Algorithm Reference - Java

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Bellman Ford

```
bellman_ford.java
Bellman-Ford algorithm for single-source shortest paths with negative edge weights.
Finds shortest paths from a source vertex to all other vertices, even with negative
edge weights. Can detect negative cycles.
Key operations:
- addEdge(u, v, weight): Add directed edge
- shortestPaths(source): Compute shortest paths, returns null if negative cycle detected
- shortestPath(source, target): Get shortest path between two vertices
Time complexity: O(VE)
Space complexity: O(V + E)
import java.util.*;
class bellman_ford {
    static class Edge {
        int from, to, weight;
        Edge(int from, int to, int weight) {
            this.from = from;
            this.to = to;
            this.weight = weight;
        }
    }
    static class BellmanFord {
        private List<Edge> edges;
        private Set<Integer> nodes;
        private static final int INF = 999999;
        BellmanFord() {
            this.edges = new ArrayList<>();
            this.nodes = new HashSet<>();
        }
        void addEdge(int u, int v, int weight) {
            edges.add(new Edge(u, v, weight));
            nodes.add(u);
            nodes.add(v);
        Map<Integer, Integer> shortestPaths(int source) {
            Map<Integer, Integer> distances = new HashMap<>();
            for (int node : nodes) {
                distances.put(node, INF);
            distances.put(source, 0);
            // Relax edges V-1 times
            for (int i = 0; i < nodes.size() - 1; i++) {</pre>
                for (Edge e : edges) {
                    if (distances.get(e.from) != INF
                            && distances.get(e.from) + e.weight < distances.get(e.to)) {
                        distances.put(e.to, distances.get(e.from) + e.weight);
                    }
                }
            }
            // Check for negative cycles
            for (Edge e : edges) {
                if (distances.get(e.from) != INF
                        && distances.get(e.from) + e.weight < distances.get(e.to)) {
                    return null;
```

```
return distances;
    }
    List<Integer> shortestPath(int source, int target) {
        Map<Integer, Integer> distances = new HashMap<>();
        Map<Integer, Integer> predecessors = new HashMap<>();
        for (int node : nodes) {
            distances.put(node, INF);
        distances.put(source, 0);
        predecessors.put(source, null);
        for (int i = 0; i < nodes.size() - 1; i++) {</pre>
            for (Edge e : edges) {
                if (distances.get(e.from) != INF
                         && distances.get(e.from) + e.weight < distances.get(e.to)) {
                     distances.put(e.to, distances.get(e.from) + e.weight);
                     predecessors.put(e.to, e.from);
                }
            }
        }
        for (Edge e : edges) {
            if (distances.get(e.from) != INF
                     && distances.get(e.from) + e.weight < distances.get(e.to)) {
                 return null;
            }
        }
        if (!predecessors.containsKey(target)) {
            return null;
        }
        List<Integer> path = new ArrayList<>();
        Integer current = target;
        while (current != null) {
            path.add(current);
            current = predecessors.get(current);
        Collections.reverse(path);
        return path;
    }
static void testMain() {
    BellmanFord bf = new BellmanFord();
    bf.addEdge(0, 1, 4);
   bf.addEdge(0, 2, 2);
bf.addEdge(1, 2, -3);
bf.addEdge(2, 3, 2);
    bf.addEdge(3, 1, 1);
    Map<Integer, Integer> result = bf.shortestPaths(0);
    assert result != null;
    assert result.get(2) == 1;
    assert result.get(3) == 3;
    List<Integer> path = bf.shortestPath(0, 3);
    assert path != null;
    assert path.get(0) == 0;
    assert path.get(path.size() - 1) == 3;
```

}

}

Bipartite Match

```
bipartite_match.java
Maximum bipartite matching using augmenting path algorithm.
Given a bipartite graph with left and right vertex sets, finds the maximum
number of edges such that no two edges share a vertex.
Key operations:
- addEdge(u, v): Add edge from left vertex u to right vertex v
- maxMatching(): Compute maximum matching size
Time complexity: O(V * E)
Space complexity: O(V + E)
import java.util.*;
class bipartite_match {
    static class BipartiteMatch {
        private int leftSize;
        private int rightSize;
        private Map<Integer, List<Integer>> graph;
        private Map<Integer, Integer> match;
        private Set<Integer> visited;
        BipartiteMatch(int leftSize, int rightSize) {
            this.leftSize = leftSize;
            this.rightSize = rightSize;
            this.graph = new HashMap<>();
            for (int i = 0; i < leftSize; i++) {</pre>
                graph.put(i, new ArrayList<>());
        }
        void addEdge(int u, int v) {
            graph.get(u).add(v);
        int maxMatching() {
            match = new HashMap<>();
            int matchingSize = 0;
            for (int u = 0; u < leftSize; u++) {
                visited = new HashSet<>();
                if (dfs(u)) {
                    matchingSize++;
            }
            return matchingSize;
        }
        private boolean dfs(int u) {
            for (int v : graph.get(u)) {
                if (visited.contains(v)) {
                    continue;
                visited.add(v);
                // If v is not matched or we can find augmenting path from match[v]
                if (!match.containsKey(v) || dfs(match.get(v))) {
                    match.put(v, u);
                    return true;
                }
            return false;
        }
```

```
Map<Integer, Integer> getMatching() {
        Map<Integer, Integer> result = new HashMap<>();
        for (Map.Entry<Integer, Integer> entry : match.entrySet()) {
             result.put(entry.getValue(), entry.getKey());
        return result;
    }
}
static void testMain() {
    BipartiteMatch b = new BipartiteMatch(3, 3);
    b.addEdge(0, 0); // 1 -> X
b.addEdge(1, 1); // 2 -> Y
b.addEdge(2, 0); // 3 -> X
    b.addEdge(0, 2); // 1 -> Z
    b.addEdge(1, 2); // 2 -> Z
    b.addEdge(2, 1); // 3 -> Y
    int matching = b.maxMatching();
    if (matching != 3) throw new AssertionError("Expected 3, got " + matching);
    Map<Integer, Integer> matches = b.getMatching();
    if (matches.size() != 3) throw new AssertionError("Expected size 3, got " + matches.size());
}
```

Convex Hull

```
convex_hull.java
Andrew's monotone chain algorithm for computing the convex hull of 2D points.
Computes the convex hull (smallest convex polygon containing all points) using
a simple and robust algorithm.
Key operations:
- convexHull(points): Returns convex hull vertices in counter-clockwise order
Time complexity: O(n \log n) dominated by sorting
Space complexity: O(n)
import java.util.*;
class convex_hull {
    static class Point implements Comparable<Point> {
        double x, y;
        Point(double x, double y) {
            this.x = x;
            this.y = y;
        }
        @Override
        public int compareTo(Point other) {
            if (this.x != other.x) {
                return Double.compare(this.x, other.x);
            return Double.compare(this.y, other.y);
        }
        @Override
        public boolean equals(Object obj) {
            if (!(obj instanceof Point)) return false;
            Point p = (Point) obj;
            return this.x == p.x && this.y == p.y;
        }
        @Override
        public int hashCode() {
            return Objects.hash(x, y);
        }
   }
    static double cross(Point o, Point a, Point b) {
        return (a.x - o.x) * (b.y - o.y) - (a.y - o.y) * (b.x - o.x);
    }
    static List<Point> convexHull(List<Point> points) {
        if (points.size() <= 1) {</pre>
            return new ArrayList<>(points);
        }
        List<Point> sorted = new ArrayList<>(points);
        Collections.sort(sorted);
        List<Point> lower = new ArrayList<>();
        for (Point p : sorted) {
            while (lower.size() >= 2
                    && cross(lower.get(lower.size() - 2), lower.get(lower.size() - 1), p) <= 0) {
                lower.remove(lower.size() - 1);
            lower.add(p);
        }
```

```
List<Point> upper = new ArrayList<>();
    for (int i = sorted.size() - 1; i >= 0; i--) {
        Point p = sorted.get(i);
        while (upper.size() >= 2
                && cross(upper.get(upper.size() - 2), upper.get(upper.size() - 1), p) <= 0) {
            upper.remove(upper.size() - 1);
        upper.add(p);
    }
    lower.remove(lower.size() - 1);
    upper.remove(upper.size() - 1);
    lower.addAll(upper);
    return lower;
}
static void testMain() {
    List<Point> pts =
            Arrays.asList(
                    new Point(0, 0),
                    new Point(1, 0),
                    new Point(1, 1),
                    new Point(0, 1),
                    new Point(0.5, 0.5));
    List<Point> hull = convexHull(pts);
    assert hull.size() == 4;
    assert hull.contains(new Point(0, 0));
    assert hull.contains(new Point(1, 0));
    assert hull.contains(new Point(1, 1));
    assert hull.contains(new Point(0, 1));
    assert !hull.contains(new Point(0.5, 0.5));
}
```

