

Local stability in structures with a standard sort

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ABSTRACT.

1. Structures with a standard sort

Let S be some Hausdorff compact topological space. We associate to S a first order structure in a language \mathcal{L}_S that has a symbol for each compact subset $C \subseteq S$ and a function symbol for each continuous functions $f : S^n \rightarrow S$. According to the context, C and f denote either the symbols of \mathcal{L}_S or their interpretation in the structure S . We write $x \in C$ for $C(x)$.

Finally, note that we could allow in \mathcal{L}_S relation symbols for all compact subsets of S^n , for any n . But this would clutter the notation adding very little to the theory.

1 Definition Let \mathcal{L} be a two sorted language. The two sorts are denoted by H and S . The language \mathcal{L} expands \mathcal{L}_H and \mathcal{L}_S with symbols sort $H^n \times S^m \rightarrow S$. An \mathcal{L} -structure is a structure of signature \mathcal{L} that interprets these symbols in equicontinuous functions. A **standard structure** is a two-sorted \mathcal{L} -structure of the form $\langle M, S \rangle$, where M is any structure of signature \mathcal{L}_H and S is fixed. Standard structures are denoted by the domain of their home sort.

We denote by \mathcal{F} the set of \mathcal{L} -formulas constructed inductively from atomic formulas of the form (i) and (ii) below using Boolean connectives \wedge, \vee ; the quantifiers \forall^H, \exists^H of sort H ; and the quantifiers \forall^S, \exists^S of sort S .

2 Definition We call atomic formulas those of the form

- i. atomic or negated atomic formulas of \mathcal{L}_H ;
- ii. $\tau \in C$, where $C \subseteq S$ is a compact set and τ is a term of sort $H^n \times S^m \rightarrow S$.

3 Remark A smaller fragment of \mathcal{F} may be of interest in some specific contexts. This is obtained by excluding equalities and inequalities from (i) and requiring that C in (ii) is a regular compact set. The discussion below applies to this smaller fragment as well.

Let $M \subseteq N$ be standard structures. We say that M is an \mathcal{F} -elementary substructure of N if the latter models all sentences in $\mathcal{F}(M)$ true in M .

A standard structure M is \mathcal{F} -saturated if it realizes every type $p(x; \xi) \subseteq \mathcal{F}(A)$, for any $A \subseteq M$ of cardinality smaller than $|M|$, that is finitely consistent in M . The following is proved in [AAVV] for signatures containing only symbols sort $H^n \times S^m \rightarrow S$ with $n \cdot m = 0$. In [Z] it is observed that, under the assumption of equicontinuity, the proof extends to the case $n \cdot m \geq 0$.

4 Theorem (Compactness) Every standard structure has a \mathcal{F} -saturated \mathcal{F} -elementary extension.

To conveniently describe classes of \mathcal{F} that only differ by the sets/predicates $C \subseteq S$ that occur in the formula, we introduce variables of new sort X . These will be used as placeholders for compact subsets of S . Let \mathcal{F}_X be defined as \mathcal{F} but replacing (ii) with

iii. $\tau(x; \xi) \in X$, where X is a variable of sort X .

Formulas in \mathcal{F}_X are denoted by $\varphi(x; \xi; X)$, where $X = X_1, \dots, X_k$ be a tuple of variables of sort X . If $C = C_1, \dots, C_k$ is a tuple of compact subsets of S then $\varphi(x; \xi; C)$, is a formula in \mathcal{F} . This we call an instance of $\varphi(x; \xi; X)$. All formulas in \mathcal{F} are instances of formulas in \mathcal{F}_X .

5 Definition Let $\varphi \in \mathcal{F}_X$ be formula – possibly with some (hidden) free variables. The pseudonegation of $\varphi \in \mathcal{F}_X$ is the formula obtained by replacing in φ the atomic formulas in \mathcal{L}_H by their negation and each connective $\wedge, \vee, \forall^H, \exists^H, \forall^S, \exists^S$ by its respective dual $\vee, \wedge, \exists^H, \forall^H, \exists^S, \forall^S$.

