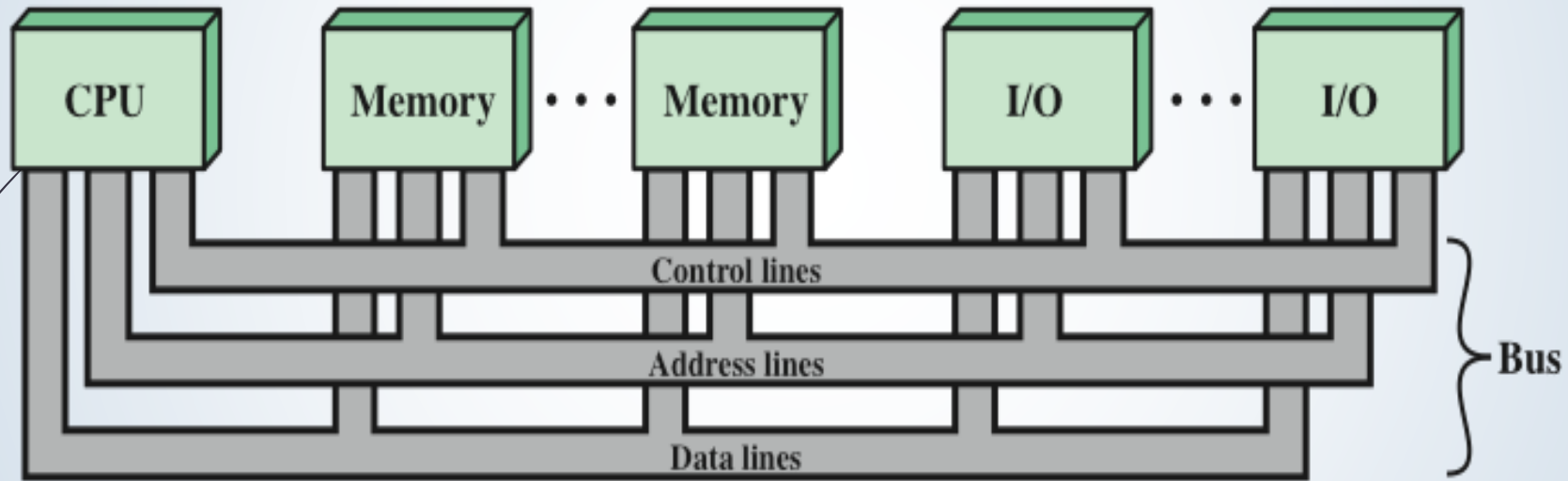
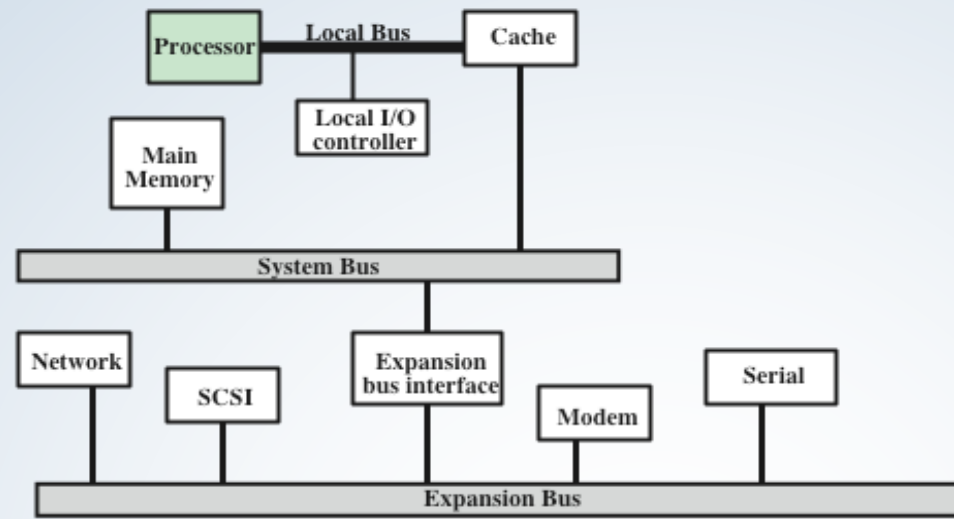
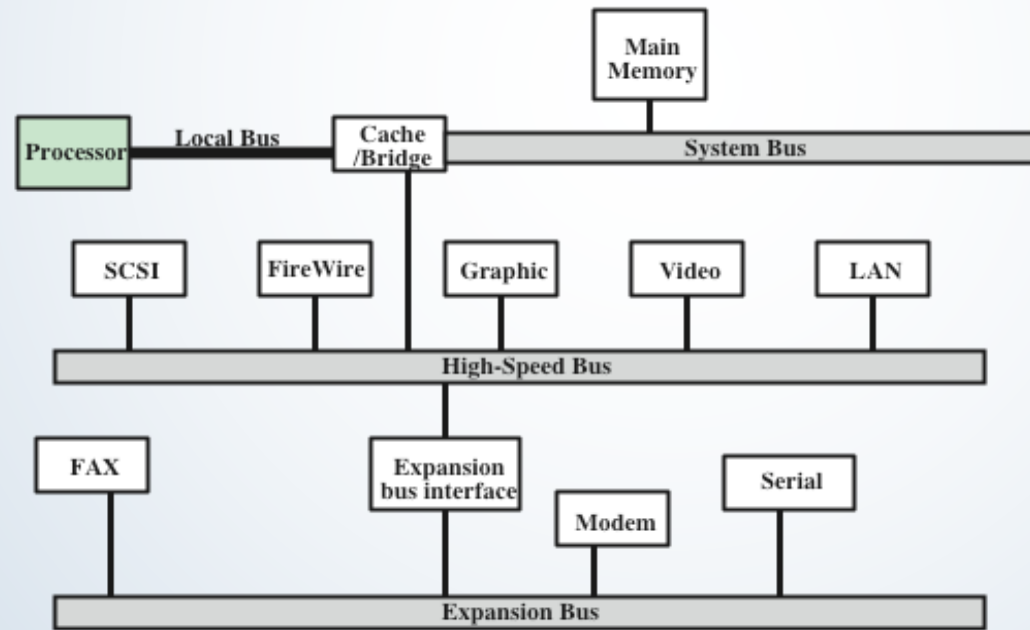


Figure 3.16 Bus Interconnection Scheme



(a) Traditional Bus Architecture



(b) High-Performance Architecture

Figure 3.17 Example Bus Configurations

B  
U  
S  
  
C  
o  
n  
f  
i  
g  
u  
r  
  
a  
t  
i  
o  
n  
s

# Elements of Bus Design

<b>Type</b>	<b>Bus Width</b>
Dedicated	Address
Multiplexed	Data
<b>Method of Arbitration</b>	<b>Data Transfer Type</b>
Centralized	Read
Distributed	Write
<b>Timing</b>	Read-modify-write
Synchronous	Read-after-write
Asynchronous	Block