Dijkstra

```
dijkstra.java
Dijkstra's algorithm for single-source shortest paths in weighted graphs.
Finds shortest paths from a source vertex to all other vertices in a graph
with non-negative edge weights.
Key operations:
- addEdge(u, v, weight): Add weighted directed edge
- shortestPaths(source): Compute shortest paths from source to all vertices
- shortestPath(source, target): Get shortest path between two vertices
Time complexity: O((V + E) \log V) with binary heap
Space complexity: O(V + E)
import java.util.*;
class dijkstra {
    static class Edge {
        int to;
        int weight;
        Edge(int to, int weight) {
            this.to = to;
            this.weight = weight;
        }
   }
    static class Node implements Comparable<Node> {
        int vertex;
        int distance;
        Node(int vertex, int distance) {
            this.vertex = vertex;
            this.distance = distance;
        }
        @Override
        public int compareTo(Node other) {
            return Integer.compare(this.distance, other.distance);
    }
    static class Dijkstra {
        private int n;
        private Map<Integer, List<Edge>> graph;
        Dijkstra(int n) {
            this.n = n;
            this.graph = new HashMap<>();
            for (int i = 0; i < n; i++) {
                graph.put(i, new ArrayList<>());
            }
        }
        void addEdge(int u, int v, int weight) {
            graph.get(u).add(new Edge(v, weight));
        }
        Map<Integer, Integer> shortestPaths(int source) {
            Map<Integer, Integer> distances = new HashMap<>();
            for (int i = 0; i < n; i++) {
                distances.put(i, Integer.MAX_VALUE);
            distances.put(source, 0);
```

```
PriorityQueue<Node> pq = new PriorityQueue<>();
    pq.offer(new Node(source, 0));
    while (!pq.isEmpty()) {
        Node current = pq.poll();
        int u = current.vertex;
        int dist = current.distance;
        if (dist > distances.get(u)) {
            continue;
        }
        for (Edge edge : graph.get(u)) {
            int v = edge.to;
            int newDist = dist + edge.weight;
            if (newDist < distances.get(v)) {</pre>
                distances.put(v, newDist);
                pq.offer(new Node(v, newDist));
            }
        }
    }
    return distances;
}
List<Integer> shortestPath(int source, int target) {
    Map<Integer, Integer> distances = new HashMap<>();
    Map<Integer, Integer> previous = new HashMap<>();
    for (int i = 0; i < n; i++) {
        distances.put(i, Integer.MAX_VALUE);
    distances.put(source, 0);
    PriorityQueue<Node> pq = new PriorityQueue<>();
    pq.offer(new Node(source, 0));
    while (!pq.isEmpty()) {
        Node current = pq.poll();
        int u = current.vertex;
        int dist = current.distance;
        if (u == target) {
            break;
        }
        if (dist > distances.get(u)) {
            continue;
        for (Edge edge : graph.get(u)) {
            int v = edge.to;
            int newDist = dist + edge.weight;
            if (newDist < distances.get(v)) {</pre>
                distances.put(v, newDist);
                previous.put(v, u);
                pq.offer(new Node(v, newDist));
            }
        }
    }
    if (!previous.containsKey(target) && target != source) {
        return null;
    List<Integer> path = new ArrayList<>();
    int current = target;
    while (current != source) {
        path.add(current);
        current = previous.get(current);
```

```
    path.add(source);
    Collections.reverse(path);

    return path;
}

static void testMain() {
    Dijkstra d = new Dijkstra(4);
    d.addEdge(0, 1, 4);
    d.addEdge(0, 2, 2);
    d.addEdge(1, 2, 1);
    d.addEdge(1, 3, 5);
    d.addEdge(2, 3, 8);

Map<Integer, Integer> distances = d.shortestPaths(0);
    assert distances.get(3) == 9;

List<Integer> path = d.shortestPath(0, 3);
    assert path.equals(Arrays.asList(0, 1, 3));
}
```

Edmonds Karp

```
edmonds_karp.java
Edmonds-Karp algorithm for computing maximum flow in a flow network.
Implementation of the Ford-Fulkerson method using BFS to find augmenting paths.
Guarantees O(V * E^2) time complexity.
Key operations:
- addEdge(u, v, capacity): Add a directed edge with given capacity
- maxFlow(source, sink): Compute maximum flow from source to sink
Space complexity: O(V^2) for adjacency matrix representation
import java.util.*;
class edmonds_karp {
    static class EdmondsKarp {
        private int n;
        private int[][] capacity;
        private int[][] flow;
        EdmondsKarp(int n) {
            this.n = n;
            this.capacity = new int[n][n];
            this.flow = new int[n][n];
        }
        void addEdge(int u, int v, int cap) {
            capacity[u][v] += cap;
        }
        int maxFlow(int source, int sink) {
            // Reset flow
            for (int i = 0; i < n; i++) {
                Arrays.fill(flow[i], 0);
            int totalFlow = 0;
            while (true) {
                // BFS to find augmenting path
                int[] parent = new int[n];
                Arrays.fill(parent, -1);
                parent[source] = source;
                Queue<Integer> queue = new LinkedList<>();
                queue.offer(source);
                while (!queue.isEmpty() && parent[sink] == -1) {
                    int u = queue.poll();
                    for (int v = 0; v < n; v++) {
                        if (parent[v] == -1 \&\& capacity[u][v] - flow[u][v] > 0) {
                            parent[v] = u;
                            queue.offer(v);
                        }
                    }
                }
                // No augmenting path found
                if (parent[sink] == -1) {
                    break;
                }
                // Find minimum residual capacity along the path
                int pathFlow = Integer.MAX_VALUE;
```

```
int v = sink;
            while (v != source) {
                int u = parent[v];
                pathFlow = Math.min(pathFlow, capacity[u][v] - flow[u][v]);
                v = u;
            }
            // Update flow along the path
            v = sink;
            while (v != source) {
                int u = parent[v];
                flow[u][v] += pathFlow;
                flow[v][u] -= pathFlow;
                v = u;
            }
            totalFlow += pathFlow;
        }
        return totalFlow;
    }
    int getFlow(int u, int v) {
        return flow[u][v];
    }
}
static void testMain() {
    EdmondsKarp e = new EdmondsKarp(4);
    e.addEdge(0, 1, 10);
    e.addEdge(0, 2, 8);
    e.addEdge(1, 2, 2);
    e.addEdge(1, 3, 5);
    e.addEdge(2, 3, 7);
    int maxFlow = e.maxFlow(0, 3);
    assert maxFlow == 12;
}
```