The pseudonegation of φ is denoted by $\sim\varphi$. Clearly, when X does not occur in φ , we have $\sim\varphi \leftrightarrow \neg\varphi$.

The following fact will be used without further mention.

6 Fact Let $\varphi(X) \in \mathcal{F}_X$ have, possibly, hidden variables of sort H and/or S . Then for every (tuples of) compact sets $C \subseteq C'$ and $\tilde{C} \cap C = \emptyset$.

$$\begin{aligned} \varphi(C) &\rightarrow \varphi(C') \\ \sim\varphi(C) &\rightarrow \neg\varphi(\tilde{C}) \end{aligned}$$

A standard structure M is \mathcal{F} -maximal if it models all $\mathcal{F}(M)$ -sentences that hold in some of its \mathcal{F} -elementary extensions.

7 Fact Let $\varphi(X) \in \mathcal{F}_X(M)$. Then the following are equivalent

1. M is \mathcal{F} -maximal
2. for every tuple C of compact subsets of S

$$M \models \varphi(C) \leftrightarrow \bigwedge \{ \varphi(C') : C' \text{ neighborhood of } C \}$$
3. for every tuple C of compact subsets of S

$$M \models \neg \varphi(C) \leftrightarrow \bigvee \{ \neg \varphi(\tilde{C}) : \tilde{C} \cap C = \emptyset \}.$$

Moreover, all of the above hold when M is \mathcal{F} -saturated.

Proof. Arguing by induction on the syntax of $\varphi(X)$ we can prove that (2) holds when M is \mathcal{F} -saturated, see [AAVW].

$2 \Leftrightarrow 3$. By normality.

$1 \Rightarrow 2$. If (2) fails, then $\neg \varphi(C) \wedge \neg \neg \varphi(\tilde{C})$ holds in M for every $\tilde{C} \cap C = \emptyset$. Let \mathcal{U} be an \mathcal{F} -saturated \mathcal{F} -elementary extension of M . As M is \mathcal{F} -maximal, $\neg \varphi(C) \wedge \neg \neg \varphi(\tilde{C})$ also holds in \mathcal{U} . Then (2) fails also in \mathcal{U} , a contradiction.

$2 \Rightarrow 1$. Let N be any \mathcal{F} -elementary extension of M . Suppose $N \models \varphi(C)$. By (2), it suffices to prove that $M \models \varphi(C')$ for every C' neighborhood of C . Suppose not. Then, by (3), $M \models \neg \varphi(\tilde{C})$ for some $\tilde{C} \cap C' = \emptyset$. Then $N \models \neg \varphi(\tilde{C})$ by \mathcal{F} -elementarity. As $\tilde{C} \cap C = \emptyset$, this contradicts the above fact. \square

2. Externally definable sets

In what follows we fix a large saturated standard structure which we denote by \mathcal{U} . Unless otherwise specified, we work inside \mathcal{U} .

Let $\varphi(x; z; X)$ be a formula in \mathcal{F}_X with the variables of sort H partitioned in two tuples $x; z$. A **global $\varphi(x; z; X)$ -type** is a maximally consistent set of formulas that are instances of $\varphi(x; b; X)$ and/or $\neg \varphi(x; b; X)$ for some $b \in \mathcal{U}^{|z|}$.