# Timing of Synchronous Bus Operations

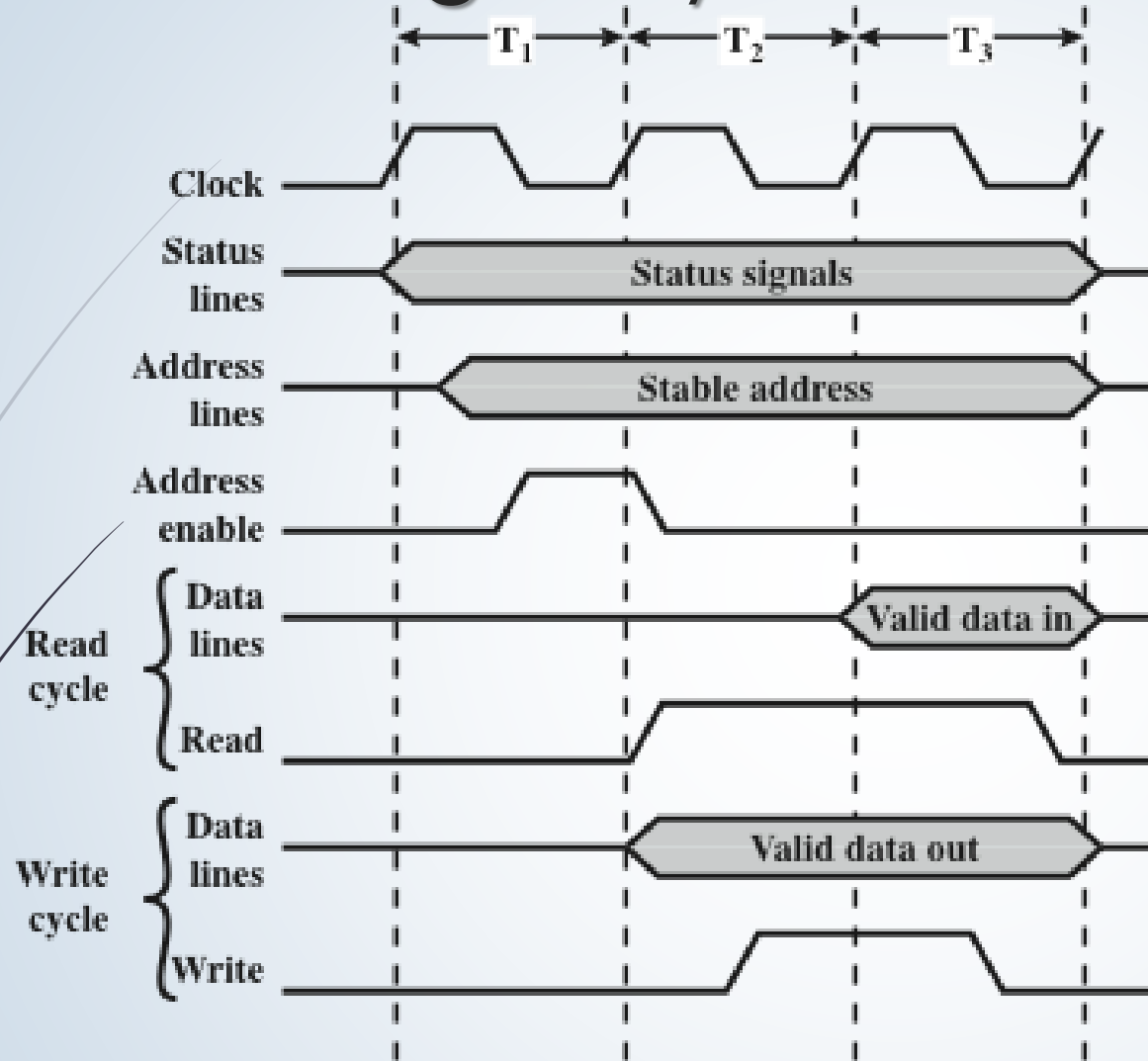
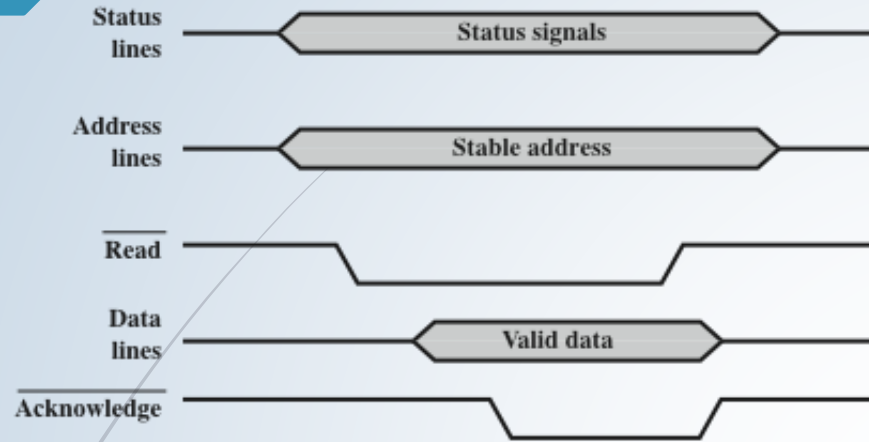


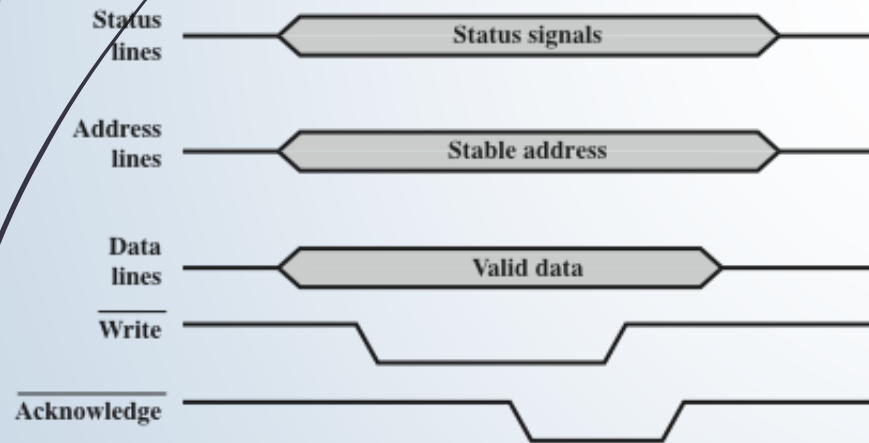
Figure 3.18 Timing of Synchronous Bus Operations



# Timing of Asynchronous Bus Operations



(a) System bus read cycle



(b) System bus write cycle

Figure 3.19 Timing of Asynchronous Bus Operations

# Point-to-Point Interconnect

Principal reason for change was the electrical constraints encountered with increasing the frequency of wide synchronous buses

At higher and higher data rates it becomes increasingly difficult to perform the synchronization and arbitration functions in a timely fashion

A conventional shared bus on the same chip magnified the difficulties of increasing bus data rate and reducing bus latency to keep up with the processors