Fenwick Tree

```
fenwick_tree.java
Fenwick Tree (Binary Indexed Tree) implementation.
A data structure for efficient prefix sum queries and point updates on an array.
Key operations:
- update(i, delta): Add delta to element at index i - O(log n)
- query(i): Get sum of elements from index 0 to i (inclusive) - O(log n)
- range_query(l, r): Get sum from index l to r (inclusive) - O(log n)
Space complexity: O(n)
Note: Uses 1-based indexing internally for simpler bit manipulation.
import java.util.function.BinaryOperator;
class fenwick_tree {
    interface Summable<T> {
        T add(T other);
        T subtract(T other);
    }
    static class FenwickTree<T> {
        private Object[] tree;
        private int size;
        private T zero;
        private BinaryOperator<T> addOp;
        private BinaryOperator<T> subtractOp;
        FenwickTree(int n, T zero, BinaryOperator<T> addOp, BinaryOperator<T> subtractOp) {
            this.size = n;
            this.zero = zero;
            this.addOp = addOp;
            this.subtractOp = subtractOp;
            this.tree = new Object[n + 1];
            for (int i = 0; i <= n; i++) {</pre>
                tree[i] = zero;
            }
        }
        // O(n log n) constructor from array
        FenwickTree(T[] values, T zero, BinaryOperator<T> addOp, BinaryOperator<T> subtractOp) {
            this(values.length, zero, addOp, subtractOp);
            for (int i = 0; i < values.length; <math>i++) {
                update(i, values[i]);
            }
        }
        // O(n) constructor from array using prefix sums
        static <T> FenwickTree<T> fromArray(
                T[] arr, T zero, BinaryOperator<T> addOp, BinaryOperator<T> subtractOp) {
            int n = arr.length;
            FenwickTree<T> ft = new FenwickTree<>(n, zero, addOp, subtractOp);
            // Compute prefix sums
            Object[] prefix = new Object[n + 1];
            prefix[0] = zero;
            for (int i = 0; i < n; i++) {
                prefix[i + 1] = addOp.apply((T) prefix[i], arr[i]);
            }
            // Build tree in O(n): each tree[i] contains sum of range [i - (i & -i) + 1, i]
            for (int i = 1; i <= n; i++) {</pre>
                int rangeStart = i - (i \& (-i)) + 1;
```

```
ft.tree[i] = subtractOp.apply((T) prefix[i], (T) prefix[rangeStart - 1]);
    }
    return ft;
}
@SuppressWarnings("unchecked")
void update(int i, T delta) {
    if (i < 0 \mid | i >= size) {
        throw new IndexOutOfBoundsException(
                "Index " + i + " out of bounds for size " + size);
    i++; // Convert to 1-based indexing
    while (i <= size) {</pre>
        tree[i] = addOp.apply((T) tree[i], delta);
        i += i & (-i);
    }
}
@SuppressWarnings("unchecked")
T query(int i) {
    if (i < 0 || i >= size) {
        throw new IndexOutOfBoundsException(
                "Index " + i + " out of bounds for size " + size);
    i++; // Convert to 1-based indexing
    T sum = zero;
    while (i > 0) {
        sum = addOp.apply(sum, (T) tree[i]);
        i -= i & (-i);
    return sum;
}
T rangeQuery(int l, int r) {
    if (l > r || l < 0 || r >= size) {
        return zero;
    if (l == 0) {
        return query(r);
    return subtractOp.apply(query(r), query(l - 1));
}
// Optional functionality (not always needed during competition)
T getValue(int i) {
    if (i < 0 || i >= size) {
        throw new IndexOutOfBoundsException(
                "Index " + i + " out of bounds for size " + size);
    if (i == 0) {
        return query(0);
    return subtractOp.apply(query(i), query(i - 1));
}
// Find smallest index >= startIndex with value > zero
// REQUIRES: all updates are non-negative, T must be comparable
@SuppressWarnings("unchecked")
Integer firstNonzeroIndex(int startIndex, java.util.Comparator<T> comparator) {
    startIndex = Math.max(startIndex, 0);
    if (startIndex >= size) {
        return null;
    }
    T prefixBefore = startIndex > 0 ? query(startIndex - 1) : zero;
    T total = query(size - 1);
    if (comparator.compare(total, prefixBefore) == 0) {
        return null;
    }
```

```
// Fenwick lower_bound: first idx with prefix_sum(idx) > prefixBefore
        int idx = 0; // 1-based cursor
        T cur = zero; // running prefix at 'idx'
        int bit = Integer.highestOneBit(size);
        while (bit > 0) {
            int nxt = idx + bit;
            if (nxt <= size) {</pre>
                T cand = addOp.apply(cur, (T) tree[nxt]);
                if (comparator.compare(cand, prefixBefore)
                        <= 0) { // move right while prefix <= target
                    cur = cand;
                    idx = nxt;
                }
            bit >>= 1;
        }
        // idx is the largest position with prefix <= prefixBefore (1-based).
        // The answer is idx (converted to 0-based).
        return idx;
    }
}
static void testMain() {
    FenwickTree<Long> f = new FenwickTree<>(5, 0L, (a, b) -> a + b, (a, b) -> a - b);
    f.update(0, 7L);
    f.update(2, 13L);
    f.update(4, 19L);
    assert f.query(4) == 39L;
    assert f.rangeQuery(1, 3) == 13L;
   // Optional functionality (not always needed during competition)
    assert f.getValue(2) == 13L;
    FenwickTree<Long> g =
            FenwickTree.fromArray(
                    new Long[] \{1L, 2L, 3L, 4L, 5L\}, 0L, (a, b) -> a + b, (a, b) -> a - b);
    assert g.query(4) == 15L;
}
```

Kmp

```
kmp.java
Knuth-Morris-Pratt (KMP) string matching algorithm.
Efficiently finds all occurrences of a pattern in a text string.
Key operations:
- computeLPS(pattern): Compute Longest Proper Prefix which is also Suffix array
- search(text, pattern): Find all starting positions where pattern occurs in text
Time complexity: O(n + m) where n is text length and m is pattern length
Space complexity: O(m) for the LPS array
import java.util.*;
class kmp {
    static int[] computeLPS(String pattern) {
        int m = pattern.length();
        int[] lps = new int[m];
        int len = 0;
        int i = 1;
        lps[0] = 0;
        while (i < m) {
            if (pattern.charAt(i) == pattern.charAt(len)) {
                len++;
                lps[i] = len;
                i++;
            } else {
                if (len != 0) {
                    len = lps[len - 1];
                } else {
                    lps[i] = 0;
                    i++;
                }
            }
        }
        return lps;
    static List<Integer> search(String text, String pattern) {
        List<Integer> result = new ArrayList<>();
        if (pattern.isEmpty()) {
            return result;
        }
        int n = text.length();
        int m = pattern.length();
        int[] lps = computeLPS(pattern);
        int i = 0; // index for text
        int j = 0; // index for pattern
        while (i < n) {
            if (text.charAt(i) == pattern.charAt(j)) {
                i++;
                j++;
            }
            if (j == m) {
                result.add(i - j);
                j = lps[j - 1];
            } else if (i < n && text.charAt(i) != pattern.charAt(j)) {</pre>
```

Kosaraju Scc

```
kosaraju_scc.java
Kosaraju's algorithm for finding strongly connected components (SCCs) in directed graphs.
A strongly connected component is a maximal set of vertices where every vertex is
reachable from every other vertex in the set. Uses two DFS passes.
Key operations:
- addEdge(u, v): Add directed edge from u to v
- findSCCs(): Find all strongly connected components
Time complexity: O(V + E)
Space complexity: O(V + E)
import java.util.*;
class kosaraju_scc {
    static class KosarajuSCC {
        private Map<Integer, List<Integer>> graph;
        private Map<Integer, List<Integer>> transpose;
        KosarajuSCC() {
            this.graph = new HashMap<>();
            this.transpose = new HashMap<>();
        }
        void addEdge(int u, int v) {
            graph.putIfAbsent(u, new ArrayList<>());
            graph.putIfAbsent(v, new ArrayList<>());
            transpose.putIfAbsent(u, new ArrayList<>());
            transpose.putIfAbsent(v, new ArrayList<>());
            graph.get(u).add(v);
            transpose.get(v).add(u);
        }
        List<List<Integer>> findSCCs() {
            // First DFS pass: compute finish order
            Set<Integer> visited = new HashSet<>();
            List<Integer> finishOrder = new ArrayList<>();
            for (int node : graph.keySet()) {
                if (!visited.contains(node)) {
                    dfs1(node, visited, finishOrder);
                }
            }
            // Second DFS pass: find SCCs on transpose graph
            visited.clear();
            List<List<Integer>> sccs = new ArrayList<>();
            for (int i = finishOrder.size() - 1; i >= 0; i--) {
                int node = finishOrder.get(i);
                if (!visited.contains(node)) {
                    List<Integer> scc = new ArrayList<>();
                    dfs2(node, visited, scc);
                    sccs.add(scc);
                }
            }
            return sccs;
        }
        private void dfs1(int node, Set<Integer> visited, List<Integer> finishOrder) {
            visited.add(node);
            if (graph.containsKey(node)) {
```

```
for (int neighbor : graph.get(node)) {
                if (!visited.contains(neighbor)) {
                    dfs1(neighbor, visited, finishOrder);
            }
        finishOrder.add(node);
    }
    private void dfs2(int node, Set<Integer> visited, List<Integer> scc) {
        visited.add(node);
        scc.add(node);
        if (transpose.containsKey(node)) {
            for (int neighbor : transpose.get(node)) {
                if (!visited.contains(neighbor)) {
                    dfs2(neighbor, visited, scc);
                }
            }
        }
    }
}
static void testMain() {
    KosarajuSCC g = new KosarajuSCC();
    g.addEdge(0, 1);
    g.addEdge(1, 2);
    g.addEdge(2, 0);
    g.addEdge(1, 3);
    g.addEdge(3, 4);
    g.addEdge(4, 5);
    g.addEdge(5, 3);
    List<List<Integer>> sccs = g.findSCCs();
    assert sccs.size() == 2;
    // Sort SCCs for comparison
    List<List<Integer>> sorted = new ArrayList<>();
    for (List<Integer> scc : sccs) {
        List<Integer> s = new ArrayList<>(scc);
        Collections.sort(s);
        sorted.add(s);
    Collections.sort(sorted, (a, b) \rightarrow Integer.compare(a.get(0), b.get(0)));
    assert sorted.get(0).equals(Arrays.asList(0, 1, 2));
    assert sorted.get(1).equals(Arrays.asList(3, 4, 5));
}
```