Let $p(x)$ be a global $\varphi(x; z; X)$ -type. We define

$$\mathcal{D}_{p, \varphi, C} = \{ b \in \mathcal{U}^z : \varphi(x; b; C) \in p \}$$

8 Fact $\neg \mathcal{D}_{p, \varphi, C} = \bigcup \{ \mathcal{D}_{p, \sim \varphi, \tilde{C}} : \tilde{C} \cap C = \emptyset \}$

Proof. (\subseteq) . Let $b \notin \mathcal{D}_{p, \varphi, C}$. Let $A \subseteq \mathcal{U}$ be such that every instance of $\varphi(x; b; X)$ and $\neg \varphi(x; b; X)$ that is consistent with $p(x) \upharpoonright A$ is finitely consistent with $p(x)$. Finally pick $a \models p(x) \upharpoonright A$. As $\neg \varphi(a, b, C)$, from Fact 7 we obtain that $\neg \varphi(a; b; \tilde{C})$ holds for some $\tilde{C} \cap C = \emptyset$. Then $\neg \varphi(x; b; \tilde{C})$ is consistent with $p(x) \upharpoonright A$. By maximality of $p(x)$, we have $\neg \varphi(x; b; \tilde{C}) \in p$. Then $b \in \mathcal{D}_{p, \sim \varphi, \tilde{C}}$.

(\supseteq) . Similar. \square

9 Fact Let $\varphi(x; z; X)$ and C be given. Then for every $\mathcal{D} \subseteq \mathcal{U}^z$ the following are equivalent

1. $\mathcal{D} = \mathcal{D}_{p, \varphi, C}$ for some global $\varphi(x; z; X)$ -type $p(x)$
2. for every finite set $B \subseteq \mathcal{U}^z$ there is an $a \in \mathcal{U}^x$ such that

$$B \cap \mathcal{D} = B \cap \varphi(a; \mathcal{U}^x; C).$$

Proof. Let $B_0 = B \setminus \mathcal{D}$ and $B_1 = B \cap \mathcal{D}$.

1 \Rightarrow 2. By the above fact, there is a $\tilde{C} \cap C = \emptyset$ such that $b \in \mathcal{D}_{p, \sim \varphi, \tilde{C}}$ for every $b \in B_0$. Let a realize the type containing $\varphi(x; b; C)$ for every $b \in B_1$ and $\sim \varphi(x; b; \tilde{C})$ for every $b \in B_0$.

2 \Rightarrow 1. There is a set $\tilde{C} \cap C = \emptyset$ such that $\sim \varphi(x; b; \tilde{C})$ for every $b \in B_0$. Let $p(x)$ be any global $\varphi(x; z; X)$ -type containing $\varphi(x; b; C)$ for every $b \in B_1$ and $\sim \varphi(x; b; \tilde{C})$ for every $b \in B_0$. Such type exists as (2) guarantees consistency. \square

Let $\mathcal{D} \subseteq \mathcal{U}^z \times S^{|X|}$. For $\alpha \in S^{|X|}$, we write \mathcal{D}_α for the set of all $b \in \mathcal{U}^z$ such that $\langle b, \alpha \rangle \in \mathcal{D}$. We say that \mathcal{D} is **externally definable** by $\varphi(x; z; X)$ if there is a global $\varphi(x; z; X)$ -type $p(x)$ such that for every α

$$\mathcal{D}_\alpha = \bigcap \{ \mathcal{D}_{p, \varphi, C} : \alpha \in C^\circ \}$$

We say that \mathcal{D} is **approximable** by $\varphi(x; z; X)$ if for every $\alpha \in S^{|X|}$ and every finite $B \subseteq \mathcal{U}^z$ there is an $a \in \mathcal{U}^{|X|}$ such that

$$B \cap \mathcal{D}_\alpha = B \cap \bigcap \{ \varphi(a; \mathcal{U}^z; C) : \alpha \in C^\circ \}$$

10 Fact Let $\varphi(x; z; X)$ be given. Then, for every $\mathcal{D} \subseteq \mathcal{U}^z \times S^{|X|}$ the following are equivalent

1. \mathcal{D} is externally definable by $\varphi(x; z; X)$
2. \mathcal{D} is approximated by $\varphi(x; z; X)$.