Has lower latency, higher data rate, and better scalability

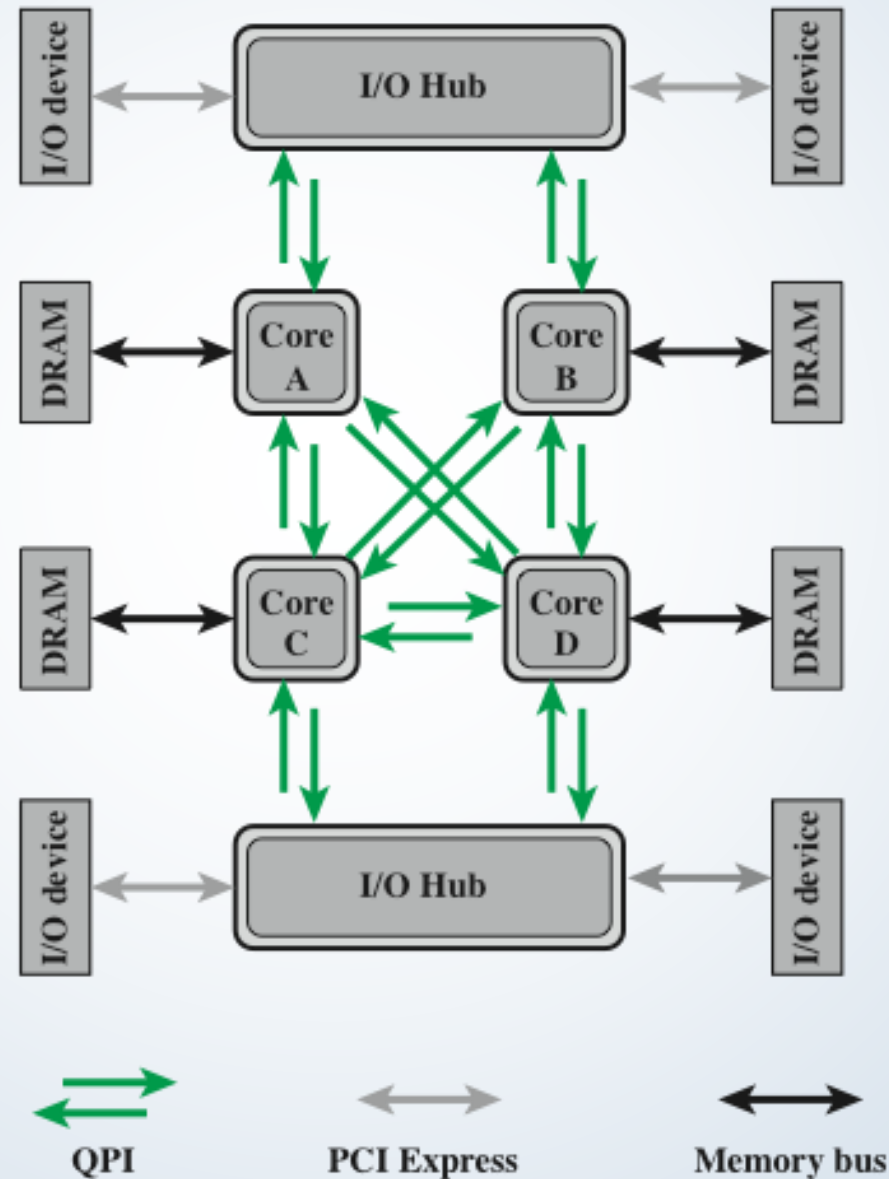


# Quick Path Interconnect

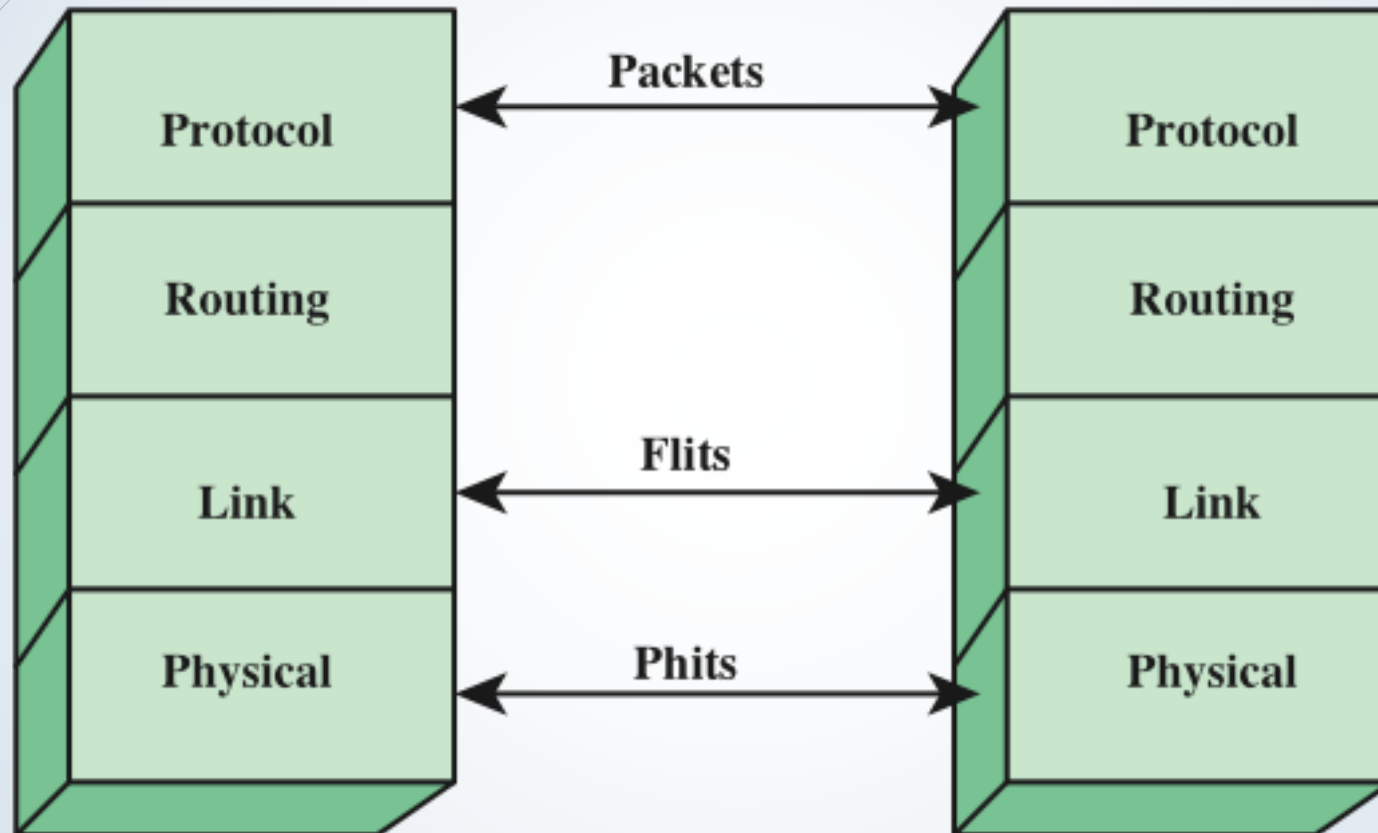
QPI

- Introduced in 2008
- Multiple direct connections
  - Direct pairwise connections to other components eliminating the need for arbitration found in shared transmission systems
- Layered protocol architecture
  - These processor level interconnects use a layered protocol architecture rather than the simple use of control signals found in shared bus arrangements
- Packetized data transfer
  - Data are sent as a sequence of packets each of which includes control headers and error control codes