Lca

```
lca.java
Lowest Common Ancestor (LCA) using Binary Lifting.
Preprocesses a tree to answer LCA queries efficiently.
Key operations:
- addEdge(u, v): Add undirected edge to tree
- build(root): Preprocess tree with given root - O(n log n)
- query(u, v): Find LCA of nodes u and v - O(\log n)
- distance(u, v): Find distance between two nodes - O(log n)
Space complexity: O(n \log n)
Binary lifting allows us to "jump" up the tree in powers of 2, enabling
efficient LCA queries.
import java.util.*;
class lca {
    static class LCA {
        private int n;
        private int maxLog;
        private Map<Integer, List<Integer>> graph;
        private int[] depth;
        private Map<Integer, Map<Integer, Integer>> up;
        LCA(int n) {
            this.n = n;
            this.maxLog = (int) Math.ceil(Math.log(n) / Math.log(2)) + 1;
            this.graph = new HashMap<>();
            this.depth = new int[n];
            this.up = new HashMap<>();
            for (int i = 0; i < n; i++) {
                graph.put(i, new ArrayList<>());
                up.put(i, new HashMap<>());
            }
        }
        void addEdge(int u, int v) {
            graph.get(u).add(v);
            graph.get(v).add(u);
        }
        void build(int root) {
            Arrays.fill(depth, 0);
            dfs(root, -1, 0);
        private void dfs(int node, int parent, int d) {
            depth[node] = d;
            if (parent != -1) {
                up.get(node).put(0, parent);
            for (int i = 1; i < maxLog; i++) {</pre>
                if (up.get(node).containsKey(i - 1)) {
                    int ancestor = up.get(node).get(i - 1);
                    if (up.get(ancestor).containsKey(i - 1)) {
                        up.get(node).put(i, up.get(ancestor).get(i - 1));
                    }
                }
            }
```

```
for (int child : graph.get(node)) {
            if (child != parent) {
                dfs(child, node, d + 1);
        }
    }
    int query(int u, int v) {
        if (depth[u] < depth[v]) {</pre>
            int temp = u;
            u = v;
            v = temp;
        }
        // Bring u to the same level as v
        int diff = depth[u] - depth[v];
        for (int i = 0; i < maxLog; i++) {
            if (((diff >> i) & 1) == 1) {
                if (up.get(u).containsKey(i)) {
                     u = up.get(u).get(i);
                }
            }
        }
        if (u == v) {
            return u;
        // Binary search for LCA
        for (int i = maxLog - 1; i >= 0; i--) {
            if (up.get(u).containsKey(i) && up.get(v).containsKey(i)) {
                int uAncestor = up.get(u).get(i);
                int vAncestor = up.get(v).get(i);
                if (uAncestor != vAncestor) {
                     u = uAncestor;
                     v = vAncestor;
                 }
            }
        return up.get(u).getOrDefault(0, u);
    }
    int distance(int u, int v) {
        int lcaNode = query(u, v);
        return depth[u] + depth[v] - 2 * depth[lcaNode];
    }
static void testMain() {
    LCA lca = new LCA(6);
    lca.addEdge(0, 1); // 1-2
lca.addEdge(0, 2); // 1-3
    lca.addEdge(1, 3); // 2-4
    lca.addEdge(1, 4); // 2-5
    lca.addEdge(2, 5); // 3-6
    lca.build(0);
    assert lca.query(3, 4) == 1; // LCA(4, 5) = 2
    assert lca.query(3, 5) == 0; // LCA(4, 6) = 1
    assert lca.distance(3, 5) == 4; // distance(4, 6) = 4
```

}

}

Polygon Area

```
polygon_area.java
Shoelace formula (Gauss's area formula) for computing the area of a polygon.
Computes the area of a simple polygon given its vertices in order (clockwise or
counter-clockwise). Works for both convex and concave polygons.
The formula: Area = 1/2 * |sum(x_i * y_i + 1) - x_i + y_i|
Time complexity: O(n) where n is the number of vertices.
Space complexity: O(1) additional space.
class polygon_area {
    static class Point {
        double x, y;
        Point(double x, double y) {
            this.x = x;
            this.y = y;
        }
    }
    static double polygonArea(Point[] vertices) {
        if (vertices.length < 3) {</pre>
            return 0.0;
        }
        int n = vertices.length;
        double area = 0.0;
        for (int i = 0; i < n; i++) {
            int j = (i + 1) \% n;
            area += vertices[i].x * vertices[j].y;
            area -= vertices[j].x * vertices[i].y;
        }
        return Math.abs(area) / 2.0;
    }
    static double polygonSignedArea(Point[] vertices) {
        if (vertices.length < 3) {</pre>
            return 0.0;
        int n = vertices.length;
        double area = 0.0;
        for (int i = 0; i < n; i++) {
            int j = (i + 1) \% n;
            area += vertices[i].x * vertices[j].y;
            area -= vertices[j].x * vertices[i].y;
        }
        return area / 2.0;
    }
    static boolean isClockwise(Point[] vertices) {
        return polygonSignedArea(vertices) < 0;</pre>
    static void testMain() {
        // Simple square with side length 2
        Point[] square = {
            new Point(0.0, 0.0), new Point(2.0, 0.0), new Point(2.0, 2.0), new Point(0.0, 2.0)
        };
        assert Math.abs(polygonArea(square) - 4.0) < 1e-9;</pre>
```

```
// Triangle with base 3 and height 4
Point[] triangle = {new Point(0.0, 0.0), new Point(3.0, 0.0), new Point(1.5, 4.0)};
assert Math.abs(polygonArea(triangle) - 6.0) < 1e-9;

// Test orientation
Point[] ccwSquare = {
    new Point(0.0, 0.0), new Point(1.0, 0.0), new Point(1.0, 1.0), new Point(0.0, 1.0)};
assert !isClockwise(ccwSquare);
}</pre>
```

Prefix Tree

```
prefix_tree.java
Prefix Tree (Trie) implementation for efficient string prefix operations.
- insert(word): Add a word to the trie - O(m) where m is word length
- search(word): Check if exact word exists - O(m)
- startsWith(prefix): Check if any word starts with prefix - O(m)
- delete(word): Remove a word from the trie - O(m)
Space complexity: O(ALPHABET_SIZE * N * M) where N is number of words and M is average length
import java.util.*;
class prefix_tree {
    static class TrieNode {
        Map<Character, TrieNode> children;
        boolean isEndOfWord;
        TrieNode() {
            children = new HashMap<>();
            isEndOfWord = false;
        }
   }
    static class PrefixTree {
        private TrieNode root;
        PrefixTree() {
            root = new TrieNode();
        }
        void insert(String word) {
            TrieNode node = root;
            for (char c : word.toCharArray()) {
                node.children.putIfAbsent(c, new TrieNode());
                node = node.children.get(c);
            }
            node.isEndOfWord = true;
        }
        boolean search(String word) {
            TrieNode node = root;
            for (char c : word.toCharArray()) {
                if (!node.children.containsKey(c)) {
                    return false;
                node = node.children.get(c);
            return node.isEndOfWord;
        }
        boolean startsWith(String prefix) {
            TrieNode node = root;
            for (char c : prefix.toCharArray()) {
                if (!node.children.containsKey(c)) {
                    return false;
                node = node.children.get(c);
            return true;
        }
        boolean delete(String word) {
            return deleteHelper(root, word, 0);
        }
```

```
private boolean deleteHelper(TrieNode node, String word, int depth) {
        if (node == null) {
            return false;
        }
        if (depth == word.length()) {
            if (!node.isEndOfWord) {
                return false;
            node.isEndOfWord = false;
            return node.children.isEmpty();
        }
        char c = word.charAt(depth);
        if (!node.children.containsKey(c)) {
            return false;
        TrieNode child = node.children.get(c);
        boolean shouldDeleteChild = deleteHelper(child, word, depth + 1);
        if (shouldDeleteChild) {
            node.children.remove(c);
            return !node.isEndOfWord && node.children.isEmpty();
        return false;
    }
}
static void testMain() {
    PrefixTree trie = new PrefixTree();
    trie.insert("cat");
    trie.insert("car");
    trie.insert("card");
    assert trie.search("car");
    assert !trie.search("ca");
    assert trie.startsWith("car");
}
```