We say that \mathcal{D} is **approximable from below** if for every $B \subseteq \mathcal{D}$ there is an $a \in \mathcal{U}^{|X|}$ such that

$$B \subseteq \varphi(a; \mathcal{U}^z; C) \subseteq \mathcal{D}.$$

3. Stable formulas

We say that $\varphi(x; z; X)$ is unstable if there are some compact sets $C \cap \tilde{C} = \emptyset$ and a sequence $\langle a_i; b_i : i < \omega \rangle$ such that for every $i < j < \omega$

$$\varphi(a_i; b_j; C) \wedge \sim \varphi(a_j; b_i; \tilde{C})$$

We say that $\varphi(x; z; X)$ is **stable** if it is not unstable.

The following fact is a characterization of stability in a context of particular interest. Let $\tau : \mathcal{U}^{x; z} \rightarrow S$ be a definable function. We say that $\tau(x; z)$ is **stable** if so is the formula $\tau(x; z) \in X$.

11 Fact Let $\tau : \mathcal{U}^{x;z} \rightarrow S$ be a definable function. Then the following are equivalent

1. the formula $\tau(x; z) \in X$ is unstable
2. for some sequence $\langle a_i; b_i : i < \omega \rangle$

$$\lim_{i \rightarrow \infty} \lim_{j \rightarrow \infty} \tau(a_i; b_j) \neq \lim_{j \rightarrow \infty} \lim_{i \rightarrow \infty} \tau(a_i; b_j)$$

3. for some sequence $\langle a_i; b_i : i < \omega \rangle$ and some neighborhood ε of the diagonal of S^2

$$\langle \tau(a_i; b_j), \tau(a_j; b_i) \rangle \notin \varepsilon$$

for every $i < j < \omega$.

Proof. $2 \Leftrightarrow 3$. Clear.

$1 \Rightarrow 3$. Let $C \cap \tilde{C} = \emptyset$ be such that $\tau(a_i; b_j) \in C$ and $\tau(a_j; b_i) \in \tilde{C}$ for every $i < j < \omega$ as given by (1). Then (3) holds for any ε contained in the complement of $C \times \tilde{C}$.

$2 \Rightarrow 1$. Let C and \tilde{C} be two disjoint intervals centered in the two limits in (2). Then (1) is witnessed by a tail of the sequence $\langle a_i; b_i : i < \omega \rangle$. \square

12 Fact Let $S = [0, 1]$. Let $\tau(x; z)$ be a stable definable function. Let $f : \mathcal{U}^{x;z} \rightarrow S$. Assume that for every finite $B \subseteq \mathcal{U}^z$ and every $\varepsilon > 0$ there is an $a \in \mathcal{U}^x$ such that

$$|f(b) - \tau(a; b)| \leq \varepsilon \quad \text{for every } b \in B.$$

Then there are some $\langle a_{i,j} : 1 \leq i, j \leq n \rangle$ such that

$$f(b) = \max_{i=1, \dots, n} \min_{j=1, \dots, n} \tau(a_{i,j}; b) \quad \text{for every } b \in \mathcal{U}^z.$$

Proof. \square

13 Lemma Let $\varphi(x; z; X)$ be stable. Let $\mathcal{D} \subseteq \mathcal{U}^{|z|}$ be approximated by $\varphi(x; z; X)$ from below. Then

$$\mathcal{D}_{p, \varphi, C} = \bigcup_{i=1}^n \varphi(a_i; \mathcal{U}^z; C)$$

for some $\langle a_i : 1 \leq i \leq n \rangle$.