# Multicore Configuration Using QPI

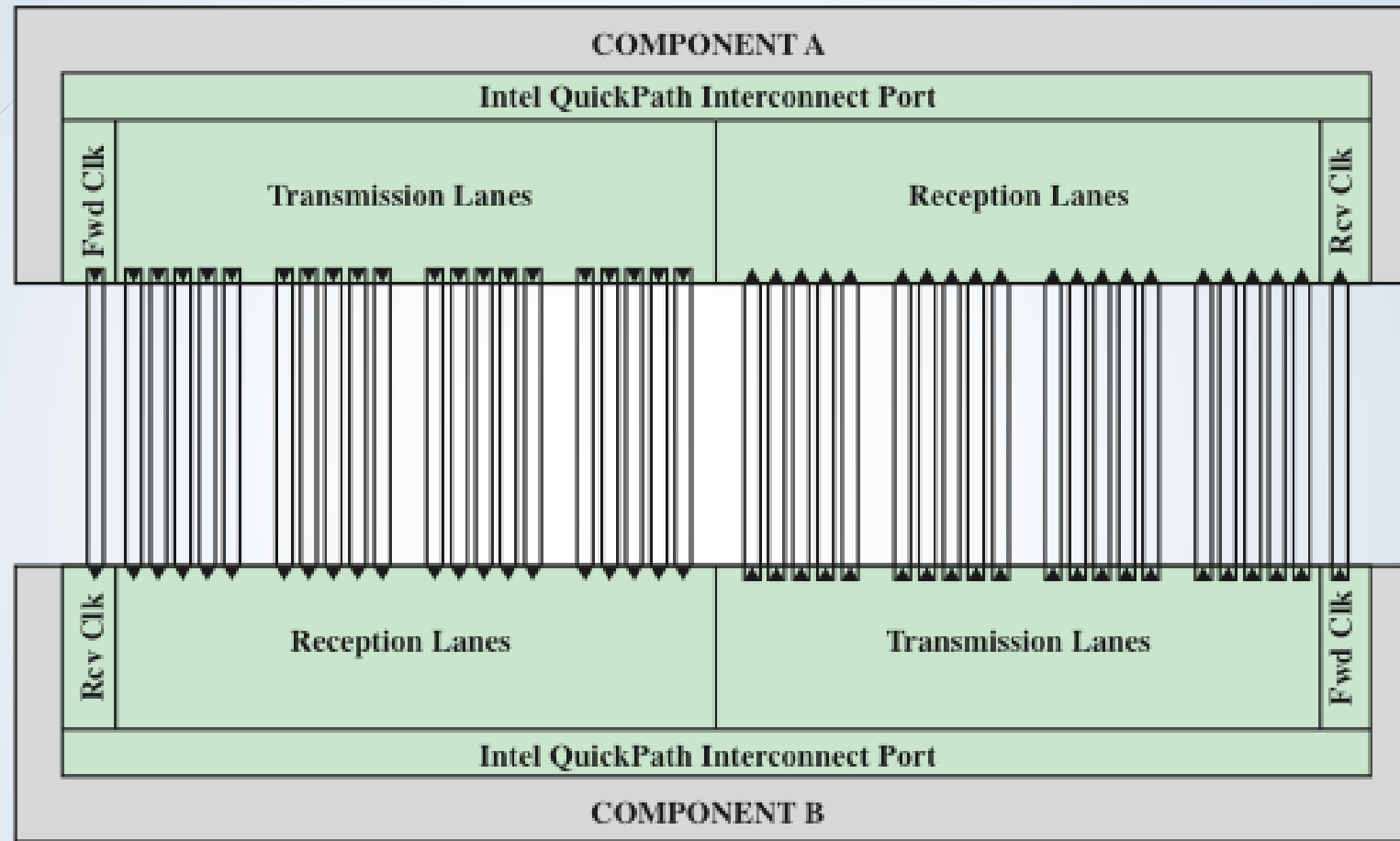


# QPI Layers



**Figure 3.21 QPI Layers**

# Physical Interface of the Intel QPI Interconnect



**Figure 3.22 Physical Interface of the Intel QPI Interconnect**



# QPI Link Layer

- ▶ Performs two key functions: *flow control* and *error control*
  - ▶ Operate on the level of the flit (flow control unit)
  - ▶ Each flit consists of a 72-bit message payload and an 8-bit error control code called a *cyclic redundancy check* (CRC)

- ▶ Flow control function
  - ▶ Needed to ensure that a sending QPI entity does not overwhelm a receiving QPI entity by sending data faster than the receiver can process the data and clear buffers for more incoming data
- ▶ Error control function
  - ▶ Detects and recovers from bit errors, and so isolates higher layers from experiencing bit errors

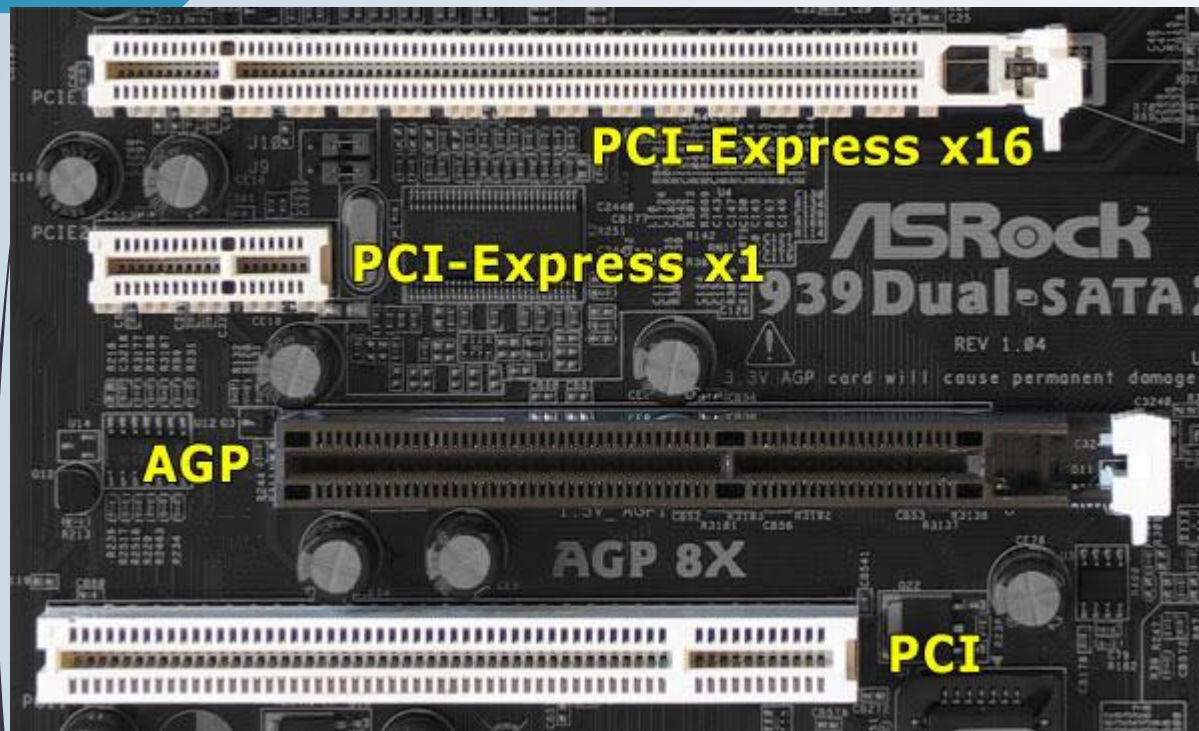
# QPI Routing and Protocol Layers

- Routing Layer
  - Used to determine the course that a packet will traverse across the available system interconnects
  - Defined by firmware and describe the possible paths that a packet can follow
- Protocol Layer
  - Packet is defined as the unit of transfer
  - One key function performed at this level is a cache coherency protocol which deals with making sure that main memory values held in multiple caches are consistent
  - A typical data packet payload is a block of data being sent to or from a cache