Priority Queue

```
priority_queue.java
Generic priority queue (min-heap) with update and remove operations.
Supports:
- push(item): Add item to heap - O(log n)
- pop(): Remove and return minimum item - O(log n)
- peek(): View minimum item without removing - O(1)
- update(old_item, new_item): Update item in heap - O(n)
- remove(item): Remove specific item - O(n)
Standard library alternatives:
- C++: std::priority_queue (basic operations only, no key-based updates/removal)
- Python: heapq module (min-heap only, no key-based updates/removal)
- Java: PriorityQueue class (basic operations only, no key-based updates/removal)
Space complexity: O(n)
import java.util.*;
class priority_queue {
    static class PriorityQueue<T extends Comparable<T>>> {
        private List<T> heap;
        PriorityQueue() {
            this.heap = new ArrayList<>();
        }
        void push(T item) {
            heap.add(item);
            siftUp(heap.size() - 1);
        }
        T pop() {
            if (heap.isEmpty()) {
                throw new IllegalStateException("Heap is empty");
            T item = heap.get(0);
            T last = heap.remove(heap.size() - 1);
            if (!heap.isEmpty()) {
                heap.set(0, last);
                siftDown(0);
            return item;
        }
        T peek() {
            if (heap.isEmpty()) {
                return null;
            return heap.get(0);
        }
        boolean contains(T item) {
            return heap.contains(item);
        }
        void update(T oldItem, T newItem) {
            int idx = heap.indexOf(oldItem);
            if (idx == -1) {
                throw new IllegalArgumentException("Item not in heap");
            heap.set(idx, newItem);
            if (newItem.compareTo(oldItem) < 0) {</pre>
                siftUp(idx);
            } else {
```

```
siftDown(idx);
        }
    }
    void remove(T item) {
        int idx = heap.indexOf(item);
        if (idx == -1) {
            throw new IllegalArgumentException("Item not in heap");
        }
        T last = heap.remove(heap.size() - 1);
        if (idx < heap.size()) {</pre>
            T oldItem = heap.get(idx);
            heap.set(idx, last);
            if (last.compareTo(oldItem) < 0) {</pre>
                siftUp(idx);
            } else {
                siftDown(idx);
            }
        }
    }
    int size() {
        return heap.size();
    }
    boolean isEmpty() {
        return heap.isEmpty();
    private void siftUp(int idx) {
        while (idx > 0) {
            int parent = (idx - 1) / 2;
            if (heap.get(idx).compareTo(heap.get(parent)) >= 0) {
            Collections.swap(heap, idx, parent);
            idx = parent;
        }
    }
    private void siftDown(int idx) {
        while (true) {
            int smallest = idx;
            int left = 2 * idx + 1;
            int right = 2 * idx + 2;
            if (left < heap.size() && heap.get(left).compareTo(heap.get(smallest)) < 0) {</pre>
                smallest = left;
            if (right < heap.size() && heap.get(right).compareTo(heap.get(smallest)) < 0) {</pre>
                smallest = right;
            if (smallest == idx) {
                break;
            Collections.swap(heap, idx, smallest);
            idx = smallest;
        }
    }
static void testMain() {
    PriorityQueue<Integer> p = new PriorityQueue<>();
    p.push(15);
    p.push(23);
    p.push(8);
    assert p.peek() == 8;
    assert p.pop() == 8;
    assert p.pop() == 15;
```

}

}

Segment Tree

```
segment_tree.java
Segment Tree for range queries and point updates.
Supports efficient range queries (sum, min, max, etc.) and point updates on an array.
Key operations:
- update(i, value): Update element at index i - O(log n)
- query(l, r): Query range [l, r] - O(log\ n)
Space complexity: O(4n) = O(n)
This implementation supports sum queries but can be modified for min/max/gcd/etc.
import java.util.*;
import java.util.function.BinaryOperator;
class segment_tree {
    static class SegmentTree<T> {
        private Object[] tree;
        private int n;
        private T zero;
        private BinaryOperator<T> combineOp;
        SegmentTree(T[] arr, T zero, BinaryOperator<T> combineOp) {
            this.n = arr.length;
            this.zero = zero;
            this.combineOp = combineOp;
            this.tree = new Object[4 * n];
            if (n > 0) {
                build(arr, 0, 0, n - 1);
            }
        }
        private void build(T[] arr, int node, int start, int end) {
            if (start == end) {
                tree[node] = arr[start];
            } else {
                int mid = (start + end) / 2;
                int leftChild = 2 * node + 1;
                int rightChild = 2 * node + 2;
                build(arr, leftChild, start, mid);
                build(arr, rightChild, mid + 1, end);
                tree[node] = combineOp.apply((T) tree[leftChild], (T) tree[rightChild]);
            }
        }
        void update(int idx, T value) {
            update(0, 0, n - 1, idx, value);
        }
        @SuppressWarnings("unchecked")
        private void update(int node, int start, int end, int idx, T value) {
            if (start == end) {
                tree[node] = value;
            } else {
                int mid = (start + end) / 2;
                int leftChild = 2 * node + 1;
                int rightChild = 2 * node + 2;
                if (idx <= mid) {</pre>
                    update(leftChild, start, mid, idx, value);
                } else {
                    update(rightChild, mid + 1, end, idx, value);
```

```
}
            tree[node] = combineOp.apply((T) tree[leftChild], (T) tree[rightChild]);
        }
    }
    T query(int l, int r) {
        if (l < 0 || r >= n || l > r) {
            throw new IllegalArgumentException("Invalid range");
        return query(0, 0, n - 1, l, r);
    }
    @SuppressWarnings("unchecked")
    private T query(int node, int start, int end, int l, int r) {
        if (r < start \mid \mid l > end) {
            return zero;
        }
        if (l <= start && end <= r) {
            return (T) tree[node];
        }
        int mid = (start + end) / 2;
        int leftChild = 2 * node + 1;
        int rightChild = 2 * node + 2;
        T leftSum = query(leftChild, start, mid, l, r);
        T rightSum = query(rightChild, mid + 1, end, l, r);
        return combineOp.apply(leftSum, rightSum);
    }
}
static void testMain() {
    Long[] arr = \{1L, 3L, 5L, 7L, 9L\};
    SegmentTree<Long> st = new SegmentTree<>(arr, 0L, (a, b) -> a + b);
    assert st.query(1, 3) == 15L;
    st.update(2, 10L);
assert st.query(1, 3) == 20L;
    assert st.query(0, 4) == 30L;
}
```