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We define $\mathcal{D}_p \subseteq \mathcal{U}^{|z|} \times S^{|x|}$ as follows

$$\mathcal{D}_p = \left\{ \langle b, \alpha \rangle : \varphi(x; b; C) \in p \text{ for every } C \text{ neighborhood of } \alpha \right\}.$$

14 Fact

$$\neg \mathcal{D}_p = \left\{ \langle b, \alpha \rangle : \sim \varphi(x; b; \tilde{C}) \in p \text{ for every } \tilde{C} \nexists \alpha \right\}.$$

Proof. Let $a \models p(x) \upharpoonright b$. Let C be the intersection of all neighborhoods of α such that $\varphi(a; b; C')$. Assume $\langle b, \alpha \rangle \notin \mathcal{D}_p$, in other words $\varphi(a; b; C)$, then C is a neighborhood of α . By Fact 7 we have $\varphi(a; b; C)$. Let $\tilde{C} \subseteq C \setminus \{\alpha\}$ be such that C is a neighborhood of \tilde{C} . Then $\neg\varphi(a; b; \tilde{C})$. Therefore by the definition of \mathcal{D}_p .

Then we can find \tilde{C} such that $\alpha \notin \tilde{C}$ and C is a neighborhood of \tilde{C} . contains two disjoint compacts Then $\sim\varphi(x; b; \tilde{C}) \in p$ by the definition of \mathcal{D}_p .

Then $\sim\varphi(x; b; \tilde{C}') \in p$ for every neighborhood \tilde{C}' of α disjoint from C' . Then for every neighborhood C of α , we have $\varphi(x; b; C) \in p$. \square

We say that \mathcal{D} is **externally defined** by $p(x)$ and $\varphi(x; z; X)$ if $\mathcal{D}_p \subseteq \mathcal{D}$ and $\tilde{\mathcal{D}}_p \subseteq \neg\mathcal{D}$.

We say that $\mathcal{D} \subseteq \mathcal{U}^{|z|} \times S^{|x|}$ is **approximated** by $\varphi(x; z; X)$ if for every finite $B \subseteq \mathcal{U}^{|z|} \times S^{|x|}$ there is an $a \in \mathcal{U}^{|x|}$ such that

$$\begin{aligned} \langle b, \alpha \rangle \in B \cap \mathcal{D} &\Rightarrow \varphi(a; b; C) \text{ for every neighborhood } C \text{ of } \alpha \\ \langle b, \alpha \rangle \in B \setminus \mathcal{D} &\Rightarrow \sim\varphi(a; b; C) \text{ for every } C \text{ such that } \alpha \notin C \end{aligned}$$

15 Fact For every $\mathcal{D} \subseteq \mathcal{U}^{|z|} \times S^{|x|}$, the following are equivalent

1. \mathcal{D} is externally definable by some global φ -type
2. \mathcal{D} is approximated by $\varphi(x; z; X)$.

Proof. (1) \Rightarrow (2). Let $B \subseteq \mathcal{U}^{|z|} \times S^{|x|}$ be finite. For $\langle b, \alpha \rangle \in B \cap \mathcal{D}$ let $C_{\langle b, \alpha \rangle}$ be some neighborhood of α . Let C be the union of all these neighborhoods. For $\langle b, \alpha \rangle \in B \setminus \mathcal{D}$ let $\tilde{C}_{\langle b, \alpha \rangle}$ be some neighborhood of α . Let \tilde{C} be the union of all these neighborhoods. By the finiteness of B , we can require that C and \tilde{C} are disjoint. Then (2) follows from the consistency of $p(x)$.

(2) \Rightarrow (1). Note that the condition of approximability asserts the finite consistency of the type $p(x)$ that is union of the following two sets of formulas

$$\begin{aligned} &\left\{ \varphi(x; b; C) : \langle b, \alpha \rangle \in \mathcal{D} \text{ and } C \text{ neighborhood of } \alpha \right\} \\ &\left\{ \sim\varphi(x; b; C) : \langle b, \alpha \rangle \notin \mathcal{D} \text{ and } \alpha \notin C \right\} \end{aligned}$$

The inclusion $\mathcal{D} \subseteq \mathcal{D}_p$ is immediate. For the converse inclusion, let $\langle b, \alpha \rangle \notin \mathcal{D}$ and let \tilde{C} be a neighborhood of α such that $\sim\varphi(x; b; \tilde{C})$. \square

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