# Peripheral Component Interconnect (PCI)

- A popular high bandwidth, processor independent bus that can function as a mezzanine or peripheral bus
- Delivers better system performance for high speed I/O subsystems
- PCI Special Interest Group (SIG)
  - Created to develop further and maintain the compatibility of the PCI specifications
- PCI Express (PCIe)
  - Point-to-point interconnect scheme intended to replace bus-based schemes such as PCI
  - Key requirement is high capacity to support the needs of higher data rate I/O devices, such as Gigabit Ethernet
  - Another requirement deals with the need to support time dependent data streams





## PCI Express Example Connectors

**x1** **BANDWIDTH**  
 Single direction: 2.5 Gbps/200 MBps  
 Dual Directions: 5 Gbps/400 MBps



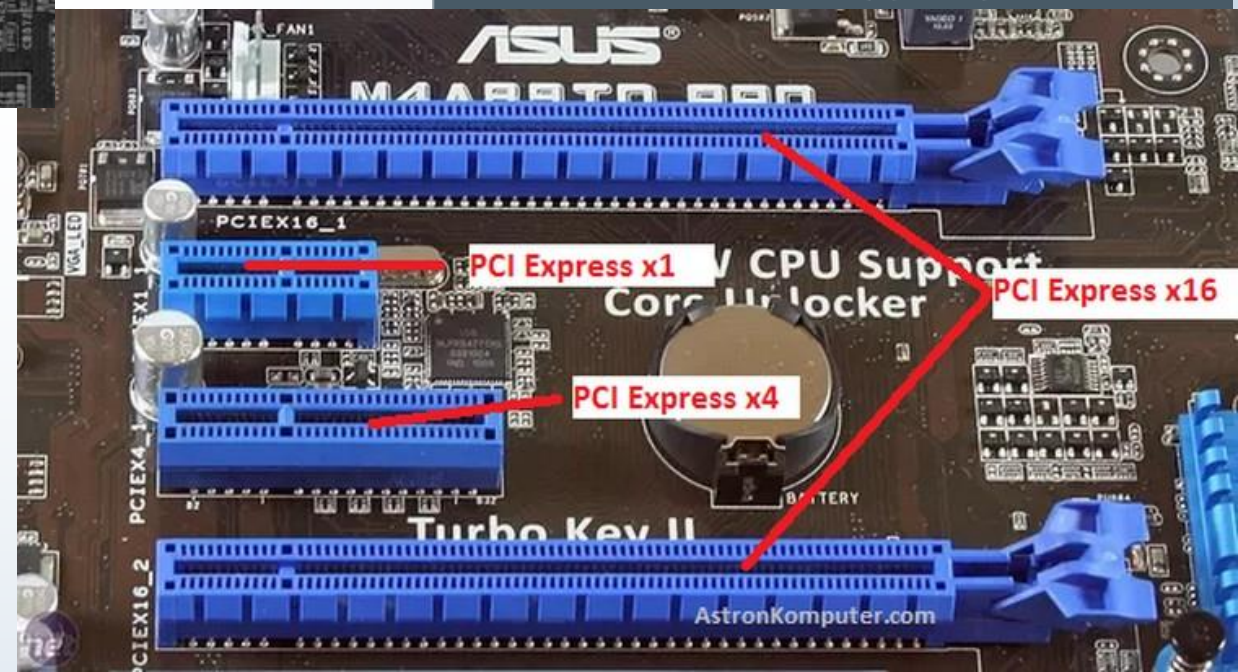
**x4** **BANDWIDTH**  
 Single direction: 10 Gbps/800 MBps  
 Dual Directions: 20 Gbps/1.6 GBps