Skiplist

```
skiplist.java
Skip list is a probabilistic data structure that maintains a sorted collection of elements.
It uses multiple levels of linked lists to achieve O(log n) average time complexity for
search, insertion, and deletion operations. Elements are inserted with randomly determined
heights, creating express lanes for faster traversal.
Standard library alternatives:
- C++: std::set / std::map (red-black tree, O(log n) guaranteed)
- Python: No built-in sorted set (use bisect module for sorted lists)
- Java: TreeSet / TreeMap (red-black tree, O(log n) guaranteed)
Time complexity: O(\log n) average for search, insert, and delete operations.
Space complexity: O(n) on average, where n is the number of elements.
import java.util.ArrayList;
import java.util.List;
import java.util.Random;
class SkipListNode<T extends Comparable<T>> {
   List<SkipListNode<T>> forward;
   SkipListNode(T value, int level) {
        this.value = value;
        this.forward = new ArrayList<>(level + 1);
        for (int i = 0; i <= level; i++) {</pre>
            this.forward.add(null);
        }
    }
}
class SkipList<T extends Comparable<T>> {
    private int maxLevel;
   private double p;
    private int level;
    private SkipListNode<T> header;
    Random random;
    SkipList(int maxLevel, double p) {
        this.maxLevel = maxLevel;
        this.p = p;
        this.level = 0;
        this.header = new SkipListNode<>(null, maxLevel);
        this.random = new Random();
    }
    SkipList() {
        this(16, 0.5);
    private int randomLevel() {
        int lvl = 0;
        while (random.nextDouble() 
            lvl++;
        }
        return lvl;
    }
    SkipList<T> insert(T value) {
        List<SkipListNode<T>> update = new ArrayList<>(maxLevel + 1);
        for (int i = 0; i <= maxLevel; i++) {</pre>
            update.add(null);
        }
```

```
SkipListNode<T> current = header;
    for (int i = level; i >= 0; i--) {
        while (current.forward.get(i) != null
                && current.forward.get(i).value.compareTo(value) < 0) {
            current = current.forward.get(i);
        }
        update.set(i, current);
    }
    int lvl = randomLevel();
    if (lvl > level) {
        for (int i = level + 1; i <= lvl; i++) {</pre>
            update.set(i, header);
        level = lvl;
    }
    SkipListNode<T> newNode = new SkipListNode<>(value, lvl);
    for (int i = 0; i <= lvl; i++) {</pre>
        newNode.forward.set(i, update.get(i).forward.get(i));
        update.get(i).forward.set(i, newNode);
    }
    return this;
}
boolean search(T value) {
    SkipListNode<T> current = header;
    for (int i = level; i >= 0; i--) {
        while (current.forward.get(i) != null
                && current.forward.get(i).value.compareTo(value) < 0) {
            current = current.forward.get(i);
        }
    }
    current = current.forward.get(0);
    return current != null && current.value.compareTo(value) == 0;
}
boolean delete(T value) {
    List<SkipListNode<T>> update = new ArrayList<>(maxLevel + 1);
    for (int i = 0; i <= maxLevel; i++) {</pre>
        update.add(null);
    }
    SkipListNode<T> current = header;
    for (int i = level; i >= 0; i--) {
        while (current.forward.get(i) != null
                && current.forward.get(i).value.compareTo(value) < 0) {
            current = current.forward.get(i);
        update.set(i, current);
    }
    current = current.forward.get(0);
    if (current == null || current.value.compareTo(value) != 0) {
        return false;
    }
    for (int i = 0; i <= level; i++) {</pre>
        if (update.get(i).forward.get(i) != current) {
            break;
        update.get(i).forward.set(i, current.forward.get(i));
    while (level > 0 && header.forward.get(level) == null) {
        level--;
    return true;
```

```
}
   // Optional functionality (not always needed during competition)
    int size() {
        int count = 0;
        SkipListNode<T> current = header.forward.get(0);
        while (current != null) {
            count++;
            current = current.forward.get(0);
        return count;
   }
    List<T> toList() {
        List<T> result = new ArrayList<>();
        SkipListNode<T> current = header.forward.get(0);
        while (current != null) {
            result.add(current.value);
            current = current.forward.get(0);
        return result;
   }
    boolean contains(T value) {
        return search(value);
    }
}
public class skiplist {
    static void testMain() {
        SkipList<Integer> sl = new SkipList<>();
        sl.random = new Random(42);
        sl.insert(10).insert(20).insert(5).insert(15);
        assert sl.search(10);
        assert sl.search(20);
        assert !sl.search(25);
        assert sl.delete(10);
        assert !sl.search(10);
        assert !sl.delete(30);
        // Optional functionality (not always needed during competition)
        SkipList<Integer> sl2 = new SkipList<>();
        sl2.random = new Random(42);
        sl2.insert(3).insert(1).insert(4).insert(1).insert(5);
        assert sl2.size() == 5;
        List<Integer> expected = List.of(1, 1, 3, 4, 5);
        assert sl2.toList().equals(expected);
        assert sl2.contains(3);
        assert !sl2.contains(7);
   }
```

Sprague Grundy

```
sprague_grundy.java
Sprague-Grundy theorem implementation for impartial games (finite, acyclic, normal-play).
The Sprague-Grundy theorem states that every impartial game is equivalent to a Nim heap
of size equal to its Grundy number (nimber). For multiple independent games,
XOR the Grundy numbers to determine the combined game value.
- GrundyEngine(moveFunction): makes it easy to plug in any game.
- grundy(state): compute nimber for a state (must be hashable).
- grundyMulti(states): XOR of nimbers for independent subgames.
- isWinningPosition(states): true iff XOR != 0.
Includes implementations for:
- Nim (single heap).
- Subtraction game (allowed moves = {1,3,4}) with period detection.
- Kayles (bowling pins) with splits into subgames via array representation.
Requirements:
- State must be hashable and canonically represented (e.g., sorted arrays).
- moveFunction must not create cycles.
import java.util.*;
import java.util.function.Function;
public class sprague_grundy {
    // Minimum EXcludant: smallest non-negative integer not occurring in 'values'
    public static int mex(Collection<Integer> values) {
        Set<Integer> s = new HashSet<>(values);
        int g = 0;
        while (s.contains(g)) {
            g++;
        }
        return g;
    }
    public static class GrundyEngine<T> {
        protected final Function<T, Collection<T>> moves;
        private final Map<T, Integer> cache = new HashMap<>();
        public GrundyEngine(Function<T, Collection<T>> moveFunction) {
            this.moves = moveFunction;
        public int grundy(T state) {
            if (cache.containsKey(state)) {
                return cache.get(state);
            }
            Collection<T> nextStates = moves.apply(state);
            if (nextStates.isEmpty()) {
                cache.put(state, 0);
                return 0;
            }
            List<Integer> nimbers = new ArrayList<>();
            for (T nextState : nextStates) {
                nimbers.add(grundy(nextState));
            int result = mex(nimbers);
            cache.put(state, result);
            return result;
        }
```

```
public int grundyMulti(Collection<T> states) {
        int result = 0;
        for (T state : states) {
            result ^= grundy(state);
        return result;
    }
    public boolean isWinningPosition(Collection<T> states) {
        return grundyMulti(states) != 0;
}
// Wrapper class for Kayles segments with proper equals/hashCode
public static class KaylesState {
    private final int[] segments;
    private final int hashCode;
    public KaylesState(int[] segments) {
        this.segments = segments.clone();
        Arrays.sort(this.segments); // Ensure canonical form
        this.hashCode = Arrays.hashCode(this.segments);
    }
    public KaylesState(List<Integer> segments) {
        this(segments.stream().mapToInt(Integer::intValue).toArray());
    public int[] getSegments() {
        return segments.clone();
    @Override
    public boolean equals(Object obj) {
        if (this == obj) return true;
        if (obj == null || getClass() != obj.getClass()) return false;
        KaylesState that = (KaylesState) obj;
        return Arrays.equals(segments, that.segments);
    }
    @Override
    public int hashCode() {
        return hashCode;
    }
    @Override
    public String toString() {
        return Arrays.toString(segments);
    }
}
// Optional functionality (not always needed during competition)
public static Integer detectPeriod(List<Integer> seq, int minPeriod, Integer maxPeriod) {
    int n = seq.size();
    if (maxPeriod == null) {
        maxPeriod = n / 2;
    for (int p = minPeriod; p <= maxPeriod; p++) {</pre>
        boolean ok = true;
        for (int i = 0; i < n; i++) {</pre>
            if (!seq.get(i).equals(seq.get(i % p))) {
                ok = false;
                break;
        if (ok) {
            return p;
        }
    }
    return null;
```

```
}
public static Collection<Integer> nimMovesSingleHeap(int n) {
    List<Integer> moves = new ArrayList<>();
    for (int k = 0; k < n; k++) {
        moves.add(k); // leave 0..n-1
    return moves;
}
public static Function<Integer, Collection<Integer>> subtractionGameMovesFactory(
        Set<Integer> allowed) {
    List<Integer> allowedSorted = new ArrayList<>(allowed);
    Collections.sort(allowedSorted);
    return n -> {
        List<Integer> moves = new ArrayList<>();
        for (int d : allowedSorted) {
            if (d <= n) {
                moves.add(n - d);
        return moves;
    };
}
public static Collection<KaylesState> kaylesMovesHelper(KaylesState state) {
    Set<KaylesState> resultSet = new HashSet<>();
    int[] segments = state.getSegments();
    for (int idx = 0; idx < segments.length; idx++) {
        int n = segments[idx];
        if (n <= 0) continue;</pre>
        // Remove one pin at position i (0..n-1)
        for (int i = 0; i < n; i++) {
            int left = i;
            int right = n - i - 1;
            List<Integer> newSeg = new ArrayList<>();
            for (int j = 0; j < idx; j++) {
                newSeg.add(segments[j]);
            if (left > 0) newSeg.add(left);
            if (right > 0) newSeg.add(right);
            for (int j = idx + 1; j < segments.length; j++) {
                newSeg.add(segments[j]);
            resultSet.add(new KaylesState(newSeg));
        }
        // Remove two adjacent pins at position i,i+1 (0..n-2)
        for (int i = 0; i < n - 1; i++) {
            int left = i;
            int right = n - i - 2;
            List<Integer> newSeg = new ArrayList<>();
            for (int j = 0; j < idx; j++) {
                newSeg.add(segments[j]);
            if (left > 0) newSeg.add(left);
            if (right > 0) newSeg.add(right);
            for (int j = idx + 1; j < segments.length; j++) {
                newSeg.add(segments[j]);
            resultSet.add(new KaylesState(newSeg));
        }
    }
    return new ArrayList<>(resultSet);
```

```
}
public static Function<KaylesState, Collection<KaylesState>> kaylesMovesFactory() {
    return sprague_grundy::kaylesMovesHelper;
public static void testMain() {
    // Test Nim with larger values
    GrundyEngine<Integer = new GrundyEngine<>(sprague_grundy::nimMovesSingleHeap);
    assert eng.grundy(42) == 42;
    assert eng.grundyMulti(Arrays.asList(17, 23, 31)) == 25; // 17^23^31 = 25 assert eng.isWinningPosition(Arrays.asList(15, 27, 36)) == true; // 15^27^36 = 48 != 0
    // Test subtraction game {1,3,4} with period 7
    GrundyEngine<Integer> eng2 =
             new GrundyEngine<>(subtractionGameMovesFactory(Set.of(1, 3, 4)));
    assert eng2.grundy(14) == 0; // 14 % 7 = 0 \rightarrow grundy = 0
    assert eng2.grundy(15) == 1; // 15 % 7 = 1 \rightarrow grundy = 1
    assert eng2.grundy(18) == 2; // 18 % 7 = 4 \rightarrow grundy = 2
    // Test Kayles
    GrundyEngine<KaylesState> eng3 = new GrundyEngine<>(kaylesMovesFactory());
    assert eng3.grundy(new KaylesState(new int[] \{7\})) == 2; // K(7) = 2
    assert eng3.grundy(new KaylesState(new int[] \{3, 5\})) == 7; // K(3)^{\wedge}K(5) = 3^{\wedge}4 = 7
}
```