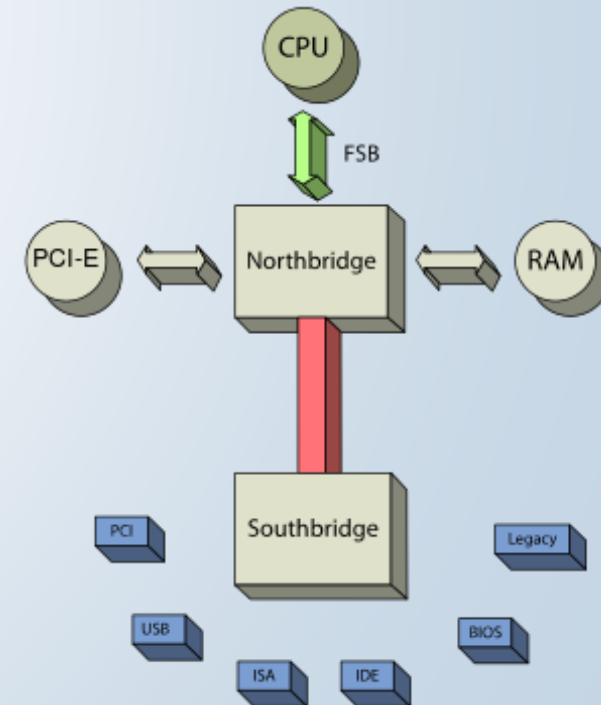
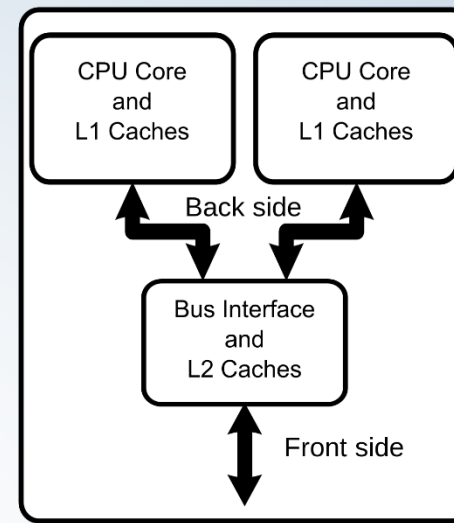
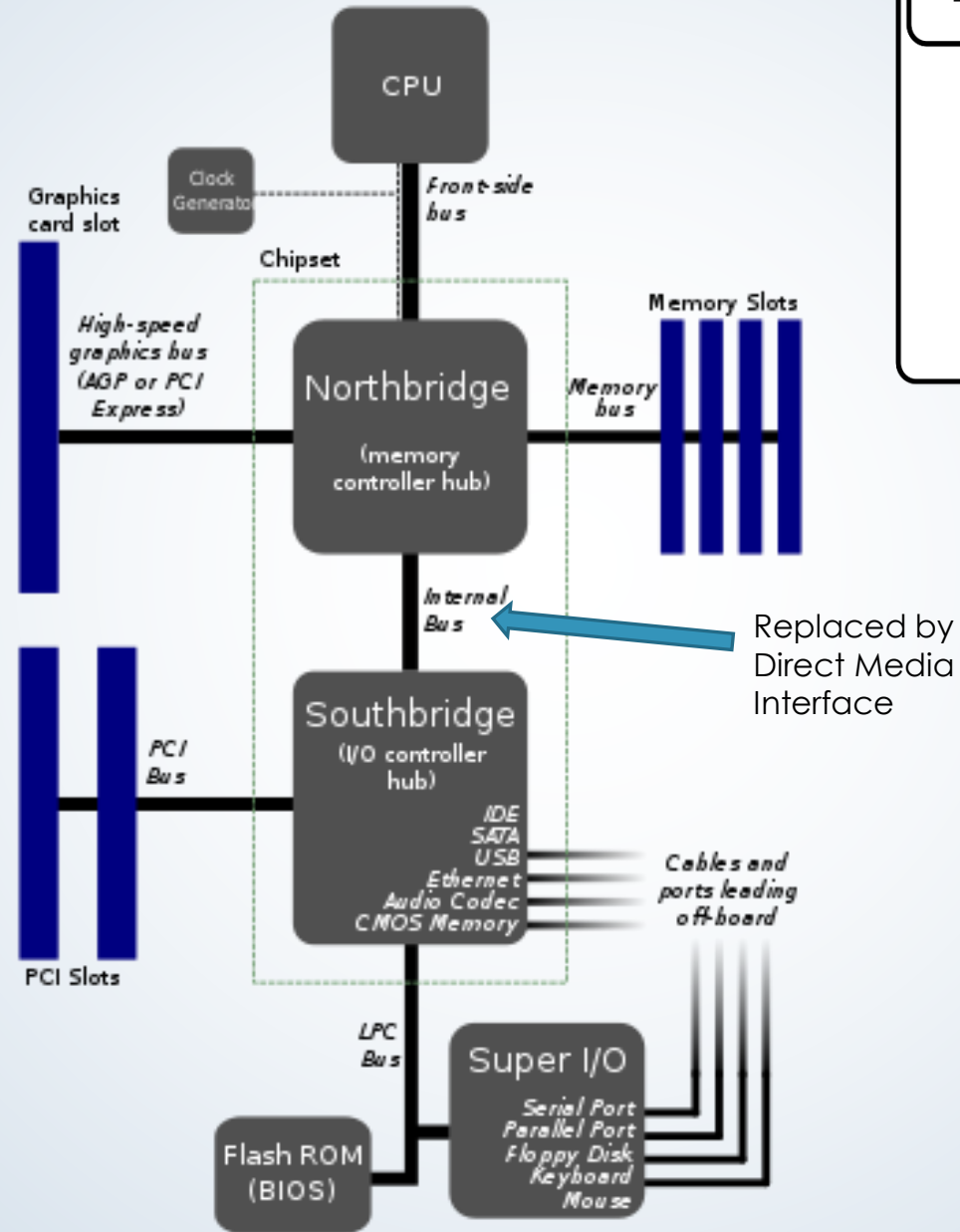
**x8** **BANDWIDTH**  
 Single direction: 20 Gbps/1.6 GBps  
 Dual Directions: 40 Gbps/3.2 GBps