Suffix Array

```
suffix_array.java
Suffix Array construction with Longest Common Prefix (LCP) array using Kasai's algorithm.
A suffix array is a sorted array of all suffixes of a string. The LCP array stores the length
of the longest common prefix between consecutive suffixes in the suffix array.
Key operations:
- Constructor: Build suffix array and LCP array for a string
- findPattern(pattern): Find all occurrences of pattern in text
Time complexity: O(n \log n) for suffix array, O(n) for LCP array
Space complexity: O(n)
import java.util.*;
class suffix_array {
    static class SuffixArray {
        private String text;
        private int n;
        private int[] sa;
        private int[] lcp;
        SuffixArray(String text) {
            this.text = text;
            this.n = text.length();
            this.sa = buildSuffixArray();
            this.lcp = buildLCPArray();
        }
        private int[] buildSuffixArray() {
            Integer[] suffixes = new Integer[n];
            for (int i = 0; i < n; i++) {
                suffixes[i] = i;
            Arrays.sort(suffixes, (a, b) -> text.substring(a).compareTo(text.substring(b)));
            int[] result = new int[n];
            for (int i = 0; i < n; i++) {
                result[i] = suffixes[i];
            return result;
        }
        private int[] buildLCPArray() {
            if (n == 0) return new int[0];
            int[] rank = new int[n];
            for (int i = 0; i < n; i++) {
                rank[sa[i]] = i;
            int[] lcp = new int[n];
            int h = 0;
            for (int i = 0; i < n; i++) {
                if (rank[i] > 0) {
                    int j = sa[rank[i] - 1];
                    while (i + h < n \&\& j + h < n \&\& text.charAt(i + h) == text.charAt(j + h)) {
                        h++;
                    lcp[rank[i]] = h;
                    if (h > 0) h--;
                }
            }
```

```
return lcp;
    }
    List<Integer> findPattern(String pattern) {
        if (pattern.isEmpty()) return new ArrayList<>();
        int m = pattern.length();
        int left = 0, right = n;
        while (left < right) {</pre>
            int mid = (left + right) / 2;
            String suffix = text.substring(sa[mid]);
            if (suffix.compareTo(pattern) < 0) {</pre>
                 left = mid + 1;
            } else {
                right = mid;
        }
        int start = left;
        left = start;
        right = n;
        while (left < right) {</pre>
            int mid = (left + right) / 2;
            String suffix = text.substring(sa[mid], Math.min(sa[mid] + m, n));
            if (suffix.compareTo(pattern) <= 0) {</pre>
                left = mid + 1;
            } else {
                right = mid;
        }
        int end = left;
        List<Integer> result = new ArrayList<>();
        for (int i = start; i < end; i++) {</pre>
            if (sa[i] + m <= n && text.substring(sa[i], sa[i] + m).equals(pattern)) {
                result.add(sa[i]);
        }
        Collections.sort(result);
        return result;
    }
    int[] getSA() {
        return sa;
    int[] getLCP() {
        return lcp;
    }
static void testMain() {
    SuffixArray sa = new SuffixArray("banana");
    assert Arrays.equals(sa.getSA(), new int[] {5, 3, 1, 0, 4, 2});
    assert Arrays.equals(sa.getLCP(), new int[] {0, 1, 3, 0, 0, 2});
    List<Integer> positions = sa.findPattern("ana");
    assert positions.equals(Arrays.asList(1, 3));
```

}

}

Topological Sort

```
topological_sort.java
Topological sorting algorithms for Directed Acyclic Graphs (DAG).
Provides two implementations:
1. Kahn's algorithm (BFS-based) - detects cycles
2. DFS-based algorithm - also detects cycles
Both return a topological ordering of vertices if the graph is a DAG,
or null if a cycle is detected.
Time complexity: O(V + E)
Space complexity: O(V + E)
import java.util.*;
class topological_sort {
    static class TopologicalSort {
        private int n;
        private Map<Integer, List<Integer>> graph;
        TopologicalSort(int n) {
            this.n = n;
            this.graph = new HashMap<>();
            for (int i = 0; i < n; i++) {
                graph.put(i, new ArrayList<>());
            }
        }
        void addEdge(int u, int v) {
            graph.get(u).add(v);
        }
        List<Integer> kahnSort() {
            int[] inDegree = new int[n];
            for (int u = 0; u < n; u++) {
                for (int v : graph.get(u)) {
                    inDegree[v]++;
                }
            Queue<Integer> queue = new LinkedList<>();
            for (int i = 0; i < n; i++) {
                if (inDegree[i] == 0) {
                    queue.offer(i);
                }
            }
            List<Integer> result = new ArrayList<>();
            while (!queue.isEmpty()) {
                int u = queue.poll();
                result.add(u);
                for (int v : graph.get(u)) {
                    inDegree[v]--;
                    if (inDegree[v] == 0) {
                        queue.offer(v);
                    }
                }
            }
            if (result.size() != n) {
                return null; // Cycle detected
            }
```

```
return result;
}
List<Integer> dfsSort() {
    Set<Integer> visited = new HashSet<>();
    Set<Integer> recStack = new HashSet<>();
    Stack<Integer> stack = new Stack<>();
    for (int i = 0; i < n; i++) {
        if (!visited.contains(i)) {
            if (!dfsVisit(i, visited, recStack, stack)) {
    return null; // Cycle detected
        }
    }
    List<Integer> result = new ArrayList<>();
    while (!stack.isEmpty()) {
        result.add(stack.pop());
    return result;
}
private boolean dfsVisit(
        int u, Set<Integer> visited, Set<Integer> recStack, Stack<Integer> stack) {
    visited.add(u);
    recStack.add(u);
    for (int v : graph.get(u)) {
        if (!visited.contains(v)) {
            if (!dfsVisit(v, visited, recStack, stack)) {
                return false;
        } else if (recStack.contains(v)) {
            return false; // Cycle detected
    }
    recStack.remove(u);
    stack.push(u);
    return true;
}
boolean hasCycle() {
    return kahnSort() == null;
}
List<Integer> longestPath(int source) {
    List<Integer> topoOrder = kahnSort();
    if (topoOrder == null) {
        return null; // Has cycle
    Map<Integer, Integer> dist = new HashMap<>();
    Map<Integer, Integer> parent = new HashMap<>();
    for (int i = 0; i < n; i++) {
        dist.put(i, Integer.MIN_VALUE);
    }
    dist.put(source, 0);
    for (int u : topoOrder) {
        if (dist.get(u) != Integer.MIN_VALUE) {
            for (int v : graph.get(u)) {
                 if (dist.get(u) + 1 > dist.get(v)) {
                     dist.put(v, dist.get(u) + 1);
                     parent.put(v, u);
                }
            }
        }
```

```
}
        // Find vertex with maximum distance
        int maxDist = Integer.MIN_VALUE;
        int endVertex = -1;
        for (int i = 0; i < n; i++) {</pre>
            if (dist.get(i) > maxDist) {
                maxDist = dist.get(i);
                endVertex = i;
            }
        }
        if (endVertex == -1 || maxDist == Integer.MIN_VALUE) {
            return Arrays.asList(source);
        }
        // Reconstruct path
        List<Integer> path = new ArrayList<>();
        int current = endVertex;
        while (current != source) {
            path.add(current);
            current = parent.get(current);
        path.add(source);
        Collections.reverse(path);
        return path;
    }
}
static void testMain() {
    TopologicalSort ts = new TopologicalSort(6);
    int[][] edges = {{5, 2}, {5, 0}, {4, 0}, {4, 1}, {2, 3}, {3, 1}};
    for (int[] edge : edges) {
        ts.addEdge(edge[0], edge[1]);
    }
    List<Integer> kahnResult = ts.kahnSort();
    List<Integer> dfsResult = ts.dfsSort();
    assert kahnResult != null;
    assert dfsResult != null;
    assert !ts.hasCycle();
    // Test with cycle
    TopologicalSort tsCycle = new TopologicalSort(3);
    tsCycle.addEdge(0, 1);
    tsCycle.addEdge(1, 2);
    tsCycle.addEdge(2, 0);
    assert tsCycle.hasCycle();
}
```

Two Sat

```
two_sat.java
2-SAT solver using Kosaraju's SCC algorithm on implication graph.
2-SAT (Boolean Satisfiability with 2 literals per clause) determines if a Boolean formula
in CNF with at most 2 literals per clause is satisfiable. Uses implication graph where
each variable x has nodes x and not-x, and clause (a OR b) creates edges not-a -> b
and not-b \rightarrow a.
Key operations:
- addClause(a, aNeg, b, bNeg): Add clause (a OR b)
- solve(): Returns assignment array if satisfiable, null otherwise
Time complexity: O(n + m) where n is variables and m is clauses
Space complexity: O(n + m)
import java.util.*;
class two_sat {
    static class TwoSAT {
        private int n;
        private List<List<Integer>> graph;
        private List<List<Integer>> transpose;
        TwoSAT(int n) {
            this.n = n;
            this.graph = new ArrayList<>();
            this.transpose = new ArrayList<>();
            for (int i = 0; i < 2 * n; i++) {
                graph.add(new ArrayList<>());
                transpose.add(new ArrayList<>());
            }
        }
        void addClause(int a, boolean aNeg, int b, boolean bNeg) {
            int aNode = 2 * a + (aNeg ? 1 : 0);
            int bNode = 2 * b + (bNeg ? 1 : 0);
            int naNode = 2 * a + (aNeg ? 0 : 1);
            int nbNode = 2 * b + (bNeg ? 0 : 1);
            graph.get(naNode).add(bNode);
            graph.get(nbNode).add(aNode);
            transpose.get(bNode).add(naNode);
            transpose.get(aNode).add(nbNode);
        }
        boolean[] solve() {
            // Kosaraju's algorithm
            boolean[] visited = new boolean[2 * n];
            List<Integer> finishOrder = new ArrayList<>();
            for (int node = 0; node < 2 * n; node++) {</pre>
                if (!visited[node]) {
                    dfs1(node, visited, finishOrder);
            }
            Arrays.fill(visited, false);
            int[] sccId = new int[2 * n];
            int currentScc = 0;
            for (int i = finishOrder.size() - 1; i >= 0; i--) {
                int node = finishOrder.get(i);
                if (!visited[node]) {
                    dfs2(node, visited, sccId, currentScc);
                    currentScc++;
```

```
// Check satisfiability
        for (int i = 0; i < n; i++) {</pre>
            if (sccId[2 * i] == sccId[2 * i + 1]) {
                 return null;
        }
        // Construct assignment
        boolean[] assignment = new boolean[n];
        for (int i = 0; i < n; i++) {
    assignment[i] = sccId[2 * i] > sccId[2 * i + 1];
        }
        return assignment;
    }
    private void dfs1(int node, boolean[] visited, List<Integer> finishOrder) {
        visited[node] = true;
        for (int neighbor : graph.get(node)) {
            if (!visited[neighbor]) {
                 dfs1(neighbor, visited, finishOrder);
        finishOrder.add(node);
    }
    private void dfs2(int node, boolean[] visited, int[] sccId, int scc) {
        visited[node] = true;
        sccId[node] = scc;
        for (int neighbor : transpose.get(node)) {
            if (!visited[neighbor]) {
                 dfs2(neighbor, visited, sccId, scc);
            }
        }
    }
}
static void testMain() {
    TwoSAT sat = new TwoSAT(2);
    sat.addClause(0, false, 1, false);
    sat.addClause(0, true, 1, false);
    sat.addClause(0, false, 1, true);
    boolean[] result = sat.solve();
    assert result != null;
    assert result[0] || result[1];
    assert !result[0] || result[1];
    assert result[0] || !result[1];
}
```

Union Find

```
union_find.java
Union-Find (Disjoint Set Union) data structure with path compression and union by rank.
Supports:
- find(x): Find the representative of the set containing x - O(\alpha(n)) amortized
- union(x, y): Merge the sets containing x and y - O(\alpha(n)) amortized
- connected(x, y): Check if x and y are in the same set - O(\alpha(n)) amortized
Space complexity: O(n)
\alpha(n) is the inverse Ackermann function, which is effectively constant for all practical values of n.
class union_find {
    static class UnionFind {
        private int[] parent;
        private int[] rank;
        UnionFind(int n) {
            parent = new int[n];
            rank = new int[n];
            for (int i = 0; i < n; i++) {
                parent[i] = i;
                rank[i] = 0;
            }
        }
        int find(int x) {
            if (parent[x] != x) {
                parent[x] = find(parent[x]); // Path compression
            return parent[x];
        }
        int union(int x, int y) {
            int rootX = find(x);
            int rootY = find(y);
            if (rootX == rootY) {
                return rootX;
            // Union by rank
            if (rank[rootX] < rank[rootY]) {</pre>
                parent[rootX] = rootY;
                merge(rootY, rootX);
                return rootY;
            } else if (rank[rootX] > rank[rootY]) {
                parent[rootY] = rootX;
                merge(rootX, rootY);
                return rootX;
            } else {
                parent[rootY] = rootX;
                merge(rootX, rootY);
                rank[rootX]++;
                return rootX;
            }
        }
        boolean connected(int x, int y) {
            return find(x) == find(y);
        }
        void merge(int root, int child) {
            // Override to define custom merge behavior when sets are united
        }
```

```
}
static class Test extends UnionFind {
    private int[] size;
    Test(int n) {
         super(n);
        size = new int[n];
for (int i = 0; i < n; i++) {</pre>
             size[i] = 1;
         }
    }
    @Override
    void merge(int root, int child) {
        size[root] += size[child];
    }
    int getSize(int x) {
        return size[find(x)];
    }
}
static void testMain() {
    Test a = new Test(3);
    int d = a.union(0, 1);
int e = a.union(d, 2);
    assert a.getSize(e) == 3;
    assert a.getSize(a.find(0)) == 3;
}
```