**x16** **BANDWIDTH**  
 Single direction: 40 Gbps/3.2 GBps  
 Dual Directions: 80 Gbps/6.4 GBps

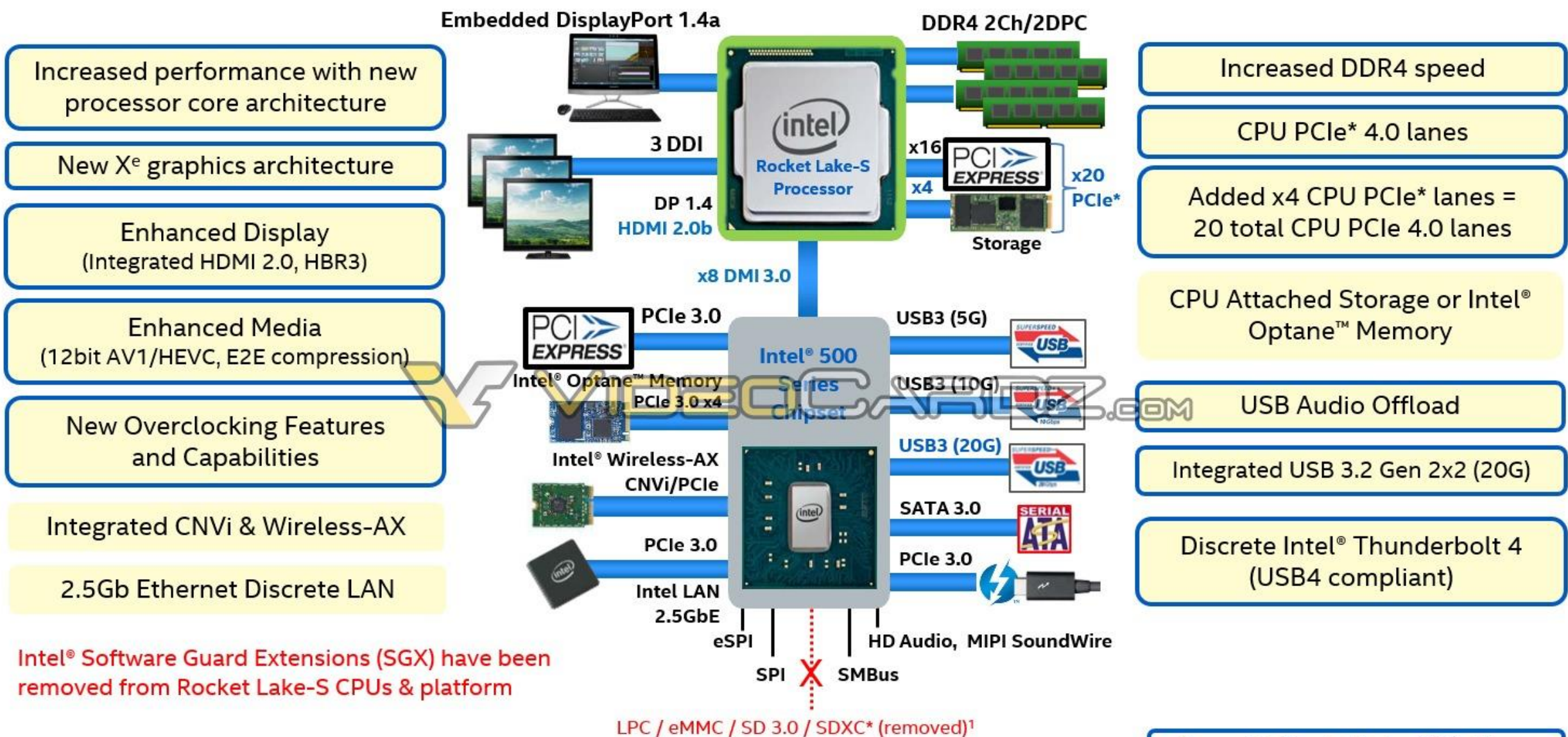


# Computer Organization





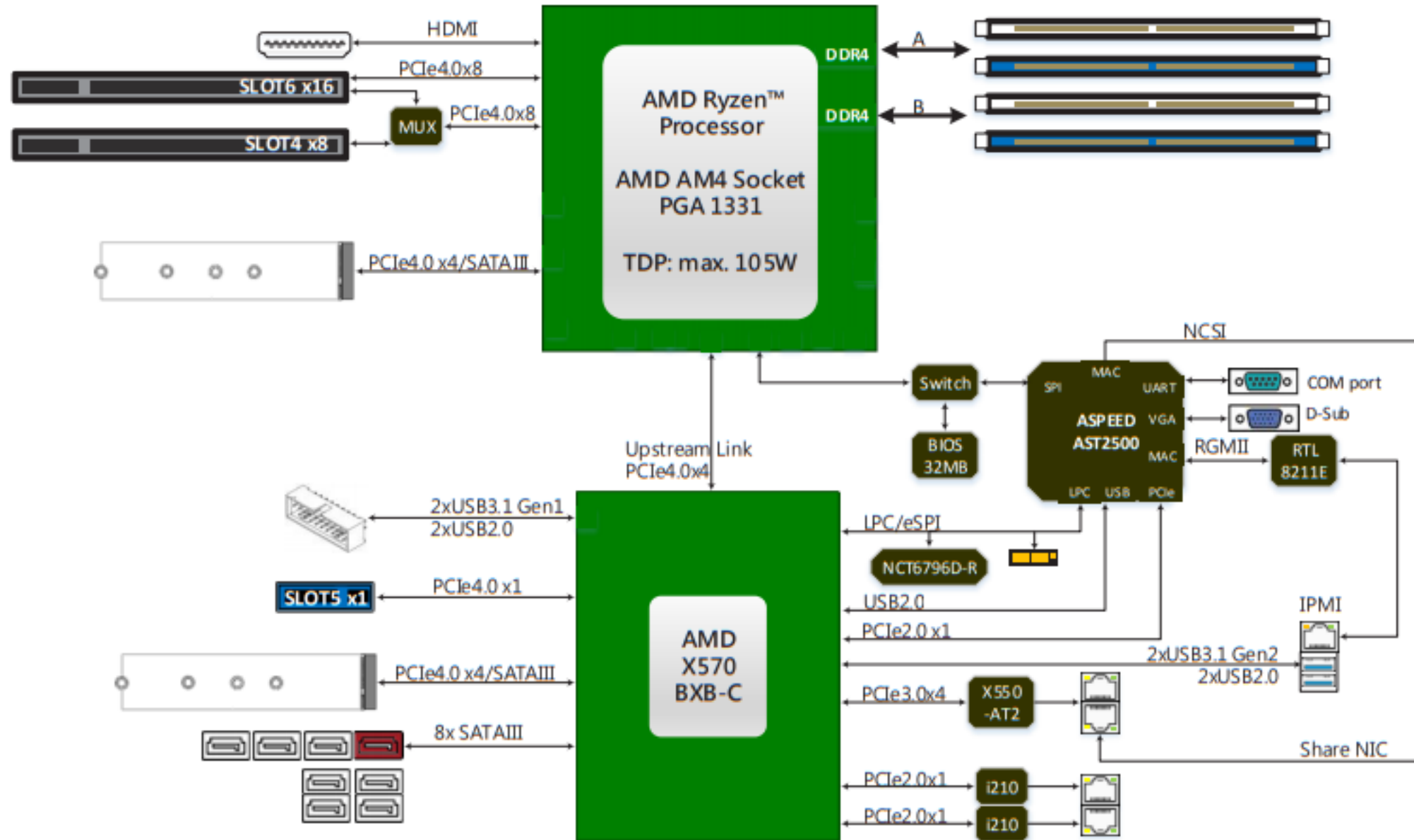
# Up to date Microarchitecture examples



New with Rocket Lake-S Platform



# Up to date Microarchitecture examples



# PCIe Configuration

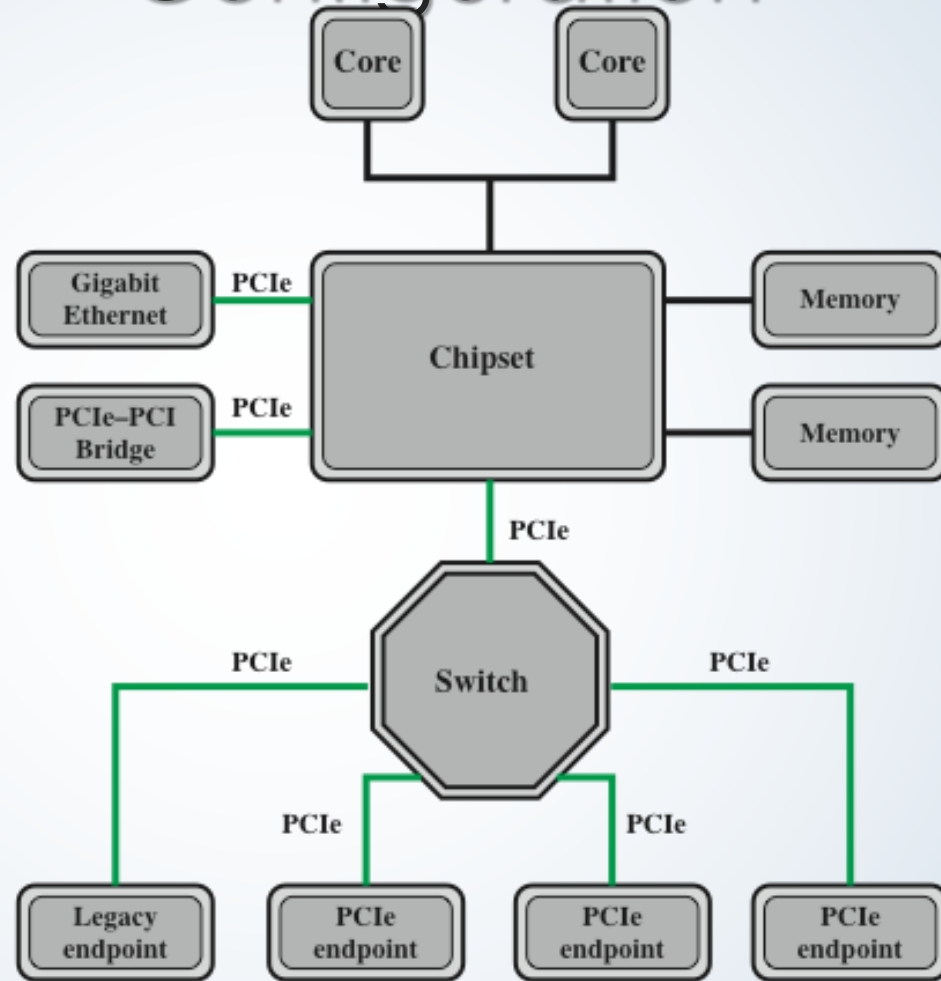
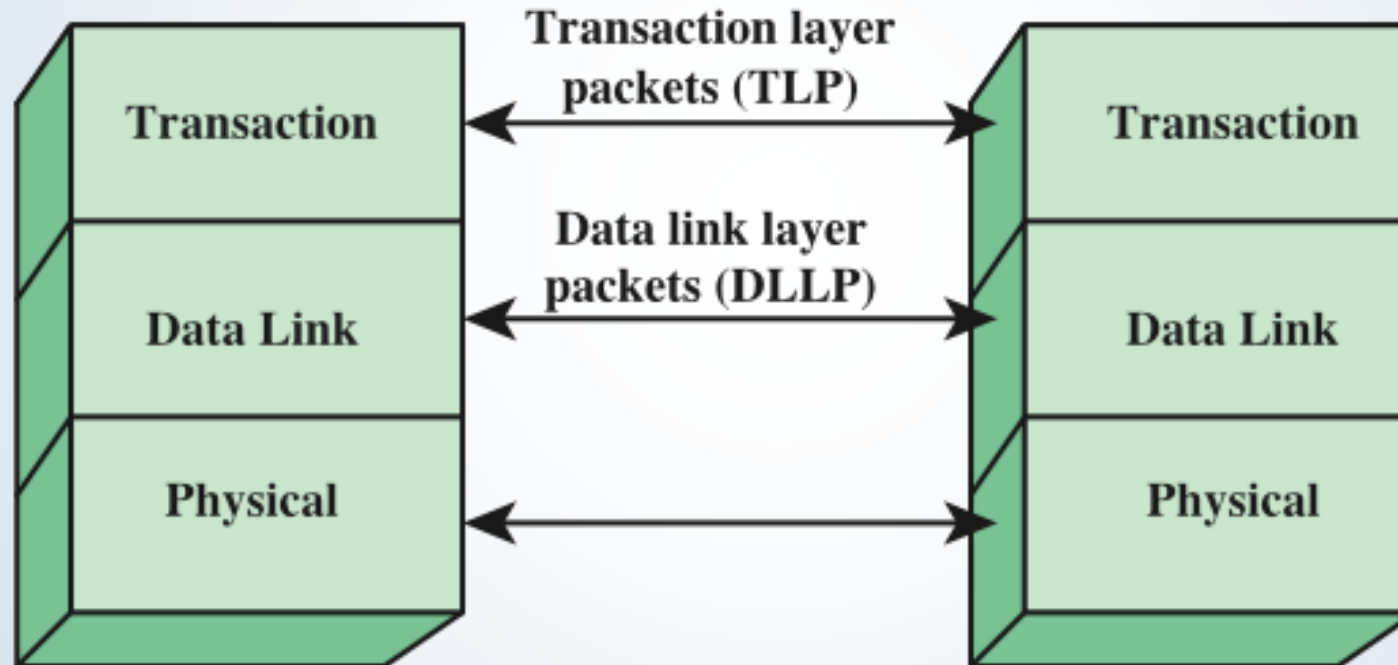
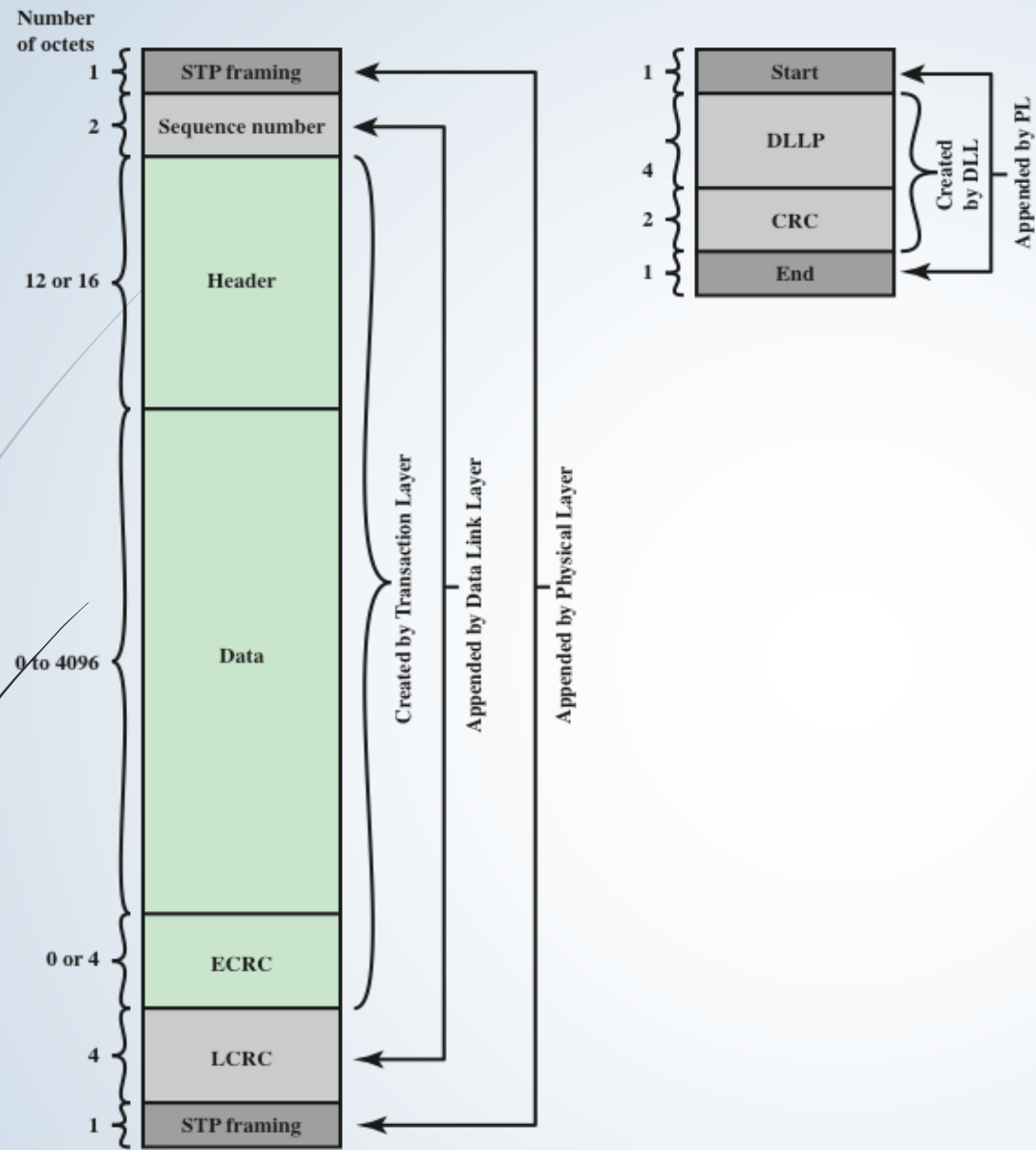


Figure 3-24 Typical Configuration Using PCIe

# PCIe Protocol Layers



**Figure 3.25 PCIe Protocol Layers**



(a) Transaction Layer Packet

(b) Data Link Layer Packet

**Figure 3.28 PCIe Protocol Data Unit Format**

# PCIe Protocol Data Unit Format

# Summary

- Computer components

- Computer function

## ➤ Chapter 3

- Instruction fetch and execute

- Interrupts

- I/O function

- Interconnection structures

- Bus interconnection

- Bus structure

- Multiple bus hierarchies

- Elements of bus design

- A Top-Level View of Computer Function and Interconnection

- Point-to-point interconnect

- QPI physical layer

- QPI link layer

- QPI routing layer

- QPI protocol layer

- PCI express

- PCI physical and logical architecture

- PCIe physical layer

- PCIe transaction layer

- PCIe data link layer